# Optimization Techniques

### Laboratory 5

Iterated Local Search, Simulated Annealing

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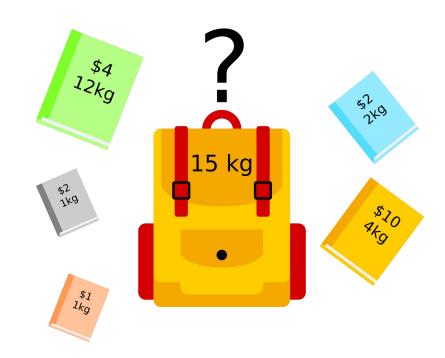
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## 0-1 Knapsack Problem

Maximize the sum of values in the bag bounded by the total weight

- Bag has capacity W
- Each item has volume w value v



## 0-1 Knapsack Problem

Assume

N = number of items

 $X = \{x_1, x_2, \dots, x_N\} = \text{set of items}$ 

 $v_i = \text{value of item } x_i$ 

 $w_i = \text{weight of item } x_i$ 

W = maximum weight the knapsack can hold

Then the 0/1 knapsack problem can be formulated as follows:



### Iterated Local Search

Stochastic Local Search method based on an embedded heuristic

#### Given solution s\*:

- Perturb s\* leading to intermediate state s'
- Perform Local Search around s', arriving at solution s'\*
- If s'\* passes the acceptance test, then s\* = s'\*
- 1. Back to step 1 until convergence

```
 \begin{array}{l} \textbf{procedure } \textit{Iterated Local Search} \\ s_0 \leftarrow \textbf{GenerateInitialSolution} \\ s^* \leftarrow \textbf{LocalSearch}(s_0) \\ \textbf{repeat} \\ s' \leftarrow \textbf{Perturbation}(s^*, \textit{history}) \\ s^{*\prime} \leftarrow \textbf{LocalSearch}(s') \\ s^* \leftarrow \textbf{AcceptanceCriterion}(s^*, s^{*\prime}, \textit{history}) \\ \textbf{until termination condition met} \\ \textbf{end} \end{array}
```

Better than random restarts as we do not completely give up on the current solution

## Simulated Annealing

#### Improves greedy search

by allowing step to a worse neighbour

- Neighbor selection transform the given solution to introduce a new solution
- Temperature scheduling
  willingness to accept a worse state
  decreased over time
- Acceptance probability
   probability of going from one state to a neighbor state

- X ← Initial configuration
- T ← Initial high temperature
- Iterate:
- 1. Do K times:
  - $1.1 E \leftarrow Eval(X)$
  - 1.2  $X' \leftarrow$  one configuration randomly selected in *Neighbors* (X)
  - 1.3  $E' \leftarrow Eval(X')$
  - $1.4 \text{ If } E' \leq E$

$$X \leftarrow X'; E \leftarrow E';$$

Else accept the move with probability  $p = e^{-(E'-E)/T}$ 

$$X \leftarrow X'$$
;  $E \leftarrow E'$ ;

2.  $T \leftarrow \alpha T$