

# Shitcoin Protocol

Didactic Cryptocoin design.

## Import packages

```
In[1]:= Needs["ProgressMapping`, FileNameJoin[{NotebookDirectory[], "ProgressMapping.wl"}]];
```

## Cryptography functions

### Description:

We provide a brief introduction to finite fields. For further information, see Chapter 3 of Koblitz [52], or the books by McEliece [61] and Lidl and Niederreiter [59].

A *finite field* consists of a finite set of elements  $F$  together with two binary operations on  $F$ , called addition and multiplication, that satisfy certain arithmetic properties. The *order* of a finite field is the number of elements in the field. There exists a finite field of order  $q$  if and only if  $q$  is a prime power. If  $q$  is a prime power, then there is essentially only one finite field of order  $q$ ; this field is denoted by  $\mathbb{F}_q$ . There are, however, many ways of representing the elements of  $\mathbb{F}_q$ . Some representations may lead to more efficient implementations of the field arithmetic in hardware or in software.

If  $q = p^m$  where  $p$  is a prime and  $m$  is a positive integer, then  $p$  is called the *characteristic* of  $\mathbb{F}_q$  and  $m$  is called the *extension degree* of  $\mathbb{F}_q$ . Most standards which specify the elliptic curve cryptographic techniques restrict the order of the underlying finite field to be an odd prime ( $q = p$ ) or a power of 2 ( $q = 2^m$ ). In §3.1, we describe the elements and the operations of the finite field  $\mathbb{F}_p$ . In §3.2, elements and the operations of the finite field  $\mathbb{F}_{2^m}$  are described, together with two methods for representing the field elements: *polynomial basis representations* and *normal basis representations*.

### 3.1 The Finite Field $\mathbb{F}_p$

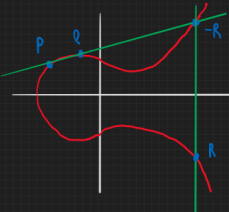
Let  $p$  be a prime number. The finite field  $\mathbb{F}_p$ , called a *prime field*, is comprised of the set of integers  $\{0, 1, 2, \dots, p-1\}$  with the following arithmetic operations:

- ADDITION: If  $a, b \in \mathbb{F}_p$ , then  $a + b = r$ , where  $r$  is the remainder when  $a + b$  is divided by  $p$  and  $0 \leq r \leq p-1$ . This is known as *addition modulo  $p$* .
- MULTIPLICATION: If  $a, b \in \mathbb{F}_p$ , then  $a \cdot b = s$ , where  $s$  is the remainder when  $a \cdot b$  is divided by  $p$  and  $0 \leq s \leq p-1$ . This is known as *multiplication modulo  $p$* .
- INVERSION: If  $a$  is a non-zero element in  $\mathbb{F}_p$ , the *inverse* of  $a$  modulo  $p$ , denoted  $a^{-1}$ , is the unique integer  $c \in \mathbb{F}_p$  for which  $a \cdot c = 1$ .

**Example 1.** (The finite field  $\mathbb{F}_{23}$ ) The elements of  $\mathbb{F}_{23}$  are  $\{0, 1, 2, \dots, 22\}$ . Examples of the arithmetic operations in  $\mathbb{F}_{23}$  are: (i)  $12 + 20 = 9$ ; (ii)  $8 \cdot 9 = 3$ ; and (iii)  $8^{-1} = 3$ .

Álgebra en curvas elípticas

$y^2 = x^3 + ax + b$



La pendiente de la recta  $\overline{PQ}$  es

$$M = \frac{y_Q - y_P}{x_Q - x_P}$$

$$y - y_P = \frac{y_Q - y_P}{x_Q - x_P} (x - x_P)$$

$$y = y_P + M(x - x_P)$$

Simplificando

$$y = Mx + d$$

La intersección ocurre en

$$(Mx + d)^2 = x^3 + ax + b$$

$$\Rightarrow x^3 - (Mx + d)^2 + ax + b = 0$$

$$x^3 - M^2x^2 - 2Md x - d^2 + ax + b = 0$$

$$x^3 - M^2x^2 + (a - 2Md)x + (b - d^2) = 0 \quad (1)$$

Se conoce de antemano que la ec (1) debe tener 3 raíces

$$(x - x_P)(x - x_Q)(x - x_R) = (x^3 - x^2x_Q - x^2x_P + x^2x_R)(x - x_R)$$

$$= x^3 - x^2x_Q - x^2x_P + x^2x_R - (x_Rx^2 - x_Qx_R - x_Px_R + x_Px_Qx_R)$$

$$= x^3 - (x_Q + x_P + x_R)x^2 + (x_Px_Q + x_Qx_R + x_Px_R)x - x_Px_Qx_R = 0 \quad (2)$$

Se puede obtener la solución para  $x_R$  igualando los términos Cuadráticos (en 1 y 2)

$$M^2 = x_Q + x_P + x_R \Rightarrow x_R = M^2 - x_P - x_Q$$

Al sustituir  $x_R$  en la ecuación de la recta y reflejar, queda

$$y_R = M(x_P - x_R) - y_P$$

Caso límite cuando Q se aproxima infinitesimalmente a P

La pendiente de la recta tangente a P se obtiene derivando  $y(x)$

$$\Rightarrow y'(x) = \frac{d}{dx} \sqrt{x^3 + ax + b} = \frac{2x^2 + a}{2\sqrt{x^3 + ax + b}}$$

$$= \frac{2x^2 + a}{2y(x)}$$

y evaluando en  $x_P$

$$\Rightarrow M = \frac{2x_P^2 + a}{2y_P}$$

y la operación suma está dada por

$$x_R = M^2 - 2x_P$$

$$y_R = M(x_P - x_R) - y_P$$

Implementation:

In[64]:=

```

EllipticCurveEvaluate[x_, a_, b_] := Sqrt[x^3 + a x + b];
EllipticCurvePointQ[{x_, y_}, a_, b_, p_] := Block[{lhs, rhs},
  lhs = Mod[y^2, p];
  rhs = Mod[x^3 + a x + b, p];

```

```

    lhs == rhs
];
EllipticCurveAdd[{px_, py_}, {qx_, qy_}] := Block[{m, rx, ry},
  If[px == qx, Return[∞]];
  m =  $\frac{qy - py}{qx - px}$ ;
  rx =  $m^2 - px - qx$ ;
  ry = m (px - rx) - py;
  {rx, ry}
];
EllipticCurveAdd[{px_, py_}, {qx_, qy_}, p_] := Block[{snum, sden, s, rx, ry},
  If[px == qx || ! CoprimeQ[qx - px, p], Return[∞]];

  snum = Mod[qy - py, p];
  sden = ModularInverse[qx - px, p];
  s = Mod[snum * sden, p];
  rx = Mod[s2 - px - qx, p];
  ry = Mod[s (px - rx) - py, p];
  {rx, ry}
];
EllipticCurveDouble[{px_, py_}, a_] := Block[{m, rx, ry},
  m =  $\frac{3 px^2 + a}{2 py}$ ;
  rx =  $m^2 - 2 px$ ;
  ry = m (px - rx) - py;
  {rx, ry}
];
EllipticCurveDouble[{px_, py_}, a_, p_] := Block[{snum, sden, s, rx, ry},
  snum = Mod[3 px2 + a, p];
  If[py == 0 || ! CoprimeQ[2 py, p], Return[∞]];
  sden = ModularInverse[2 py, p];
  s = Mod[snum * sden, p];
  rx = Mod[s2 - 2 px, p];
  ry = Mod[s (px - rx) - py, p];
  {rx, ry}
];

ComputeCyclicGroup[G_, a_, p_] := Block[{G2},
  G2 = EllipticCurveDouble[G, a, p];
  Prepend[NestWhileList[EllipticCurveAdd[G, #, p] &, G2, # != ∞ &], G]
];
MultiplicationPath[G_, a_, p_, k_] := Block[{G2},
  G2 = EllipticCurveDouble[G, a, p];
  Prepend[NestList[EllipticCurveAdd[G, #, p] &, G2, k - 2], G]
];
MultiplyBasePoint[G_, a_, p_, k_] := Block[{kBinary, P},
  kBinary = Drop[IntegerDigits[k, 2], 1];
  P = G;
  Do[
    P = EllipticCurveDouble[P, a, p];

```

```

    If[kBinary[[i]] == 1, P = EllipticCurveAdd[P, G, p]];
    ,
    {i, 1, Length[kBinary]}
];
Return[P];
];
MultiplyBasePointHex[{hexGx_, hexGy_}, aHex_, pHex_, kHex_] :=
Block[{Gx, Gy, k, Px, Py, a, p},
  Gx = FromDigits[hexGx, 16];
  Gy = FromDigits[hexGy, 16];
  k = FromDigits[kHex, 16];
  a = FromDigits[aHex, 16];
  p = FromDigits[pHex, 16];
  {Px, Py} = MultiplyBasePoint[{Gx, Gy}, a, p, k];
  {ToHex[Px, 64], ToHex[Py, 64]}
];

```

## Example operations

```

In[38]:= k = 5;
a = 2;
p = 17;

In[41]:= G = {5, 1};

In[43]:= G2 = EllipticCurveDouble[G, a, p]
Out[43]= {6, 3}

In[44]:= G3 = EllipticCurveAdd[G, G2, p]
Out[44]= {10, 6}

In[45]:= MultiplyBasePoint[G, a, p, 10]
Out[45]= {7, 11}

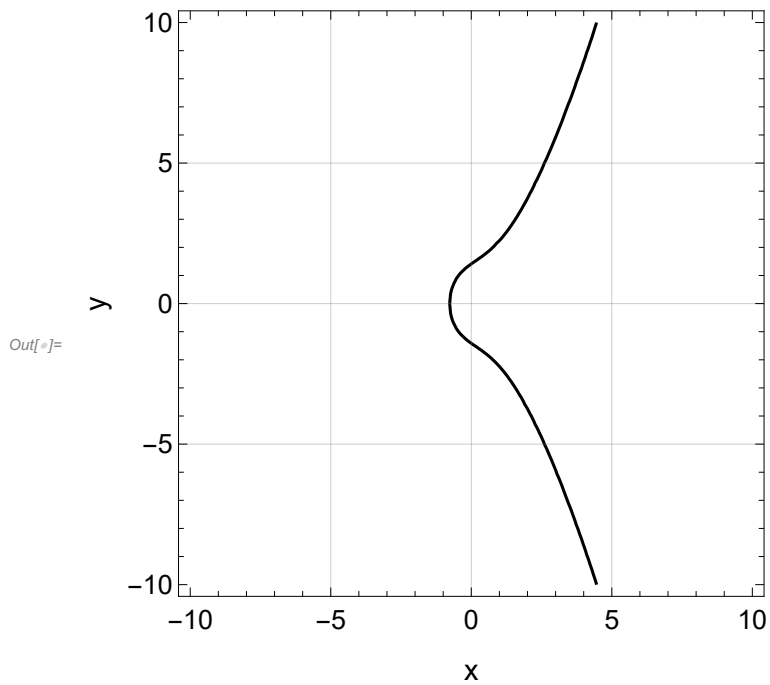
```

## Curves over the Reals Field

```

In[ ]:= a = 2;
b = 2;
ContourPlot[y^2 == x^3 + a x + b, {x, -10, 10}, {y, -10, 10},
  PlotTheme -> "Monochrome",
  Frame -> True,
  BaseStyle -> FontSize -> 14,
  GridLines -> Automatic,
  FrameLabel -> {Style["x", 15], Style["y", 15]}
]

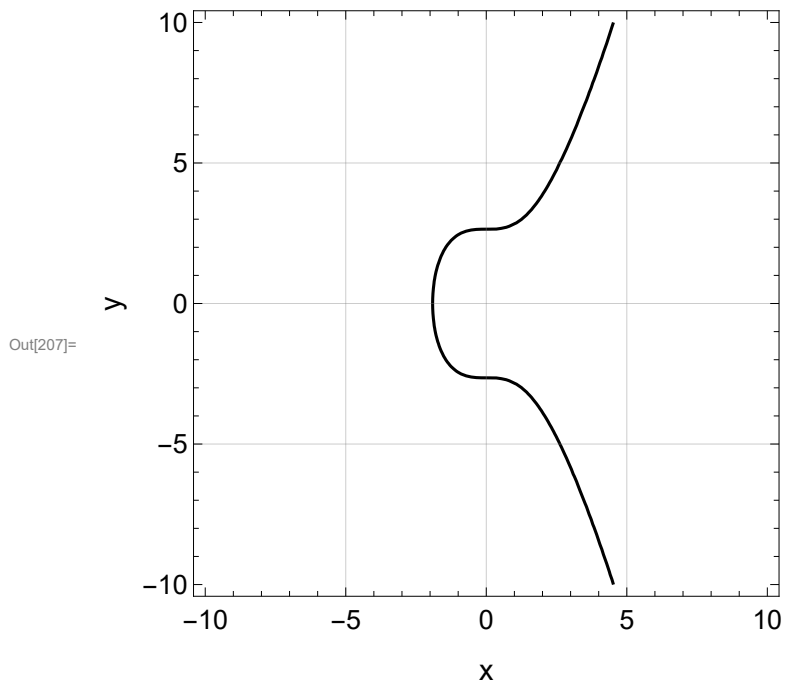
```



```

In[205]:= a = 0;
b = 7;
ContourPlot[y^2 == x^3 + a x + b, {x, -10, 10}, {y, -10, 10},
  PlotTheme -> "Monochrome",
  Frame -> True,
  BaseStyle -> FontSize -> 14,
  GridLines -> Automatic,
  FrameLabel -> {Style["x", 15], Style["y", 15]}
]

```



```

In[219]:= G = {2, N@EllipticCurveEvaluate[2, a, b]}

```

```

Out[219]= {2, 3.87298}

```

```

In[220]:= G2 = EllipticCurveDouble[G, a]

```

```

Out[220]= {-1.6, 1.70411}

```

```

In[226]:= G3 = EllipticCurveAdd[G, G2]

```

```

Out[226]= {-0.037037, -2.64574}

```

```

In[227]:= G4 = EllipticCurveAdd[G, G3]

```

```

Out[227]= {8.27769, -23.9622}

```

```

In[228]:= G5 = EllipticCurveAdd[G, G4]

```

```

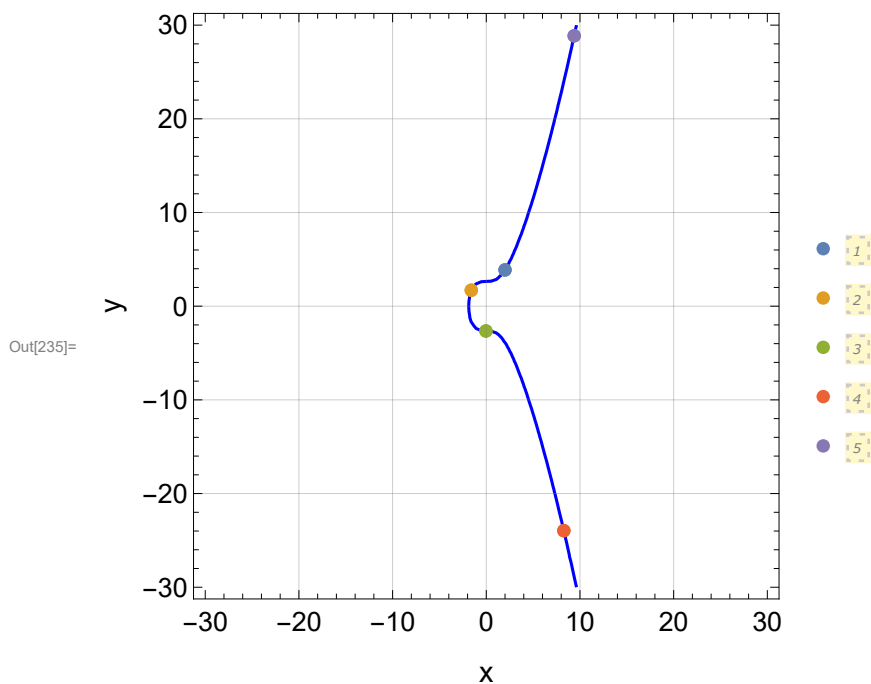
Out[228]= {9.38258, 28.8613}

```

```

In[235]:= Show[
  ContourPlot[y^2 == x^3 + a x + b, {x, -30, 30}, {y, -30, 30},
    PlotTheme -> "Monochrome",
    Frame -> True,
    BaseStyle -> FontSize -> 14,
    GridLines -> Automatic,
    FrameLabel -> {Style["x", 15], Style["y", 15]},
    ContourStyle -> Blue
  ],
  ListPlot[{G}, {G2}, {G3}, {G4}, {G5}],
    PlotStyle -> PointSize[Large],
    PlotLegends -> Automatic
]
]

```



## Curves over Finite Field

```

In[89]:= CalculateCurvePoints[a_, b_, p_] := Flatten[ProgressTable[
  If[EllipticCurvePointQ[{x, y}, a, b, p], {x, y}, Nothing], {x, 0, p}, {y, 0, p}], 1];

```

$\mathbb{F}_{17}$

```

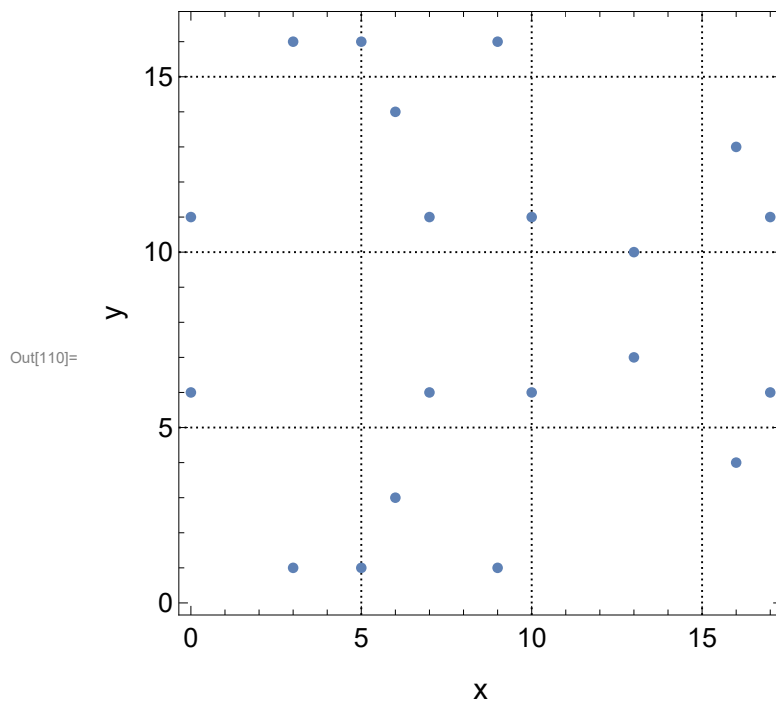
In[99]:= k = 5;
a = 2;
b = 2;
p = 17;
G = {5, 1};

```

```
In[105]:= GGroup = ComputeCyclicGroup[G, a, p];
Thread[Range[Length[GGroup]] → GGroup]
```

```
Out[106]:= {1 → {5, 1}, 2 → {6, 3}, 3 → {10, 6}, 4 → {3, 1}, 5 → {9, 16}, 6 → {16, 13}, 7 → {0, 6},
8 → {13, 7}, 9 → {7, 6}, 10 → {7, 11}, 11 → {13, 10}, 12 → {0, 11}, 13 → {16, 4},
14 → {9, 1}, 15 → {3, 16}, 16 → {10, 11}, 17 → {6, 14}, 18 → {5, 16}, 19 → ∞}
```

```
In[109]:= curveSet = CalculateCurvePoints[a, b, p];
ListPlot[curveSet,
  AspectRatio → 1,
  PlotTheme → "Monochrome",
  Frame → True,
  BaseStyle → FontSize → 14,
  GridLines → Automatic,
  FrameLabel → {Style["x", 15], Style["y", 15]}
]
```

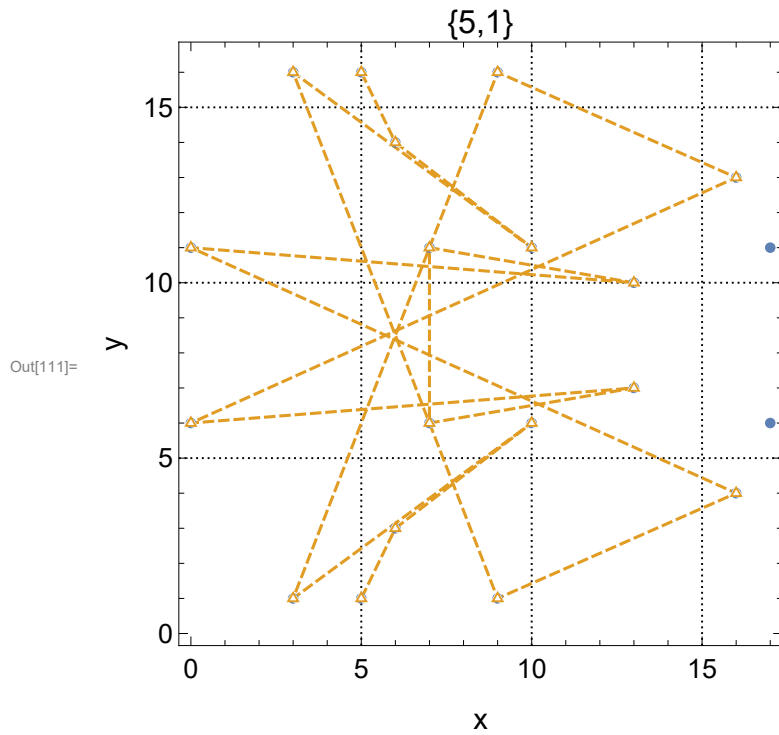




```

In[111]:= ListPlot[{curveSet, GGroup},
  AspectRatio → 1,
  Joined → {False, True},
  PlotTheme → "Monochrome",
  Frame → True,
  BaseStyle → FontSize → 14,
  GridLines → Automatic,
  FrameLabel → {Style["x", 15], Style["y", 15]},
  PlotLabel → "{5,1}"
]

```

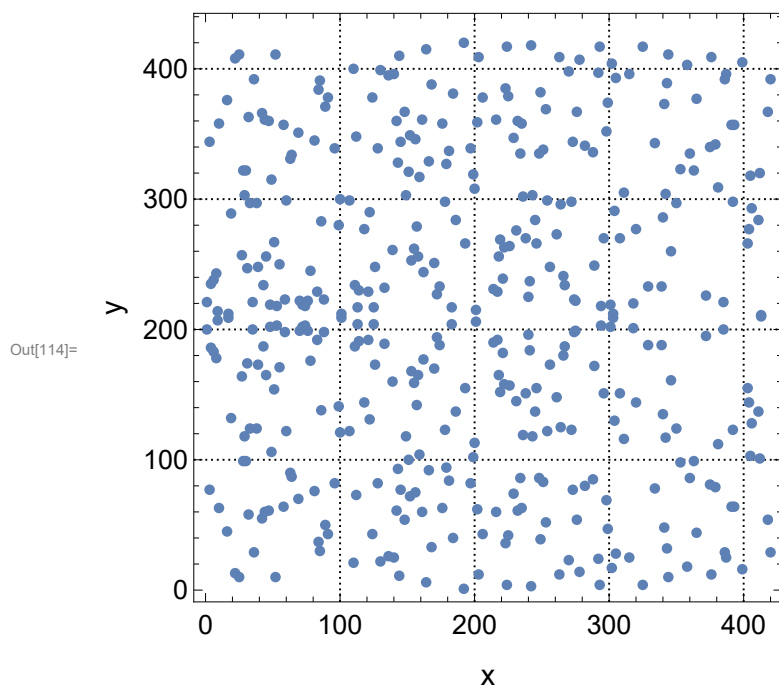


$\mathbb{F}_{421}$ 

```

In[112]:= p = 421;
curveSet = CalculateCurvePoints[a, b, p];
ListPlot[curveSet,
  AspectRatio → 1,
  PlotTheme → "Monochrome",
  Frame → True,
  BaseStyle → FontSize → 14,
  GridLines → Automatic,
  FrameLabel → {Style["x", 15], Style["y", 15]}
]

```



```

In[53]:= GGroup = ComputeCyclicGroup[{6, 183}, a, p];

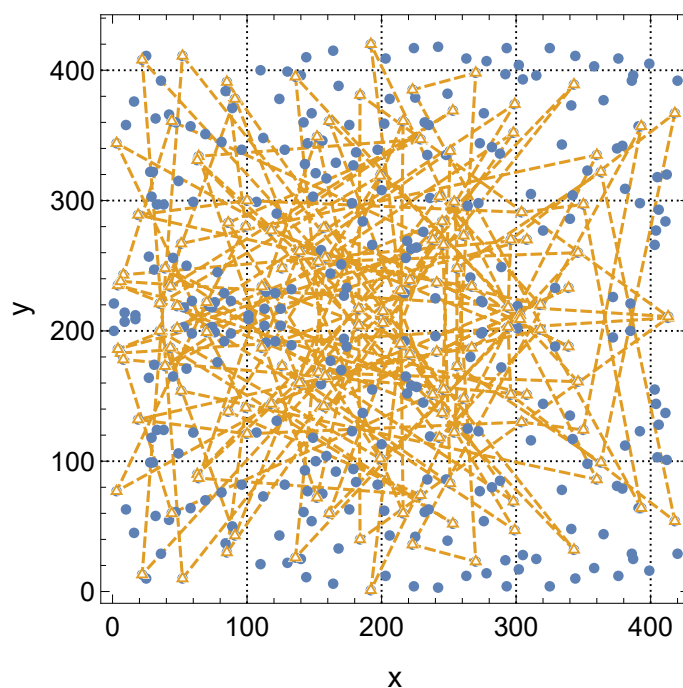
```

```

In[54]:= ListPlot[{curveSet, GGroup},
  AspectRatio -> 1,
  Joined -> {False, True},
  PlotTheme -> "Monochrome",
  Frame -> True,
  BaseStyle -> FontSize -> 14,
  GridLines -> Automatic,
  FrameLabel -> {Style["x", 15], Style["y", 15]}
]

```

Out[54]=



# ECDSA signature and verification

## Description:

### ECDSA Sign

The ECDSA signing algorithm ([RFC 6979](#)) takes as input a message *msg* + a private key *privKey* and produces as output a **signature**, which consists of pair of integers  $\{r, s\}$ . The **ECDSA signing** algorithm is based on the [ElGamal signature scheme](#) and works as follows (with minor simplifications):

1. Calculate the message **hash**, using a cryptographic hash function like SHA-256:  $h = \text{hash}(msg)$
2. Generate securely a **random** number *k* in the range  $[1..n-1]$ 
  - In case of **deterministic-ECDSA**, the value *k* is HMAC-derived from *h* + *privKey* (see [RFC 6979](#))
3. Calculate the random point  $R = k * G$  and take its x-coordinate:  $r = R.x$
4. Calculate the signature proof:  $s = k^{-1} * (h + r * privKey) \pmod n$ 
  - The modular inverse  $k^{-1} \pmod n$  is an integer, such that  $k * k^{-1} \equiv 1 \pmod n$
5. Return the **signature**  $\{r, s\}$ .

The calculated **signature**  $\{r, s\}$  is a pair of integers, each in the range  $[1..n-1]$ . It encodes the random point  $R = k * G$ , along with a proof *s*, confirming that the signer knows the message *h* and the private key *privKey*. The proof *s* is by idea verifiable using the corresponding *pubKey*.

**ECDSA signatures** are **2 times longer** than the signer's **private key** for the curve used during the signing process. For example, for 256-bit elliptic curves (like `secp256k1`) the ECDSA signature is 512 bits (64 bytes) and for 521-bit curves (like `secp521r1`) the signature is 1042 bits.

### ECDSA Verify Signature

The algorithm to **verify a ECDSA signature** takes as input the signed message *msg* + the signature  $\{r, s\}$  produced from the signing algorithm + the public key *pubKey*, corresponding to the signer's private key. The output is boolean value: **valid** or **invalid** signature. The **ECDSA signature verify** algorithm works as follows (with minor simplifications):

1. Calculate the message **hash**, with the same cryptographic hash function used during the signing:  $h = \text{hash}(msg)$
2. Calculate the modular inverse of the signature proof:  $s^{-1} \pmod n$
3. Recover the random point used during the signing:  $R' = (h * s^{-1}) * G + (r * s^{-1}) * pubKey$
4. Take from *R'* its x-coordinate:  $r' = R'.x$
5. Calculate the signature validation **result** by comparing whether  $r' == r$

The general idea of the signature verification is to **recover the point *R'*** using the public key and check whether it is same point *R*, generated randomly during the signing process.

## How Does it Work?

The **ECDSA signature**  $\{r, s\}$  has the following simple explanation:

- The signing **signing** encodes a random point  $R$  (represented by its x-coordinate only) through elliptic-curve transformations using the private key **privKey** and the message hash  $h$  into a number  $s$ , which is the **proof** that the message signer knows the private key **privKey**. The signature  $\{r, s\}$  cannot reveal the private key due to the difficulty of the **ECDLP problem**.
- The **signature verification** decodes the proof number  $s$  from the signature back to its original point  $R$ , using the public key **pubKey** and the message hash  $h$  and compares the x-coordinate of the recovered  $R$  with the  $r$  value from the signature.

## The Math behind the ECDSA Sign / Verify

Read this section **only if you like math**. Most developer may skip it.

How does the above sign / verify scheme work? It is not obvious, but let's play a bit with the equations.

The equation behind the recovering of the point  $R'$ , calculated during the **signature verification**, can be transformed by replacing the **pubKey** with **privKey** \*  $G$  as follows:

$$\begin{aligned} R' &= (h * s^{-1}) * G + (r * s^{-1}) * \text{pubKey} = \\ &= (h * s^{-1}) * G + (r * s^{-1}) * \text{privKey} * G = \\ &= (h + r * \text{privKey}) * s^{-1} * G \end{aligned}$$

If we take the number  $s = k^{-1} * (h + r * \text{privKey}) \pmod{n}$ , calculated during the signing process, we can calculate  $s^{-1} = s^{-1} \pmod{n}$  like this:

$$\begin{aligned} s^{-1} &= s^{-1} \pmod{n} = \\ &= (k^{-1} * (h + r * \text{privKey}))^{-1} \pmod{n} = \\ &= k * (h + r * \text{privKey})^{-1} \pmod{n} \end{aligned}$$

Now, replace  $s^{-1}$  in the point  $R'$ .

$$\begin{aligned} R' &= (h + r * \text{privKey}) * s^{-1} * G = \\ &= (h + r * \text{privKey}) * k * (h + r * \text{privKey})^{-1} \pmod{n} * G = \\ &= k * G \end{aligned}$$

The final step is to **compare** the point  $R'$  (decoded by the **pubKey**) with the point  $R$  (encoded by the **privKey**). The algorithm in fact compares only the x-coordinates of  $R'$  and  $R$ : the integers  $r'$  and  $r$ .

It is expected that  $r' == r$  if the signature is **valid** and  $r' \neq r$  if the signature or the message or the public key is incorrect.

Implementation:

$$\ln[1\ 1]:=$$

```

ECDSASign[G_, a_, p_, m_, k_, hash_, privateKey_] := Block[{x, y, r, s},
  {x, y} = MultiplyBasePoint[G, a, p, k];
  r = x;
  s = Mod[ModularInverse[k, m] * (hash + r * privateKey), m];
  {r, s}
];

ECDSSASignHex[{hexGx_, hexGy_}, aHex_, pHex_, mHex_, kHex_, hashHex_, privateKeyHex_] :=
Block[{Gx, Gy, a, p, m, k, hash, privateKey, r, s},
  Gx = FromDigits[hexGx, 16];
  Gy = FromDigits[hexGy, 16];
  a = FromDigits[aHex, 16];
  p = FromDigits[pHex, 16];
  m = FromDigits[mHex, 16];
  k = FromDigits[kHex, 16];
  hash = FromDigits[hashHex, 16];
  privateKey = FromDigits[privateKeyHex, 16];
  {r, s} = ECDSASign[{Gx, Gy}, a, p, m, k, hash, privateKey];

  ToHex[r, 64] <> ToHex[s, 64]
];

ECDSAValidate[G_, a_, p_, m_, {r_, s_}, hash_, publicKey_] := Block[{w, u1, u2, X},
  If[r < 1 || r > m - 1, Return[False]];
  If[s < 1 || s > m - 1, Return[False]];

  w = ModularInverse[s, m];
  u1 = Mod[hash * w, m];
  u2 = Mod[r * w, m];
  X = EllipticCurveAdd[
    MultiplyBasePoint[G, a, p, u1], MultiplyBasePoint[publicKey, a, p, u2], p];
  Mod[First[X], m] == r
];

ECDSAValidateHex[{hexGx_, hexGy_}, aHex_, pHex_,
mHex_, signature_, hashHex_, {publicKeyXHex_, publicKeyYHex_}] :=
Block[{rHex, sHex, Gx, Gy, a, p, m, r, s, hash, publicKeyX, publicKeyY},
  {rHex, sHex} = StringPartition[signature, 64];
  Gx = FromDigits[hexGx, 16];
  Gy = FromDigits[hexGy, 16];
  a = FromDigits[aHex, 16];
  p = FromDigits[pHex, 16];
  m = FromDigits[mHex, 16];
  r = FromDigits[rHex, 16];
  s = FromDigits[sHex, 16];
  hash = FromDigits[hashHex, 16];
  publicKeyX = FromDigits[publicKeyXHex, 16];
  publicKeyY = FromDigits[publicKeyYHex, 16];

  ECDSAValidate[{Gx, Gy}, a, p, m, {r, s}, hash, {publicKeyX, publicKeyY}]
];

```

### Example:

```

In[115]:= a = 1;
          b = 4;
          p = 23;
          G = {0, 2};
          m = Length[ComputeCyclicGroup[G, a, p]];

In[120]:= privateKey = 5
Out[120]= 5

In[121]:= publicKey = MultiplyBasePoint[G, a, p, privateKey]
Out[121]= {7, 20}

In[ ]:= messageHash = 17;

In[ ]:= {r, s} = ECDSASign[G, a, p, m, k, messageHash, privateKey]
Out[ ]:= {18, 11}

In[ ]:= ECDSAValidate[G, a, p, m, {r, s}, messageHash, publicKey]
Out[ ]:= True

```

### Visualization:

```

In[125]:= path = MultiplicationPath[G, a, p, privateKey]
Out[125]= {{0, 2}, {13, 12}, {11, 9}, {1, 12}, {7, 20}}

In[123]:= curveSet = CalculateCurvePoints[a, b, p];

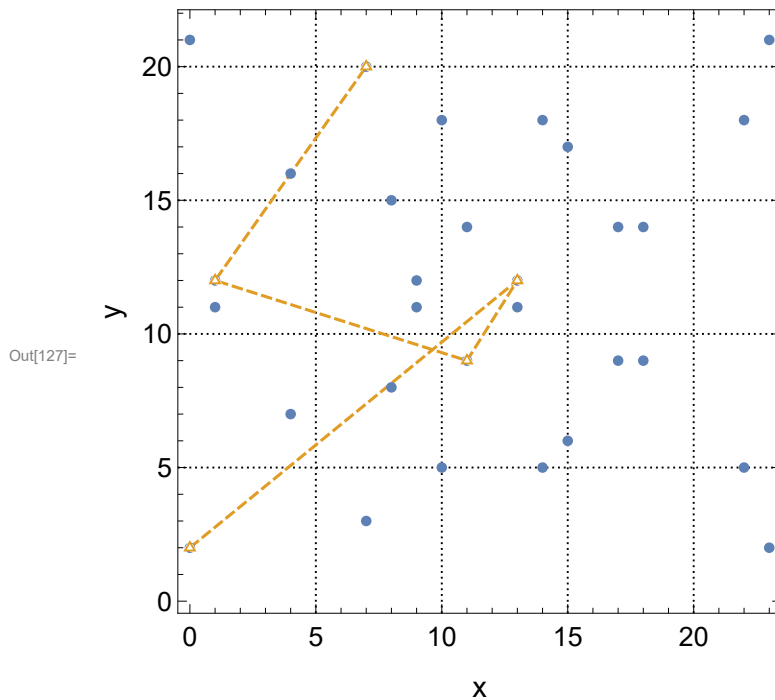
```



```

In[127]:= ListPlot[{curveSet, path},
  AspectRatio → 1,
  Joined → {False, True},
  PlotTheme → "Monochrome",
  Frame → True,
  BaseStyle → FontSize → 14,
  GridLines → Automatic,
  FrameLabel → {Style["x", 15], Style["y", 15]}
]

```



## Network simulation

secp256k1 parameters

```

In[139]:= p = "FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFC2F";
m = "FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFEBAAEDCE6AF48A03BBFD25E8CD0364141";
a = "0000000000000000000000000000000000000000000000000000000000000000";
b = "000000000000000000000000000000000000000000000000000000000000007";
G = { "79BE667EF9DCBBAC55A06295CE870B07029BFCDB2DCE28D959F2815B16F81798",
      "483ADA7726A3C4655DA4FBFC0E1108A8FD17B448A68554199C47D08FFB10D4B8" };

```

## Key Generation

The **ECDSA key-pair** consists of:

- **private key** (integer): *privKey*
- **public key** (EC point): *pubKey* = *privKey* \* *G*

The **private key** is generated as a **random integer** in the range  $[0 \dots n-1]$ . The public key *pubKey* is a point on the elliptic curve, calculated by the EC point multiplication: *pubKey* = *privKey* \* *G* (the private key, multiplied by the generator point *G*).

In[15]:=

```
ToHex[n_, pad_ : 64] := IntegerString[n, 16, pad];
Concatenate[l_] := Apply[Join, l];
GenerateNodes[G_, a_, m_, n_] :=
  Block[{mdec, privateKeys, publicKeys, privateKeysPadded, publicKeysPadded},
    mdec = FromDigits[m, 16];
    privateKeys = Table[ToHex[RandomInteger[{2, mdec - 1}], 64], n];
    publicKeys = Map[StringJoin[MultiplyBasePointHex[G, a, m, #]] &, privateKeys];

    MapThread[<| "PublicKey" -> #1, "PrivateKey" -> #2 |> &, {publicKeys, privateKeys}]
  ];
```

```
In[145]:= networkNodes = GenerateNodes[G, a, p, 10]
```

```
Out[145]= {<|PublicKey →
936864b882a618a25708fd32e3287cac10d289a27b22cd130d3042e8e25e05cf74ea37cdd13550a4c2ce4\
2c6d0e222d3410c4a00ddd932261e8e3b8ccd65a68f, PrivateKey →
39629ecaccd9f717c1f515f7415f5d13ca2468be1de6c6966ce21f3af4c8029f|>, <|PublicKey →
61249a16b3f56e3db642e5ca162beb0b64c3bee9e8768902c1482f6580e25de7a26590bb8cb37bbb4eb1f\
dc9a49d3e645a5ca480101030c732609033ccbc2b8b, PrivateKey →
1946334d1e5cd7bb60a83ee3b6170eab392be71a862d56e5c7ef72a4db647101|>, <|PublicKey →
5e6ac138d7041676c8ccea562f5de44071aedd90cf18e393b4288efdb8bb08048a73fee8baf07e577fd\
2d99588749360dd0c46ad953967cf6775232a0ea8be, PrivateKey →
05dd2f218f631dafc92c8e8042501a9084de22a8495030f32012c695ee78fe76|>, <|PublicKey →
8eb2357242224513080b264505b3a68d4db1dc36e1cfe0a50f0bd90a03f81de9376940dac7b6e6092ce7d\
54250b70e6b5282f605a9a81803005f2703c8734ed2, PrivateKey →
401d188c38232c4612c94312684d1b5627ab22bd2f29279f86565d41c00608a0|>, <|PublicKey →
8fdefc4cf847de6659a79c1641fa801aef8310b7b493cd5c810b4c3b5d58a99f43a298d156d9b71ee7b45\
84c5c79e89dace77d4b96d1e6634c9fcdcd289b2b0c, PrivateKey →
9fda40e27f6d01d9b428d9467621ee76970478c4582fc27d0ed89561a0ff82cd|>, <|PublicKey →
dded0656896dbaef475215738f5120b0f477c8f3110a65f5b30dedcb684ce600987642476c187fb88bef7\
02061cb2a5aa78426f92397d4b59f290de923dd1ddc, PrivateKey →
bf5fda59b79d949beb61d292235de6288b733bd2bd6b0856e76e992e24568b3f|>, <|PublicKey →
24203ad958810688c6690f64b2ab410aa246b2fc00cd26565ac503604562e87173227191b2b54eb2abf86\
bd0711e77ffb74f7d371531f8b121e54e7ad3591569, PrivateKey →
be8e3c235f19007081120a6f697ca8db0fb6eeb1e85996c775b797f6a784f7df|>, <|PublicKey →
9bcfa371088022e83287d55892844fb08284bb5e402123e1bc608c7ab4dbf664cab4065c5bec6a127fb6c\
e11c3a8424f038aef5a799a6b725362eb23e2459d2, PrivateKey →
63a9b4f9f1c7550b233d9734aa38a7c3000f0970cf38b49820698bdcccc4330a|>, <|PublicKey →
94698ebc7e686035d1cf40c503e15622bb4a0d9c82d7a2e40a09e3a5e69327a9e8fd48281d666dcde3165\
2a9392fae483b02b620ae9ea634f8151bbb3428f42b, PrivateKey →
f983987749daa34f40f20be9e8c30dd58916760ee8eab3934589f1389d326ee0|>, <|PublicKey →
6798e006757a961c3394d90c1b1ab8fedce2ab563b68d8d17f65a729fa9aae25433e4d34f661f30e5206a\
f049e154aa6018470f707969dba926605cc14c5e4bd,
PrivateKey → c91bf857ddc61c2fdb4c1934e2a9ef4797f652f907a3b79ea4da204b935575fd|>}
```

In[133]:=

```

Transaction[sender_, receiver_, units_] := <|
  "Sender" → sender, "Receiver" → receiver, "Units" → units|>;
SimulateTransactions[networkNodes_, n_] := Block[{transactors},
  transactors =
    Take[Partition[RandomSample[networkNodes[[All, "PublicKey"]]], 2], n];
  Map[Transaction[First[#], Last[#], ToHex[RandomInteger[{1, 255}], 2]] &,
    transactors]
];

TransactionBlock[transactions_, prevBlockSignature_] :=
  prevBlockSignature <> StringJoin[Concatenate[Map[Values, transactions]]];
RandomNonceHex[m_] := IntegerString[RandomInteger[{1, FromDigits[m, 16]}], 16];

MineBlock[G_, a_, p_, m_, hash_, privateKey_, difficulty_ : 2] :=
  Block[{i, signature, leadingDigits},
    i = 1;
    signature = "11";
    leadingDigits = StringJoin[ConstantArray["0", difficulty]];

    While[StringTake[signature, 2] ≠ leadingDigits,
      signature = ECDSASignHex[G, a, p, m, RandomNonceHex[m], hash, privateKey];
      i++;
    ];
    {i, signature}
  ];

```

In[146]:= transactions = SimulateTransactions[networkNodes, 3]

Out[146]=

```

{<|Sender →
  94698ebc7e686035d1cf40c503e15622bb4a0d9c82d7a2e40a09e3a5e69327a9e8fd48281d666dcde3165\
  2a9392fae483b02b620ae9ea634f8151bbb3428f42b, Receiver →
  8eb2357242224513080b264505b3a68d4db1dc36e1cfe0a50f0bd90a03f81de9376940dac7b6e6092ce7d\
  54250b70e6b5282f605a9a81803005f2703c8734ed2, Units → 6d|>, <|Sender →
  5e6ac138d7041676c8cceaf562f5de44071aeddff90cf18e393b4288efdb8bb08048a73fee8baf07e577fd\
  2d99588749360dd0c46ad953967cf6775232a0ea8be, Receiver →
  dded0656896dbaeef475215738f5120b0f477c8f3110a65f5b30dedcb684ce600987642476c187fb88bef7\
  02061cb2a5aa78426f92397d4b59f290de923dd1ddc, Units → e0|>, <|Sender →
  9bcfa371088022e83287d55892844fb08284bb5e402123e1bc608c7ab4dbf664cab4065c5bec6a127fb6c\
  e11c3a8424f038aeff5a799a6b725362eb23e2459d2, Receiver →
  61249a16b3f56e3db642e5ca162beb0b64c3bee9e8768902c1482f6580e25de7a26590bb8cb37bbb4eb1f\
  dc9a49d3e645a5ca480101030c732609033ccbc2b8b, Units → 0d|> }

```

```
In[151]: block = TransactionBlock(transactions, ToHex[0, 128])
```

[illegible]

```
In[152]:= hash = Hash[block, "SHA256", "HexString"]
```

```
Out[152]= 2d61c71a084e1f019be8695470bd70e17da372ae80df9b3ffc247560923c7531
```

## Simulation of nodes

```

ln[158]:=
AddRewardToBlock[transactions_, publicKey_] := Append[
    transactions, <|"Sender" → ToHex[0, 128], "Receiver" → publicKey, "Units" → "01"|>];
SimulateNode[prevBlockSignature_, transactions_, privateKey_, publicKey_,
    {G_, a_, p_, m_}] := Block[{transactionsWithReward, block, hash, publicKeyParts,
    transactionsWithReward = AddRewardToBlock[RandomSample[transactions], publicKey];
    block = TransactionBlock[transactionsWithReward, prevBlockSignature];
    hash = Hash[block, "SHA256", "HexString"];
    publicKeyParts = StringPartition[publicKey, 64];

    {block, publicKeyParts, MineBlock[G, a, p, m, hash, privateKey, 2]}
];

```

```
In[165]:= nodesSimulation = ProgressMap[SimulateNode[ToHex[0, 128], transactions,
#["PrivateKey"], #["PublicKey"], {G, a, p, m}] &, networkNodes];
```

```
In[166]:= firstMine = First[MinimalBy[nodesSimulation[[All, 3]], First]];
```

```
In[167]:= firstMiner = FirstCase[nodesSimulation, { , , firstMine}];
```



```
In[183]:= nodesSimulation = ProgressMap[SimulateNode[prevBlockSignature,  
    transactions, #["PrivateKey"], #["PublicKey"], {G, a, p, m}] &, networkNodes];  
  
In[184]:= firstMine = First[MinimalBy[nodesSimulation[[All, 3]], First]];  
firstMiner = FirstCase[nodesSimulation, {_, _, firstMine}];  
newBlock2 = <|"TransactionBlock" → First[firstMiner],  
    "Signature" -> Last[Last[firstMiner]], "PublicKey" → firstMiner[[2]]|>  
  
Out[186]= <|TransactionBlock →  
007d7f5ee6e52b4f41d9d5229ec38ddf971db0af00a3c9c46c3115a90378f60ee9d2ee31e60358967b6fa2:  
1a19f7069a5b28dec501d1fc42a7cdf69d98f770b9dded0656896dbaef475215738f5120b0f477c8f:  
3110a65f5b30dedcb684ce600987642476c187fb88bef702061cb2a5aa78426f92397d4b59f290de9:  
23dd1ddc5e6ac138d7041676c8cceaf562f5de44071aeddf90cf18e393b4288efdb8bb08048a73fee:  
8baf07e577fd2d99588749360dd0c46ad953967cfc6775232a0ea8be3a6798e006757a961c3394d90c:  
1b1ab8fedce2ab563b68d8d17f65a729fa9aae25433e4d34f661f30e5206af049e154aa6018470f70:  
7969dba926605cc14c5e4bd61249a16b3f56e3db642e5ca162beb0b64c3bee9e8768902c1482f6580:  
e25de7a26590bb8cb37bbb4eb1fdc9a49d3e645a5ca480101030c732609033ccbc2b8b698fdefc4cf:  
847de6659a79c1641fa801aef8310b7b493cd5c810b4c3b5d58a99f43a298d156d9b71ee7b4584c5c:  
79e89dace77d4b96d1e6634c9fcdcd289b2b0c936864b882a618a25708fd32e3287cac10d289a27b2:  
2cd130d3042e8e25e05cf74ea37cdd13550a4c2ce42c6d0e222d3410c4a00ddd932261e8e3b8ccd65:  
a68fee000000000000000000000000000000000000000000000000000000000000000000000000:  
000000000000000000000000000000000000000000000000000000000000000000000000000000:  
a68d4db1dc36e1cfe0a50f0bd90a03f81de9376940dac7b6e6092ce7d54250b70e6b5282f605a9a81:  
803005f2703c8734ed201, Signature →  
00ade813fa79a3f34805c227c02df4b7ce02e8854b64c084735a1073cde85f99542cbdc161e9e14305ec2e:  
172b7fbdde80e48fc9de6cfe62b0e98112972408cd,  
PublicKey → {8eb2357242224513080b264505b3a68d4db1dc36e1cfe0a50f0bd90a03f81de9,  
376940dac7b6e6092ce7d54250b70e6b5282f605a9a81803005f2703c8734ed2}|>
```

## Iteration 3

```

In[187]:= transactions = SimulateTransactions[networkNodes, 3]

Out[187]:= {<| Sender →
8fdefc4cf847de6659a79c1641fa801aef8310b7b493cd5c810b4c3b5d58a99f43a298d156d9b71ee7b45\
84c5c79e89dace77d4b96d1e6634c9fcddc289b2b0c, Receiver →
61249a16b3f56e3db642e5ca162beb0b64c3bee9e8768902c1482f6580e25de7a26590bb8cb37bbb4eb1f\
dc9a49d3e645a5ca480101030c732609033ccbc2b8b, Units → de |>, <| Sender →
936864b882a618a25708fd32e3287cac10d289a27b22cd130d3042e8e25e05cf74ea37cdd13550a4c2ce4\
2c6d0e222d3410c4a00ddd932261e8e3b8ccd65a68f, Receiver →
24203ad958810688c6690f64b2ab410aa246b2fc00cd26565ac503604562e87173227191b2b54eb2abf86\
bd0711e77ffb74f7d371531f8b121e54e7ad3591569, Units → 08 |>, <| Sender →
9bcfa371088022e83287d55892844fb08284bb5e402123e1bc608c7ab4dbf664cab4065c5bec6a127fb6c\
e11c3a8424f038aef5a799a6b725362eb23e2459d2, Receiver →
8eb2357242224513080b264505b3a68d4db1dc36e1cfe0a50f0bd90a03f81de9376940dac7b6e6092ce7d\
54250b70e6b5282f605a9a81803005f2703c8734ed2, Units → f0 |> }

In[188]:= prevBlockSignature = newBlock2["Signature"];

In[189]:= nodesSimulation = ProgressMap[SimulateNode[prevBlockSignature,
transactions, #["PrivateKey"], #["PublicKey"], {G, a, p, m}] &, networkNodes];

```

