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Contents

Foreword

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1 Scope

The present document specifies the coding, multiplexing and mapping to physical channels for E-UTRA.

2 References

The following documents contain provisions which, through reference in this text, constitute provisions of the present document.

- References are either specific (identified by date of publication, edition number, version number, etc.) or non-specific.
- For a specific reference, subsequent revisions do not apply.
- For a non-specific reference, the latest version applies. In the case of a reference to a 3GPP document (including a GSM document), a non-specific reference implicitly refers to the latest version of that document *in the same Release as the present document*.

[1] 3GPP TR 21.905: "Vocabulary for 3GPP Specifications".

[2] 3GPP TS 36.211: "Evolved Universal Terrestrial Radio Access (E-UTRA); Physical channels and modulation".

[3] 3GPP TS 36.213: "Evolved Universal Terrestrial Radio Access (E-UTRA); Physical layer procedures".

[4] 3GPP TS 36.306: "Evolved Universal Terrestrial Radio Access (E-UTRA); User Equipment (UE) radio access capabilities".

3 Definitions, symbols and abbreviations

3.1 Definitions

For the purposes of the present document, the terms and definitions given in [1] and the following apply. A term defined in the present document takes precedence over the definition of the same term, if any, in [1].

Definition format

<defined term>*: <definition>.*

3.2 Symbols

For the purposes of the present document, the following symbols apply:

$N_{\mathrm{RB}}^{\mathrm{DL}}$	Downlink bandwidth configuration, expressed in number of resource blocks [2]
$N_{\mathrm{RB}}^{\mathrm{UL}}$	Uplink bandwidth configuration, expressed in number of resource blocks [2]
$N_{\mathrm{symb}}^{\mathrm{PUSCH}}$	Number of SC-FDMA symbols carrying PUSCH in a subframe
$N_{\mathrm{symb}}^{\mathrm{UL}}$	Number of SC-FDMA symbols in an uplink slot
N_{SRS}	Number of SC-FDMA symbols used for SRS transmission in a subframe (0 or 1).

3.3 Abbreviations

For the purposes of the present document, the following abbreviations apply:

BCH	Broadcast channel
CFI	Control Format Indicator
CP	Cyclic Prefix
DCI	Downlink Control Information
DL-SCH	Downlink Shared channel
FDD	Frequency Division Duplexing
HI	HARQ indicator
MCH	Multicast channel
PBCH	Physical Broadcast channel
PCFICH	Physical Control Format Indicator channel
PCH	Paging channel
PDCCH	Physical Downlink Control channel
PDSCH	Physical Downlink Shared channel
PHICH	Physical HARQ indicator channel
PMCH	Physical Multicast channel
PRACH	Physical Random Access channel
PUCCH	Physical Uplink Control channel
PUSCH	Physical Uplink Shared channel
RACH	Random Access channel
SRS	Sounding Reference Signal
TDD	Time Division Duplexing
UCI	Uplink Control Information
UL-SCH	Uplink Shared channel

4 Mapping to physical channels

4.1 Uplink

Table 4.1-1 specifies the mapping of the uplink transport channels to their corresponding physical channels. Table 4.1-2 specifies the mapping of the uplink control channel information to its corresponding physical channel.

Table 4.1-1

TrCH	Physical Channel
UL-SCH	PUSCH
RACH	PRACH

Table 4.1-2

Control information	Physical Channel
UCI	PUCCH, PUSCH

4.2 Downlink

Table 4.2-1 specifies the mapping of the downlink transport channels to their corresponding physical channels. Table 4.2-2 specifies the mapping of the downlink control channel information to its corresponding physical channel.

Table 4.2-1

TrCH	Physical Channel
DL-SCH	PDSCH
BCH	PBCH
PCH	PDSCH
MCH	PMCH

Table 4.2-2

Control information	Physical Channel
CFI	PCFICH
HI	PHICH
DCI	PDCCH

5 Channel coding, multiplexing and interleaving

Data and control streams from/to MAC layer are encoded /decoded to offer transport and control services over the radio transmission link. Channel coding scheme is a combination of error detection, error correcting, rate matching, interleaving and transport channel or control information mapping onto/splitting from physical channels.

5.1 Generic procedures

This section contains coding procedures which are used for more than one transport channel or control information type.

5.1.1 CRC calculation

Denote the input bits to the CRC computation by $a_0, a_1, a_2, a_3, ..., a_{A-1}$, and the parity bits by $p_0, p_1, p_2, p_3, ..., p_{L-1}$. A is the size of the input sequence and L is the number of parity bits. The parity bits are generated by one of the following cyclic generator polynomials:

- $g_{CRC24A}(D) = [D^{24} + D^{23} + D^{18} + D^{17} + D^{14} + D^{11} + D^{10} + D^{7} + D^{6} + D^{5} + D^{4} + D^{3} + D + 1]$ and;
- $g_{CRC24B}(D) = [D^{24} + D^{23} + D^{6} + D^{5} + D + 1]$ for a CRC length $L = 24$ and;
- $g_{CRC16}(D) = [D^{16} + D^{12} + D^{5} + 1]$ for a CRC length $L = 16$.
- $g_{CRC8}(D) = [D^{8} + D^{7} + D^{4} + D^{3} + D + 1]$ for a CRC length of $L = 8$.

The encoding is performed in a systematic form, which means that in GF(2), the polynomial:

$$a_0 D^{A+23} + a_1 D^{A+22} + ... + a_{A-1} D^{24} + p_0 D^{23} + p_1 D^{22} + ... + p_{22} D^1 + p_{23}$$

yields a remainder equal to 0 when divided by the corresponding length-24 CRC generator polynomial, $g_{CRC24A}(D)$ or $g_{CRC24B}(D)$, the polynomial:

$$a_0 D^{A+15} + a_1 D^{A+14} + ... + a_{A-1} D^{16} + p_0 D^{15} + p_1 D^{14} + ... + p_{14} D^1 + p_{15}$$

yields a remainder equal to 0 when divided by $g_{CRC16}(D)$, and the polynomial:

$$a_0 D^{A+7} + a_1 D^{A+6} + ... + a_{A-1} D^{8} + p_0 D^{7} + p_1 D^{6} + ... + p_{6} D^1 + p_{7}$$

yields a remainder equal to 0 when divided by $g_{CRC8}(D)$.

The bits after CRC attachment are denoted by $b_0, b_1, b_2, b_3, ..., b_{B-1}$, where $B = A + L$. The relation between a_k and b_k is:

$b_k = a_k$ for k = 0, 1, 2, …, A-1

$b_k = p_{k-A}$ for $k = A, A+1, A+2, ..., A+L-1$.

5.1.2 Code block segmentation and code block CRC attachment

The input bit sequence to the code block segmentation is denoted by $b_0, b_1, b_2, b_3, ..., b_{B-1}$, where $B > 0$. If B is larger than the maximum code block size Z, segmentation of the input bit sequence is performed and an additional CRC sequence of $L = 24$ bits is attached to each code block. The maximum code block size is:

- $Z = 6144$.

If the number of filler bits F calculated below is not 0, filler bits are added to the beginning of the first block.

Note that if $B < 40$, filler bits are added to the beginning of the code block.

The filler bits shall be set to <*NULL*> at the input to the encoder.

Total number of code blocks C is determined by:

if $B \le Z$

$L = 0$

Number of code blocks: $C = 1$

$B' = B$

else

$L = 24$

Number of code blocks: $C = \lceil B/(Z-L) \rceil$.

$B' = B + C \cdot L$

end if

The bits output from code block segmentation, for $C \ne 0$, are denoted by $c_{r0}, c_{r1}, c_{r2}, c_{r3}, ..., c_{r(K_r-1)}$, where r is the code block number, and K_r is the number of bits for the code block number r.

Number of bits in each code block (applicable for $C \ne 0$ only):

First segmentation size: K_+ = minimum K in table 5.1.3-3 such that $C \cdot K \ge B'$

if $C = 1$

the number of code blocks with length K_+ is $C_+ = 1$, $K_- = 0$, $C_- = 0$

else if $C > 1$

Second segmentation size: K_- = maximum K in table 5.1.3-3 such that $K < K_+$

$\Delta_K = K_+ - K_-$

Number of segments of size K_-: $C_- = \left\lfloor \frac{C \cdot K_+ - B'}{\Delta_K} \right\rfloor$.

Number of segments of size K_+: $C_+ = C - C_-$.

end if

Number of filler bits: $F = C_+ \cdot K_+ + C_- \cdot K_- - B'$

for k = 0 to F-1 -- Insertion of filler bits

$c_{0k} = <NULL>$

end for

$k = F$

$s = 0$

for r = 0 to C-1

if $r < C_-$

$K_r = K_-$

else

$K_r = K_+$

end if

while $k < K_r - L$

$c_{rk} = b_s$

$k = k + 1$

$s = s + 1$

end while

if C >1

The sequence $c_{r0}, c_{r1}, c_{r2}, c_{r3}, ..., c_{r(K_r - L - 1)}$ is used to calculate the CRC parity bits $p_{r0}, p_{r1}, p_{r2}, ..., p_{r(L-1)}$ according to subclause 5.1.1 with the generator polynomial $g_{CRC24B}(D)$. For CRC calculation it is assumed that filler bits, if present, have the value 0.

while $k < K_r$

$c_{rk} = p_{r(k+L-K_r)}$

$k = k + 1$

end while

end if

$k = 0$

end for

5.1.3 Channel coding

The bit sequence input for a given code block to channel coding is denoted by $c_0, c_1, c_2, c_3, ..., c_{K-1}$, where K is the number of bits to encode. After encoding the bits are denoted by $d_0^{(i)}, d_1^{(i)}, d_2^{(i)}, d_3^{(i)}, ..., d_{D-1}^{(i)}$, where D is the number of encoded bits per output stream and i indexes the encoder output stream. The relation between c_k and $d_k^{(i)}$ and between K and D is dependent on the channel coding scheme.

The following channel coding schemes can be applied to TrCHs:

- tail biting convolutional coding;
- turbo coding.

Usage of coding scheme and coding rate for the different types of TrCH is shown in table 5.1.3-1. Usage of coding scheme and coding rate for the different control information types is shown in table 5.1.3-2.

The values of D in connection with each coding scheme:

- tail biting convolutional coding with rate 1/3: $D = K$;
- turbo coding with rate 1/3: $D = K + 4$.

The range for the output stream index i is 0, 1 and 2 for both coding schemes.

Table 5.1.3-1: Usage of channel coding scheme and coding rate for TrCHs

TrCH	Coding scheme	Coding rate
UL-SCH	Turbo coding	1/3
DL-SCH		
PCH		
MCH		
BCH	Tail biting convolutional coding	1/3

Table 5.1.3-2: Usage of channel coding scheme and coding rate for control information

Control Information	Coding scheme	Coding rate
DCI	Tail biting convolutional coding	1/3
CFI	Block code	1/16
HI	Repetition code	1/3
UCI	Block code	variable
	Tail biting convolutional coding	1/3

5.1.3.1 Tail biting convolutional coding

A tail biting convolutional code with constraint length 7 and coding rate 1/3 is defined.

The configuration of the convolutional encoder is presented in figure 5.1.3-1.

The initial value of the shift register of the encoder shall be set to the values corresponding to the last 6 information bits in the input stream so that the initial and final states of the shift register are the same. Therefore, denoting the shift register of the encoder by $s_0, s_1, s_2, ..., s_5$, then the initial value of the shift register shall be set to

$$s_i = c_{(K-1-i)}$$

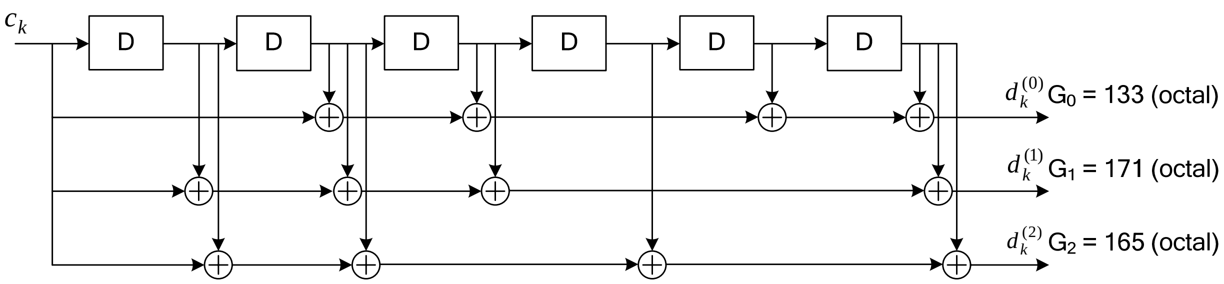

Figure 5.1.3-1: Rate 1/3 tail biting convolutional encoder

The encoder output streams $d_k^{(0)}$, $d_k^{(1)}$ and $d_k^{(2)}$ correspond to the first, second and third parity streams, respectively as shown in Figure 5.1.3-1.

5.1.3.2 Turbo coding

5.1.3.2.1 Turbo encoder

The scheme of turbo encoder is a Parallel Concatenated Convolutional Code (PCCC) with two 8-state constituent encoders and one turbo code internal interleaver. The coding rate of turbo encoder is 1/3. The structure of turbo encoder is illustrated in figure 5.1.3-2.

The transfer function of the 8-state constituent code for the PCCC is:

$$G(D) = \left[1, \frac{g_1(D)}{g_0(D)}\right],$$

where

$$g_0(D) = 1 + D^2 + D^3,$$

$$g_1(D) = 1 + D + D^3.$$

The initial value of the shift registers of the 8-state constituent encoders shall be all zeros when starting to encode the input bits.

The output from the turbo encoder is

$$d_k^{(0)} = x_k$$

$$d_k^{(1)} = z_k$$

$$d_k^{(2)} = z'_k$$

for $k = 0,1,2,..., K-1$.

If the code block to be encoded is the 0-th code block and the number of filler bits is greater than zero, i.e., $F > 0$, then the encoder shall set c_k, = 0, k = 0,…,(F-1) at its input and shall set $d_k^{(0)} = <NULL>$, k = 0,…,(F-1) and $d_k^{(1)} = <NULL>$, k = 0,…,(F-1) at its output.

The bits input to the turbo encoder are denoted by $c_0, c_1, c_2, c_3,..., c_{K-1}$, and the bits output from the first and second 8-state constituent encoders are denoted by $z_0, z_1, z_2, z_3,..., z_{K-1}$ and $z'_0, z'_1, z'_2, z'_3,..., z'_{K-1}$, respectively. The bits output from the turbo code internal interleaver are denoted by $c'_0, c'_1,..., c'_{K-1}$, and these bits are to be the input to the second 8-state constituent encoder.

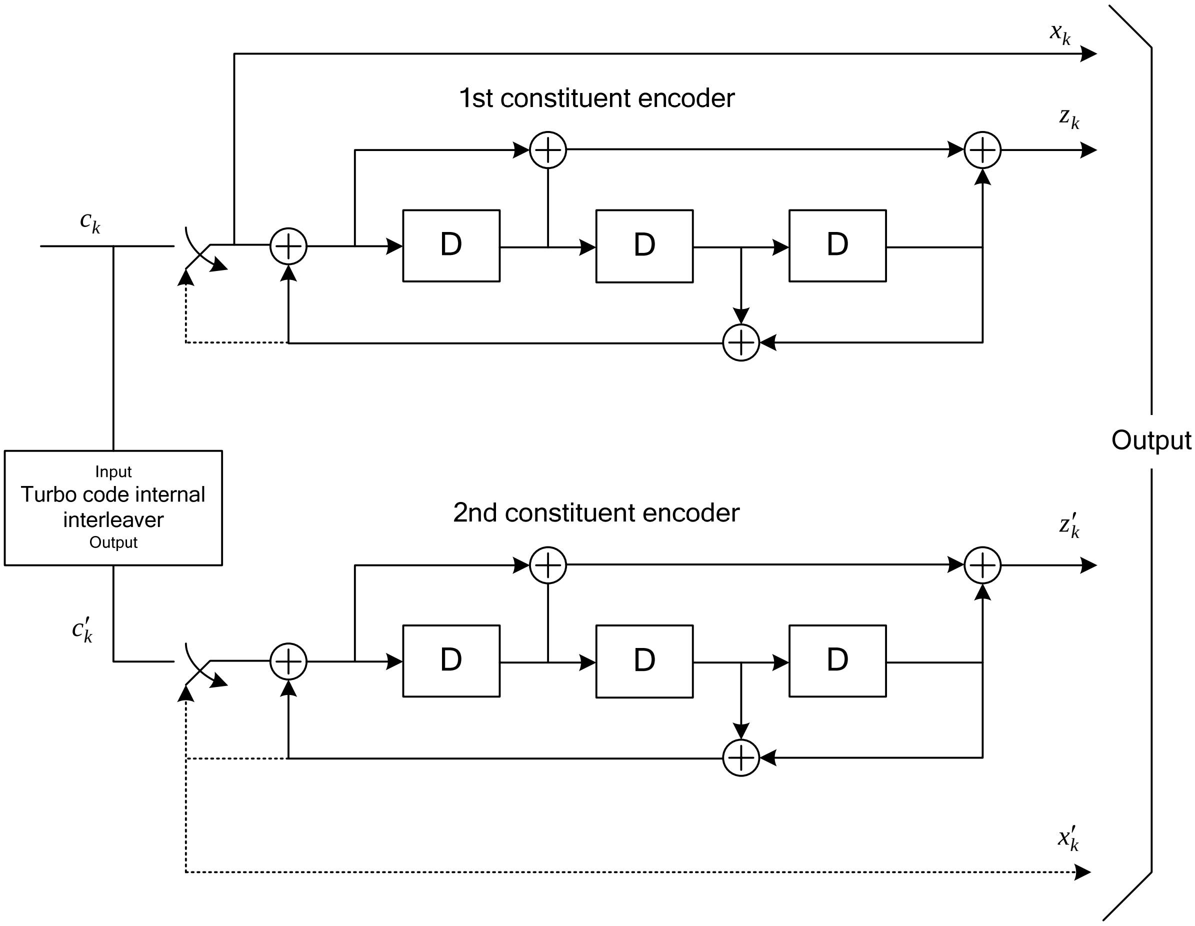

Figure 5.1.3-2: Structure of rate 1/3 turbo encoder (dotted lines apply for trellis termination only)

5.1.3.2.2 Trellis termination for turbo encoder

Trellis termination is performed by taking the tail bits from the shift register feedback after all information bits are encoded. Tail bits are padded after the encoding of information bits.

The first three tail bits shall be used to terminate the first constituent encoder (upper switch of figure 5.1.3-2 in lower position) while the second constituent encoder is disabled. The last three tail bits shall be used to terminate the second constituent encoder (lower switch of figure 5.1.3-2 in lower position) while the first constituent encoder is disabled.

The transmitted bits for trellis termination shall then be:

$$d_K^{(0)} = x_K\,,\ d_{K+1}^{(0)} = z_{K+1}\,,\ d_{K+2}^{(0)} = x'_K\,,\ d_{K+3}^{(0)} = z'_{K+1}$$

$$d_K^{(1)} = z_K\,,\ d_{K+1}^{(1)} = x_{K+2}\,,\ d_{K+2}^{(1)} = z'_K\,,\ d_{K+3}^{(1)} = x'_{K+2}$$

$$d_K^{(2)} = x_{K+1}\,,\ d_{K+1}^{(2)} = z_{K+2}\,,\ d_{K+2}^{(2)} = x'_{K+1}\,,\ d_{K+3}^{(2)} = z'_{K+2}$$

5.1.3.2.3 Turbo code internal interleaver

The bits input to the turbo code internal interleaver are denoted by $c_0, c_1, ..., c_{K-1}$, where K is the number of input bits. The bits output from the turbo code internal interleaver are denoted by $c'_0, c'_1, ..., c'_{K-1}$.

The relationship between the input and output bits is as follows:

$c'_i = c_{\Pi(i)}$, i=0, 1,…, (K-1) where the relationship between the output index i and the input index $\Pi(i)$ satisfies the following quadratic form:

$$\Pi(i) = \left(f_1 \cdot i + f_2 \cdot i^2\right) \bmod K$$

The parameters f_1 and f_2 depend on the block size K and are summarized in Table 5.1.3-3.

Table 5.1.3-3: Turbo code internal interleaver parameters

i	K_i	f_1	f_2	*i*	K_i	f_1	f_2	*i*	K_i	f_1	f_2	*i*	K_i	f_1	f_2
1	40	3	10	48	416	25	52	95	1120	67	140	142	3200	111	240
2	48	7	12	49	424	51	106	96	1152	35	72	143	3264	443	204
3	56	19	42	50	432	47	72	97	1184	19	74	144	3328	51	104
4	64	7	16	51	440	91	110	98	1216	39	76	145	3392	51	212
5	72	7	18	52	448	29	168	99	1248	19	78	146	3456	451	192
6	80	11	20	53	456	29	114	100	1280	199	240	147	3520	257	220
7	88	5	22	54	464	247	58	101	1312	21	82	148	3584	57	336
8	96	11	24	55	472	29	118	102	1344	211	252	149	3648	313	228
9	104	7	26	56	480	89	180	103	1376	21	86	150	3712	271	232
10	112	41	84	57	488	91	122	104	1408	43	88	151	3776	179	236
11	120	103	90	58	496	157	62	105	1440	149	60	152	3840	331	120
12	128	15	32	59	504	55	84	106	1472	45	92	153	3904	363	244
13	136	9	34	60	512	31	64	107	1504	49	846	154	3968	375	248
14	144	17	108	61	528	17	66	108	1536	71	48	155	4032	127	168
15	152	9	38	62	544	35	68	109	1568	13	28	156	4096	31	64
16	160	21	120	63	560	227	420	110	1600	17	80	157	4160	33	130
17	168	101	84	64	576	65	96	111	1632	25	102	158	4224	43	264
18	176	21	44	65	592	19	74	112	1664	183	104	159	4288	33	134
19	184	57	46	66	608	37	76	113	1696	55	954	160	4352	477	408
20	192	23	48	67	624	41	234	114	1728	127	96	161	4416	35	138
21	200	13	50	68	640	39	80	115	1760	27	110	162	4480	233	280
22	208	27	52	69	656	185	82	116	1792	29	112	163	4544	357	142
23	216	11	36	70	672	43	252	117	1824	29	114	164	4608	337	480
24	224	27	56	71	688	21	86	118	1856	57	116	165	4672	37	146
25	232	85	58	72	704	155	44	119	1888	45	354	166	4736	71	444
26	240	29	60	73	720	79	120	120	1920	31	120	167	4800	71	120
27	248	33	62	74	736	139	92	121	1952	59	610	168	4864	37	152
28	256	15	32	75	752	23	94	122	1984	185	124	169	4928	39	462
29	264	17	198	76	768	217	48	123	2016	113	420	170	4992	127	234
30	272	33	68	77	784	25	98	124	2048	31	64	171	5056	39	158
31	280	103	210	78	800	17	80	125	2112	17	66	172	5120	39	80
32	288	19	36	79	816	127	102	126	2176	171	136	173	5184	31	96
33	296	19	74	80	832	25	52	127	2240	209	420	174	5248	113	902
34	304	37	76	81	848	239	106	128	2304	253	216	175	5312	41	166
35	312	19	78	82	864	17	48	129	2368	367	444	176	5376	251	336
36	320	21	120	83	880	137	110	130	2432	265	456	177	5440	43	170
37	328	21	82	84	896	215	112	131	2496	181	468	178	5504	21	86
38	336	115	84	85	912	29	114	132	2560	39	80	179	5568	43	174
39	344	193	86	86	928	15	58	133	2624	27	164	180	5632	45	176
40	352	21	44	87	944	147	118	134	2688	127	504	181	5696	45	178
41	360	133	90	88	960	29	60	135	2752	143	172	182	5760	161	120
42	368	81	46	89	976	59	122	136	2816	43	88	183	5824	89	182
43	376	45	94	90	992	65	124	137	2880	29	300	184	5888	323	184
44	384	23	48	91	1008	55	84	138	2944	45	92	185	5952	47	186
45	392	243	98	92	1024	31	64	139	3008	157	188	186	6016	23	94
46	400	151	40	93	1056	17	66	140	3072	47	96	187	6080	47	190
47	408	155	102	94	1088	171	204	141	3136	13	28	188	6144	263	480

5.1.4 Rate matching

5.1.4.1 Rate matching for turbo coded transport channels

The rate matching for turbo coded transport channels is defined per coded block and consists of interleaving the three information bit streams $d_k^{(0)}$, $d_k^{(1)}$ and $d_k^{(2)}$, followed by the collection of bits and the generation of a circular buffer as depicted in Figure 5.1.4-1. The output bits for each code block are transmitted as described in subclause 5.1.4.1.2.

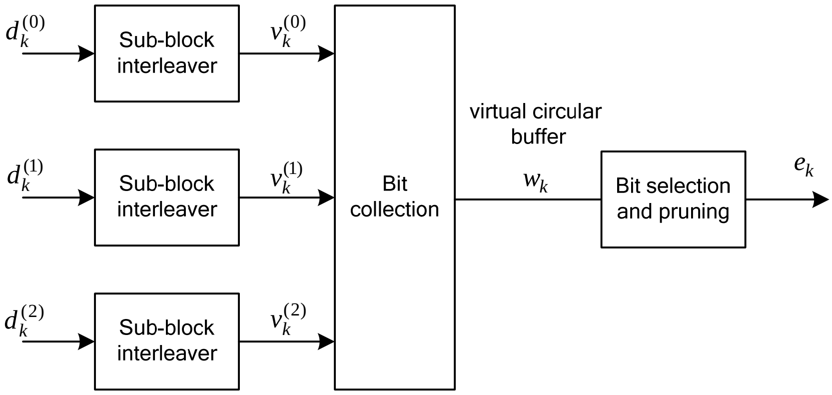

Figure 5.1.4-1. Rate matching for turbo coded transport channels

The bit stream $d_k^{(0)}$ is interleaved according to the sub-block interleaver defined in subclause 5.1.4.1.1 with an output sequence defined as $v_0^{(0)}, v_1^{(0)}, v_2^{(0)}, \ldots, v_{K_\Pi-1}^{(0)}$ and where K_Π is defined in subclause 5.1.4.1.1.

The bit stream $d_k^{(1)}$ is interleaved according to the sub-block interleaver defined in subclause 5.1.4.1.1 with an output sequence defined as $v_0^{(1)}, v_1^{(1)}, v_2^{(1)}, \ldots, v_{K_\Pi-1}^{(1)}$.

The bit stream $d_k^{(2)}$ is interleaved according to the sub-block interleaver defined in subclause 5.1.4.1.1 with an output sequence defined as $v_0^{(2)}, v_1^{(2)}, v_2^{(2)}, \ldots, v_{K_\Pi-1}^{(2)}$.

The sequence of bits e_k for transmission is generated according to subclause 5.1.4.1.2.

5.1.4.1.1 Sub-block interleaver

The bits input to the block interleaver are denoted by $d_0^{(i)}, d_1^{(i)}, d_2^{(i)}, \ldots, d_{D-1}^{(i)}$, where D is the number of bits. The output bit sequence from the block interleaver is derived as follows:

(1) Assign $C_{subblock}^{TC} = 32$ to be the number of columns of the matrix. The columns of the matrix are numbered 0, 1, 2,…, $C_{subblock}^{TC} - 1$ from left to right.

(2) Determine the number of rows of the matrix $R_{subblock}^{TC}$, by finding minimum integer $R_{subblock}^{TC}$ such that:

$$D \le \left(R_{subblock}^{TC} \times C_{subblock}^{TC}\right)$$

The rows of rectangular matrix are numbered 0, 1, 2,…, $R_{subblock}^{TC} - 1$ from top to bottom.

(3) If $\left(R_{subblock}^{TC} \times C_{subblock}^{TC}\right) > D$, then $N_D = \left(R_{subblock}^{TC} \times C_{subblock}^{TC} - D\right)$ dummy bits are padded such that y_k = <*NULL*> for k = 0, 1,…, N_D - 1. Then, write the input bit sequence, i.e. $y_{N_D+k} = d_k^{(i)}$, k = 0, 1,…, D-1, into the $\left(R_{subblock}^{TC} \times C_{subblock}^{TC}\right)$ matrix row by row starting with bit y_0 in column 0 of row 0:

$$\begin{bmatrix} y_0 & y_1 & y_2 & \cdots & y_{C_{subblock}^{TC}-1} \\ y_{C_{subblock}^{TC}} & y_{C_{subblock}^{TC}+1} & y_{C_{subblock}^{TC}+2} & \cdots & y_{2C_{subblock}^{TC}-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ y_{(R_{subblock}^{TC}-1)\times C_{subblock}^{TC}} & y_{(R_{subblock}^{TC}-1)\times C_{subblock}^{TC}+1} & y_{(R_{subblock}^{TC}-1)\times C_{subblock}^{TC}+2} & \cdots & y_{(R_{subblock}^{TC}\times C_{subblock}^{TC}-1)} \end{bmatrix}$$

For $d_k^{(0)}$ and $d_k^{(1)}$:

(4) Perform the inter-column permutation for the matrix based on the pattern $\langle P(j) \rangle_{j\in\{0,1,\ldots,C_{subblock}^{TC}-1\}}$ that is shown in table 5.1.4-1, where P(j) is the original column position of the j-th permuted column. After permutation of the columns, the inter-column permuted $\left(R_{subblock}^{TC} \times C_{subblock}^{TC}\right)$ matrix is equal to

$$\begin{bmatrix} y_{P(0)} & y_{P(1)} & y_{P(2)} & \cdots & y_{P(C_{subblock}^{TC}-1)} \\ y_{P(0)+C_{subblock}^{TC}} & y_{P(1)+C_{subblock}^{TC}} & y_{P(2)+C_{subblock}^{TC}} & \cdots & y_{P(C_{subblock}^{TC}-1)+C_{subblock}^{TC}} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ y_{P(0)+(R_{subblock}^{TC}-1)\times C_{subblock}^{TC}} & y_{P(1)+(R_{subblock}^{TC}-1)\times C_{subblock}^{TC}} & y_{P(2)+(R_{subblock}^{TC}-1)\times C_{subblock}^{TC}} & \cdots & y_{P(C_{subblock}^{TC}-1)+(R_{subblock}^{TC}-1)\times C_{subblock}^{TC}} \end{bmatrix}$$

(5) The output of the block interleaver is the bit sequence read out column by column from the inter-column permuted $\left(R_{subblock}^{TC} \times C_{subblock}^{TC}\right)$matrix. The bits after sub-block interleaving are denoted by $v_0^{(i)}, v_1^{(i)}, v_2^{(i)}, \ldots, v_{K_\Pi-1}^{(i)}$, where $v_0^{(i)}$ corresponds to $y_{P(0)}$, $v_1^{(i)}$ to $y_{P(0)+C_{subblock}^{TC}}$ … and $K_\Pi = \left(R_{subblock}^{TC} \times C_{subblock}^{TC}\right)$.

For $d_k^{(2)}$:

(4) The output of the sub-block interleaver is denoted by $v_0^{(2)}, v_1^{(2)}, v_2^{(2)}, \ldots, v_{K_\Pi-1}^{(2)}$, where $v_k^{(2)} = y_{\pi(k)}$ and where

$$\pi(k) = \left(P\left(\left\lfloor \frac{k}{R_{subblock}^{TC}} \right\rfloor \right) + C_{subblock}^{TC} \times \left(k \bmod R_{subblock}^{TC} \right) + 1 \right) \bmod K_\Pi$$

The permutation function P is defined in Table 5.1.4-1.

Table 5.1.4-1 Inter-column permutation pattern for sub-block interleaver

Number of columns $C_{subblock}^{TC}$	**Inter-column permutation pattern** $< P(0), P(1), \ldots, P(C_{subblock}^{TC}-1) >$
32	< 0, 16, 8, 24, 4, 20, 12, 28, 2, 18, 10, 26, 6, 22, 14, 30, 1, 17, 9, 25, 5, 21, 13, 29, 3, 19, 11, 27, 7, 23, 15, 31 >

5.1.4.1.2 Bit collection, selection and transmission

The circular buffer of length $K_w = 3K_\Pi$ for the r-th coded block is generated as follows:

$w_k = v_k^{(0)}$ for k = 0,…, $K_\Pi - 1$

$w_{K_\Pi+2k} = v_k^{(1)}$ for k = 0,…, $K_\Pi - 1$

$w_{K_\Pi+2k+1} = v_k^{(2)}$ for k = 0,…, $K_\Pi - 1$

Denote the soft buffer size for the transport block by N_{IR} bits and the soft buffer size for the r-th code block by N_{cb} bits. The size N_{cb} is obtained as follows, where C is the number of code blocks computed in subclause 5.1.2:

- $N_{cb} = \min\left(\left\lfloor \frac{N_{IR}}{C} \right\rfloor, K_w\right)$ for downlink turbo coded transport channels
- $N_{cb} = K_w$ for uplink turbo coded transport channels

where N_{IR} is equal to:

$$N_{IR} = \left\lfloor \frac{N_{soft}}{K_{\mathrm{MIMO}} \cdot \min\left(M_{\mathrm{DL\_HARQ}}, M_{\mathrm{limit}}\right)} \right\rfloor$$

where:

N_{soft} is the total number of soft channel bits [4].

K_{MIMO} is equal to 2 if the UE is configured to receive PDSCH transmissions based on transmission modes 3 or 4 as defined in Section 7.1 in [3], 1 otherwise.

$M_{\mathrm{DL\_HARQ}}$ is the maximum number of DL HARQ processes (8 for FDD; 4, 6, 7, 9, 10, 12 or 15 for TDD depending on the UL/DL configuration defined in [2]).

M_{limit} is a constant equal to 9.

Denoting by E the rate matching output sequence length for the r-th coded block, and rv_{idx} the redundancy version number for this transmission (rv_{idx} = 0, 1, 2 or 3), the rate matching output bit sequence is e_k, k = 0,1,..., $E-1$.

Define by G the total number of bits available for the transmission of one transport block.

Set $G' = G/\left(N_L \cdot Q_m\right)$ where Q_m is equal to 2 for QPSK, 4 for 16QAM and 6 for 64QAM, and where

- N_L is equal to 1 for transport blocks mapped onto one transmission layer, i.e., single-antenna, 1-layer spatial multiplexing, both transport blocks for 2-layer spatial multiplexing, or the first transport block for 3-layer spatial multiplexing, and
- N_L is equal to 2 for transport blocks mapped onto two or four transmission layers, i.e., 2-layer transmit diversity, the second transport block for 3-layer spatial multiplexing, both transport blocks for 4-layer spatial multiplexing, or 4-layer transmit diversity.

Set $\gamma = G' \bmod C$, where C is the number of code blocks computed in subclause 5.1.2.

if $r \le C - \gamma - 1$

set $E = N_L \cdot Q_m \cdot \lfloor G'/C \rfloor$

else

set $E = N_L \cdot Q_m \cdot \lceil G'/C \rceil$

end if

Set $k_0 = R_{subblock}^{TC} \cdot \left(2 \cdot \left\lceil \frac{N_{cb}}{8R_{subblock}^{TC}} \right\rceil \cdot rv_{idx} + 2\right)$, where $R_{subblock}^{TC}$ is the number of rows defined in subclause 5.1.4.1.1.

Set k = 0 and j = 0

while { $k < E$ }

if $w_{(k_0+j)\bmod N_{cb}} \neq <NULL>$

$e_k = w_{(k_0+j)\bmod N_{cb}}$

$k = k + 1$

end if

$j = j + 1$

end while

5.1.4.2 Rate matching for convolutionally coded transport channels and control information

The rate matching for convolutionally coded transport channels and control information consists of interleaving the three bit streams, $d_k^{(0)}$, $d_k^{(1)}$ and $d_k^{(2)}$, followed by the collection of bits and the generation of a circular buffer as depicted in Figure 5.1.4-2. The output bits are transmitted as described in subclause 5.1.4.2.2.

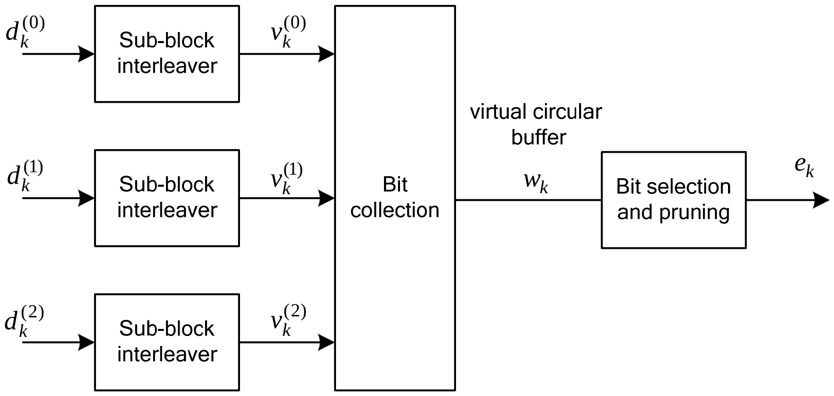

Figure 5.1.4-2. Rate matching for convolutionally coded transport channels and control information

The bit stream $d_k^{(0)}$ is interleaved according to the sub-block interleaver defined in subclause 5.1.4.2.1 with an output sequence defined as $v_0^{(0)}, v_1^{(0)}, v_2^{(0)},..., v_{K_\Pi-1}^{(0)}$ and where K_Π is defined in subclause 5.1.4.2.1.

The bit stream $d_k^{(1)}$ is interleaved according to the sub-block interleaver defined in subclause 5.1.4.2.1 with an output sequence defined as $v_0^{(1)}, v_1^{(1)}, v_2^{(1)},..., v_{K_\Pi-1}^{(1)}$.

The bit stream $d_k^{(2)}$ is interleaved according to the sub-block interleaver defined in subclause 5.1.4.2.1 with an output sequence defined as $v_0^{(2)}, v_1^{(2)}, v_2^{(2)},..., v_{K_\Pi-1}^{(2)}$.

The sequence of bits e_k for transmission is generated according to subclause 5.1.4.2.2.

5.1.4.2.1 Sub-block interleaver

The bits input to the block interleaver are denoted by $d_0^{(i)}, d_1^{(i)}, d_2^{(i)},..., d_{D-1}^{(i)}$, where D is the number of bits. The output bit sequence from the block interleaver is derived as follows:

(1) Assign $C_{subblock}^{CC} = 32$ to be the number of columns of the matrix. The columns of the matrix are numbered 0, 1, 2,…, $C_{subblock}^{CC} - 1$ from left to right.

(2) Determine the number of rows of the matrix $R_{subblock}^{CC}$, by finding minimum integer $R_{subblock}^{CC}$ such that:

$$D \le \left(R_{subblock}^{CC} \times C_{subblock}^{CC}\right)$$

The rows of rectangular matrix are numbered 0, 1, 2,…, $R_{subblock}^{CC} - 1$ from top to bottom.

(3) If $\left(R_{subblock}^{CC} \times C_{subblock}^{CC}\right) > D$, then $N_D = \left(R_{subblock}^{CC} \times C_{subblock}^{CC} - D\right)$ dummy bits are padded such that y_k = <*NULL*> for k = 0, 1,…, N_D - 1. Then, write the input bit sequence, i.e. $y_{N_D+k} = d_k^{(i)}$, k = 0, 1,…, D-1, into the $\left(R_{subblock}^{CC} \times C_{subblock}^{CC}\right)$ matrix row by row starting with bit y_0 in column 0 of row 0:

$$\begin{bmatrix} y_0 & y_1 & y_2 & \cdots & y_{C_{subblock}^{CC}-1} \\ y_{C_{subblock}^{CC}} & y_{C_{subblock}^{CC}+1} & y_{C_{subblock}^{CC}+2} & \cdots & y_{2C_{subblock}^{CC}-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ y_{(R_{subblock}^{CC}-1)\times C_{subblock}^{CC}} & y_{(R_{subblock}^{CC}-1)\times C_{subblock}^{CC}+1} & y_{(R_{subblock}^{CC}-1)\times C_{subblock}^{CC}+2} & \cdots & y_{(R_{subblock}^{CC}\times C_{subblock}^{CC}-1)} \end{bmatrix}$$

(4) Perform the inter-column permutation for the matrix based on the pattern $\langle P(j)\rangle_{j\in\{0,1,\dots,C_{subblock}^{CC}-1\}}$ that is shown in table 5.1.4-2, where P(j) is the original column position of the j-th permuted column. After permutation of the columns, the inter-column permuted $\left(R_{subblock}^{CC} \times C_{subblock}^{CC}\right)$ matrix is equal to

$$\begin{bmatrix} y_{P(0)} & y_{P(1)} & y_{P(2)} & \cdots & y_{P(C_{subblock}^{CC}-1)} \\ y_{P(0)+C_{subblock}^{CC}} & y_{P(1)+C_{subblock}^{CC}} & y_{P(2)+C_{subblock}^{CC}} & \cdots & y_{P(C_{subblock}^{CC}-1)+C_{subblock}^{CC}} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ y_{P(0)+(R_{subblock}^{CC}-1)\times C_{subblock}^{CC}} & y_{P(1)+(R_{subblock}^{CC}-1)\times C_{subblock}^{CC}} & y_{P(2)+(R_{subblock}^{CC}-1)\times C_{subblock}^{CC}} & \cdots & y_{P(C_{subblock}^{CC}-1)+(R_{subblock}^{CC}-1)\times C_{subblock}^{CC}} \end{bmatrix}$$

(5) The output of the block interleaver is the bit sequence read out column by column from the inter-column permuted $\left(R_{subblock}^{CC} \times C_{subblock}^{CC}\right)$ matrix. The bits after sub-block interleaving are denoted by $v_0^{(i)}, v_1^{(i)}, v_2^{(i)},\dots, v_{K_\Pi-1}^{(i)}$, where $v_0^{(i)}$ corresponds to $y_{P(0)}$, $v_1^{(i)}$ to $y_{P(0)+C_{subblock}^{CC}}$ … and $K_\Pi = \left(R_{subblock}^{CC} \times C_{subblock}^{CC}\right)$

Table 5.1.4-2 Inter-column permutation pattern for sub-block interleaver

Number of columns $C_{subblock}^{CC}$	**Inter-column permutation pattern** $< P(0), P(1),\dots, P(C_{subblock}^{CC}-1) >$
32	< 1, 17, 9, 25, 5, 21, 13, 29, 3, 19, 11, 27, 7, 23, 15, 31, 0, 16, 8, 24, 4, 20, 12, 28, 2, 18, 10, 26, 6, 22, 14, 30 >

This block interleaver is also used in interleaving PDCCH modulation symbols. In that case, the input bit sequence consists of PDCCH symbol quadruplets [2].

5.1.4.2.2 Bit collection, selection and transmission

The circular buffer of length $K_w = 3K_\Pi$ is generated as follows:

$w_k = v_k^{(0)}$ for k = 0,…, $K_\Pi - 1$

$w_{K_\Pi+k} = v_k^{(1)}$ for k = 0,…, $K_\Pi - 1$

$w_{2K_\Pi+k} = v_k^{(2)}$ for k = 0,…, $K_\Pi - 1$

Denoting by E the rate matching output sequence length, the rate matching output bit sequence is e_k , k = 0,1,..., $E-1$.

Set k = 0 and j = 0

while { $k < E$ }

if $w_{j \bmod K_w} \neq <NULL>$

$e_k = w_{j \bmod K_w}$

$k = k + 1$

end if

$j = j + 1$

end while

5.1.5 Code block concatenation

The input bit sequence for the code block concatenation and channel interleaving block are the sequences e_{rk}, for $r = 0,..., C-1$ and $k = 0,..., E_r - 1$. The output bit sequence from the code block concatenation and channel interleaving block is the sequence f_k for $k = 0,..., G-1$.

The code block concatenation consists of sequentially concatenating the rate matching outputs for the different code blocks. Therefore,

Set $k = 0$ and $r = 0$

while $r < C$

Set $j = 0$

while $j < E_r$

$f_k = e_{rj}$

$k = k + 1$

$j = j + 1$

end while

$r = r + 1$

end while

5.2 Uplink transport channels and control information

5.2.1 Random access channel

The sequence index for the random access channel is received from higher layers and is processed according to [2].

5.2.2 Uplink shared channel

Figure 5.2.2-1 shows the processing structure for the UL-SCH transport channel. Data arrives to the coding unit in form of a maximum of one transport block every transmission time interval (TTI). The following coding steps can be identified:

- Add CRC to the transport block
- Code block segmentation and code block CRC attachment
- Channel coding of data and control information

- Rate matching
- Code block concatenation
- Multiplexing of data and control information
- Channel interleaver

The coding steps for UL-SCH transport channel are shown in the figure below.

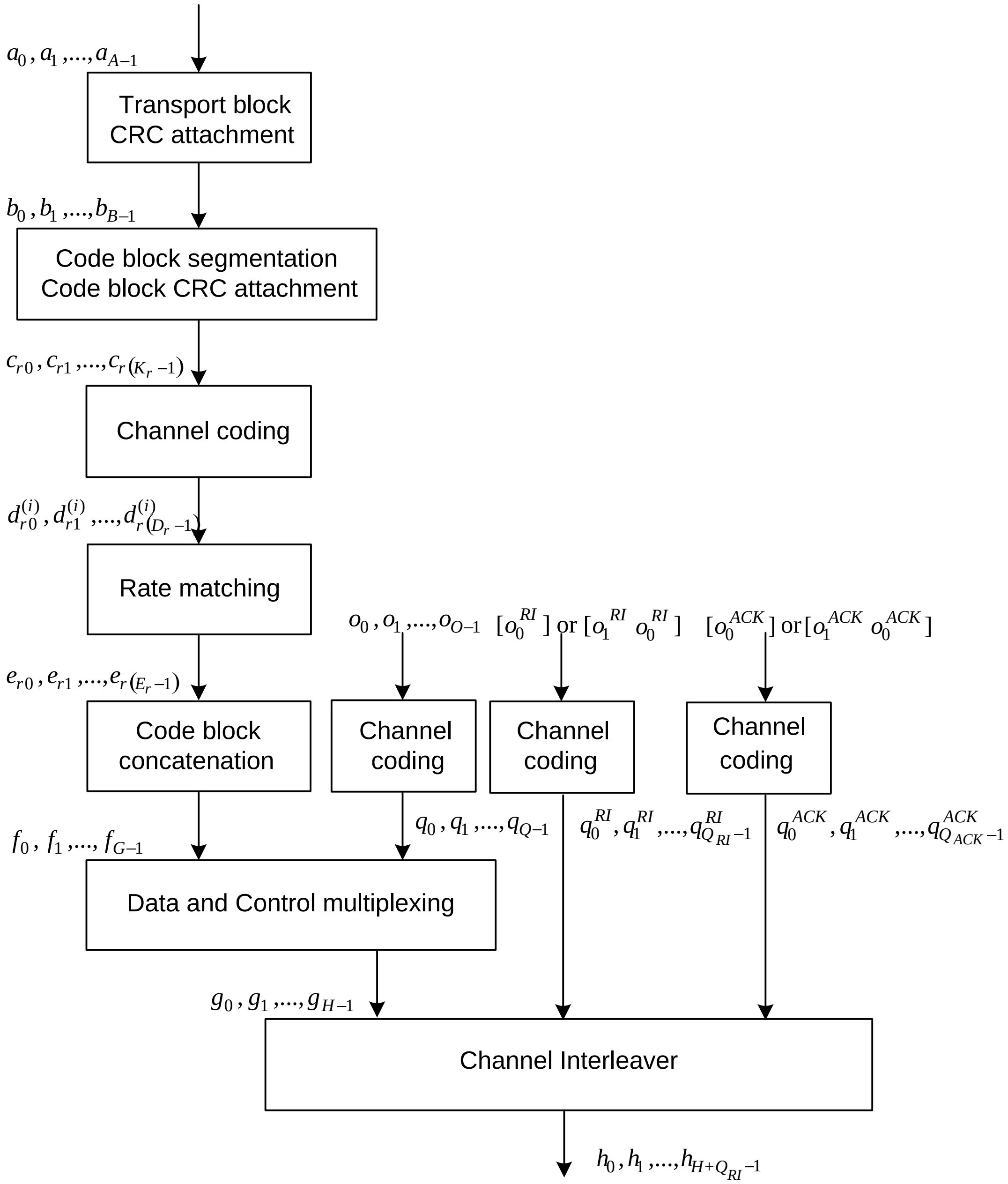

Figure 5.2.2-1: Transport channel processing for UL-SCH

5.2.2.1 Transport block CRC attachment

Error detection is provided on UL-SCH transport blocks through a Cyclic Redundancy Check (CRC).

The entire transport block is used to calculate the CRC parity bits. Denote the bits in a transport block delivered to layer 1 by $a_0, a_1, a_2, a_3, ..., a_{A-1}$, and the parity bits by $p_0, p_1, p_2, p_3, ..., p_{L-1}$. A is the size of the transport block and L is the number of parity bits.

The parity bits are computed and attached to the UL-SCH transport block according to subclause 5.1.1 setting L to 24 bits and using the generator polynomial $g_{CRC24A}(D)$.

5.2.2.2 Code block segmentation and code block CRC attachment

The bits input to the code block segmentation are denoted by $b_0, b_1, b_2, b_3, ..., b_{B-1}$ where B is the number of bits in the transport block (including CRC).

Code block segmentation and code block CRC attachment are performed according to subclause 5.1.2.

The bits after code block segmentation are denoted by $c_{r0}, c_{r1}, c_{r2}, c_{r3}, ..., c_{r(K_r-1)}$, where r is the code block number and K_r is the number of bits for code block number r.

5.2.2.3 Channel coding of UL-SCH

Code blocks are delivered to the channel coding block. The bits in a code block are denoted by $c_{r0}, c_{r1}, c_{r2}, c_{r3}, ..., c_{r(K_r-1)}$, where r is the code block number, and K_r is the number of bits in code block number r. The total number of code blocks is denoted by C and each code block is individually turbo encoded according to subclause 5.1.3.2.

After encoding the bits are denoted by $d_{r0}^{(i)}, d_{r1}^{(i)}, d_{r2}^{(i)}, d_{r3}^{(i)}, ..., d_{r(D_r-1)}^{(i)}$, with $i = 0,1,$ and 2 and where D_r is the number of bits on the i-th coded stream for code block number r, i.e. $D_r = K_r + 4$.

5.2.2.4 Rate matching

Turbo coded blocks are delivered to the rate matching block. They are denoted by $d_{r0}^{(i)}, d_{r1}^{(i)}, d_{r2}^{(i)}, d_{r3}^{(i)}, ..., d_{r(D_r-1)}^{(i)}$, with $i = 0,1,$ and 2, and where r is the code block number, i is the coded stream index, and D_r is the number of bits in each coded stream of code block number r. The total number of code blocks is denoted by C and each coded block is individually rate matched according to subclause 5.1.4.1.

After rate matching, the bits are denoted by $e_{r0}, e_{r1}, e_{r2}, e_{r3}, ..., e_{r(E_r-1)}$, where r is the coded block number, and where E_r is the number of rate matched bits for code block number r.

5.2.2.5 Code block concatenation

The bits input to the code block concatenation block are denoted by $e_{r0}, e_{r1}, e_{r2}, e_{r3}, ..., e_{r(E_r-1)}$ for $r = 0,..., C-1$ and where E_r is the number of rate matched bits for the r-th code block.

Code block concatenation is performed according to subclause 5.1.5.

The bits after code block concatenation are denoted by $f_0, f_1, f_2, f_3, ..., f_{G-1}$, where G is the total number of coded bits for transmission excluding the bits used for control transmission, when control information is multiplexed with the UL-SCH transmission.

5.2.2.6 Channel coding of control information

Control data arrives at the coding unit in the form of channel quality information (CQI and/or PMI), HARQ-ACK and rank indication. Different coding rates for the control information are achieved by allocating different number of coded symbols for its transmission. When control data are transmitted in the PUSCH, the channel coding for HARQ-ACK, rank indication and channel quality information $o_0, o_1, o_2, ..., o_{O-1}$ is done independently.

The number of coded symbols for HARQ-ACK and rank indicator is determined by

$$Q' = \left\lceil \frac{\frac{O}{Q_m R}}{10^{\frac{-\Delta_{offset}^{PUSCH}}{10}}} \right\rceil$$

where O is the number of ACK/NACK bits or rank indicator bits and R is the code rate given by,

$$R = \frac{\sum_{r=0}^{C-1} K_r}{Q_m \cdot M_{\mathrm{sc}}^{\mathrm{PUSCH}} \cdot N_{\mathrm{symb}}^{\mathrm{PUSCH}}}$$

where $M_{\mathrm{sc}}^{\mathrm{PUSCH}}$ is the scheduled bandwidth for uplink transmission, expressed as a number of subcarriers in [2].

For HARQ-ACK information $Q_{ACK} = Q'$ and [$\Delta_{offset}^{PUSCH} = \Delta_{offset}^{HARQ-ACK}$], where $\Delta_{offset}^{HARQ-ACK}$ is signalled by higher layer.

For rank indication $Q_{RI} = Q'$ and [$\Delta_{offset}^{PUSCH} = \Delta_{offset}^{RI}$], where Δ_{offset}^{RI} is signalled by higher layer.

For HARQ-ACK information

- If HARQ-ACK consists of 1-bit of information, i.e., $[o_0^{ACK}]$, it is first encoded according to Table 5.2.2-1.
- If HARQ-ACK consists of 2-bits of information, i.e., $[o_0^{ACK}\ o_1^{ACK}]$, it is first encoded according to Table 5.2.2-2 where $o_2^{ACK} = (o_0^{ACK} \oplus o_1^{ACK})$ and where '$\oplus$' represents XOR operation.

Table 5.2.2-1: Encoding of 1-bit HARQ-ACK

Q_m	Encoded HARQ-ACK
2	$[o_0^{ACK}\ \mathrm{x}]$
4	$[o_0^{ACK}\ \mathrm{x\ x\ x}]$
6	$[o_0^{ACK}\ \mathrm{x\ x\ x\ x\ x}]$

Table 5.2.2-2: Encoding of 2-bit HARQ-ACK

Q_m	Encoded HARQ-ACK
2	$[o_0^{ACK}\ o_1^{ACK}\ o_2^{ACK}\ o_0^{ACK}\ o_1^{ACK}\ o_2^{ACK}]$
4	$[o_0^{ACK}\ o_1^{ACK}\ \mathrm{x\ x}\ o_2^{ACK}\ o_0^{ACK}\ \mathrm{x\ x}\ o_1^{ACK}\ o_2^{ACK}\ \mathrm{x\ x}]$
6	$[o_0^{ACK}\ o_1^{ACK}\ \mathrm{x\ x\ x\ x}\ o_2^{ACK}\ o_0^{ACK}\ \mathrm{x\ x\ x\ x}\ o_1^{ACK}\ o_2^{ACK}\ \mathrm{x\ x\ x\ x}]$

The "x" in Table 5.2.2-1 and 5.2.2-2 are placeholders for [2] to scramble the HARQ-ACK bits in a way that maximizes the Euclidean distance of the modulation symbols carrying HARQ-ACK information.

The bit sequence $q_0^{ACK}, q_1^{ACK}, q_2^{ACK}, ..., q_{Q_{ACK}-1}^{ACK}$ is obtained by concatenation of multiple encoded HARQ-ACK blocks where Q_{ACK} is the total number of coded bit for all the encoded HARQ-ACK blocks. The last concatenation of the encoded HARQ-ACK block may be partial so that the total bit sequence length is equal to Q_{ACK}. The vector sequence output of the channel coding for HARQ-ACK information is denoted by $\underline{q}_0^{ACK}, \underline{q}_1^{ACK}, ..., \underline{q}_{Q'_{ACK}-1}^{ACK}$, where $Q'_{ACK} = Q_{ACK} / Q_m$, and is obtained as follows:

Set i ,k to 0

while $i < Q_{ACK}$

$\underline{q}_k^{ACK} = [q_i^{ACK} \ldots q_{i+Q_m-1}^{ACK}]^T$

$i = i + Q_m$

$k = k + 1$

end while

For rank indication (RI)

- If RI consists of 1-bit of information, i.e., $[o_0^{RI}]$, it is first encoded according to Table 5.2.2-3.
- If RI consists of 2-bits of information, i.e., $[o_0^{RI} \ o_1^{RI}]$, it is first encoded according to Table 5.2.2-4 where $o_2^{RI} = (o_0^{RI} + o_1^{RI}) \bmod 2$.

Table 5.2.2-3: Encoding of 1-bit RI

Q_m	Encoded RI
2	$[o_0^{RI} \ \mathrm{x}]$
4	$[o_0^{RI} \ \mathrm{x\ x\ x}]$
6	$[o_0^{RI} \ \mathrm{x\ x\ x\ x\ x}]$

Table 5.2.2-4: Encoding of 2-bit RI

Q_m	Encoded RI
2	$[o_0^{RI} \ o_1^{RI} \ o_2^{RI} \ o_0^{RI} \ o_1^{RI} \ o_2^{RI}]$
4	$[o_0^{RI} \ o_1^{RI} \ \mathrm{x\ x} \ o_2^{RI} \ o_0^{RI} \ \mathrm{x\ x} \ o_1^{RI} \ o_2^{RI} \ \mathrm{x\ x}]$
6	$[o_0^{RI} \ o_1^{RI} \ \mathrm{x\ x\ x\ x} \ o_2^{RI} \ o_0^{RI} \ \mathrm{x\ x\ x\ x} \ o_1^{RI} \ o_2^{RI} \ \mathrm{x\ x\ x\ x}]$

The “x” in Table 5.2.2-3 and 5.2.2-4 are placeholders for [2] to scramble the RI bits in a way that maximizes the Euclidean distance of the modulation symbols carrying rank information.

The bit sequence $q_0^{RI}, q_1^{RI}, q_2^{RI}, \ldots, q_{Q_{RI}-1}^{RI}$ is obtained by concatenation of multiple encoded RI blocks where Q_{RI} is the total number of coded bit for all the encoded RI blocks. The last concatenation of the encoded RI block may be partial so that the total bit sequence length is equal to Q_{RI}. The vector sequence output of the channel coding for rank information is denoted by $\underline{q}_0^{RI}, \underline{q}_1^{RI}, \ldots, \underline{q}_{Q'_{RI}-1}^{RI}$, where $Q'_{RI} = Q_{RI} / Q_m$, and is obtained as follows:

Set i ,k to 0

while $i < Q_{RI}$

$\underline{q}_k^{RI} = [q_i^{RI} \ldots q_{i+Q_m-1}^{RI}]^T$

$i = i + Q_m$

$k = k + 1$

end while

For channel quality control information (CQI and/or PMI)

The number of coded symbols for channel quality information is determined by

$$Q = \left\lceil \frac{\dfrac{O}{Q_m R}}{10^{\frac{-\Delta_{offset}^{PUSCH}}{10}}} \right\rceil$$

where O is the number of CQI and CRC bits, and [$\Delta_{offset}^{PUSCH} = \Delta_{offset}^{CQI}$], where Δ_{offset}^{CQI} is signalled by higher layer.

- If the payload size is less than or equal to 11 bits, the channel coding of the channel quality information is performed according to subclause 5.2.2.6.4 with input sequence $o_0, o_1, o_2, ..., o_{O-1}$.
- For payload sizes greater than 11 bits, the CRC attachment, channel coding and rate matching of the channel quality information is performed according to subclauses 5.1.1, 5.1.3.1 and 5.1.4.2, respectively. The input bit sequence to the CRC attachment is $o_0, o_1, o_2, ..., o_{O-1}$ and the CRC length is $L = 8$. The output bit sequence of the CRC attachment operation is the input bit sequence to the channel coding operation. The output bit sequence of the channel coding operation is the input bit sequence to the rate matching operation.

The output sequence for the channel coding of channel quality information is denoted by $q_0, q_1, q_2, q_3, ..., q_{Q-1}$.

5.2.2.6.1 Channel quality information formats for wideband CQI reports

Table 5.2.2.6.1-1 shows the fields and the corresponding bit widths for the channel quality information feedback for wideband reports for PDSCH transmissions over closed-loop spatial multiplexing. N in Table 5.2.2.6.1-1 is defined in subclause 7.2 [3].

Table 5.2.2.6.1-1: Fields for channel quality information (CQI) feedback for wideband CQI reports (closed loop spatial multiplexing PDSCH transmission)

Field	Bitwidth			
	2 antenna ports		4 antenna ports	
	Rank = 1	Rank = 2	Rank = 1	Rank > 1
Wideband CQI codeword 0	4	4	4	4
Wideband CQI codeword 1	0	4	0	4
Precoding matrix indication	$2N$	N	$4N$	$4N$

The channel quality bits in Table 5.2.2.6.1-1 form the bit sequence $o_0, o_1, o_2, ..., o_{O-1}$ with o_0 corresponding to the first bit of the first field in the table, o_1 corresponding to the second bit of the first field in the table, and o_{O-1} corresponding to the last bit in the last field in the table. The field of PMI shall be in the increasing order of the subband index [3]. The first bit of each field corresponds to MSB and the last bit LSB.

5.2.2.6.2 Channel quality information formats for higher layer configured subband CQI reports

Table 5.2.2.6.2-1 shows the fields and the corresponding bit widths for the channel quality information feedback for higher layer configured report for PDSCH transmissions over single antenna port, transmit diversity and open loop spatial multiplexing. N in Table 5.2.2.6.2-1 is defined in subclause 7.2 [3].

Table 5.2.2.6.2-1: Fields for channel quality information (CQI) feedback for higher layer configured subband CQI reports (single antenna port, transmit diversity and open loop spatial multiplexing PDSCH transmission)

Field	Bitwidth
Wide-band CQI codeword	4
Subband differential CQI	$2N$

Table 5.2.2.6.2-2 shows the fields and the corresponding bit widths for the channel quality information feedback for

higher layer configured report for PDSCH transmissions over closed loop spatial multiplexing. N in Table 5.2.2.6.2-2 is defined in subclause 7.2 [3].

Table 5.2.2.6.2-2: Fields for channel quality information (CQI) feedback for higher layer configured subband CQI reports (closed loop spatial multiplexing PDSCH transmission)

Field	Bitwidth			
	2 antenna ports		4 antenna ports	
	Rank = 1	Rank = 2	Rank = 1	Rank > 1
Wide-band CQI codeword 0	4	4	4	4
Subband differential CQI codeword 0	$2N$	$2N$	$2N$	$2N$
Wide-band CQI codeword 1	0	4	0	4
Subband differential CQI codeword 1	0	$2N$	0	$2N$
Precoding matrix indication	2	1	4	4

The channel quality bits in Table 5.2.2.6.2-1 through Table 5.2.2.6.2-2 form the bit sequence $o_0, o_1, o_2, ..., o_{O-1}$ with o_0 corresponding to the first bit of the first field in each of the tables, o_1 corresponding to the second bit of the first field in each of the tables, and o_{O-1} corresponding to the last bit in the last field in each of the tables. The field of the PMI and subband differential CQI shall be in the increasing order of the subband index [3]. The first bit of each field corresponds to MSB and the last bit LSB.

5.2.2.6.3 Channel quality information formats for UE selected subband CQI reports

Table 5.2.2.6.3-1 shows the fields and the corresponding bit widths for the channel quality information feedback for UE selected subband CQI for PDSCH transmissions over single antenna port, transmit diversity and open loop spatial multiplexing. L in Table 5.2.2.6.3-1 is defined in subclause 7.2 [3].

Table 5.2.2.6.3-1: Fields for channel quality information (CQI) feedback for UE selected subband CQI reports (single antenna port, transmit diversity and open loop spatial multiplexing PDSCH transmission)

Field	Bitwidth
Wide-band CQI codeword	4
Subband differential CQI	2
Position of the M selected subbands	L

Table 5.2.2.6.3-2 shows the fields and the corresponding bit widths for the channel quality information feedback for UE selected subband CQI for PDSCH transmissions over closed loop spatial multiplexing. L in Table 5.2.2.6.3-2 is defined in subclause 7.2 [3].

Table 5.2.2.6.3-2: Fields for channel quality information (CQI) feedback for UE selected subband CQI reports (closed loop spatial multiplexing PDSCH transmission)

Field	Bitwidth			
	2 antenna ports		4 antenna ports	
	Rank = 1	Rank = 2	Rank = 1	Rank > 1
Wide-band CQI codeword 0	4	4	4	4
Subband differential CQI codeword 0	2	2	2	2
Wide-band CQI codeword 1	0	4	0	4
Subband differential CQI codeword 1	0	2	0	2
Position of the M selected subbands	L	L	L	L
Precoding matrix indication	4	2	8	8

The channel quality bits in Table 5.2.2.6.3-1 through Table 5.2.2.6.3-2 form the bit sequence $o_0, o_1, o_2, ..., o_{O-1}$ with o_0 corresponding to the first bit of the first field in each of the tables, o_1 corresponding to the second bit of the first field in each of the tables, and o_{O-1} corresponding to the last bit in the last field in each of the tables. The field of PMI shall be in the increasing order of the subband index [3], wideband PMI followed by the PMI for the M selected subband. The first bit of each field corresponds to MSB and the last bit LSB.

5.2.2.6.4 Channel coding for CQI/PMI information in PUSCH

The channel quality bits input to the channel coding block are denoted by $o_0, o_1, o_2, o_3, ..., o_{O-1}$ where O is the number of bits. The number of channel quality bits depends on the transmission format as indicated in subclause 5.2.3.3.1 for wideband reports and in subclause 5.2.3.3.2 for UE-selected subbands reports.

The channel quality indication is first coded using a (32, O) block code. The code words of the (32, O) block code are a linear combination of the 11 basis sequences denoted $M_{i,n}$ and defined in Table 5.2.2.6.4-1.

Table 5.2.2.6.4-1: Basis sequences for (32, *O*) code

i	$M_{i,0}$	$M_{i,1}$	$M_{i,2}$	$M_{i,3}$	$M_{i,4}$	$M_{i,5}$	$M_{i,6}$	$M_{i,7}$	$M_{i,8}$	$M_{i,9}$	$M_{i,10}$
0	1	1	0	0	0	0	0	0	0	0	1
1	1	1	1	0	0	0	0	0	0	1	1
2	1	0	0	1	0	0	1	0	1	1	1
3	1	0	1	1	0	0	0	0	1	0	1
4	1	1	1	1	0	0	0	1	0	0	1
5	1	1	0	0	1	0	1	1	1	0	1
6	1	0	1	0	1	0	1	0	1	1	1
7	1	0	0	1	1	0	0	1	1	0	1
8	1	1	0	1	1	0	0	1	0	1	1
9	1	0	1	1	1	0	1	0	0	1	1
10	1	0	1	0	0	1	1	1	0	1	1
11	1	1	1	0	0	1	1	0	1	0	1
12	1	0	0	1	0	1	0	1	1	1	1
13	1	1	0	1	0	1	0	1	0	1	1
14	1	0	0	0	1	1	0	1	0	0	1
15	1	1	0	0	1	1	1	1	0	1	1
16	1	1	1	0	1	1	1	0	0	1	0
17	1	0	0	1	1	1	0	0	1	0	0
18	1	1	0	1	1	1	1	1	0	0	0
19	1	0	0	0	0	1	1	0	0	0	0
20	1	0	1	0	0	0	1	0	0	0	1
21	1	1	0	1	0	0	0	0	0	1	1
22	1	0	0	0	1	0	0	1	1	0	1
23	1	1	1	0	1	0	0	0	1	1	1
24	1	1	1	1	1	0	1	1	1	1	0
25	1	1	0	0	0	1	1	1	0	0	1
26	1	0	1	1	0	1	0	0	1	1	0
27	1	1	1	1	0	1	0	1	1	1	0
28	1	0	1	0	1	1	1	0	1	0	0
29	1	0	1	1	1	1	1	1	1	0	0
30	1	1	1	1	1	1	1	1	1	1	1
31	1	0	0	0	0	0	0	0	0	0	0

The encoded CQI/PMI block is denoted by $b_0, b_1, b_2, b_3, ..., b_{B-1}$ where $B = 32$ and

$$b_i = \sum_{n=0}^{O-1} \left(o_n \cdot M_{i,n}\right) \bmod 2 \quad \text{where } i = 0, 1, 2, \ldots, B\text{-1}.$$

The output bit sequence $q_0, q_1, q_2, q_3, ..., q_{Q-1}$ is obtained by circular repetition of the encoded CQI/PMI block as follows

$$q_i = b_{(i \bmod B)} \quad \text{where } i = 0, 1, 2, \ldots, Q\text{-1}.$$

5.2.2.7 Data and control multiplexing

The control and data multiplexing is performed such that HARQ-ACK information is present on both slots and is mapped to resources around the demodulation reference signals. In addition, the multiplexing ensures that control and data information are mapped to different modulation symbols.

The inputs to the data and control multiplexing are the coded bits of the control information denoted by $q_0, q_1, q_2, q_3,..., q_{Q-1}$ and the coded bits of the UL-SCH denoted by $f_0, f_1, f_2, f_3,..., f_{G-1}$. The output of the data and control multiplexing operation is denoted by $\underline{g}_0, \underline{g}_1, \underline{g}_2, \underline{g}_3,..., \underline{g}_{H'-1}$, where $H = (G+Q)$ and $H' = H / Q_m$, and where $\underline{g}_i$, $i = 0,..., H'-1$ are column vectors of length Q_m. H is the total number of coded bits allocated for UL-SCH data and CQI/PMI data.

Denote the number of SC-FDMA symbols per subframe for PUSCH transmission by $N_{\text{symb}}^{\text{PUSCH}} = \left(2 \cdot \left(N_{\text{symb}}^{\text{UL}} - 1\right) - N_{SRS}\right)$.

The control information and the data shall be multiplexed as follows:

Set *i*, *j*, *k* to 0

while $j < Q$ -- first place the control information

$\underline{g}_k = [q_j \; ... \; q_{j+Q_m-1}]^T$

$j = j + Q_m$

$k = k + 1$

end while

while $i < G$ -- then place the data

$\underline{g}_k = [f_i \; ... \; f_{i+Q_m-1}]^T$

$i = i + Q_m$

$k = k + 1$

end while

5.2.2.8 Channel interleaver

The channel interleaver described in this subclause in conjunction with the resource element mapping for PUSCH in [2] implements a time-first mapping of modulation symbols onto the transmit waveform while ensuring that the HARQ-ACK information is present on both slots in the subframe and is mapped to resources around the uplink demodulation reference signals.

The input to the channel interleaver are denoted by $\underline{g}_0, \underline{g}_1, \underline{g}_2,..., \underline{g}_{H'-1}$, $\underline{q}_0^{RI}, \underline{q}_1^{RI}, \underline{q}_2^{RI},..., \underline{q}_{Q'_{RI}-1}^{RI}$ and $\underline{q}_0^{ACK}, \underline{q}_1^{ACK}, \underline{q}_2^{ACK},..., \underline{q}_{Q'_{ACK}-1}^{ACK}$. The number of modulation symbols in the subframe is given by $H'' = H' + Q'_{RI}$ .. The output bit sequence from the channel interleaver is derived as follows:

(1) Assign $C_{mux} = N_{\text{symb}}^{\text{PUSCH}}$ to be the number of columns of the matrix. The columns of the matrix are numbered 0, 1, 2,…, $C_{mux} - 1$ from left to right.

(2) The number of rows of the matrix is $R_{mux} = \left(H'' \cdot Q_m\right) / C_{mux}$ and we define $R'_{mux} = R_{mux} / Q_m$.

The rows of the rectangular matrix are numbered 0, 1, 2,…, $R_{mux} - 1$ from top to bottom.

(3) If rank information is transmitted in this subframe, the vector sequence $\underline{q}_0^{RI}, \underline{q}_1^{RI}, \underline{q}_2^{RI}, ..., \underline{q}_{Q'_{RI}-1}^{RI}$ is written onto the columns indicated by Table 5.2.2.8-1, and by sets of Q_m rows starting from the last row and moving upwards according to the following pseudocode.

Set i, j to 0.

Set r to $R'_{mux} - 1$

while $i < Q'_{RI}$

$c_{RI} = \text{Column Set}(j)$

$\underline{y}_{r \times C_{mux} + c_{RI}} = \underline{q}_i^{RI}$

$i = i + 1$

$r = R'_{mux} - 1 - \lfloor i/4 \rfloor$

$j = (j+3) \bmod 4$

end while

(4) Write the input vector sequence, i.e., $\underline{y}_k = \underline{g}_k$ for k = 0, 1,…, $H'-1$, into the $(R_{mux} \times C_{mux})$ matrix by sets of Q_m rows starting with the vector $\underline{y}_0$ in column 0 and rows 0 to $(Q_m - 1)$ and skipping the matrix entries that are already occupied:

$$\begin{bmatrix} \underline{y}_0 & \underline{y}_1 & \underline{y}_2 & \cdots & \underline{y}_{C_{mux}-1} \\ \underline{y}_{C_{mux}} & \underline{y}_{C_{mux}+1} & \underline{y}_{C_{mux}+2} & \cdots & \underline{y}_{2C_{mux}-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \underline{y}_{(R'_{mux}-1)\times C_{mux}} & \underline{y}_{(R'_{mux}-1)\times C_{mux}+1} & \underline{y}_{(R'_{mux}-1)\times C_{mux}+2} & \cdots & \underline{y}_{(R'_{mux}\times C_{mux}-1)} \end{bmatrix}$$

(5) If HARQ-ACK information is transmitted in this subframe, the vector sequence $\underline{q}_0^{ACK}, \underline{q}_1^{ACK}, \underline{q}_2^{ACK}, ..., \underline{q}_{Q'_{ACK}-1}^{ACK}$ is written onto the columns indicated by Table 5.2.2.8-2, and by sets of Q_m rows starting from the last row and moving upwards according to the following pseudocode. Note that this operation overwrites some of the channel interleaver entries obtained in step (4).

Set i, j to 0.

Set r to $R'_{mux} - 1$

while $i < Q'_{ACK}$

$c_{ACK} = \text{ColumnSet}(j)$

$\underline{y}_{r \times C_{mux} + c_{ACK}} = \underline{q}_i^{ACK}$

$i = i + 1$

$r = R'_{mux} - 1 - \lfloor i/4 \rfloor$

$j = (j+3) \bmod 4$

end while

Where ColumnSet is given in Table 5.2.2.8-1 and indexed left to right from 0 to 3.

(6) The output of the block interleaver is the bit sequence read out column by column from the $(R_{mux} \times C_{mux})$ matrix. The bits after channel interleaving are denoted by $h_0, h_1, h_2, ..., h_{H+Q_{RI}-1}$.

Table 5.2.2.8-1: Column set for Insertion of rank information

CP configuration	Column Set
Normal	{1, 4, 7, 10}
Extended	{0, 3, 5, 8}

Table 5.2.2.8-2: Column set for Insertion of HARQ-ACK information

CP configuration	Column Set
Normal	{2, 3, 8, 9}
Extended	{1, 2, 6, 7}

5.2.3 Uplink control information on PUCCH

Data arrives to the coding unit in form of indicators for measurement indication, scheduling request and HARQ acknowledgement.

Three forms of channel coding are used, one for the channel quality information (CQI), another for HARQ-ACK (acknowledgement) and scheduling request and another for combination of channel quality information (CQI) and HARQ-ACK.

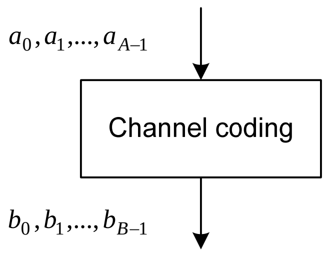

Figure 5.2.3-1: Processing for UCI

5.2.3.1 Channel coding for UCI HARQ-ACK

The HARQ acknowledgement bits are received from higher layers. Each positive acknowledgement (ACK) is encoded as a binary '0' and each negative acknowledgement (NAK) is encoded as a binary '1'. The HARQ-ACK bits are processed according to [2].

5.2.3.2 Channel coding for UCI scheduling request

The scheduling request indication is received from higher layers and is processed according to [2].

5.2.3.3 Channel coding for UCI channel quality information

The channel quality bits input to the channel coding block are denoted by $a_0, a_1, a_2, a_3, ..., a_{A-1}$ where A is the number of bits. The number of channel quality bits depends on the transmission format as indicated in subclause 5.2.3.3.1 for wideband reports and in subclause 5.2.3.3.2 for UE-selected subbands reports.

The channel quality indication is coded using a (20, *A*) code. The code words of the (20, *A*) code are a linear combination of the 13 basis sequences denoted $M_{i,n}$ and defined in Table 5.2.3.3-1.

Table 5.2.3.3-1: Basis sequences for (20, *A*) code

i	$M_{i,0}$	$M_{i,1}$	$M_{i,2}$	$M_{i,3}$	$M_{i,4}$	$M_{i,5}$	$M_{i,6}$	$M_{i,7}$	$M_{i,8}$	$M_{i,9}$	$M_{i,10}$	$M_{i,11}$	$M_{i,12}$
0	1	1	0	0	0	0	0	0	0	0	1	1	0
1	1	1	1	0	0	0	0	0	0	1	1	1	0
2	1	0	0	1	0	0	1	0	1	1	1	1	1
3	1	0	1	1	0	0	0	0	1	0	1	1	1
4	1	1	1	1	0	0	0	1	0	0	1	1	1
5	1	1	0	0	1	0	1	1	1	0	1	1	1
6	1	0	1	0	1	0	1	0	1	1	1	1	1
7	1	0	0	1	1	0	0	1	1	0	1	1	1
8	1	1	0	1	1	0	0	1	0	1	1	1	1
9	1	0	1	1	1	0	1	0	0	1	1	1	1
10	1	0	1	0	0	1	1	1	0	1	1	1	1
11	1	1	1	0	0	1	1	0	1	0	1	1	1
12	1	0	0	1	0	1	0	1	1	1	1	1	1
13	1	1	0	1	0	1	0	1	0	1	1	1	1
14	1	0	0	0	1	1	0	1	0	0	1	0	1
15	1	1	0	0	1	1	1	1	0	1	1	0	1
16	1	1	1	0	1	1	1	0	0	1	0	1	1
17	1	0	0	1	1	1	0	0	1	0	0	1	1
18	1	1	0	1	1	1	1	1	0	0	0	0	0
19	1	0	0	0	0	1	1	0	0	0	0	0	0

After encoding the bits are denoted by $b_0, b_1, b_2, b_3, ..., b_{B-1}$ where $B = 20$ and with

$$b_i = \sum_{n=0}^{A-1} \left(a_n \cdot M_{i,n} \right) \bmod 2$$ where *i* = 0, 1, 2, …, *B*-1.

5.2.3.3.1 Channel quality information formats for wideband reports

Table 5.2.3.3.1-1 shows the fields and the corresponding bit widths for the channel quality information feedback for wideband reports for PDSCH transmissions over a single antenna port, transmit diversity or with open loop spatial multiplexing.

Table 5.2.3.3.1-1: UCI fields for channel quality information (CQI) feedback for wideband reports (single antenna port, transmit diversity or open loop spatial multiplexing PDSCH transmission)

Field	Bitwidth
Wide-band CQI	4

Table 5.2.3.3.1-2 shows the fields and the corresponding bit widths for the channel quality and precoding matrix information feedback for wideband reports for PDSCH transmissions with closed loop spatial multiplexing.

Table 5.2.3.3.1-2: UCI fields for channel quality and precoding information (CQI/PMI) feedback for wideband reports (closed loop spatial multiplexing PDSCH transmission)

Field	Bitwidths			
	2 antenna ports		4 antenna ports	
	Rank = 1	Rank = 2	Rank = 1	Rank > 1
Wide-band CQI	4	4	4	4
Spatial differential CQI	0	3	0	3
Precoding matrix indication	2	1	4	4

Table 5.2.3.3.1-3 shows the fields and the corresponding bit widths for the rank indication feedback for wideband reports for PDSCH transmissions for open and closed loop spatial multiplexing.

Table 5.2.3.3.1-3: UCI fields for rank indication (RI) feedback for wideband reports

Field	Bitwidths		
	2 antenna ports	4 antenna ports	
		Max 2 layers	Max 4 layers
Rank indication	1	1	2

The channel quality bits in Table 5.2.3.3.1-1 through Table 5.2.3.3.1-3 form the bit sequence $a_0, a_1, a_2, a_3, ..., a_{A-1}$ with a_0 corresponding to the first bit of the first field in each of the tables, a_1 corresponding to the second bit of the first field in each of the tables, and a_{A-1} corresponding to the last bit in the last field in each of the tables. The first bit of each field corresponds to MSB and the last bit LSB.

5.2.3.3.2 Channel quality information formats for UE-selected sub-band reports

Table 5.2.3.3.2-1 shows the fields and the corresponding bit widths for the sub-band channel quality information feedback for UE-selected sub-band reports for PDSCH transmissions over a single antenna port, transmit diversity or with open loop spatial multiplexing.

Table 5.2.3.3.2-1: UCI fields for channel quality information (CQI) feedback for UE-selected sub-band reports (single antenna port, transmit diversity or open loop spatial multiplexing PDSCH transmission)

Field	Bitwidth
Sub-band CQI	4
Sub-band label	1 or 2

Table 5.2.3.3.2-2 shows the fields and the corresponding bit widths for the sub-band channel quality information feedback for UE-selected sub-band reports for PDSCH transmissions with closed loop spatial multiplexing.

Table 5.2.3.3.2-2: UCI fields for channel quality information (CQI) feedback for UE-selected sub-band reports (closed loop spatial multiplexing PDSCH transmission)

Field	Bitwidths			
	2 antenna ports		4 antenna ports	
	Rank = 1	Rank = 2	Rank = 1	Rank > 1
Sub-band CQI	4	4	4	4
Spatial differential CQI	0	3	0	3
Sub-band label	1 or 2	1 or 2	1 or 2	1 or 2

Table 5.2.3.3.2-3 shows the fields and the corresponding bit widths for the wide-band channel quality and precoding matrix information feedback for UE-selected sub-band reports for PDSCH transmissions with closed loop spatial multiplexing.

Table 5.2.3.3.2-3: UCI fields for channel quality and precoding information (CQI/PMI) feedback for UE-selected sub-band reports (closed loop spatial multiplexing PDSCH transmission)

Field	Bitwidths			
	2 antenna ports		4 antenna ports	
	Rank = 1	Rank = 2	Rank = 1	Rank > 1
Wide-band CQI	4	4	4	4
Spatial differential CQI	0	3	0	3
Precoding matrix indication	2	1	4	4

Table 5.2.3.3.2-4 shows the fields and the corresponding bit widths for the rank indication feedback for UE-selected sub-band reports for PDSCH transmissions for open and closed loop spatial multiplexing.

Table 5.2.3.3.2-4: UCI fields for rank indication (RI) feedback for UE-selected sub-band reports

Field	Bitwidths		
	2 antenna ports	4 antenna ports	
		Max 2 layers	Max 4 layers
Rank indication	1	1	2

The channel quality bits in Table 5.2.3.3.2-1 through Table 5.2.3.3.2-4 form the bit sequence $a_0, a_1, a_2, a_3, ..., a_{A-1}$ with a_0 corresponding to the first bit of the first field in each of the tables, a_1 corresponding to the second bit of the first field in each of the tables, and a_{A-1} corresponding to the last bit in the last field in each of the tables. The first bit of each field corresponds to MSB and the last bit LSB.

5.2.3.4 Channel coding for UCI channel quality information and HARQ-ACK

This section defines the channel coding scheme for the simultaneous transmission of channel quality information and HARQ-ACK information in a subframe.

When normal CP is used for uplink transmission, the channel quality information is coded according to subclause 5.2.3.3 with input bit sequence $a'_0, a'_1, a'_2, a'_3, ..., a'_{A'-1}$ and output bit sequence $b'_0, b'_1, b'_2, b'_3, ..., b'_{B'-1}$, where $B' = 20$. The HARQ acknowledgement bits are denoted by a''_0 in case one HARQ acknowledgement bit or a''_0, a''_1 in case two HARQ acknowledgement bits are reported per subframe. Each positive acknowledgement (ACK) is encoded as a binary '0' and each negative acknowledgement (NAK) is encoded as a binary '1'.

The output of this channel coding block for normal CP is denoted by $b_0, b_1, b_2, b_3, ..., b_{B-1}$, where

$$b_i = b'_i, i = 0, ..., B' - 1$$

In case one HARQ acknowledgement bit is reported per subframe:

$$b_{B'} = a''_0 \text{ and } B = (B' + 1)$$

In case two HARQ acknowledgement bits are reported per subframe:

$$b_{B'} = a''_0, b_{B'+1} = a''_1 \text{ and } B = (B' + 2)$$

When extended CP is used for uplink transmission, the channel quality information and the HARQ-ACK acknowledgement bits are jointly coded. The HARQ acknowledgement bits are denoted by a''_0 in case one HARQ acknowledgement bit or $[a''_0, a''_1]$ in case two HARQ acknowledgement bits are reported per subframe.

The channel quality information denoted by $a'_0, a'_1, a'_2, a'_3, ..., a'_{A'-1}$ is multiplexed with the HARQ acknowledgement bits to yield the sequence $a_0, a_1, a_2, a_3, ..., a_{A-1}$ as follows

$$a_i = a'_i, i = 0, ..., A' - 1$$

and

$a_{A'} = a''_0$ and $A = (A' + 1)$ in case one HARQ-acknowledgement bit is reported per subframe, or

$a_{A'} = a''_0$, $a_{(A'+1)} = a''_1$ and $A = (A'+2)$ in case two HARQ-acknowledgement bits are reported per subframe.

The sequence $a_0, a_1, a_2, a_3, ..., a_{A-1}$ is encoded according to section 5.2.3.3 to yield the output bit sequence $b_0, b_1, b_2, b_3, ..., b_{B-1}$ where $B = 20$.

5.2.4 Uplink control information on PUSCH without UL-SCH data

When control data are sent via PUSCH without UL-SCH data, the following coding steps can be identified:

- Channel coding of control information
- Control information mapping
- Channel interleaver

5.2.4.1 Channel coding of control information

Control data arrives at the coding unit in the form of channel quality information (CQI and/or PMI), HARQ-ACK and rank indication. Different coding rates for the control information are achieved by allocating different number of coded symbols for its transmission. The channel coding and rate matching of the control data is performed according to subclause 5.2.2.6. The coded output sequence for channel quality information is denoted by $q_0, q_1, q_2, q_3, ..., q_{Q-1}$, the coded vector sequence output for HARQ-ACK is denoted by $\underline{q}_0^{ACK}, \underline{q}_1^{ACK}, \underline{q}_2^{ACK}, ..., \underline{q}_{Q'_{ACK}-1}^{ACK}$ and the coded vector sequence output for rank indication is denoted by $\underline{q}_0^{RI}, \underline{q}_1^{RI}, \underline{q}_2^{RI}, ..., \underline{q}_{Q'_{RI}-1}^{RI}$.

5.2.4.2 Control information mapping

The input are the coded bits of the channel quality information denoted by $q_0, q_1, q_2, q_3, ..., q_{Q-1}$. The output is denoted by $\underline{g}_0, \underline{g}_1, \underline{g}_2, \underline{g}_3, ..., \underline{g}_{H'-1}$, where $H = Q$ and $H' = H / Q_m$, and where $\underline{g}_i$, $i = 0, ..., H'-1$ are column vectors of length Q_m.

Denote the number of SC-FDMA symbols per subframe for PUSCH transmission by $N_{\text{symb}}^{\text{PUSCH}} = \left(2 \cdot \left(N_{\text{symb}}^{\text{UL}} - 1\right) - N_{SRS}\right)$.

The control information shall be mapped as follows:

Set *j*, *k* to 0

while $j < Q$

$\underline{g}_k = [q_j \; ... \; q_{j+Q_m-1}]^T$

$j = j + Q_m$

$k = k + 1$

end while

5.2.4.3 Channel interleaver

The vector sequences $\underline{g}_0, \underline{g}_1, \underline{g}_2, ..., \underline{g}_{H'-1}$, $\underline{q}_0^{RI}, \underline{q}_1^{RI}, \underline{q}_2^{RI}, ..., \underline{q}_{Q'_{RI}-1}^{RI}$ and $\underline{q}_0^{ACK}, \underline{q}_1^{ACK}, \underline{q}_2^{ACK}, ..., \underline{q}_{Q'_{ACK}-1}^{ACK}$ are channel interleaved according subclause 5.2.2.8. The bits after channel interleaving are denoted by $h_0, h_1, h_2, ..., h_{H+Q_{RI}-1}$.

5.3 Downlink transport channels and control information

5.3.1 Broadcast channel

Figure 5.3.1-1 shows the processing structure for the BCH transport channel. Data arrives to the coding unit in form of a maximum of one transport block every transmission time interval (TTI) of 40ms. The following coding steps can be identified:

- Add CRC to the transport block
- Channel coding
- Rate matching

The coding steps for BCH transport channel are shown in the figure below.

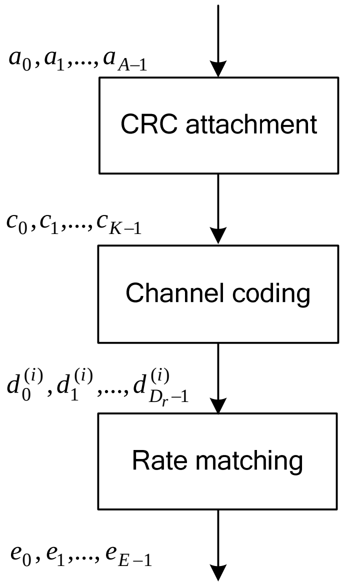

Figure 5.3.1-1: Transport channel processing for BCH

5.3.1.1 Transport block CRC attachment

Error detection is provided on BCH transport blocks through a Cyclic Redundancy Check (CRC).

The entire transport block is used to calculate the CRC parity bits. Denote the bits in a transport block delivered to layer 1 by $a_0, a_1, a_2, a_3, ..., a_{A-1}$, and the parity bits by $p_0, p_1, p_2, p_3, ..., p_{L-1}$. A is the size of the transport block and L is the number of parity bits.

The parity bits are computed and attached to the BCH transport block according to subclause 5.1.1 setting L to 16 bits. After the attachment, the CRC bits are scrambled according to the eNode-B transmit antenna configuration with the sequence $x_{ant,0}, x_{ant,1}, ..., x_{ant,15}$ as indicated in Table 5.3.1.1-1 to form the sequence of bits $c_0, c_1, c_2, c_3, ..., c_{K-1}$ where

$c_k = a_k$ for k = 0, 1, 2, …, A-1

$c_k = \left(p_{k-A} + x_{ant,k-A}\right) \bmod 2$ for k = A, A+1, A+2,..., A+15.

Table 5.3.1.1-1: CRC mask for PBCH

Number of transmit antenna ports at eNode-B	PBCH CRC mask $< x_{ant,0}, x_{ant,1}, ..., x_{ant,15} >$
1	<0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0>
2	<1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1>
4	<0, 1, 0, 1, 0, 1, 0, 1, 0, 1, 0, 1, 0, 1, 0, 1>

5.3.1.2 Channel coding

Information bits are delivered to the channel coding block. They are denoted by $c_0, c_1, c_2, c_3, ..., c_{K-1}$, where K is the number of bits, and they are tail biting convolutionally encoded according to subclause 5.1.3.1.

After encoding the bits are denoted by $d_0^{(i)}, d_1^{(i)}, d_2^{(i)}, d_3^{(i)}, ..., d_{D-1}^{(i)}$, with $i = 0,1,$ and 2 , and where D is the number of bits on the i-th coded stream, i.e., $D = K$.

5.3.1.3 Rate matching

A tail biting convolutionally coded block is delivered to the rate matching block. This block of coded bits is denoted by $d_0^{(i)}, d_1^{(i)}, d_2^{(i)}, d_3^{(i)}, ..., d_{D-1}^{(i)}$, with $i = 0,1,$ and 2 , and where i is the coded stream index and D is the number of bits in each coded stream. This coded block is rate matched according to subclause 5.1.4.2.

After rate matching, the bits are denoted by $e_0, e_1, e_2, e_3, ..., e_{E-1}$, where E is the number of rate matched bits.

5.3.2 Downlink shared channel, Paging channel and Multicast channel

Figure 5.3.2-1 shows the processing structure for the DL-SCH, PCH and MCH transport channels. Data arrives to the coding unit in form of a maximum of one transport block every transmission time interval (TTI). The following coding steps can be identified:

– Add CRC to the transport block

– Code block segmentation and code block CRC attachment

– Channel coding

– Rate matching

– Code block concatenation

The coding steps for DL-SCH, PCH and MCH transport channels are shown in the figure below.

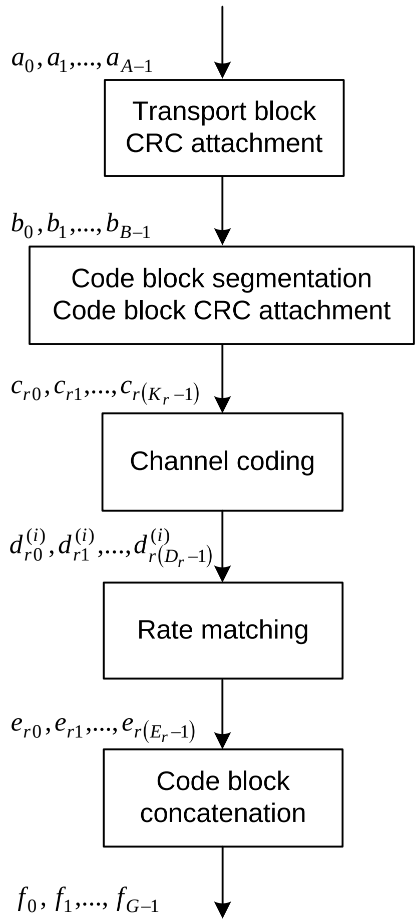

Figure 5.3.2-1: Transport channel processing for DL-SCH, PCH and MCH

5.3.2.1 Transport block CRC attachment

Error detection is provided on transport blocks through a Cyclic Redundancy Check (CRC).

The entire transport block is used to calculate the CRC parity bits. Denote the bits in a transport block delivered to layer 1 by $a_0, a_1, a_2, a_3, ..., a_{A-1}$, and the parity bits by $p_0, p_1, p_2, p_3, ..., p_{L-1}$. A is the size of the transport block and L is the number of parity bits.

The parity bits are computed and attached to the transport block according to subclause 5.1.1 setting L to 24 bits and using the generator polynomial $g_{CRC24A}(D)$.

5.3.2.2 Code block segmentation and code block CRC attachment

The bits input to the code block segmentation are denoted by $b_0, b_1, b_2, b_3, ..., b_{B-1}$ where B is the number of bits in the transport block (including CRC).

Code block segmentation and code block CRC attachment are performed according to subclause 5.1.2.

The bits after code block segmentation are denoted by $c_{r0}, c_{r1}, c_{r2}, c_{r3}, ..., c_{r(K_r-1)}$, where r is the code block number and K_r is the number of bits for code block number r.

5.3.2.3 Channel coding

Code blocks are delivered to the channel coding block. They are denoted by $c_{r0}, c_{r1}, c_{r2}, c_{r3}, ..., c_{r(K_r-1)}$, where r is the code block number, and K_r is the number of bits in code block number r. The total number of code blocks is denoted by C and each code block is individually turbo encoded according to subclause 5.1.3.2.

After encoding the bits are denoted by $d_{r0}^{(i)}, d_{r1}^{(i)}, d_{r2}^{(i)}, d_{r3}^{(i)}, ..., d_{r(D_r-1)}^{(i)}$, with $i = 0,1,$ and 2 , and where D_r is the number of bits on the *i*-th coded stream for code block number *r*, i.e. $D_r = K_r + 4$.

5.3.2.4 Rate matching

Turbo coded blocks are delivered to the rate matching block. They are denoted by $d_{r0}^{(i)}, d_{r1}^{(i)}, d_{r2}^{(i)}, d_{r3}^{(i)}, ..., d_{r(D_r-1)}^{(i)}$, with $i = 0,1,$ and 2 , and where *r* is the code block number, *i* is the coded stream index, and D_r is the number of bits in each coded stream of code block number *r*. The total number of code blocks is denoted by *C* and each coded block is individually rate matched according to subclause 5.1.4.1.

After rate matching, the bits are denoted by $e_{r0}, e_{r1}, e_{r2}, e_{r3}, ..., e_{r(E_r-1)}$, where *r* is the coded block number, and where E_r is the number of rate matched bits for code block number *r*.

5.3.2.5 Code block concatenation

The bits input to the code block concatenation block are denoted by $e_{r0}, e_{r1}, e_{r2}, e_{r3}, ..., e_{r(E_r-1)}$ for $r = 0,..., C-1$ and where E_r is the number of rate matched bits for the *r*-th code block.

Code block concatenation is performed according to subclause 5.1.5.2.

The bits after code block concatenation are denoted by $f_0, f_1, f_2, f_3,..., f_{G-1}$, where *G* is the total number of coded bits for transmission.

5.3.3 Downlink control information

A DCI transports downlink or uplink scheduling information, or uplink power control commands for one MAC ID. The MAC ID is implicitly encoded in the CRC.

Figure 5.3.3-1 shows the processing structure for the DCI. The following coding steps can be identified:

- Information element multiplexing
- CRC attachment
- Channel coding
- Rate matching

The coding steps for DCI are shown in the figure below.

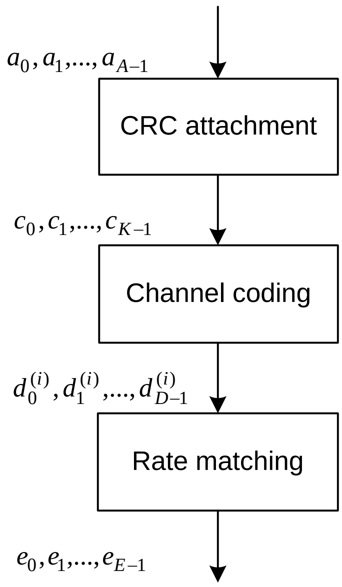

Figure 5.3.3-1: Processing for DCI

5.3.3.1 DCI formats

The ordering of the information bits in the payload description below starts with a_0 and ends with a_{A-1}.

The information fields are multiplexed according to the order they are listed in each DCI format. The first bit of each information field corresponds to MSB.

5.3.3.1.1 Format 0

DCI format 0 is used for the scheduling of PUSCH.

The following information is transmitted by means of the DCI format 0:

- Flag for format0/format1A differentiation – 1 bit
- Hopping flag – 1 bit
- Resource block assignment and hopping resource allocation – $\lceil \log_2(N_{\mathrm{RB}}^{\mathrm{UL}}(N_{\mathrm{RB}}^{\mathrm{UL}}+1)/2) \rceil$ bits
  - For PUSCH hopping:
    - N_{UL_hop} bits are used to obtain the value of $\tilde{n}_{PRB}(i)$ as indicated in subclause [8.4] of [3]
    - $\left(\lceil \log_2(N_{\mathrm{RB}}^{\mathrm{UL}}(N_{\mathrm{RB}}^{\mathrm{UL}}+1)/2) \rceil - N_{\mathrm{UL\_hop}}\right)$ bits provide the resource allocation of the first slot in the UL subframe
  - For non-hopping PUSCH:
    - $\left(\lceil \log_2(N_{\mathrm{RB}}^{\mathrm{UL}}(N_{\mathrm{RB}}^{\mathrm{UL}}+1)/2) \rceil\right)$ bits provide the resource allocation of the first slot in the UL subframe
- Modulation and coding scheme and redundancy version – 5 bits
- New data indicator – 1 bit
- TPC command for scheduled PUSCH – 2 bits
- Cyclic shift for DM RS – 3 bits

- UL index (2 bits, this field just applies to TDD operation)
- CQI request – 1 bit

If the number of information bits in format 0 is less than for format 1A, zeros shall be appended to format 0 until the payload size equals that of format 1A.

5.3.3.1.2 Format 1

DCI format 1 is used for the scheduling of one PDSCH codeword.

The following information is transmitted by means of the DCI format 1:

- Resource allocation header (resource allocation type 0 / type 1) – 1 bit
- Resource block assignment:
  - For resource allocation type 0 [3],
    - $\lceil N_{\mathrm{RB}}^{\mathrm{DL}} / P \rceil$ bits provide the resource allocation
  - For resource allocation type 1 [3],
    - $\lceil \log_2(P) \rceil$ bits of this field are used as a header specific to this resource allocation type to indicate the selected resource blocks subset
    - 1 bit indicates a shift of the resource allocation span
    - $\left(\lceil N_{\mathrm{RB}}^{\mathrm{DL}} / P \rceil - \lceil \log_2(P) \rceil - 1 \right)$ bits provide the resource allocation

  where the value of P depends on the number of DL resource blocks as indicated in subclause [7.1.1] of [3]
- Modulation and coding scheme – 5 bits
- HARQ process number – 3 bits (FDD) , 4 bits (TDD)
- New data indicator – 1 bit
- Redundancy version – 2 bits
- TPC command for PUCCH – 2 bits
- Downlink Assignment Index (this field just applies to TDD operation) – 2 bits

5.3.3.1.3 Format 1A

DCI format 1A is used for the compact scheduling of one PDSCH codeword. DCI format 1A can be used for scheduling PDSCH to UEs configured in any transmission mode.

The following information is transmitted by means of the DCI format 1A:

- Flag for format0/format1A differentiation – 1 bit
- Localized/Distributed VRB assignment flag – 1 bit
- Resource block assignment – $\lceil \log_2(N_{\mathrm{RB}}^{\mathrm{DL}}(N_{\mathrm{RB}}^{\mathrm{DL}}+1)/2) \rceil$ bits
  - For localized VRB:

    $\lceil \log_2(N_{\mathrm{RB}}^{\mathrm{DL}}(N_{\mathrm{RB}}^{\mathrm{DL}}+1)/2) \rceil$ bits provide the resource allocation
  - For distributed VRB:

- For $N_{RB}^{DL} < 50$
  - $\lceil \log_2(N_{RB}^{DL}(N_{RB}^{DL}+1)/2) \rceil$ bits provide the resource allocation
- For $N_{RB}^{DL} \geq 50$
  - 1 bit indicates if $N_{gap} = N_{gap,1}$ or $N_{gap} = N_{gap,2}$
  - $(\lceil \log_2(N_{RB}^{DL}(N_{RB}^{DL}+1)/2) \rceil - 1)$ bits provide the resource allocation

- Modulation and coding scheme – 5bits
- HARQ process number – 3 bits (FDD) , 4 bits (TDD)
- New data indicator – 1 bit
- Redundancy version – 2 bits
- TPC command for PUCCH – 2 bits
- Downlink Assignment Index (this field just applies to TDD operation) – 2 bits

If the number of information bits in format 1A is less than for format 0, zeros shall be appended to format 1A until the payload size equals that of format 0.

5.3.3.1.4 Format 1C

DCI format 1C is used for very compact assignments of DL-SCH assignments. The DL-SCH transmission shall always use QPSK modulation when format 1C is used.

The following information is transmitted by means of the DCI format 1C:

- Resource block assignment [5 bits]
- Modulation and coding scheme – 3 bits
- Redundancy version – 2 bits

5.3.3.1.5 Format 2

DCI format 2 is used for scheduling PDSCH to UEs configured in spatial multiplexing mode.

The following information is transmitted by means of the DCI format 2:

In general:

- Resource allocation header (resource allocation type 0 / type 1) – 1 bit
- Resource block assignment:
  - For resource allocation type 0 [3],
    - $\lceil N_{RB}^{DL} / P \rceil$ bits provide the resource allocation
  - For resource allocation type 1 [3],
    - $\lceil \log_2(P) \rceil$ bits of this field are used as a header specific to this resource allocation type to indicate the selected resource blocks subset
    - 1 bit indicates a shift of the resource allocation span
    - $(\lceil N_{RB}^{DL} / P \rceil - \lceil \log_2(P) \rceil - 1)$ bits provide the resource allocation

where the value of P depends on the number of DL resource blocks as indicated in subclause [7.1.1] of [3]

- TPC command for PUCCH – 2 bits

- Downlink Assignment Index (this field just applies to TDD operation) – 2 bits

- HARQ process number - 3 bits (FDD), 4 bits (TDD)

- HARQ swap flag – 1 bit

For codeword 1:

- Modulation and coding scheme – 5 bits

- New data indicator – 1 bit

- Redundancy version – 2 bits

For codeword 2:

- Modulation and coding scheme – 5 bits

- New data indicator – 1 bit

- Redundancy version – 2 bits

Precoding information – number of bits as specified in Table 5.3.3.1.5-1

A codeword field is also used for determining whether the corresponding codeword is enabled or disabled. The interpretation of the precoding information field depends on the number of enabled codewords according to Table 5.3.3.1.5-2, Table 5.3.3.1.5-3 and Table 5.3.3.1.5-4. Note that PMI indicates which codebook index is used in Table 6.3.4.2.3-1 or Table 6.3.4.2.3-2 of [2]. The combination of a single enabled codeword and RI=2 in Table 5.3.3.1.5-3 and Table 5.3.3.1.5-4 is only supported for retransmission of the corresponding HARQ process.

For the open-loop spatial multiplexing transmission mode with 2 antenna ports, the precoding information field is not present. The number of transmission layers, RI, is equal to 2 if both codewords are enabled; and is equal to 1 if codeword 1 is enabled while codeword 2 is disabled.

Table 5.3.3.1.5-1: Number of bits for precoding information

Number of antenna ports at eNode-B	Transmission mode	
	Closed-loop spatial multiplexing	Open-loop spatial multiplexing
2	3	0
4	6	2

Table 5.3.3.1.5-2: Content of precoding information field for 2 antenna ports and closed-loop spatial multiplexing transmission mode

One codeword: Codeword 1 enabled, Codeword 2 disabled	Two codewords: Codeword 1 enabled, Codeword 2 enabled

Bit field mapped to index	Message	Bit field mapped to index	Message
0	RI=1: transmit diversity	0	RI=2: PMI corresponding to precoder matrix $\frac{1}{2}\begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$
1	RI=1: PMI corresponding to precoding vector $\begin{bmatrix} 1 & 1 \end{bmatrix}^T / \sqrt{2}$	1	RI=2: PMI corresponding to precoder matrix $\frac{1}{2}\begin{bmatrix} 1 & 1 \\ j & -j \end{bmatrix}$
2	RI=1: PMI corresponding to precoder vector $\begin{bmatrix} 1 & -1 \end{bmatrix}^T / \sqrt{2}$	2	RI=2: Precoding according to the latest PMI report on PUSCH
3	RI=1: PMI corresponding to precoder vector $\begin{bmatrix} 1 & j \end{bmatrix}^T / \sqrt{2}$	3	reserved
4	RI=1: PMI corresponding to precoder vector $\begin{bmatrix} 1 & -j \end{bmatrix}^T / \sqrt{2}$	4	reserved
5	RI=1: Precoding according to the latest PMI report on PUSCH, if RI=2 reported, use 1st column of all precoders implied by reported PMI(s) and reported RI	5	reserved
6	RI=1: Precoding according to the latest PMI report on PUSCH, if RI=2 reported, use 2nd column of all precoders implied by reported PMI(s) and reported RI	6	reserved
7	reserved	7	reserved

Table 5.3.3.1.5-3: Content of precoding information field for 4 antenna ports and closed-loop spatial multiplexing transmission mode

One codeword: Codeword 1 enabled, Codeword 2 disabled		Two codewords: Codeword 1 enabled, Codeword 2 enabled	
Bit field mapped to index	**Message**	**Bit field mapped to index**	**Message**
0	RI=1: transmit diversity	0	RI=2: PMI=0
1	RI=1: PMI=0	1	RI=2: PMI=1
2	RI=1: PMI=1	⋮	⋮
⋮	⋮	15	RI=2: PMI=15
16	RI=1: PMI=15	16	RI=2: Precoding according to the latest PMI report on PUSCH using the precoder(s) indicated by the reported PMI(s)
17	RI=1: Precoding according to the latest PMI report on PUSCH using the precoder(s) indicated by the reported PMI(s)	17	RI=3: PMI=0
18	RI=2: PMI=0	18	RI=3: PMI=0
19	RI=2: PMI=1	19	RI=3: PMI=1
⋮	⋮	⋮	⋮
33	RI=2: PMI=15	32	RI=3: PMI=15
34	RI=2: Precoding according to the latest PMI report on PUSCH using the precoder(s) indicated by the reported PMI(s)	33	RI=3: Precoding according to the latest PMI report on PUSCH using the precoder(s) indicated by the reported PMI(s)
35 – 63	reserved	34	RI=4: PMI=0
		35	RI=4: PMI=1
		⋮	⋮
		49	RI=4: PMI=15
		50	RI=4: Precoding according to the latest PMI report on PUSCH using the precoder(s) indicated by the reported PMI(s)
		51 – 63	Reserved

Table 5.3.3.1.5-4: Content of precoding information field for 4 antenna ports and open-loop spatial multiplexing transmission mode

One codeword: Codeword 1 enabled, Codeword 2 disabled		Two codewords: Codeword 1 enabled, Codeword 2 enabled	
Bit field mapped to index	**Message**	**Bit field mapped to index**	**Message**
0	RI=1: transmit diversity	0	RI=2: precoder cycling with large delay CDD
1	RI=2: precoder cycling with large delay CDD	1	RI=3: precoder cycling with large delay CDD
2	reserved	2	RI=4: precoder cycling with large delay CDD
3	reserved	3	reserved

5.3.3.1.6 Format 3

DCI format 3 is used for the transmission of TPC commands for PUCCH and PUSCH with 2-bit power adjustments.

The following information is transmitted by means of the DCI format 3:

- TPC command number 1, TPC command number 2,…, TPC command number N

where $N = \left\lfloor \frac{L_{\text{format 0}}}{2} \right\rfloor$, and where $L_{\text{format 0}}$ is equal to the payload size of format 0 before CRC attachment.

5.3.3.1.7 Format 3A

DCI format 3A is used for the transmission of TPC commands for PUCCH and PUSCH with single bit power adjustments.

The following information is transmitted by means of the DCI format 3A:

- TPC command number 1, TPC command number 2,…, TPC command number M

where $M = L_{\text{format 0}}$, and where $L_{\text{format 0}}$ is equal to the payload size of format 0 before CRC attachment.

5.3.3.2 CRC attachment

Error detection is provided on DCI transmissions through a Cyclic Redundancy Check (CRC).

The entire PDCCH payload is used to calculate the CRC parity bits. Denote the bits of the PDCCH payload by $a_0, a_1, a_2, a_3, ..., a_{A-1}$, and the parity bits by $p_0, p_1, p_2, p_3, ..., p_{L-1}$. A is the PDCCH payload size and L is the number of parity bits.

The parity bits are computed and attached according to subclause 5.1.1 setting L to 16 bits, resulting in the sequence $b_0, b_1, b_2, b_3, ..., b_{B-1}$, where $B = A + L$.

In the case where UE transmit antenna selection is not configured or applicable, after attachment, the CRC parity bits are scrambled with the UE identity $x_{ue,0}, x_{ue,1}, ..., x_{ue,15}$ to form the sequence of bits $c_0, c_1, c_2, c_3, ..., c_{B-1}$. The relation between c_k and b_k is:

$c_k = b_k$ for k = 0, 1, 2, …, A-1

$$c_k = \left(b_k + x_{ue,k-A}\right) \bmod 2 \quad \text{for } k = A, A+1, A+2,..., A+15.$$

In the case where UE transmit antenna selection is configured and applicable, after attachment, the CRC parity bits are scrambled with the Antenna selection mask $x_{AS,0}, x_{AS,1},..., x_{AS,15}$ as indicated in Table 5.3.3.2-1 and the UE identity $x_{ue,0}, x_{ue,1},..., x_{ue,15}$ to form the sequence of bits $c_0, c_1, c_2, c_3,..., c_{B-1}$. The relation between c_k and b_k is:

$$c_k = b_k \quad \text{for } k = 0, 1, 2, \ldots, A\text{-}1$$

$$c_k = \left(b_k + x_{ue,k-A} + x_{AS,k-A}\right) \bmod 2 \quad \text{for } k = A, A+1, A+2,..., A+15.$$

Table 5.3.3.2-1: UE transmit antenna selection mask

UE transmit antenna selection	Antenna selection mask $< x_{AS,0}, x_{AS,1},..., x_{AS,15} >$
UE port 0	<0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0>
UE port 1	<0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1>

5.3.3.3 Channel coding

Information bits are delivered to the channel coding block. They are denoted by $c_0, c_1, c_2, c_3,..., c_{K-1}$, where K is the number of bits, and they are tail biting convolutionally encoded according to subclause 5.1.3.1.

After encoding the bits are denoted by $d_0^{(i)}, d_1^{(i)}, d_2^{(i)}, d_3^{(i)},..., d_{D-1}^{(i)}$, with $i = 0,1, \text{and } 2$, and where D is the number of bits on the i-th coded stream, i.e., $D = K$.

5.3.3.4 Rate matching

A tail biting convolutionally coded block is delivered to the rate matching block. This block of coded bits is denoted by $d_0^{(i)}, d_1^{(i)}, d_2^{(i)}, d_3^{(i)},..., d_{D-1}^{(i)}$, with $i = 0,1, \text{and } 2$, and where i is the coded stream index and D is the number of bits in each coded stream. This coded block is rate matched according to subclause 5.1.4.2.

After rate matching, the bits are denoted by $e_0, e_1, e_2, e_3,..., e_{E-1}$, where E is the number of rate matched bits.

5.3.4 Control format indicator

Data arrives each subframe to the coding unit in the form of an indicator for the time span, in units of OFDM symbols, of the DCI in that subframe. The CFI takes values CFI = 1, 2 or 3. For system bandwidths $N_{\text{RB}}^{\text{DL}} > [10]$, the span of the DCI in units of OFDM symbols, 1, 2 or 3, is given by the CFI. For system bandwidths $N_{\text{RB}}^{\text{DL}} \le [10]$, the span of the DCI in units of OFDM symbols, 2, 3 or 4, is given by CFI+1.

The coding flow is shown in Figure 5.3.4-1.

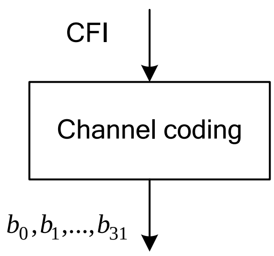

Figure 5.3.4-1 Coding for CFI

5.3.4.1 Channel coding

The control format indicator is coded according to Table 5.3.4-1.

Table 5.3.4-1: CFI codewords

CFI	**CFI codeword** $< \mathbf{b_0, b_1, \ldots, b_{31}} >$
1	<0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1>
2	<1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0>
3	<1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1>
4 (Reserved)	<0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0>

5.3.5 HARQ indicator

Data arrives to the coding unit in form of indicators for HARQ acknowledgement.

The coding flow is shown in Figure 5.3.5-1.

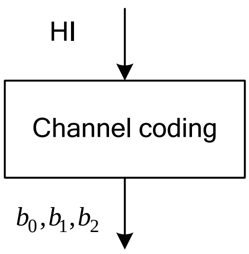

Figure 5.3.5-1 Coding for HI

5.3.5.1 Channel coding

The HARQ indicator is coded according to Table 5.3.5-1, where for a positive acknowledgement HI = 0 and for a negative acknowledgement HI = 1.

Table 5.3.5-1: HI codewords

HI	**HI codeword** $< \mathbf{b_0, b_1, b_2} >$
0	< 0,0,0 >
1	< 1,1,1 >

Annex A (informative): Change history

Change history							
Date	**TSG #**	**TSG Doc.**	**CR**	**Rev**	**Subject/Comment**	**Old**	**New**
2006-09					Skeleton		0.0.0
2006-10					Updated skeleton	0.0.0	0.0.1
2006-10					Endorsed skeleton	0.0.1	0.1.0
2006-11					Added TC. Added Broadcast, Paging and MBMS transport channels in Table 4.2-1.	0.1.0	0.1.1
2006-11					Endorsed v 0.2.0	0.1.1	0.2.0
2006-12					Added CC. Added type of coding for each transport channel or control information.	0.2.0	0.2.1
2007-01					Editor’s version	0.2.1	0.2.2
2007-01					Endorsed v 0.3.0	0.2.2	0.3.0
2007-02					Added QPP turbo Interleaver description.	0.3.0	0.3.1
2007-02					Editor’s version	0.3.1	0.3.2
2007-02					Endorsed v 0.4.0	0.3.2	0.4.0
2007-02					Added CRC details for PDSCH, PDCCH and PUSCH. Added QPP turbo-interleaver parameters. Set Z to 6144. Added details on code block segmentation.	0.4.0	0.4.1
2007-02					Editor’s version	0.4.1	0.4.2
2007-03	RAN#35	RP-070170			For information at RAN#35	0.4.2	1.0.0
2007-03					Editor’s version	1.0.0	1.0.1
2007-03					Editor’s version	1.0.1	1.1.0
2007-05					Editor’s version	1.1.0	1.1.1
2007-05					Editor’s version	1.1.1	1.1.2
2007-05					Editor’s version	1.1.2	1.2.0
2007-06					Added circular buffer rate matching for PDSCH and PUSCH. Miscellaneous changes.	1.2.0	1.2.1
2007-06					Editor’s version	1.2.1	1.2.2
2007-07					Editor’s version	1.2.2	1.2.3
2007-07					Endorsed by email following decision taken at RAN1#49b	1.2.3	1.3.0
2007-08					Editor’s version including decision from RAN1#49bis.	1.3.0	1.3.1
2007-08					Editor’s version	1.3.1	1.3.2
2007-08					Editor’s version	1.3.2	1.4.0
2007-09					Editor’s version with decisions from RAN1#50	1,4.0	1,4,1
2007-09					Editor’s version	1.4.1	1.4.2
10/09/07	RAN#37	RP-070730	-	-	For approval at RAN#37	1.4.2	2.0.0
12/09/07	RAN_37	RP-070730	-	-	Approved version	2.0.0	8.0.0
28/11/07	RAN_38	RP-070949	0001	-	Update of 36.212	8.0.0	8.1.0
05/03/08	RAN_39	RP-080145	0002	-	Update to 36.212 incorporating decisions from RAN1#51bis and RAN1#52	8.1.0	8.2.0
28/05/08	RAN_40	RP-080433	0003	-	Joint coding of CQI and ACK on PUCCH	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0004	1	ACK insertion into PUSCH	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0005	1	Introduction of format 1C	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0006	1	Miscellaneous fixes to 36.212	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0008	1	On multiplexing scheme for indicators	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0009	1	On the soft buffer split of MIMO and TDD	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0010	-	Resource assignment field for distributed VRB	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0011	-	Clarifying the use of the different DCI formats	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0012	1	Clarifying the value of N_L	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0013	-	Payload size for DCI formats 3 and 3A	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0014	-	Coding of ACK on PUSCH	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0015	1	Coding of RI on PUSCH and mapping	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0016	-	CRC for control information on PUSCH	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0017	-	Introduction of Downlink Assignment Index	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0018	-	Coding of CQI/PMI on PUSCH coming from PUCCH	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0019	-	Simultaneous transmission of aperiodic CQI and UL control	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0020	-	Encoding of antenna indicator on DCI format 0	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0021	-	PDCCH coverage in narrow bandwidths	8.2.0	8.3.0
28/05/08	RAN_40	RP-080433	0022	-	Closed-loop and open-loop spatial multiplexing	8.2.0	8.3.0
28/05/08	RAN_40	RP-080457	0023	-	Formula for linkage between PUSCH MCS and amount of resources used for control	8.2.0	8.3.0