

# **CSE 1201**

## **Data Structure**

**Week 2**

**Lecture 3**

### **Chapter 2:**

### **Arrays, Pointers and Records**

# Introduction

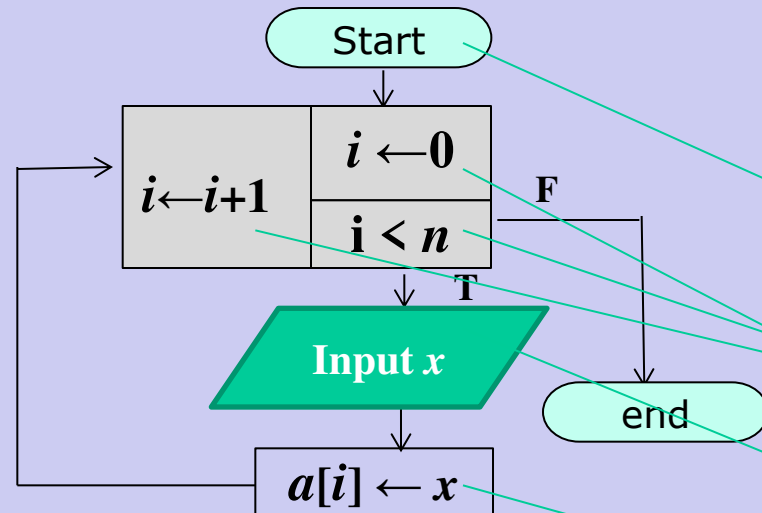
- Arrays
  - Structures of related data items
  - Static entity (same size throughout program)
- A few types
  - Pointer-based arrays (C-like)
  - Arrays as objects (C++)

# Introduction

- Array
  - Consecutive group of memory locations
  - Same name and type (**int**, **char**, etc.)
- To refer to an element
  - Specify array name and position number (index)
  - Format: arrayname[ position number ]
  - First element at position 1 (0 for C)
- N-element array **c**
  - $$\mathbf{c[0]}, \mathbf{c[1]} \dots \mathbf{c[n-1]}$$
  - Nth element as position N-1

# Array Creation

## Topic 1: Write an Algorithm to insert $n$ elements in an array



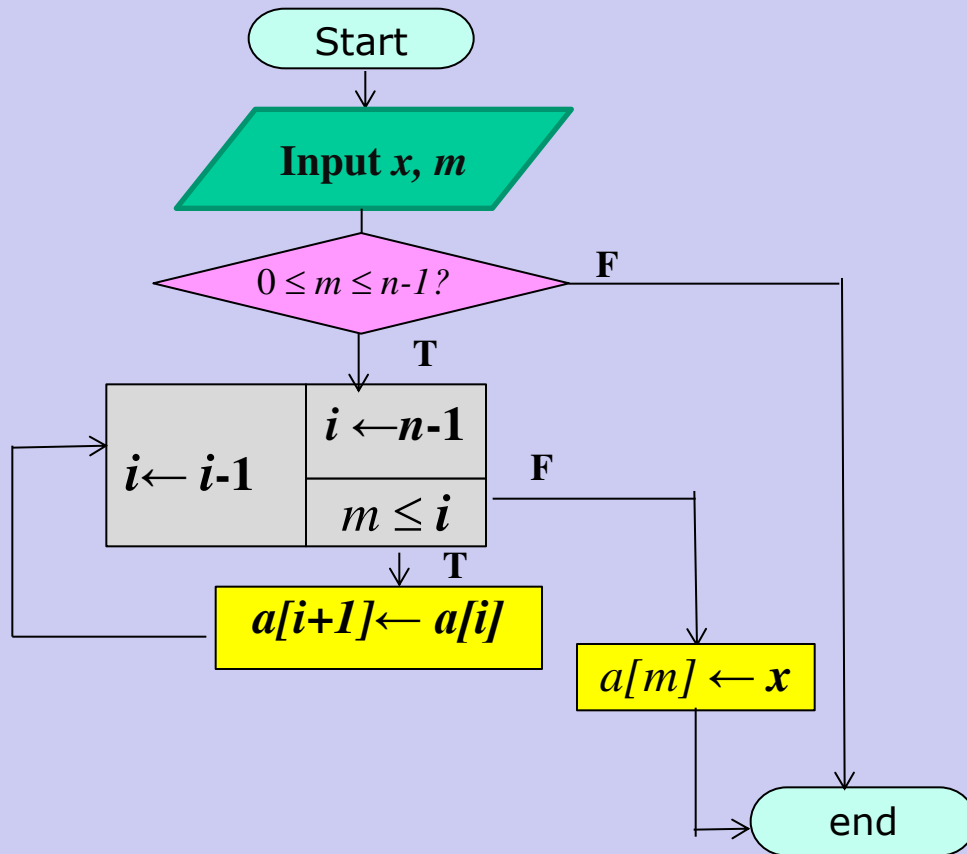
$n$ : total elements  
 $x$ : input variable  
Array Elements:  $a[0] \dots a[n-1]$

```
int main(){
    int i, a[100], n;
    n=10;
    for(i=0;i<n;i++){
        scanf("%d",&a[i]);
        a[i]=x;
    }
}
```

$a[]$	10	32	45		
	0	1	2	3	....

# Array Insertion

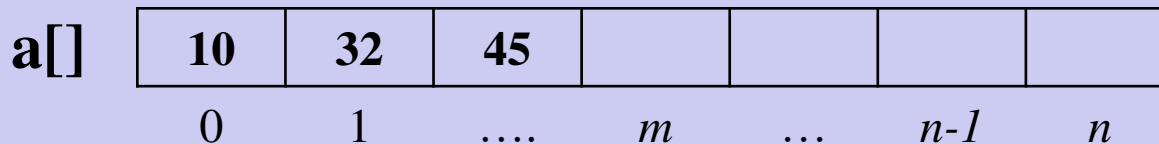
Topic : Insert a new element at index  $m$ .



$n$ : total elements  
 $m$ : index  $0 \leq m \leq n-1$   
 $x$ : input variable for new data  
Array Elements:  $a[0] \dots a[n-1]$

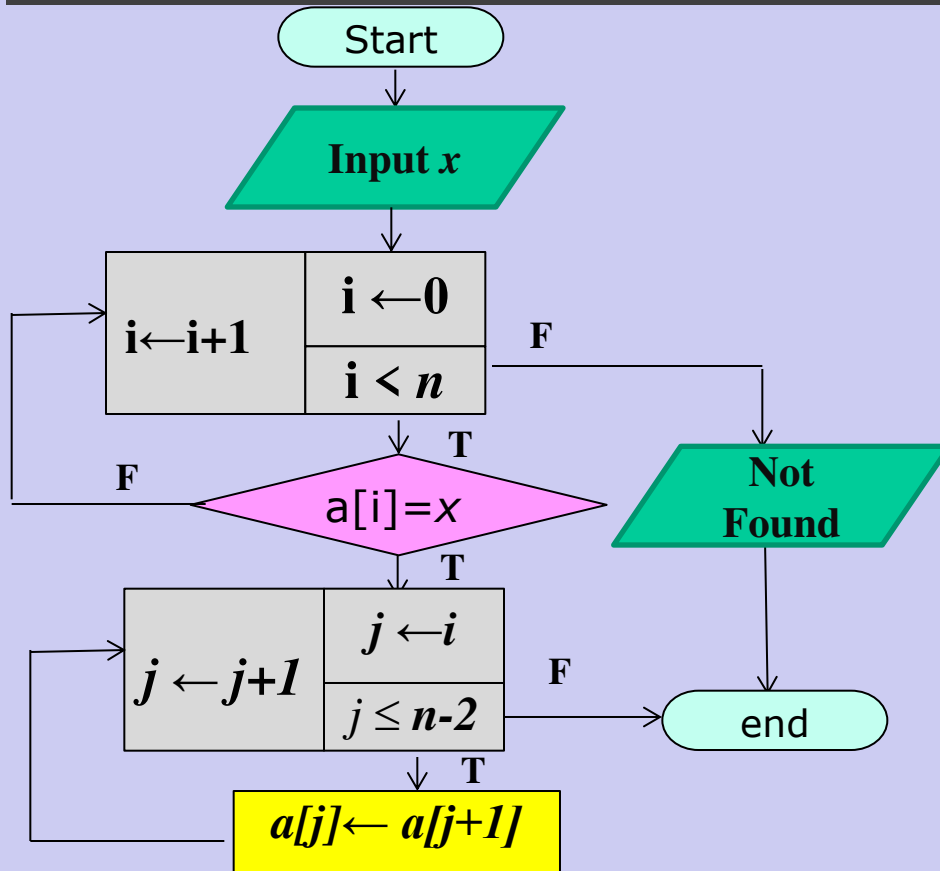
## Shifting Required

all elements from index  $m$  to  $(n-1)$  are needed to be shifted to index  $(m+1)$  to  $n$  respectively. Total  $(n-m)$  elements to be shifted.



# Array Deletion

Topic : Delete an specific element  $x$



$n$ : total elements

$m$ : index  $0 \leq m \leq n$

$x$ : input variable to be deleted

Array Elements:  $a[0] \dots a[n-1]$

## Shifting Required

all elements from index  $(m+1)$  to  $(n-1)$  are needed to be shifted from index  $m$  to  $(n-2)$  respectively. Total  $(n-m-1)$  elements to be shifted.

$a[]$	10	32	45				
	0	1	....	$m$	...	$n-1$	$n$

# Bubble Sort

# Sorting

- **Sorting takes an unordered collection and makes it an ordered one.**

1	2	3	4	5	6
77	42	35	12	101	5



1	2	3	4	5	6
5	12	35	42	77	101



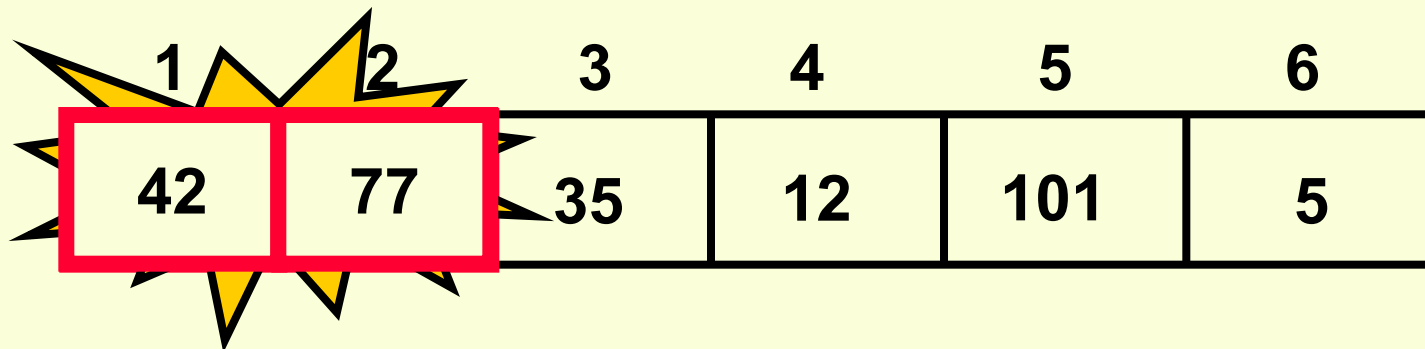
# "Bubbling Up" the Largest Element

- Traverse a collection of elements
  - Move from the front to the end
  - “Bubble” the **largest value** to the end using **pair-wise comparisons and swapping**

1	2	3	4	5	6
77	42	35	12	101	5

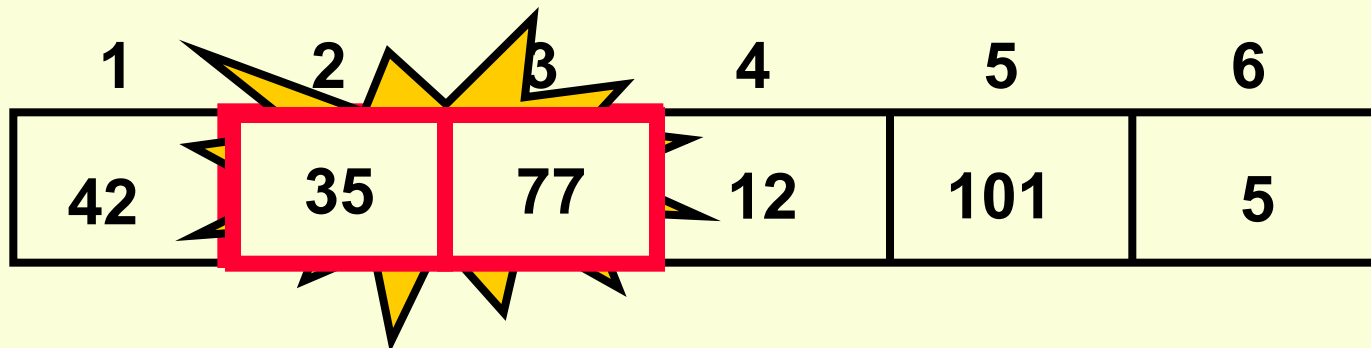
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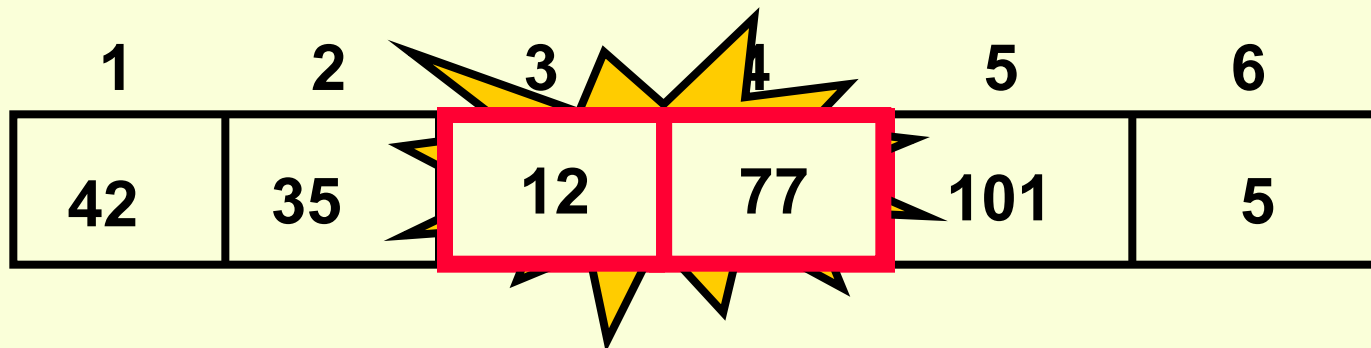
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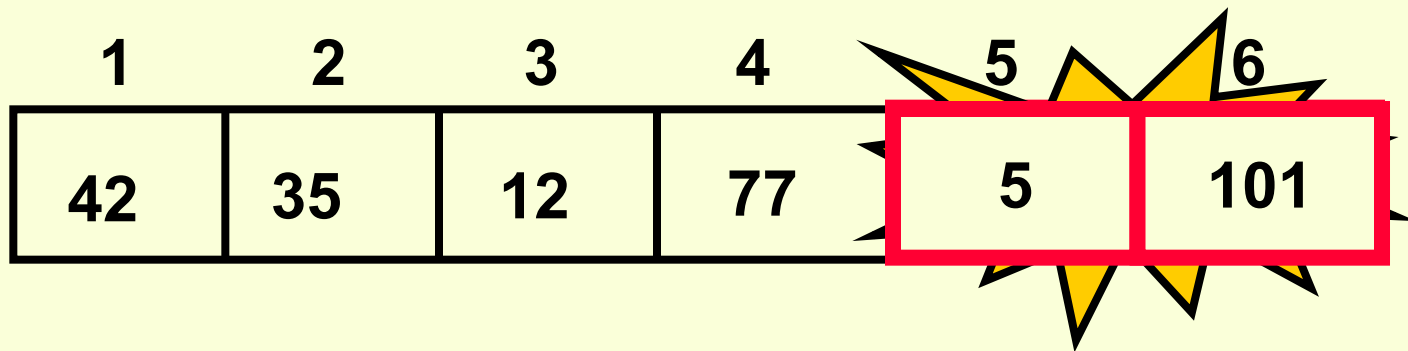
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1	2	3	4	5	6
42	35	12	77	101	5

No need to swap

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1	2	3	4	5	6
42	35	12	77	5	101

Largest value correctly placed

## Items of Interest

- Notice that only the largest value is correctly placed
- All other values are still out of order
- So we need to **repeat this process**

1	2	3	4	5	6
42	35	12	77	5	101

Largest value correctly placed



# Repeat “Bubble Up” How Many Times?

- If we have  $N$  elements...
- And if each time we bubble an element, we place it in its correct location...
- Then we repeat the “bubble up” process  $N - 1$  times.
- This guarantees we'll correctly place all  $N$  elements.

# “Bubbling” All the Elements

Diagram illustrating the first pass of bubble sort on an array of 6 elements. The elements are compared and swapped as they bubble towards the end of the array. The elements being compared or swapped are highlighted in red.

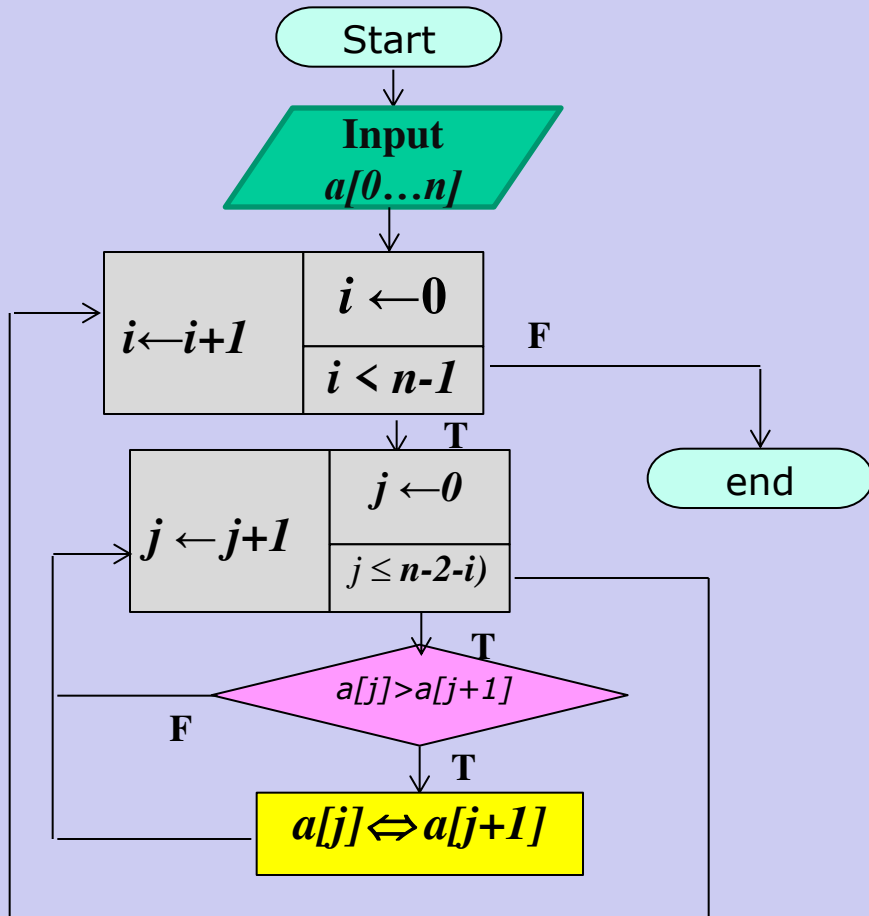
1	2	3	4	5	6
42	35	12	77	5	101
35	12	42	5	77	101
12	35	5	42	77	101
12	5	35	42	77	101
5	12	35	42	77	101

# Reducing the Number of Comparisons

1	2	3	4	5	6
77	42	35	12	101	5
1	2	3	4	5	6
42	35	12	77	5	101
1	2	3	4	5	6
35	12	42	5	77	101
1	2	3	4	5	6
12	35	5	42	77	101
1	2	3	4	5	6
12	5	35	42	77	101

# Bubble Sort

**Topic 5: Write an algorithm to sort an array using Bubble Sort.**



$n$ : total elements

Array Elements:  $a[0].....a[n-1]$

## Complexity Analysis of Bubble Sort

In Bubble Sort,  $n-1$  comparisons will be done in the 1st pass,  $n-2$  in 2nd pass,  $n-3$  in 3rd pass and so on. So the total number of comparisons will be,

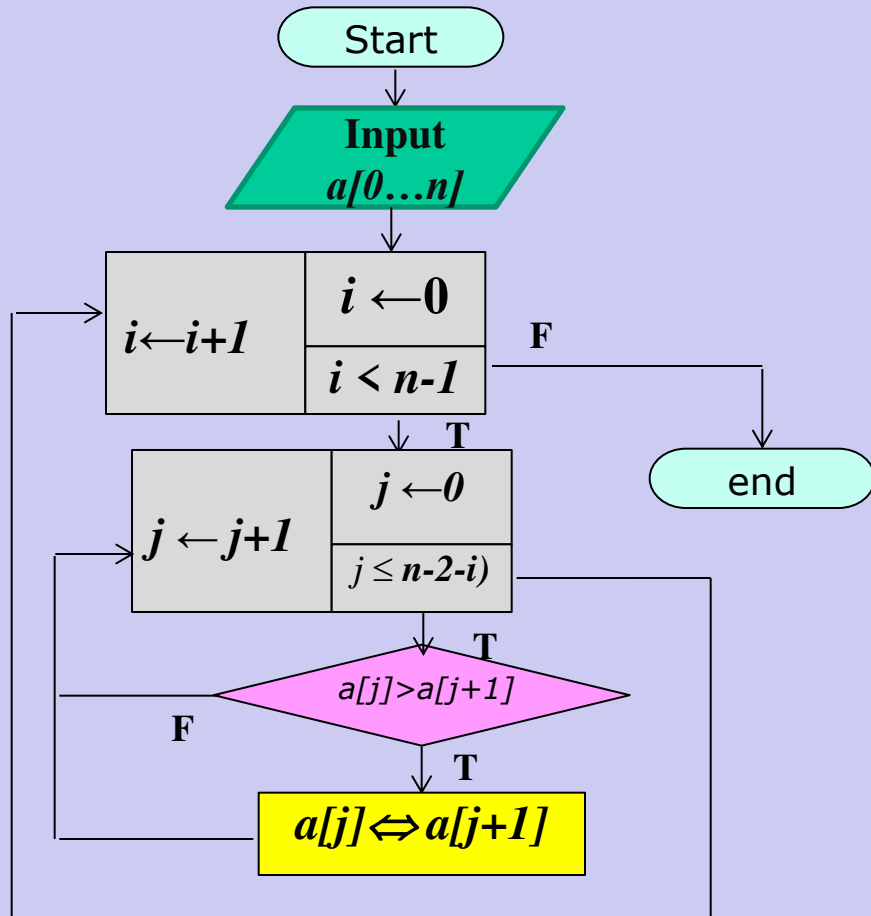
**Output:**

Sum =  $(n-1) + (n-2) + (n-3) + ..... + 3 + 2 + 1$

Sum =  $n(n-1)/2$  i.e  **$O(n^2)$**

# Bubble Sort

**Topic 5: Write an algorithm to sort an array using Bubble Sort.**



$n$ : total elements

Array Elements:  $a[0].....a[n-1]$

```
#include <stdio.h>
#include <stdlib.h>
int main()
{
    int a[6]={10,3,41,12,77,21};
    int n=6;
    int i,j,t;
    for(i=0;i<n-1;i++)
        for(j=0;j<=n-2-i;j++)
            if(a[j]>a[j+1])
            {
                t=a[j];a[j]=a[j+1];a[j+1]=t;
            }
    for(i=0;i<n;i++)
        printf("%d ",a[i]);
    return 0;
}
```

**Corresponding C program**

# Two-dimensional Arrays

## Matrix

→ In computer programming, a **matrix** can be defined with a 2-dimensional array. Any array with 'm' columns and 'n' rows represents a mXn matrix.

## Sparse Matrix

→ There may be a situation in which a matrix contains more number of ZERO values than NON-ZERO values. Such matrix is known as sparse matrix

## Example

consider a matrix of size 100 X 100 containing only 10 non-zero elements. In this matrix, only 10 spaces are filled with non-zero values and remaining spaces of matrix are filled with zero. That means, totally we allocate  $100 \times 100 \times 2 = 20000$  bytes of space to store this integer matrix. And to access these 10 non-zero elements we have to make scanning for 10000 times.

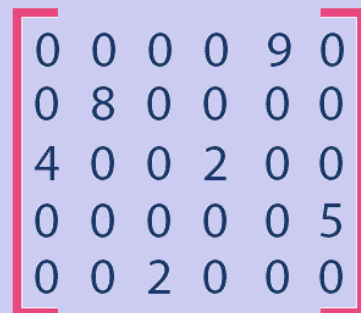
# Sparse Matrix

## Representation

- Triplet representation
- Linked representation

## Triplet Representation

In this representation, we consider only non-zero values along with their row and column index values. In this representation, the 0<sup>th</sup> row stores total rows, total columns and total non-zero values in the matrix. consider a matrix of size 5 X 6 containing 6 number of non-zero values. This matrix can be represented as shown in the image.

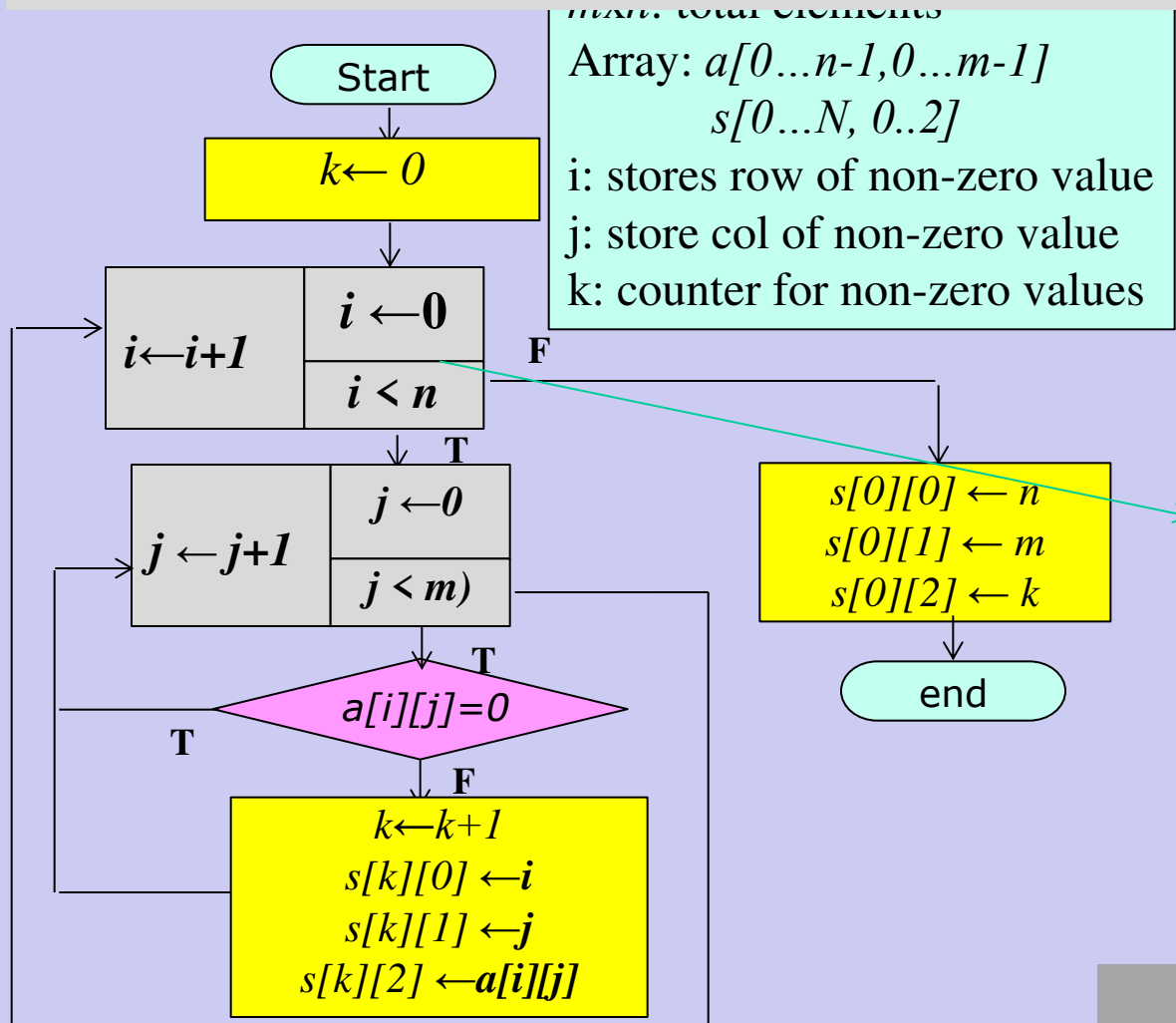


0	0	0	0	9	0
0	8	0	0	0	0
4	0	0	2	0	0
0	0	0	0	0	5
0	0	2	0	0	0

Rows	Columns	Values
5	6	6
0	4	9
1	1	8
2	0	4
2	2	2
3	5	5
4	2	2

# Sparse Matrix

## Topic 5: Algorithm for Triplex Representation



```
Int main(){
    int a[5][6], s[100][3];
    int k=0,n=5,m=6,i,j;
    for(i=0;i<n;i++){
        for(j=0;j<m;j++){
            if(a[i][j]!=0){
                k++;
                s[k][0]=i;
                s[k][1]=j;
                s[k][2]=a[i][j];
            }
        }
    }
    s[0][0]=n; s[0][1]=m; s[0][2]=k;
}
```

Corresponding C program



# Searching

**Search: locate an item in a list of data/information.**

**Two approaches will be discussed...**

**1. Linear or Sequential Search:**

- **Searches sequentially for an element.**
- **Starts from the first element.**

**2. Binary Search:**

- **Searches an element by dividing the sorted elements in a list into two sub-list**
- **Starts with the middle element.**

# Linear Search

**Algorithm:**

**Input:** Array, #elements, item (to search)

**Start with the element at index = 0**

**Step 1:** Compare the element at *index* with item. If its equal to item then return *index* with status “Found” otherwise go to step 2.

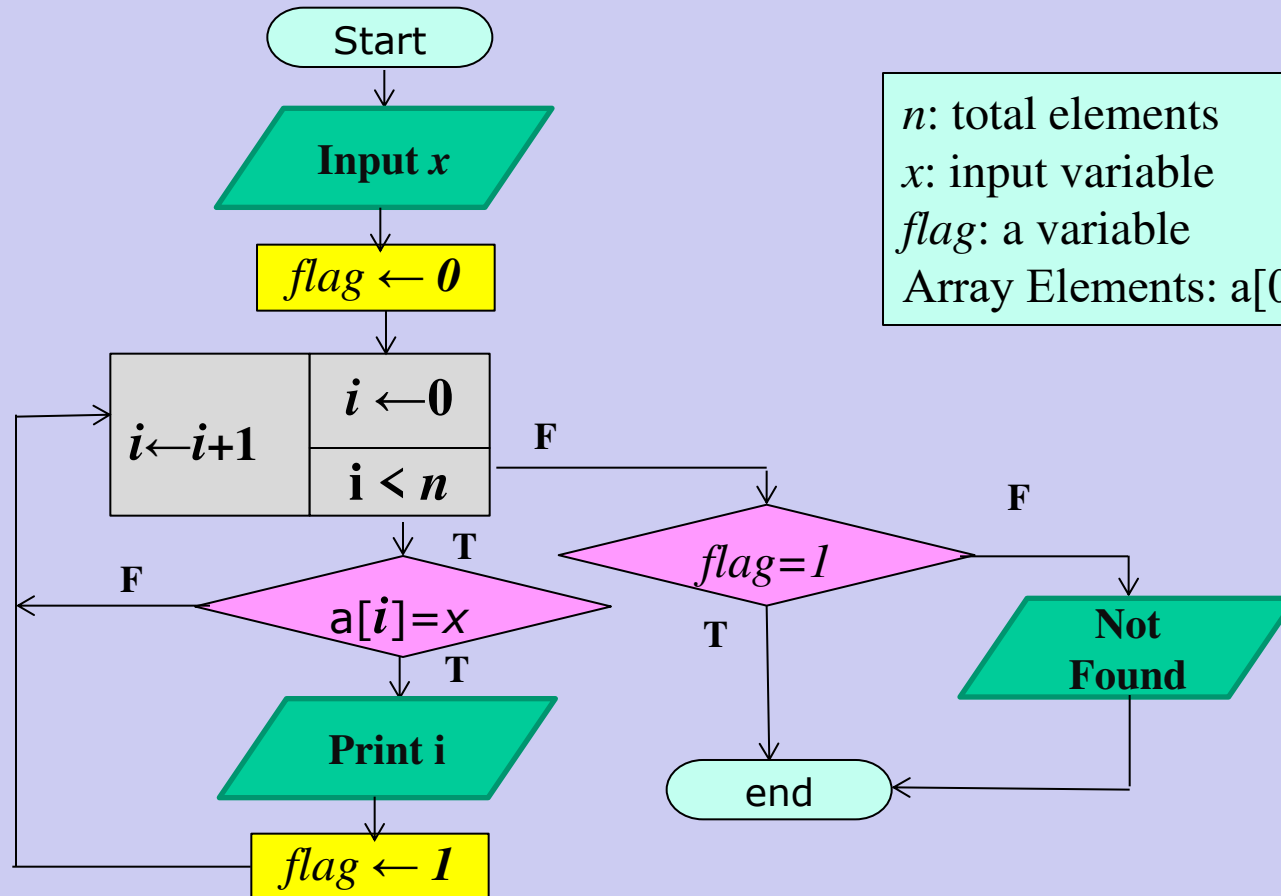
**Step 2:** Increase index by 1. If index is less than #elements go to step 1 otherwise return -1 with status “Not found”.

<i>index</i>	<i>index</i>	<i>index</i>	<i>index</i>	<i>index</i>	<i>index</i>				
0	1	2	3	4	5	6	7	8	9
30	62	20	100	77	15	90	52	88	45

<i>value</i>	15
<i>found</i>	true
<i>position</i>	5

# Linear Searching

## Topic 3: Modifications of previous algorithm



$n$ : total elements

$x$ : input variable

$flag$ : a variable

Array Elements:  $a[0] \dots a[n-1]$

### Complexity

$O(1)$  for best-case

$O(n)$  for worst-case

# Binary Search

Suppose an array `ax[]` has 10 elements

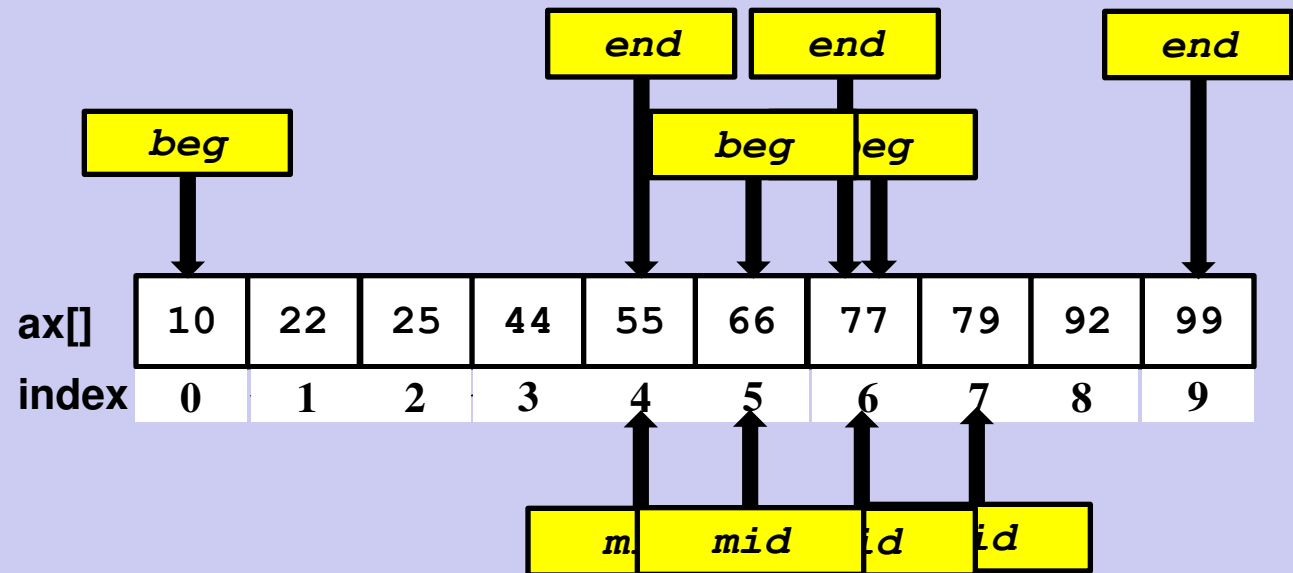
value	65
-------	----

beg	5
-----	---

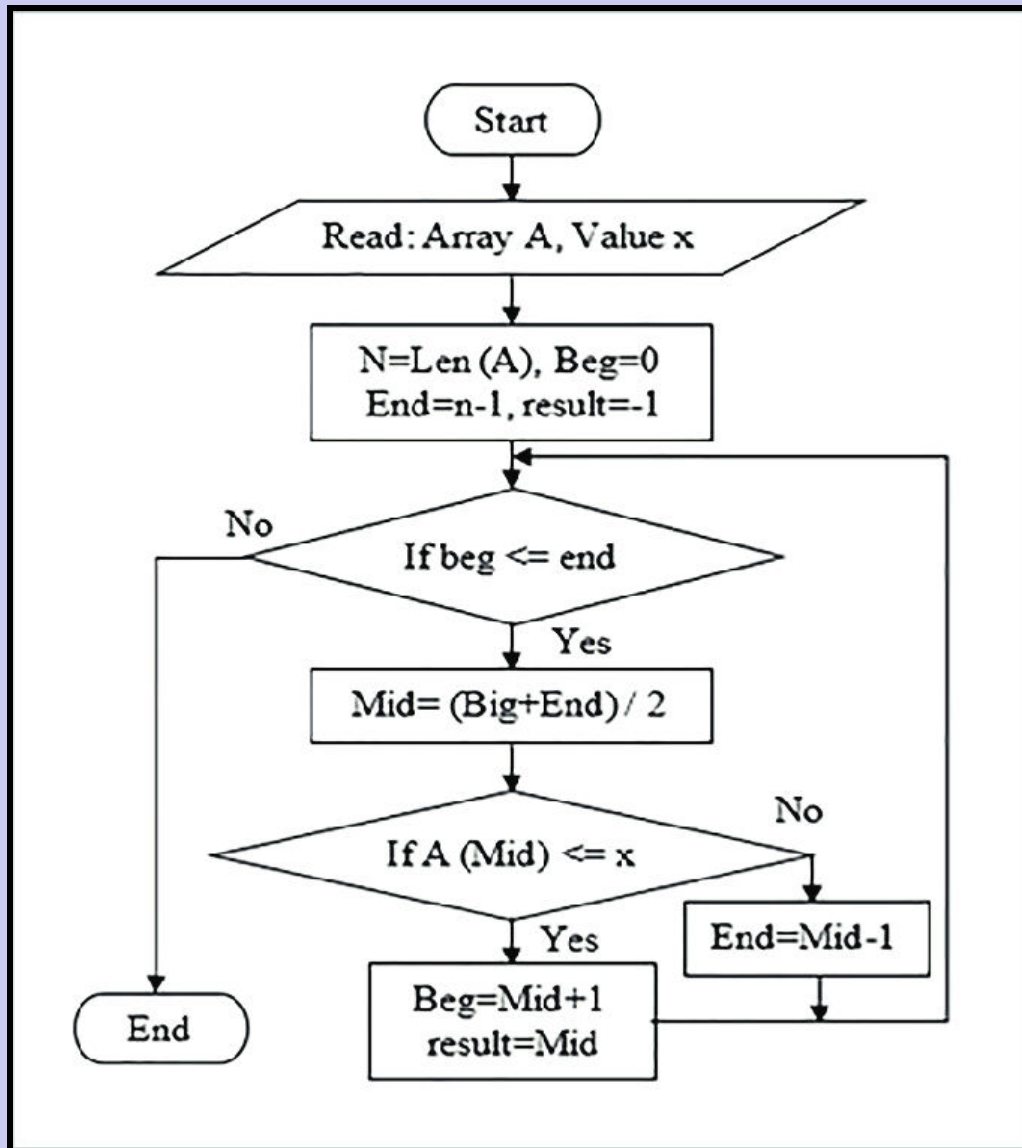
end	4
-----	---

mid	6
-----	---

not found
-----------



# Flowchart: Binary Search



At Iteration 1:

*Length of array =  $n$*

At Iteration 2:

*Length of array =  $n/2$*

At Iteration 3:

*Length of array =  $(n/2)/2 = n/2^2$*

Therefore, after Iteration  $k$ :

*Length of array =  $n/2^k$*

Also, we know that after  $k$  iterations, the **length of the array becomes**

**1** Therefore, the Length of the array

$$n/2^k = 1$$

$$\Rightarrow n = 2^k$$

Applying log function on both sides:

$$\Rightarrow \log_2 n = \log_2 2^k$$

$$\Rightarrow \log_2 n = k * \log_2 2$$

As  $(\log_a(a) = 1)$  Therefore,  $k = \log_2(n)$

**Time Complexity =  $O(\log_2 n)$**

# Binary Search Program

```
#include <iostream>

using namespace std;

int binarySearch(int array[], int x, int low, int high)
{
    // Repeat until the pointers low and high meet
    // each other

    while (low <= high) {
        int mid = low + (high - low) / 2;

        if (array[mid] == x)
            return mid;

        if (array[mid] < x)
            low = mid + 1;
        else
            high = mid - 1;
    }

    return -1;
}
```

```
int main(void) {
    int array[] = {3, 4, 5, 6, 7, 8, 9};
    int x = 7;

    //int n = sizeof(array) / sizeof(array[0]);
    n=7;

    int result = binarySearch(array, x, 0, n - 1);
    if (result == -1)
        printf("Not found");
    else
        printf("Element is found at index %d",
            result);
}
```



Quiz Time

Let's have  
some fun!

**Q1: An array is what kind of data structure?**

**A**

**Linear**

**B**

**Non-linear**

**C**

**Sorted**

**D**

**Sequential**

**30**



Q2: If a new element is needed to be stored at index 5 of 100 elements array then how many elements should be shifted?

- A** 5
- B** 100
- C** 95
- D** 96

30

Q3: If an element is needed to be deleted with index 10 of 100 elements array then how many elements should be shifted?

**A** 10

**B** 90

**C** 89

**D** 110

**30**

Q4: How many times bubble up is repeated in Bubble Sort algorithm with 10 elements?

**A** 10

**B** 11

**C** 9

**D** 12

30



Q5: The complexity of Bubble Sort algorithm is \_\_\_\_\_?

- A**  $O(n)$
- B**  $O(n-1)$
- C**  $O(2n)$
- D**  $O(n^2)$

30

Q6: The sparse matrix has more \_\_\_\_\_ values?

**A**

zero

**B**

non-zero

**C**

positive

**D**

negative

30

Q7: Which searching can be performed with unsorted data?

**A**

Binary Search

**B**

Linear Search

**C**

**D**

30



Q9: The complexity of Binary search is \_\_\_\_\_?

- A**  $O(n)$
- B**  $O(2^n)$
- C**  $O(n^2)$
- D**  $O(\log_2 n)$

30

# Q10: Identify the following \_\_\_\_\_.

A

G

B

M

C

R

D

D

First Commercial Processor	Price US\$60
Launched	November 15, 1971; 50 years ago
Discontinued	1981
Common manufacturer(s)	•Intel
Performance	
Max. CPU clock rate	740-750 kHz
Data width	4 bits
Address width	12 bits (multiplexed)

30



# Assignments

**Prob 1: Write an algorithm to insert an element after a specific element.**

**Prob 2: Write an algorithm to delete all the multiple matching elements.**

**Prob 3: Write an algorithm to split an array using a particular condition.**

**Prob 4: Write an algorithm to merge two sorted arrays into one sorted array.**

**Prob 5: Write an algorithm to multiply to matrices.**

**Prob 6: Write C programs for Creating Triplex form of Sparse Matrix.**

**End of Chapter 2**