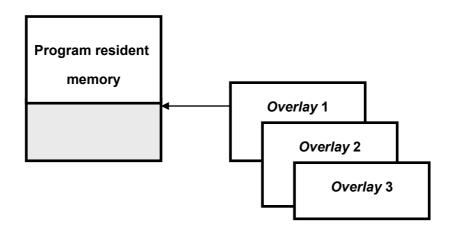
Memory management

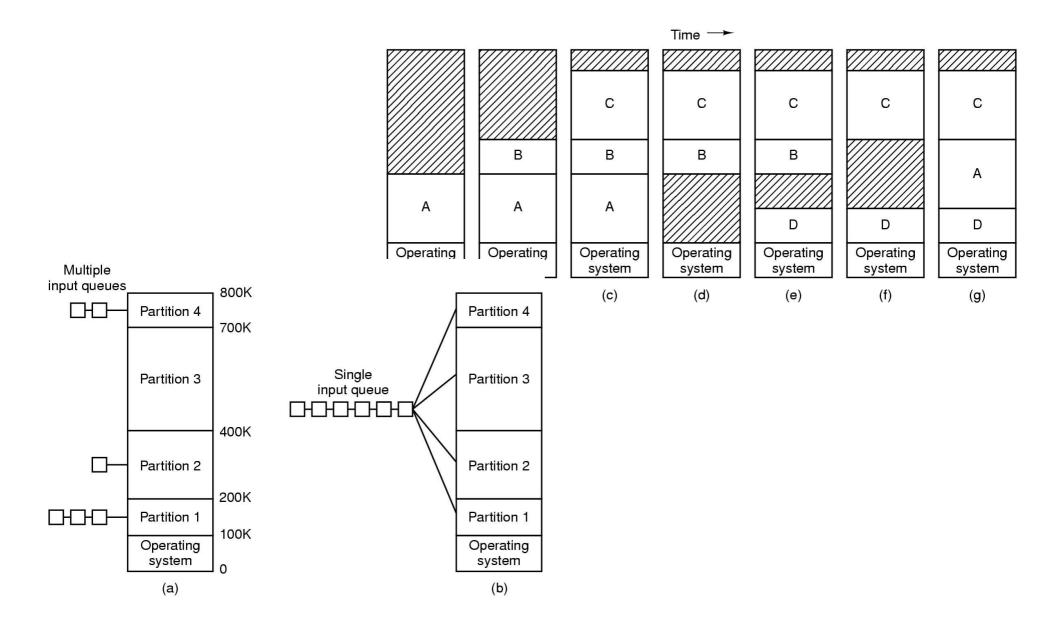
- Single process
 - All memory assigned to the process
 - Addresses defined at compile time
- Multiple processes. How to:
 - assign memory
 - manage addresses?
 - manage relocation?
 - manage program grow?

manage program grow?

- overlay
 - Statically defined memory zone
 - loaded on demand
 - By the program
 - may be replaced



Memory management



```
for(i = 0; i < 1000; i++){
                           .L3:
                           movl-4(%rbp), %eax
   vect[i] = c;
                           cltq
                           movzbl-5(%rbp), %edx
                           movb %dl, -1008(%rbp,%rax)
                           addl $1, -4(%rbp)
                           .L2:
                           cmpl$999, -4(%rbp)
                          jle.L3
```

Virtual memory

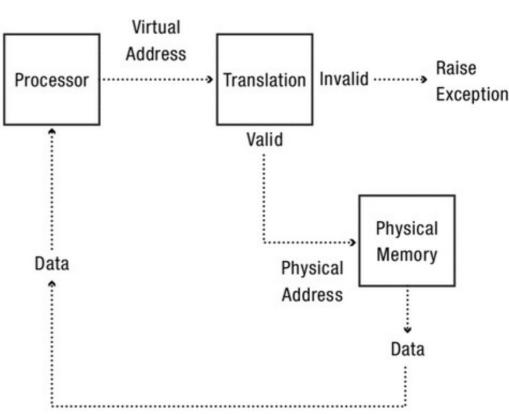
- Program addresses
 - are independent on the physical location
- memory manager

- translates addresses

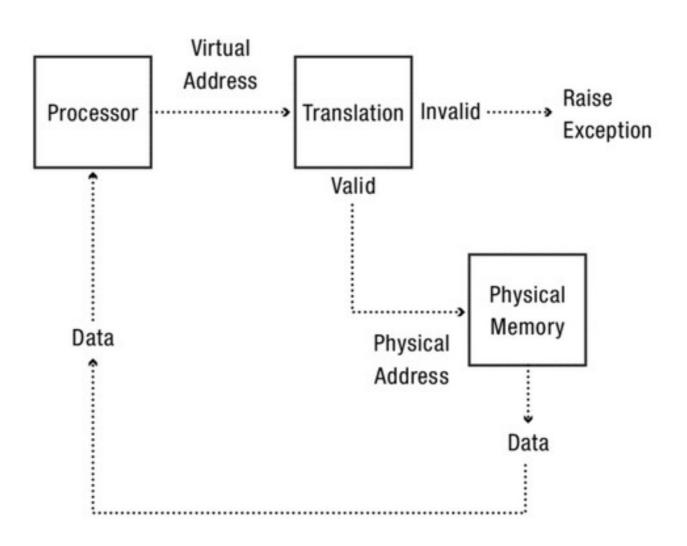
virtual -> physical

verifies permitions

Address Translation



Address Translation Concept



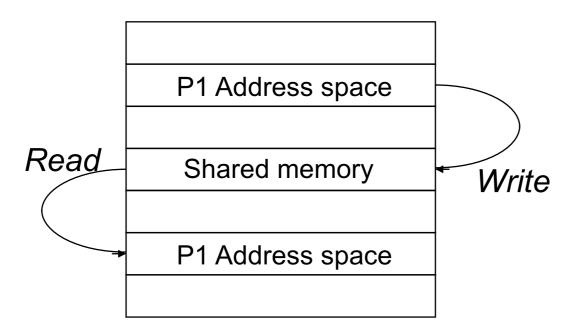
Adress translation

- Process isolation
- Interprocess communication
- Shared code
- Program debugging
- Efficient I/O
- Memory mapped files
- Virtual memory

- Checkpointing
- Process migration
- Information flow control
- DSM

Shared memory

- Processes use regular variables/vector to communicate
 - Available in shared memory
- Should be explicitly created.
 - Each process has its own address space



Shared memory Advantages

Random Access

 you can update a small piece in the middle of a data structure, rather than the entire structure

Efficiency

- unlike message queues and pipes,
 - copy data between user memory
 → kernel memory
- shared memory is directly accessed
- Shared memory resides in the user process memory
 - Is shared among other processes

Process comunication client server write() read(),mq_receive(), write(),mq_send(), **≜**read() or msgsnd() or magrev () process kerne input output (pipe, FIFO, or file file message queue) shared memory client server client address space server address space proces kerne input output file file

Shared memory Disadvantages

- No automatic synchronization
 - In pipes or message queues
 - Programmer has to provide synchronization
 - Semaphores or signals.
- Pointers are only valid within a given process.
 - Pointer offsets cannot be assumed to be valid across inter-process boundaries.
 - This complicates the sharing of linked lists or binary trees.
- Variables are "produced by the compiler"
 - Names can not be used to access shared memory

Shared memory in *NIX

- System V shared memory
 - Original shared memory mechanism, still widely used
 - Sharing between unrelated processes
- Shared mappings mmap
 - Shared file mappings
 - Sharing between unrelated processes, backed by filesystem
 - Shared anonymous mappings
 - Sharing between related processes only (related via fork())
- POSIX shared memory
 - Sharing between unrelated processes, without overhead of filesystem I/O
 - Intended to be simpler and better than older APIs

Shared memory in *NIX

- Programming steps
 - Define Shared data structure
 - Creation of memory segment
 - Configuration
 - Assignment to address
 - Access
 - Disconnection
 - Destruction

In multiple processes

System V shared memory

- Shared memory operations
 - shmget
 - allocates a shared memory segment
 - shmctl
 - allows the user to receive information on a shared memory segment,
 - set the owner, group, and permissions of a shared memory segment,
 - destroy a segment
 - shmat
 - attaches the shared memory segment (identified by shmid) to the address space of the calling process
 - shmdt
 - detaches the shared memory segment (located at the address specified by shmaddr) from the address space of the calling process

System V shared memory

- int shmget(key_t key, size_t size, int shmflg);
 - key known by all processes
 - Flags IPC_CREAT | 0666

- void *shmat(int shmid, const void *shmaddr, int shmflg);
 - Shmid returned by shmget
 - shmaddr NULL or other address
 - Shmflg SHM_EXEC SHM_RDONLY

System V shared memory

```
char * shm;
• key = 5678;

    /* Create the segment */

if ((shmid = shmget(key, SHMSZ, IPC CREAT | 0666)) < 0) {</li>
      perror("shmget"); exit(1);

    /*Now we attach the segment to our data space.*/

if ((shm = shmat(shmid, NULL, 0)) == (char *) -1) {
    perror("shmat"); exit(1);
• }

    Repeated on several processes

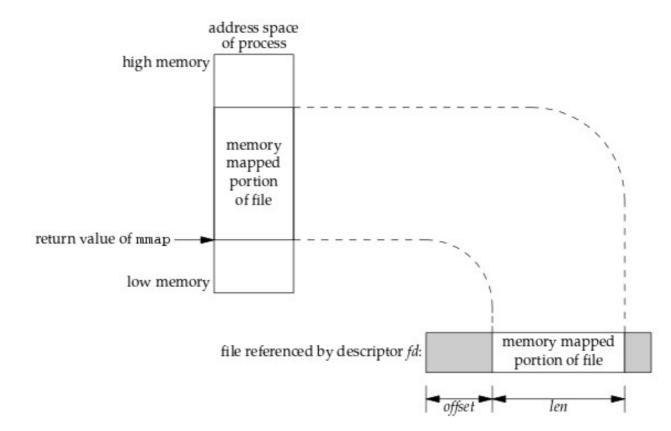
    Key must be known

    Followed by fork
```

Shm value is shared

Memory mapped files

- Access a file contect
 - With memory access operations
 - Assignments e accesses



Memory mapped files

```
    void *mmap(void * addr, size_tlen ,
        int prot, int flags,
        int fd, off_toffset);
```

- Prot -
 - PROT READ PROT WRITE PROT EXEC PROT NONE
- Flags
 - MAP_SHARED MAP_PRIVATE MAP_FIXED
- int munmap(void * addr, size_t len);

Memory access

- After mmap a variable contains a pointer to region
 - Programmer can access that memory as a
 - Pointer to variable
 - Vector
- Is it possible to create linked lists in shared memory?
 - No. mmap in different processes returns different addresses

Memory mapped files

- Int * ptr;
- fd = open(argv[1], O_RDWR | O_CREAT, FILE_MODE);
- ptr = Mmap(NULL, sizeof(int), PROT_READ | PROT_WRITE, MAP_SHARED, fd, 0);
- Close(fd);

Mmap and fork

- Open could be avoided
 - Parent and son share the address to memory
 - Memory is shared among processes
- 4.4BSD introduced anonymous memory sharing
 - Flag MAP_SHARED|MAP_ANON
 - Fd -1
- ptr = Mmap(NULL, sizeof(int),

```
PROT_READ | PROT_WRITE, MAP_SHARED | MAP_ANON, -1, 0);
```

POSIX shared memory

- Mmap
 - File system incurrs overhead
- Sharing between unrelated processes, without overhead of filesystem I/O
- Intended to be simpler and better than older APIs
- New function to create memory regions

POSIX shared memory

- Create/opens a shared memory space
 - fd mem = shm open("/myregion", /*region name*/
 - O_CREAT | O_RDWR, 0600);
 - Memory regions are created with size 0
- Assign a size
 - ftrucate (fd_mem, sizeof(int))
 - If the object has already been sized by another process, you can get its size with the fstat function
- A global region has been created
 - Data stored is kernel persistent
 - But still inaccessible by processes → use mmap with fd mem

```
    int fd = shm open(memname,

                    O CREAT O TRUNC O RDWR,
 0666);
• if (fd == -1)
       error and die("shm open");
int r = ftruncate(fd, sizeof(int));
• if (r != 0)
   error and die("ftruncate");
int *v int = mmap(0, sizeof(int),
          PROT READ | PROT WRITE, MAP SHARED,
```

fd, 0);

Shutdown

- Close the shared memory object
 - close(fd_mem)
- Unmap the shared memory object:
 - munmap (pointer, SHM_SIZE);
 - The address become free to be used.
- At this stage if other processes open and map a memory region
 - They can access data previously written
- To remove permanently the shared memory object:
 - shm_unlink (SHARED_MEMORY_NAME);
 - The object is effectively deleted after the last call to shm_unlink

Synchronization on shared memory

- Multiple processe accesing same variable
- Requirtes synchronization

- Mutexes / condition variables
- Semaphores

Mutexes across processes

• shmm->lock is in shared memory