

I/O Devices - Disk (ch. 36+37)

Operating Systems

Based on: Three Easy Pieces by Arpaci-Dusseau

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Hard Disk Drives

- Main form of persistent data storage
- File system technology: predicated on their behavior

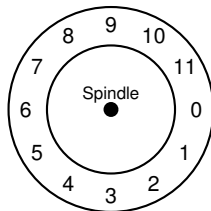
How do modern hard-disk drives store data?
What is the interface?

The Interface

- Consists of sectors (512-byte blocks)
 - Numbered 0 to $n - 1$ (the drive **address space**)
 - Each can be read or written
- Multi-sector operations are possible
 - Many file systems read or write 4KB at a time
 - Only guarantee: single 512-byte block write is **atomic**
 - i.e., will completely entirely or not at all
 - **Torn write**: only portion of a larger write complete
- Common assumption: sequential access is the fastest

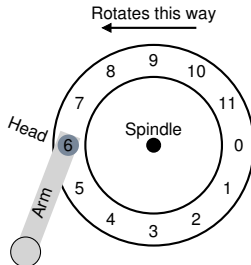
Basic Geometry

- A **platter**
 - Circular surface on which data is stored
 - Two sides, each called a **surface**
- A disk has one or more platters
 - Bound together around the **spindle**
 - Connected to a motor that spins the platters
 - Fixed rate of **rotations per minute (RPM)**

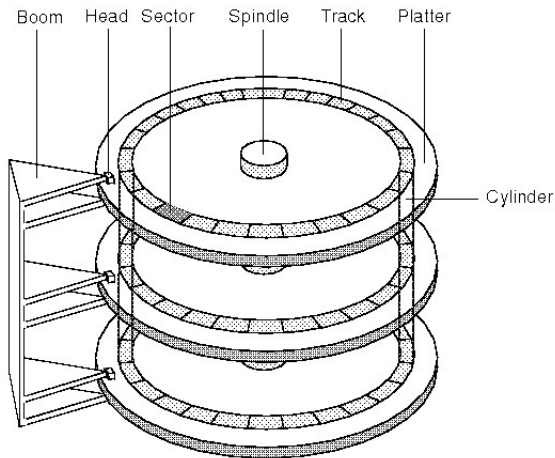


Basic Geometry

- Data is encoded in **tracks**
 - Concentric circles of sectors
 - Single surface contains thousands of tracks
- Read and write accomplished by **disk head**
 - One per surface
 - Attached to a **disk arm**
 - Moves across surface

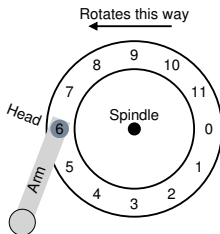


Basic Geometry



Single-track Latency

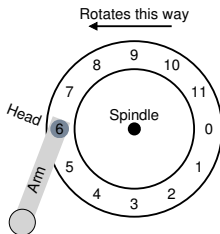
- Single track, with 12 sectors
- **Rotational delay**: wait for desired sector to reach disk head:



- Full rotational delay is R
 - Wait for sector 0?

Single-track Latency

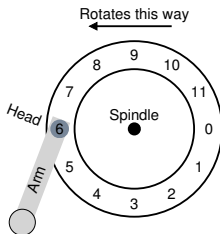
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 - Worst-case request?

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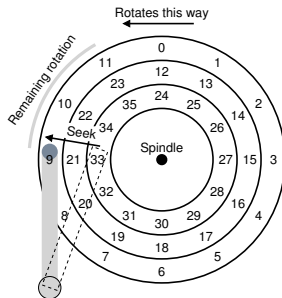
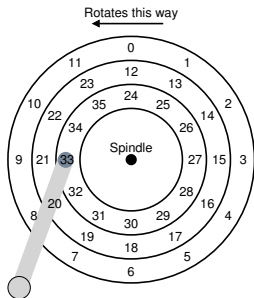
- Single track, with 12 sectors
- **Rotational delay**: wait for desired sector to reach disk head:



- Full rotational delay is R
 - Wait for sector 0? $\frac{R}{2}$
 - Worst-case request? sector 5 ($\frac{11R}{12}$)

Multiple Tracks

- **Seek:** move disk arm to the correct track
 - Costly disk operation, along with rotation
 - Acceleration: disk arm gets moving
 - Coasting: moving at full speed
 - Deceleration: arm slows down
 - Settling: head carefully positioned over correct track
 - **Settling time:** often quite significant, e.g., 0.5 to 2 ms



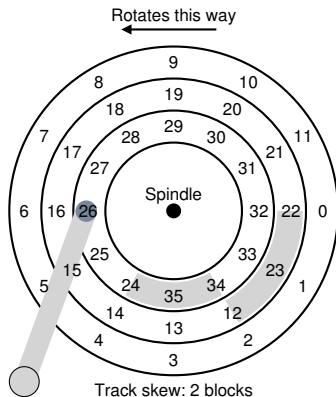
I/O Time

- Seek
- Wait for rotational delay
- **Transfer**: data is read from or written to surface

Other Details

- **Track skew**

- Switching tracks \rightarrow time to reposition the head
- Without skew, desired next block already rotated \rightarrow have to wait almost entire rotational delay



- **Multi-zoned** disk drives
 - Outer tracks have more sectors than inner tracks
 - Disk is organized into multiple zones
 - Each zone has the same number of sectors per track
 - Outer zones have more sectors than inner zones

- **Cache**

- Hold data read from or written to disk (8 to 64 MB)
- Quickly respond to requests
- e.g., read all sectors on a track and cache in memory
- **Write-through**
 - Acknowledge write when it's written to disk
- **Writeback**
 - Acknowledge write when data is in cache
 - Faster but dangerous: consistency issues (order not guaranteed)

I/O Time

$$T_{I/O} = T_{seek} + T_{rotation} + T_{transfer}$$

- Rate of I/O: $R_{I/O} = \frac{Size_{transfer}}{T_{I/O}}$
- **Random** workload
 - Small 4KB reads to random locations
- **Sequential** workload
 - Read 100MB of consecutive sectors

- Two example modern disks (Seagate):

	Cheetah 15K.5	Barracuda
Capacity	300 GB	1 TB
RPM	15,000	7,200
Average Seek	4 ms	9 ms
Max Transfer	125 MB/s	105 MB/s
Platters	4	4
Cache	16 MB	16/32 MB
Connects via	SCSI	SATA

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- On the Cheetah:

- $T_{seek} =$
- $T_{rotation} =$
- Random $T_{transfer} =$
- Seq. $T_{transfer} =$

- On the Barracuda:

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- $T_{seek} = 9ms$
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- On the Cheetah:

- $T_{seek} = 4ms$
- $T_{rotation} = 2ms$
- Random $T_{transfer} = 30\mu s$
- Seq. $T_{transfer} = 800ms$

- On the Barracuda:

- $T_{seek} = 9ms$
- $T_{rotation} = 4.2ms$
- Random $T_{transfer} = 38\mu s$
- Seq. $T_{transfer} = 950ms$

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	Cheetah 15K.5	Barracuda
$T_{I/O}$ Random		
$R_{I/O}$ Random		
$T_{I/O}$ Sequential		
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	Cheetah 15K.5	Barracuda
$T_{I/O}$ Random	6 ms	13.2 ms
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$T_{I/O}$ Random	6 ms	13.2 ms
$R_{I/O}$ Random	0.66 MB/s	0.31 MB/s
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$T_{I/O}$ Random	6 ms	13.2 ms
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$T_{I/O}$ Sequential	806 ms	963 ms
$R_{I/O}$ Sequential		

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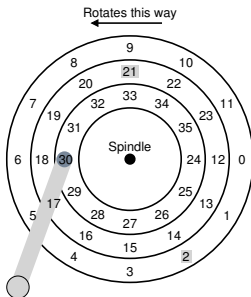
- Use disks sequentially!**

- **Disk scheduler**

- OS examines requests and decides which to schedule next
- Can make a good guess how long a job will take
 - By estimating seek and rotation delay
- Greedily pick least time to service first

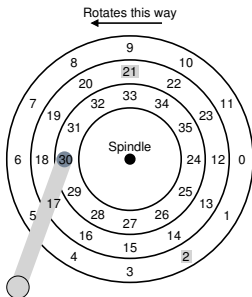
Disk Scheduling

- **Shortest Seek Time First (SSTF)**
 - Order I/O requests by track
 - Pick request on nearest track



Disk Scheduling

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 - Order I/O requests by track
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- Issue request to 21, then issue request to 2

Disk Scheduling

- Drive geometry not available to host OS
 - Sees an array of blocks
 - Solution?

Disk Scheduling

- Drive geometry not available to host OS
 - Sees an array of blocks
 - Solution? implement **nearest-block-first (NBF)**
- More fundamental problem?

Disk Scheduling

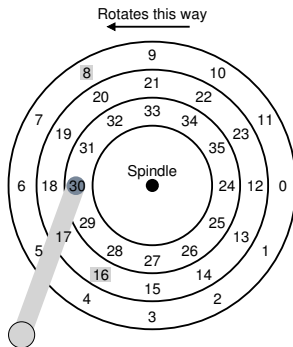
- Drive geometry not available to host OS
 - Sees an array of blocks
 - Solution? implement **nearest-block-first (NBF)**
- More fundamental problem?
- **Starvation**
 - Steady stream of requests to inner track
 - Other tracks ignored completely

Disk Scheduling

- **Elevator**

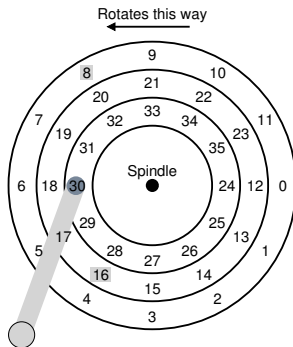
- Service requests in order across the tracks (back and forth)
- **Sweep**: single pass across the disk
 - Request of already-serviced track is queued until the next sweep
- **F-SCAN**: freeze queue when doing a sweep
 - Prevents starvation of far-away requests
- **C-SCAN**: sweep from outer-to-inner
 - Resets at outer track to begin again
 - Instead of both directions (favors middle tracks)
- Still problematic: ignores rotation

Disk Scheduling



- Schedule sector 16 or sector 8 next?

Disk Scheduling



- Schedule sector 16 or sector 8 next? it depends
 - Seek time much higher than rotational delay → SSTF
 - Seek faster than rotation → service request 8

Disk Scheduling

- Modern drives: seek and rotation times roughly equivalent
- **Shortest Positioning Time First (SPTF)**
 - Difficult to implement in OS
 - Usually performed inside a drive
 - OS picks best few requests and issues all to disk
- **I/O merging**
 - Series of requests - sectors 33, 8, then 34
 - OS merges 33 and 34 into a single two-block request
- **Work-conserving**
 - Wait before issuing I/O to disk
 - New and “better” request may arrive

Case Study

- IDE disk
 - Four types of registers:
 - Control, command block, status, and error
 - Available at specific “I/O addresses”
 - Using `in` and `out` instructions

Case Study

• The IDE interface:

Control Register:

Address 0x3F6 = 0x08 (0000 1RE0): R=reset,
E=0 means "enable interrupt"

Command Block Registers:

Address 0x1F0 = Data Port
Address 0x1F1 = Error
Address 0x1F2 = Sector Count
Address 0x1F3 = LBA low byte
Address 0x1F4 = LBA mid byte
Address 0x1F5 = LBA hi byte
Address 0x1F6 = 1B1D TOP4LBA: B=LBA, D=drive
Address 0x1F7 = Command/status

Status Register (Address 0x1F7):

7	6	5	4	3	2	1	0
BUSY	READY	FAULT	SEEK	DRQ	CORR	IDDEX	ERROR

Error Register (Address 0x1F1): (check when ERROR==1)

7	6	5	4	3	2	1	0
BBK	UNC	MC	IDNF	MCR	ABRT	T0NF	AMNF

BBK = Bad Block

UNC = Uncorrectable data error

MC = Media Changed

IDNF = ID mark Not Found

MCR = Media Change Requested

ABRT = Command aborted

T0NF = Track 0 Not Found

AMNF = Address Mark Not Found

Case Study

- **Wait for device to be ready:** read Status Register (0x1F7) until READY and not BUSY
- **Write parameters to command registers:** write the sector count, logical block address (LBA) of the sectors to be accessed, and drive number (master=0x00 or slave=0x10, as IDE permits just two drives) to command registers (0x1F2-0x1F6)
- **Start the I/O:** by issuing read/write to command register. Write READ—WRITE command to command register (0x1F7)
- **Data transfer (for writes):** wait until drive status is READY and DRQ (drive request for data); write data to data port
- **Handle interrupts:** in the simplest case, handle an interrupt for each sector transferred; more complex approaches allow batching and thus one final interrupt when the entire transfer is complete
- **Error handling:** after each operation, read the status register. If the ERROR bit is on, read the error register for details

xv6 ide driver: wait

```
static int ide_wait_ready() {  
    while (((int) r = inb(0x1f7)) & IDE_BSY) ||  
        !(r & IDE_DRDY));  
}
```

xv6 ide driver: start

```
static void ide_start_request(struct buf *b) {
    ide_wait_ready();
    outb(0x3f6, 0); // generate interrupt
    outb(0x1f2, 1); // how many sectors?
    outb(0x1f3, b->sector & 0xff); // LBA
    outb(0x1f4, (b->sector >> 8) & 0xff); // ...
    outb(0x1f5, (b->sector >> 16) & 0xff); // ...
    outb(0x1f6, 0xe0 | ((b->dev&1)<<4) |
               ((b->sector >> 24) & 0x0f));
    if (b->flags & B_DIRTY) {
        outb(0x1f7, IDE_CMD_WRITE); // WRITE
        outsl(0x1f0, b->data, 512/4); //
    } else {
        outb(0x1f7, IDE_CMD_READ); // this is a READ
                                   (no data)
    }
}
```

xv6 ide driver: rw

```
void ide_rw(struct buf *b) {
    acquire(&ide_lock);
    for (struct buf **pp = &ide_queue; *pp;
         pp = &(*pp)->qnext);

    *pp = b;

    if (ide_queue == b)
        ide_start_request(b); // send req to disk
    while ((b->flags & (B_VALID|B_DIRTY)) != B_VALID)
        sleep(b, &ide_lock); // wait for completion
    release(&ide_lock);
}
```

xv6 ide driver: isr

```
void ide_intr() {
    struct buf *b;
    acquire(&ide_lock);
    if (!(b->flags & B_DIRTY) &&
        ide_wait_ready() >= 0)
        insl(0x1f0, b->data, 512/4);
    b->flags |= B_VALID;
    b->flags &= ~B_DIRTY;
    wakeup(b); // wake waiting process
    if ((ide_queue = b->qnext) != 0) //
        ide_start_request(ide_queue);
    release(&ide_lock);
}
```

Summary (Hard Disk Drives)

- 512-byte sectors
 - **Platter** with two **surfaces**, bound around the **spindle**
 - Fixed rate of **RPM**
 - Data encoded in **tracks**, read and write by **disk head**
- **Rotational delay**: wait for sector to reach head
- **Seek**: move disk arm to correct track
 - Acceleration → coasting → deceleration → settling
- **I/O time**: seek → wait for rotational delay → transfer
- **Cache** holds read/write data
 - **Write-through**: acknowledge on write to disk
 - **Writeback**: acknowledge when data is in cache
- Disk scheduling
 - **SSTF, NBF, Elevator (sweep, F-SCAN, C-SCAN), SPTF**
 - **I/O merging**: merge requests for consecutive sectors
 - **Work-conserving**: wait before issuing I/O to disk