## **KVM**

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# Kernel modules

#### Linux kernel modules

#### Module = kernel code that can be loaded/unloaded at runtime

- Allows to add/remove kernel features while system is running
- Modules have **full privileges** and control of the system → buggy modules may crash the kernel!
- Make it easy to develop drivers without rebooting
- Help keep kernel image size to a minimum
- Help reduce boot time: avoid spending time initializing devices and kernel features that will only be needed later
- Modules installed in /lib/modules/<kernel\_version>/kernel and have the .ko extension

### **Loading modules**

To load a single module without its dependencies:

```
sudo insmod <module_path>.ko
```

To load a module with its dependencies:

```
sudo modprobe <module_name>
```

- modprobe reads
  /lib/modules/<kernel\_version>/modules.dep.bin to
  determine:
  - each module's location (path)
  - each module's dependencies

### Module utilities

 To get information about a module (parameters, license, description, dependencies, etc.):

```
modinfo <module_name>
modinfo <module_path>.ko
```

• To display all loaded modules (see /proc/modules):

```
lsmod
```

■ To remove a module (and its depedencies with -r):

```
rmmod <module_name>
```

# **KVM**

### What is KVM?

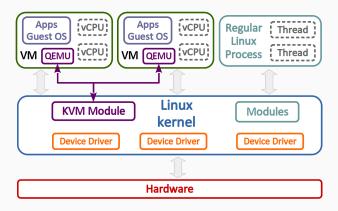
- KVM stands for Kernel based Virtual Machine
- Linux kernel module providing hardware-assisted virtualization
- Provide a virtualization API for hypervisors
- Requires Intel VT-x or AMD-V
- Originally, KVM virtualized only CPU and memory
  - devices (I/O) had to be emulated by QEMU
- Nowadays, KVM supports device virtualization (PIC and probably more)
- Being part of Linux, KVM is open-source software

#### **Evolution of KVM**

- Introduced to make VT-x/AMD-V available to user space
  - expose virtualization features securely
  - interface: /dev/kvm
- Quickly merged into Linux mainline
  - available since kernel 2.6.20 (2006)
  - from first LKML posting to kernel merge: only 3 months!
  - 7300 lines of C code! (as of Linux 5.8.12)
- Evolved significantly since 2006
  - ported to other architectures: RAM, s390, PowerPC, IA64
  - became recognized & driving part of Linux kernel
  - quick support of latest virtualization features

#### KVM model

- User processes can create VMs → VMs are just user space processes
- Virtual CPUs (vCPUs) mapped to kernel threads



### Architectural benefits of KVM model

- Proximity of guest and user space hypervisor
  - lacksquare Only one address space switch: guest  $\leftrightarrow$  host
  - Both run in user space  $\rightarrow$  lighter context switch
- Massive Linux kernel reuse
  - Memory management
  - Scheduler
  - I/O stacks, power management, host CPU hot-plugging, etc.
- Massive Linux user space reuse
  - Network configuration
  - Handling of VM images
  - Logging, tracing, debugging

# KSM (Kernel Samepage Merging)

- Linux kernel feature that deduplicates "identical" pages found across user processes
  - benefit: massive gain in memory usage!
  - drawback: security (theoretical)
- Spawn ksmd daemon which inspects pages for possible merges
- Kernel must be compiled with CONFIG\_KSM=y (> 2.6.32)
- KSM controled<sup>1</sup> by writing to /sys/kernel/mm/ksm/run:
  - 0: stop ksmd from running but keep merged pages
  - 1: run ksmd
  - 2: stop ksmd and unmerge all pages currently merged
- Article on KSM "Increasing memory density by using KSM"<sup>2</sup>

<sup>&</sup>lt;sup>1</sup>https://www.kernel.org/doc/Documentation/vm/ksm.txt

<sup>&</sup>lt;sup>2</sup>https://www.kernel.org/doc/ols/2009/ols2009-pages-19-28.pdf

#### Can a host run KVM?

• Check for hardware virtualization support (Intel or AMD):

```
$ lscpu|grep Flags|grep "vmx\|svm"
Flags: fpu vme de pse tsc msr pae mce cx8 apic sep
   mtrr pge mca cmov pat pse36 clflush dts acpi mmx
  fxsr sse sse2 ss ht tm pbe syscall nx pdpe1gb rdtscp
  lm constant_tsc arch_perfmon pebs bts rep_good nopl
   xtopology nonstop_tsc cpuid aperfmperf pni
  pclmulqdq dtes64 monitor ds_cpl vmx smx est tm2
  ssse3 sdbg fma cx16 xtpr pdcm ...
```

• Check the kvm module is loaded in the kernel:

```
$ lsmod|grep kvm
kvm_intel 282624 0
kvm 663552 1 kvm_intel
```

### **Accessing KVM**

- To access the KVM device, /dev/kvm, one must either (depends on the Linux distro):
  - be in the kvm group<sup>3</sup>
  - have the proper ACL permissions<sup>4</sup>
- Typical examples of KVM API use:
  - QEMU when launched with -enable-kvm
  - Any other Linux-based hypervisor using KVM
  - Any application using /dev/kvm, typically a custom VMM

 $<sup>^3</sup>$ https://linuxize.com/post/how-to-add-user-to-group-in-linux/

<sup>&</sup>lt;sup>4</sup>https://www.redhat.com/sysadmin/linux-access-control-lists

# **KVM API**

#### Overview

- Device /dev/kvm provides access to the KVM API
  - kvm module must be loaded!
- Requests performed through ioctl calls
- Provide 3 types of resources, accessed by file descriptors:
  - system (KVM): VM creation, memory mapping, etc.
  - VM: vCPU creation, interrupts, etc.
  - vCPU : access to registers, etc.

### Basic example

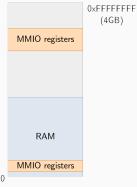
```
#include ux/kvm.h>
const uint8_t code[] = { // Write 0x42 at I/O adress 0x80:
 0x66, 0xba, 0x80, 0x00, // mov dx,0x80
 0xb0, 0x42, // mov al, 0x42
 0xee
                        // out dx, al
1:
int main(int argc, char **argv) {
   // Open KVM
   int kvmfd = open("/dev/kvm", O_RDWR | O_CLOEXEC);
   // Set up a virtual machine
   int vmfd = ioctl(kvmfd, KVM_CREATE_VM, (unsigned long)0);
   // Get some page-aligned memory and copy code to it
   void *m = mmap(0, 0x1000, PROT_READ|PROT_WRITE, MAP_SHARED|MAP_ANONYMOUS,
        -1, 0);
   memcpv(m. code, sizeof(code)):
   // Create a virtual CPU #0
   int vcpufd = ioctl(vmfd, KVM_CREATE_VCPU, (unsigned long)0));
   // Run the VCPU #0
   while (1) {
       ioctl(vcpufd, KVM_RUN, 0);
```

### **Devices: MMIO vs PMIO**

- Memory-Mapped I/O devices (MMIO)
  - device registers are mapped into the physical memory space
  - RAM and devices share the same address space!
  - read/write to/from these devices happen exactly like memory
  - all instructions dealing with memory operands can interact with these devices
    - e.g. mov instruction (x86)  $\rightarrow$  mov al, [42]
- Port-Mapped I/O devices (PMIO)
  - device registers are mapped into a specific memory space,
     distinct from the physical memory space
  - require **specific** instructions to access these devices
    - e.g. in and out instructions (x86)  $\rightarrow$  in al,42

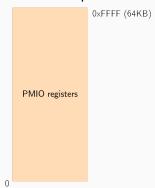
### **Device address spaces**

### Physical Memory Space



 MMIO address space, accessed using regular memory instructions (mov)

#### PMIO Address Space



 Distinct memory space, only accessed using specific assembly instructions (in, out)

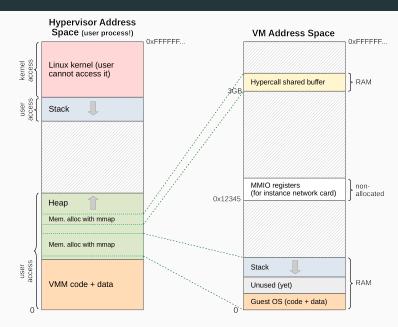
### KVM workflow from the VMM's perspective

- 1. Create a KVM device
- 2. Create a VM
- 3. Allocate RAM for the VM
- 4. Map allocated RAM into the VM's address space
- 5. Load guest OS binary blob into the VM's RAM
- 6. Create a vCPU and setup its registers
- Run the vCPU on the guest OS' code until a VMexit is triggered
  - handle VMexit to emulate the guest OS' expected behavior
  - resume vCPU execution

#### vCPU execution and VMexits

- vCPU execution happens at native speed, without interruption until guest OS code generates a VMexit
- A VMexit occurs when guest OS code:
  - reads from or writes to an I/O port (PMIO)
  - writes to a physical address that has no mapping (MMIO)
  - triggers some special operation (including errors)
- After a VMexit is handled, the VMM resumes the vCPU's execution

# VM memory mapping example



## (1) Create a KVM device

 To obtain a file descriptor on the kvm device and check the stable version of the API is available:

```
int kvmfd = open("/dev/kvm", O_RDWR | O_CLOEXEC);
if (kvmfd < 0) err(1, "%s", "/dev/kvm");

int version = ioctl(kvmfd, KVM_GET_API_VERSION, NULL);
if (version < 0) err(1, "KVM_GET_API_VERSION");
if (version != KVM_API_VERSION) err(1, "Unsupported version of the KVM API");</pre>
```

## (2) Create a VM

To obtain a file descriptor on a newly created VM:

```
int vmfd = ioctl(kvmfd, KVM_CREATE_VM, 0);
if (vmfd < 0) err(1, "KVM_CREATE_VM");</pre>
```

## (3) Allocate RAM for the VM

- Memory allocated for the guest is made out of pages
- Pages are 4KB in size and aligned to 4KB
- Must use mmap to allocate pages (not malloc!)

```
// Alloc 256KB for the guest
u_int page_count = 64;
u_int ram_size = 4096*page_count;
uint8_t *mem = mmap(NULL, ram_size, PROT_READ|
    PROT_WRITE, MAP_SHARED|MAP_ANONYMOUS, -1, 0);
if (!mem) err(1, "Allocating guest memory");
```

## (4) Map allocated RAM into the VM's address space

- Define where in the VM, the memory is physically mapped
- Below code maps mem at physical address 0 in the VM
- Each memory mapping (region in KVM lingo) must be associated to a different slot, here 0

```
struct kvm_userspace_memory_region memreg = {
    .slot = 0,
    .guest_phys_addr = 0,
    .memory_size = ram_size,
    .userspace_addr = (uint64_t)mem,
    .flags = 0
};
if (ioctl(vmfd, KVM_SET_USER_MEMORY_REGION, &memreg)
    < 0) err(1, "KVM_SET_USER_MEMORY_REGION");</pre>
```

## (5) Load guest OS into VM's RAM

- Guest OS must be "loaded" into the guest address space
- Simplest guest OS = mini bare-metal OS
- Simply perform the following operations:
  - 1. read the binary file generated from compiling/linking the guest OS' code + data
  - 2. read the file into the pages allocated for the guest RAM
  - 3. easiest is to load it at address 0 in guest physical memory
    - later on, we'll set the CPU instruction pointer to 0 as well

## (6) Create a vCPU

- A vCPU is referenced through a file descriptor, vcpufd below
- The vCPU is represented as a memory-mapped file
  - the memory-mapped area is a <a href="mailto:kvm\_run">kvm\_run</a> structure, <a href="mailto:run">run</a> below

## (6) Setup the vCPU registers (1/2)

• Initialize code segment register cs:

# (6) Setup the vCPU registers (2/2)

- Initialize instruction pointer rip to point to the beginning of the OS' code (0)
- Initialize stack pointer rsp to point to top of the RAM (ram\_size)
- Initialize flags register rflags (bit 1 is reserved: must be 1)

## (7) Run the vCPU

- Use the KVM\_RUN ioctl on the vCPU file descriptor to run it
- This ioctl blocks until the vCPU triggers a VMexit!

```
while (1)) {
 int status = ioctl(vcpufd, KVM_RUN, NULL);
 if (status < 0) {
      err(1, "VMM: KVM_RUN");
  switch (run->exit_reason) {
      case KVM_EXIT_IO:
      case KVM_EXIT_MMIO:
      case KVM_EXIT_HLT:
     case KVM_EXIT_SHUTDOWN:
      case KVM_EXIT_FAIL_ENTRY:
      case KVM_EXIT_INTERNAL_ERROR:
      default:
          fprintf(stderr, "KVM error");
```

## (8) Handle VMexits to emulate desired behavior

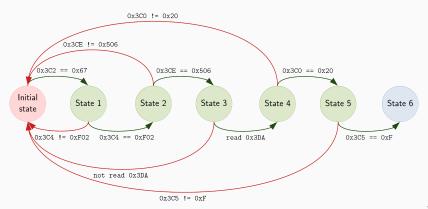
- There are many types of VMexits, but these 2 are especially interesting:
  - KVM\_EXIT\_IO: encountered when guest issues an I/O (in or out) machine instruction
  - KVM\_EXIT\_MMIO: encountered when guest writes to an address that has no mapping (i.e. no RAM)
- Note that KVM\_EXIT\_IO is significantly faster than KVM\_EXIT\_MMIO (according to KVM documentation)
- VMM must then emulate the desired behavior expected by the guest OS

#### **Device emulation**

- Emulating a real device allows a guest OS implementing a driver for the real hardware to use it
- OS implements drivers to support various physical devices, for instance:
  - VGA display, mouse, keyboard, SATA drive, etc.
- Device drivers write and read to specific device registers
  - either MMIO or PMIO or both
- How does an hypervisor emulate a device?

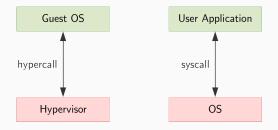
#### Device emulation: state machine

```
// Code excerpt of init. sequence for VGA X-mode 400x300
// This code uses PMIO registers only
outb(0x3C2, 0x67);
outw(0x3C4, 0x0F02); // enable writing to all planes
outw(0x3CE, 0x0506); // graphic mode
inb(0x3DA);
outb(0x3C0, 0x20); // enable video
outb(0x3C5, 0x0F);
```



## **Hypercalls**

- Mechanism for the guest to request the help of the hypervisor
- Similar to a system call between an application and an OS



- Benefits
  - much simpler drivers → no need to emulate the real hardware!
  - much better performance!

## Hypercalls: simple implementation example

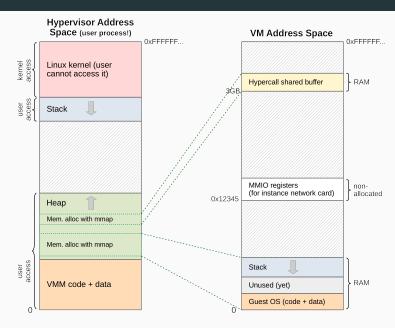
#### Guest

- write hypercall arguments in shared memory (buffer) between guest ↔ hypervisor
- write hypercall request to an unused port, e.g. OxABBA

### Hypervisor

- when KVM\_EXIT\_IO encountered, check whether it's an hypercall
- if hypercall, then extract its arguments and emulate the expected behavior
  - possibly write an output to the shared buffer

## Hypercalls: shared buffer



# Gracefully exit the VM (1/2)

### Case 1: guest OS explicitely stops the CPU

- Guest: executes the hlt machine instruction to stop the CPU
- Hypervisor: hlt triggers the KVM\_EXIT\_HLT VMexit
- Do not forget to deallocate the VM's resources!

## Gracefully exit the VM (2/2)

### Case 2: execution of the guest OS never stops (infinite loop)

- → KVM\_RUN ioctl never stops either!
- How to handle a graceful exit initiated on the host?
  - run the vCPU in a dedicated thread
  - configure the thread to be asynchronously interruptible with:

```
pthread_setcanceltype(PTHREAD_CANCEL_ASYNCHRONOUS, NULL);
```

when the VM must be stopped, cancel the vCPU thread with:

```
pthread_cancel(vcpu_thread);
```

Do not forget to deallocate the VM's resources!

## Cleanup the VM's data structures

- Using munmap, free all memory regions allocated with mmap
- Close the KVM device using the file descriptor previously returned by open(/dev/kvm, ...)

#### Resources

- KVM http://linux-kvm.org
- KVM API reference https://www.kernel.org/doc/html/latest/virt/kvm/api.html
- Using the KVM API https://lwn.net/Articles/658511/
- Kernel module (Arch Linux documentation) https://wiki.archlinux.org/index.php/Kernel\_module