

Quantum Parallelism

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1 Quantum Parallelism



Let us now assume that we have an **unknown** function f with the following domain and range:

$$f : \{0, 1\}^n \rightarrow \{0, 1\}$$

By unknown I mean that **we do not know** any mathematical description of f , we only have access to a black box that, given a binary string \vec{x} , it (the black box) produces the value $f(\vec{x})$.

Hence, we are able to produce ordered pairs $(\vec{x}, f(\vec{x}))$, one at a time.



Let us suppose that we want to know the exact rules of operation of f .

Using classical computers, we would have to compute $f(\vec{x})$ for all the potential inputs $\vec{x} = 0, 1, \dots, 2^n - 1$. We could to this job either by

- Serial substitution, which would take a very large amount of time for large n , or
- An exponential number of computers, each one ready to compute $f(\vec{x})$ for a given value of \vec{x} .



Instead, we could use quantum mechanics in order to build a (macro) quantum gate which would be able to do this job in a single step.



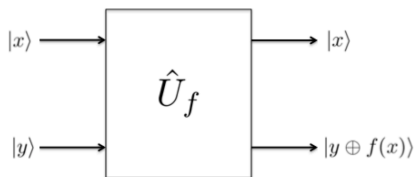
Let us start with the basic case:

$$f : \{0, 1\} \rightarrow \{0, 1\}$$



Let us postulate the existence of the following gate \hat{U}_f :

$$\hat{U}_f : |x\rangle|y\rangle \rightarrow |x\rangle|y \oplus f(x)\rangle$$



XOR Operation		
X	Y	$X \oplus Y$
0	0	0
0	1	1
1	0	1
1	1	0

$$|x, y\rangle \rightarrow |x, y \oplus f(x)\rangle$$

$$|0, 0\rangle \rightarrow |0, 0 \oplus f(0)\rangle$$

$$|0, 1\rangle \rightarrow |0, 1 \oplus f(0)\rangle$$

$$|1, 0\rangle \rightarrow |1, 0 \oplus f(1)\rangle$$

$$|1, 1\rangle \rightarrow |1, 1 \oplus f(1)\rangle$$

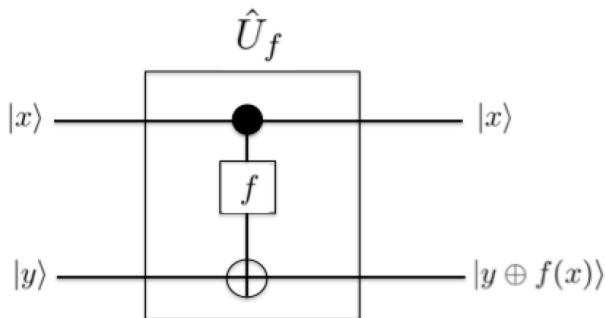


Now, note that:

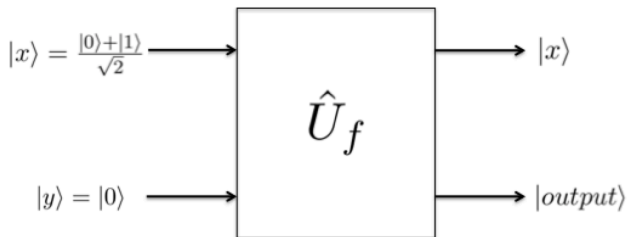
- $f(0) = 0$ or $f(0) = 1$.
 - Suppose $f(0) = 0$. Then, $0 \oplus f(0) = 0 \oplus 0 = 0 = f(0)$.
Consequently, $0 \oplus f(0) = f(0)$ for $f(0) = 0$.
 - Suppose $f(0) = 1$. Then, $0 \oplus f(0) = 0 \oplus 1 = 1 = f(0)$.
Consequently, $0 \oplus f(0) = f(0)$ for $f(0) = 1$.
 - **Thus, $0 \oplus f(0) = f(0)$ regardless the value of $f(0)$.**
- $f(1) = 0$ or $f(1) = 1$.
 - Suppose $f(1) = 0$. Then, $0 \oplus f(1) = 0 \oplus 0 = 0 = f(1)$.
Consequently, $0 \oplus f(1) = f(1)$ for $f(1) = 0$.
 - Suppose $f(1) = 1$. Then, $0 \oplus f(1) = 0 \oplus 1 = 1 = f(1)$.
Consequently, $0 \oplus f(1) = f(1)$ for $f(1) = 1$.
 - **Thus, $0 \oplus f(1) = f(1)$ regardless the value of $f(1)$.**



Operator \hat{U}_f can be created via a variation of the C-NOT gate:



Let us now run \hat{U}_f with $|x\rangle = \frac{|0\rangle+|1\rangle}{\sqrt{2}}$ and $|y\rangle = |0\rangle$ as input, i.e.



$$\begin{aligned}\hat{U}_f|x\rangle|y\rangle &= \hat{U}_f\left[\frac{|0\rangle + |1\rangle}{\sqrt{2}} \otimes |0\rangle\right] \\&= \hat{U}_f\left[\frac{|00\rangle + |10\rangle}{\sqrt{2}}\right] \\&= \frac{1}{\sqrt{2}}\left[\hat{U}_f|00\rangle + \hat{U}_f|10\rangle\right] \\&= \frac{1}{\sqrt{2}}\left[|0\rangle|0 \oplus f(0)\rangle + |1\rangle|0 \oplus f(1)\rangle\right] \\&= \frac{1}{\sqrt{2}}\left[|0\rangle|f(0)\rangle + |1\rangle|f(1)\rangle\right].\end{aligned}$$

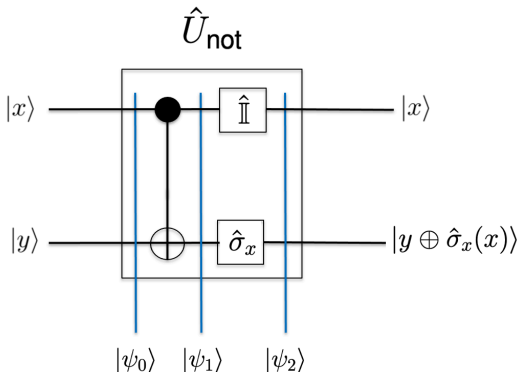
This equation comprises both $f(0)$ and $f(1)$ after a *single run of the gate*.

So, gate \hat{U}_f solves the problem of evaluating f , $\forall x \in D_f$ in one single step.



Example - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Let us now present a detailed example of quantum parallelism:
parallel computation of the logical **NOT** gate.



Example - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Let $|x\rangle = |0\rangle$ and $|y\rangle = |0\rangle$ then

$$|\psi_0\rangle = |0\rangle \otimes |0\rangle$$

$$\begin{aligned} |\psi_1\rangle &= \hat{C}_{\text{not}}|00\rangle \\ &= (|00\rangle\langle 00| + |01\rangle\langle 01| + |10\rangle\langle 11| + |11\rangle\langle 10|)|00\rangle \\ &= |00\rangle \end{aligned}$$

$$\begin{aligned} |\psi_2\rangle &= (\hat{\mathbb{I}} \otimes \hat{\sigma}_x)|0\rangle \otimes |0\rangle \\ &= (|0\rangle\langle 0| + |1\rangle\langle 1|)|0\rangle \otimes (|0\rangle\langle 1| + |1\rangle\langle 0|)|0\rangle \\ &= |01\rangle \end{aligned}$$

So,

$$|\psi_2\rangle = |01\rangle = |0\rangle \otimes \hat{\sigma}_x|0\rangle = |0\rangle|f(0)\rangle$$



Example - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Let $|x\rangle = |1\rangle$ and $|y\rangle = |0\rangle$ then

$$|\psi_0\rangle = |1\rangle \otimes |0\rangle$$

$$\begin{aligned} |\psi_1\rangle &= \hat{C}_{\text{not}}|10\rangle \\ &= (|00\rangle\langle 00| + |01\rangle\langle 01| + |10\rangle\langle 11| + |11\rangle\langle 10|)|10\rangle \\ &= |11\rangle \end{aligned}$$

$$\begin{aligned} |\psi_2\rangle &= (\hat{\mathbb{I}} \otimes \hat{\sigma}_x)|1\rangle \otimes |1\rangle \\ &= (|0\rangle\langle 0| + |1\rangle\langle 1|)|1\rangle \otimes (|0\rangle\langle 1| + |1\rangle\langle 0|)|1\rangle \\ &= |10\rangle \end{aligned}$$

So,

$$|\psi_2\rangle = |10\rangle = |1\rangle \otimes \hat{\sigma}_x|1\rangle = |1\rangle|f(1)\rangle$$



Example - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Let $|x\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}}$ and $|y\rangle = |0\rangle$ then

$$|\psi_0\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}} \otimes |0\rangle$$

$$|\psi_0\rangle = \frac{|00\rangle + |10\rangle}{\sqrt{2}}$$

$$|\psi_1\rangle = \hat{C}_{\text{not}} \frac{|00\rangle + |10\rangle}{\sqrt{2}}$$

$$= (|00\rangle\langle 00| + |01\rangle\langle 01| + |10\rangle\langle 11| + |11\rangle\langle 10|) \left(\frac{|00\rangle + |10\rangle}{\sqrt{2}} \right)$$

$$= \frac{|00\rangle + |11\rangle}{\sqrt{2}}$$



Example - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

$$\begin{aligned}
 |\psi_2\rangle &= (\hat{\mathbb{I}} \otimes \hat{\sigma}_x) \left(\frac{|00\rangle + |11\rangle}{\sqrt{2}} \right) \\
 &= \frac{1}{\sqrt{2}} \left(\hat{\mathbb{I}}|0\rangle \otimes \hat{\sigma}_x|0\rangle + \hat{\mathbb{I}}|1\rangle \otimes \hat{\sigma}_x|1\rangle \right) \\
 &= \frac{1}{\sqrt{2}} (|01\rangle + |10\rangle) \\
 &= \frac{|01\rangle + |10\rangle}{\sqrt{2}}
 \end{aligned}$$

So,

$$|\psi_2\rangle = \frac{|01\rangle + |10\rangle}{\sqrt{2}} = \frac{|0\rangle|\hat{\sigma}_x(0)\rangle + |1\rangle|\hat{\sigma}_x(1)\rangle}{\sqrt{2}} = \frac{|0\rangle|f(0)\rangle + |1\rangle|f(1)\rangle}{\sqrt{2}}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Let us now simulate \hat{U}_{not} on Matlab.

We start by remembering the Kronecker product which is a matrix representation of the tensor product.

Suppose we have operators \hat{A} and \hat{B} that have matrix representations given by $A = (a_{ij})$, $B = (b_{ij})$, two matrices of order n , respectively (for simplicity, we take square matrices as the general definition of the Kronecker product of two non-square matrices can be straightforwardly deduced). So,

$$A = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ \vdots & \vdots & \vdots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{pmatrix}$$

$$B = \begin{pmatrix} b_{11} & b_{12} & \dots & b_{1n} \\ \vdots & \vdots & \vdots & \vdots \\ b_{n1} & b_{n2} & \dots & b_{nn} \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

The Kronecker product of matrices A and B , **i.e., a matrix representation of the tensor product $\hat{A} \otimes \hat{B}$** , is given by

$$\begin{aligned}
 A \otimes B &= \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & & \vdots \\ a_{n1} & \cdots & a_{nn} \end{pmatrix} \otimes \begin{pmatrix} b_{11} & \cdots & b_{1n} \\ \vdots & & \vdots \\ b_{n1} & \cdots & b_{nn} \end{pmatrix} = \begin{pmatrix} a_{11}B & \cdots & a_{1n}B \\ \vdots & & \vdots \\ a_{n1}B & \cdots & a_{nn}B \end{pmatrix} \\
 &= \begin{pmatrix} a_{11} \begin{pmatrix} b_{11} & \cdots & b_{1n} \\ b_{21} & \cdots & b_{2n} \\ \vdots & & \vdots \\ b_{n1} & \cdots & b_{nn} \end{pmatrix} & \cdots & a_{1n} \begin{pmatrix} b_{11} & \cdots & b_{1n} \\ b_{21} & \cdots & b_{2n} \\ \vdots & & \vdots \\ b_{n1} & \cdots & b_{nn} \end{pmatrix} \\ \vdots & & \vdots \\ a_{n1} \begin{pmatrix} b_{11} & \cdots & b_{1n} \\ b_{21} & \cdots & b_{2n} \\ \vdots & & \vdots \\ b_{n1} & \cdots & b_{nn} \end{pmatrix} & \cdots & a_{nn} \begin{pmatrix} b_{11} & \cdots & b_{1n} \\ b_{21} & \cdots & b_{2n} \\ \vdots & & \vdots \\ b_{n1} & \cdots & b_{nn} \end{pmatrix} \end{pmatrix}
 \end{aligned}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

For example, let

$$\hat{H} = \frac{1}{\sqrt{2}}(|0\rangle\langle 0| + |0\rangle\langle 1| + |1\rangle\langle 0| - |1\rangle\langle 1|)$$

be the Hadamard operator.

Write

$$\hat{H} \otimes \hat{H}$$

in matrix notation.



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

$$\begin{aligned}
 H \otimes H &= \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \otimes \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \\
 &= \frac{1}{2} \begin{pmatrix} 1 \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} & 1 \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \\ 1 \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} & -1 \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \end{pmatrix} \\
 &= \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix}
 \end{aligned}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Let us now introduce the matrix form of operators and kets involved in the computation of U_{not} .

$$\mathbb{I} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} ; \quad \sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} ; \quad C_{\text{not}} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} ;$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Also,

$$\mathbb{I} \otimes \sigma_x = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \otimes \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} & 0 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \\ 0 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} & 1 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Moreover,

$$U_{\text{not}} = (\mathbb{I} \otimes \sigma_x) C_{\text{not}} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Furthermore,

$$|0\rangle \otimes |0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\ 0 \begin{pmatrix} 1 \\ 0 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$|0\rangle \otimes |1\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \begin{pmatrix} 0 \\ 1 \end{pmatrix} \\ 0 \begin{pmatrix} 0 \\ 1 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Finally,

$$|1\rangle \otimes |0\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\ 1 \begin{pmatrix} 1 \\ 0 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

$$|1\rangle \otimes |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \begin{pmatrix} 0 \\ 1 \end{pmatrix} \\ 1 \begin{pmatrix} 0 \\ 1 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

In summary,

$$U_{\text{not}} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

$$|0\rangle \otimes |0\rangle = |00\rangle = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$|0\rangle \otimes |1\rangle = |01\rangle = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix}$$

$$|1\rangle \otimes |0\rangle = |10\rangle = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

$$|1\rangle \otimes |1\rangle = |11\rangle = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Now, we know the following:

1. For $|x\rangle = |0\rangle$ and $|y\rangle = |0\rangle$

$$\begin{aligned}\hat{U}_{\text{not}}|x\rangle|y\rangle &= \hat{U}_{\text{not}}|0\rangle|0\rangle = |0\rangle|f(0)\rangle = |0\rangle|1\rangle \text{ i.e.,} \\ \hat{U}_{\text{not}}|00\rangle &= |01\rangle\end{aligned}$$

2. For $|x\rangle = |1\rangle$ and $|y\rangle = |0\rangle$

$$\begin{aligned}\hat{U}_{\text{not}}|x\rangle|y\rangle &= \hat{U}_{\text{not}}|1\rangle|0\rangle = |1\rangle|f(1)\rangle = |1\rangle|0\rangle \text{ i.e.,} \\ \hat{U}_{\text{not}}|10\rangle &= |10\rangle\end{aligned}$$

How do these two results look like in matrix notation?

Based on our previous (and precious) calculations, we can tell the following:



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

So,

1.

$$U_{\text{not}}|00\rangle = |01\rangle \Rightarrow \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix}$$

2.

$$U_{\text{not}}|10\rangle = |10\rangle \Rightarrow \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Now, let us see the parallel power of U_{not} .

Let $|x\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}} = \frac{|0\rangle}{\sqrt{2}} + \frac{|1\rangle}{\sqrt{2}}$ and $|y\rangle = |0\rangle$. So, the total initial state is given by

$$|x\rangle \otimes |y\rangle = |xy\rangle = \frac{|00\rangle}{\sqrt{2}} + \frac{|10\rangle}{\sqrt{2}}$$

In matrix notation, $|xy\rangle$ has the following form:

$$|xy\rangle = \frac{|00\rangle}{\sqrt{2}} + \frac{|10\rangle}{\sqrt{2}} = \frac{|00\rangle + |10\rangle}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} + \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix} = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ 0 \\ \frac{1}{\sqrt{2}} \\ 0 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

We also know that

$$\hat{U}_{\text{not}} \left(\frac{|00\rangle + |10\rangle}{\sqrt{2}} \right) = \frac{|0\rangle|f(0)\rangle + |1\rangle|f(1)\rangle}{\sqrt{2}} = \frac{|01\rangle + |10\rangle}{\sqrt{2}}$$

being the matrix representation of $\frac{|01\rangle + |10\rangle}{\sqrt{2}}$ given by

$$\frac{|01\rangle + |10\rangle}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} + \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \\ 0 \end{pmatrix}$$



Matlab Simulation - Quantum Parallelism and the NOT logical gate ($\hat{\sigma}_x$)

Finally, the computation of $\hat{U}_{\text{not}} \left(\frac{|00\rangle + |10\rangle}{\sqrt{2}} \right)$ in matrix notation is given by

$$U_{\text{not}} \left(\frac{|00\rangle + |10\rangle}{\sqrt{2}} \right) = \frac{|01\rangle + |10\rangle}{\sqrt{2}} \Rightarrow$$

$$\begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} \frac{1}{\sqrt{2}} \\ 0 \\ \frac{1}{\sqrt{2}} \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \\ 0 \end{pmatrix}$$

Let us see how this works in Matlab!



Generalized Quantum Parallelism

Generalized Quantum Parallelism

$$f : \{0, 1\}^n \rightarrow \{0, 1\}$$



Let us now focus on

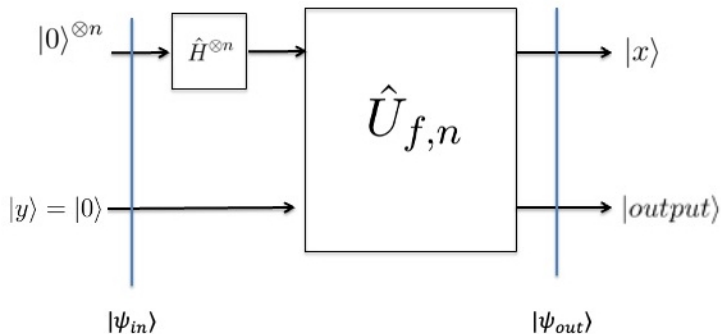
$$f : \{0, 1\}^n \rightarrow \{0, 1\}$$

We define $\hat{U}_{f,n}$, a generalization of \hat{U}_f , as follows:

$$\hat{U}_{f,n}|\vec{x}\rangle|y\rangle \rightarrow |\vec{x}\rangle|y \oplus f(\vec{x})\rangle$$

where $\vec{x} \in \{0, 1\}^n$ is an n -bit string and $f(\vec{x}) \in \{0, 1\}$





$$|\psi_{in}\rangle = \hat{H}^{\otimes n} \otimes \hat{\mathbb{I}} \left(|0\rangle^{\otimes n} \otimes |0\rangle \right) = \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |\vec{x}\rangle \otimes |0\rangle$$



$$\begin{aligned} |\psi_{out}\rangle &= \hat{U}_{f,n} |\psi_{in}\rangle \\ &= \hat{U}_{f,n} \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |\vec{x}\rangle \otimes |0\rangle \\ &= \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} \hat{U}_{f,n} |\vec{x}\rangle \otimes |0\rangle \\ &= \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |\vec{x}\rangle \otimes |0 \oplus f(\vec{x})\rangle \\ &= \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |\vec{x}\rangle \otimes |f(\vec{x})\rangle \end{aligned}$$

Therefore, f has been evaluated $\forall x \in D_f$ in a single step!

