

Monads and their applications

Dr. Daniel Schäppi's course lecture notes

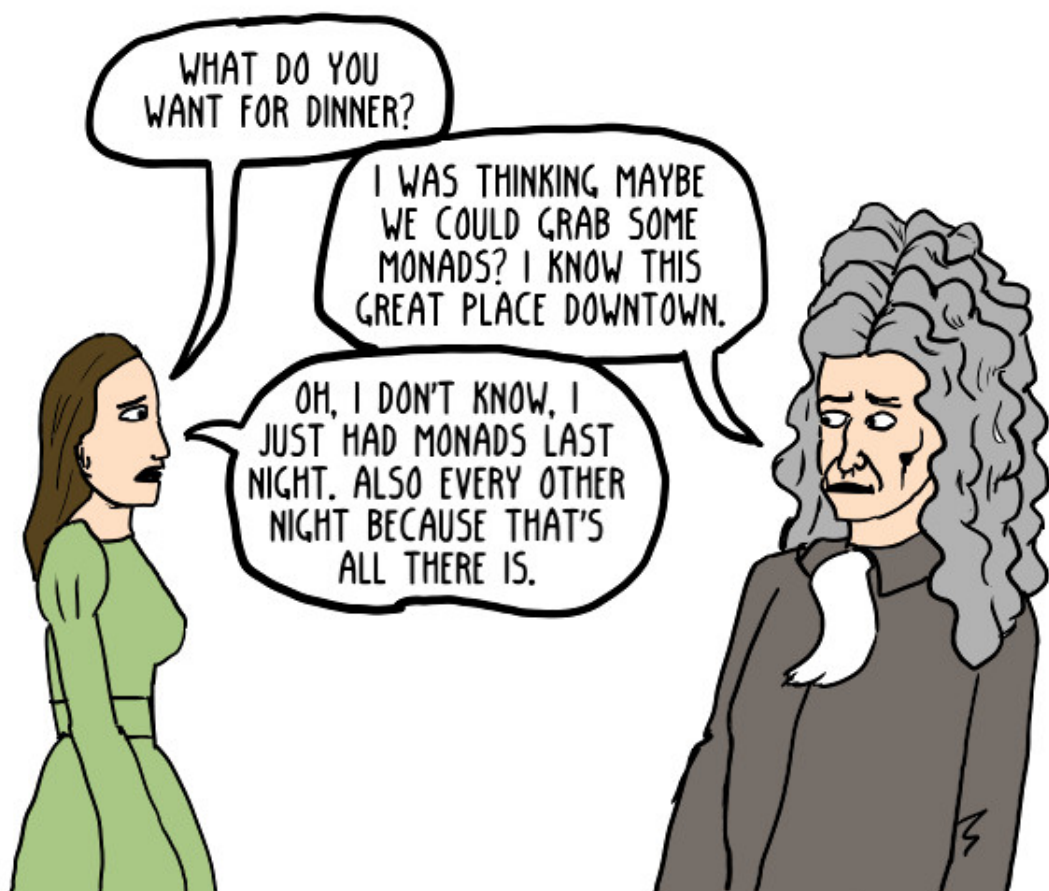
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Categorical preliminaries

Definition 0.0.1 (Categories). A *category* \mathcal{C} consists of:

1. a collection of objects $\text{Ob}(\mathcal{C})$;
2. a collection of arrows $\text{Ar}(\mathcal{C})$;
3. two maps $\text{dom}, \text{cod}: \text{Ar}(\mathcal{C}) \rightarrow \text{Ob}(\mathcal{C})$;
4. a map $\text{id}_-: \text{Ob}(\mathcal{C}) \rightarrow \text{Ar}(\mathcal{C})$ with $\text{dom}(\text{id}_c) = c = \text{cod}(\text{id}_c)$;
5. for every $f, g \in \text{Ar}(\mathcal{C})$ such that $\text{cod}(f) = \text{dom}(g)$ a unique composite morphism gf such that $\text{cod}(gf) = \text{cod}(g)$, $\text{dom}(gf) = \text{dom}(f)$.

This data has to satisfy the following axioms

1. given $f \in \text{Ar}(\mathcal{C})$, $c = \text{dom}(f)$ and $c' = \text{cod}(f)$, $\text{id}_{c'} f = f = f \text{id}_c$, that is the composition is unital;
2. given a composable triple $f, g, h \in \text{Ar}(\mathcal{C})$, $h(gf) = (hg)f$, that is the composition is associative.

An arrow f such that $c = \text{dom}(f)$ and $c' = \text{cod}(f)$ is denoted $f: c \rightarrow c'$.

Definition 0.0.2 (Functors).

Definition 0.0.3 (Full functors, faithful functor).

Definition 0.0.4 (Natural transformations).

Definition 0.0.5 (Equivalent functors).

Definition 0.0.6 (Representable Functors).

Definition 0.0.7 (Whiskering).

Definition 0.0.8 (Horizontal and vertical composition of nat.transf.).

Definition 0.0.9 (adjunctions).

Lemma 0.0.10 (Yoneda).

Proof.

□

We will denote by \mathcal{Y} (the kana for “Yo”) the Yoneda embedding $\mathcal{C} \hookrightarrow \mathbf{Set}^{\mathcal{C}^{\text{op}}}$.

1 Monads and algebras

1.1 Introduction

Throughout mathematics we encounter structures defined by some action morphisms. Here we give some examples.

Example 1.1.1. Given a group G , we may consider a G -set X described by an action map $G \times X \rightarrow X$.

Example 1.1.2. Given an abelian group M and a ring R , we can get an R -module M by fixing a group homomorphism $R \otimes_{\mathbb{Z}} M \rightarrow M$.

Example 1.1.3. Given a monoid M in **Set**, we get a map $\prod_{k=1}^n M \rightarrow M$, $(m_1, \dots, m_n) \mapsto ((\dots((m_1 m_2) m_3) \dots) m_{n-1}) m_n$. This induces an action map from $W(M) = \prod_{n \in \mathbb{N}} \prod_{k=1}^n M$, the set of words on M , to M .

Example 1.1.4. Given a set X , let $\mathcal{U}X$ be the set of ultrafilters on it. Any compact T2 topology on X allows us to see each ultrafilter as a system of neighborhoods of a unique point in X , hence it gives us a unique map $\mathcal{U}X \rightarrow X$ sending each ultrafilter to the respective point.

Example 1.1.5. Given a directed graph $D = (V, E, E \stackrel{s}{\rightharpoonup}_t V)$, we can create its free category FD , where the objects are the vertices and $FD(v, w) = \{\text{finite paths } v \rightarrow \dots \rightarrow w\}$. We set id_v to be the path of length 0, while composition is just the concatenation of paths.

In particular, if D is the directed graph with $V = \{0, \dots, n\}$ and an edge $j \rightarrow k$ if and only if $k = j + 1$, we have $FD \cong [n]$.

If $D = \{*\}$ and $E = \{* \rightarrow *\}$, then $FD(*, *) \cong \mathbb{N}$.

Given a small category \mathcal{C} , we may consider the underlying graph $U\mathcal{C} = D$ with $V = \text{Ob}(\mathcal{C})$, $E = \text{Ar}(\mathcal{C})$, $s = \text{dom}$ and $t = \text{cod}$. We get then an action map $UFU\mathcal{C} \rightarrow U\mathcal{C}$ sending a finite path to its composite. This map is a morphism of directed graphs.

Notice that we always have a category \mathcal{C} and some functor $T: \mathcal{C} \rightarrow \mathcal{C}$ with an action map $T\mathcal{C} \rightarrow \mathcal{C}$. How can we see all of these examples as specific instances of a general phenomenon?

Definition 1.1.6. A *monad* on a category \mathcal{C} is a triple (T, μ, η) where $T: \mathcal{C} \rightarrow \mathcal{C}$ is a functor, while $\mu: T^2 \Rightarrow T$ and $\eta: \text{id}_{\mathcal{C}} \Rightarrow T$ are natural transformations such that the following diagrams commute.

$$\begin{array}{ccc} T^3 & \xrightarrow{T\mu} & T^2 \\ \mu T \downarrow & & \downarrow \mu \\ T^2 & \xrightarrow{\mu} & T \end{array} \qquad \begin{array}{ccccc} T & \xrightarrow{\eta T} & T^2 & \xleftarrow{T\eta} & T \\ & \searrow \text{id}_T & \downarrow \mu & \swarrow \text{id}_T & \\ & & T & & \end{array}$$

μ is called the *multiplicative map*, while η is the *unit* of T .

The commutativity of the first diagram is equivalent to stating that the following two diagrams are equal.

$$\begin{array}{ccc}
& \mathcal{C} & \xrightarrow{T} \mathcal{C} \\
T \nearrow & \Downarrow \mu & \searrow T \\
\mathcal{C} & \xrightarrow{T} & \mathcal{C}
\end{array}
=
\begin{array}{ccc}
& \mathcal{C} & \xrightarrow{T} \mathcal{C} \\
T \nearrow & \Downarrow \mu & \searrow T \\
\mathcal{C} & \xrightarrow{T} & \mathcal{C}
\end{array}$$

On the other hand, the second diagram can be rephrased as follows:

$$\begin{array}{ccc}
& \mathcal{C} & \\
\Downarrow \eta \nearrow & & \searrow T \\
\mathcal{C} & \xrightarrow{T} & \mathcal{C}
\end{array}
=
T \left(\begin{array}{c} \mathcal{C} \\ \text{=} \\ \mathcal{C} \end{array} \right) T
=
\begin{array}{ccc}
& \mathcal{C} & \\
T \nearrow & \Downarrow \mu & \searrow T \\
\mathcal{C} & \xrightarrow{T} & \mathcal{C}
\end{array}$$

A monad naturally defines other algebraic structures, which we now introduce.

Definition 1.1.7. Given a monad (T, μ, η) , a T -algebra or T -module is a pair (a, α) , where $a \in \text{Ob}(\mathcal{C})$ and $\alpha: Ta \rightarrow a$ is such that the following diagrams commute.

$$\begin{array}{ccc}
T^2a & \xrightarrow{T\alpha} & Ta \\
\mu_a \downarrow & & \downarrow \alpha \\
Ta & \xrightarrow{\alpha} & a
\end{array}
\quad
\begin{array}{ccc}
a & \xrightarrow{\eta_a} & Ta \\
& \searrow & \downarrow \alpha \\
& & a
\end{array}$$

Definition 1.1.8. A morphism of T -algebras $(a, \alpha) \rightarrow (b, \beta)$ is a morphism $f: a \rightarrow b$ such that the following diagram commutes:

$$\begin{array}{ccc}
Ta & \xrightarrow{Tf} & Tb \\
\alpha \downarrow & & \downarrow \beta \\
a & \xrightarrow{f} & b
\end{array}$$

T -algebras form a category $T\text{-Alg}$, which has a natural forgetful functor $U^T: T\text{-Alg} \rightarrow \mathcal{C}$.

We now show how to recover the examples previously given with this language.

Example 1.1.9.

$$\begin{aligned}
T &= G \times -: \mathbf{Set} \rightarrow \mathbf{Set} \\
\mu_A &: G \times (G \times A) \rightarrow G \times A \\
&\quad (g, (h, a)) \mapsto (gh, a) \\
\eta_A &: A \rightarrow G \times A \\
&\quad a \mapsto (e, a)
\end{aligned}$$

is a monad and (A, α) is a T -algebra if and only if A is a G -set. It follows that $T\text{-Alg} \cong G\text{-Set}$.

Example 1.1.10. Given a ring R , $T = R \otimes_{\mathbb{Z}} -: \mathbf{Ab} \rightarrow \mathbf{Ab}$ is a monad when considered with the following natural transformations:

$$\begin{aligned}
\mu_- &: R \otimes_{\mathbb{Z}} (R \otimes_{\mathbb{Z}} -) \cong (R \otimes_{\mathbb{Z}} R) \otimes_{\mathbb{Z}} - \Rightarrow R \otimes_{\mathbb{Z}} - \\
\eta_- &: - \cong \mathbb{Z} \otimes_{\mathbb{Z}} - \Rightarrow R \otimes_{\mathbb{Z}} -
\end{aligned}$$

We have that $(R \otimes_{\mathbb{Z}} -)\text{-Alg} \cong \mathbf{Mod}_R$.

Example 1.1.11. Consider $W: \mathbf{Set} \rightarrow \mathbf{Set}$ given by $WX = \coprod_{n \in \mathbb{N}} \coprod_{k=1}^n X$. Multiplication $\mu_X: WWX \rightarrow WX$ is given by concatenation of words, while the unit $\eta_X: X \rightarrow WX$ is just $x \mapsto (x)$. With this, $W\text{-Alg} \cong \text{Mon}(\mathbf{Set})$.

Example 1.1.12. The functor \mathcal{U} defined in Example 1.1.4, equipped with suitable natural transformations, is a monad on \mathbf{Set} and $\mathcal{U}\text{-Alg} \cong \mathbf{CHTop}$, the category of compact T2 spaces.

Example 1.1.13. The free-forgetful adjunction $F \dashv U$ between categories and directed graphs induces a monad on the latter, with $UF\text{-Alg} \cong \mathbf{Cat}$.

1.2 Monadic functors

Now that we have introduced these structures, our aim is to characterize *monadic functors*, that is functors $U: \mathcal{A} \rightarrow \mathcal{C}$ which are equivalent to $U^T: T\text{-Alg} \rightarrow \mathcal{C}$ for some monad (T, μ, η) on \mathcal{C} .

First of all, notice that U^T is faithful by construction, hence U must be faithful, but more is true.

Lemma 1.2.1. The functor U^T is conservative, that is if $U^T f$ is an isomorphism then f is an isomorphism of T -algebras.

Proof. Suppose that g is the inverse of $f: a \rightarrow b$ and f is a morphism $(a, \alpha) \rightarrow (b, \beta)$. We only need to prove that the square on the left commutes, that is $g\beta = \alpha Tg$:

$$\begin{array}{ccccc} Tb & \xrightarrow{Tg} & Ta & \xrightarrow{Tf} & Tb \\ \beta \downarrow & & \alpha \downarrow & & \downarrow \beta \\ b & \xrightarrow{g} & a & \xrightarrow{f} & b \end{array}$$

We see that $fg\beta = \beta$ and $f\alpha Tg = \beta T f T g = \beta T(fg) = \beta T \text{id}_b = \beta$, hence the thesis. \square

Remark 1.2.2. Notice that the forgetful functor $U: \mathbf{Top} \rightarrow \mathbf{Set}$ can't be monadic since it does not reflect isomorphisms. However, if we restrict it to the full subcategory of \mathbf{Top} spanned by compact Hausdorff spaces we indeed obtain a monadic functor.

Proposition 1.2.3. The functor $U^T: T\text{-Alg} \rightarrow \mathcal{C}$ has a left adjoint $F^T: \mathcal{C} \rightarrow T\text{-Alg}$ such that $F^T c = (Tc, \mu_c)$, $F^T f: (Tc, \mu_c) \xrightarrow{Tf} (Td, \mu_d)$ and $U^T F^T = T$. Furthermore, the unit of this adjunction is given by $\gamma_c = \eta_c: c \rightarrow U^T F^T c = Tc$ and the counit has components $\epsilon_{(a, \alpha)} = \alpha: (Ta, \mu_a) \rightarrow (a, \alpha)$.

Proof. (i) To show that (Tc, μ_c) is a T -algebra we need the following diagrams to be commutative.

$$\begin{array}{ccc} T^3 c & \xrightarrow{T\mu_c} & T^2 c \\ \mu_{Tc} \downarrow & & \downarrow \mu_c \\ T^2 c & \xrightarrow{\mu_c} & Tc \end{array} \quad \begin{array}{ccc} Tc & \xrightarrow{\eta_{Tc}} & T^2 c \\ & \searrow & \downarrow \mu_c \\ & & Tc \end{array}$$

These are exactly the associativity and one of the unit laws for (T, μ, η) .

- (ii) For every $f: c \rightarrow c'$, Tf is a morphism of algebras $(Tc, \mu_c) \rightarrow (Tc', \mu_{c'})$. The diagram

$$\begin{array}{ccc} T^2c & \xrightarrow{T^2f} & T^2c' \\ \mu_c \downarrow & & \downarrow \mu_{c'} \\ Tc & \xrightarrow{Tf} & Tc' \end{array}$$

is commutative because of the naturality of μ . Hence F^T is defined on morphisms. It is a functor by the functoriality of T .

- (iii) The unit is natural by assumption. We claim that $\epsilon_{(a, \alpha)} = \alpha$ is a morphism of algebras

$$F^T U^T(a, \alpha) = F^T a = (Ta, \mu_a) \rightarrow \text{id}_{T\text{-Alg}}(a, \alpha) = (a, \alpha)$$

and ϵ is a natural transformation $F^T U^T \Rightarrow \text{id}_{T\text{-Alg}}$. Let's check it. We know that α is a morphism of algebras if and only if

$$\begin{array}{ccc} T^2a & \xrightarrow{T\alpha} & Ta \\ \mu_a \downarrow & & \downarrow \alpha \\ Ta & \xrightarrow{\alpha} & a \end{array}$$

is commutative. But this is one of the two T -algebra axioms! Moreover, to prove that ϵ is natural, we need to show that

$$\begin{array}{ccc} (Ta, \mu_a) & \xrightarrow{\alpha = \epsilon_{(a, \alpha)}} & (a, \alpha) \\ Tf \downarrow & & \downarrow f \\ (Tb, \mu_b) & \xrightarrow{\beta = \epsilon_{(b, \beta)}} & (b, \beta) \end{array}$$

is commutative, but this is the axiom for f to be a morphism of T -algebras!

- (iv) It remains to check the two triangular identities $\epsilon F^T \cdot F^T \eta = \text{id}_{F^T}$ and $U^T \epsilon \cdot \eta U^T = \text{id}_{U^T}$. These are to be checked on the components at c and (a, α) , respectively.

$$\begin{array}{ccc} (Tc, \mu_c) & \xrightarrow{T\eta_c} & (T^2c, \mu_{Tc}) \\ & \searrow & \downarrow \mu_{Tc} \\ & & (Tc, \mu_c) \end{array} \quad \begin{array}{ccc} a & \xrightarrow{\eta_a} & Ta \\ & \searrow & \downarrow \alpha \\ & & a \end{array}$$

The commutativity of these diagrams is ensured by the second unit law for a monad and the unit law for the T -algebra (a, α) , respectively. \square

Definition 1.2.4. Algebras of the form (Tc, μ_c) are called *free T -algebras*.

Thanks to the proposition above we can prove that, given a monad T we can always find an adjunction that generates it. Actually, the converse holds too.

Proposition 1.2.5. If $U: \mathcal{D} \rightarrow \mathcal{C}$ has a left adjoint F with unit η and counit ϵ , then $(UF, U\epsilon F, \eta)$ is a monad on \mathcal{C} . Also, if (T, μ, η) is a monad on \mathcal{C} , then $(U^T F^T, U^T \epsilon F^T, \eta) = (T, \mu, \eta)$.

Proof. Let us check the axioms. First of all, the associativity holds due to the following equations.

$$\begin{aligned}
& \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \\
& \quad \quad \quad \begin{array}{c} \mathcal{C} \\ \nearrow U \quad \searrow F \\ \Downarrow \epsilon \end{array} \\
& = \\
& \quad \quad \quad \begin{array}{c} \mathcal{C} \\ \nearrow U \quad \searrow F \\ \Downarrow \epsilon \end{array} \\
& = \\
& \quad \quad \quad \begin{array}{c} \mathcal{C} \\ \nearrow U \quad \searrow F \\ \Downarrow \epsilon \end{array}
\end{aligned}$$

Unit laws:

$$\begin{array}{c} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \\ \quad \quad \quad \begin{array}{c} \mathcal{C} \\ \nearrow U \quad \searrow F \\ \Downarrow \epsilon \end{array} \end{array}$$

is equal to 1_{UF} , since $\epsilon F \cdot F\eta = 1_F$ by one of the triangular identities of the adjunction $F \dashv U$. Furthermore,

$$\begin{array}{c} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \xrightarrow{F} \mathcal{D} \xrightarrow{\quad} \mathcal{C} \\ \quad \quad \quad \begin{array}{c} \mathcal{C} \\ \nearrow U \quad \searrow F \\ \Downarrow \epsilon \end{array} \end{array}$$

is equal to 1_{UF} . This follows from the explicit description of the unit and the counit of the adjunction $F^T \dashv U^T$, in fact $U^T \epsilon F^T c = U^T \epsilon_{(Tc, \eta c)} = \mu_c$. \square

Example 1.2.6 (Interesting adjunction, boring monad). Let us consider the adjunction $U: \mathbf{Top} \rightleftarrows \mathbf{Set}: \text{Disc} = F$, whose left adjoint assigns to every set X the discrete topological space $FX = (X, 2^X)$. It's immediate to see that $UF X = X$, hence $UF = \text{id}_{\mathbf{Set}}$. How many natural transformations $\text{id}_{\mathbf{Set}} = UF \xrightarrow{\alpha} UF = \text{id}_{\mathbf{Set}}$ are there? We know that $\text{id}_{\mathbf{Set}} \cong \text{Hom}(*, -)$, so $\text{Nat}(\text{id}_{\mathbf{Set}}, \text{id}_{\mathbf{Set}}) \cong \text{Nat}(\text{Hom}(*, -), \text{Hom}(*, -)) \cong \text{Hom}(*, *) = \{\text{id}_*\}$ by Yoneda, hence $\alpha = \text{id}$. Therefore $(UF, U\epsilon F, \eta) = (\text{id}_{\mathbf{Set}}, \text{id}, \text{id})$.

Example 1.2.7. If S is a set, $\mathbf{Set}(S, -): \mathbf{Set} \rightarrow \mathbf{Set}$ is right adjoint to $S \times -: \mathbf{Set} \rightarrow \mathbf{Set}$, so we get a monad $X \mapsto \mathbf{Set}(S, S \times X)$. This is called *the state monad* and is important in Computer Science.

There is always a comparison morphism $\mathcal{D} \xrightarrow{\bar{U}} UF\text{-Alg}$ such that

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{\bar{U}} & UF\text{-Alg} \\ & \searrow U & \swarrow U^{UF} \\ & \mathcal{C} & \end{array}$$

commutes. We set $\bar{U}f = (Ud, UFUd \xrightarrow{U\epsilon_d} Ud) = (Ud, U\epsilon_d)$. More specifically, for a given functor $G: \mathcal{D} \rightarrow \mathcal{C}$ we can ask what do we need to get an equivalence $\bar{G}: \mathcal{D} \rightarrow T\text{-Alg}$. To get there, we will need a few more definitions and lemmas.

1.3 The category of T -actions

Just like a monad (T, μ, η) defines a category $T\text{-Alg}$, it also allows us to construct another category from functors $\mathcal{D} \rightarrow \mathcal{C}$.

Definition 1.3.1. Given a monad (T, μ, η) on a category \mathcal{C} and fixed another category \mathcal{D} , a T -action on a functor $G: \mathcal{D} \rightarrow \mathcal{C}$ is a natural transformation $\gamma: TG \Rightarrow G$ such that the diagrams

$$\begin{array}{ccc} T^2G & \xrightarrow{T\gamma} & TG \\ \mu G \downarrow & & \downarrow \gamma \\ TG & \xrightarrow{\gamma} & G \end{array} \quad \begin{array}{ccc} G & \xrightarrow{\eta G} & TG \\ & \searrow & \downarrow \gamma \\ & & G \end{array}$$

commute.

A morphism of T -actions $(G, \gamma) \xrightarrow{\varphi} (K, \kappa)$ is a natural transformation $\varphi: G \Rightarrow K$ such that

$$\begin{array}{ccc} TG & \xrightarrow{T\varphi} & TK \\ \gamma \downarrow & & \downarrow \kappa \\ G & \xrightarrow{\varphi} & K \end{array}$$

commutes.

Up to size, T -actions and their morphisms assemble into a category $T\text{-Act}(\mathcal{D})$.

Example 1.3.2. The functor $U^T: T\text{-Alg} \rightarrow \mathcal{C}$ has a T -action given by $(U^T, \alpha: TU^T \Rightarrow U^T)$, where $\alpha_{(b, \beta)} := \beta: Tb \rightarrow b$.

Example 1.3.3. Given an adjunction $F \dashv U: \mathcal{C} \rightleftarrows \mathcal{D}$ with unit $\eta: \text{id}_{\mathcal{C}} \Rightarrow UF$ and counit $\epsilon: FU \Rightarrow \text{id}_{\mathcal{D}}$, we get a monad on $(UF, U\epsilon F, \eta)$ on \mathcal{C} . We have then a UF -action $U\epsilon: UFU \Rightarrow U$, where the axioms follow from the triangular identities and the naturality of $U\epsilon$.

Proposition 1.3.4. (U^T, α) is the universal T -action, that is for any category \mathcal{D} the functor $\text{Cat}(\mathcal{D}, T\text{-Alg}) \rightarrow T\text{-Act}(\mathcal{D})$ sending G to $(U^T G, \alpha G)$ and $\beta: G \Rightarrow H$ to $U^T \beta: (U^T G, \alpha G) \Rightarrow (U^T H, \alpha H)$ is an isomorphism of categories.

Proof. In other words, for every T -action (G, γ) there exists a unique lift $\overline{G}: \mathcal{D} \rightarrow T\text{-Alg}$ such that $(U^T \overline{G}, \alpha \overline{G}) = (G, \gamma)$ and for every $\phi: (G, \gamma) \Rightarrow (K, \kappa)$ there is a unique $\overline{\phi}: \overline{G} \Rightarrow \overline{K}$ with $U^T \overline{\phi} = \phi$.

It is enough to set $\overline{G}d := (Gd, \gamma_d)$ on objects, $\overline{G}f := Gf$ on morphisms, $\overline{\phi}_d := \phi_d$ and check the axioms.

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{\exists! \overline{G}} & T\text{-Alg} \\ & \searrow G & \downarrow U^T \\ & & \mathcal{C} \end{array}$$

□

Following the construction in this proof, from the last example we get the comparison functor for the adjunction $F \dashv U$. In particular, $\overline{U}d = (Ud, U\epsilon_d)$. Furthermore, this means that $U: \mathbf{Top} \rightarrow \mathbf{Set}$ factors through identities.

1.4 Limits and colimits in the category of algebras

We have shown that the forgetful functor $U^T: T\text{-Alg} \rightarrow \mathcal{C}$ is a right adjoint, and as such it preserves limits. However, more is true.

Proposition 1.4.1. For any monad (T, μ, η) on \mathcal{C} , the forgetful functor $U^T: T\text{-Alg} \rightarrow \mathcal{C}$ strictly creates limits.

Proof. This statement means that, for any diagram $D: I \rightarrow T\text{-Alg}$ such that $U^T D: I \rightarrow \mathcal{C}$ has a limit (l, κ_i) in \mathcal{C} , there is a unique T -algebra structure $\lambda: Tl \rightarrow l$ such that κ_i is a morphism of T -algebras for all $i \in I$ and this makes $((l, \lambda), \kappa_i)$ into a limit of D .

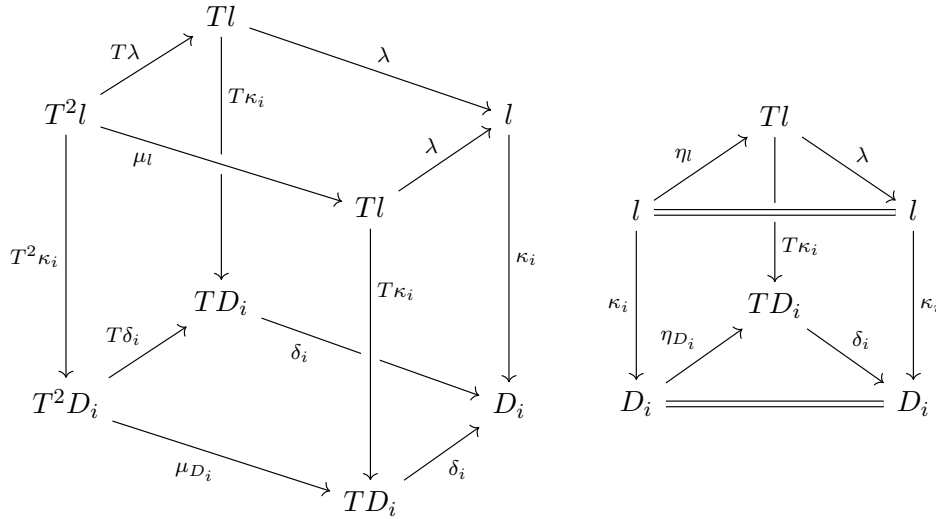
Now we begin the proof.

First of all, remember that $D\phi: D_i \rightarrow D_j$ is a morphism of T -algebras for all $\phi: i \rightarrow j$ by assumption, hence the morphisms $\delta_i T\kappa_i: Tl \rightarrow D_i$ define a cone over D , where δ_i is the T -algebra structure on D_i . By the universal property of the limit, there is a unique morphism $\lambda: Tl \rightarrow l$ making the following diagram commute for all i .

$$\begin{array}{ccc} Tl & \xrightarrow{T\kappa_i} & TD_i \\ \lambda \downarrow & & \downarrow \delta_i \\ l & \xrightarrow{\kappa_i} & D_i \end{array}$$

This tells us that, if the limit $((l, \lambda), \kappa_i)$ of D exists, it is unique. We have to check that (l, λ) is a T -algebra.

Notice that for all i all of the faces of the following diagrams, except for possibly the top ones, commute.



Since the κ_i are jointly monic, the upper face commutes and therefore (l, λ) is a T -algebra. It remains to check that $((l, \lambda), \kappa_i)$ factors every other cone over D .

Let $\gamma_i: (x, \zeta) \rightarrow (D_i, \delta_i)$ be a cone over D . Then, there is a unique $f: x \rightarrow l$ in \mathcal{C} such that $\kappa_i f = \gamma_i$. We only have to show that f is a morphism of T -algebras $(x, \zeta) \rightarrow (l, \lambda)$.

Consider the following diagram and notice that the outer square, the one on the right and the two triangles commute, hence the square on the left commutes as well since the κ_i are jointly monic.

$$\begin{array}{ccccc}
& & T\gamma_i & & \\
& \nearrow & & \searrow & \\
Tx & \xrightarrow{Tf} & Tl & \xrightarrow{T\kappa_i} & TD_i \\
\downarrow \zeta & & \downarrow \lambda & & \downarrow \delta_i \\
x & \xrightarrow{f} & l & \xrightarrow{\kappa_i} & D_i \\
& \searrow & & \nearrow & \\
& & \gamma_i & &
\end{array}$$

□

A similar statement holds for colimits.

Proposition 1.4.2. Given a monad (T, μ, η) on \mathcal{C} , the forgetful functor $U^T: T\text{-Alg} \rightarrow \mathcal{C}$ strictly creates any colimit preserved by both T and T^2 .

Proof. Similarly to the dual situation, this means that for any diagram $D: I \rightarrow T\text{-Alg}$ such that $U^T D: I \rightarrow \mathcal{C}$ has a colimit (c, κ_i) preserved by both T and T^2 , there is a unique T -algebra structure $\lambda: Tc \rightarrow c$ such that κ_i is a morphism of T -algebras for all $i \in I$. This makes $((c, \lambda), \kappa_i)$ into a colimit of D .

The proof is essentially dual to the one given earlier, in the sense that we find again a unique $\lambda: Tc \rightarrow c$ using the universal property of the colimit $(Tc, T\kappa_i)$ of TD .

$$\begin{array}{ccc}
TD_i & \xrightarrow{T\kappa_i} & Tc \\
\delta_i \downarrow & & \downarrow \lambda \\
D_i & \xrightarrow{\kappa_i} & c
\end{array}$$

To check that (c, λ) is an algebra we use the universal property of $(T^2c, T^2\kappa_i)$, for μ , and the one of (c, κ_i) , for η . □

Remark 1.4.3. The same statements hold for monadic functors, except for the fact that they might not create limits and colimits strictly since they are just equivalent to a U^T .

Remark 1.4.4. If T is a monad on a complete category \mathcal{C} , then $T\text{-Alg}$ is complete. If \mathcal{C} is cocomplete and T is cocontinuous, then $T\text{-Alg}$ is cocomplete.

Example 1.4.5. Let \mathcal{C} be a small category. There is a cocontinuous monad on the category of $\text{Ob}(\mathcal{C})$ -indexed collections of sets whose category of algebras is the functor category $[\mathcal{C}, \mathbf{Set}]$. The underlying endofunctor of this monad is defined as

$$\begin{aligned}
T: [\text{Ob}(\mathcal{C}), \mathbf{Set}] &\rightarrow [\text{Ob}(\mathcal{C}), \mathbf{Set}] \\
(X_c)_{c \in \mathcal{C}} &\mapsto \left(\coprod_{d \in \mathcal{C}} \mathcal{C}(d, c) \times X_d \right)_{c \in \mathcal{C}}
\end{aligned}$$

Since $[\text{Ob}(\mathcal{C}), \mathbf{Set}]$ is complete and cocomplete, so is $[\mathcal{C}, \mathbf{Set}]$ (with limits and colimits computed pointwise).

2 Beck's monadicity theorem

The final ingredient we need is the observation that T -algebras admit canonical presentations using free algebras.

2.1 Split coequalizers and Beck's theorem

Example 2.1.1. Pick an epi $F \twoheadrightarrow G$ in the category of groups **Grp**, where F is a free group. The kernel of this homomorphism defines a (normal) subgroup K of F , giving rise to the sequence $K \hookrightarrow F \twoheadrightarrow G$. We can take another epi $F' \twoheadrightarrow K$, with F' again a free group. Therefore G is the cokernel of some morphism $F' \rightarrow F$. This argument applies to rings, algebras etc.

It is natural to ask if we can do this systematically for general T -algebras. Given (a, α) in $T\text{-Alg}$, we have $F^T U^T(a, \alpha) \rightarrow (a, \alpha)$ i.e. $(Ta, \mu_a) \xrightarrow{\alpha} (a, \alpha)$. A candidate¹ for F' would be $F^T U^T(Ta, \mu_a) = (T^2a, \mu_{Ta})$. What are the “elements” of Ta ? Notice that

$$(T^2a, \mu_{Ta}) \xrightarrow[\mu_a]{T\alpha} (Ta, \mu_a) \xrightarrow{\alpha} (a, \alpha)$$

is a well defined diagram in $T\text{-Alg}$, with $\alpha\mu_a = \alpha T\alpha$. Moreover, this is a coequalizer. We can use Proposition 1.4.2 to prove it, so that we need to check whether U^T sends the diagram above into a coequalizer preserved by T and T^2 . In \mathcal{C} , we get the diagram

$$\begin{array}{ccccc} T^2a & \xrightarrow[\mu_a]{T\alpha} & Ta & \xrightarrow{\alpha} & a \\ \eta_{Ta} \swarrow & & \nwarrow \eta_a & & \\ & & & & \end{array}$$

in which the following equations hold true: $\alpha T\alpha = \alpha\mu_a$, $\alpha\eta_a = \text{id}_a$, $\mu_a\eta_{Ta} = \text{id}_{Ta}$ and $\eta_a\alpha = T\alpha\eta_{Ta}$ by naturality. It is a particular case of a more general concept.

Definition 2.1.2. A *split coequalizer* is a diagram of the form

$$\begin{array}{ccccc} a & \xrightarrow[\underset{t}{\curvearrowright}]{\underset{g}{\curvearrowleft}} f & b & \xrightarrow[\underset{s}{\curvearrowleft}]{\underset{h}{\curvearrowright}} & c \end{array}$$

so that $hf = hg$, $hs = \text{id}_c$, $gt = \text{id}_b$, and $ft = sh$.

Proposition 2.1.3. In the above situation,

$$a \xrightarrow[\underset{g}{\curvearrowright}]{\underset{f}{\curvearrowleft}} b \xrightarrow{h} c$$

is a coequalizer. In particular, any functor preserves this coequalizer.

¹Think about free groups: in that case we take words on Ta .

Proof. Take $k: b \rightarrow d$ such that $kf = kg$ and define $\bar{k} := ks$. Then we have

$$\bar{k}h = ksh = kft = kgt = k.$$

Uniqueness is clear since h is a (split) epi. □

T and T^2 preserve split coequalizers, so they preserve our coequalizer.

Corollary 2.1.4. Let T be a monad on \mathcal{C} and (a, α) a T -algebra. Then

$$(T^2a, \mu_{Ta}) \xrightarrow[\mu_a]{T\alpha} (Ta, \mu_a) \xrightarrow{\alpha} (a, \alpha)$$

is a coequalizer in $T\text{-Alg}$, which $U^T: T\text{-Alg} \rightarrow \mathcal{C}$ sends to a split coequalizer in \mathcal{C} .

Proof. We have already observed that the second statement holds, so that $\text{coeq}(U^T(T\alpha), U^T(\mu_a))$ is preserved by T and T^2 . Hence there exists a unique lift of the (split) coequalizer in \mathcal{C} to a coequalizer in $T\text{-Alg}$. □

Results like the previous one inspire us to look at the parallel pairs of morphisms in a category which are sent to split coequalizers or, to say it better, to a parallel pair of morphisms that can be extended to a split coequalizer diagram. This kind of pairs will be of crucial importance in the following.

Definition 2.1.5. Let $U: \mathcal{D} \rightarrow \mathcal{C}$ be a functor. A pair of morphisms $f, g: d \rightrightarrows d'$ in \mathcal{D} is *U-split* if $Uf, Ug: Ud \rightrightarrows Ud'$ is part of a split coequalizer in \mathcal{C} .

Remark 2.1.6. $T\alpha, \mu_a: (T^2a, \mu_{Ta}) \rightrightarrows (Ta, \mu_a)$ is a U^T -split pair. Moreover, $T\text{-Alg}$ has coequalizers of U^T -split pairs and U^T preserves them. Hence, functors equivalent to U^T satisfy three conditions:

1. they have a left adjoint;
2. they are conservative;
3. U -split pairs have coequalizers which are preserved by U .

Theorem 1 (Beck). Let $U: \mathcal{D} \rightarrow \mathcal{C}$ be a right adjoint to $F: \mathcal{C} \rightarrow \mathcal{D}$. Let $(T = UF, U\epsilon F, \eta)$ be the induced monad and $\bar{U}: \mathcal{D} \rightarrow T\text{-Alg}$ be the comparison functor.

1. If \mathcal{D} has coequalizers of U -split pairs, then \bar{U} has a left adjoint $\bar{F}: T\text{-Alg} \rightarrow \mathcal{D}$;
2. if, in addition, U preserves coequalizers of U -split pairs, the unit $\bar{\eta}: \text{id}_{T\text{-Alg}} \Rightarrow \bar{U}\bar{F}$ is an isomorphism;
3. if U is also conservative, then \bar{U} is an equivalence of categories.

Proof. 1. For each free T -algebra (Ta, μ_a) we have

$$\begin{aligned} T\text{-Alg}((Ta, \mu_a), \bar{U}-) &= T\text{-Alg}(F^T a, \bar{U}-) \\ &\cong \mathcal{C}(a, U^T \bar{U}-) \\ &= \mathcal{C}(a, U-) \\ &\cong \mathcal{D}(Fa, -) \end{aligned}$$

therefore the value of \overline{F} at (Ta, μ_a) has to be Fa . Since every T -algebra is a coequalizer of free algebras which is preserved by U^T , we may define $\overline{F}(a, \alpha)$ as the coequalizer of a pair of morphisms $F Ta \rightrightarrows Fa$. We write this as $FUFU^T(a, \alpha) \rightrightarrows FU^T(a, \alpha)$. Consider the following pair of morphisms of functors

$$FUFU^T \xrightleftharpoons[\epsilon FU^T]{F\alpha} FU^T$$

in the functor category $[T\text{-Alg}, \mathcal{D}]$. We claim that this pair has a coequalizer and $\overline{F}: T\text{-Alg} \rightarrow \mathcal{D}$ is left adjoint to \overline{U} . Note that the pair of morphisms just above becomes split after the composition with $U: \mathcal{D} \rightarrow \mathcal{C}$. In fact

$$\begin{array}{ccc} U F U F U^T & \xrightleftharpoons[U \epsilon F U^T]{U F \alpha} & U F U^T \xrightarrow{\alpha} U^T \\ \nwarrow \eta U F U^T & & \nwarrow \eta U^T \end{array}$$

is a split coequalizer in $[T\text{-Alg}, \mathcal{C}]$, given that it holds pointwise since $UF = T$. Let us denote by $\beta: FU^T \rightarrow \overline{F}$ the colimit (computed pointwise) of the pair $F\alpha, \epsilon FU^T: FUFU^T \rightrightarrows FU^T$. Precomposing this pair with \overline{U} and recalling that $\alpha\overline{U} = U\epsilon$, $U^T\overline{U} = U$, we get the pair

$$FUFU \xrightleftharpoons[\epsilon FU]{FU\epsilon} FU,$$

which is coequalized by $\epsilon: FU \Rightarrow \text{id}_{\mathcal{D}}$.

$$\begin{array}{ccc} FUFU & \xrightleftharpoons[\epsilon FU]{FU\epsilon} & FU \xrightarrow{\beta\overline{U}} \overline{F}\overline{U} \\ & \searrow \epsilon & \downarrow \exists! \bar{\epsilon} \\ & & \text{id}_{\mathcal{D}} \end{array}$$

Since $\overline{F}\overline{U}$ is the coequalizer of the diagram above, there exists a unique $\bar{\epsilon}: \overline{F}\overline{U} \Rightarrow \text{id}_{\mathcal{D}}$ such that $\bar{\epsilon} \cdot \beta\overline{U} = \epsilon$. To get the unit $\bar{\eta}: \text{id}_{T\text{-Alg}} \Rightarrow \overline{U}\overline{F}$ we need to describe a morphism of T -actions $(U^T, \alpha) \rightarrow (U^T\overline{U}\overline{F}, \alpha\overline{U}\overline{F})$. We claim that the natural transformation induced by the universal property of the split coequalizer in the first row

$$\begin{array}{ccccc} U F U F U^T & \xrightleftharpoons[U \epsilon F U^T]{U F \alpha} & U F U^T & \xrightarrow{\alpha} & U^T \\ \parallel & & \parallel & & \downarrow \exists! \bar{\eta} \\ U^T \overline{U} F U F U^T & \xrightleftharpoons[U^T \overline{U} \epsilon F U^T]{U^T \overline{U} F \alpha} & U^T \overline{U} F U^T & \xrightarrow{U^T \overline{U} \beta} & U^T \overline{U} \overline{F} \end{array}$$

is a morphism of T -actions².

Unraveling what this means, we have to check that the diagram

$$\begin{array}{ccc} U F a & \xrightarrow{U F \bar{\eta}(a, \alpha)} & U F U \overline{F}(a, \alpha) \\ \alpha \downarrow & & \downarrow U \epsilon_{\overline{F}(a, \alpha)} \\ a & \xrightarrow{\bar{\eta}(a, \alpha)} & U \overline{F}(a, \alpha) \end{array}$$

²In fact, this tells us that the morphism $\bar{\eta}_{(a, \alpha)}: a \rightarrow U^T \overline{U} \overline{F}(a, \alpha)$ in \mathcal{C} lifts uniquely to a morphism of T -algebras $\bar{\eta}_{(a, \alpha)}: (a, \alpha) \rightarrow \overline{U} \overline{F}(a, \alpha)$.

is commutative. We know that $\bar{\eta}\alpha = U\beta$ by the definition of $\bar{\eta}$. Moreover, α is a split epi in \mathcal{C} , hence we can precompose with $UF\alpha$ (again a split epi) and check the commutativity of the resulting diagram. We get the diagram

$$\begin{array}{ccccc}
 & & & & U\epsilon_{Fa} \\
 & & & & \curvearrowright \\
 UFUFa & & & & \\
 \downarrow UF\alpha & & & & \downarrow U\epsilon_{Fa} \\
 & & UFU\beta_{(a,\alpha)} \text{ nat. of } \epsilon & & \\
 & & \downarrow & & \\
 UFa & \xrightarrow{UF\bar{\eta}_{(a,\alpha)}} & UFUF(a,\alpha) & & UFa \\
 \downarrow \alpha & \searrow U\beta_{(a,\alpha)} & \downarrow U\epsilon_{\bar{F}(a,\alpha)} & & \downarrow U\beta_{(a,\alpha)} \\
 a & \xrightarrow{\bar{\eta}_{(a,\alpha)}} & U\bar{F}(a,\alpha) & \xleftarrow{U\beta_{(a,\alpha)}} &
 \end{array}$$

The definition of β as a coequalizer implies that $\beta_{(a,\alpha)}F\alpha = \beta_{(a,\alpha)}\epsilon_{Fa}$, so we get the natural transformation $\bar{\eta}: \text{id}_{T\text{-Alg}} \Rightarrow \bar{U}\bar{F}$. The only thing left to do is checking the triangular identities, which is left to the reader.

2. If U preserves coequalizers of U -split pairs, both $U\bar{F}$ and U^T are coequalizers of the above diagram, hence $\bar{\eta}$ is an isomorphism.
3. From the triangular identities, $\bar{U}\bar{\epsilon}\cdot\bar{\eta}\bar{U} = \text{id}_{\bar{U}}$, hence $\bar{U}\bar{\epsilon}$ is an isomorphism. Being $U^T\bar{U} = U$ conservative, $\bar{\epsilon}$ is an isomorphism as well. \square

Definition 2.1.7. A pair $f, g: c \rightrightarrows d$ in a category \mathcal{C} is *reflexive* if there exists a common section $i: d \rightarrow c$, that is $fi = gi = \text{id}_d$.

A coequalizer of a reflexive pair is a *reflexive coequalizer*.

Remark 2.1.8. To give a cone of a reflexive pair it is enough to give a map $h: d \rightarrow x$ such that $hf = hg$, hence $\text{colim}(c \rightrightarrows d) \cong \text{colim}(c \rightrightarrows d)$.

Proposition 2.1.9. In Beck's monadicity theorem it suffices for (1) that coequalizers of reflexive U -split pairs exist, while in (2) and (3) we only need for them to be preserved.

Proof. The pair

$$FUFUT \xrightleftharpoons[\epsilon_{FUT}]{F\alpha} FUT$$

has $F\eta U^T$ as common section. In fact, $\alpha \cdot \eta U^T = \text{id}_{UT}$ by the unit law of the T -action $\alpha: TU^T \Rightarrow U^T$ and $\epsilon F \cdot F\eta = \text{id}_F$ by the triangular identities. \square

Example 2.1.10. Let \mathcal{A} and \mathcal{B} be small categories, \mathcal{C} a category which is both complete and cocomplete, and $G: \mathcal{A} \rightarrow \mathcal{B}$ a functor. The restriction along G , G^* , has both adjoints, given by left and right Kan extensions. Notice that the induced monad is cocontinuous since G^* is a left adjoint. Moreover, G^* is conservative if G is essentially surjective, thus any essentially surjective functor G induces a monadic adjunction as follows:

$$\begin{array}{ccc}
 [\mathcal{B}, \mathcal{C}] & \xrightarrow{\bar{G}^*} & G^* \text{Lan}_G\text{-Alg} \\
 \downarrow G^* & \Downarrow \cong & \downarrow \\
 & & [\mathcal{A}, \mathcal{C}]
 \end{array}$$

We are going to show why reflexive coequalizers are useful, but first we need some preliminary definitions and results.

Definition 2.1.11. A functor $F: \mathcal{J} \rightarrow \mathcal{J}$ between small categories is called *final* if for any diagram $D: \mathcal{J} \rightarrow \mathcal{C}$ the comparison morphism $\text{colim}_{\mathcal{J}} DF \rightarrow \text{colim}_{\mathcal{J}} D$ is an isomorphism whenever both colimits exist.

Proposition 2.1.12. Let $F: \mathcal{J} \rightarrow \mathcal{J}$ be a functor between small categories. The following are equivalent:

- (i) F is final;
- (ii) the unique isomorphism

$$\begin{array}{ccc} \mathcal{J}^{\text{op}} & \xrightarrow{F^{\text{op}}} & \mathcal{J}^{\text{op}} \\ & \Downarrow \cong & \\ * & & * \\ & \searrow & \swarrow \\ & \mathbf{Set} & \end{array}$$

exhibits $*$ as $\text{Lan}_{F^{\text{op}}} *$;

- (iii) for each $j \in \mathcal{J}$, the category $(j \downarrow F)$ is connected.

Proof. (ii) \iff (iii) We have $\text{Lan}_{F^{\text{op}}} *(j) \cong \text{colim}_{(j \downarrow F)} *$ by the formula for Kan extensions. A colimit of $(j \downarrow F) \rightarrow \mathbf{Set}$, $(\phi, j') \mapsto *$ is terminal if and only if $(j \downarrow F)$ is connected, hence the thesis.

- (ii) \implies (i) Let $D: \mathcal{J} \rightarrow \mathcal{C}$ be a diagram. We can then write $\text{Cocone}(D, -)$ as follows:

$$\text{Cocone}(D, X) \cong \text{Nat}(*, \mathcal{C}(D-, X)) \cong [\mathcal{J}^{\text{op}}, \mathbf{Set}](*, \mathcal{C}(D-, X))$$

By definition of left Kan extension, we also have

$$[\mathcal{J}^{\text{op}}, \mathbf{Set}](*, \mathcal{C}(DF-, X)) \cong [\mathcal{J}^{\text{op}}, \mathbf{Set}](\text{Lan}_{F^{\text{op}}} *, \mathcal{C}(D-, X))$$

If $\text{Lan}_{F^{\text{op}}} * \cong *$, this shows that $\text{colim}_{\mathcal{J}} DF \cong \text{colim}_{\mathcal{J}} D$.

- (i) \implies (iii) Left as an exercise. □

Definition 2.1.13. A small category \mathcal{J} is *sifted* if the diagonal $\Delta: \mathcal{J} \rightarrow \mathcal{J} \times \mathcal{J}$ is final. A colimit is sifted if the domain category is.

Example 2.1.14. For any filtered category \mathcal{J} , the category $((i, i') \downarrow \Delta)$ is again filtered for any $(i, i') \in \mathcal{J} \times \mathcal{J}$ and hence connected, thus filtered colimits are sifted.

Example 2.1.15. Coequalizers are not sifted. Indeed, their indexing category $\mathcal{J} = \{1 \rightrightarrows 0\}$ is such that $((0, 1) \downarrow \Delta)$ is not connected. However, reflexive coequalizers are sifted. Checking it for yourself may be a tedious yet useful exercise.

Example 2.1.16. Coproducts and initial objects are not sifted, for their slice categories may be either empty or have several connected components.

Example 2.1.17. Pushouts are not sifted.

Proposition 2.1.18. If $F: \mathcal{A} \times \mathcal{B} \rightarrow \mathcal{C}$ is a functor preserving reflexive coequalizers in each variable, that is for any $a \in \mathcal{A}$, $b \in \mathcal{B}$ the functors $F(a, -): \mathcal{B} \rightarrow \mathcal{C}$ and $F(-, b): \mathcal{A} \rightarrow \mathcal{C}$ preserve reflexive coequalizers, then F preserves reflexive coequalizers as well.

Proof. We need to check that, given a reflexive coequalizer $a_0 \rightrightarrows a_1 \rightarrow a_2$ in \mathcal{A} and $b_0 \rightrightarrows b_1 \rightarrow b_2$ in \mathcal{B} , the diagonal of the following diagram is a coequalizer diagram in \mathcal{C} .

$$\begin{array}{ccccc}
 F(a_0, b_0) & \rightrightarrows & F(a_1, b_0) & \longrightarrow & F(a_2, b_0) \\
 \Downarrow & \nearrow & \Downarrow & & \Downarrow \\
 F(a_0, b_1) & \rightrightarrows & F(a_1, b_1) & \longrightarrow & F(a_2, b_1) \\
 \downarrow & & \downarrow & \searrow & \downarrow \\
 F(a_0, b_2) & \rightrightarrows & F(a_1, b_2) & \longrightarrow & F(a_2, b_2)
 \end{array}$$

From general facts, $F(a_2, b_2)$ is the colimit of the square in the top left. We may prove this using the sections, however in this case we can use the fact that, under certain hypothesis, given a diagram $D: \mathcal{J} \times \mathcal{J} \rightarrow \mathcal{C}$ we have $\text{colim}_{\mathcal{J} \times \mathcal{J}} D \cong \text{colim}_{\mathcal{J}} D\Delta$. Specifically, we want this to hold when $\mathcal{J} = \{a_0 \rightrightarrows a_1\}$. But the category $((i, i') \downarrow \Delta)$ is connected for any $i \in \mathcal{J}$, hence Δ is final and we have the thesis. \square

Example 2.1.19. The functor $\mathbf{Set} \times \mathbf{Set} \xrightarrow{- \times -} \mathbf{Set}$ satisfies the hypothesis of the theorem since \mathbf{Set} is cartesian closed, hence $X \mapsto X \times X$ preserves reflexive coequalizers by the proposition. This shows that $\mathbf{Set}(X, -) \cong \Pi_{x \in X} \mathbf{Set}(*, -)$ preserves reflexive coequalizers if X is finite, hence the functor $\mathbf{Set}(X, -): \mathbf{Set} \rightarrow \mathbf{Set}$ is monadic with $T = \mathbf{Set}(X, X \times -)$ for X finite.

$$\begin{array}{ccc}
 \mathbf{Set} & \xrightarrow{\cong} & T\text{-Alg} \\
 \searrow \mathbf{Set}(X, -) & & \swarrow U^T \\
 & \mathbf{Set} &
 \end{array}$$

We also have the following result.

Proposition 2.1.20. If $F: \mathcal{B} \times \mathcal{C} \rightarrow \mathcal{D}$ is a functor preserving sifted colimits in each variable, then it preserves them as a functor $\mathcal{B} \times \mathcal{C} \rightarrow \mathcal{D}$.

Remark 2.1.21. This proposition is false if we drop the siftedness condition, for under the functor $\mathbf{Set} \times \mathbf{Set} \xrightarrow{- \times -} \mathbf{Set}$ we have $(* + *) \times (* + *) \not\cong * + * \cong (* \times *) + (* \times *)$

2.2 Algebraic Theories and Finitary Monads

Many objects in algebra can be described as sets X with some finitary operations $X^n \xrightarrow{m_i} X$ subject to a list of axioms involving the m_i and their products.

Example 2.2.1. (Commutative) monoids and groups, rings and, fixed a ring R , R -modules, where for each $r \in R$ we specify an operation $X \xrightarrow{r} X$.

We can construct categories whose objects are sets paired with operations fulfilling the axioms and functions commuting with the operations as morphisms. These are called models for single-sorted finitary theories.

Proposition 2.2.2. The forgetful functor from a category \mathcal{C} of models of a single-sorted finitary theory to **Set** creates sifted colimits.

Proof. Notice that the n -fold product $\mathbf{Set} \xrightarrow{(-)^n} \mathbf{Set}$ factors as $\mathbf{Set} \xrightarrow{\Delta} \mathbf{Set}^n \xrightarrow{-\times\cdots\times-} \mathbf{Set}$, hence it preserves sifted colimits. This gives us unique candidates for operations on the colimit such that the cocone in **Set** is a morphism of models. These operations satisfy the axioms because the domain is in each case again of the form colim^n .

$$\begin{array}{ccccc} X^n & \rightrightarrows & Y^n & \longrightarrow & Z^n \\ \downarrow m_i & & \downarrow m_i & & \downarrow \exists! m_i \\ X & \rightrightarrows & Y & \longrightarrow & Z \end{array}$$

□

Corollary 2.2.3. Let \mathcal{C}, \mathcal{D} be categories of models of single sorted finitary theories. Any functor $U: \mathcal{C} \rightarrow \mathcal{D}$ which commutes with the forgetful functor to **Set** preserves reflexive coequalizers. If U has a left adjoint, then it is monadic.

Proof. The forgetful functors $\mathcal{C} \rightarrow \mathbf{Set}, \mathcal{D} \rightarrow \mathbf{Set}$ are both conservative, hence U is as well. The statement about reflexive coequalizers was just proved and the last claim follows from Beck's theorem. □

Example 2.2.4. The forgetful functor $\mathbf{CAlg}_R \rightarrow \mathbf{Mod}_R$ is monadic for any commutative ring R . The same goes for the forgetful functors from **Grp**, **Ab**, \mathbf{Mod}_R , **Rng**, **CRng**, **Ring** and **CRing** to **Set**.

Definition 2.2.5. A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is finitary if it commutes with filtered colimits. A monad is finitary if its underlying endofunctor T is finitary.

Remark 2.2.6. Asking for the underlying endofunctor T to be finitary is equivalent to asking for its forgetful functor U^T to be finitary.

Example 2.2.7. All the forgetful functors from single sorted finitary theories which have a left adjoint (fact: all of them do) come from finitary monads on **Set**.

In general, preserving filtered colimits is a weaker condition than preserving sifted colimits. However, for endofunctors on **Set** the two coincide.

In order to prove this, we need to know how such an endofunctor is determined by its action on finite sets and the idea is to check its behaviour on finite sets, since every set is a directed union of its finite subsets.

We write $[\mathcal{C}, \mathcal{D}]_{\text{fin}}$ for the full subcategory of $[\mathcal{C}, \mathcal{D}]$ whose objects are finitary functors, $\mathbf{Set}_{\text{fin}}$ for the category of finite sets.

Theorem 2.2.8. Restriction along the inclusion $K: \mathbf{Set}_{\text{fin}} \rightarrow \mathbf{Set}$ induces an equivalence $[\mathbf{Set}, \mathbf{Set}]_{\text{fin}} \rightarrow [\mathbf{Set}_{\text{fin}}, \mathbf{Set}]$ whose inverse send F to $\text{Lan}_K F$.

Before proving this theorem, we present some consequences.

Corollary 2.2.9. Any finitary functor $\mathbf{Set} \rightarrow \mathbf{Set}$ preserves sifted colimits. In particular, if $T: \mathbf{Set} \rightarrow \mathbf{Set}$ is the underlying endofunctor of a finitary monad, then $U^T: T\text{-Alg} \rightarrow \mathbf{Set}$ creates sifted colimits.

Proof. Recall that $\text{Lan}_K: [\mathbf{Set}_{\text{fin}}, \mathbf{Set}] \rightarrow [\mathbf{Set}, \mathbf{Set}]_{\text{fin}}$ is a left adjoint and on both sides colimits are computed pointwise. Moreover, functors preserving colimits of a given class are closed under formation of pointwise colimits, hence it is enough to check on a generating set.

If $\text{Lan}_K F_i$ preserves sifted colimits for all $F_i: \mathbf{Set}_{\text{fin}} \rightarrow \mathbf{Set}$, then the functor $\text{colim Lan}_K F_i \cong \text{Lan}_K \text{colim } F_i$ preserves sifted colimits.

Any functor $F: \mathbf{Set}_{\text{fin}} \rightarrow \mathbf{Set}$ is a colimit of representable functors $\mathbf{Set}_{\text{fin}}(X, -)$. Indeed, it is enough to consider $(\mathcal{J} \downarrow F) \rightarrow [\mathbf{Set}_{\text{fin}}, \mathbf{Set}]$, $(\mathbf{Set}_{\text{fin}}(X, -) \Rightarrow F) \mapsto \mathbf{Set}_{\text{fin}}(X, -)$ and notice that this being a colimit diagram essentially follows from Yoneda as $F \cong \text{colim}_{(\mathcal{J} \downarrow F)} \text{dom}$.

By the previous theorem, it is enough to check now that the functor $\text{Lan}_K \mathbf{Set}_{\text{fin}}(X, -)$ preserves sifted colimits.

Observe the following diagram:

$$\begin{array}{ccc} \mathbf{Set}_{\text{fin}} & \xrightarrow{K} & \mathbf{Set} \\ \mathbf{Set}_{\text{fin}}(X, -) \searrow & \xRightarrow{\quad} & \swarrow \text{Lan}_K \mathbf{Set}_{\text{fin}}(X, -) \\ & \mathbf{Set} & \end{array}$$

A natural transformation $\text{Lan}_K \mathbf{Set}_{\text{fin}}(X, -) \Rightarrow G$ by definition is equivalent to a natural transformation $\mathbf{Set}_{\text{fin}}(X, -) \Rightarrow GK$, which by Yoneda is equivalent to a map $*$ $\rightarrow GKX$, which again by Yoneda corresponds to a natural transformation $\mathbf{Set}(KX, -) \Rightarrow G$, hence $\text{Lan}_K \mathbf{Set}_{\text{fin}}(X, -) \cong \mathbf{Set}(KX, -)$.

We only have to check now that $\mathbf{Set}(KX, -) \cong \prod_{x \in X} \mathbf{Set}(*, -)$ preserves sifted colimits, but this is just the functor $Y \mapsto \prod_{x \in X} Y$, which as we know commutes with sifted colimits. \square

Proposition 2.2.10. A category \mathcal{C} with sifted colimits is cocomplete if and only if it has an initial object and binary coproducts.

Proof. One implication is obvious. For the other one, notice that we get finitary coproducts immediately. An infinite coproduct can be written as a filtered colimit of finite coproducts. We are only missing coequalizers. If $f, g: a \rightrightarrows b$ is a pair, then $f + \text{id}, g + \text{id}: a + b \rightrightarrows b$ in in_b (where in_b is the inclusion of b in the coproduct) is a reflexive pair hence the coequalizer $k = \text{coeq}(a + b \rightrightarrows b)$ exists and coequalizes the original pair. The universal $h: b \rightarrow k$ is also a coequalizer of $f, g: a \rightrightarrows b$. \square

Theorem 2.2.11. Let X be a set, $T: \prod_{x \in X} \mathbf{Set} \rightarrow \prod_{x \in X} \mathbf{Set}$ be a finitary monad. Then $T\text{-Alg}$ is cocomplete.

Proof. $T\text{-Alg}$ has always an initial object, namely the free algebra of the initial object $(T \prod_{x \in X} \emptyset, \mu_{\prod_{x \in X} \emptyset})$. Similarly, for $a, b \in \prod_{x \in X} \mathbf{Set}$, using that F^T is a left adjoint (hence it preserves colimits) we find that $(T(a \coprod b), \mu_{a \coprod b})$ is a coproduct of (Ta, μ_a) and (Tb, μ_b) . Therefore $T\text{-Alg}$ has coproducts of free algebras. We want to check that we have binary coproducts of T -algebras (a, α) and (b, β) . We have reflective coequalizers

$$\begin{array}{ccc} (T^2 a, \mu_{Ta}) & \begin{array}{c} \xrightarrow{T\alpha} \\ \xleftarrow{T\eta_a} \\ \xrightarrow{\mu_a} \end{array} & (Ta, \mu_a) \xrightarrow{\alpha} (a, \alpha) \\ (T^2 b, \mu_{Tb}) & \begin{array}{c} \xrightarrow{T\beta} \\ \xleftarrow{T\eta_b} \\ \xrightarrow{\mu_b} \end{array} & (Tb, \mu_b) \xrightarrow{\beta} (b, \beta) \end{array}$$

so we get a new reflective pair by taking coproducts of the free algebras

$$(T^2a, \mu_{Ta}) + (T^2b, \mu_{Tb}) \xrightleftharpoons[\mu_a + \mu_b]{T\alpha + T\beta} (Ta, \mu_a) + (Tb, \mu_b)$$

From the corollary, T preserves sifted colimits, hence $T\text{-Alg}$ has reflexive coequalizers. Then the diagram above has a coequalizer, which is a coproduct of (a, α) and (b, β) . \square

Remark 2.2.12. This shows that **Ab**, **Grp**, **Rng** etc. are cocomplete.

Remark 2.2.13. We only used the fact that T preserves sifted colimits, hence a monad on a cocomplete category \mathcal{C} preserving sifted colimits has a cocomplete category of algebras.

Remark 2.2.14. In fact, we only need that reflexive coequalizers and filtered colimits exist in $T\text{-Alg}$.

2.3 Dense generators

The aim of this section is to prove the theorem about finitary endofunctors of **Set**. We want to identify “nice” generating subcategories like $\mathbf{Set}_{\text{fin}} \rightarrow \mathbf{Set}$.

Definition 2.3.1. Let $K: \mathcal{A} \rightarrow \mathcal{C}$ be the inclusion of a full subcategory or, equivalently, a fully faithful functor. We define the *restricted Yoneda functor* $\tilde{K}: \mathcal{C} \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$ as the functor sending $c \in \mathcal{C}$ to $\mathcal{C}(K-, c)$.

The *canonical cocone* on the domain functor

$$\begin{aligned} \text{dom}: (K \downarrow c) &\rightarrow \mathcal{C} \\ (a, \varphi) &\mapsto Ka \end{aligned}$$

has components $\varphi: \text{dom}(a, \varphi) = Ka \xrightarrow{\varphi} c$.

Definition 2.3.2. A colimit of a diagram $D: I \rightarrow \mathcal{C}$ is *K-absolute* if it is preserved by $\tilde{K}: \mathcal{C} \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$.

Definition 2.3.3. A full subcategory/fully faithful functor $K: \mathcal{A} \rightarrow \mathcal{C}$ is called *dense* if \tilde{K} is fully faithful.

Theorem 2.3.4. Let \mathcal{C} be a cocomplete category, \mathcal{A} a small category and $K: \mathcal{A} \rightarrow \mathcal{C}$ a fully faithful functor. Then \tilde{K} has a left adjoint given by $\text{Lan}_{\mathcal{A}} K: [\mathcal{A}^{\text{op}}, \mathbf{Set}] \rightarrow \mathcal{C}$.

$$\begin{array}{ccc} \mathcal{A} & \xrightarrow{\mathcal{A}} & [\mathcal{A}^{\text{op}}, \mathbf{Set}] \\ & \searrow K & \swarrow \text{Lan}_{\mathcal{A}} K \\ & \mathcal{C} & \end{array}$$

Moreover, the following conditions are equivalent:

1. \tilde{K} is fully faithful;
2. for every $c \in \mathcal{C}$, the canonical cocone on $\text{dom}: (\mathcal{A} \downarrow c) \rightarrow \mathcal{C}$ exhibits c as colimit of the diagram $\text{dom}: (K \downarrow c) \rightarrow \mathcal{C}$;

3. every object is a K -absolute colimit of a diagram of the form $\mathcal{J} \xrightarrow{D} \mathcal{A} \xrightarrow{K} \mathcal{C}$;
4. there exists some family of diagrams $D_i: \mathcal{J}_i \rightarrow \mathcal{C}$ which have K -absolute colimits and \mathcal{C} is the closure³ of \mathcal{A} under the colimits of D_i ;
5. the counit of $\text{Lan}_{\mathcal{J}} K \dashv \tilde{K}$ is an isomorphism.

Proof. We have $\text{Lan}_{\mathcal{J}} K(F) = \text{colim}_{(\mathcal{J} \downarrow F)} Ka$, hence

$$\begin{aligned}
 \mathcal{C}(\text{Lan}_{\mathcal{J}} K(F), c) &\cong \lim_{(\mathcal{J} \downarrow F)} \mathcal{C}(Ka, c) \\
 &\stackrel{\text{Yoneda}}{\cong} \lim_{(\mathcal{J} \downarrow F)} [\mathcal{A}^{\text{op}}, \mathbf{Set}] (\mathcal{A}(-, a), \mathcal{C}(K-, c)) \\
 &\cong [\mathcal{A}^{\text{op}}, \mathbf{Set}] \left(\text{colim}_{(\mathcal{J} \downarrow F)} \mathcal{A}(-, a), \mathcal{C}(K-, c) \right) \\
 &\cong [\mathcal{A}^{\text{op}}, \mathbf{Set}] (F, \mathcal{C}(K-, c))
 \end{aligned}$$

hence $\text{Lan}_{\mathcal{J}} K \dashv \tilde{K}$. It remains to check the implications.

- (1) \implies (2) If \tilde{K} is fully faithful, it suffices to check that the image of the canonical cocone under \tilde{K} is a colimit. Since K is fully faithful, this image is precisely the diagram $\mathcal{A}(-, a) \cong \mathcal{C}(K-, Ka) \Rightarrow \mathcal{C}(K-, c)$. Then

$$\text{colim}_{(K \downarrow c)} \tilde{K} \text{ dom} = \text{colim}_{(\mathcal{J} \downarrow \mathcal{C}(K-, c))} \mathcal{A}(-, a) \cong \mathcal{C}(K-, c).$$

This proves 2. and the fact that the colimit of $(K \downarrow c) \rightarrow \mathcal{C}$ is preserved by $\tilde{K}: \mathcal{C} \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$.

- (2) \implies (3) As we just observed, colim dom is \tilde{K} -absolute and $\text{dom}: (K \downarrow c) \rightarrow \mathcal{C}$ factors through \mathcal{A} .

- (3) \implies (4) Clear.

- (4) \implies (5) Let \mathcal{B} be full subcategory spanned by the elements $b \in \mathcal{C}$ such that $\epsilon_b: \text{Lan}_{\mathcal{J}} K \tilde{K}(b) \rightarrow b$ is an isomorphism. It is closed under K -absolute colimits since they are preserved by \tilde{K} (by definition), by the left adjoint $\text{Lan}_{\mathcal{J}} K$ and by $\text{id}_{\mathcal{C}}$. It remains to check that the counit at $a \in \mathcal{A}$ is an isomorphism. But $\tilde{K}a = \mathcal{C}(K-, Ka) \cong \mathcal{A}(-, a)$, so $\text{Lan}_{\mathcal{J}} K(\mathcal{A}(-, a)) \cong a$, where the last isomorphism is a consequence of the fact that $(\mathcal{J} \downarrow \mathcal{A}(-, a))$ has a terminal object $\text{id}: \mathcal{A}(-, a) \Rightarrow \mathcal{A}(-, a)$.

- (5) \implies (1) Any right adjoint whose counit is an isomorphism is fully faithful. \square

Remark 2.3.5. Notice that the first four points of the Theorem 2.3.4 are equivalent even when \mathcal{C} is not cocomplete. In fact, we also have the implication (4) \implies (1) without assumptions on the cocompleteness of \mathcal{C} .

³That is, the smallest full subcategory \mathcal{B} of \mathcal{C} which contains \mathcal{A} and which satisfies the following

if some diagram $D_i: \mathcal{J}_i \rightarrow \mathcal{C}$ factors through \mathcal{B} , then $\text{colim } D_i \in \mathcal{B}$

is \mathcal{C} itself.

Proof. We give an alternative proof of the implication (4) \implies (1).

Consider the (full) subcategory \mathcal{B} of \mathcal{C} with objects

$$\{b \mid \tilde{K}_{Ka,b}: \mathcal{C}(Ka, b) \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}] (\tilde{K}Ka, \tilde{K}b) \text{ is bijective } \forall a \in \mathcal{A}\}$$

Since $\tilde{K}Ka = \mathcal{A}(-, a)$, by Yoneda, the target is given by $\text{ev}_a \circ \tilde{K}(b)$. Therefore it preserves K -absolute colimits. The domain is also equal to $\text{ev}_a \circ \tilde{K}(b)$, by definition. Hence it preserves K -absolute colimits and contains Ka' for every a' , therefore $\mathcal{B} = \mathcal{C}$. Consider now

$$\mathcal{B}' = \{b' \mid \mathcal{C}(b', b) \xrightarrow{\tilde{K}} [\mathcal{A}^{\text{op}}, \mathbf{Set}] (\tilde{K}b', \tilde{K}b) \text{ is bijective}\}$$

This is closed under all K -absolute colimits and contains Ka by the above argument, hence it is all of \mathcal{C} . \square

Definition 2.3.6. A fully faithful $K: \mathcal{A} \rightarrow \mathcal{C}$ is *dense* if \tilde{K} is fully faithful. A collection of diagrams $\{D_j: \mathcal{J}_j \rightarrow \mathcal{C}\}$ s.t. \mathcal{C} is the closure of K under colimits of D_j and $\text{colim } D_j$ is K -absolute is a *density presentation*.

Remark 2.3.7. The definition of density makes sense for arbitrary K , but the implication (4) \implies (1) does not work in general.

Example 2.3.8. 1. $\mathcal{J}: \mathcal{A} \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$ is dense. In fact, we have

$$\tilde{\mathcal{J}}(G) = [\mathcal{A}^{\text{op}}, \mathbf{Set}] (\mathcal{J} -, G) \xrightarrow{\text{Yoneda}} G,$$

thus $\tilde{\mathcal{J}} \cong \text{id}$ preserves all colimits. The colimit indexed by the slice $(\mathcal{J} \downarrow F)$ gives a density presentation.

2. $K: \mathbf{Set}_{\text{fin}} \hookrightarrow \mathbf{Set}$ is dense: S finite implies $\mathbf{Set}(S, -)$ preserves sifted, hence filtered, colimits. Filtered (and sifted) colimits are K -absolute, and a density presentation for K can be found once one writes an arbitrary set as union of its finite subset.
3. $K: \{*\} \subseteq \mathbf{Set}$ is dense: $\tilde{K}: \mathbf{Set} \rightarrow [*, \mathbf{Set}] \cong \mathbf{Set}$, $X \mapsto \mathbf{Set}(*, X) \cong X$, hence we find that \tilde{K} preserves all colimits. We can use coproducts for density presentation.
4. $\{k\} \subset \mathbf{Vect}_k$ is not dense even though every vector space is a coproduct of copies of k .

Definition 2.3.9. Given *any* functor $F: \mathcal{A} \rightarrow \mathcal{D}$, we can talk about the *restricted Yoneda embedding* $\tilde{F}: \mathcal{D} \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$, sending $d \in \mathcal{D}$ to $\mathcal{D}(F-, d)$.

Proposition 2.3.10. Let $K: \mathcal{A} \rightarrow \mathcal{C}$ be fully faithful and $F: \mathcal{A} \rightarrow \mathcal{D}$ any functor. Suppose there exists $L: \mathcal{C} \rightarrow \mathcal{D}$ and bijections $\mathcal{D}(Lc, d) \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}] (\mathcal{C}(K-, c), \mathcal{D}(F-, d))$ natural both in c and d . Then there is an isomorphism $\eta: F \xrightarrow{\sim} L \circ K$ exhibiting L as left Kan extension of F along K .

Proof. The presheaf $\mathcal{C}(K-, c)$ is the colimit of the canonical cone on $(\mathcal{J} \downarrow \mathcal{C}(K-, c))$. By Yoneda, this is equivalent to $(K \downarrow c)$ with objects $(a \in \mathcal{A}, \varphi: Ka \rightarrow c)$ and the evident morphisms. If $c \cong Ka$, this has (a, id_a) as terminal object. Therefore the colimit is $\mathcal{C}(K-, Ka) \cong \mathcal{A}(-, a)$. Moreover, the formula above in this case gives

$$\mathcal{D}(LKa, d) \cong [\mathcal{A}^{\text{op}}, \mathbf{Set}] (\mathcal{A}(-, c), \mathcal{D}(F-, d)) \cong \mathcal{D}(Fa, d).$$

This shows that $L \circ K \cong F$ naturally. Then $Lc = \text{colim}_{(K \downarrow c)} Fa$, by the classical formula for left Kan extensions. \square

Definition 2.3.11. Any such Kan extension is called pointwise.

Theorem 2.3.12. Let $K: \mathcal{A} \rightarrow \mathcal{C}$ be a fully faithful dense functor with density presentation $\{D_j: \mathcal{J}_j \rightarrow \mathcal{C}\}_{j \in \mathcal{J}}$. Let $F: \mathcal{A} \rightarrow \mathcal{D}$ be a functor and assume that \mathcal{D} has colimits of shape $\mathcal{J}_j \ \forall j \in \mathcal{J}$. Then the pointwise left Kan extension of F along K exists. In particular, the unit $F \Rightarrow \text{Lan}_K F \circ K$ is an isomorphism.

Proof. We want a L as in the above proposition $\mathcal{D}(Lc, d) \cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](\mathcal{C}(K-, c), \mathcal{D}(F-, d))$ natural in c, d . This simply says that $\forall c \in \mathcal{C}$ the functor $[\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K}c, \tilde{F}-): \mathcal{D} \rightarrow \mathbf{Set}$ is representable. Take $\mathcal{B} := \{b \in \mathcal{C} \mid [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K}b, \tilde{F}-) \text{ is representable}\}$. If $b = Ka$, then $\tilde{K}b = \mathcal{C}(K-, Ka) \cong \mathcal{A}(-, a)$. Therefore

$$\begin{aligned} [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K}Ka, \tilde{F}d) &\cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](\mathcal{A}(-, a), \tilde{F}d) \\ &\stackrel{\text{Yoneda}}{\cong} \tilde{F}d(a) \\ &\cong \mathcal{D}(Fa, d) \end{aligned}$$

so it is represented by Fa and $Ka \in \mathcal{B}$. Furthermore \mathcal{B} is closed under \tilde{K} -absolute colimits of shape \mathcal{J}_j . Let $D: \mathcal{J}_j \rightarrow \mathcal{B}$ be a diagram such that $\text{colim}_{\mathcal{J}_j} D$ exists in \mathcal{C} and is preserved by \tilde{K} . We claim that $\text{colim}_{\mathcal{J}_j} D \in \mathcal{B}$.

$$\begin{aligned} [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K} \text{colim}_{\mathcal{J}_j} D, \tilde{F}d) &\cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](\text{colim}_{\mathcal{J}_j} \tilde{K}D, \tilde{F}d) \\ &\cong \lim_{\mathcal{J}_j} [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K}D, \tilde{F}d) \\ &\cong \lim_{\mathcal{J}_j} \mathcal{D}(LD, d) \\ &\cong \mathcal{D}(\text{colim}_{\mathcal{J}_j} LD, d) \end{aligned}$$

Hence $\mathcal{B} = \mathcal{C}$ and we get the functor $L \cong \text{Lan}_K F$. \square

Lemma 2.3.13. On the same conditions as before, pointwise Kan extensions along K preserve K -absolute colimits.

Proof. By definition, $\mathcal{D}(\text{Lan}_K F(c), d) \cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K}c, \tilde{F}d)$. Let $\text{colim}_{i \in \mathcal{J}_j} D_i$ be K -absolute. Then the claim is shown by the following chain of isomorphisms.

$$\begin{aligned} \mathcal{D}(\text{Lan}_K F(\text{colim } D_i), d) &\cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K} \text{colim } D_i, \tilde{F}d) \\ &\cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](\text{colim } \tilde{K}D_i, \tilde{F}d) \\ &\cong \lim [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{K}D_i, \tilde{F}d) \\ &\cong \lim \mathcal{D}(\text{Lan}_K F(D_i), d) \\ &\cong \mathcal{D}(\text{colim } \text{Lan}_K F(D_i), d). \end{aligned} \quad \square$$

Theorem 2.3.14. Let $K: \mathcal{A} \rightarrow \mathcal{C}$ be fully faithful, φ a class of colimit shapes and assume there exists a density presentation for K with colimits of shape $\mathcal{J}_j \in \varphi$. Let \mathcal{D} be a category with colimits of shape φ . We write $\varphi\text{-}\mathbf{Cocts}(\mathcal{C}, \mathcal{D})$ for the category of functors $\mathcal{C} \rightarrow \mathcal{D}$ which preserve φ -colimits. If all colimits of shape φ are K -absolute, then

$$[\mathcal{A}, \mathcal{D}] \xrightleftharpoons[\text{K}^*]{\text{Lan}_K} \varphi\text{-}\mathbf{Cocts}(\mathcal{C}, \mathcal{D})$$

is an equivalence.

Proof. The existence of Lan_K is guaranteed by the fact that \mathcal{D} has φ -colimits and K has density presentation with colimits of shape $\mathbb{J}_j \in \varphi$. By definition, Lan_K is left adjoint to $K^*: [\mathcal{C}, \mathcal{D}] \rightarrow [\mathcal{A}, \mathcal{D}]$ and by the lemma it lands in $\varphi\text{-Cocts}(\mathcal{C}, \mathcal{D})$. Then Lan_K is a left adjoint to the restriction. The unit is an isomorphism since $\text{Lan}_K F$ is pointwise, so it suffice to check that K^* is conservative. Let $G, H: \mathcal{C} \rightarrow \mathcal{D}$ be φ -cocts, $\alpha: G \Rightarrow H$ natural transformation such that αK is an isomorphism. Then, $\{c \mid \alpha_c \text{ is an isomorphism}\}$ contains Ka for every a and is closed under colimits of density presentations. Hence α is an isomorphism. \square

Corollary 2.3.15. For $K: \mathbf{Set}_{\text{fin}} \hookrightarrow \mathbf{Set}$, we get that

$$[\mathbf{Set}_{\text{fin}}, \mathbf{Set}] \begin{array}{c} \xrightarrow{\text{Lan}_K} \\ \perp \\ \xleftarrow{K^*} \end{array} [\mathbf{Set}, \mathbf{Set}]_{\text{fin}}$$

is an equivalence.

Definition 2.3.16. Let Φ be a class of colimit shapes and \mathcal{A} be a small category. We write $\Phi(\mathcal{A})$ for the closure of the representable presheaves in $[\mathcal{A}^{\text{op}}, \mathbf{Set}]$ under Φ -colimits. We have a functor $\mathfrak{J}: \mathcal{A} \rightarrow \Phi(\mathcal{A})$.

Remark 2.3.17. By construction, there exists a density presentation for $\mathfrak{J}: \mathcal{A} \rightarrow \Phi(\mathcal{A})$ consisting of Φ -colimits. This follows from $\tilde{\mathfrak{J}}: \Phi(\mathcal{A}) \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$ being simply the inclusion.

Theorem 2.3.18. Let Φ be a class of small colimit shapes, \mathcal{A} a small category. Then $\mathcal{A} \xrightarrow{\mathfrak{J}} \Phi(\mathcal{A})$ is the *free cocompletion* of \mathcal{A} under Φ -colimits, that is

$$[\mathcal{A}, \mathcal{C}] \begin{array}{c} \xrightarrow{\text{Lan}_{\mathfrak{J}}} \\ \perp \\ \xleftarrow{\mathfrak{J}^*} \end{array} \Phi\text{-Cocts}(\Phi(\mathcal{A}), \mathcal{C})$$

is an equivalence for every Φ -cocomplete \mathcal{C} . In particular, if Φ is the class of all small colimit shapes, then $\Phi(\mathcal{A}) = [\mathcal{A}^{\text{op}}, \mathbf{Set}]$ is the free cocompletion of \mathcal{A} . In this case, every $L \cong \text{Lan}_{\mathfrak{J}} F \in \mathbf{Cocts}([\mathcal{A}^{\text{op}}, \mathbf{Set}], \mathcal{C})$ has a right adjoint: $\text{Lan}_{\mathfrak{J}} F \dashv \tilde{F}$.

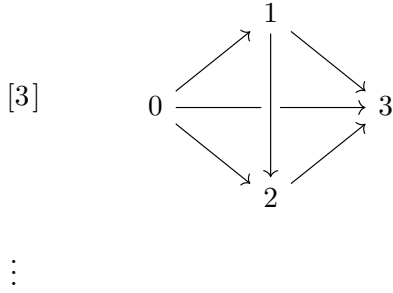
Proof. Since $\mathfrak{J}: \mathcal{A} \rightarrow \Phi(\mathcal{A})$ has density presentation consisting of Φ -colimits, the two equivalences follow from the previous theorem. To see that L has right adjoint, it suffices to check $L \cong \text{Lan}_{\mathfrak{J}} F \dashv \tilde{F}$:

$$\mathcal{C}(\text{Lan}_{\mathfrak{J}} F(G), c') \cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](\tilde{\mathfrak{J}}(G), \tilde{F}c') \cong [\mathcal{A}^{\text{op}}, \mathbf{Set}](G, \tilde{F}c')$$

since $\text{Lan}_{\mathfrak{J}} \Delta_{\bullet}: F$ is pointwise. \square

Example 2.3.19. Let Δ be the category of finite non-empty ordinals $[0], [1], \dots$ and order preserving maps. We have a functor $\Delta_{\bullet}: \Delta \rightarrow \mathbf{Top}$, sending $[n]$ to the standard (geometric) n -simplex Δ_n so we get an adjunction $\text{Lan}_{\mathfrak{J}} \Delta_{\bullet}: [\mathcal{A}^{\text{op}}, \mathbf{Set}] \rightleftarrows \mathbf{Top}: \Delta_{\bullet}^*$. $\text{Lan}_{\mathfrak{J}} \Delta_{\bullet}$ is called the *geometric realization* and $\Delta_{\bullet}^* =: \text{Sing}(-)$ is called the *singular complex*. In pictures:

$$[2] \quad \begin{array}{ccc} & 1 & \\ \nearrow & & \searrow \\ 0 & \xrightarrow{\quad} & 2 \end{array}$$



$[\Delta^{\text{op}}, \mathbf{Set}]$ is called the category of *simplicial sets* and it is denoted by \mathbf{sSet} . $\text{Lan}_{\mathbf{J}} \Delta_{\bullet}$ is denoted by $|\cdot|$.

2.4 Locally presentable categories

Definition 2.4.1. Let κ be a regular cardinal⁴. A category \mathcal{C} is κ -*filtered* if any diagram in \mathcal{C} of size $< \kappa$ has a cocone. Equivalently, if it is non-empty and for any set of objects $\{x_i\}$ of cardinality $< \kappa$ there exists $x \in \mathcal{C}$ and $x_i \rightarrow x$ such that

$$\begin{array}{ccc}
& \xrightarrow{\quad} & \\
x_i & \xrightarrow{\quad} & x_j \longrightarrow x \\
& \xrightarrow{\quad} & \\
& < \kappa &
\end{array}$$

is coequalizing. If $\kappa = \aleph_0$ a κ -filtered category is just a filtered one.

Definition 2.4.2. An object $c \in \mathcal{C}$ is called κ -*presentable* if $\mathcal{C}(c, -)$ preserves κ -filtered colimits. If $\kappa = \aleph_0$, c is called *finitely presentable*.

Definition 2.4.3. A category \mathcal{A} is called κ -*accessible* if there exists a small subcategory \mathcal{A}_0 of κ -presentable objects such that \mathcal{A} is the closure of \mathcal{A}_0 under κ -filtered colimits. The category \mathcal{A} is called *locally κ -presentable* if it is κ -accessible and cocomplete.

Example 2.4.4. Take $\mathcal{A} = \mathbf{Set}$, $\mathcal{A}_0 = \mathbf{Set}_{\text{fin}}$ and $S \in \mathbf{Set}_{\text{fin}}$. Then

$$\mathbf{Set}(S, -) = \prod_{|S| \text{ finite}} \mathbf{Set}(*, -)$$

preserves sifted, hence filtered, colimits. Thus \mathbf{Set} is locally finitely presentable.

Remark 2.4.5. For \mathcal{A}_0 as in the definition, $K: \mathcal{A}_0 \rightarrow \mathcal{A}$ the inclusion, we find that κ -filtered colimits are K -absolute.

$$\begin{array}{ccc}
\mathcal{A} & \xrightarrow{\tilde{K}} & [\mathcal{A}_0^{\text{op}}, \mathbf{Set}] \xrightarrow{\text{ev}_a} \mathbf{Set} \\
a' & \mapsto & \mathcal{A}(K-, a') \mapsto \mathcal{A}(a, a')
\end{array}$$

so that $\mathcal{A}_0 \xrightarrow{K} \mathcal{A}$ has density presentation consisting of κ -filtered colimits. Thus $\tilde{K}: \mathcal{A} \rightarrow [\mathcal{A}_0^{\text{op}}, \mathbf{Set}]$ is fully faithful and preserves κ -filtered colimits.

Definition 2.4.6. A functor is κ -*accessible* if it commutes with κ -filtered colimits. We write $[\mathcal{A}, \mathcal{B}]_{\kappa}$ for the subcategory of κ -accessible functors ($[\mathcal{A}, \mathcal{B}]_{\text{fin}}$ if $\kappa = \aleph_0$).

⁴Namely, a union of $< \kappa$ sets of cardinality $< \kappa$ has cardinality $< \kappa$.

It follows that

$$[\mathcal{A}_0, \mathcal{B}] \begin{array}{c} \xrightarrow{\text{Lan}_K} \\ \perp \\ \xleftarrow{K^*} \end{array} [\mathcal{A}, \mathcal{B}]_\kappa$$

if $\mathcal{A}, \mathcal{A}_0$ are as above and \mathcal{B} has κ -filtered colimits.

Remark 2.4.7. A κ -accessible category \mathcal{A} is locally κ -presentable if and only if $\mathcal{A} \xrightarrow{\tilde{K}} [\mathcal{A}_0^{\text{op}}, \mathbf{Set}]$ has a left adjoint for any choice of $\mathcal{A}_0 \subseteq \mathcal{A}$ that defines it as the closure under κ -filtered colimits.

Definition 2.4.8. A monad (T, μ, η) is said to have rank κ if T is a κ -accessible endofunctor.

Definition 2.4.9. A category \mathcal{A} is called *accessible* (resp. *locally presentable*) if it is κ -accessible (resp. locally κ -presentable) for some regular cardinal κ . A functor is accessible if it is κ -accessible for some κ . A monad has rank if it is accessible.

Our next goal is to prove that if (T, μ, η) is an accessible monad on a locally presentable category, then $T\text{-Alg}$ is locally presentable.

Proposition 2.4.10. Let \mathcal{A} be κ -accessible, $\mathcal{A}_0 \subseteq \mathcal{A}$ be the small subcategory of κ -presentable objects such that \mathcal{A} is the closure of \mathcal{A}_0 under κ -filtered colimits, and (T, μ, η) be a monad of rank κ on \mathcal{A} . Then $\mathcal{B} = \{(Ta_0, \mu_{a_0}) \mid a_0 \in \mathcal{A}_0\}$ is a dense generator of $T\text{-Alg}$.

Proof. First note that (Ta_0, μ_{a_0}) is κ -presentable. Indeed, we have $T\text{-Alg}((Ta_0, \mu_{a_0}), -) \cong \mathcal{A}(a_0, U^T -)$ and U^T creates all colimits that T preserves, in particular κ -filtered ones. Hence, LHS preserves κ -filtered colimits and for this reason (Ta_0, μ_{a_0}) is κ -presentable. Let Φ_1 be the class of κ -filtered diagrams. Writing $K: \mathcal{B} \rightarrow T\text{-Alg}$ for the inclusion, we have just shown that Φ_1 -colimits are K -absolute. Note that the closure of \mathcal{B} under Φ_1 contains all free algebras (Ta, μ_a) since \mathcal{A} is the closure of \mathcal{A}_0 under Φ_1 -colimits. Let Φ_2 be the class of diagrams $(T^2a, \mu_{Ta}) \rightrightarrows (Ta, \mu_a)$ for all $(a, \alpha) \in T\text{-Alg}$. The closure under $\Phi_1 \cup \Phi_2$ is clearly all of $T\text{-Alg}$, so we just need to show that Φ_2 -colimits are K -absolute, that is preserved by each $T\text{-Alg}((Ta_0, \mu_{a_0}), -)$. Since $T\text{-Alg}((Ta_0, \mu_{a_0}), -) \cong \mathcal{A}(a_0, U^T -)$ and U^T sends a coequalizer in question to a split coequalizer, the colimit is indeed K -absolute. \square

Example 2.4.11. *Finite* free groups, abelian groups, commutative rings etc. form dense generators of **Grp**, **Ab**, **CRng** etc.

Recall that the category of T -algebras of a finitary monad $T: \mathbf{Set} \rightarrow \mathbf{Set}$ is cocomplete. We would like to know that $T\text{-Alg}$ is locally finitely presentable. This result can be proved using the following fact (which is in turn an easy consequence of the result about the commutativity of κ -filtered colimits with κ -small limits in \mathbf{Set}):

κ -presentable objects are closed under κ -small colimits.

Proposition 2.4.12. Let \mathcal{C} be a cocomplete category which has a *small* dense subcategory consisting of κ -presentable objects. Then \mathcal{C} is locally κ -presentable.

Proof. Let \mathcal{A}' be the closure of \mathcal{A} under κ -small colimits. This is constructed as follows: $\mathcal{A}_0 = \mathcal{A}$. For any ordinal i we set

$$\mathcal{A}_{i+1} = \{\text{colimits of } \kappa\text{-small diagrams in } \mathcal{A}_i\}$$

and for a limit-ordinal λ we set $\mathcal{A}_\lambda = \bigcup_{\mu < \lambda} \mathcal{A}_\mu$. This terminates when $\lambda = \kappa$, so \mathcal{A}_κ is the colimit closure and thus small. From the above mentioned fact we know that \mathcal{A}' consists of κ -presentable objects. Since it contains \mathcal{A} , the inclusion $K: \mathcal{A}' \rightarrow \mathcal{C}$ is dense. If \mathcal{A}' is dense, then each object in \mathcal{C} is a colimit of $(\mathcal{A}' \downarrow c)$ which is a κ -filtered diagram by construction. \square

Corollary 2.4.13. For each finitary monad T on **Set**, the category $T\text{-Alg}$ is locally finitely presentable. Moreover if T is a monad of rank κ on a locally κ -presentable category, then $T\text{-Alg}$ is locally κ -presentable if and only if it is cocomplete.

Theorem 2.4.14. Let \mathcal{J} be a filtered category and

$$X: \mathcal{J} \rightarrow \mathbf{Set}, \quad i \mapsto X_i$$

a diagram and $(X_i \xrightarrow{n_i} X)_i$ a cocone. Then $(X_i \rightarrow X)_i$ is a colimit cocone if and only if

- i) For all $x \in X$ there exists an $i \in \mathcal{J}$ and an $\tilde{x} \in X_i$ such that $x = n_i(\tilde{x})$.
- ii) If $x, y \in X_i$ satisfy $n_i(x) = n_i(y)$, then there is some $\phi: i \rightarrow j$ such that $X_\phi(x) = X_\phi(y)$.
(Informally: “all equalities that hold in X hold in some X_j .”)

Proof. Given any other cone $\lambda_i: X_i \rightarrow Y$ we define $f: X \rightarrow Y$ by sending $x \rightarrow \lambda_i(\tilde{x})$ for any \tilde{x} in i). This is well defined by ii) and filteredness. It only remains to show that there exists such a cocone. Take $X = (\coprod X_i) / \sim$ with $(x, i) \sim (y, j)$ if there is some diagram $i \xrightarrow{\phi} k \xleftarrow{\psi} j$ in \mathcal{J} with $X_\phi(x) = X_\psi(y)$. \square

Corollary 2.4.15. In **Set** filtered colimits commute with finite limits and κ -filtered colimits commute with κ -small limits.

Proof. Check that a levelwise equalizer of cones satisfying i) and ii) above still satisfies i) and ii). This can be done by chasing through the following diagram

$$\begin{array}{ccccc} X_i & \hookrightarrow & Y_i & \rightrightarrows & Z_i \\ \downarrow & & \downarrow & & \downarrow \\ X & \hookrightarrow & Y & \rightrightarrows & Z \end{array}$$

For $< \kappa$ -fold products, we use κ -filteredness to extend ii) to any set of pairs of size $< \kappa$. Then check that the product of cones satisfying i) and ii) of size $< \kappa$ still satisfies i) and ii). \square

Corollary 2.4.16. In any category, κ -presentable objects are closed under κ -small colimits.

Proof. Let $C: \mathcal{J} \rightarrow \mathcal{C}$ be a filtered diagram and $D: \mathcal{I} \rightarrow \mathcal{C}$ a κ -small diagram of κ -presentable objects D_J .

$$\begin{aligned} \mathcal{C}(\operatorname{colim}_{\mathcal{J}} D_j, \operatorname{colim}_{\mathcal{I}} C_i) &\cong \lim_{\mathcal{J}} \mathcal{C}(D_j, \operatorname{colim}_{\mathcal{I}} C_i) \\ &\cong \lim_{\mathcal{J}} (\operatorname{colim}_{\mathcal{I}} \mathcal{C}(D_j, C_i)) \\ \text{Explicit check in Set} &\cong \operatorname{colim}_{\mathcal{J}} (\lim_{\mathcal{I}} \mathcal{C}(D_j, C_i)) \\ &\cong \operatorname{colim}_{\mathcal{J}} \mathcal{C}(\operatorname{colim}_{\mathcal{I}} D_j, C_i). \end{aligned} \quad \square$$

Proposition 2.4.17. Each object in a locally presentable category is λ -presentable for $\lambda \gg 0$.

Proof. Let \mathcal{C} be locally κ -presentable, choose a small dense subcategory \mathcal{A} of κ -presentable objects. So, any object $c \in \mathcal{C}$ we have is a colimit of $\operatorname{dom}: (\mathcal{A} \downarrow c) \rightarrow \mathcal{C}$. Choose λ such that $\lambda > \kappa$ and $\lambda > |\operatorname{Arr}(\mathcal{A}/c)|$. \square

The characterization of filtered colimits in **Set** gives the following characterization of finitely presentable objects: a is finitely presentable if for all filtered colimits $k_i: c_i \rightarrow c$ in \mathcal{C} and all $f: a \rightarrow c$ there exists a factorization

$$\begin{array}{ccc} & c_i & \\ f' \nearrow & & \searrow k_i \\ a & \xrightarrow{f} & c \end{array}$$

and any two such lifts f', f'' satisfying $k_i \circ f' = k_i \circ f''$ become equal after composing with some $c_\phi: c_i \rightarrow c_j$.

Corollary 2.4.18. Let \mathcal{C} be a locally κ -presentable category. We have that κ -filtered colimits commute with κ -small limits in \mathcal{C} .

Proof. Choose a small dense subcategory $\mathcal{A} \subset \mathcal{C}$ of κ -presentable objects. The inclusion $K: \mathcal{A} \rightarrow \mathcal{C}$ induces a fully faithful functor $\tilde{K}: \mathcal{C} \rightarrow [\mathcal{A}^{\text{op}}, \mathbf{Set}]$ with left adjoint $\text{Lan}_{\tilde{K}} K$, hence it preserves all limits. This implies that \mathcal{C} is complete as a reflective subcategory of the complete category $[\mathcal{A}^{\text{op}}, \mathbf{Set}]$. Moreover, \tilde{K} preserves κ -filtered colimits, given that $\text{ev}_a \circ \tilde{K} = \mathcal{C}(Ka, -)$ hence this reduces the problem to limits and κ -filtered colimits in $[\mathcal{A}^{\text{op}}, \mathbf{Set}]$, where both are computed levelwise. \square

Proposition 2.4.19. Let \mathcal{C}, \mathcal{D} be locally κ -presentable, $\lambda \geq \kappa$ a regular cardinal. Then the category $[\mathcal{C}, \mathcal{D}]_\lambda$ of λ -accessible functors and natural transformations is locally small, cocomplete and the inclusion $[\mathcal{C}, \mathcal{D}]_\lambda \rightarrow [\mathcal{C}, \mathcal{D}]$ preserves colimits. In fact, $[\mathcal{C}, \mathcal{D}]_\lambda$ is locally presentable.

Proof. The category \mathcal{C}_λ of λ -presentable objects in \mathcal{C} is essentially small and each $(\mathcal{C}_\lambda \downarrow c)$ is λ -filtered, so $\mathcal{C}_\lambda \hookrightarrow \mathcal{C}$ is dense with density presentation consisting of λ -filtered colimits. From a general fact, the left adjoint of

$$[\mathcal{C}_\lambda, \mathcal{D}] \xrightleftharpoons[\text{K}^*]{\text{Lan}_K} [\mathcal{C}, \mathcal{D}]$$

induces an equivalence onto its essential image, which is precisely $[\mathcal{C}, \mathcal{D}]_\lambda$. In other words, $[\mathcal{C}, \mathcal{D}]_\lambda \cong [\mathcal{C}_\lambda, \mathcal{D}]$ is locally small and the inclusion preserves all colimits. Furthermore, $[\mathcal{C}, \mathcal{D}]_\lambda$ is locally presentable since $[\mathcal{C}_\lambda, \mathcal{D}]$ is locally κ -presentable. \square

Corollary 2.4.20. The category of accessible functors $[\mathcal{C}, \mathcal{D}]_{\text{acc}}$ is closed under small colimits in $[\mathcal{C}, \mathcal{D}]$.

Proof. This is clear, since $[\mathcal{C}, \mathcal{D}]_{\text{acc}} = \bigcup_\lambda [\mathcal{C}, \mathcal{D}]_\lambda$. \square

The following theorem about dense functor has already been secretly used previously. Let us prove it once and for all.

Theorem 2.4.21. Consider two small categories \mathcal{A} and \mathcal{A}' and two fully faithful functors $\mathcal{A} \xrightarrow{P} \mathcal{A}' \xrightarrow{J} \mathcal{C}$. If the composite $K = JP$ is dense, then both P and J are dense.

Proof. It is immediate for P , since $\tilde{P} = \tilde{K}|_{\mathcal{A}'}$. Let us show J is dense. Note that we have

$$\begin{aligned} \mathcal{C}(Jd, c) &\xrightarrow[\sim]{\tilde{JP}} [\mathcal{A}^{\text{op}}, \mathbf{Set}](\mathcal{C}(JP-, Jd), \mathcal{C}(JP-, c)) \\ &\xrightarrow[\sim]{(J_{P-, d})^*} [\mathcal{A}^{\text{op}}, \mathbf{Set}](\mathcal{A}'(P-, d), \mathcal{C}(JP-, c)) \end{aligned}$$

where the first isomorphism holds because JP is dense and the second one because J is fully faithful. Consequently J is the pointwise left Kan extension of JP along P :

$$\begin{array}{ccc} \mathcal{A} & \xrightarrow{P} & \mathcal{A}' \\ JP \searrow & \xRightarrow{\sim} & \swarrow J \\ & \mathcal{C} & \end{array}$$

Since the left Kan extension is pointwise, it is preserved by any cocontinuous functor out of \mathcal{C} . In particular, for every $c \in \mathcal{C}$ we can apply $\mathcal{C}(-, c): \mathcal{C} \rightarrow \mathbf{Set}^{\text{op}}$ and we get that

$$\begin{array}{ccc} \mathcal{A}^{\text{op}} & \xrightarrow{P^{\text{op}}} & (\mathcal{A}')^{\text{op}} \\ \mathcal{C}(JP-, c) \searrow & \xleftarrow{\sim} & \swarrow \mathcal{C}(J-, c) \\ & \mathbf{Set} & \end{array}$$

is a right Kan extension for every $c \in \mathcal{C}$. In particular, each $\alpha: \mathcal{C}(J-, c) \Rightarrow \mathcal{C}(J-, c')$ is uniquely given by $\alpha P^{\text{op}}: \mathcal{C}(JP-, c) \Rightarrow \mathcal{C}(JP-, c')$. By density of $JP = K$, αP^{op} must be of the form g_* for a unique $g: c \rightarrow c'$. By uniqueness, $\alpha = \mathcal{C}(J-, g)$, hence \tilde{J} is full. Moreover, \tilde{JP} is equal to the composition $\mathcal{C} \xrightarrow{\tilde{J}} [(\mathcal{A}')^{\text{op}}, \mathbf{Set}] \xrightarrow{(P^{\text{op}})^*} [\mathcal{A}^{\text{op}}, \mathbf{Set}]$ and then \tilde{J} is also faithful. \square

2.5 Cocompleteness of categories of algebras

The goal of this section is to show that, if T is a monad on a locally presentable category \mathcal{C} and T has rank (it is accessible), then $T\text{-Alg}$ is cocomplete and thus locally presentable (the last bit is a consequence of a previous result).

There exists a single construction which admits the following as special cases:

- free monad on an endofunctor,
- free monad on a pointed endofunctor,
- free monoid on an object in a monoidal category,
- orthogonal factorization system generated by a set of morphisms,
- reflectiveness of a small orthogonality class,
- cocompletion of $T\text{-Alg}$ for suitable monads T ,
- existence of colimits of diagrams of accessible monads.

This was observed by G. M. Kelly in *A unified treatment of transfinite constructions for free algebras, free monoids, colimits, associated sheaves, and so on* (1980), which is “hard to read” but simplifies greatly in the context of locally presentable categories.

Throughout this section we will work with locally presentable categories and accessible functors.

Kelly’s main observation is that all obvious constructions can be reduced to the case of algebras for a well pointed endofunctor.

Definition 2.5.1. Let $S: \mathcal{C} \rightarrow \mathcal{C}$ be a functor. We call S *pointed* if there exists $\sigma: \text{id}_{\mathcal{C}} \Rightarrow S$. The pair (S, σ) is *well pointed* if $S\sigma = \sigma S: S \Rightarrow S^2$.

Definition 2.5.2. Given a pointed endofunctor (S, σ) , a (S, σ) -algebra is a pair (a, α) , where $\alpha: Sa \rightarrow a$ is a morphism in \mathcal{C} s.t. $\alpha\sigma_a = \text{id}_a$ (basically a monad without multiplication gives an example). A morphism of algebras $(a, \alpha) \rightarrow (b, \beta)$ is a morphism $f: a \rightarrow b$ in \mathcal{C} such that

$$\begin{array}{ccc} Sa & \xrightarrow{Sf} & Sb \\ \alpha \downarrow & & \downarrow \beta \\ a & \xrightarrow{f} & b \end{array}$$

commutes. We write $(S, \sigma)\text{-Alg}$ for the resulting category and $U^S: (S, \sigma)\text{-Alg} \rightarrow \mathcal{C}$ for the forgetful functor.

Lemma 2.5.3. If (S, σ) is a well pointed endofunctor, then there exists at most one algebra structure for any object and it exists if and only if σ_a is invertible, in which case $\alpha = \sigma_a^{-1}$. Moreover, $U^S: (S, \sigma)\text{-Alg} \rightarrow \mathcal{C}$ is fully faithful⁵. In other words, $(S, \sigma)\text{-Alg}$ is isomorphic to the full subcategory of \mathcal{C} given by $\{a \in \mathcal{C} \mid \sigma_a \text{ is invertible}\}$.

Proof. For fixed $(a, \alpha) \in (S, \sigma)\text{-Alg}$, the diagram

$$\begin{array}{ccc} Sa & \xrightarrow{\sigma_{Sa}} & S^2a \\ \alpha \downarrow & & \downarrow S\alpha \\ a & \xrightarrow{\sigma_a} & Sa \end{array}$$

commutes by the naturality of σ . Since S is well pointed, this implies $\sigma_a\alpha = S\alpha \cdot \sigma_{Sa} = S\alpha \cdot S\sigma_a = S(\alpha\sigma_a) = S(\text{id}_a) = \text{id}_{Sa}$, therefore $\alpha = \sigma_a^{-1}$. On the other hand, if σ_a is invertible then (a, σ_a^{-1}) is a (S, σ) -algebra.

If $f: a \rightarrow b$ is any morphism, and both σ_a and σ_b are invertible, then

$$\begin{array}{ccc} Sa & \xrightarrow{Sf} & Sb \\ \sigma_a^{-1} \downarrow & & \downarrow \sigma_b^{-1} \\ a & \xrightarrow{f} & b \end{array}$$

commutes by naturality of σ , so U^S is full (being faithful by construction). It follows that $U^S: (S, \sigma)\text{-Alg} \rightarrow \{a \in \mathcal{C} \mid \sigma_a \text{ is invertible}\}$ is bijective on objects and fully faithful, so it is an isomorphism. \square

Lemma 2.5.4. Let (S, σ) be a pointed endofunctor, then $U^S: (S, \sigma)\text{-Alg} \rightarrow \mathcal{C}$ is monadic if and only if it has a left adjoint.

Proof. U^S is conservative and creates all colimits preserved by S . In particular, it preserves coequalizers of U^S -split pairs. \square

Definition 2.5.5. For an endofunctor $F: \mathcal{C} \rightarrow \mathcal{C}$ (or a pointed endofunctor (S, σ)), we say that the *algebraically free monad* on F (respectively (S, σ)) exists if $U^F: F\text{-Alg} \rightarrow \mathcal{C}$ (or $U^S: (S, \sigma)\text{-Alg} \rightarrow \mathcal{C}$) has a left adjoint.

We will denote by **Ord** the category of ordinals.

⁵For instance, it is easy to show that if the unit of a monad satisfies the condition of well pointedness, then the monad is idempotent.

Theorem 2.5.6. Let \mathcal{C} be a category with colimits of chains (that is the domain of the diagram is an ordinal). Let (S, σ) be a well pointed endofunctor such that S preserves κ -filtered colimits. Then, the algebraically free monad on (S, σ) exists. In particular, $\{c \in \mathcal{C} \mid \sigma_c \text{ is an isomorphism}\}$ is a reflective subcategory.

Proof. For a given object $c \in \mathcal{C}$ we define a functor $S^\bullet c: \mathbf{Ord} \rightarrow \mathcal{C}$ by setting $S^0 c := c$, while $S^{\lambda+1} c := S(S^\lambda c)$, with $S^\lambda c \rightarrow S^{\lambda+1} c$ given by $\sigma_{S^\lambda c}$ for $\lambda \in \mathbf{Ord}$. Given a limit ordinal μ , we set $S^\mu c = \text{colim}_{\lambda < \mu} S^\lambda c$.

We claim that $S^\kappa c$ lies in $(S, \sigma)\text{-Alg}$, that is $\sigma_{S^\kappa c}$ is an isomorphism. We will prove this by constructing an inverse $\alpha: S(S^\kappa c) \rightarrow S^\kappa c$.

Since S is κ -accessible, $S^{\kappa+1} c = S(S^\kappa c) = \text{colim}_{\lambda < \kappa} S(S^\lambda c)$. We construct a cocone on $S(S^\bullet c)$ by considering the maps $l_{\lambda+1}: S(S^\lambda c) = S^{\lambda+1} c \rightarrow S^\kappa c$ exhibiting $S^\kappa c$ as a colimit.

$$\begin{array}{ccc} S(S^\lambda c) & = & S^{\lambda+1} c \xrightarrow{\sigma_{S^\lambda c} = \sigma_{S^{\lambda+1} c}} S(S^{\lambda+1} c) \\ \downarrow l_{\lambda+1} & & \downarrow l_{\lambda+2} \\ S^\kappa c & \xlongequal{\quad} & S^\kappa c \end{array}$$

Well pointedness gives us the upper equality and the diagram commutes, hence we get a cocone culminating in $S^\kappa c$, which will then factor uniquely through the cocone culminating in $S(S^\kappa c)$ as $\alpha: S(S^\kappa c) \rightarrow S^\kappa c$. By construction, the following diagram commutes and $l_{\mu+1} \sigma_{S^\mu c} = l_\mu$.

$$\begin{array}{ccccc} S^\mu c & \xrightarrow{\sigma_{S^\mu c}} & S(S^\mu c) & & \\ \downarrow l_\mu & & \downarrow S l_\mu & \searrow l_{\mu+1} & \\ S^\kappa c & \xrightarrow{\sigma_{S^\kappa c}} & S(S^\kappa c) & \xrightarrow{\alpha} & S^\kappa c \end{array}$$

Passing to the colimit, this implies that $\alpha \sigma_{S^\kappa c} = \text{id}_{S^\kappa c}$ because the l_μ on the left and $l_{\mu+1} \sigma_{S^\mu c} = l_\mu$ become identities, hence (S^κ, α) is indeed a (S, σ) -algebra.

We now claim that $l_0: c \rightarrow S^\kappa c$ defines a reflection into the full subcategory given by $\mathcal{B} := \{c \in \mathcal{C} \mid \sigma_c \text{ is an isomorphism}\}$. Firstly, we have shown that $S^\kappa c \in \mathcal{B}$, hence we only need $l_0^*: \mathcal{B}(S^\kappa c, b) \rightarrow \mathcal{C}(c, b)$ to be a bijection for all $b \in \mathcal{B}$.

Since representable functors $\mathcal{C}(-, b)$ send colimits to limits, this immediately reduces to the following: given $b \in \mathcal{B}$, $c \in \mathcal{C}$, the map $\sigma_c^*: \mathcal{C}(Sc, b) \rightarrow \mathcal{C}(c, b)$ is a bijection.

Using well pointedness, we can write the inverse to σ_c^* as $\mathcal{C}(c, b) \rightarrow \mathcal{C}(c, b)$, $f \mapsto \sigma_b^{-1} S f$. \square

Theorem 2.5.7. Let \mathcal{C} be a cocomplete category, $F: \mathcal{C} \rightarrow \mathcal{C}$ an endofunctor. The slice category $(F \downarrow \mathcal{C})$ is cocomplete. Moreover, all colimits preserved by F are computed pointwise, that is $\text{colim}_{\mathcal{J}}(a_i, b_i, \alpha_i: Fa_i \rightarrow b_i) = (\text{colim}_{\mathcal{J}} a_i, \text{colim}_{\mathcal{J}} b_i, \text{colim}_{\mathcal{J}} \alpha_i: F(\text{colim}_{\mathcal{J}} a_i) \rightarrow \text{colim}_{\mathcal{J}} b_i)$.

Proof. Giving a diagram $D: \mathcal{J} \rightarrow (F \downarrow \mathcal{C})$ amounts to giving diagrams $a_\bullet: \mathcal{J} \rightarrow \mathcal{C}$, $b_\bullet: \mathcal{J} \rightarrow \mathcal{C}$ and a natural transformation $\alpha_\bullet: Fa_\bullet \Rightarrow b_\bullet$.

Giving a cocone on this with vertex $(c, d, \gamma: Fc \rightarrow d)$ is equivalent to giving morphisms $\text{colim}_{\mathcal{J}} a_i \rightarrow c$, $\text{colim}_{\mathcal{J}} b_i \rightarrow c$ such that the following diagram commutes for all i .

$$\begin{array}{ccccc} Fa_i & \longrightarrow & F(\text{colim}_{\mathcal{J}} a_i) & \longrightarrow & Fc \\ \downarrow & & \downarrow & & \downarrow \gamma \\ b_i & \longrightarrow & \text{colim}_{\mathcal{J}} b_i & \longrightarrow & d \end{array}$$

Equivalently, we can give a morphism $\text{colim}_{\mathcal{I}} a_i \rightarrow c$ and a morphism from the pushout p to d making the following diagram commute.

$$\begin{array}{ccccc} \text{colim}_{\mathcal{I}} Fa_i & \longrightarrow & F(\text{colim}_{\mathcal{I}} a_i) & \longrightarrow & Fc \\ \downarrow \text{colim}_{\mathcal{I}} \alpha_i & & \downarrow & & \downarrow \gamma \\ \text{colim}_{\mathcal{I}} b_i & \longrightarrow & p & \longrightarrow & d \end{array}$$

We have then the colimit $(\text{colim}_{\mathcal{I}} a_i, p, F(\text{colim}_{\mathcal{I}} a_i) \rightarrow p)$ in $(F \downarrow \mathcal{C})$. In particular, if F preserves this colimit, then the top map $F(\text{colim}_{\mathcal{I}} a_i) \rightarrow Fc$ is an isomorphism, in which case we may take $p = d$ and the identity as the map from p to d . \square

Proposition 2.5.8. If in the theorem above \mathcal{C} is locally presentable and F is accessible, then $(F \downarrow \mathcal{C})$ is locally presentable.

Proof. There exists a regular cardinal κ such that \mathcal{C} is locally κ -presentable and $F(\mathcal{C}_\lambda) \subset \mathcal{C}_\kappa$, with F λ -accessible and $\lambda \leq \kappa$. We claim that the full subcategory $\mathcal{A} := \{(a, b, \alpha: Fa \rightarrow b) \mid a \in \mathcal{C}_\lambda, b \in \mathcal{C}_\kappa\}$ is dense and consists of κ -presentable objects in $(F \downarrow \mathcal{C})$.

The fact that it consists of κ -presentable objects follows from the facts that κ -filtered colimits in $(F \downarrow \mathcal{C})$ are computed pointwise.

To prove density, we want that for each $(c, d, \gamma: Fc \rightarrow d)$ the canonical cocone $(\mathcal{A} \downarrow (c, d, \gamma))$ exhibits (c, d, γ) as a colimit. In the arrow category $[[1], \mathcal{C}]$, $Fc \rightarrow d$ is a colimit of all κ -presentable pairs $c_0, c_1 \in \mathcal{C}_\lambda$.

$$\begin{array}{ccc} c_0 & \longrightarrow & Fc \\ \downarrow & & \downarrow \\ c_1 & \longrightarrow & d \end{array}$$

We need to check that the natural functor $(\mathcal{A} \downarrow (c, d, \gamma)) \rightarrow ([1], \mathcal{C}]_\lambda \downarrow \gamma)$ is final.

Check for yourself that the category $([1], \mathcal{C}]_\lambda \downarrow \gamma)$ we are considering is actually filtered and specifically can always find a pair of morphisms completing the following commutative diagram, where $Fa \rightarrow b$ comes from \mathcal{A} .

$$\begin{array}{ccccc} & & Fa & & \\ & \nearrow & \downarrow & \searrow & \\ c_0 & \longrightarrow & Fc & & \\ \downarrow & & \downarrow b & & \downarrow \\ c_1 & \longrightarrow & d & & \end{array}$$

The codomains form a colimit diagram in \mathcal{C} , hence we are left with checking that the domains form a colimit diagram as well. To do this, we use the fact that $(\mathcal{C}_\lambda \downarrow c) \rightarrow \mathcal{C}$ has colimit c and an argument similar to the previous one. \square

2.6 Algebraically free Monads on a pointed Endofunctor

Let $T: \mathcal{C} \rightarrow \mathcal{C}$ a κ -accessible endofunctor, with \mathcal{C} cocomplete. As we have already shown, the category $(T \downarrow \mathcal{C})$ is cocomplete and κ -filtered colimits in $(T \downarrow \mathcal{C})$ are computed objectwise.

Given a natural transformation $\alpha: T' \rightarrow T$, we get an adjunction

$$(T' \downarrow \mathcal{C}) \begin{array}{c} \xrightarrow{\alpha_!} \\ \perp \\ \xleftarrow{\alpha^*} \end{array} (T \downarrow \mathcal{C})$$

where $\alpha^*(a, b, Ta \xrightarrow{\gamma} b) = (a, b, T'a \rightarrow Ta \xrightarrow{\gamma} b)$ and $\alpha_!$ is given by the pushout

$$\alpha_!(a, b, \beta) = (Ta \xrightarrow{\gamma} c) \quad \begin{array}{ccc} T'a & \xrightarrow{\beta} & b \\ \downarrow \alpha_a & & \downarrow \\ Ta & \xrightarrow{\gamma} & c \end{array}$$

If T, T' are κ -accessible, then α^* is κ -accessible. If we apply this to the case $T' = \text{id}_{\mathcal{C}}$, $\alpha = \tau: \text{id}_{\mathcal{C}} \rightarrow T$, then $(\text{id}_{\mathcal{C}} \downarrow \mathcal{C} = \mathcal{C}^{[1]})$ is the arrow category, and τ^* sends $(a, b, Ta \xrightarrow{\alpha} b)$ to $a \rightarrow Ta \rightarrow b$.

Now use exercise 5.1.

Proposition 2.6.1. If

$$\mathcal{C} \begin{array}{c} \xrightarrow{U} \\ \top \\ \xleftarrow{F} \end{array} \mathcal{D}$$

is an adjunction, (S, σ) is a well-pointed endofunctor on \mathcal{D} and \mathcal{C} has pushouts, then

$$\begin{array}{ccc} FU & \xrightarrow{F\sigma U} & FSU \\ \downarrow \epsilon & & \downarrow \\ \text{id}_{\mathcal{C}} & \xrightarrow{\sigma'} & S' \end{array}$$

defines a well-pointed endofunctor (S', σ') on \mathcal{C} such that the square

$$\begin{array}{ccc} (S', \sigma')\text{-Alg} & \xrightarrow{\bar{U}} & (S, \sigma)\text{-Alg} \\ \downarrow & & \downarrow \\ \mathcal{C} & \xrightarrow{U} & \mathcal{D} \end{array}$$

is a pullback. □

Theorem 2.6.2. Let \mathcal{C} be a cocomplete category, (T, τ) a well-pointed endofunctor on \mathcal{C} with T κ -accessible. Then $(T, \tau)\text{-Alg}$ is a reflexive subcategory of $(T \downarrow \mathcal{C})$ and the algebraically free monad on (T, τ) exists. In particular, by reflexivity, $(T, \tau)\text{-Alg}$ is cocomplete.

Proof. Notice that the functor

$$(T, \tau)\text{-Alg} \rightarrow (T \downarrow \mathcal{C}), \quad (a, \alpha) \mapsto (a, a, Ta \xrightarrow{\alpha} a), \quad f \mapsto (f, f)$$

is fully faithful. It is clearly faithful, and if $(f, g): (a, a, \alpha) \rightarrow (b, b, \beta)$ is a morphism in $(T \downarrow \mathcal{C})$, then we have a commutative diagram

$$\begin{array}{ccccc} a & \xrightarrow{\tau_a} & Ta & \xrightarrow{\alpha} & a \\ \downarrow f & & \downarrow Tf & & \downarrow g \\ b & \xrightarrow{\tau_b} & Tb & \xrightarrow{\beta} & b \end{array}$$

Since (a, α) and (b, β) are algebras, we have $\beta\tau_b = \text{id}_b$ and $\alpha\tau_a = \text{id}_a$. It immediately follows that $f = g$. Moreover, the essential image of this functor is

$$\{(a, b, \gamma) \in (T \downarrow \mathcal{C}) \mid \gamma\tau_a \text{ is an isomorphism}\}$$

Apply the previous proposition to the pullback

$$\begin{array}{ccc} (T, \tau)\text{-Alg} & \longrightarrow & \text{Iso}(\mathcal{C}) \\ \downarrow & & \downarrow \\ (T \downarrow \mathcal{C}) & \xrightarrow{\tau^*} & \mathcal{C}^{[1]} \\ (a, b, \gamma) & \longmapsto & \gamma\tau_a \end{array}$$

and the well-pointed endofunctor $S: \mathcal{C}^{[1]} \rightarrow \mathcal{C}^{[1]}$ given by $S(c \rightarrow d) = \text{id}_d$ with $(S, \sigma)\text{-Alg} = \text{Iso}(\mathcal{C})$. We obtain (S', σ') such that $(T, \tau)\text{-Alg} \cong (S', \sigma')\text{-Alg}$. Since the pushout from the previous proposition consists of κ -accessible functors (here we use that τ^* is κ -accessible), (S', σ') is κ -accessible. It follows that $(S', \sigma')\text{-Alg}$ is reflexive in $(T \downarrow \mathcal{C})$, as claimed.

Thus $(T, \tau)\text{-Alg}$ is cocomplete. Note that the forgetful functor $U^T: (T, \tau)\text{-Alg} \rightarrow \mathcal{C}$ factors as

$$\begin{array}{ccc} (T, \tau)\text{-Alg} & \longrightarrow & (T \downarrow \mathcal{C}) \xrightarrow{\text{dom}} \mathcal{C} \\ (a, \alpha) & \longmapsto & (a, a, \alpha) \mapsto a \end{array}$$

and $(T, \tau)\text{-Alg} \rightarrow \mathcal{C}$ has a left adjoint. We only need to a left adjoint to dom which is given by $c \mapsto (c, Tc, \text{id}_{Tc})$. \square

Theorem 2.6.3. Let \mathcal{C} a cocomplete category and F a κ -accessible endofunctor. The category $F\text{-Alg}$ is cocomplete and the algebraically free monad on F exists, that is the functor $U^F: F\text{-Alg} \rightarrow \mathcal{C}$ has a left adjoint.

Proof. Let T be the coproduct $F + \text{id}_{\mathcal{C}}$ and $\tau: \text{id}_{\mathcal{C}} \rightarrow F + \text{id}_{\mathcal{C}}$ the inclusion. Then $(T, \tau)\text{-Alg} \cong F\text{-Alg}$ is an isomorphism which is compatible with the forgetful functors. \square

For example, we can easily prove the following

Proposition 2.6.4. Let \mathcal{C} a κ -presentable category and $T: \mathcal{C} \rightarrow \mathcal{C}$ κ -accessible. Then $(T \downarrow \mathcal{C})$ is locally κ -presentable.

Proof. Consider the functor $\mathcal{C} \times \mathcal{C} \xrightarrow{F} \mathcal{C} \times \mathcal{C}$, $(a, b) \mapsto (\emptyset, Ta)$. Then $F\text{-Alg} \cong (T \downarrow \mathcal{C})$ and U^F is κ -accessible. Since U^F is monadic, the free objects on the κ -presentable objects form a dense generating set consisting of κ -presentable objects in $F\text{-Alg} \cong (T \downarrow \mathcal{C})$. \square

Remark 2.6.5. An analysis of the construction of (S', σ') in the proof of the previous theorem shows that $S': (T \downarrow \mathcal{C}) \rightarrow (T \downarrow \mathcal{C})$ sends $(a, b, \alpha: Ta \rightarrow b)$ to $(b, c, \gamma: Tb \rightarrow c)$ where

$$T^2a \xrightarrow[\tau_{Ta}]{T\tau_a} T^2a \xrightarrow{T\alpha} Tb \xrightarrow{\gamma} c \quad (1)$$

is a coequalizer diagram in \mathcal{C} . Notice that γ is a coequalizer of $Ta \cdot T\tau_a$ and $T\alpha \cdot \tau_{Ta}$ (see the exercises for more details).

Proposition 2.6.6. Let (S, σ) be a well-pointed endofunctor on \mathcal{C} and let $L: \mathcal{C} \rightarrow \mathcal{C}$ be a functor. If $\pi: S \rightarrow L$ is a natural transformation such that $\pi_c: Sc \rightarrow Lc$ is epic for all $c \in \mathcal{C}$, then $(L, \pi \cdot \sigma)$ is a well-pointed endofunctor and $(L, \pi \sigma)\text{-Alg}$ is equivalent to the full subcategory of $(S, \sigma)\text{-Alg}$ on objects (a, α) such that $\pi_a: Sa \rightarrow La$ is an isomorphism.

Proof. Exercise. □

Now let (T, η, μ) be a monad on a cocomplete category \mathcal{C} and assume T is κ -accessible. We define an endofunctor $L: (T \downarrow \mathcal{C}) \rightarrow (T \downarrow \mathcal{C})$ as follows: Given $(a, b, \alpha: Ta \rightarrow b)$ we set $L(a, b, \alpha) = (b, d, \gamma: Tb \rightarrow d)$ with γ defined by the following pushout in \mathcal{C} .

$$\begin{array}{ccc} T^2a & \xrightarrow{\mu_a} & Ta \\ T\alpha \downarrow & & \downarrow \delta \\ Tb & \xrightarrow{\gamma} & d \end{array}$$

Using this construction we can finally prove the following.

Theorem 2.6.7. Let \mathcal{C} be a complete category, (T, η, μ) a monad over \mathcal{C} with T κ -accessible. Then $T\text{-Alg}$ is reflexive in $(T \downarrow \mathcal{C})$ and cocomplete.

Proof. Recall that $(T, \eta)\text{-Alg}$ is reflexive in $(T \downarrow \mathcal{C})$ and we have a well-pointed endofunctor given by S' described in 1. Remember the functor L just defined.

Since μ_a has a section $T\eta_a$, it is the coequalizer of id_a and $T\eta_a \cdot \mu_a$.

$$\begin{array}{ccccc} T^2a & \xlongequal{\quad} & T^2a & \xrightarrow{T\alpha} & Tb & \xrightarrow{\gamma} & d \\ & \searrow \mu_a & \nearrow T\eta_a & & & & \\ & & Ta & & & & \end{array} \quad (2)$$

Given that $\mu_a \cdot T\eta_a = \text{id}_{Ta}$, we have $\delta = \delta \cdot \mu_a \cdot T\eta_a = \gamma \cdot T\alpha \cdot T\eta_a$. Moreover, γ coequalizes $T\alpha \cdot T\eta_a$ and $T\alpha \cdot \eta_{Ta}$, hence there exists a unique $\pi: c \rightarrow d$ making the following diagram commute.

$$\begin{array}{ccc} Tb & \xrightarrow{\beta} & c \\ \parallel & & \downarrow \pi \\ Tb & \xrightarrow{\gamma} & d \end{array}$$

This defines a natural transformation $(\text{id}, \pi): S(a, b, \alpha) \rightarrow L(a, b, \alpha)$, where the components are epimorphisms because γ is a coequalizer of the diagram in 1 and therefore an epimorphism.

We get then a well-pointed endofunctor $(L, \pi \sigma')$ over $(T \downarrow \mathcal{C})$ with $(L, \pi \sigma')\text{-Alg}$ equivalent to the full subcategory of $(S', \sigma')\text{-Alg}$ given by the objects b such that (id_b, π) is an isomorphism. We also have an equivalence $(T, \eta)\text{-Alg} \rightarrow (S', \sigma')\text{-Alg}$, $(a, \alpha) \mapsto (a, a, \alpha)$, hence we get that $(S', \sigma')\text{-Alg}$ is isomorphic to the full subcategory of $(T, \eta)\text{-Alg}$ given by $\{(a, \alpha) \in (T, \eta)\text{-Alg} \mid (\text{id}_a, \pi): S'(a, a, \alpha) \rightarrow L(a, a, \alpha) \text{ is an isomorphism}\}$.

In this case, the coequalizer of 1 is actually $\alpha: Ta \rightarrow a$, hence our π looks as follows.

$$\begin{array}{ccc} Ta & \xrightarrow{\alpha} & a \\ \parallel & & \downarrow \pi \\ Ta & \xrightarrow{\gamma} & d \end{array}$$

Having π invertible is then equivalent to α being the coequalizer of 2, where $b = a$. If it is a coequalizer diagram, $\alpha\mu_a = \alpha T\alpha$, which implies that (a, α) is a T -algebra. Conversely, if (a, α) is a T -algebra, then this is a split coequalizer in \mathcal{C} . It follows that $T\text{-Alg}$ is equivalent to $(L, \pi\sigma')\text{-Alg}$.

L is accessible since T is and κ -filtered colimits in $(T \downarrow \mathcal{C})$ are computed as in \mathcal{C} , thus $(L, \pi\sigma')\text{-Alg} \rightarrow (T \downarrow \mathcal{C})$ has a left adjoint and therefore $T\text{-Alg} \rightarrow (T \downarrow \mathcal{C})$, $(a, \alpha) \mapsto (a, a, \alpha)$ is fully faithful and has a left adjoint. \square

We have the following result as a consequence.

Theorem 2.6.8. Given a locally κ -presentable category \mathcal{C} and a monad (T, η, μ) of rank κ , $T\text{-Alg}$ is locally κ -presentable.

Proof. We have shown that $\{(Ta, \mu_a) \mid a \in \mathcal{C}_\kappa\}$ is a dense generating system of κ -presentable objects, hence the claim follows from the fact that $T\text{-Alg}$ is cocomplete. \square

3 Monads in 2-category theory

3.1 Monads are monadic

Given an endofunctor $F: \mathcal{C} \rightarrow \mathcal{C}$, an algebraically free monad on F exists if $U^F: F\text{-Alg} \rightarrow \mathcal{C}$ has a left adjoint $F^F: \mathcal{C} \rightarrow F\text{-Alg}$. We write then $T(F) = (U^F F^F, \eta, U^F \epsilon F^F)$ for the resulting monad. From Beck's theorem, we know that $J: F\text{-Alg} \rightarrow T(F)\text{-Alg}$, $(a, \alpha) \mapsto (U^F(a, \alpha) = a, U^F \epsilon_{(a, \alpha)})$ is an equivalence of categories. We also have a natural transformation $\psi: F \Rightarrow T(F)$ corresponding via adjunction to $\alpha: F U^F \Rightarrow U^F$. This gives us a functor $\psi^*: T(F)\text{-Alg} \rightarrow F\text{-Alg}$, $(a, \alpha) \mapsto (a, \alpha \psi_a)$ such that $\psi^* J = \text{id}_{F\text{-Alg}}$. We have the following result.

Proposition 3.1.1. In the described situation, ψ^* is an isomorphism of categories.

Proof. We still have to show that J is surjective on objects, which follows from the fact that both U^F and $U^{T(F)}$ are isofibrations and the fact that two $T(F)$ -algebras isomorphic via id_a are equal. \square

Definition 3.1.2. A *morphism of monads* $(T, \eta, \mu) \rightarrow (T', \eta', \mu')$ over a category \mathcal{C} is a natural transformation $\phi: T \Rightarrow T'$ making the following diagrams commute.

$$\begin{array}{ccc} T^2 & \xrightarrow{\phi^2} & (T')^2 \\ \Downarrow \mu & & \Downarrow \mu' \\ T & \xrightarrow{\phi} & T' \end{array} \quad \begin{array}{ccc} \text{id}_{\mathcal{C}} & \xrightarrow{\eta} & T \\ & \searrow \eta' & \Downarrow \phi \\ & & T' \end{array}$$

The first diagram is equivalent to equating the following two.

$$\begin{array}{ccc} & \mathcal{C} & \\ T \nearrow & & \searrow T \\ \mathcal{C} & \xrightarrow{T} & \mathcal{C} \\ & \Downarrow \phi & \\ & T' & \end{array} = \begin{array}{ccc} & \mathcal{C} & \\ T \nearrow & & \searrow T \\ \mathcal{C} & \xrightarrow{T} & \mathcal{C} \\ & \Downarrow \mu' & \\ & T' & \end{array}$$

The second diagram amounts saying that the following two are equal.

$$\begin{array}{ccc} & \text{id}_{\mathcal{C}} & \\ & \Downarrow \eta & \\ \mathcal{C} & \xrightarrow{T} & \mathcal{C} \\ & \Downarrow \phi & \\ & T' & \end{array} = \begin{array}{ccc} & \text{id}_{\mathcal{C}} & \\ & \Downarrow \eta' & \\ \mathcal{C} & \xrightarrow{T} & \mathcal{C} \\ & \Downarrow \phi & \\ & T' & \end{array}$$

We denote the category of monads over \mathcal{C} by $\text{Mnd}(\mathcal{C})$.

Proposition 3.1.3. The functor

$$\begin{aligned} \mathbf{Mnd}(\mathcal{C})^{\text{op}} &\rightarrow (\mathbf{CAT} \downarrow \mathcal{C}) \\ (T, \eta, \mu) &\mapsto (U^F: T\text{-Alg} \rightarrow \mathcal{C}) \\ \phi &\mapsto \left(\begin{array}{l} \phi^*: T'\text{-Alg} \rightarrow T\text{-Alg} \\ (a, \alpha) \mapsto (a, \alpha\phi_a) \end{array} \right) \end{aligned}$$

is fully faithful.

Proof. Prove by yourself that this is a functor. Consider then two monads T, T' over \mathcal{C} . Giving a functor making the following diagram commute amounts to giving an action on $U^{T'}$, that is $\rho: TU^{T'} \Rightarrow U^{T'}$.

$$\begin{array}{ccc} T'\text{-Alg} & \xrightarrow{\quad} & T\text{-Alg} \\ & \searrow U^{T'} & \swarrow U^T \\ & \mathcal{C} & \end{array}$$

Using the adjunction $F^{T'} \dashv U^{T'}$, this corresponds to a unique natural transformation $T \Rightarrow U^{T'}F^{T'} = T'$. Notice that the T -action axioms for ρ correspond precisely to axioms for morphisms of monads, hence we are done. \square

Proposition 3.1.4 (Algebraically free monads are free). Let $F: \mathcal{C} \rightarrow \mathcal{C}$ be an endofunctor such that the algebraically free monad $T(F)$ exists. Then, for every monad T over \mathcal{C} , the natural transformation $\psi: F \Rightarrow T(F)$ induces a bijection $\psi^*: \mathbf{Mnd}(\mathcal{C}) \rightarrow [\mathcal{C}, \mathcal{C}](F, T)$.

Proof. Consider the following commutative diagram. We want to prove that the horizontal arrows and the one on the right are bijection, which will give us the thesis.

— Placeholder —

The diagram commutes by Yoneda, while the top map are bijections by the previous proposition and the one on the right is a bijection given by composing with $\psi^*: T(F)\text{-Alg} \xrightarrow{\sim} F\text{-Alg}$.

Notice that giving the following commutative diagram amounts to giving a natural transformation $\rho: FG \Rightarrow G$ without requiring any additional property.

$$\begin{array}{ccc} \mathcal{A} & \xrightarrow{\quad} & F\text{-Alg} \\ & \searrow G & \swarrow U^F \\ & \mathcal{C} & \end{array}$$

Here the natural transformations $FU^T \Rightarrow U^T$ correspond bijectively to natural transformations $F \Rightarrow U^TF^T = T$ by adjunction, hence the bottom map is bijective as well. \square

Theorem 3.1.5. Let \mathcal{C} be a locally κ -presentable category. We write $\mathbf{Mnd}_\kappa(\mathcal{C})$ for the full subcategory of $\mathbf{Mnd}(\mathcal{C})$ given by κ -accessible monads. Then, the forgetful functor $U: \mathbf{Mnd}_\kappa(\mathcal{C}) \rightarrow [\mathcal{C}, \mathcal{C}]_\kappa$ is monadic and κ -accessible. In particular, $\mathbf{Mnd}_\kappa(\mathcal{C})$ is locally κ -presentable.

Proof. We already have seen, that for a κ -accessible F , the algebraic free monad $T(F)$ exists and is κ -accessible. So since $\phi: F \Rightarrow T(F)$ is a universal morphism to a monad, the functor $T(-): [\mathcal{C}, \mathcal{C}]_\kappa \rightarrow \mathbf{Mnd}_\kappa(\mathcal{C})$ is left adjoint to U . It is also easy to see, that U is conservative, since the inverse of a natural transformation, which is a morphism of monads, is a morphism of monads again. It thus remains to show, that U creates coequalizers of U -split reflexive pairs. Note first that for any $F \in [\mathcal{C}, \mathcal{C}]_\kappa$ both pre- and post-composition with F preserve coequalizers of such pairs. In particular those preserved by the functors $[\mathcal{C}, \mathcal{C}]_\kappa \rightarrow [\mathcal{C}, \mathcal{C}]_\kappa$ given by the assignments $F \mapsto F \circ F$ and $F \mapsto F \circ F \circ F$. Now for a coequalizer

$$UT_1 \rightrightarrows UT_2 \longrightarrow T'$$

of such a U -split reflexive pair, the diagram

$$UT_1 \circ UT_1 \rightrightarrows UT_2 \circ UT_2 \longrightarrow T' \circ T'$$

is a coequalizer diagram as well. So we get a unique natural transformation $\mu': T' \circ T' \Rightarrow T'$. Using the same argument, one can identify $T' \circ T' \circ T'$ as a coequalizer and verify the associativity. Also the unit $\eta: \text{id} \Rightarrow T'$ can be constructed this way, such that the unit law holds. Thus we have constructed a coequalizer in $\mathbf{Mnd}_\kappa(\mathcal{C})$ and U is monadic. To see the second claim, it remains to show, that U preserves κ -filtered colimits. Note again, that for any $F \in [\mathcal{C}, \mathcal{C}]_\kappa$ pre-composition with F preserves all colimits, since those are computed point wise, and post-composition preserves κ -filtered colimits, since F is κ -accessible. Thus, since κ -filtered colimits are sifted, the functor $[\mathcal{C}, \mathcal{C}]_\kappa \times [\mathcal{C}, \mathcal{C}]_\kappa \rightarrow [\mathcal{C}, \mathcal{C}]_\kappa$ given by composition preserves such. So again we conclude, that the functors given by the assignments $F \mapsto F \circ F$ and $F \mapsto F \circ F \circ F$ preserve κ -filtered colimits and use the same arguments as above to see, that U creates κ -filtered colimits. \square

Proposition 3.1.6. Let \mathcal{C} be locally κ -presentable. Then the functor

$$(-)\text{-Alg}: \mathbf{Mnd}_\kappa(\mathcal{C})^{\text{op}} \rightarrow (\mathbf{CAT} \downarrow \mathcal{C})$$

preserves limits. That is, sends colimits in $\mathbf{Mnd}_\kappa(\mathcal{C})$ to limits.

The following proof uses a generalisation of the endomorphism monad $\langle b, b \rangle$ on an object $b \in \mathcal{C}$, which was discussed in the exercise class. We will prove the details in use later, but give an idea why the statement holds.

Proof. Recall, that this endomorphism monad was given by $\text{Ran}_b b: \mathcal{C} \rightarrow \mathcal{C}$. Similarly, given two objects $a, b \in \mathcal{C}$, one can construct a monad $\langle a, b \rangle$ given by the right Kan extension of b along a .

$$\begin{array}{ccc} \mathcal{C} & & \\ & \swarrow a & \\ & \Rightarrow & * \\ & \searrow b & \\ \mathcal{C} & & \end{array}$$

Furthermore a morphism $f: a \rightarrow b$ gives natural transformations $\langle a, f \rangle: \langle a, a \rangle \rightarrow \langle a, b \rangle$ and $\langle f, b \rangle: \langle b, b \rangle \rightarrow \langle a, b \rangle$. Now, just by unveiling the definitions, one checks, that giving a morphism of monads $T \Rightarrow \langle a, a \rangle$ is equivalent to giving a T -algebra structure on a . Moreover the pullback

$$\begin{array}{ccc}
\langle f, f \rangle & \longrightarrow & \langle b, b \rangle \\
\downarrow & & \downarrow \\
\langle a, a \rangle & \longrightarrow & \langle a, b \rangle
\end{array}$$

induced by a morphism $f: a \rightarrow b$, also gives a monad, such that, for fixed T -algebra structures $T \Rightarrow \langle a, a \rangle$ and $T \Rightarrow \langle b, b \rangle$, there exists a morphism of monads $T \Rightarrow \langle f, f \rangle$, if and only if f is a morphism of T -algebras. Now, we claim there also exist κ -accessible monads $\langle a, a \rangle_\kappa$ and $\langle f, f \rangle_\kappa$, having the property, that for any κ -accessible monad T , there is a natural isomorphism

$$\mathbf{Mnd}_\kappa(\mathcal{C})(T, \langle a, a \rangle_\kappa) \cong \mathbf{Mnd}(\mathcal{C})(T, \langle a, a \rangle)$$

(similarly for $\langle f, f \rangle_\kappa$). The construction of those will be given later in the lecture. Unveiling the constructions given above, one sees that giving an object in $(\text{colim } T_i)\text{-Alg}$ is equivalent to giving a compatible system of T_i -algebra structures on a fixed object $a \in \mathcal{C}$. This shows $(\text{colim } T_i)\text{-Alg} \cong \lim(T_i\text{-Alg})$ in $(\mathbf{CAT} \downarrow \mathcal{C})$. \square

Now we can use this proposition to construct monads via presentations.

Example 3.1.7. Let \mathcal{E} be a locally presentable cartesian closed category, that is there is a product functor $- \times X: \mathcal{E} \rightarrow \mathcal{E}$, which has a right adjoint for all $X \in \mathcal{E}$. Now we start with the endofunctor $F_1: \mathcal{E} \rightarrow \mathcal{E}$ given by the assignment $X \mapsto X \times X + X$. So the category of $F_1\text{-Alg}$ is given by the following data

$$T(F_1)\text{-Alg} \cong F_1\text{-Alg} = \{ (X, m, e): m: X \times X \rightarrow X, e: * \rightarrow X \}$$

satisfying no axioms. Furthermore we take another endofunctor $F_2: \mathcal{E} \rightarrow \mathcal{E}$ given by $X \mapsto X \times X \times X + X + X$ and obtain the data

$$T(F_2)\text{-Alg} \cong F_2\text{-Alg} = \{ (X, f_1, f_2, f_3): f_1: X \times X \times X \rightarrow X, f_2, f_3: X \rightarrow X \}$$

Note that F_1 and F_2 both are κ -accessible by the usual sifted colimit argument, thus the algebraically free monads $T(F_1)$ and $T(F_2)$ exist and are κ -accessible. We now define two functors

$$F_1\text{-Alg} \begin{array}{c} \xrightarrow{G_1} \\ \xrightarrow{G_2} \end{array} F_2\text{-Alg}$$

by the formulas

$$\begin{aligned}
G_1(X, m, e) &= (X, m \circ m \times X, m \circ e \times X, m \circ X \times e) \\
G_2(X, m, e) &= (X, m \circ X \times m, \text{id}_X, \text{id}_X)
\end{aligned}$$

By fullness and faithfulness of $(-)\text{-Alg}$, we get monad morphisms $\varphi_1, \varphi_2: T(F_2) \rightarrow T(F_1)$, inducing G_1 and G_2 up to isomorphism ($G_1 = \varphi_1^*, G_2 = \varphi_2^*$). Now the coequalizer T_{mon} of φ_1 and φ_2 has algebras isomorphic to the equalizer of G_1 and G_2 . So it gives the data of a monoid object in \mathcal{E}

$$T_{\text{mon}}\text{-Alg} = \{ (X, m, e): m \circ m \times X = m \circ X \times e, m \circ e \times X = \text{id}_X, m \circ X \times e = \text{id}_X \}$$

Example 3.1.8. In the exercise classes, we will see **Cat**, the category of small categories and functors, is locally finitely presentable. Let $\mathcal{D} = \{0 \rightrightarrows 1\}$ be the category with two objects and two non trivial parallel morphisms. We now want to show, that small categories with chosen coequalizers are monadic. First we take the endofunctor $F_1: \mathbf{Cat} \rightarrow \mathbf{Cat}$ given by $\mathcal{C} \mapsto [\mathcal{C}, \mathcal{D}]$ and obtain

$$F_1\text{-Alg} = \{(\mathcal{C}, l: [\mathcal{C}, \mathcal{D}] \rightarrow \mathcal{C})\}$$

We now want to say, that such a l is left adjoint to the constant diagram functor $c: \mathcal{C} \rightarrow [\mathcal{C}, \mathcal{D}]$. For this we need unit η and counit ε . To construct them we use the arrow category $[1]$, since to give a natural transformation

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{f} & \mathcal{D} \\ & \Downarrow & \\ \mathcal{C} & \xrightarrow{g} & \mathcal{D} \end{array}$$

it amounts to give a homotopy

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{f} & \mathcal{D} \\ \downarrow & & \searrow \\ \mathcal{C} \times [1] & \longrightarrow & \mathcal{D} \\ \uparrow & & \nearrow \\ \mathcal{C} & \xrightarrow{g} & \mathcal{D} \end{array}$$

so we can get the counit $\varepsilon: l \circ c \rightarrow \text{id}$ using the endofunctor $F_2(\mathcal{C}) = \mathcal{C} \times [1]$. Since the unit has to be of the form $\eta: [\mathcal{D}, \mathcal{C}] \times [1] \rightarrow [\mathcal{D}, \mathcal{C}]$, there has to be done a bit more. But the functor $[\mathcal{D}, -]$ is right adjoint to $- \times \mathcal{D}$ and via this adjunction such a morphism corresponds to $[\mathcal{D}, \mathcal{C}] \times [1] \times \mathcal{D} \rightarrow \mathcal{C}$. So we can get the unit using the endofunctor given by $F_3(\mathcal{C}) = [\mathcal{D}, \mathcal{C}] \times [1] \times \mathcal{D}$. Now we can express, having "same" natural transformations with desired source and target via equalizers of $(-)\text{-Alg}$. Also the Δ -identities then can be expressed in a second coequalizer step. But in this construction the following problem appears: Since the morphisms preserve the chosen coequalizers on the nose, they will rarely arise in nature.

3.2 Symmetric monoidal categories

Definition 3.2.1. A *monoidal category* is a tuple $(\mathcal{V}, \otimes, I, \alpha, \lambda, \rho)$, where \mathcal{V} is a category, $\otimes: \mathcal{V} \times \mathcal{V} \rightarrow \mathcal{V}$ a functor, $I \in \mathcal{V}$ an object, $\alpha: (- \otimes -) \otimes - \Rightarrow - \otimes (- \otimes -)$, $\lambda: I \otimes - \Rightarrow \text{id}$ and $\rho: - \otimes I \Rightarrow \text{id}$ natural isomorphisms such that for every $W, X, Y, Z \in \mathcal{V}$ the diagrams

$$\begin{array}{ccccc} & & (W \otimes X) \otimes (Y \otimes Z) & & \\ \alpha_{W \otimes X, Y, Z} \nearrow & & & \searrow \alpha_{W, X, Y \otimes Z} & \\ ((W \otimes X) \otimes Y) \otimes Z & & & & W \otimes (X \otimes (Y \otimes Z)) \\ \alpha_{W, X, Y} \otimes \text{id}_Z \searrow & & & \nearrow \text{id}_W \otimes \alpha_{X, Y, Z} & \\ (W \otimes (X \otimes Y)) \otimes Z & \xrightarrow{\alpha_{W, X \otimes Y, Z}} & W \otimes ((X \otimes Y) \otimes Z) & & \\ & & & & \\ & & (X \otimes I) \otimes Y & \xrightarrow{\alpha_{X, I, Y}} & X \otimes (I \otimes Y) \\ & \searrow \rho_X \otimes \text{id}_Y & & \swarrow \text{id}_X \otimes \lambda_Y & \\ & & X \otimes Y & & \end{array}$$

commute. We call \otimes the *tensor product*, I the *unit object* or *tensor unit*, α the *associator*, λ the *left unitor* and ρ the *right unitor*.

Example 3.2.2. We now list some monoidal categories.

1. If \mathcal{E} is a category with finite products, then $(\mathcal{E}, \times, *)$ is monoidal, with α , λ and ρ induced by the universal property. Instances of this are **Set**, **Cat**, **Grp**, **sSet**, **Top**, **CGTop** and **CGHTop**.
2. $(\mathbf{Ab}, \otimes_{\mathbb{Z}}, \mathbb{Z})$ and, given a commutative ring R , $(\mathbf{Mod}_R, \otimes_R, R)$ and $(\mathbf{dgMod}_R, \otimes_R, R)$.
3. The order $\overline{\mathbb{R}}_+ = [0, \infty]$ with $\otimes = +$, $I = 0$.
4. A monoid in **Cat** or **CAT** is a monoidal category such that α , λ and ρ are identities. This is the case of $[\mathcal{C}, \mathcal{C}]$, $[\mathcal{C}, \mathcal{C}]_{\kappa}$, $\Phi\text{-Cocts}[\mathcal{C}, \mathcal{C}]$.

We mention without proof the following fundamental theorem.

Theorem 3.2.3 (Mac Lane). Any diagram built from \otimes , I , α , λ , ρ and their iterations is commutative.

Given a word of objects tensored among them, any two choices of bracketing are uniquely isomorphic. This result is plausible if \otimes is derived from an universal property as in (1) – (3) and clear if \mathcal{V} is strict, like in (3) and (4), while the general proof uses a rewriting argument which can be found in *Categories for the Working Mathematician*.

Definition 3.2.4. A *lax monoidal functor* from \mathcal{V} to \mathcal{W} is a triple (F, ϕ_0, ϕ) , where $F: \mathcal{V} \rightarrow \mathcal{W}$ is a functor, $\phi_0: I_{\mathcal{W}} \rightarrow FI_{\mathcal{V}}$ a morphism and $\phi: \otimes_{\mathcal{W}} \circ (F \times F) \Rightarrow F \circ \otimes_{\mathcal{V}}$ a natural transformation such that for all $X, Y, Z \in \mathcal{V}$ the diagrams

$$\begin{array}{ccc}
 (FX \otimes_{\mathcal{W}} FY) \otimes_{\mathcal{W}} FZ & \xrightarrow{\alpha^{\mathcal{W}}} & FX \otimes_{\mathcal{W}} (FY \otimes_{\mathcal{W}} FZ) \\
 \downarrow \phi_{X,Y} \otimes_{\mathcal{W}} FZ & & \downarrow FX \otimes_{\mathcal{W}} \phi_{X,Z} \\
 F(X \otimes_{\mathcal{V}} Y) \otimes_{\mathcal{W}} FZ & & FX \otimes_{\mathcal{W}} F(Y \otimes_{\mathcal{V}} Z) \\
 \downarrow \phi_{X \otimes_{\mathcal{V}} Y, Z} & & \downarrow \phi_{X,Y} \otimes_{\mathcal{W}} Z \\
 F((X \otimes_{\mathcal{V}} Y) \otimes_{\mathcal{V}} Z) & \xrightarrow{F\alpha^{\mathcal{V}}} & F(X \otimes_{\mathcal{V}} (Y \otimes_{\mathcal{V}} Z))
 \end{array}$$

$$\begin{array}{ccc}
 I_{\mathcal{W}} \otimes_{\mathcal{W}} FX & \xrightarrow{\phi_0 \otimes_{\mathcal{W}} FX} & FI_{\mathcal{V}} \otimes_{\mathcal{W}} FX \\
 \downarrow \lambda_X^{\mathcal{W}} & & \downarrow \phi_{I_{\mathcal{V}}, X} \\
 FX & \xleftarrow{F\lambda_X^{\mathcal{V}}} & F(I_{\mathcal{V}} \otimes_{\mathcal{V}} X)
 \end{array}
 \quad
 \begin{array}{ccc}
 FX \otimes_{\mathcal{W}} I_{\mathcal{W}} & \xrightarrow{FX \otimes_{\mathcal{W}} \phi_0} & FI_{\mathcal{V}} \otimes_{\mathcal{W}} FX \\
 \downarrow \rho_X^{\mathcal{W}} & & \downarrow \phi_{X, I_{\mathcal{V}}} \\
 FX & \xleftarrow{F\rho_X^{\mathcal{V}}} & F(I_{\mathcal{V}} \otimes_{\mathcal{V}} X)
 \end{array}$$

are commutative.

If we reverse the direction of ϕ_0 and ϕ we get *oplax monoidal functors*.

A *strong (strict) monoidal functor* is a lax monoidal functor such that ϕ_0 and ϕ are isomorphisms (respectively identities).

A *monoidal natural transformation* from (F, ϕ_0, ϕ) to (G, ψ_0, ψ) is a natural transformation $\gamma: F \Rightarrow G$ such that the diagrams

$$\begin{array}{ccc}
 FX \otimes_{\mathcal{W}} FY & \xrightarrow{\gamma_X \otimes_{\mathcal{W}} \gamma_Y} & GX \otimes_{\mathcal{W}} GY \\
 \downarrow \phi_{X,Y} & & \downarrow \psi_{X,Y} \\
 F(X \otimes_{\mathcal{V}} Y) & \xrightarrow{\gamma_{X \otimes_{\mathcal{V}} Y}} & G(X \otimes_{\mathcal{V}} Y)
 \end{array}
 \quad
 \begin{array}{ccc}
 & I_{\mathcal{W}} & \\
 \phi_0 \swarrow & & \searrow \psi_0 \\
 FI_{\mathcal{V}} & \xrightarrow{\gamma_{I_{\mathcal{V}}}} & GI_{\mathcal{V}}
 \end{array}$$

commute.

Proposition 3.2.5. Lax monoidal functors compose and monoidal natural transformations whisker.

Proposition 3.2.6. There is a finitary monad T on \mathbf{Cat} such that $T\text{-Alg}$ is the category of monoidal categories and strict monoidal functors.

Proof. We can write down a presentation of this monad using the finitary endofunctors $X \mapsto X \times X$, $X \mapsto X \times X \times X \times [1]$. \square

Example 3.2.7. Given a locally small monoidal category \mathcal{V} , the functor $\mathcal{V}(I, -): \mathcal{V} \rightarrow \mathbf{Set}$ is lax monoidal, with $\phi_0: \{*\} \rightarrow \mathcal{V}(I, I)$, $*$ $\mapsto \text{id}_I$ and $\phi_{X,Y}$ sending $(f, g) \in \mathcal{V}(I, X) \times \mathcal{V}(I, Y)$ to $(f \otimes g) \circ \lambda_I^{-1} = (f \otimes g) \circ \rho_I^{-1}: I \xrightarrow{\sim} I \otimes I \rightarrow X \otimes Y$. It is universally denoted by $V: \mathcal{V} \rightarrow \mathbf{Set}$ and, if \mathcal{V} has coproducts, it has a left adjoint given by $F: \mathbf{Set} \rightarrow \mathcal{V}$, $S \mapsto \amalg_S I$.

Assuming for simplicity that \mathcal{V} is cocomplete, it is easy to show that F is strong monoidal if \otimes preserves colimits in each variable by using that \mathbf{Set} is the free cocomplete category on $\{*\}$.

The previous example is an instance of a more general phenomenon, as shown by the following result.

Theorem 3.2.8. Let $F: \mathcal{V} \rightarrow \mathcal{W}$ be a left adjoint to U . If \mathcal{V}, \mathcal{W} are monoidal, F, U lax and η, ϵ monoidal natural transformations, then F is strong monoidal. Conversely, if (F, ϕ_0, ϕ) is strong monoidal and U is any right adjoint, then

$$\begin{array}{ccc}
 I_{\mathcal{V}} \xrightarrow{\eta_{I_{\mathcal{V}}}} UFI_{\mathcal{V}} & UX \otimes_{\mathcal{V}} UY \xrightarrow{\psi_{X,Y}} U(X \otimes_{\mathcal{V}} Y) & \\
 \searrow \psi_0 & \downarrow U(\epsilon_{X \otimes_{\mathcal{V}} Y}) & \uparrow \eta_{UX \otimes_{\mathcal{W}} UY} \\
 & UI_{\mathcal{W}} & \\
 & UF(UX \otimes_{\mathcal{V}} UY) \xrightarrow{U\phi_{X,Y}^{-1}} U(FUX \otimes_{\mathcal{W}} FUY) &
 \end{array}$$

define a lax monoidal structure on U which is unique with the property that η, ϵ are monoidal.

Proof. Exercise. \square

Example 3.2.9. Given a homomorphism of commutative rings $R \rightarrow S$, then $S \otimes_R - \dashv U$, where U is the restriction on scalars, and $- \otimes_R S: \mathbf{Mod}_R \rightarrow \mathbf{Mod}_S$ is a monoidal adjunction.

The free module functor $\mathbf{Set} \rightarrow \mathbf{Mod}_R$ is strong monoidal by the previous example.

Remark 3.2.10. The last example still holds if we substitute to \mathbf{Mod}_R any cocomplete monoidal category \mathcal{V} with $- \otimes V, V \otimes -$ cocontinuous.

Definition 3.2.11. A monoid in a monoidal category \mathcal{V} is a triple (M, m, u) where $m: M \otimes M \rightarrow M$ is the *multiplication*, $u: I \rightarrow M$ the *unit* and the diagrams

$$\begin{array}{ccccc}
 & (M \otimes M) \otimes M & \xrightarrow{\alpha_{M,M,M}} & M \otimes (M \otimes M) & \\
 m \otimes M \swarrow & & & & \searrow M \otimes m \\
 M \otimes M & & & & M \otimes M \\
 & m \searrow & & m \swarrow & \\
 & M & & M &
 \end{array}$$

$$\begin{array}{ccccc}
I \otimes M & \xrightarrow{u \otimes M} & M \otimes M & \xleftarrow{M \otimes u} & M \otimes I \\
& \searrow \lambda_M & \downarrow m & \swarrow \rho_M & \\
& & M & &
\end{array}$$

commute.

Morphisms of monoids are maps $f: M \rightarrow M'$ such that $m' \circ (f \otimes f) = fm$, $fu = u'$.

We write $\mathbf{Mon}(\mathcal{V})$ for the category of monoids over \mathcal{V} .

Remark 3.2.12. If \mathcal{V} is additive, monoids are often called algebras as well because $\mathbf{Mon}(\mathbf{Mod}_R) = \mathbf{Alg}_R$, $\mathbf{Mon}(\mathbf{dgMod}_R) = \mathbf{dgAlg}_R$.

Example 3.2.13. For R a commutative ring, $\mathbf{-Alg}_R$ is locally finitely presentable

Proposition 3.2.14. If $F \dashv U: \mathcal{V} \rightarrow \mathbf{W}$ is a monoidal adjunction (F strong, U monoidal), then $F \dashv U$ lifts to an adjunction of monoids

$$\begin{array}{ccc}
& \xrightarrow{\quad \overline{F} \quad} & \\
\mathbf{Mon}(\mathcal{V}) & \perp & \mathbf{Mon}(\mathbf{W}) \\
& \xleftarrow{\quad \overline{U} \quad} & \\
\downarrow & & \downarrow \\
\mathcal{V} & \xrightarrow{\quad F \quad} & \mathbf{W} \\
& \perp & \\
& \xleftarrow{\quad U \quad} &
\end{array}$$

where $\overline{F}(M, m, n) = (FM, Fm \circ \phi_{M,M}^F, Fn \circ \phi_0^F)$ and $\overline{U}(M', m', n') = (UM, Um' \circ \phi_{U_M, M}, Un' \circ \phi_{U_0})$.

Proof. The axioms for lax monoidal functors show that these are indeed monads, naturality of ϕ^F, ϕ^U shows that Ff is a monoid morphism if f is. It follows that $\overline{F}, \overline{U}$ are indeed functors. Axioms for monoidal transformations show η, ϵ are monoid morphisms. Hence $\overline{F} \dashv \overline{U}$. \square

Example 3.2.15. i) A ring homomorphism $R \rightarrow S$ of commutative rings induces base

$$\text{change functor } \mathbf{Alg}_R \begin{array}{c} \xrightarrow{S \otimes_R -} \\ \xleftarrow{\text{forget}} \end{array} \mathbf{Alg}_S \quad \text{This works equally for commutative algebras, etc.}$$

ii) Let \mathbf{C} locally κ -presentable. Then $[\mathbf{C}, \mathbf{C}]_\kappa \rightarrow [\mathbf{C}, \mathbf{C}]$ is strict monoidal and a left adjoint, so it lifts to a left adjoint $\mathbf{Mod}_\kappa(\mathbf{C}) \rightarrow \mathbf{Mod}(\mathbf{C})$. Why is it left adjoint? Let $K: \mathbf{C}_\kappa \rightarrow \mathbf{C}$ be the inclusion. Since Φ is cocontinuous we have

$$\begin{array}{ccc}
[\mathbf{C}, \mathbf{C}] & \xrightarrow{\quad} & [\mathbf{C}, \mathbf{C}] \\
\uparrow \text{Lan}_K \simeq & \nearrow \text{Lan}_K & \uparrow \\
[\mathbf{C}_\kappa, \mathbf{C}] & \xrightarrow{\quad K^* \quad} &
\end{array}$$

Hence the inclusion is indeed a left adjoint.

Remark 3.2.16. This completes the proof that $(-)\text{-Alg}: \mathbf{Mnd}_\kappa(\mathbf{C})^{\text{op}} \rightarrow \mathbf{Cat}/\mathbf{C}$ is full and faithful and sends colimits of monads to limits of categories.

Let $(\mathcal{V}, \otimes, I)$ be a monoidal category such that $x \otimes -$ and $- \otimes x$ preserve coproducts for all $x \in \mathcal{V}$. The category of \mathcal{V} -matrices with index set S is $[S \times S, \mathcal{V}] = \prod_{S \times S} \mathcal{V}$ and denoted by $\text{Mat}(\mathcal{V}, S)$. There is a natural monoidal structure on $\text{Mat}(\mathcal{V}, S)$ given by matrix multiplication:

$$M(x, y)_{(x, y) \in S^2} \otimes N(x, y)_{(x, y) \in S^2} = \left(\sum_{z \in S} M(z, y) \otimes N(x, z) \right)_{(x, y) \in S^2}.$$

and unit $(I_{x, y})_{(x, y) \in S^2}$ where $I_{x, y} = I$ if $x = y$ and $I_{x, y} = \emptyset$, else. The (α, λ, ρ) on $\text{Mat}(\mathcal{V}, S)$ are induced by the ones on \mathcal{V} and via the universal properties coproducts.

Definition 3.2.17. A \mathcal{V} -category \mathcal{A} with object set S is a monoid in $\text{Mat}(\mathcal{V}, S)$ i.e. for each pair $(a, b) \in S^2$ an object $\mathcal{A}(a, b) \in \mathcal{V}$, called the \mathcal{V} -object of homomorphisms. Moreover we have units $I \xrightarrow{\text{id}_a} \mathcal{A}(a, a)$ and composition homomorphisms

$$\mathcal{A}(b, c) \otimes \mathcal{A}(a, b) \xrightarrow{C_{a, b, c}} \mathcal{A}(a, c)$$

such that the following axioms hold:

$$\begin{array}{ccc} (\mathcal{A}(c, d) \otimes \mathcal{A}(b, c)) \otimes \mathcal{A}(a, b) & \xrightarrow{\cong} & \mathcal{A}(c, d) \otimes (\mathcal{A}(b, c) \otimes \mathcal{A}(a, b)) \\ \downarrow c \otimes \mathcal{A}(a, b) & & \downarrow \mathcal{A}(c, d) \otimes c \\ \mathcal{A}(b, d) \otimes \mathcal{A}(a, b) & & \mathcal{A}(c, d) \otimes \mathcal{A}(a, c) \\ & \searrow c \quad \swarrow c & \\ & \mathcal{A}(a, d) & \end{array}$$

and the two unit axioms hold. This is in fact the definition if \mathcal{V} does not have all coproducts or $x \otimes -$ does not preserve them.

Example 3.2.18. • $\mathcal{V} = \mathbf{Top} \rightsquigarrow$ topological categories

- $\mathcal{V} = \mathbf{Ab} \rightsquigarrow$ additive categories
- $\mathcal{V} = \mathbf{Mod}_R \rightsquigarrow$ linear categories
- $\mathcal{V} = \mathbf{Mod}_R^{\text{d.g.}} \rightsquigarrow$ graded categories
- $\mathcal{V} = \mathbf{Cat} \rightsquigarrow$ 2-categories
- $\mathcal{V} = n\text{-}\mathbf{Cat} \rightsquigarrow$ strict $(n + 1)$ -categories

Proposition 3.2.19. If \mathcal{V} is presentable and both $x \otimes -$ and $- \otimes x$ are κ -accessible for all $x \in \mathbf{D}$ and preserve coproducts, then the category of \mathcal{V} -categories with fixed object set S is locally κ -presentable.

Proof. It suffices to show that the matrices are κ -accessible in each variable, which follows from the fact, that colimits commute. \square

Definition 3.2.20. Let \mathcal{V} be a monoidal category with coproducts, and such that $x \otimes -$ and $- \otimes x$ preserve coproducts for each $x \in \mathcal{V}$. Let $f: S \rightarrow T$ be a map of sets. We write

$$\begin{aligned} f^*: \text{Mat}(\mathcal{V}, T) &\rightarrow \text{Mat}(\mathcal{V}, S) \\ (T \times T \xrightarrow{M} \mathcal{V}) &\mapsto (S \times S \xrightarrow{f \times f} T \times T \xrightarrow{M} \mathcal{V}) \end{aligned}$$

and $f_*: \text{Mat}(\mathcal{V}, S) \rightarrow \text{Mat}(\mathcal{V}, T)$ for its left adjoint. It sends M to

$$(f_*M)(a, b) = \sum_{\{(x, y): fx=a, fy=b\}} M(x, y).$$

Proposition 3.2.21. The left adjoint is strong monoidal.

Definition 3.2.22. A \mathcal{V} -functor $(S, \mathcal{A}) \rightarrow (T, \mathcal{B})$ is a pair $(F, (F_{a,b})_{(a,b) \in S^2})$ where $F: S \rightarrow T$ is a function and $(F_{a,b})_{(a,b) \in S^2}$ is a monoid morphism $\mathcal{A} \rightarrow F^*\mathcal{B}$, that is $F_{a,b}: \mathcal{A}(a, b) \rightarrow \mathcal{B}(Fa, Fb)$ is a morphism in \mathcal{V} for any $a, b \in S$.

We denote the resulting category of small \mathcal{V} -categories by $\mathcal{V}\text{-Cat}$. We may also define $\mathcal{V}\text{-CAT}$ as the category of all \mathcal{V} -categories. We may define the category of \mathcal{V} -graphs analogously. We have an obvious forgetful functor from \mathcal{V} -categories to \mathcal{V} -graphs.

Example 3.2.23. We now want to see, what happens if we take $\mathcal{V} = \mathbf{Cat}$. That is unveiling the data of a 2-category. First we have a set/class of objects $\text{Ob}(\mathcal{K})$, called 0-cells, and for any two $A, B \in \text{Ob}(\mathcal{K})$, we have a category $\mathcal{K}(A, B)$. We call the objects of $\mathcal{K}(A, B)$ 1-cells from A to B and denote them $f: A \rightarrow B$. The morphisms in $\mathcal{K}(A, B)$ are called 2-cells, denoted by

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ & \Downarrow \alpha & \\ A & \xrightarrow{g} & B \end{array}$$

The categorical structure of $\mathcal{K}(A, B)$ now tells us, that we have vertical compositions

$$\begin{array}{ccc} A & \xrightarrow{\quad} & B \\ & \Downarrow \alpha & \\ A & \xrightarrow{\quad} & B \\ & \Downarrow \beta & \\ A & \xrightarrow{\quad} & B \end{array}$$

which is associative and has a unit $\text{id}_f: f \Rightarrow f$. We also have identities $\text{id}_A: * \rightarrow \mathcal{K}(A, A)$, written $\text{id}_A: A \rightarrow A$, and composition functors $\mathcal{K}(B, C) \times \mathcal{K}(A, B) \rightarrow \mathcal{K}(A, C)$ which give in particular horizontal composition of 1-cells $A \rightarrow B \rightarrow C$ and whiskering operations

$$(B \xrightarrow{g} C, A \xrightarrow{f} B) \longmapsto A \xrightarrow{gf} C$$

$$\begin{array}{ccc} B & \xrightarrow{g} & C \\ \Downarrow \alpha & & \\ B & \xrightarrow{h} & C \end{array} \longmapsto \begin{array}{ccc} A & \xrightarrow{gf} & C \\ \Downarrow \alpha_f & & \\ A & \xrightarrow{hf} & C \end{array}$$

and similarly on the other side. To say that this defines a functor, means that these operations satisfy the interchange law: That is, given a diagram of the form

$$\begin{array}{ccccc} A & \xrightarrow{f} & B & \xrightarrow{h} & C \\ & \Downarrow \alpha & & \Downarrow \beta & \\ A & \xrightarrow{g} & B & \xrightarrow{k} & C \end{array}$$

we have $\beta_g \cdot h_\alpha = k_\alpha \cdot \beta_f$ or in pictures

$$\begin{array}{ccc} \cdot \xrightarrow{\quad} \cdot & \xrightarrow{\quad} & \cdot \xrightarrow{\quad} \cdot \\ \Downarrow \psi & & \Downarrow \psi \\ \cdot \xrightarrow{\quad} \cdot & \xrightarrow{\quad} & \cdot \xrightarrow{\quad} \cdot \end{array} = \begin{array}{ccc} \cdot \xrightarrow{\quad} \cdot & \xrightarrow{\quad} & \cdot \xrightarrow{\quad} \cdot \\ \Downarrow \psi & & \Downarrow \psi \\ \cdot \xrightarrow{\quad} \cdot & \xrightarrow{\quad} & \cdot \xrightarrow{\quad} \cdot \end{array}$$

This follows from the fact that giving a functor $F: \mathcal{C} \times \mathcal{D} \rightarrow \mathcal{E}$ amounts to giving compatible sets of functors $F(-, d)$ and $F(c, -)$ for all $(c, d) \in \mathcal{C} \times \mathcal{D}$.

Example 3.2.24. Examples of 2-categories are

- (i) **Cat** and **CAT** with small or locally small categories as 0-cells, functors as 1-cells and natural transformations as 2-cells.
- (ii) Locally κ -presentable categories, κ -accessible functors and natural transformations.
- (iii) Monoidal categories, lax monoidal functors and monoidal natural transformations.
- (iv) For \mathcal{V} a monoidal category, $\mathcal{V}\text{-Cat}$ and $\mathcal{V}\text{-CAT}$ are 2-categories. 0-cells are small/locally small \mathcal{V} -categories, 1-cells are \mathcal{V} -functors and 2-cells are \mathcal{V} -natural transformations.

Definition 3.2.25. Let \mathcal{A} and \mathcal{B} be two \mathcal{V} -categories and $F, G: \mathcal{A} \rightarrow \mathcal{B}$ be two \mathcal{V} -functors. Then a \mathcal{V} -natural transformation $F \Rightarrow G$ is a collection of morphisms $(\alpha_A: I \rightarrow \mathcal{B}(FA, GA))_{A \in \mathcal{A}}$ in \mathcal{V} (note that this collection can be indexed by a class of objects), such that for all objects A, B in \mathcal{A} the diagram

$$\begin{array}{ccccc}
 & I \otimes \mathcal{A}(A, B) & \xrightarrow{\alpha_B \otimes F} & \mathcal{B}(FB, GB) \otimes \mathcal{B}(FA, FB) & \\
 \nearrow \lambda^{-1} & & & & \searrow \circ \\
 \mathcal{A}(A, B) & & & & \mathcal{B}(FA, GB) \\
 \searrow \rho^{-1} & & & & \nearrow \circ \\
 & \mathcal{A}(B, B) \otimes I & \xrightarrow{G \otimes \alpha_A} & \mathcal{B}(GA, GB) \otimes \mathcal{B}(FA, GA) &
 \end{array}$$

is commutative. We then first define the whiskering operations. So consider the diagram

$$\mathcal{A}' \xrightarrow{K} \mathcal{A} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \end{array} \mathcal{B} \xrightarrow{L} \mathcal{B}'$$

then we define $(L, \alpha)_{A \in \mathcal{A}}$ via the composition

$$I \xrightarrow{\alpha_A} \mathcal{B}(FA, GA) \xrightarrow{L} \mathcal{B}'(LFB, LGB)$$

and $(\alpha, K)_{A' \in \mathcal{A}'}$ by $\alpha_{KA'}: I \rightarrow \mathcal{B}(FKA', GKA')$. Clearly α_K is a \mathcal{V} -natural transformation $FK \Rightarrow GK$. To see this for L_α compare the needed diagram with

$$\begin{array}{ccccc}
 \mathcal{B}(FB, GB) \otimes \mathcal{B}(FA, FB) & \xrightarrow{L \otimes L} & \mathcal{B}'(LFB, LGB) \otimes \mathcal{B}'(LFA, LFB) & & \\
 \searrow \circ & & & \searrow \circ & \\
 & \mathcal{B}(FA, GB) & \xrightarrow{L} & \mathcal{B}'(LFA, LGB) & \\
 \nearrow \circ & & & \nearrow \circ & \\
 \mathcal{B}(GA, GB) \otimes \mathcal{B}(FA, GA) & \xrightarrow{L \otimes L} & \mathcal{B}'(LGA, LGB) \otimes \mathcal{B}'(LFA, LGA) & &
 \end{array}$$

Now given a diagram of the form

$$\begin{array}{ccc}
 & F & \\
 \mathcal{A} & \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \\ \Downarrow \beta \\ \xrightarrow{H} \end{array} & \mathcal{B} \\
 & H &
 \end{array}$$

we define the vertical composition $(\beta \cdot \alpha)_A$ by

$$I \cong I \otimes I \xrightarrow{\beta_A \otimes \alpha_A} \mathcal{B}(GA, HA) \otimes \mathcal{B}(FA, GA) \xrightarrow{\circ} \mathcal{B}(FA, HA)$$

with unit natural transformation $\text{id}_{FA}: I \rightarrow \mathcal{B}(FA, FA)$. We leave it to the reader, to check associativity and the interchange law.

Example 3.2.26. For $\mathcal{V} = \mathbf{Set}$, we get precisely the 2-categories of categories, functors and natural transformations. So what happens if we take **Cat** or even **CAT**. 0-cells are 2-categories. 1-cells are 2-functors or **Cat**-functors, that is: Given two 2-categories $\mathcal{K}, \mathcal{K}'$, we have an assignment $\text{Ob}(\mathcal{K}) \rightarrow \text{Ob}(\mathcal{K}')$ of the form $A \mapsto FA$ and for any two objects A, B in $\text{Ob}(\mathcal{K})$ there is a functor $F_{A,B}: \mathcal{K}(A, B) \rightarrow \mathcal{K}'(A, B)$

$$A \begin{array}{c} \xrightarrow{f} \\ \Downarrow \alpha \\ \xrightarrow{g} \end{array} B \longmapsto FA \begin{array}{c} \xrightarrow{Ff} \\ \Downarrow F\alpha \\ \xrightarrow{Fg} \end{array} FB$$

to say that this defines a functor is exactly to say that this assignment respects vertical composition. The first \mathcal{V} -functor axiom says that $F(\text{id}_A) = \text{id}_{FA}$ and the second, that the diagram

$$\begin{array}{ccc} \mathcal{K}(B, C) \times \mathcal{K}(A, B) & \xrightarrow{\circ} & \mathcal{K}(A, C) \\ F \times F \downarrow & & \downarrow F \\ \mathcal{K}'(FB, FC) \times \mathcal{K}'(FA, FB) & \xrightarrow{\circ} & \mathcal{K}'(FA, FC) \end{array}$$

commutes, so F preserves the whiskering operation. What are 2-natural transformations?

$$\mathcal{K} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \end{array} \mathcal{K}'$$

For all objects A in $\text{Ob}(\mathcal{K})$ we have a morphism $\alpha_A: * \rightarrow \mathcal{K}'(FA, GA)$, i.e. a 1-cell $\alpha_A: FA \rightarrow GA$, such that the **Cat**-naturality axioms hold. That is the diagram

$$\begin{array}{ccccc} & & \mathcal{K}'(FB, GB) \times \mathcal{K}'(FA, FB) & & \\ & \nearrow \alpha_B \times F & & \searrow \circ & \\ \mathcal{K}(A, B) & & & & \mathcal{K}'(FA, GB) \\ & \searrow G \times \alpha_A & & \nearrow \circ & \\ & & \mathcal{K}'(GA, GB) \times \mathcal{K}'(FA, GA) & & \end{array}$$

commutes. On objects that say, that for any 1-cell $f: A \rightarrow B$ the diagram

$$\begin{array}{ccc} FA & \xrightarrow{Ff} & FB \\ \alpha_A \downarrow & & \downarrow \alpha_B \\ GA & \xrightarrow{Gf} & GB \end{array}$$

commutes and on morphisms it says that for all 2-cells $\varphi: f \Rightarrow g$, we have

$$\begin{array}{ccc} \begin{array}{ccc} FA & \xrightarrow{Ff} & FB \\ \Downarrow F\varphi & & \\ \alpha_A \downarrow & & \downarrow \alpha_B \\ GA & \xrightarrow{Gg} & GB \end{array} & = & \begin{array}{ccc} FA & \xrightarrow{Ff} & FB \\ \alpha_A \downarrow & & \downarrow \alpha_B \\ GA & \xrightarrow{Gf} & GB \\ \Downarrow G\varphi & & \\ & & \end{array} \end{array}$$

Such transformations are called (strict) 2-natural transformations. The constructed categories will be denoted by **2-Cat** and **2-CAT**.

Definition 3.2.27. Let \mathcal{V} be a monoidal category. A \mathcal{V} -monad is a monad in $\mathcal{V}\text{-Cat}$ or $\mathcal{V}\text{-CAT}$. In other words, a \mathcal{V} -monad on a \mathcal{V} -category \mathcal{C} (a 0-cell in $\mathcal{V}\text{-CAT}$) is a \mathcal{V} -functor $T: \mathcal{C} \rightarrow \mathcal{C}$ equipped with \mathcal{V} -natural transformations μ and η filling the usual pasting diagrams.

The goal of the following section is to use them to define new \mathcal{V} -categories from old ones and develop enriched category theory. Namely, we will construct a new \mathcal{V} -category of T -algebras out of a \mathcal{V} -monad T . For this we need the underlying *ordinary* or *unenriched* category of a \mathcal{V} -category \mathcal{C} . We have a lax monoidal functor $V: \mathcal{V} \rightarrow \mathbf{Set}$ which induces the functor

$$\begin{aligned} \mathcal{V}\text{-CAT} &\longrightarrow \mathbf{Set}\text{-CAT} = \mathbf{CAT} \\ \mathcal{C} &\longmapsto V_* \mathcal{C} \end{aligned}$$

Here $V_* \mathcal{C}$ has the same object class as \mathcal{C} and $V_* \mathcal{C}(a, b) = V(\mathcal{C}(a, b))$. For the composition we use the lax monoidal structure of \mathcal{V} , i.e. the morphisms in $V_* \mathcal{C}$ from a to b are given by morphisms $I \xrightarrow{f} \mathcal{C}(a, b)$. The composition of the morphisms $f: a \rightarrow b$ and $g: b \rightarrow c$ is the morphism defined as

$$g \circ f = I \cong I \otimes I \xrightarrow{g \otimes f} \mathcal{C}(b, c) \otimes \mathcal{C}(a, b) \xrightarrow{\circ} \mathcal{C}(a, c).$$

We write \mathcal{C}_0 for the underlying unenriched category of \mathcal{C} . It would thus be good if $(T\text{-Alg})_0 \cong T_0\text{-Alg}$ in the sense we defined before. Hence the objects should be T_0 -algebras, i.e. pairs (A, α) , $\alpha: TA \rightarrow A \cong I \xrightarrow{\alpha} \mathcal{C}(TA, A)$ such that the two algebra axioms hold.

Example 3.2.28. A few examples of the action of the functor $\mathcal{V}\text{-CAT} \rightarrow \mathbf{CAT}$:

1. when $\mathcal{V} = \mathbf{Mod}_R$ we just forget the additive structure of the hom-sets;
2. if $\mathcal{V} = \mathbf{Top}$ we forget the topology;
3. for $\mathcal{V} = \mathbf{dgMod}_R$ we consider cycles of degree zero;
4. if $\mathcal{V} = (\mathbf{sSet}, \times)$ then $V_* \mathcal{C}$ forgets all the simplices in $\mathcal{C}(a, b)$ except the 0-simplices, i.e. the vertices;
5. if $\mathcal{V} = [0, \infty]$, a \mathcal{V} -enriched category is a metric space (in the sense of Lawvere) and the composition is given by the triangle inequality. The functor above just sees the poset of real numbers as a set.

From now on we assume that \mathcal{V} has equalizers.

Proposition 3.2.29. Let \mathcal{C} be a \mathcal{V} -category and (T, η, μ) a \mathcal{V} -monad on \mathcal{C} . For algebras $(A, \alpha), (B, \beta) \in T_0\text{-Alg}$ let

$$\begin{array}{ccccc} & & \mathcal{C}(TA, TB) & & \\ & \nearrow T & & \nwarrow \beta_* & \\ T\text{-Alg}((A, \alpha), (B, \beta)) & \xrightarrow{U} \mathcal{C}(A, B) & \xrightarrow{\alpha^*} & \mathcal{C}(TA, B) & \end{array}$$

be an equalizer in \mathcal{V} , where β_* is the composition

$$\mathcal{C}(TA, TB) \cong I \otimes \mathcal{C}(TA, TB) \xrightarrow{\beta \otimes \text{id}} \mathcal{C}(TB, B) \otimes \mathcal{C}(TA, TB) \xrightarrow{\circ} \mathcal{C}(TA, B)$$

and similarly $\alpha^* = \circ \cdot \text{id} \otimes \alpha$. Then there is a unique way to define a structure of \mathcal{V} -category with objects equal to $T_0\text{-Alg}$ and hom-object $T\text{-Alg}((A, \alpha), (B, \beta)) \in \mathcal{V}$ such that U becomes a \mathcal{V} -functor.

Proof. For the identities note that $\text{id}_A: I \rightarrow \mathcal{C}(A, A)$ equalizes the two arrows, if $(A, \alpha) = (B, \beta)$. Namely, we have $\text{id}_A \alpha = \alpha T \text{id}_A$ thus we get a factorization

$$\begin{array}{ccc} T\text{-Alg}((A, \alpha), (A, \alpha)) & \xrightarrow{U} & \mathcal{C}(A, A) \\ & \nwarrow \text{dashed} & \nearrow \text{id}_A \\ & I & \end{array}$$

$\exists! \text{id}_{(A, \alpha)}$

and by U being regular monic we have to define $\text{id}_{(A, \alpha)}$ as this dashed arrow if we want U to be a \mathcal{V} -functor. Similarly we want to define composition s.t. the diagram

$$\begin{array}{ccc} T\text{-Alg}((B, \beta), (C, \gamma)) \otimes T\text{-Alg}((A, \alpha), (B, \beta)) & \xrightarrow{U \otimes U} & \mathcal{C}(B, C) \otimes \mathcal{C}(A, C) \\ \downarrow & & \downarrow \circ \\ T\text{-Alg}((A, \alpha), (C, \gamma)) & \xrightarrow{U} & \mathcal{C}(A, C) \end{array}$$

commutes. So one has to check that $\circ \cdot U \otimes U$ equalizes the two arrows defining the equalizer at the bottom of the diagram. One checks that this is the case by translating the usual proof that morphisms of T -algebras compose first into a proof just using the hom-sets (not their elements) and then into a proof in the monoidal category \mathcal{V} . It remains to check that this defines a \mathcal{V} -category and that U is indeed a \mathcal{V} -functor. The first follows from the fact that each $U_{((A, \alpha), (B, \beta))}$ is a monomorphism and the fact that \mathcal{C} is a \mathcal{V} -category. By design, the diagrams above are exactly the \mathcal{V} -functor axioms for U . This also shows uniqueness. \square

Example 3.2.30. If \mathcal{V} = “sets with structure” we just get the corresponding substructure on the morphism set, e.g. subspace topology, submodules etc.

1. If G is a topological group then $G \times -: \mathbf{Top} \rightarrow \mathbf{Top}$ is a **Top**-monad if **Top** is a cartesian closed category of topological spaces such as compactly generated weak hausdorff spaces. From that we get the topological category of G -spaces.
2. If $\mathcal{V} = \mathbf{Ab}$ and R is a ring, then **Ab** is an **Ab**-category since we can add morphisms of abelian groups and this is \mathbb{Z} -bilinear. Moreover $R \otimes_{\mathbb{Z}} -: \mathbf{Ab} \rightarrow \mathbf{Ab}$ is an additive monad. $T\text{-Alg}$ is simply **Mod** $_R$ with addition of R -module homomorphisms.
3. For $\mathcal{V} = \mathbf{Cat}$ we have 2-monads T in **2-CAT**. From a 2-category K with a 2-monad $T: K \rightarrow K$ we get a new 2-category $T\text{-Alg}$. 0-cells are elements of $T_0\text{-Alg}$, namely pairs (A, α) s.t.

$$\begin{array}{ccc} T^2 A & \xrightarrow{T\alpha} & T A \\ \mu_A \downarrow & & \downarrow \alpha \\ T A & \xrightarrow{\alpha} & A \end{array} \quad \begin{array}{ccc} A & \xrightarrow{\eta_A} & T A \\ & \searrow & \downarrow \alpha \\ & & A \end{array}$$

are commutative. A 1-cell in $T\text{-Alg}$ is simply a morphism in $T_0\text{-Alg}$, that is, a 1-cell $A \xrightarrow{f} B$ in K s.t. the diagram

$$\begin{array}{ccc} T A & \xrightarrow{Tf} & T B \\ \alpha \downarrow & & \downarrow \beta \\ A & \xrightarrow{f} & B \end{array}$$

is commutative. A 2-cell in $T\text{-Alg}$ is a morphism in the equalizer

$$T\text{-Alg}((A, \alpha), (B, \beta)) \xrightarrow{U} K(A, B) \begin{array}{c} \xrightarrow{T} K(TA, TB) \\ \xrightarrow{\alpha^*} K(TA, B) \end{array} \xrightarrow{\beta_*}$$

i.e. a 2-cell $\varphi: f \Rightarrow g$ such that

$$\begin{array}{ccc} TA & \xrightarrow{T\varphi} & TB \\ \downarrow \alpha & \Downarrow & \downarrow \\ A & \xrightarrow{g} & B \end{array} = \begin{array}{ccc} TA & \longrightarrow & TB \\ \downarrow & & \downarrow \\ A & \xrightarrow{\varphi} & B \end{array}$$

In this case we can talk about pseudomorphisms and lax/oplax morphisms. For lax morphisms those are the squares

$$\begin{array}{ccc} TA & \xrightarrow{Tf} & TB \\ \downarrow & \swarrow \bar{f} & \downarrow \\ A & \xrightarrow{f} & B \end{array}$$

(\bar{f} is an isomorphism in the pseudo case) subject to some axioms.

We get four 2-categories from T : the original $T\text{-Alg}$ is called $T\text{-Alg}_S$ and its 1-cells are the *strict* morphisms of algebras. We have non full inclusions

$$\begin{array}{ccc} & & T\text{-Alg}_L \\ & \nearrow & \\ T\text{-Alg}_S & \longrightarrow & T\text{-Alg}_P \\ & \searrow & \\ & & T\text{-Alg}^L \end{array}$$

There is a 2-monad T on \mathbf{Cat} s.t. $T\text{-Alg}_P = \mathbf{Mon}_{str}(\mathbf{Cat})$ with strong morphisms.

Now we want to define opposite \mathcal{V} -categories and \mathcal{V} -functors of several variables using the tensor product of \mathcal{V} -categories. This is similar to the product of categories $\mathcal{C} \times \mathcal{D}$, however there $(f, f') \cdot (g, g') = (fg, f'g')$, which changes the order of f' and g . The following definition allows us to fix this.

Definition 3.2.31. Let \mathcal{V} be a monoidal category. A *braiding* on \mathcal{V} is a natural isomorphism

$$\begin{array}{ccc} \mathcal{V} \times \mathcal{V} & \xrightarrow{\tau} & \mathcal{V} \times \mathcal{V} \\ \otimes \searrow & \xrightarrow[\gamma]{\sim} & \swarrow \otimes \\ & \mathcal{V} & \end{array},$$

where τ is the switch functor and for all $A, B, C \in \mathcal{V}$ the diagram

$$\begin{array}{ccccc}
 & & A \otimes (B \otimes C) & \xrightarrow{\gamma} & (B \otimes C) \otimes A \\
 & \nearrow \alpha & & & \searrow \alpha \\
 (A \otimes B) \otimes C & & & & B \otimes (C \otimes A) \\
 & \searrow \gamma \otimes \text{id} & & & \nearrow \text{id} \otimes \gamma \\
 & & (B \otimes A) \otimes C & \xrightarrow{\alpha} & B \otimes (A \otimes C)
 \end{array}$$

and the one obtained by inverting the α commute.

A braiding is called a *simmetry* if $\gamma_{A,B} \circ \gamma_{B,A} = \text{id}$ for all $A, B \in \mathcal{V}$.

Remark 3.2.32. If γ is a simmetry, then either one of the above hexagons implies the other. Moreover, the diagram

$$\begin{array}{ccc}
 I \otimes A & \xrightarrow{\gamma_{I,A}} & A \otimes I \\
 \searrow \lambda_A & & \swarrow \rho_A \\
 & A &
 \end{array}$$

commutes.

Example 3.2.33.

- (i) If $\mathcal{E} = (\mathcal{E}, \times, *)$ is cartesian, then the switch $\tau: \mathcal{E} \times \mathcal{E} \rightarrow \mathcal{E}$ defines a simmetry. This is the case of **Set**, categories of presheaves, **sSet**, **Top**, **Cat**, etc.
- (ii) If R is a unital commutative ring, \mathbf{Mod}_R is a symmetric monoidal category with $\gamma: M \otimes_R N \rightarrow N \otimes_R M$ the canonical isomorphism.
- (iii) If A is an abelian group, write $A\text{-}\mathbf{Mod}_R$ for the category of A -graded R -modules, that is $\Pi_A \mathbf{Mod}_R$. We have a functor $A\text{-}\mathbf{Mod}_R \times A\text{-}\mathbf{Mod}_R \rightarrow A\text{-}\mathbf{Mod}_R$, $((V_i), (W_i)) \mapsto (\bigoplus_{i+j=k} V_i \otimes_R W_j)$. This has a monoidal structure with unit R concentrated in degree $0 \in A$. There are obvious choices for λ and ρ .

A *normalized 3-cocycle* on A with values in R^\times is a function $h: A \times A \times A \rightarrow R^\times$ such that, for any tuple $(i, j, k, l) \in A^4$

$$\begin{aligned}
 h(i, 0, j) &= 1, \\
 h(j, k, l) \cdot h(i, j + l, k) \cdot h(i, j, k) &= h(i, j, k + l) \cdot h(i + j, k, l).
 \end{aligned}$$

We define $\alpha^h: ((U_\bullet) \otimes (V_\bullet)) \otimes (W_\bullet) \rightarrow (U_\bullet) \otimes ((V_\bullet) \otimes (W_\bullet))$ on the component given by the triple (i, j, k) as

$$\begin{aligned}
 (U_i \otimes V_j) \otimes W_k &\rightarrow U_i \otimes (V_j \otimes W_k) \\
 (a \otimes b) \otimes c &\mapsto h(i, j, k) \cdot a \otimes (b \otimes c)
 \end{aligned}$$

The two axioms for normalized 3-cocycles say precisely that this is a monoidal structure on $A\text{-}\mathbf{Mod}_R$. Also, by considering modules concentrated in a single degree, one finds that all associators are of this form.

A *normalized abelian 3-cocycle* is an arrow h as above plus a map $c: A \times A \rightarrow R^\times$ such that, for any triple $(i, j, k) \in A^3$,

$$\begin{aligned}
 h(j, k, i) \cdot c(i, j + k) \cdot h(i, j, k) &= c(i, k) \cdot h(j, i, k) \cdot c(i, j), \\
 h(k, i, j)^{-1} \cdot c(i + j, k) \cdot h(i, j, k)^{-1} &= c(i, k) \cdot h(i, j, k)^{-1} \cdot c(j, k).
 \end{aligned}$$

Given such (h, c) , we get a braiding defined by

$$\begin{aligned} V_i \otimes W_j &\rightarrow W_j \otimes V_i \\ a \otimes b &\mapsto c(i, j) \cdot b \otimes a \end{aligned}$$

and this is a symmetry if and only if $c(i, j) = c(j, i) = 1$ for all i, j .

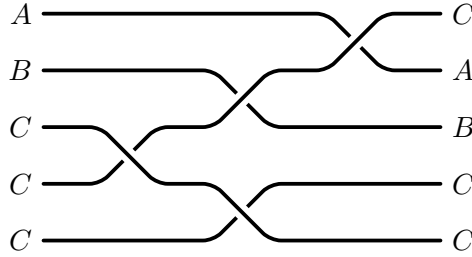
If we take the constant map $h \equiv 1$, then the two axioms say exactly that c is bilinear.

For $A = \mathbb{Z}$, given any $u \in R^\times$ we can define $c(i, j) = u^{|ij|}$, which gives a symmetry if $u^2 = 1$. In particular, if R is a domain we have $R^\times = \{\pm 1\}$ and the condition is always satisfied.

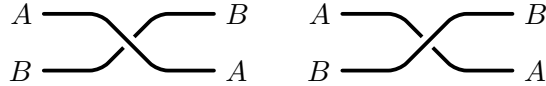
There are then only two symmetric monoidal structures on $\mathbb{Z}\text{-}\mathbf{Mod}_R$, the trivial one and the one given by $V_i \otimes W_j \rightarrow W_j \otimes V_i$, $a \otimes b \mapsto (-1)^{|ij|} \cdot b \otimes a$. The latter is the symmetry given by the *Koszul sign rule*.

(iv) The symmetry given by the Koszul sign rule lifts to a symmetry on \mathbf{dgMod}_R , while the trivial one does not.

Remark 3.2.34. The coherence theorem for braided (symmetric) monoidal categories does not say that “all diagrams commute”, in particular $\gamma_{X,X}: X \otimes X \rightarrow X \otimes X$ in general is not the identity on $X \otimes X$. Instead, it tells us that the morphism is completely given by a labelled braid, for example



completely describes a map $A \otimes B \otimes C \otimes C \otimes C \rightarrow C \otimes A \otimes B \otimes C \otimes C$. If γ is a symmetry, then only the permutation of the objects matters, hence the following braids induce the same morphism.



Definition 3.2.35. Let \mathcal{V} and \mathcal{W} be braided monoidal categories, $F: \mathcal{V} \rightarrow \mathcal{W}$ a lax/strong/strict monoidal functor. We call F a *braided lax/strong/strict monoidal functor* if the diagram

$$\begin{array}{ccc} FA \otimes_{\mathcal{W}} FB & \xrightarrow{\gamma_{\mathcal{W}}} & FB \otimes_{\mathcal{W}} FA \\ \downarrow \phi & & \downarrow \phi \\ F(A \otimes_{\mathcal{V}} B) & \xrightarrow{F\gamma_{\mathcal{V}}} & F(B \otimes_{\mathcal{V}} A) \end{array}$$

commutes.

A *braided natural transformation* is just a monoidal natural transformation between braided monoidal functors.

If \mathcal{V} and \mathcal{W} are braided symmetric monoidal categories, then the braided functors and natural transformations are also called symmetric.

Example 3.2.36.

- (i) If $\phi: R \rightarrow S$ is a map of commutative rings, then $S \otimes_R -: \mathbf{Mod}_R \rightarrow \mathbf{Mod}_S$ is a symmetric strong monoidal functor.
- (ii) If A is an abelian group, (h, c) a normalized abelian 3-cocycle on A with values in R^\times and $\phi: R \rightarrow S$ a ring homomorphism, we have that $S \otimes_R -: A\text{-}\mathbf{Mod}_R^{(h,c)} \rightarrow A\text{-}\mathbf{Mod}_S^{(\phi h, \phi c)}$ is a braided strong monoidal functor. In particular, base change is a symmetric strong monoidal functor for both the trivial and the Koszul symmetry on $\mathbb{Z}\text{-}\mathbf{Mod}_R$.
- (iii) If F is a braided strong monoidal left adjoint, then the right adjoint is braided lax monoidal.

Definition 3.2.37. Let \mathcal{V} be a braided monoidal category. A monoid (M, μ, η) in \mathcal{V} is *commutative* if

$$\begin{array}{ccc} M \otimes M & \xrightarrow{\gamma_{M,M}} & M \otimes M \\ & \searrow \mu & \swarrow \mu \\ & M & \end{array}$$

commutes.

A *morphism of commutative monoids* is just a morphism of monoids.

Remark 3.2.38. In general, a lax monoidal functor will not lift to commutative monoids, but a braided lax monoidal functor will. It follows that we have a 2-functor $\mathbf{CMon}: \mathbf{BrMonCAT} \rightarrow \mathbf{CAT}$, $\mathcal{V} \mapsto \mathbf{CMon}(\mathcal{V})$, sending braided monoidal categories to their categories of commutative monoids.

Theorem 3.2.39. If \mathcal{V} is a locally presentable monoidal category with $- \otimes -$ cocontinuous in both variables, then $\mathbf{CMon}(\mathcal{V}) \rightarrow \mathcal{V}$ is monadic and accessible. Also, $\mathbf{CMon}(\mathcal{V})$ is locally κ -presentable if \mathcal{V} is.

Proof. Adapt the one for all monoids with an additional action. □

Definition 3.2.40. Let \mathcal{V} be a braided monoidal category. Define the *opposite* of a \mathcal{V} -category \mathcal{A} by:

- $\text{Ob}(\mathcal{A}^{\text{op}}) = \text{Ob}(\mathcal{A})$,
- $\mathcal{A}^{\text{op}}(A, B) = \mathcal{A}(B, A)$,
- $\text{id}_{\mathcal{A}}$ the same morphism as for \mathcal{A} and
- composition by the diagram

$$\begin{array}{ccc} \mathcal{A}^{\text{op}}(B, C) \otimes \mathcal{A}^{\text{op}}(A, B) & \xlongequal{\quad} & \mathcal{A}(C, B) \otimes \mathcal{A}(B, A) \\ \downarrow \circ_{\mathcal{A}^{\text{op}}} & & \downarrow \gamma \\ \mathcal{A}^{\text{op}}(A, C) = \mathcal{A}(C, A) & \xleftarrow{\circ_{\mathcal{A}}} & \mathcal{A}(B, A) \otimes \mathcal{A}(C, B) \end{array}$$

We want to talk about \mathcal{V} -functors of “several variables.” For this we need $\mathcal{A} \otimes \mathcal{B}$.

Definition 3.2.41. Let \mathcal{A}, \mathcal{B} be \mathcal{V} -categories. Define the \mathcal{V} -category $\mathcal{A} \otimes \mathcal{B}$ by

- $\text{Ob}(\mathcal{A} \otimes \mathcal{B}) = \text{Ob}(\mathcal{A}) \times \text{Ob}(\mathcal{B})$
- $(\mathcal{A} \otimes \mathcal{B})((A, B), (A', B')) = \mathcal{A}(A, A') \otimes \mathcal{B}(B, B')$

- identities: $I \cong I \otimes I \xrightarrow{\text{id}_A \otimes \text{id}_B} \mathcal{A} \otimes \mathcal{B}((A, B), (A, B))$ and
- compositions

$$\begin{array}{c}
(\mathcal{A}(A', A'') \otimes \mathcal{B}(B', B'')) \otimes (\mathcal{A}(A, A') \otimes \mathcal{B}(B, B')) \\
\downarrow \text{isomorphism built from } \gamma\text{'s} \\
\mathcal{A}(A', A'') \otimes \mathcal{A}(A, A') \otimes \mathcal{B}(B', B'') \otimes \mathcal{B}(B, B') \\
\downarrow \circ \otimes \circ \\
\mathcal{A}(A, A'') \otimes \mathcal{B}(B, B'')
\end{array}$$

Note: The first isomorphism is unique, if \mathcal{V} is symmetric.

The final ingredient for Yoneda is enrichment of \mathcal{V} over itself. For this we need an internal Hom-functor.

Definition 3.2.42. A monoidal category \mathcal{V} is *closed monoidal* if for any $X \in \mathcal{V}$ the functors $X \otimes -$ and $- \otimes X$ have right adjoints $[X, -]_l$ and $[-, X]_r$. We denote the unit and counit by coev and ev respectively. For example

$$\text{ev}^X: [X, Y]_r \otimes X \xrightarrow{(\text{ev}^X)_Y} Y.$$

Remark 3.2.43. If \mathcal{V} is braided, we have $- \otimes X \cong X \otimes -$, so $[-, -]_l$ exists if and only $[-, -]_r$ does and they are isomorphic. We simply write $[-, -] = [-, -]_r$ in this case. In other words $\mathcal{V}(X \otimes Y, Z) \cong \mathcal{V}(X, [Y, Z])$.

Remark 3.2.44. If $- \otimes X$ and $X \otimes -$ have right adjoints, the monoidal natural transformations may *not* define a braiding γ ! We need more compatibility.

Proposition 3.2.45. Let \mathcal{V} be a right-closed (that is $[-, -]_r$ exists) monoidal category. Then the morphisms

$$([Y, Z]_r \otimes [X, Y]_r) \rightarrow [X, Z]_r \quad \text{and} \quad I \rightarrow [X, X]_r$$

corresponding to

$$\begin{array}{ccc}
([Y, Z]_r \otimes [X, Y]_r) \otimes X & \longrightarrow & Z \\
\alpha \downarrow & & \uparrow \text{ev}^Y \\
[Y, Z]_r \otimes ([X, Y]_r \otimes X) & \xrightarrow{\text{id} \otimes \text{ev}^X} & [Y, Z]_r \otimes Y
\end{array}
\quad \text{and} \quad I \otimes X \xrightarrow{\lambda_X} X$$

give a \mathcal{V} -category structure on $\text{Ob}(\mathcal{V})$ with underlying category isomorphic to \mathcal{V} .

Proof. The proof is slightly tedious; We refer to Kelly's book.

A more abstract argument is possible if \mathcal{V} is locally presentable and biclosed. Then we have a monoidal left adjoint

$$\begin{array}{ccc}
\mathcal{V} & \xrightarrow{\quad \top \quad} & [\mathcal{V}, \mathcal{V}]_\kappa \xrightarrow{\quad \top \quad} [\mathcal{V}, \mathcal{V}] \\
X & \longmapsto & - \otimes X
\end{array}$$

Use the $[\mathcal{V}, \mathcal{V}]$ -enrichment on $\text{Ob}(\mathcal{V})$ given by $\langle \mathcal{V}, \mathcal{W} \rangle$ (previous exercise). Pull this back along the right adjoint and get $R\langle \mathcal{V}, \mathcal{W} \rangle = [V, W]_r$. \square

Example 3.2.46. 1) If \mathcal{V} is a category of “sets with structure,” that is if $V: \mathcal{V} \rightarrow \mathbf{Set}$ is monadic for example $\mathcal{V} = \mathbf{Mod}_R, \mathbf{Ab}$ or $\mathcal{V} = \mathbf{Top}_{\text{CGWH}}$, then $[-, -]$ is just the obvious structure on Hom-sets of \mathcal{V} . Specifically if M, N are R -modules, then $\text{Hom}_R(M, N)$ has the natural R -modules structure.

2) For $\mathcal{V} = \mathbf{Cat}$, $[A, B]$ is just the category of functors from A to B . Note that this is not just structure on the Hom-sets, we need the additional data of natural transformations.

3) Even more involved: $\mathbf{dgMod}_R, \mathbf{sSet}$:

If we have all these structures, that is a symmetric monoidal closed category, we can define a \mathcal{V} -functor

$$\mathcal{C}(-, -): \mathcal{C}^{\text{op}} \otimes \mathcal{C} \rightarrow \mathcal{V}$$

as follows. In order to do this, we use the following way of constructing functors $\mathcal{A} \otimes \mathcal{B} \rightarrow \mathcal{C}$:

Proposition 3.2.47. To give a \mathcal{V} -functor $T: \mathcal{A} \otimes \mathcal{B} \rightarrow \mathcal{C}$ amounts to giving families of functors $T(A, -): \mathcal{B} \rightarrow \mathcal{C}$ and $(T(-, B): \mathcal{A} \rightarrow \mathcal{C})$ such that

- i) On objects $T(A, -)(B) = T(-, B)(A)$ which we denote by $T(A, B)$.
- ii) For all $A, A' \in \mathcal{A}$ and $B, B' \in \mathcal{B}$ the diagram

$$\begin{array}{ccc}
 & \xrightarrow{T(-, B') \otimes T(A, -)} & \\
 \mathcal{A}(A, A') \otimes \mathcal{B}(B, B') & & \mathcal{C}(T(A, B'), T(A', B')) \otimes \mathcal{C}(T(A, B), T(A', B')) \\
 \downarrow \gamma & & \downarrow \circ \\
 & & \mathcal{C}(T(A, B), T(A', B')) \\
 \mathcal{B}(B, B') \otimes \mathcal{A}(A, A') & & \uparrow \circ \\
 & & \mathcal{C}(T(A', B), T(A', B')) \otimes \mathcal{C}(T(A, B), T(A', B)) \\
 & \xrightarrow{T(A', -) \otimes T(-, B)} &
 \end{array}$$

commutes. This means that “it does not matter in which way we compose.” In this case, $T_{(A, B), (A', B')}$ is given by the now well defined composite in the above diagram.

Proof. Long exercise. □

Now we need to define $\mathcal{C}(c, -): \mathcal{C} \rightarrow \mathcal{V}$, $\mathcal{C}(-, c): \mathcal{C}^{\text{op}} \rightarrow \mathcal{V}$. But $\mathcal{C}(-, c)$ is just $\mathcal{C}^{\text{op}}(c, -)$, so we only need to prove that the covariant one is a well defined functor. On objects we define $\mathcal{C}(c, -)(c') = \mathcal{C}(c, c') \in \mathcal{V}$. The action on morphisms is given by the morphism

$$\mathcal{C}(c, -)_{c', c'': \mathcal{C}(c', c'') \rightarrow [\mathcal{C}(c, c'), \mathcal{C}(c', c'')]$$

corresponding under adjunction to the composition

$$\mathcal{C}(c', c'') \otimes \mathcal{C}(c, c') \xrightarrow{\circ} \mathcal{C}(c, c'')$$

The diagram in the above proposition commutes by adjunction since composition is associative.

Example 3.2.48. 1) For \mathcal{V} = “sets with structure,” the functor $\mathcal{C}(-, -)$ simply defines a lift

$$\begin{array}{ccc} & & \mathcal{V} \\ & \nearrow \mathcal{C}(-, -) & \downarrow V \\ \mathcal{C}^{\text{op}} \otimes \mathcal{C} & & \mathbf{Set} \\ & \searrow \mathcal{C}_0(-, -) & \end{array}$$

we remember that $\mathcal{C}(-, -)$ is an R -module, a topological space, etc.

- 2) For $\mathcal{C} = \mathbf{Cat}$, we get $\mathcal{K}(-, -): \mathcal{K}^{\text{op}} \times \mathcal{K} \rightarrow \mathbf{Cat}$, $(x, y) \mapsto \mathcal{K}(x, y)$. a 2-functor, where the action on 1-cells is given by whiskering on either side.
- 3) For \mathcal{V} itself, we get $\mathcal{V}(-, -): \mathcal{V}^{\text{op}} \otimes \mathcal{V} \rightarrow \mathcal{V}$, $(V, W) \mapsto [V, W]$. The underlying set of this is $\mathcal{V}(I, [V, W]) \cong \mathcal{V}(V, W)$. To avoid confusion, we will write $\mathcal{V}_0(V, W)$ for the *set* of morphisms in \mathcal{V} .

Proposition 3.2.49. There is a \mathcal{V} -functor $- \otimes -: \mathcal{V} \otimes \mathcal{V} \rightarrow \mathcal{V}$ which on Hom-objects is the morphism

$$[X, X'] \otimes [Y, Y'] \rightarrow [X \otimes Y, X' \otimes Y']$$

corresponding by adjunction to the morphism

$$([X, X'] \otimes [Y, Y']) \otimes (X \otimes Y) \cong ([X, X'] \otimes X) \otimes ([Y, Y'] \otimes Y) \xrightarrow{\text{ev}^X \otimes \text{ev}^Y} X' \otimes Y'.$$

For this functor, α, λ, ρ are \mathcal{V} -natural transformations. Moreover for each X , the maps ev^X and coev^X are \mathcal{V} -natural, so $- \otimes X$ is a left adjoint to the functor $[X, -]$ in $\mathcal{V}\text{-CAT}$

Proof. By adjunction. Straightforward, but tedious (see Kelley). □

Remark 3.2.50. One can check that all “reasonable” morphisms built from the canonical ones are \mathcal{V} -natural. For example, if $f: A \rightarrow B$ is a morphism in \mathcal{A}_0 , we get \mathcal{V} -natural transformations

$$\mathcal{A}(f, -): \mathcal{A}(B, -) \Rightarrow \mathcal{A}(A, -) \quad \text{and} \quad \mathcal{A}(-, f): \mathcal{A}(-, A) \Rightarrow \mathcal{A}(-, B)$$

defined by applying

$$\mathcal{A}_0^{\text{op}} \times \mathcal{A}_0 \rightarrow (\mathcal{A}^{\text{op}} \otimes \mathcal{A}) \xrightarrow{\mathcal{A}(-, -)_0} \mathcal{V}_0$$

to (f, id) and (id, f) respectively.

For details—or more precisely—a good list of instructions on how to proceed efficiently can be found in Kelley’s “Basic concepts,” sections 1.7 and 1.8.

3.3 The weak Yoneda Lemma

Remark 3.3.1. With the morphism just defined, we can express \mathcal{V} -naturality of $\alpha: F \Rightarrow G$ where $F, G: \mathcal{C} \rightarrow \mathcal{D}$ by saying that the following diagram commutes for all $c, c' \in \mathcal{C}$:

$$\begin{array}{ccc} \mathcal{C}(c, c') & \xrightarrow{F} & \mathcal{D}(Fc, Fc') \\ \downarrow G & & \downarrow \mathcal{D}(Fc, \alpha_{c'}) \\ \mathcal{D}(Gc, Gc') & \xrightarrow{\mathcal{D}(\alpha_c, Gc')} & \mathcal{D}(Fc, Gc') \end{array}$$

Theorem 3.3.2 (Weak Yoneda lemma). Let \mathcal{V} be a symmetric monoidal category, \mathcal{A} a \mathcal{V} -category, $F: \mathcal{A} \rightarrow \mathcal{V}$ a \mathcal{V} -functor, $A \in \mathcal{A}$ and $\alpha: \mathcal{A}(A, -) \Rightarrow F$ a \mathcal{V} -natural transformation. Consider the map

$$\phi(A): I \xrightarrow{\text{id}_A} \mathcal{A}(A, -) \xrightarrow{\alpha_A} FA.$$

The assignment

$$\begin{aligned} \phi: \mathcal{V}\text{-Nat}(\mathcal{A}(A, -), F) &\rightarrow \mathcal{V}_0(I, FA) \\ A &\mapsto \phi(A) \end{aligned}$$

is a bijection, where the inverse map ψ is given by sending $\eta: I \rightarrow FA$ to the \mathcal{V} -natural transformation

$$\mathcal{A}(A, B) \xrightarrow{F_{A,B}} [FA, FB] \xrightarrow{[\eta, \text{id}_{FA}]} [I, FA] \cong FA.$$

Proof. \mathcal{V} -naturality follows from the general principle previously mentioned that “all” morphisms coming from the monoidal structure are \mathcal{V} -natural. Since $F_{A,A}(\text{id}_A) = \text{id}_{FA}$, we get $\phi\psi = \text{id}$ by construction. We still have to prove that $\psi\phi = \text{id}$.

Consider the diagram

$$\begin{array}{ccccccc} \mathcal{A}(A, B) & \xrightarrow{\mathcal{A}(A, -)} & [\mathcal{A}(A, A), \mathcal{A}(A, B)] & \xrightarrow{[\text{id}_A, I]} & [I, \mathcal{A}(A, B)] & \xrightarrow{\sim} & \mathcal{A}(A, B) \\ \downarrow F & & \downarrow [I, \alpha_B] & & \downarrow [I, \alpha_B] & & \downarrow \alpha_B \\ [FA, FB] & \xrightarrow{[\alpha_A, I]} & [\mathcal{A}(A, A), FA] & \xrightarrow{[\text{id}_A, I]} & [I, FB] & \xrightarrow{\sim} & FA \end{array} \quad ,$$

$\xrightarrow{\quad [\phi(B), I] \quad}$

where the left and right squares on extremes commute by \mathcal{V} -naturality, while for the one in the middle we consider the functor $[-, -]: \mathcal{V}^{\text{op}} \otimes \mathcal{V} \rightarrow \mathcal{V}$.

The claim follows by checking that the composition is the identity. \square

Theorem 3.3.3 (Parametrized Yoneda). Let $T: \mathcal{B}^{\text{op}} \otimes \mathcal{A} \rightarrow \mathcal{V}$ be a \mathcal{V} -functor and suppose that for all $B \in \mathcal{B}$ there exists a $KB \in \mathcal{A}$ and a \mathcal{V} -natural isomorphism $\alpha_B: \mathcal{A}(KB, -) \xrightarrow{\sim} T(B, -)$. Then there is a unique way to define $K_{B,C}: \mathcal{B}(B, C) \rightarrow \mathcal{A}(KB, KC)$ in \mathcal{V} such that K is a \mathcal{V} -functor and $(\alpha_B)_A: \mathcal{A}(KB, A) \rightarrow T(B, A)$ is \mathcal{V} -natural in both variables as a \mathcal{V} -functor $\mathcal{B}^{\text{op}} \otimes \mathcal{A} \rightarrow \mathcal{V}$.

Proof. One checks that \mathcal{V} -naturality of $(\alpha_A)_B$ amounts to commutativity of the diagram

$$\begin{array}{ccc} \mathcal{B}(B, C) & \xrightarrow{K_{B,C}} & \mathcal{A}(KB, KC) \xrightarrow{T(B, -)} [T(B, KB), T(B, KC)] \\ \downarrow T(-, KC) & & \searrow \sim \downarrow [\eta_B, I] \\ [T(C, KC), T(B, KC)] & \xrightarrow{[\eta_C, I]} & [I, T(B, KC)] \end{array} \quad ,$$

$(\alpha_B)_{KC}$

where $\eta_B = \phi(\alpha_B)$ and the triangle commutes by Yoneda. Since $(\alpha_B)_{KC}$ is an isomorphism there exists a unique candidate $K_{B,C}$ and one only has to check that it works. \square

Remark 3.3.4. This is really useful as a way of constructing \mathcal{V} -functors via universal properties and representability results.

3.4 Weighted colimits and enriched presheaf categories

We want to define the \mathcal{V} -category $[\mathcal{A}^{\text{op}}, \mathcal{V}]$ of \mathcal{V} -presheaves or \mathcal{V} -functors for \mathcal{A} small. We will do this by defining a suitable \mathcal{V} -monad T such that $[\mathcal{A}^{\text{op}}, \mathcal{V}] := T\text{-Alg}$. We want our category to have at least coproducts and equalizers, so from now on we assume that \mathcal{V} is a (co)complete, symmetric monoidal and closed. Such an object is called *cosmos*, after Bénabou cosmos.

Definition 3.4.1. Let \mathcal{C} be a \mathcal{V} -category, $(C_j)_{j \in J} \in \text{Ob}(\mathcal{C})^J$ a family of objects in \mathcal{C} . We say that a collection $j_j: C_j \rightarrow C$ exhibits C as a \mathcal{V} -coproduct of the $(C_j)_{j \in J}$ if

$$\mathcal{C}(j_j, D): \mathcal{C}(C, D) \rightarrow \mathcal{C}(C_j, D)$$

is a product diagram in \mathcal{V}_0 for all $D \in \mathcal{C}$.

Similarly, we define a \mathcal{V} -coequalizer

$$A \rightrightarrows B \rightarrow C$$

by requiring that

$$\mathcal{C}(C, D) \rightarrow \mathcal{C}(B, D) \rightrightarrows \mathcal{C}(A, D)$$

is a coequalizer in \mathcal{V}_0 .

Dualizing the definitions, we find the notions of *product* and *equalizer*.

Remark 3.4.2. If we apply $\mathcal{V}_0(A, -): \mathcal{V}_0 \rightarrow \mathbf{Set}$, we see that \mathcal{V} -coproducts and \mathcal{V} -coequalizers are in particular coproducts and coequalizers in \mathcal{V}_0 .

Example 3.4.3. For $\mathcal{V} = \mathbf{Cat}$, a \mathcal{V} -coequalizer also has a 2-dimensional universal property, that is given one $A \rightrightarrows B \rightarrow C$ and a 2-cell from B to D there is a unique 2-cell from C to D making the following diagram commute.

For enriched categories, there is an important third kind of colimit called *copower* or *tensor*.

Definition 3.4.4. Let \mathcal{C} be a \mathcal{V} -category, $V \in \mathcal{V}$, $C \in \mathcal{C}$. We say that the copower of C by V exists if the \mathcal{V} -functor $[V, \mathcal{C}(C, -)]: \mathcal{C} \rightarrow \mathcal{V}$ is representable by some object $V \odot C \in \mathcal{C}$, the copower, that is we have a \mathcal{V} -natural isomorphism $\mathcal{C}(V \otimes C, -) \Rightarrow [V, \mathcal{C}(C, -)]$.

Dualizing the definition, we find the notion of *power* or *cotensor*.

Remark 3.4.5. By parametrized Yoneda, we get a \mathcal{V} -functor

$$- \odot -: \mathcal{V} \otimes \mathcal{C} \rightarrow \mathcal{C}$$

if all copowers exist. It is associative up to coherent isomorphism, so it defines a kind of weak action of \mathcal{V} on \mathcal{C} .

Example 3.4.6. For $\mathcal{V} = \mathcal{C} = \mathbf{Cat}$, $C \times [1]$ is the copower of $C \in \mathcal{C}$ by $[1] \in \mathcal{V}$. Indeed, we have a pair of bijective correspondences inducing the one we want as follows:

$$C \times [1] \rightarrow D \quad \leftrightarrow \quad C \begin{array}{c} \xrightarrow{\quad} \\ \Downarrow \alpha \\ \xrightarrow{\quad} \end{array} C \quad \leftrightarrow \quad C \rightarrow D^{[1]}.$$

Proposition 3.4.7. A \mathcal{V} -category \mathcal{C} has \mathcal{V} -coproducts and \mathcal{V} -coequalizers if \mathcal{C}_0 has coproducts and coequalizers and these are preserved by the functor $\mathcal{C}_0(-, D): \mathcal{C}_0 \rightarrow \mathcal{V}^{\text{op}}$ for every $D \in \mathcal{C}$.

Proof. It follows from the definition. \square

Corollary 3.4.8. The \mathcal{V} -categories \mathcal{V} and \mathcal{V}^{op} have all \mathcal{V} -coproducts and \mathcal{V} -coequalizers.

Proof. We need to check that $[-, V]_0: \mathcal{C}_0 \rightarrow \mathcal{V}_0^{\text{op}}$ preserves coproducts and coequalizers, but this is just $[-, V]: \mathcal{V}_0 \rightarrow \mathcal{V}_0^{\text{op}}$ and we have $[-, V] \dashv [-, V]: \mathcal{V}_0^{\text{op}} \rightarrow \mathcal{V}_0$ because

$$\mathcal{V}_0(X, [Y, Z]) \cong \mathcal{V}_0(X \otimes Y, Z) \cong \mathcal{V}_0(Y \otimes X, Z) \cong \mathcal{V}_0(Y, [X, Z]).$$

For \mathcal{V}^{op} , we need to check that $[V, -]_0: \mathcal{V}_0 \rightarrow \mathcal{V}_0^{\text{op}}$ preserves the limits, which follows from $- \otimes V \dashv [V, -]$. \square

Proposition 3.4.9. The \mathcal{V} -category \mathcal{V} has all powers and copowers given by $[V, C]$ and $V \otimes C$ respectively.

Proof. We need \mathcal{V} -natural isomorphisms $[V, [C, D]] \cong [V \otimes C, D]$, which follows from the fact that we have a \mathcal{V} -adjunction $- \otimes C \dashv [C, -]$. Similarly, use the symmetry isomorphism to get a \mathcal{V} -natural isomorphism $[D, [V, C]] \cong [V, [D, C]]$. \square

Definition 3.4.10. A \mathcal{V} -category \mathcal{C} is \mathcal{V} -cocomplete if it has all \mathcal{V} -coequalizers, \mathcal{V} -coproducts and copowers. If it satisfies the dual conditions, then it is \mathcal{V} -complete.

Example 3.4.11. If $(\mathcal{C}_j)_{j \in J}$ is a family of \mathcal{V} -(co)complete \mathcal{V} -categories, then $\prod_{j \in J} \mathcal{C}_j$ is a \mathcal{V} -(co)complete \mathcal{V} -category.

Proposition 3.4.12. If \mathcal{C} has powers (cotensors), then \mathcal{C} is cocomplete, if and only if \mathcal{C}_0 is cocomplete and \mathcal{C} has copowers.

Proof. " \Rightarrow " : We have already seen this.

" \Leftarrow " : We need to show, that ordinary coequalizers and coproducts in \mathcal{C}_0 are automatically \mathcal{V} -coequalizers and \mathcal{V} -coproducts. We will just check the case of coequalizers and leave the other case for the reader. We know, that we have a natural bijection of sets between equalizers

$$\mathcal{C}_0(K, D) \xrightarrow{\cong} \text{Eq}(\mathcal{C}_0(B, D) \xrightarrow{f^*} \mathcal{C}_0(A, D) \xrightarrow{g^*})$$

and coequalizers

$$A \xrightarrow[f]{g} B \xrightarrow{k} K$$

in \mathcal{C}_0 . For each $E \in \mathcal{C}$ and $V \in \mathcal{V}_0$, we thus have

$$\begin{array}{ccccc} \mathcal{C}_0(K.E^V) & \longrightarrow & \mathcal{C}_0(B, E^V) & \rightrightarrows & \mathcal{C}_0(A, E^V) \\ \cong \downarrow & & \downarrow \cong & & \downarrow \cong \\ \mathcal{V}_0(V, \mathcal{C}(K, E)) & \longrightarrow & \mathcal{V}_0(V, \mathcal{C}(B, E)) & \rightrightarrows & \mathcal{V}_0(V, \mathcal{C}(A, E)) \end{array}$$

and since this holds for all $V \in \mathcal{V}_0$, this implies that $\mathcal{C}(K, E) \cong \text{Eq}(f^*, g^*)$ in \mathcal{V} , as claimed. \square

Corollary 3.4.13. \mathcal{V} is complete and cocomplete as \mathcal{V} -category.

Remark 3.4.14. Note, that the existence of powers for a strong generating set suffices.

Definition 3.4.15. We say, that a \mathcal{V} -functor $F: \mathcal{C} \rightarrow \mathcal{D}$ preserves (certain) \mathcal{V} -coequalizers or \mathcal{V} -coproducts, if $F_0: \mathcal{C}_0 \rightarrow \mathcal{D}_0$ preserves coequalizers or coproducts. $[V]$

To talk about preservation of copowers, we need a canonical comparison morphism $\bar{F}: V \odot FC \rightarrow F(V \odot C)$, which we define to be

$$\begin{array}{ccc} I & \longrightarrow & \mathcal{D}(V \odot FC, F(V \odot C)) \\ \eta \downarrow & & \uparrow \cong \\ [V, \mathcal{C}(C, V \odot C)] & \longrightarrow & [V, \mathcal{D}(FC, F(V \odot C))] \end{array}$$

where the lower horizontal morphism is $[V, F]$ and η corresponds via weak Yoneda to the \mathcal{V} -natural isomorphism $\mathcal{C}(V \odot C, -) \cong [V, \mathcal{C}(C, -)]$.

Definition 3.4.16. We say, that F preserves the copower $V \odot C$, if $\bar{F}: V \odot FC \rightarrow F(V \odot C)$ is an isomorphism in \mathcal{D}_0 .

Lemma 3.4.17. Let \mathcal{C} be a \mathcal{V} -category and $\mathcal{B} \subset \mathcal{C}$ the full subcategory, generated by those object $B \in \mathcal{C}$ such that $V \odot B$ exists (in \mathcal{C}) for all $V \in \mathcal{V}$. Then \mathcal{B} is closed in \mathcal{C} under \mathcal{V} -coequalizers and \mathcal{V} -coproducts.

Proof. Let

$$A \rightrightarrows^f_g B \xrightarrow{k} K$$

be a \mathcal{V} -coequalizer, such that $B, A \in \mathcal{B}$. We need to show, that the \mathcal{V} -coequalizer of

$$V \odot A \rightrightarrows^{V \odot f}_{V \odot g} V \odot B$$

is given by $V \odot K$. Indeed we have an induced isomorphism

$$\begin{array}{ccccc} \mathcal{C}(V \odot K, D) & \xrightarrow{eq} & \mathcal{C}(V \odot, D) & \rightrightarrows & \mathcal{C}(V \odot A, D) \\ \cong \downarrow \exists! & & \downarrow \cong & & \downarrow \cong \\ [V, \mathcal{C}(K, D)] & \xrightarrow{eq} & [V, \mathcal{C}(B, D)] & \rightrightarrows & [V, \mathcal{C}(A, D)] \end{array}$$

and the proof for coproducts is similar. \square

Theorem 3.4.18. Let \mathcal{C} be a complete \mathcal{V} -category and let $T: \mathcal{C} \rightarrow \mathcal{C}$ be a \mathcal{V} -monad. Then $T\text{-Alg}$ is complete and $U: T\text{-Alg} \rightarrow \mathcal{C}$ preserves power, \mathcal{V} -products and \mathcal{V} -equalizers. If \mathcal{C} is also cocomplete, then $T\text{-Alg}$ is cocomplete, if and only if the underlying unriched category $(T\text{-Alg})_0 \cong T_0\text{-Alg}$ is cocomplete.

Proof. We know that $T_0\text{-Alg}$ is complete, so we need to show that equalizers and products are \mathcal{V} -equalizers and \mathcal{V} -products. But hom-objects are defined as equalizers in \mathcal{V}

$$T\text{-Alg}((A, a), (B, b)) \longrightarrow \mathcal{C}(A, B) \rightrightarrows \mathcal{C}(TA, B)$$

and we thus get the claim for \mathcal{V} -equalizers and \mathcal{V} -products, since equalizers and products commute with equalizers in \mathcal{V}_0 . We will leave the claim for powers as an exercise. But once we have the powers, we get from the cocompleteness of $T_0\text{-Alg}$, that $T\text{-Alg}$ has \mathcal{V} -coequalizers and \mathcal{V} -coproducts. So it remains to show, that $T\text{-Alg}$ has copowers. For this we use the lemma above. Since every object is a coequalizer of free algebras, hence a \mathcal{V} -coequalizer, it suffices to check this for free algebras, i.e. algebras is the image of the left \mathcal{V} -adjoint $F: \mathcal{C} \rightarrow T\text{-Alg}$. So we are done if we can show, that left \mathcal{V} -adjoints preserve copowers. This follows from the next proposition. \square

Proposition 3.4.19. Left \mathcal{V} -adjoint preserve \mathcal{V} -coequalizers, \mathcal{V} -coproducts and copowers.

Proof. Let $F: \mathcal{C} \rightarrow \mathcal{D}$ be a left \mathcal{V} -adjoint $U \vdash F$. The claims all follow as in the unenriched case. For copowers we have the isomorphisms

$$\mathcal{D}(F(V \odot C), D) \cong \mathcal{C}(V \odot C, UD) \cong [V, \mathcal{C}(C, UD)] \cong [V, \mathcal{D}(FC, D)]$$

and one checks, that this is the counit morphism if the target has copowers. \square

We are now ready to define enriched presheaf categories. Let \mathcal{A} be a small \mathcal{V} -category. Then $\prod_{A \in \text{Ob}(\mathcal{A})} \mathcal{V}$ is clearly a complete and cocomplete \mathcal{V} -category with everything computed point wise. We define the \mathcal{V} -monad for presheafs by

$$T((FA)_{A \in \text{Ob}(\mathcal{A})}) = (\coprod_{A \in \text{Ob}(\mathcal{A})} \mathcal{A}(B, A) \odot FA)_{B \in \text{Ob}(\mathcal{A})}$$

with unit given by identities and multiplication given by composition. A T -algebra is thus a collection $(FA) \in \prod_{A \in \text{Ob}(\mathcal{A})} \mathcal{V}$ with action $\coprod \mathcal{A}(B, A) \odot FA \rightarrow FB$, which amounts precisely to a \mathcal{V} -functor $\mathcal{A}(B, A) \rightarrow [FA, FB]$ i.e. a \mathcal{V} -functor $\mathcal{A}^{\text{op}} \rightarrow \mathcal{V}$.

Definition 3.4.20. We write $[\mathcal{A}^{\text{op}}, \mathcal{V}]$ for $T\text{-Alg}$ and call it the \mathcal{V} -category of \mathcal{V} -presheafs on \mathcal{A} . By construction we have $[\mathcal{A}^{\text{op}}, \mathcal{V}]_0 = \mathcal{V}\text{-CAT}(\mathcal{A}^{\text{op}}, \mathcal{V})$.

Remark 3.4.21. :

- (1) The same construction works for any cocomplete \mathcal{V} -category \mathcal{C} and we get \mathcal{V} -categories $[\mathcal{A}^{\text{op}}, \mathcal{C}]$ and $[\mathcal{A}, \mathcal{C}]$.
- (2) The statement " T is a \mathcal{V} -monad" actually needs to be checked. It can be done using Kelly (1.7,1.8) and the universal properties of \coprod and \odot (see also later exercise).
- (3) We have as enriched Yoneda lemma basically by definition: Taking the free algebra of the collection $(I_A)_{B \in \text{Ob}(\mathcal{A})}$, given by I if $B = A$ and \emptyset else, is precisely $\mathcal{A}(-, A)$. So we get isomorphisms

$$[\mathcal{A}^{\text{op}}, \mathcal{V}](\mathcal{A}(-, A), F) \cong T\text{-Alg}((F(I_A)), F) \cong \prod \mathcal{V}(I_A, (FB)_{B \in \mathcal{A}}) \cong FA$$

- (4) The hom-object is by definition the equalizer

$$[\mathcal{A}^{\text{op}}, \mathcal{V}](F, G) \longrightarrow \prod_A [FA, GA] \rightrightarrows \prod_{A, B} [\mathcal{A}(A, B) \odot FA, GB]$$

Proposition 3.4.22. The \mathcal{V} -categories $[\mathcal{A}^{\text{op}}, \mathcal{C}]$ and $[\mathcal{A}, \mathcal{C}]$ are complete (resp. cocomplete), if \mathcal{A} is small and \mathcal{C} is complete (resp. cocomplete).

Proof. This follows, since T_0 is cocontinuous. \square

Definition 3.4.23. Given a \mathcal{V} -category \mathcal{C} and a \mathcal{V} -functor $K: \mathcal{A} \rightarrow \mathcal{C}$, where \mathcal{A} is small, we have a natural T -action on the \mathcal{V} -functor $\mathcal{C} \rightarrow \prod_{\text{Ob}(\mathcal{A})} \mathcal{V}$ given by the assignment $C \mapsto \mathcal{C}(KA, C)$. Now we write

$$\text{Hom}_{\mathcal{A}}(K, -): \mathcal{C} \rightarrow [\mathcal{A}^{\text{op}}, \mathcal{V}]$$

for the induced \mathcal{V} functor, given by sending C to $\mathcal{C}(K-, C)$. This is also written as \tilde{K} .

Definition 3.4.24. Given a \mathcal{V} -presheaf $W: \mathcal{A}^{\text{op}} \rightarrow \mathcal{V}$ and a \mathcal{V} -functor $K: \mathcal{A} \rightarrow \mathcal{C}$, we say that the W -weighted colimit of K exists if

$$[\mathcal{A}^{\text{op}}, \mathcal{V}](W, \text{Hom}_{\mathcal{A}}(K, -))$$

is representable. A representing object is denoted by $W \odot_{\mathcal{A}} K$, so we have

$$\mathcal{C}(W \odot_{\mathcal{A}} K, C) \cong [\mathcal{A}^{\text{op}}, \mathcal{V}](W, \text{Hom}_{\mathcal{A}}(K, C))$$