15-751*, Fall 2016:

A Theorist's Toolkit

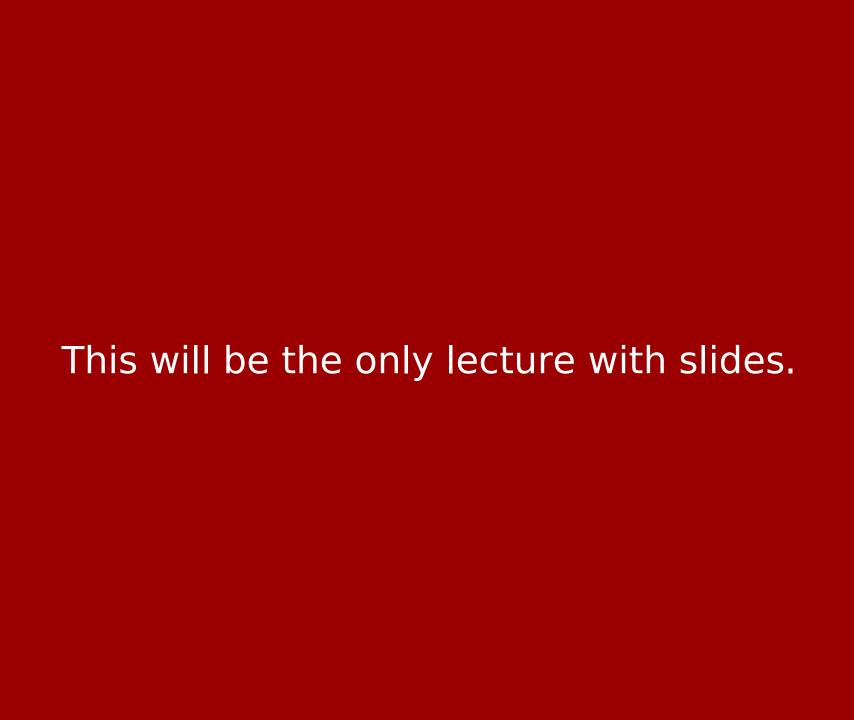
Lecture 1



Venkat Guruswami



Ryan O'Donnell



Today's Lecture

Part 1: What this course is about.

(Answer: CS theory and math.)

Part 2: How to present mathematics.

Part 3: How to do mathematics.

Part 1: What this course is about.

A Theorist's Toolkit

A Theorist's Toolkit

The indefinite article.

Moving on.

A Theorist's Toolkit

Theoretical Computer Science (TCS)

= Algorithms and Computational Complexity

A Theorist's Toolkit

A brief introduction to a number of topics and tools

– especially mathematical ones – that arise in TCS research.

An approximate syllabus

Asymptotics Central Limit Theorem **Chernoff Bounds Computational Models** Spectral Graph Theory Fields and polynomials Error-correcting codes Derandomization **Linear Programming** Semidefinite Programming

Constraint Satisfaction Problems Communication Complexity Information Theory High-dimensional Geometry Streaming & Property Testing Analysis of Boolean Functions **Proof Complexity Quantum Computation** Cryptography Probabilistically Checkable Proofs

More breadth than depth...

The "background knowledge" that it helps to know for...

- Reading TCS papers
- Attending TCS talks
- Doing TCS research

Who is this class for?

- Anyone who might be interested in doing TCS research.
- Prime target audience member (?):
 First-year graduate student with some interest in CS theory.
- "Mathematical maturity" is a must.

Logistics

http://www.cs.cmu.edu/~15751 for links to homework, Piazza (sign up!), schedule

HOMEWORK will be the main form of evaluation.

Approximately 8, typically due/out on Mondays.

LaTeX solutions mandatory.

Homework 1 is released, due Sept. 19!

Grades: 80% homework

15% take-home final

5% participation

Part 2: How to present mathematics.

Writing math: LaTeX

You must become fluent in LaTeX.

If you are very unfamiliar with LaTeX:

- Start with Overleaf.com or ShareLaTeX.com
- Use Google & tex.stackexchange.com to find out how to do anything & everything
- Maybe get "Short Math Guide for LaTeX"

LaTeX workflow

Select a good text editor that understands you are writing in LaTeX.

Essential features....

- syntax highlighting
- hotkey to compile/display
- "synchronization"/"roundtripping"
- text/reference autocompletion
- block commenting

Produce **pdf** output (command: PDFLatex).

Nobody uses .dvi or .ps these days!

Popular LaTeX editors

Browser-based: (good for collaboration... or solo!)

Overleaf.com, ShareLaTeX.com

Windows: (typically also need MiKTeX)

WinEdt, TeXnicCenter

Mac:

Texpad, TeXShop

Linux:

Kile

All three:

- Texmaker, TeXstudio
- Sublime + LaTeXing (or LaTeXTools) plugin
- Emacs + AUCTeX + RefTeX

LaTeX workflow

Create a stub .tex file, a lifetime .sty file, and a lifetime .bib file.

LaTeX — stub .tex file example

```
\documentclass[11pt] {article}
\usepackage{odonnell}
\begin{document}
\title{}
\author{Ryan O'Donnell\thanks{odonnell@cs.cmu.edu}}
\date{\today}
\maketitle
%\begin{abstract}
%\end{abstract}
%\section{}
%\bibliographystyle{alpha}
%\bibliography{odonnell}
```

refers to odonnell.sty, my lifetime .sty file

refers to odonnell.bib, my lifetime .bib file

\end{document}

LaTeX — lifetime .sty file

```
\usepackage{fixltx2e,amsmath,amssymb,amsthm,amsfonts,bbm,graphicx,fullpage}
\usepackage[colorlinks,citecolor=blue,bookmarks=true]{hyperref}
\theoremstyle{plain}
\newtheorem{theorem}{Theorem}[section]
\newtheorem{lemma}[theorem]{Lemma}
\newtheorem{corollary}[theorem]{Corollary}
\newtheorem{proposition}[theorem]{Proposition}
\theoremstyle{definition}
\newtheorem{definition}[theorem]{Definition}
\newtheorem{remark}[theorem]{Remark}
% for ``commenting out'' chunks of text
\newcommand{\ignore}[1]{}
% for notes on the text
\newcommand{\ryansays}[1]{{\color{red}{\tiny [Ryan: #1]}}}
% macros
\newcommand{\R}{{\mathbbm R}}
\newcommand{\eps}{\epsilon}
```

Start yours today!

Use .bib file management software:

e.g., JabRef, Zotero, Mendeley, Papers

```
@STRING{proc = {Proceedings of the}}
@STRING{focs = {IEEE Symposium on Foundations of Computer Science}}
@STRING{focs12 = proc # { 53rd } # ann # { } # focs}
@STRING{focs13 = proc # { 54th } # ann # { } # focs}
@STRING{stoc = {ACM Symposium on Theory of Computing}}
@ARTICLE {AGHP92,
   author = {Alon, Noga and Goldreich, Oded and H{\aa}stad, Johan and Peralta, Ren{\'e}},
    title = \{\text{Simple constructions of almost } \{\$k\$\} - \text{wise independent random variables} \}
  journal = {Random Structures \& Algorithms},
   volume = \{3\},
      doi = \{10.1002/rsa.3240030308\},
      url = \{http://dx.doi.org/10.1002/rsa.3240030308\},\
@ONLINE{CLRS13.
      author = {Chan, Siu On and Lee, James and Raghavendra, Prasad and Steurer, David},
       title = {Approximate constraint satisfaction requires large {LP} relaxations},
        date = \{2013-09-03\},
        note = \{arXiv:1309.0563\},
      eprint = \{1309.0563\}
    author = {Mossel, Elchanan and R{\'a}cz, Mikl{\'o}s},
     title = {A quantitative {G}ibbard--{S}atterthwaite theorem without neutrality},
     pages = \{1041 - -1060\},
 publisher = {ACM},
       doi = \{10.1145/2213977.2214071\},
       url = \{http://dx.doi.org/10.1145/2213977.2214071\},
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Please make your .bib entries high quality!

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[AKK+08] Sanjeev Arora, Subhash Khot, Alexandra Kolla, David Steurer, Madhur Tulsiani, and Nisheeth K. Vishnoi, *Unique games on expanding constraint graphs are easy: extended abstract*, STOC (Richard E. Ladner and Cynthia Dwork, eds.), ACM, 2008, pp. 21–28.

[Alo86] Noga Alon, Eigenvalues and expanders, Combinatorica 6 (1986), no. 2, 83–96.

[AM85] N. Alon and V. D. Milman, λ_1 , isoperimetric inequalities for graphs, and superconcentrators, Journal of Combinatorial Theory. Series B **38** (1985), 73–88.

[AR98] Yonatan Aumann and Yuval Rabani, An O(log k) approximate min-cut max-flow theorem and approximation algorithm, SIAM Journal on Computing 27 (1998), no. 1, 291–301.

[ARV04] Sanjeev Arora, Satish Rao, and Umesh Vazirani, Expander flows, geometric embeddings and graph partitioning, Proceedings of the thirty-sixth annual ACM Symposium on Theory of Computing (STOC-04) (New York), ACM Press, June 13–15 2004, pp. 222–231.

[Che70] J. Cheeger, A lower bound on smallest eigenvalue of a laplacian, Problems in Analysis (Papers dedicated to Salomon Bochner) (1970), 195–199.

[DKSV06] Nikhil R. Devanur, Subhash Khot, Rishi Saket, and Nisheeth K. Vishnoi, *Integrality gaps for sparsest cut and minimum linear arrangement problems*, STOC, 2006, pp. 537–546.

[FS02] Feige and Schechtman, On the optimality of the random hyperplane rounding technique for MAX CUT, RSA: Random Structures Algorithms 20 (2002).

[GMPT07] Konstantinos Georgiou, Avner Magen, Toniann Pitassi, and Iannis Tourlakis, *Integrality gaps of* 2 - o(1) for vertex cover SDPs in the lovész-schrijver hierarchy, FOCS, IEEE Computer Society, 2007, pp. 702–712.

[Kho02] Subhash Khot, On the power of unique 2-prover 1-round games, STOC, 2002, pp. 767–775.

This is gross.
You may as well

have a spelling mistake in your title.

Please make your .bib entries high quality!

- Have a standard for citing proceedings (FOCS, STOC, etc.) and arXiv and ECCC.
- Get capitalization correct: {B}races needed
- Put in people's first and last names—with the diacritcs!
- Use math mode for math parts of titles

Where to get .bib entries:

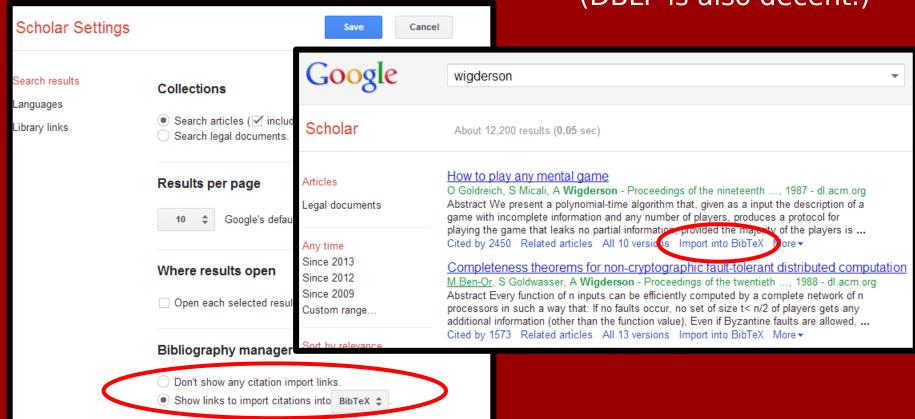
1. Always try www.ams.org/mrlookup first. (The best entries, and you don't need ams.org access.)

AMERICAN MATHEMATICAL SOCIETY						
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		{SIAM Journal on Computing},				
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		{9851005 (electronic)},				
		{0097-5397},				
		{60G50 (52A20 68W20)},				
		{2203735 (2007h:60041)},				
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DOI	=	{10.1137/S009753970544727X},				
URL	=	{http://dx.doi.org/10.1137/S009753970544727X},				
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l .						

Where to get .bib entries:

2. Failing that, scholar.google.com.

(DBLP is also decent.)



LaTeX — version control

Version control software a must: will save you heartache and help you collaborate with others.

Browser editors like Overleaf & ShareLaTeX have some built in.

Else, Dropbox or similar are the bare minimum.

If sophisticated & can convince your coauthors:
Git (bitbucket, SourceTree, gitobox)

DON'T

\$ < U, V > \$

"quotes"

$$$$ f(x) = x^2 $$$$

\$ log(1+x) \$

```
\[ (\frac{ax+b}{cy})^2 \]
```

<u>DO</u>

```
$ \langle U, V \rangle $
```

``quotes''

```
f(x) = x^2
```

```
\[
  \left(\frac{ax+b}{cy}\right)^2
\]
```

DON'T

\begin{eqnarray} y &=& (x+1)^2 \\ &=& x^2+2x+1 \end{eqnarray}

If A is a matrix, then

Lemma \ref{lem:big} is
due to Blum \cite{Blu99}

assuming- as we do - that the Birch-Swinnerton-Dyer Conjecture holds

<u>DO</u>

```
\begin{align}
  y &= (x+1)^2 \\
    &= x^2+2x+1
\end{align}
```

If \$A\$ is a matrix, then

Lemma~\ref{lem:big} is
due to Blum~\cite{Blu99}

assuming---as we do---that the Birch--Swinnerton-Dyer Conjecture holds

DON'T DO

```
one party, e.g. Alice, is one party, e.g.\ Alice, is \[ a-b=2 ~~~~~~ a+b=4 a-b=2 \qquad a+b=4
```

```
\] \]
```

we execute ALG(x) we execute textnormal(ALG)(x)

I could go on. When in doubt, look up the correct thing to do at tex.stackexchange.com!

Writing mathematics well

This is a challenging, lifelong skill.

If I had to give two pieces of advice...

1. This is math, so it has to be 100% correct.

that said,

2. Take pity on your poor reader; help them out.

LaTeX — drawing

DON'T BE LAZY: include figures to help the reader.

```
\usepackage{graphicx}
...
\includegraphics{mypicture.png}
```

Was that so hard?

Works with .jpg, .png, .pdf, .eps

To draw figs: *Inkscape, TikZ, ...* but there's a learning curve.

Recommendation: draw figs with your presentation software (PowerPoint, Keynote, ...), since you have to learn it anyway...

Presentation software

If you write a paper, you'll have to make a talk.

To make a talk, you'll need PowerPoint/Keynote/Beamer.

Any of these is fine, but you'll still suffer the

"drawing figures" challenge with Beamer.

It's not "hip", but **become a hacker** in one of these.

Learn to integrate **beautiful math equations**:

PowerPoint: Built-in Equation Editor is now good!

Alternative: IguanaTex

Keynote: LaTeXiT (I'm told)

Beamer: Automatic

Presenting math well



I like Kayvon Fatahalian's tips:

http://www.cs.cmu.edu/~kayvonf/misc/cleartalktips.pdf

Part 3: How to **do** mathematics/TCS.

Finding papers

Use Google Scholar.

- Use CMU credentials (VPN) to get journal issues online.
- Alternative: check the author's home page.
- Books/some older journal articles can be found in the actual physical science library, in Wean.
- For books, first use Google Books / Amazon's "Read Inside" feature to try to find what you want.
- There are other methods.
- All else fails: Interlibrary loan is not too slow (https://illiad.library.cmu.edu/illiad/illiad.dll)

Finding papers

If you look at a paper, even briefly:

- 1. Check its "cited by" link on Google Scholar
- 2. Save a local copy.

Maintain a giant folder of saved papers! Even better: Add to Jabref/Mendeley/Zotero...

Use a consistent naming convention.

E.g., nisan-wigderson-log-rank-conj.pdf
This will save you 100's of hours, lifetime.

How to find papers to read

Papers citing / cited by the paper you're reading

Proceedings of recent FOCS/STOC/SODA/CCC

Google Scholar Alerts

http://scirate.com

Stay au courant

Read TCS blogs: http://feedworld.net/toc

Follow your favorite researchers on Twitter

Attend talks at CMU:

Theory Lunch, every Wed. noon, GHC6115 ACO Seminar, every Thu. 3:30, Wean 8220

CMU's theory-announce mailing list:

https://mailman.srv.cs.cmu.edu/mailman/listinfo/theory-announce

CMU's TCS YouTube channel:

https://www.youtube.com/channel/UCWFp4UWNiOv71j0sPbdNiqw

Stay au courant

Watch videos of talks from elsewhere:

- https://sites.google.com/site/plustcs/
- https://video.ias.edu/csdm
- https://simons.berkeley.edu/videos
- http://www.birs.ca/videos/2016
- (e.g., Venkat's talk there from 2 days ago...)
- https://www.youtube.com/user/MicrosoftResearch
- http://theory.cs.princeton.edu/videos.html
- Many conferences (e.g., some STOC/FOCS)
 have videos of the talks available online
- Some researchers' home pages have links to their talks

Streetfighting Mathematics

Q: What is the next number in the series?

1, 2, 5, 20, 125, 1070, ???

A: Just look it up at oeis.org

(Online Encyclopedia of Integer Sequences)

1, 2, 5, 20, 125, 1070

Michael Somos, Nov 16 2002

Search

Hints

+20

(Greetings from The On-Line Encyclopedia of Integer Sequences!)

Search: seq:1,2,5,20,125,1070

Displaying 1-1 of 1 result found. page 1 Sort: relevance | references | number | modified | created | Format: long | short | data A076795 Partial sums of (2n-1)!!0, 1, 2, 5, 20, 125, 1070, 11465, 146600, 2173625, 36633050, 691362125, 14440672700, 330674815925, 8236528396550, 221694575073425, 6411977928702800, 198310761891213425, 6530970632654064050 (list; graph; refs; listen; history; text; internal format) OFFSET 0.3 LINKS Vincenzo Librandi, Table of n, a(n) for n = 0..300FORMULA E.g.f.: $exp(x) * Integral \{t=0, x\} exp(-t)/sqrt(1-2*t) dt$. a(n)=a(n-1)(2n-2)-a(n-2)(2n-3). $a(n) \sim 1/sqrt(2)/n * 2^n * (n/e)^n$. G.f.: A(x)=x/(1-x)*(1+x/(U(0)-x)), where U(k)=(2*k+1)*x+1-x+1(2*k+3)*x/U(k+1); (continued fraction Euler's 1st kind, 1-step). -Sergei N. Gladkovskii, Jun 27 2012 G.f.: x/(1-x)/Q(0), where Q(k)=1-x*(k+1)/Q(k+1); (continued fraction). -Sergei N. Gladkovskii, May 19 2013 G.f.: G(0)*x/(1-x), where G(k) = 1 - x*(k+1)/(x*(k+1) - 1/G(k+1)); (continued fraction). - Sergei N. Gladkovskii, Aug 04 2013 MATHEMATICA Join [{0}, Accumulate [Table [(2n-1)!!, {n, 0, 20}]]] (* Harvey P. Dale, Jan 27 2013 *) PROG (PARI) a(n)=if(n<0, 0, sum(k=0, n-1, (2*k)!/k!/2^k)) CROSSREFS Cf. A001147. KEYWORD nonn AUTHOR

Q: What are the Stirling numbers of the second kind?

What is the explicit formula for them?

A: Look it up on Wikipedia

Stirling numbers of the second kind

From Wikipedia, the free encyclopedia

In mathematics, particularly in combinatorics, a Stirling number of the second **kind** is the number of ways to partition a set of *n* objects into *k* non-empty subsets and is denoted by S(n,k) or ${n \choose k}$.[1] Stirling numbers of the second kind occur in the field of mathematics called combinatorics and the study of partitions.

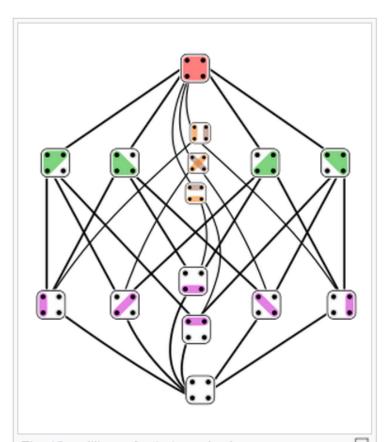
Stirling numbers of the second kind are one of two kinds of Stirling numbers, the other kind being called Stirling numbers of the first kind. Mutually inverse (finite or infinite) triangular matrices can be formed by arranging the Stirling numbers of the first respectively second kind according to the parameters n, k.

Contents [show]

Definition [edit source | edit beta]

The Stirling numbers of the second kind $\binom{n}{k}$ count the number of ways to partition a set of n labelled objects into k nonempty unlabelled subsets. Equivalently, they count the number of different equivalence relations with precisely k equivalence classes that can be defined on an n element set. Obviously,

$$\binom{n}{1} = \binom{n}{n} = 1.$$



The 15 partitions of a 4-element set ordered in a Hasse diagram

Q: What is 0.601907230197?

(This question based on a true story.)



A: Look it up at Inverse Symbolic Calculator



Standard lookup results for 0.601907230197

Best guess: BesK(1,1)

BesK(1,1)-exp(-Pi)^GAM(1/12)	6019072301972343
BesK(1,1)	6019072301972345
BesselK(1,1)	
Besk(T,T)+exp(-Pi)^GAM(1/12)	6019072301972347



Q: What is the Bessel K function?

A: Look it up on Wikipedia

An anecdote



Ryan Williams had an awesome CMU PhD thesis. I read the first draft. Its #1 theorem was:

Theorem: Any alg. for SAT using $n^{o(1)}$ space requires time $\widetilde{\Omega}(n^c)$, where c is the largest root of $c^3-c^2-2c+1=0$; i.e., $c\approx 1.801$.

I had my computer calculate a few more digits: $c \approx 1.801937736$.

Plugged it into Inverse Symbolic Calculator...



Standard lookup results for 1.801937736

Best guess: 1+2*x-x^2-x^3

rootofx^3-x^2-2*x+1;	1801937735804838
F(2/7,5/7;1/2;3/4)	
-3-3*x+6*x^2-6*x^3+3*x^5	
1+2*x-x^2-x^3	
cos(P#4/7)/ccs(P#4/8)	
cos(Pi*1/7)+cos(Pi*1/7)	

An anecdote

I let him know, and now his famous theorem reads:

Theorem: Any alg. for SAT using $n^{o(1)}$ space requires time $\widetilde{\Omega}(n^{2\cos(\pi/7)})$.

Q: Let K, L ⊆ Rⁿ be closed, bounded, convex sets with smooth boundary.
Does KuL have piecewise-smooth boundary?

A: Well, I didn't know, but it's the kind of question where you just **know** that some expert in analysis knows the answer.

Ask on mathoverflow.net.



Users

Baddes

Unanswered

If K and L are compact convex sets with smooth boundary, does their union have piecewise-smooth boundary?



Clarification: by "piecewise", I mean a finite number of pieces.



I'm sure this must be true, but my search for a citation was in vain (although I did learn the new term "polyconvex").



Thanks!

math**overflow**



add comment

start a bounty

2 Answers

active

votes



I don't think this is true. Suppose one of the sets is essentially $\{(x,y):y\geq x^2\}$ in the plane (cut off in some smooth way at the top, to make it compact). And suppose the other one is the same except that the parabolic lower boundary has been replaced by the graph of something like $y = x^2 + e^{-1/x^2} \sin(1/x)$ In other words add a fierce oscillation but damped so strongly that the region above the curve is still convex (i.e., d^2y/dx^2 remains positive). (I haven't done the arithmetic to make sure my e^{-1/x^2} damping is sufficient; if it isn't, then replace it by a more vigorous damping.) The union of the two convex sets will have infinitely many corners, where $\sin(1/x)$ is 0.

share edit flag

answered Nov 10 '10 at 17:16

oldest

Tagged dg.differential-geometry × 2347 convexity smoothness analytic-geometry integral-geometry Asked 2 Years Ago Viewed 293 Times Active 2 Years Ago

Ask Question

Stack Exchange

We're 5 years old! Share your story.

Community Bulletin

Who gets the reputation? Best of MathOverflow meta Long-term archiving of MathOverflow meta Two tags for partially ordered sets

Love this site?

Stackexchange sites

Mathoverflow.net:

For research-level questions about math.

CSTheory.stackexchange.com,

CS.stackexchange.com,

Two places for questions about TCS.

math.stackexchange.com:

For help with math questions at any level. (Do not post your homework here!!!)

tex.stackexchange.com:

For any questions about LaTeX.

Q: What's the 4^{th} -order Taylor series for arcsin(x)?

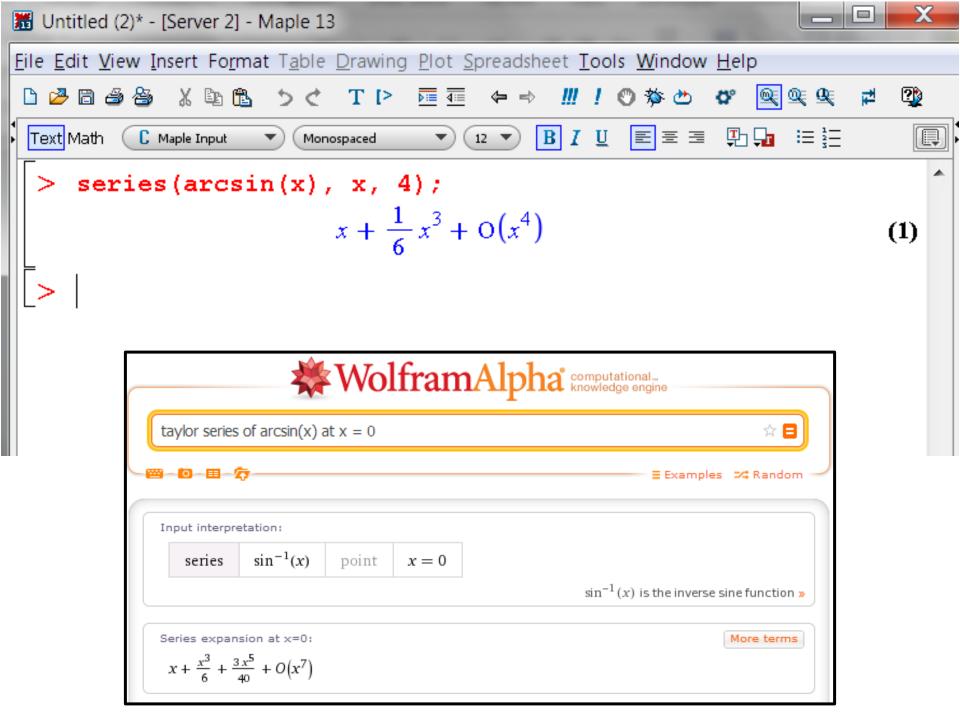
A: Ask Maple/Mathematica/Sage.

(The first two are equally awesome.

Sage is free, and is based around python.

Maple/Mathematica freely accessible at CMU.

For quick things, use wolframalpha.com.)



What else are Maple/Mathematica good for?

A: Everything. Use liberally.

Finding eigenvalues, SVDs Plotting functions Testing numerical conjectures Writing code Solving linear progs (symbolically, too) Gröbner bases Simplifying complicated expressions Finite field arithmetic Generating random numbers Solving differential equations **Explicit computations** Integrating (symbolically/numerically) Finding roots of equations Visualizing graphs Solving systems of equations Empirically checking inequalities Maximizing/minimizing expressions Quadratic programming Outputting LaTeX of expressions Curve fitting Asymptotics and Taylor series Inverting matrices (symbolically, too) Testing primality/irreducibility Differentiating

What else are Maple/Mathematica good for?

A: Everything. Use liberally.

Basically, if it's a math problem, and you think someone in history ever thought of using a computer to do it, then Maple/Mathematica can do it.

PS: You should also learn Matlab.

Often better for numerical things.

Streetfighting Mathematics

an example

Q: Suppose p(x) is a polynomial of degree ≤ k which is bounded in [-1,+1] for x ∈ [-1,+1].
 What is the largest p'(0) can be?

Remark: This question actually comes up from time to time in analysis of boolean functions.

You can probably solve it with judicious Googling.

Also appropriate for math.stackexchange.com, **if** you put in a reasonable effort first.

Let's solve it using streetfighting mathematics.

Q: Suppose p(x) is a polynomial of degree ≤ k which is bounded in [-1,+1] for x ∈ [-1,+1].
What is the largest p'(0) can be?

Let's think about
$$k = 3$$
, say, so

$$p(x) = a + bx + cx^2 + dx^3$$

For each value of x, e.g. x = .2, we have a **constraint**:

$$-1 \le a + .2b + .04c + .008d \le +1$$

We want to maximize b

Q: Suppose p(x which is bound what is the

We have infinitely many constraints, but probably not much changes if we just take some random 5000 of them.

Let's think about
$$k = 3$$
, say, so

$$p(x) = a + bx + cx^2 + dx^3$$

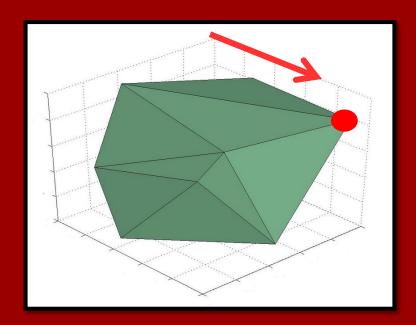
For each value of x, e.g. x = .2, we have a **constraint**:

$$-1 \le a + .2b + .04c + .008d \le +1$$

We want to maximize b

So say we have 10,000 linear inequalities over the variables a, b, c, d; they form some polytope in ${\bf R}^4$.

We want to maximize b.



This is a "Linear Program" (a topic we'll study).

Maple/Mathematica/Matlab can solve it.

```
> with(Optimization): with(RandomTools):
> deg := 3; numPts := 5000:
> P := x -> add(c[i]*x^i, i = 0..deg):

> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
> Constrs := map(y -> (P(y) >= -1, P(y) <= 1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
> BestPoly := subs(op(Soln[2]), P(x));
> plot(BestPoly, x = -1..1);
```

```
deg = 3
> P := x \rightarrow add(c[i] *x^i, i = 0..deq):
> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
> Constrs := map(y \rightarrow (P(y) >= -1, P(y) \leq -1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
                                                             BestDeriv := 3.00056057358036
> BestPoly := subs(op(Soln[2]), P(x));
> plot(BestPoly, x = -1..1);
                                                            -0.5
                                                                                   0.5
```

deg = 3: looks like maximizer is $p(x) = 3x - 4x^3$

```
deg := 1
> P := x \rightarrow add(c[i]*x^i, i = 0..deg):
> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
> Constrs := map(y \rightarrow (P(y) >= -1, P(y) \leq -1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
                                                              BestDeriv := 1.00041785758203
> BestPoly := subs(op(Soln[2]), P(x));
                                                 BestPoly := 0.000239437959177035980 + 1.00041785758202950 x
> plot(BestPoly, x = -1..1);
                                                                       0.5
                                                             -0.5
                                                                                     0.5
                                                                       -0.5
```

deg = 1: p(x) = 1x

```
deg = 2
> P := x \rightarrow add(c[i] *x^i, i = 0..deq):
> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
\geq Constrs := map(y \Rightarrow (P(y) \geq= -1, P(y) \leq= 1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
                                                               BestDeriv := 1.00154659899863
> BestPoly := subs(op(Soln[2]), P(x));
                                        BestPoly := 0.500012915154317117 + 1.00154659899863407 x - 0.501560741214689410 x^2
> plot(BestPoly, x = -1..1);
                                                                                       0.5
                                                                        -0.5
```

deg = 2: $p(x) = .5 + \mathbf{1}x - .5x^2$

```
deg := 3
> P := x \rightarrow add(c[i]*x^i, i = 0..deg):
> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
> Constrs := map(y \rightarrow (P(y) >= -1, P(y) \leq= 1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
                                                               BestDeriv := 3.00056057358036
> BestPoly := subs(op(Soln[2]), P(x));
                          BestPoly = -0.000426544117962384708 + 3.00056057358035799 x + 0.00170592447557217143 x^2 - 4.00223867084451790 x^3
> plot(BestPoly, x = -1..1);
                                                                        0.5
                                                              -0.5
                                                                                      0.5
```

(2)

deg = 3: $p(x) = 3x - 4x^3$

```
deg = 4
> P := x \rightarrow add(c[i]*x^i, i = 0..deg):
> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
> Constrs := map(y -> (P(y) >= -1, P(y) <= 1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
                                                              BestDeriv := 3.00028125652306
> BestPoly := subs(op(Soln[2]), P(x));
              BestPoly = 0.0000121482101641112857 + 3.00028125652306343 x - 0.000326369013471309970 x^2 - 4.00112504781797629 x^3 + 0.00111128812486230256 x^4
> plot(BestPoly, x = -1..1);
                                                                       0.5
                                                             -0.5
                                                                                     0.5
```

(2)

(3)

deg = 4: $p(x) = 3x-4x^3$ again (!)

```
(1)
                                                                        deg = 5
> P := x \rightarrow add(c[i] *x^i, i = 0..deg):
> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
> Constrs := map(y -> (P(y) >= -1, P(y) <= 1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
                                                               BestDeriv := 5.00120532339474
                                                                                                                                                          (2)
> BestPoly := subs(op(Soln[2]), P(x));
     BestPoly = -0.00106623559767311328 + 5.00120532339473822 x + 0.0127965087812789788 x^2 - 20.0144708379618130 x^3 - 0.0170841289591253838 x^4 + 16.0192856795011309 x^5
                                                                                                                                                          (3)
> plot(BestPoly, x = -1..1);
                                                                         0.5
                                                               -0.5
                                                                                       0.5
```

deg = 5: $p(x) = 5x - 20x^3 + 16x^5$

```
deg = 6
> P := x \rightarrow add(c[i]*x^i, i = 0..deg):
> ConstrPts := [seq(2*Generate(float(range = 0..1, method = uniform))-1, i=1..numPts)]:
\geq Constrs := map(y \Rightarrow (P(y) \Rightarrow -1, P(y) \leq 1), ConstrPts):
> Soln := LPSolve(c[1], Constrs, maximize):
> BestDeriv := Soln[1];
                                                                  BestDeriv := 5.00176146095726
> BestPoly := subs(op(Soln[2]), P(x));
BestPoly = 0.000368034916638261978 + 5.00176146095725738 x - 0.00321856542933066266 x^2 - 20.0210030986764060 x^3 - 0.00827062048655928528 x^4 + 16.0279861868147115 x^5
    + 0.0188502914009799974 x^6
> plot(BestPoly, x = -1..1);
                                                                            0.5 -
                                                                 -0.5
                                                                                           0.5
```

(2)

deg = 6: $p(x) = 5x - 20x^3 + 16x^5$ again

Summary:

Except for weird anomaly at degree 2, looks like degree 2k optimizer is the same as the degree 2k-1 optimizer.

(In fact, that's true; can you see why?)

So let's focus on odd degree.

$$p_1(x) = 1x$$

$$p_3(x) = 3x-4x^3$$

$$p_5(x) = 5x-20x^3+16x^5$$

Largest p'(0) seems to equal degree, but now what?

Summary:

Except for weird anomalooks like degree 2k opt as the degree 2k-1 opt

Try typing these coefficients into oeis.org

So let's focus on odd degree.

$$p_1(x) = 1x$$

$$p_3(x) = 3x-4x^3$$

$$p_5(x) = 5x-20x^3+16x^5$$

Largest p'(0) seems to equal degree, but now what?

1,3,-4,5,-20,16 Search Hints

(Greetings from The On-Line Encyclopedia of Integer Sequences!)

Search: seq:1,3,-4,5,-20,16

Displaying 1-1 of 1 result found.

page 1

Sort: relevance | references | number | modified | created | Formation | short | data

A084930 Triangle of coefficients of Chebyshev polynomials T {2n+1} (x).

+10

1, -3, 4, 5, -20, 16, -7, 56, -112, 51, 0, 102, 132, -576, 256, -11, 220, -1232, 2816, -2816, 1024, 13, -364, 2912, -9984, 16640, -13312, 4096, -15, 560, -6048, 28800, -70400, 92160, -61440, 16384, 17, -816, 11424, -71808, 239360, -452608, 487424, -278528, 65536, -19, 1140, -20064, 160512, -695552 (list; table; graph; refs; listen; history; text; internal format) OFFSET 0.2

COMMENTS

Comment from Herb Conn, HCR 83, Box 93, Custer, SD 57730, Jan 28 2005: "Letting x = 2 Cos 2A, we have the following trigonometric identities: "Sin 3A = 3*Sin A - 4*Sin^3 A

"Sin 5A = 5*Sin A - 20*Sin^3 A + 16*Sin^5 A

"Sin 7A = 7*Sin A - 56*Sin^3 A + 112*Sin^5 A - 64*Sin^7 A

"Sin 9A = 9*Sin A - 120*Sin^3 A + 432*Sin^5 A - 576*Sin^7 A + 256*Sin^9 A, etc."

Cayley (1876) states "Write sin u = x, then we have sin u = x, [...] sin $3u = 3x - 4x^3$, [...] sin $5u = 5x - 20x^3 + 16x^5$, [...]". Since $T_n(\cos(u)) = \cos(nu)$ for all integer n, $\sin(u) = \cos(u-pi/2)$, and $\sin(u+kpi) = (-1)^k \sin(u)$ it follows that $T_n(\sin(u)) = (-1)^n((n-1)/2)$ sin(nu) for all odd integer n. – Michael Somos, Jun 19 2012

REFERENCES

A. Cayley, On an Expression for 1 +- sin(2p+1)u in Terms of sin u, Messenger of Mathematics, 5 (1876), pp. 7-8 = Mathematical Papers Vol. 10, n. 630, pp. 1-2.

Theodore J. Rivlin, Chebyshev polynomials: from approximation theory to algebra and number theory, 2. ed., Wiley, New York, 1990. p. 37, eq. (1.96) and p. 4, eq.(1.10).

LINKS

Table of n, a(n) for n=0..49.

- M. Abramowitz and I. A. Stegun, eds., <u>Handbook of Mathematical Functions</u>, National Bureau of Standards, Applied Math. Series 55, Tenth Printing, 1972 [alternative scanned copy].
- M. Abramowitz and I. A. Stegun, eds., <u>Handbook of Mathematical Functions</u>, National Bureau of Standards Applied Math. Series 55, Tenth Printing, 1972, p. 795.

Index entries for sequences related to Chebyshev polynomials.

FORMULA

Alternate rows of A008310.



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~-Български Ö



Chebyshev polynomials

From Wikipedia, the free encyclopedia

Not to be confused with discrete Chebyshev polynomials.



It has been suggested that *Dickson polynomial* be merged into this article. (Discuss) *Proposed since September 2011*.

In mathematics the **Chebyshev polynomials**, named after Pafnuty Chebyshev,^[1] are a sequence of orthogonal polynomials which are related to de Moivre's formula and which can be defined recursively. One usually distinguishes between **Chebyshev polynomials of the first kind** which are denoted T_n and **Chebyshev polynomials of the second kind** which are denoted T_n . The letter T is used because of the alternative transliterations of the name *Chebyshev* as *Tchebycheff*, *Tchebyshev* (French) or *Tschebyschow* (German).

The Chebyshev polynomials T_n or U_n are polynomials of degree n and the sequence of Chebyshev polynomials of either kind composes a polynomial sequence.

Chebyshev polynomials are polynomials with the largest possible leading coemerch, but subject to the condition that their absolute value is bounded on the interval b. 1. They are also the extremal polynomials for many other properties.^[2]

Chebyshev polynomials are important in approximation theory because the roots of the Chebyshev polynomials of the first kind, which are also called Chebyshev nodes, are used as nodes in polynomial interpolation. The resulting interpolation polynomial minimizes the problem of Runge's phenomenon and provides an approximation that is close to the polynomial of best approximation to a continuous function under the maximum norm. This approximation leads directly to the method of Clenshaw–Curtis quadrature.

In the study of differential equations they arise as the solution to the Chebyshev differential equations

$$(1 - x^2)y'' - xy' + n^2y = 0$$

Notes [edit source | edit beta]

- ^ Chebyshev polynomials were first presented in: P. L. Chebyshev (1854) "Théorie des mécanismes connus sous le nom de parallélogrammes,"
 <u>Mémoires de Sarans en angers presentes à l'Académie de Saint-Pétersbourg</u>, vol. 7, pages 539–586.
- A Rivlin, Theodore J. The Chebyshev polynomials. Pure and Applied Mathematics. Wiley-Interscience [John Wiley & Sons], New York-London-Sydney, 1974. Chapter 2, "Extremal Properties", pp. 56--123.
- 3. A Jeroen Demeyer Diopriantine Cets Co., Polynomial Diago and Hilbortic Touth Problem for Fronting Problem for Fronting Problems (2007), p.70.
- 4. A a b c Boyd, John P. (2001). Chebyshev and Fourier Spectral Methods [1] (second ed.). Dover. ISBN 0-486-41183-4.
- 5. ^ Chebyshev Interpolation: An Interactive Tour 🗗

References [edit source | edit beta]

- Abramowitz, Milton; Stegun, Irene A., eds. (1965), "Chapter 22"
 ^๗, Handbook of Mathematical Functions with Formulas, Graphs, and Mathematical Tables, New York: Dover, p. 773, ISBN 978-0486612720, MR 0167642 ^๗.
- Dette, Holger (1995), A Note on Some Peculiar Nonlinear Extremal Phenomena of the Chebyshev Polynomials **4**, *Proceedings of the Edinburgh Mathematical Society* **38**, 343-355
- Eremenko, A.; Lempert, L. (1994), An Extremal Problem For Polynomials , Proceedings of the American Mathematical Society,
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- Suetin, P.K. (2001), "C/c021940" &, in Hazewinkel, Michiel, Encyclopedia of Mathematics, Springer, ISBN 978-1-55608-010-4

External links [edit source | edit beta]

- Module for Chebyshev Polynomials by John H. Mathews
- Numerical Computing with Functions: The Chebfun Project



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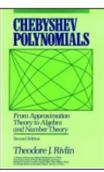
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Chebyshev polynomials: from approximation theory to algebra and number theory



Theodore J. Rivlin

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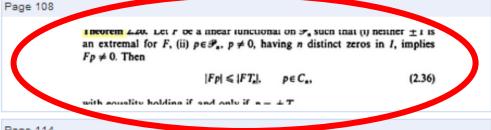
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Page 114

3 Bounding the approximate degree of the majority function

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We will require the following well-known fact, which follows from, e.g., [Riv90, Theorem 2.20].

Proposition 3.1. Let p(x) be a polynomial of degree at most k which satisfies $|p(x)| \le 1$ whenever $|x| \le 1$. Then $|p'(0)| \le 2\lceil \frac{k}{2} \rceil - 1 \le k$.

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References

[Riv90] Theodore J. Rivlin. Chebyshev Polynomials. John Wiley & Sons, New York, second edition, 1990.

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