



# Data Structure & Algorithms

*Trainer: Nilesh Ghule*



# Course Introduction

- ✓ ① Backtracking
- ✓ ② Greedy method
- ✓ ③ Dynamic programming
- ✓ ④ Divide & Conquer

## • Data Structure and Algorithms

- Data Structures: Linked list, Stack, Queue, Binary search tree, Heap, Graph. Hash Table
- Algorithms: Sorting, Searching, Stack, Queue & Linked list applications, Graph algorithms.

## • Course Goals

- Implement each DS & Algorithms from scratch.
- Understand complexity of algorithms. ✓

## • Course Schedules 15 days \* 4 hrs

- 16<sup>th</sup> Feb 2023 to 4<sup>th</sup> Mar 2023 (\*) sth/6th (reserved)
- Mon-Sat: Lecture – ~~5~~:00 PM to ~~8~~:00 PM

## • Resource sharing 4 pm 8 pm

- <https://github.com/nilesh-g/dsa-06> ✓
- Recorded videos will be available for 7 days.  
<http://students.sunbeamapps.org>

## • Course Format

- Participants are encouraged to code alongside (copy code from code-sharing utility in student portal).
- Post your queries in chat box (on logical end of each topic).
- Practice assignments will be shared. They are optional. If any doubts, share on WA group (possibly with screenshot). Faculty members or peers can help.

## • Programming language

- DS & Algorithms are language independent.
- Classroom coding will be in Java (use IDE of your choice).
- Will share C++/Python codes at the end of session.
- Language pre-requisites?



# Course Pre-requisites

## Java

*vars/const  
loops  
if, switch*

- Language Funda
- Methods
- Class & Object
- static members
- Arrays
- Collections

*L ArrayList  
& HashMap*

## Python

- Language Funda
- Functions
- Class & Object
- Collections

## C++

- Language Funda
- Functions
- Class & Object
- Friend class
- Arrays
- Pointers

## C ✓

- Language Funda
- Functions
- Structures
- Arrays
- Pointers



# Data Structure

- Data Structure

- Organizing data in memory
- Processing the data

} *efficiently*

- Common data structures

- Array ✓
- Linked List ✓
- Stack ✓
- Queue ✓
- Hash Table ✓

- Advanced data structures

- Tree ✓
- Heap ✓
- Graph ✓



# Data Structure

- Data Structure
  - Organizing data in memory
  - Processing the data
- Common data structures
  - Array
  - Linked List
  - Stack
  - Queue
  - HashTable
- Advanced data structures
  - Tree
  - Heap
  - Graph

## • Asymptotic analysis

- It is not exact analysis
- Big' O notation

approx measure w.r.t. few params.

time analysis: msec or usec  
speed/clock, cores, ...

space analysis: bytes  
architecture

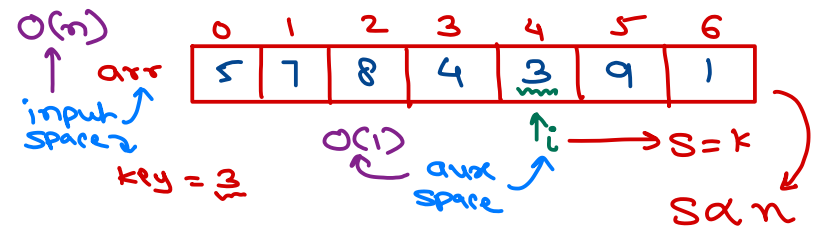
gcc, vs, arm - 4  
x86 gcc  
TC, avr - 2  
gcc  
int

## • Space complexity

- Unit space to store the data (Input space) and additional space to process the data (Auxiliary space).

- $O(1)$ ,  $O(n)$ ,  $O(n^2)$

$S=k$   $S \propto n$   $S \propto n^2$



## • Time complexity

- Unit time required to complete any algorithm.
- Approximate measure of time required to complete any algorithm.
- Depends on loops in the algorithm.
- $O(n^3)$ ,  $O(n^2)$ ,  $O(n \log n)$ ,  $O(n)$ ,  $O(\log n)$ ,  $O(1)$

num of iterations

$T=k$   
 $T \propto \log n$   
 $T \propto n$

# Time complexity

- Write a program to calculate factorial of given number.  $(n)$

```
res = 1;
for (i = 1; i <= n; i++)
    res = res * i;
print(res);
```

$T \propto n$   
 $\downarrow$   
 $O(n)$

- Print 2-D matrix of  $n \times n$ .

```
for (i = 0; i < n; i++) {
    for (j = 0; j < n; j++) {
        print(arr[i][j]);
    }
}
```

$\leftarrow n$  times  
 $\leftarrow n$  itrs

total itrs  
 $n \times n$   
 $T \propto n^2$   
 $\downarrow$   
 $O(n^2)$

- Print given number into binary.

$11 = (1011)_2$

11	
5	1
2	1
1	0
0	1

$\uparrow$  3

```
while (n > 0) {
    print(n % 2);
    n = n / 2;
}
```

$\boxed{2^{\text{itr}} = n}$

$\log_2 \text{itr} = \log_2 n$   
 $\text{itr} \log_2 2 = \log_2 n$   
 $\text{itr} = \frac{\log_2 n}{\log_2 2}$

- Print table of given number.  $(n)$

```
for (i = 1; i <= 10; i++)
    print(n * i);
```

$T = k$   
 $\downarrow$   
 $O(1)$

# Space complexity

aux space

$\text{res}$   $i$

$\boxed{\phantom{0000}}$   $\boxed{\phantom{0000}}$

$S = k \rightarrow O(1)$

$i$   $j$

$\boxed{\phantom{0000}}$   $\boxed{\phantom{0000}}$

$S = k \rightarrow O(1)$

$T \propto \text{itrs}$   
 $T \propto \frac{\log n}{\log 2}$

$S = k \rightarrow O(1)$

$T \propto \log n$   
 $O(\log n)$

$i$

$\boxed{\phantom{0000}}$

$S = k \rightarrow O(1)$



# Linear Search

key = 77 → found at index 5

- Find a number in a list of given numbers (random order).

arr

0	1	2	3	4	5	6	7	8
88	33	66	99	11	77	22	55	11

i

- Time complexity
  - Worst case
  - Best case
  - Average case



# Linear Search

- Find a number in a list of given numbers (random order).

key = 70 → not found

-1	0	1	2	3	4	5	6	7	8
arr	88	33	66	99	11	77	22	55	11

- Time complexity
  - Worst case
  - Best case
  - Average case

```
int linSearch (arr, key) {  
    for (i=0; i<n; i++) {  
        if (arr[i] == key) {  
            return i;  
        }  
    }  
    return -1;  
}
```





# Linear Search

- Find a number in a list of given numbers (random order).

-1	0	1	2	3	4	5	6	7	8
arr	88	33	66	99	11	77	22	55	11

if element is repeated, the first occurrence index will return.

```
int linSearch(arr, key) {  
    for(i=0; i<n; i++) {  
        if(arr[i] == key) {  
            return i;  
        }  
    }  
    return -1;  
}
```

- Time complexity

- Worst case

key = 70 → n itrs  
↓  
max itrs  
 $T \propto n$   
 $O(n)$

- Best case

key = 88 → single itr  
↓  
minimum itrs  
 $T = k$   
 $O(1)$

- Average case

random element find  
 $T \propto \frac{n}{2}$   
 $T \propto n$   
 $O(n)$



# Binary Search

key = 88

✓ Possible only if array is already sorted.

```
int binSearch(arr, key) {
```

```
    l = 0;
```

```
    r = n - 1;
```

```
    while (l <= r) {
```

```
        m = (l + r) / 2;
```

```
        if (key == arr[m])
```

```
            return m;
```

```
        if (key > arr[m])
```

```
            l = m + 1;
```

```
        else
```

```
            r = m - 1;
```

```
    }  
    return -1;
```

<u>0</u>	1	2	3	4	5	6	7	<u>8</u>
11	22	33	44	55	66	77	88	99

l r  
m

$$2^{\text{itr}} \approx n$$

$$\log 2^{\text{itr}} = \log n$$

$$\text{itr} \log 2 = \log n$$

$$\text{itr} = \frac{\log n}{\log 2}$$

$$T \propto \frac{\log n}{\log 2}$$

$$T \propto \log n \rightarrow O(\log n)$$

best case  
Finding middle  
single itr  
 $T = k \rightarrow O(1)$

avg/worst case  
 $\rightarrow$  not found.



*Thank you!*

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