



# Data Structure & Algorithms

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# Stack and Queue

- Stack & Queue are utility data structures.
- Can be implemented using array or linked lists.
- Usually time complexity of stack & queue operations is  $O(1)$ .
- Stack is Last-In-First-Out structure.
- Stack operations
  - push()
  - pop()
  - peek()
  - isEmpty()
  - isFull()\*

- Simple queue is First-In-First-Out structure.
- Queue operations
  - push()
  - pop()
  - peek()
  - isEmpty()
  - isFull()\*
- Queue types
  - Linear queue ✓
  - Circular queue ✓
  - Deque ✓
  - Priority queue ✓



# Linear Queue

rear = push new element  
front = pop an element

<del>11</del>	<del>22</del>	<del>33</del>	44	55	66
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$f$   $r$

arr[6]  
front -  
rear -

peek:  
return arr[f+1];

init:  
 $f = -1;$   
 $r = -1;$

is empty:  
 $f == r$

push:  
 $r++;$   
 $arr[r] = val;$

pop:  
 $f++;$   
 ~~$arr[f] = 0;$~~

is full:  
 $r == \text{MAX} - 1$

-1

0	1	2	3	4	5

$f = r$

-1

0	1	2	3	4	5
<del>10</del>	<del>20</del>	<del>30</del>	<del>40</del>		

$f < r$

-1

0	1	2	3	4	5
10	20	30	40	50	60

$r$



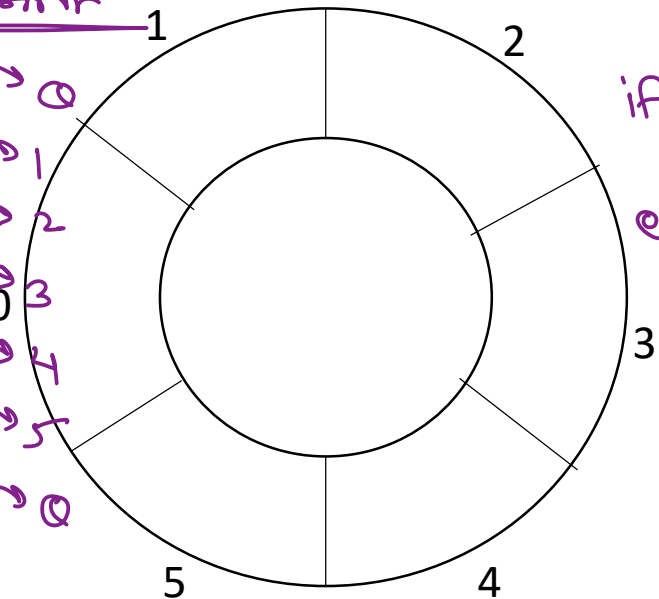
# Circular Queue

- In linear queue (using array) when rear reaches last index, further elements cannot be added, even if space is available due to deletion of elements from front. Thus space utilization is poor.
- Circular queue allows adding elements at the start of array if rear reaches last index and space is free at the start of the array.
- Thus rear and front can be incremented in circular fashion i.e. 0, 1, 2, 3, ..., n-1. So they are said to be circular queue.  $0, 1, 2, \dots$
- However queue full and empty conditions become tricky.



Circular increment of rear & front

$(-1 + 1) \% 6 \rightarrow 0$   
 $(0 + 1) \% 6 \rightarrow 1$   
 $(1 + 1) \% 6 \rightarrow 2$   
 $(2 + 1) \% 6 \rightarrow 3$   
 $(3 + 1) \% 6 \rightarrow 4$   
 $(4 + 1) \% 6 \rightarrow 5$   
 $(5 + 1) \% 6 \rightarrow 0$



if( $r == \text{max} - 1$ )  
     $r = 0$ ;  
else  
     $r++$ ;



# Circular Queue

init:  
 $r = -1;$   
 $f = -1$

peek:  
 $i = (f + 1) \% \text{size};$   
 $\text{return arr}[i];$

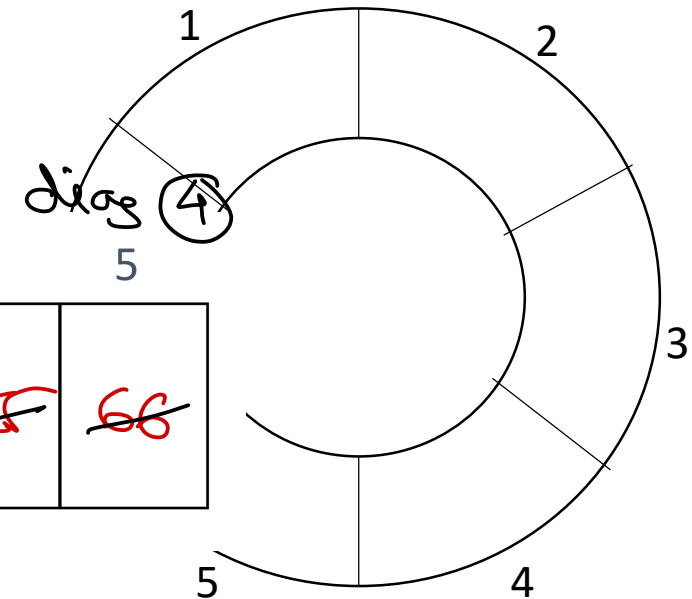
push:  
 $r = (r + 1) \% \text{size};$   
 $\text{arr}[r] = \text{val};$

pop:  
 $f = (f + 1) \% \text{size};$   
 $\text{if } (f == r) \{$   
     $f = -1;$   
     $r = -1;$   
 $\}$

isEmpty:  
 $(r == f) \ \&\& \ (r == -1)$



$r$   $f$



# Circular Queue

diag ②



f

r

diag ③



f

queue full.

$(r == f) \ \&\& \ (r != -1)$

diag ①



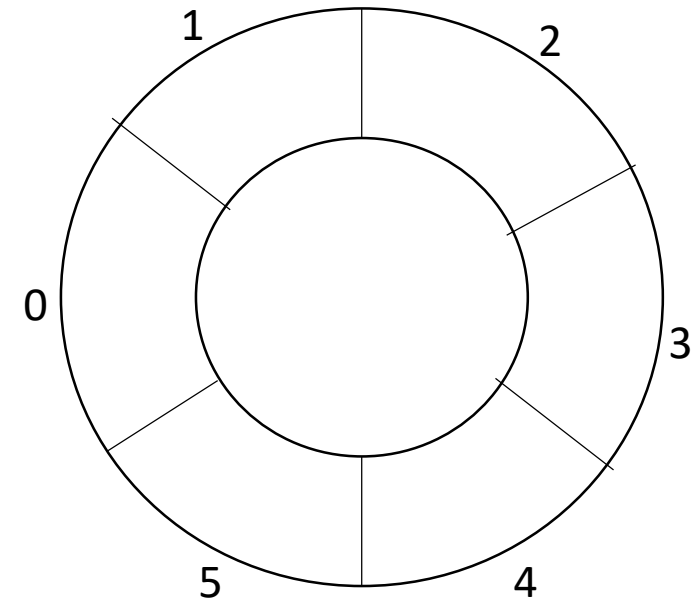
f

queue full.

r

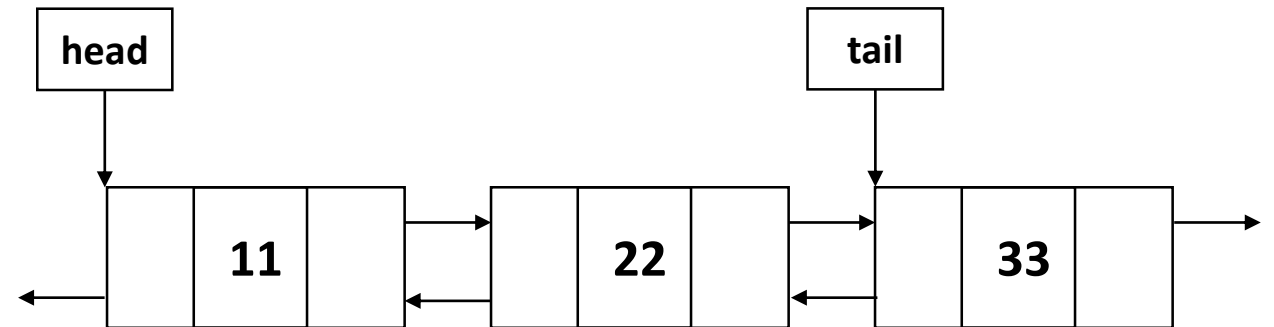
$(f == -1) \ \&\& \ (r == \text{size} - 1)$

OR



# DeQueue

- In double ended queue, values can be added or deleted from front end or rear end.



# Priority queue

- In priority queue, element with highest priority is removed first.

Doesn't follow FIFO.

Typically push/pop time complexity is not  $O(1)$ .

Efficient way of implementing priority queue  
is heap data structures - time complexity is  $O(\log n)$ .





# Stack

init:

$top = -1;$

push:

$top++;$

$arr[top] = val;$

pop:

$top--;$

peek:

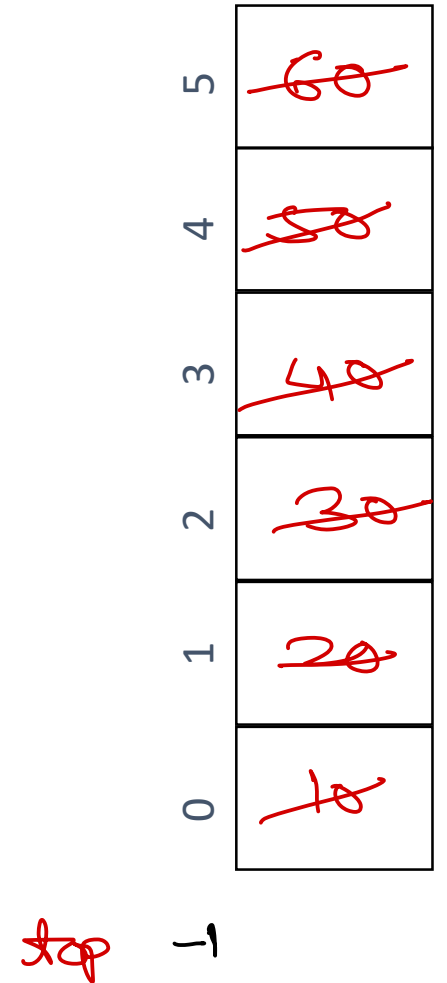
$return arr[top];$

is full:

$top == size - 1$

is empty:

$top == -1$



# Stack / Queue in Java collections

- class java.util.Stack<E>

- E push(E);
- E pop();
- E peek();
- boolean isEmpty();

- interface java.util.Queue<E>

- boolean offer(E e);
- E poll();
- E peek();
- boolean isEmpty();



## Expression Notations

$A + B \rightarrow$  infix

$+ A B \rightarrow$  prefix

$A B + \rightarrow$  postfix

## Operator Priorities

$( ) \leftarrow$  Highest

$\$$

$* /$

$+ - \leftarrow$  lowest

$$\textcircled{7} \quad \textcircled{6} \quad \textcircled{8} \quad \textcircled{3} \quad \textcircled{1} \quad \textcircled{2} \quad \textcircled{4} \quad \textcircled{9} \quad \textcircled{5}$$
$$A + B * C - (D + E * F / G - H) + I \$ J$$

$$A + B * C - (D + \underline{E F *} / G - H) + I \$ J$$

$$A + B * C - (D + \underline{E F * G /} - H) + I \$ J$$

$$A + B * C - (\underline{D E F * G / +} - H) + I \$ J$$

$$A + B * C - \underline{D F F * G / + H -} + I \$ J$$

$$A + \underline{B * C} - \underline{D F F * G / + H -} + \underline{I J \$}$$

$$\underline{A + B C *} - \underline{D F F * G / + H -} + \underline{I J \$}$$

$$\underline{A B C * +} - \underline{D F F * G / + H -} + \underline{I J \$}$$

$$\underline{A B C * + D F F * G / + H -} - + \underline{I J \$}$$

$$\underline{A B C * + D F F * G / + H - - I J \$ +}$$

$$\underline{+ - + A * B C - + D / * E F G H \$ I J}$$


To solve any expr  
programmatically.

step 1: Convert given  
infix into postfix or  
prefix.

step 2: solve that  
postfix or prefix.



# Infix to Postfix

•  $5 + 9 - 4 * (8 - 6 / 2) + 1 * (7 - 3)$  

$5 9 + 4 8 6 2 / - * - 1 7 3 - * +$

- ① traverse infix expr left to right.
- ② if symbol is operand, then append to postfix.
- ③ if symbol is operator, push it on stack.
  - ↳ if priority of topmost operator  $\geq$  priority of cur operator  
pop it from stack & append to postfix.
- ④ when all symbols from infix are completed, pop all operators from stack and append to postfix.
- ⑤ if cur sym is  $($ , push on stack.
- ⑥ if cur sym is  $)$ , pop all operators from stack and append to postfix until  $($  is found.  
Also pop and discard  $($ .




# Infix to Prefix

•  $5 + 9 - 4 * ( 8 - 6 / 2 ) + 1 \% ( 7 - 3 )$






*Thank you!*

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