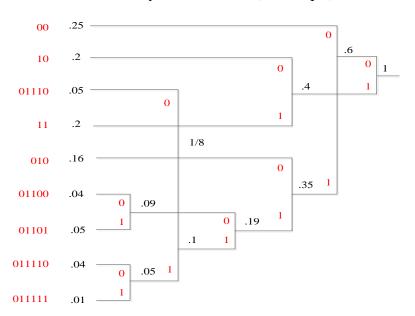
## Problem Set 4

## **Suggested Solutions**

- 1. Let  $X = \{x_1, x_2, x_3\}$  be a ternary random variable with probability distribution  $\{0.5, 0.4, 0.1\}$ .
  - a) Construct a binary Huffman code for X. Calculate the code's expected length  $\overline{L}_1$  and code efficiency  $\eta_1 = H(X)/\overline{L}_1$ .
  - b) Construct a binary Huffman code for two i.i.d. copies  $X^2$  of X, calculate the code's expected length  $\bar{L}_2$  and code efficiency  $\eta_2 = \frac{H(X^2)}{\bar{L}_2} = \frac{2H(X)}{\bar{L}_2}$ .
  - c) Make a comparison between  $\eta_1$  and  $\eta_2$ . If a binary Huffman code is used for n i.i.d. copies  $X^n$ , what is the asymptotic value of  $\eta_n = \frac{H(X^n)}{\overline{L}_n} = \frac{nH(X)}{\overline{L}_n}$  when  $n \to \infty$ ?
  - **Sol.** a) Binary Huffman code:  $\{0, 10, 11\}$ , expected length  $\overline{L}_1 = 1*.5+2*.4+2*.1 = 1.5$ , code efficiency  $\eta_1 = H(0.5, 0.4, 0.1)/1.5 = 1.361/1.5 = .9073$ .
  - b) The probability distribution of  $X^2$  is (.25, .2, .05, .2, .16, .04, .05, .04, .01). The construction of a binary Huffman code (not unique) is



 $\bar{L}_2$  = .5+.4+.25+.4+.48+.2+.25+.24+.06 = 2.78, code efficiency  $\eta_2$  = 2\*H(0.5, 0.4, 0.1)/2.78 = 2.722/2.78 = .9791.

- c)  $\eta_n$  tends to 1 with  $n \to \infty$ .
- 2. Consider 4 different codes:

```
{000, 10, 00, 11};
{100, 101, 0, 11};
{01, 100, 011, 00, 111, 1010, 1011, 1101};
{01, 111, 011, 00, 010, 110}
```

- a) Which ones do not satisfy Kraft inequality?
- b) Is each of these codes qualified as a prefix code? If not, please explain.

**Sol.** a) 3<sup>rd</sup> one does not satisfy Kraft inequality.

- b) 1<sup>st</sup> (including 000 and 00), 3<sup>rd</sup> (including 01 and 011), 4<sup>th</sup> (including 01 and 011) are not prefix codes.
- 3. For the following three codes, which ones cannot be constructed by Huffman coding?
  - a)  $\{0, 10, 11\};$
  - b) {00, 01, 10, 110};
  - c)  $\{01, 10\};$
- **Sol.** b) and c) cannot be constructed by Huffman coding, because code b) can be reduced to another uniquely decodable code {00,01,10,11}, which has a shorter expected length, and c) can be reduced to another uniquely decodable code {0, 1}.

1}, and 
$$p_{Y|X} = \begin{bmatrix} 1/2 & 1/4 & 1/4 \\ 1/4 & 1/4 & 1/2 \end{bmatrix}$$
. Assume  $P(X=0) = p_0$ ,  $P(X=1) = 1 - p_0$ .

- a) Calculate H(Y|X);
- b) Find the distribution of Y;
- c) What is the value of  $p_0$  that maximizes H(Y)?
- d) What is the capacity C of this channel?

 $= 5/4 - (3/4)\log 3 = 0.0613.$ 

$$\begin{aligned} &\textbf{Sol. a)} \ H(Y|X) = p_0 H(Y \mid X=0) + (1-p_0) H(Y \mid X=1) \\ &= p_0 H(1/2, 1/4, 1/4) + (1-p_0) H(1/4, 1/4, 1/2) = H(1/2, 1/4, 1/4) = 3/2. \\ &b) \ p(Y=0) = (1/2) p_0 + (1/4) (1-p_0) = (1/4) (p_0+1), \\ &p(Y=1) = (1/4) p_0 + (1/4) (1-p_0) = 1/4, \ p(Y=2) = (1/4) ((1-p_0)+1). \\ &c) \ When \ p(Y=0) = p(Y=2), \ i.e., \ when \ p_0 = 1/2, \ H(Y) \ is \ maximized. \\ &d) \ C = max_{p0} \ I(X;Y) = max_{p0} H(Y) - H(Y \mid X) \\ &= H(3/8, 1/4, 3/8) - 3/2 = 2*(3/8) log(8/3) + (1/4) log4 - 3/2 \end{aligned}$$