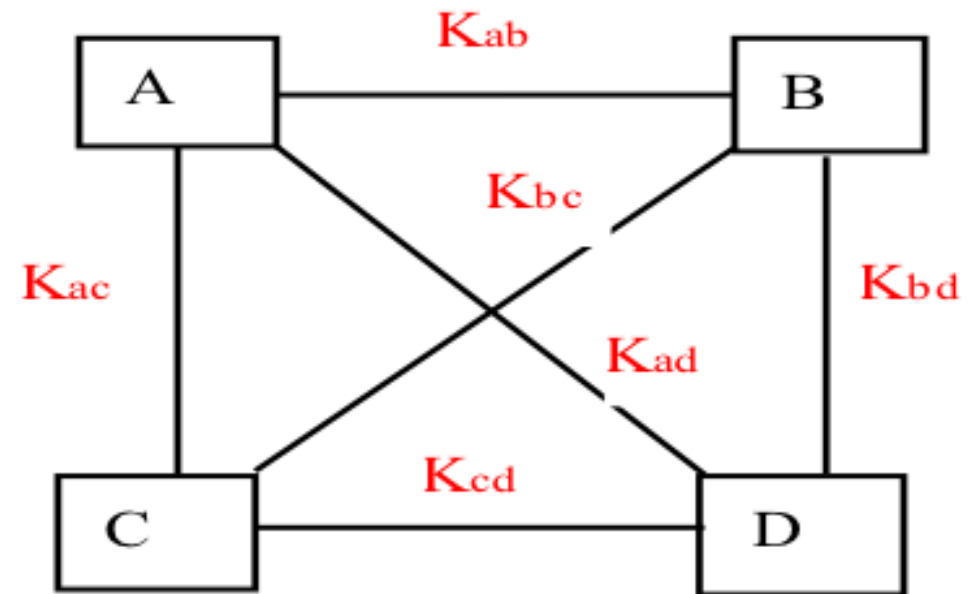


## Symmetric Key Distribution without using PKC

- ❑ Symmetric key distribution using symmetric key encryption - Needham-Schroeder Protocol.
- ❑ This protocol is widely used in single sign on (SSO) solutions, e.g. window domain authentication, Kerberos.

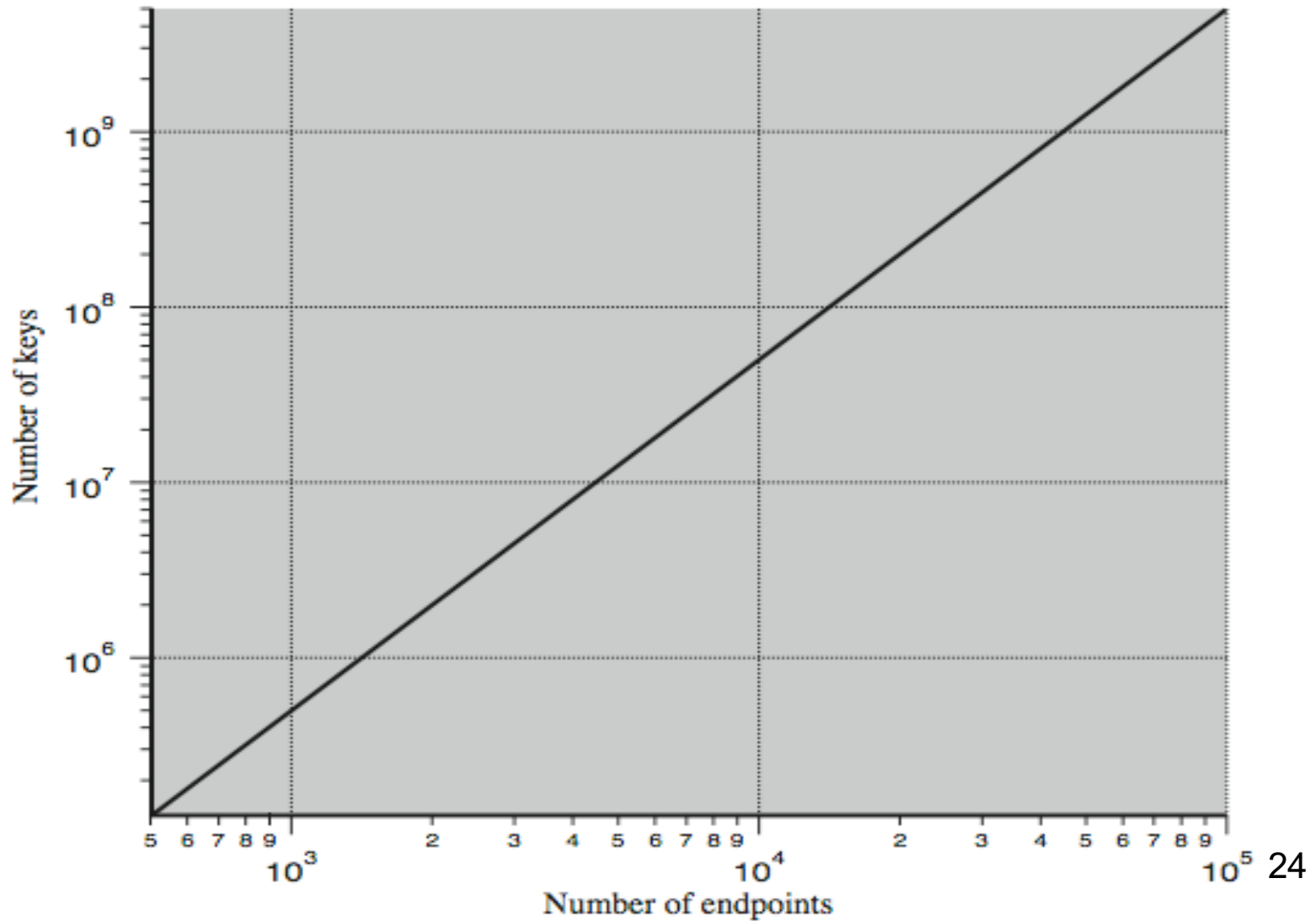
## Distribution without using PKC – Approach-One

- ❑ *Approach One*: Given  $n$  users (parties/nodes) to communicate with each other, the system needs  $n(n-1)/2$  keys.
- ❑ As  $n$  increases, the number of keys becomes untenable for everyone.
- ❑ The  $n^2$  problem!



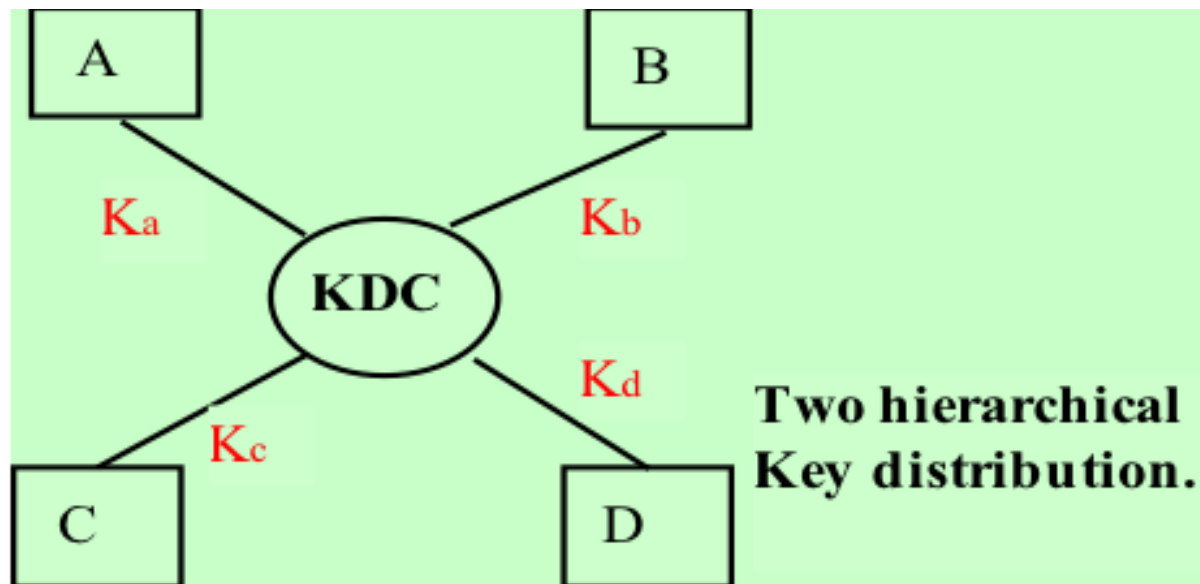
**One hierarchical Key distribution.**

## Distribution without using PKC - Scalability problem



## Distribution without using PKC – Approach-Two

- *Approach Two*: use a key distribution centre (*KDC*) or security server.
  - A key hierarchy, e.g. two hierarchical approach - *master keys* (*long-term keys*) and *session keys* (*valid just for one session*).



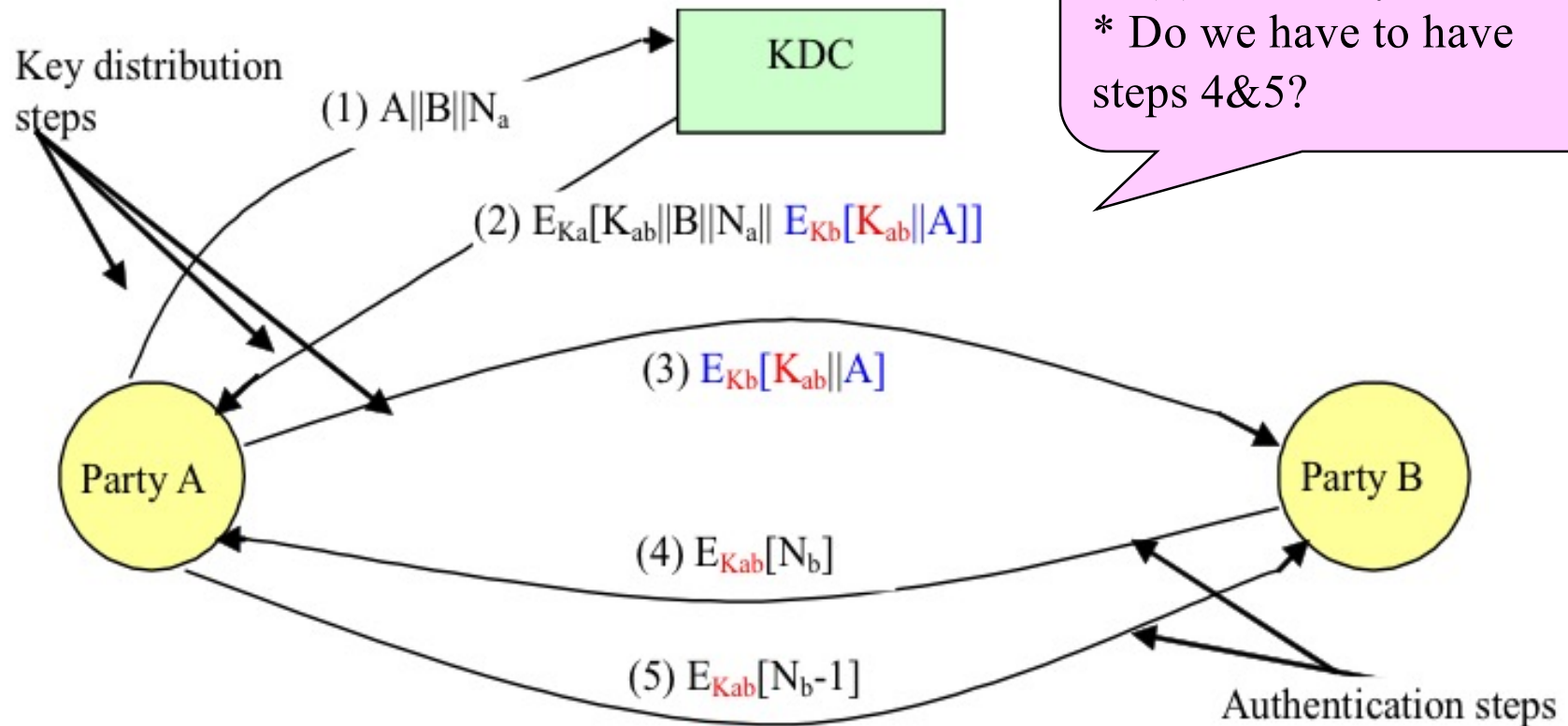
## Distribution without using PKC – Approach-Two

- ❑ A unique **master key**, shared between a pair of user/*KDC*, is for session key distribution.
- ❑ A session key is to secure a particular session.
- ❑ **Benefit** of using Approach Two
  - Reduces the scale of the problem - reduces the  $n^2$  problem to an  $n$  problem, thus making the system more scalable.
- ❑ **But:**
  - The need to trust the intermediaries - KDC.
    - KDC has enough information to impersonate anyone to anyone. If it is compromised, all the resources in the system are vulnerable.
  - KDC is a single point of failure.
  - KDC may be a performance bottleneck.

## Distribution without using PKC - Needham-Schroeder Protocol

- ❑ The Needham-Schroeder is a **key distribution protocol**.
- ❑ It uses Approach-Two. That is:
  - both parties,  $A$  and  $B$ , shares a secret key with the KDC,  $K_a$  and  $K_b$  ;
  - $A$  and  $B$  wishes to establish a secure communication channel, i.e. establish a shared one-time session key  $K_{ab}$ , for use between  $A$  and  $B$  in this session.
- ❑  $N_a$ ,  $N_b$  are nonces (random challenges), generated by  $A$  and  $B$  respectively, to keep messages fresh.

## Distribution without using PKC - Needham-Schroeder Protocol



## Distribution without using PKC - Needham-Schroeder Protocol

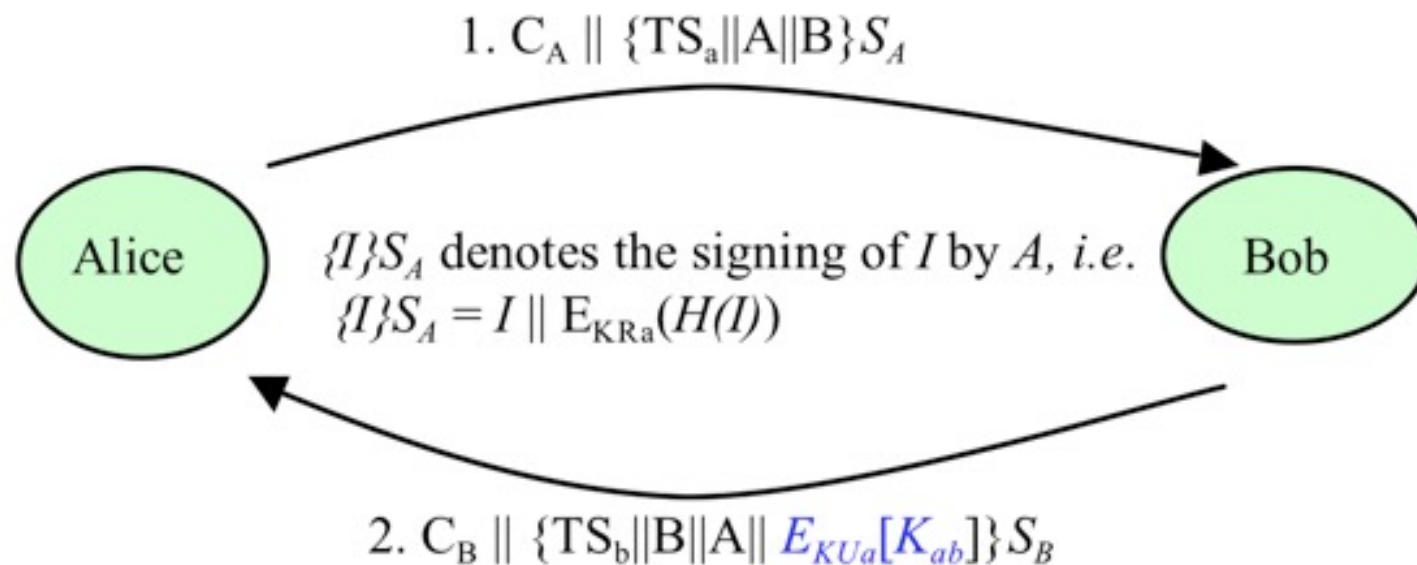
- (1) *A* sends a request to *KDC* for a session key to establish a secure channel with *B*.
  - (2) *KDC* generate a random number  $K_{ab}$ , and replies with the response containing
    - session key  $K_{ab}$ .
    - original request enables *A* matching the response with the request.
    - an item (the session key and *A*' s identity) which only *B* can view.
  - (3) *A* forwards the item to *B*.
- At this point, the session key is securely delivered to A and B, and they may begin secure communication.*
- (4) *B* sends a nonce  $N_b$  to *A* encrypted using the new session key.
  - (5) *A* responds with  $N_b-1$ .

*Steps (4) & (5) assure B that the message received in (3) was not a replay, i.e. to authenticate A.*



## Symmetric Key Distribution using PKC – Two passes

- ❑ Secret key distribution with mutual authentication using public key cipher + **timestamps**.  $C_A$  and  $C_B$  are, respectively, Alice's and Bob's certificates.



## Symmetric Key Distribution using PKC - Three passes

- ❑ Symmetrical key distribution with mutual authentication using public key cipher + **nonces** (random numbers).
- ❑ In both of these two protocols, entity authentication is done by using digital signatures.

