

1. (a) Proof: $1^\circ \|\vec{x}\|_\infty = \max_{i \in [1, n]} |x_i| \geq 0$ ✓

You need to justify $\|x\|_\infty = 0 \Leftrightarrow x = 0$.

$2^\circ \forall \vec{x} \in \mathbb{R}^n, \alpha \in \mathbb{R}$.

$$\|\alpha \vec{x}\|_\infty = \max_{i \in [1, n]} |\alpha x_i| = \max_{i \in [1, n]} |\alpha| \cdot |x_i| = |\alpha| \cdot \max_{i \in [1, n]} |x_i| = |\alpha| \cdot \|\vec{x}\|_\infty$$
 ✓

$3^\circ \|\vec{x} + \vec{y}\|_\infty = \max_{i \in [1, n]} |x_i + y_i| \leq \max_{i \in [1, n]} |x_i| + \max_{j \in [1, n]} |y_j| = \|\vec{x}\|_\infty + \|\vec{y}\|_\infty$ ✓

Over all, $\|\vec{x}\|_\infty = \max_{i \in [1, n]} |x_i|$ is indeed a norm on \mathbb{R}^n .

(b) Proof: As $\forall i \in [1, n] (i \in \mathbb{N}^+)$, $0 \leq |x_i| \leq \max_{i \in [1, n]} |x_i|$

and as $\exists j \in [1, n] (j \in \mathbb{N}^+)$, $x_j = \max_{i \in [1, n]} |x_i|$ ✓

so $\max_{i \in [1, n]} |x_i| \leq \sum_{i=1}^n |x_i| \leq \sum_{j=1}^n \left(\max_{i \in [1, n]} |x_i| \right) = n \cdot \max_{i \in [1, n]} |x_i|$

i.e. $\|\vec{x}\|_\infty \leq \|\vec{x}\|_1 \leq n \cdot \|\vec{x}\|_\infty$

(c) Proof:

$$\begin{aligned} |\vec{x}^T \vec{y}| &= \left| \sum_{i=1}^n x_i y_i \right| \leq \sum_{i=1}^n |x_i y_i| \leq \sum_{i=1}^n \left(|x_i| \cdot \max_{j \in [1, n]} |y_j| \right) \\ &= \left(\sum_{i=1}^n |x_i| \right) \cdot \max_{j \in [1, n]} |y_j| \\ &= \|\vec{x}\|_1 \cdot \|\vec{y}\|_\infty \end{aligned}$$
 ✓

2. (a) Proof: (I denote the reason "as $\|A\vec{x}\|_2$ is the 2-norm on \mathbb{R}^m " as "①")

$1^\circ \|A\|_2 = \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} \|A\vec{x}\|_2 \stackrel{\text{①}}{\geq} 0$ ✓

You need to justify $\|A\|_2 = 0 \Leftrightarrow A = 0$.

$2^\circ \|\alpha A\|_2 = \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} \|(\alpha A)\vec{x}\|_2 = \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} \|\alpha(A\vec{x})\|_2$

$\stackrel{\text{①}}{=} \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} \alpha \cdot \|A\vec{x}\|_2$

$= \alpha \cdot \|A\|_2$ ✓

3° Assume $B^{m \times n}$,

thus $\|A+B\|_2 = \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} \|(A+B)\vec{x}\|_2 = \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} \|A\vec{x} + B\vec{x}\|_2$

$\stackrel{\text{①}}{\leq} \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} (\|A\vec{x}\|_2 + \|B\vec{x}\|_2)$

$\leq \max_{\vec{x} \in \mathbb{R}^n, \|\vec{x}\|_2=1} \|A\vec{x}\|_2 + \max_{\vec{y} \in \mathbb{R}^n, \|\vec{y}\|_2=1} \|B\vec{y}\|_2$

$= \|A\|_2 + \|B\|_2$ ✓

Over all, $\|\cdot\|_2$ is indeed a norm on $\mathbb{R}^{m \times n}$.

(b) Proof: $\forall \vec{x} \in \mathbb{R}^n$, $\vec{x} = \|\vec{x}\|_2 \cdot \vec{\alpha}$ where $\|\vec{\alpha}\|_2 = 1$ and $\|\vec{x}\|_2 \in \mathbb{R}$, $\|\vec{x}\|_2 \geq 0$
 thus, $\|A\vec{x}\|_2 = \|A \cdot \|\vec{x}\|_2 \vec{\alpha}\|_2 = \|\|\vec{x}\|_2 \cdot (A\vec{\alpha})\|_2 = \|\vec{x}\|_2 \cdot \|A\vec{\alpha}\|_2$
 $\leq \|\vec{x}\|_2 \cdot \max_{\vec{y} \in \mathbb{R}^n, \|\vec{y}\|_2=1} \|A\vec{y}\|_2$
 $= \|\vec{x}\|_2 \cdot \|A\|_2$

(c) Proof: $\|AB\|_2 = \max_{\vec{x} \in \mathbb{R}^p, \|\vec{x}\|_2=1} \|(AB)\vec{x}\|_2 = \max_{\vec{x} \in \mathbb{R}^p, \|\vec{x}\|_2=1} \|A(B\vec{x})\|_2$
 $\leq \|A\|_2 \cdot \|B\vec{x}\|_2$, $\vec{x} \in \mathbb{R}^p$ and $\|\vec{x}\|_2=1$ (as we proved in (b).)
 $\leq \|A\|_2 \cdot \|B\|_2 \cdot \|\vec{x}\|_2$
 $= \|A\|_2 \cdot \|B\|_2$

3. Proof:

Part 1: Firstly prove: $\|A\|_2 = \max_{\|\vec{x}\|_2=1} \|A\vec{x}\|_2 = \sqrt{\lambda_{\max}(A^T A)}$

Given a real symmetric matrix $P^{n \times n}$, there must be an orthogonal matrix O , s.t.

$$O^T P O = O^{-1} P O = \Lambda = \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{bmatrix}, \text{ where } \lambda_i \text{ is a eigenvalue of } P \text{ and columns of } O \text{ are eigenvectors of corresponding eigenvalues.}$$

Thus, we can prove:

$$\begin{aligned} \max_{\|\vec{x}\|_2=1} \vec{x}^T P \vec{x} & \xrightarrow[\text{thus } \vec{x} = O\vec{y} \text{ and } \|\vec{y}\|_2=1]{\text{let } \vec{y} = O^T \vec{x} = \begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix}} \max_{\|\vec{y}\|_2=1} \vec{y}^T (O^T P O) \vec{y} \\ &= \max_{\|\vec{y}\|_2=1} \vec{y}^T \Lambda \vec{y} \\ &= \max_{\|\vec{y}\|_2=1} \sum_{i=1}^n (\lambda_i \cdot y_i^2) \\ &\leq \lambda_{\max} \cdot \max_{\|\vec{y}\|_2=1} \sum_{i=1}^n y_i^2 \\ &= \lambda_{\max} \quad (\lambda_{\max} \text{ is the largest eigenvalue of } P) \end{aligned}$$

Therefore, when \vec{x} is an eigenvector of λ_{\max} , $\max_{\|\vec{x}\|_2=1} \vec{x}^T P \vec{x} = \lambda_{\max}$.

As $A^T A$ is a real symmetric matrix from $\mathbb{R}^{n \times n}$,

$$\begin{aligned} \sqrt{\lambda_{\max}(A^T A)} &= \sqrt{\max_{\|\vec{x}\|_2=1} \vec{x}^T (A^T A) \vec{x}} = \sqrt{\max_{\|\vec{x}\|_2=1} (A\vec{x})^T (A\vec{x})} = \sqrt{\max_{\|\vec{x}\|_2=1} \|A\vec{x}\|_2^2} \\ &= \max_{\|\vec{x}\|_2=1} \|A\vec{x}\|_2 = \|A\|_2 \end{aligned}$$

Part 2: Through the calculation of $A^T A$, it's obvious that

$$\|A\|_F = \sqrt{\text{tr}(A^T A)}$$

$$\text{thus } \|A\|_F = \sqrt{\text{tr}(A^T A)} = \sqrt{\sum_{i=1}^n \lambda_i}, \text{ where } \lambda_i \text{ is an eigenvalue of } A^T A.$$

$$\text{since } \lambda_{\max}(A^T A) \leq \sum_{i=1}^n \lambda_i \leq n \cdot \lambda_{\max}(A^T A)$$

$$\text{thus } \sqrt{\lambda_{\max}(A^T A)} \leq \sqrt{\sum_{i=1}^n \lambda_i} \leq \sqrt{n \cdot \lambda_{\max}(A^T A)}$$

$$\text{i.e. } \|A\|_2 \leq \|A\|_F \leq \sqrt{n} \cdot \|A\|_2$$

4. (a) $M_n \vec{d} = \left[\begin{array}{c} a_n \\ \vdots \\ a_n \end{array} \right] \}$ n -dim col-vector

$= a_n \cdot \vec{d}$, ✓ $d^T M_n$? -2

thus a_n is an eigenvalue of M_n and \vec{d} is an eigenvector of $\lambda = a_n$.

(b) Assume $M_n = [m_{ij}]$, $M_n^2 = [c_{ij}]$

thus $c_{ij} = \sum_{k=1}^n m_{ik} m_{kj}$

for any row of M_n^2 , the sum of the row is:

$$\sum_{j=1}^n \left(\sum_{k=1}^n m_{ik} m_{kj} \right) = \sum_{k=1}^n \left(m_{ik} \cdot \sum_{j=1}^n m_{kj} \right)$$

$$= \sum_{k=1}^n (m_{ik} \cdot a_n)$$

$$= a_n \cdot a_n = a_n^2$$
 ✓

Similarly, for any column of M_n^2 , the sum of the column is

$$\sum_{i=1}^n \left(\sum_{k=1}^n m_{ik} m_{kj} \right) = \sum_{k=1}^n \left(m_{kj} \cdot \sum_{i=1}^n m_{ik} \right)$$

$$= \sum_{k=1}^n (m_{kj} \cdot a_n) = a_n^2$$

Proved.

(c) Proof: According to Schur decomposition, $\forall A \in \mathbb{R}^{n \times n}$, A can be decomposed as $A = Q H Q^T$,

where Q is an orthogonal matrix and H is a triangular matrix.

As $M_n \in \mathbb{R}^{n \times n}$, thus we assume $M_n = Q H Q^T = Q H Q^{-1}$,

where $Q^{-1} = Q^T$ and H is triangular.

so $H = Q^T M_n Q$ ⁽¹⁾, we assume that $H = [h_{ij}]$, $Q = [q_{ij}]$, $M_n = [m_{ij}]$,
 $Q^T M_n = [z_{ij}]$

by calculating (1), we have $h_{ij} = \sum_{k=1}^n z_{ik} \cdot q_{kj}$

$$= \sum_{k=1}^n \left[q_{kj} \cdot \left(\sum_{l=1}^n q_{li} \cdot m_{kl} \right) \right]$$

$$= \sum_{l=1}^n \sum_{k=1}^n (q_{lj} \cdot q_{li} \cdot m_{kl})$$

$$\begin{aligned}
\text{thus, } |h_{ij}| &= \left| \sum_{l=1}^n \sum_{k=1}^n (q_{lj} \cdot q_{ki} \cdot m_{kl}) \right| \\
&\leq \sum_l \sum_k (m_{kl} \cdot |q_{lj} \cdot q_{ki}|) \quad (m_{kl} > 0) \\
&\leq \frac{1}{2} \sum_l \sum_k [m_{kl} \cdot (q_{lj}^2 + q_{ki}^2)] \\
&= \frac{1}{2} \left[\left(\sum_l q_{lj}^2 \cdot \sum_k m_{kl} \right) + \left(\sum_k q_{ki}^2 \cdot \sum_l m_{kl} \right) \right] \\
&= \frac{1}{2} \left[\left(\sum_l q_{lj}^2 \cdot a_n \right) + \left(\sum_k q_{ki}^2 \cdot a_n \right) \right] \quad \text{i.e. } |h_{ij}| \leq a_n \\
&= \frac{a_n}{2} \left(\sum_l q_{lj}^2 + \sum_k q_{ki}^2 \right) = \frac{a_n}{2} \times (1+1) = a_n \quad \text{--- denoted as (2)}
\end{aligned}$$

Because $M_n \sim H$, so M_n and H share the same eigenvalues, and as H is triangular, diagonal entries of H are eigenvalues of H .

So diagonal entries of H are eigenvalues of M_n .

As we have proved in (2), all entries of $H \leq a_n$, and as we showed in problem (a), a_n is an eigenvalue of M_n , thus we can infer that a_n is the largest eigenvalue of M_n .

As we proved in question 3, we have $\|M_n\|_2 = \sqrt{\lambda_{\max}(M_n^T M_n)}$

$\xleftarrow{\text{largest eigenvalue}} = \sigma_{\max}(M_n) \quad \text{why?}$
 $= a_n$

5. Proof: Order the m numbers from smallest to largest, and denote them as $a_{s_1}, a_{s_2}, \dots, a_{s_m}$.

\downarrow \downarrow
 smallest largest

Part 1: If we assume the number which minimizes this summation is in the range of $(-\infty, a_{s_1}) \cup (a_{s_m}, +\infty)$, and denote the summation in this situation as $d \Rightarrow$ Then we can always find that the 2 endpoints of the interval $[a_{s_1}, a_{s_m}]$, i.e. a_{s_1} & a_{s_m} , can make the summation smaller than d .

Thus, we can deny the assumption above and infer that the number we try to find is in the range of $[a_{s_1}, a_{s_m}]$.

Part 2: When the target number is in $[a_{s_1}, a_{s_m}]$, the summation can be expressed as:

$$\begin{aligned}
\sum_{i=1}^m |a_i - b| &= \sum_{i=1}^m |a_{s_i} - b| = (|a_{s_1} - b| + |a_{s_m} - b|) + \sum_{i=2}^{m-1} |a_{s_i} - b| \\
&= (a_{s_m} - a_{s_1}) + \sum_{i=2}^{m-1} |a_{s_i} - b|, \quad (b \in [a_{s_1}, a_{s_m}])
\end{aligned}$$

where $(a_m - a_1)$ is constant and the optimization problem is changed into minimizing the summation: $\sum_{i=2}^{m-1} |a_i - b|$

Part 3:

Exactly the same as the discussion above, we can infer that the target number is not in the range of $[a_1, a_2) \cup (a_{m-1}, a_m]$, i.e. we narrow the potential interval into $[a_2, a_{m-1}]$.

If we do this procedure iteratively, we can finally determine that the number we want is located in the range denoted as T :

$$T = \begin{cases} [a_{\lfloor \frac{m}{2} \rfloor}, a_{\lfloor \frac{m}{2} \rfloor + 1}], & \text{if } m \text{ is even} \\ [a_{\lfloor \frac{m}{2} \rfloor}, a_{\lfloor \frac{m}{2} \rfloor + 1}], & \text{if } m \text{ is odd} \end{cases}$$

Part 4:

1° if m is even,

$\forall b \in T$, b can minimize $\sum_{i=1}^m |a_i - b|$ as $\sum_{i=1}^{\frac{m}{2}} (a_{s, m+1-i} - a_{s_i})$

let $b = \frac{a_{s, \frac{m}{2}} + a_{s, \frac{m}{2} + 1}}{2}$ i.e. b is the median.

2° if m is odd,

$$\forall b \in T, \sum_{i=1}^m |a_i - b| = \underbrace{\sum_{i=1}^{\lfloor \frac{m}{2} \rfloor} (a_{s, m+1-i} - a_{s_i})}_{\text{constant}} + |a_{s, \lfloor \frac{m}{2} \rfloor + 1} - b|$$

obviously if we let $b = a_{s, \lfloor \frac{m}{2} \rfloor + 1}$ i.e. b is the median, then the summation is minimized.

Over all, a median of a_1, a_2, \dots, a_m minimizes $\sum_{i=1}^m |a_i - b|$ over all $b \in \mathbb{R}$.

6. (a) Because in the K-means algorithm, we calculate \vec{z}_j using

$$\begin{pmatrix} z_{j1} \\ z_{j2} \\ \vdots \\ z_{jn} \end{pmatrix} = \frac{1}{|G_j|} \cdot \sum_{i \in G_j} \begin{pmatrix} x_{i1} \\ x_{i2} \\ \vdots \\ x_{in} \end{pmatrix}$$

since $x_{i1}, x_{i2}, \dots, x_{in}$ are all nonnegative and $|G_j| \geq 1$,

We can know $z_{j1}, z_{j2}, \dots, z_{jn}$ are all nonnegative, i.e. all \vec{z}_j are also nonnegative.

(b) Suppose $\vec{x}_i = \begin{pmatrix} x_{i1} \\ x_{i2} \\ \vdots \\ x_{in} \end{pmatrix}$ where $\sum_{k=1}^n x_{ik} = 1$

$$\text{since } \vec{z}_j = \frac{1}{|G_j|} \sum_{i \in G_j} \vec{x}_i = \begin{pmatrix} z_{j1} \\ z_{j2} \\ \vdots \\ z_{jn} \end{pmatrix} \checkmark$$

$$\text{thus } z_{jk} = \frac{1}{|G_j|} \sum_{i \in G_j} x_{ik} \quad (k = 1, 2, \dots, n) \checkmark$$

$$\begin{aligned} \text{thus } \sum_{k=1}^n z_{jk} &= \sum_{k=1}^n \left(\frac{1}{|G_j|} \sum_{i \in G_j} x_{ik} \right) = \frac{1}{|G_j|} \left(\sum_{i \in G_j} \sum_{k=1}^n x_{ik} \right) \\ &= \frac{1}{|G_j|} \left(\sum_{i \in G_j} 1 \right) \\ &= \frac{1}{|G_j|} \cdot |G_j| = 1. \end{aligned} \checkmark$$

As we have explained in problem (a), because \vec{x}_i are nonnegative, \vec{z}_j are also nonnegative.

In conclusion, all \vec{z}_j are also proportions.

(c) $(\vec{z}_j)_i$ means the proportion of vectors which has a value of 1 in i -th entry
in all vectors of Group j . ✓