

Mini Lecture 14 (addendum) Milk!!!!

Randal Burns
Parallel Computing for Data Science

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Synchronization

A look inside the critical section

Two common goals for synchronization

- Contention:
 - How to resolve the conflicts that result from multiple processes trying to access shared resources?
- Cooperation:
 - An action by one process may enable another action by another process
 - In such cases, processes should coordinate their actions



Why is synchronization hard?

- Design an algorithm for purchasing milk between two roommates
- Steps:
 - Arrive home
 - Look in fridge for milk
 - Leave for grocery
 - Buy milk
 - Arrive home with purchased milk

Mini-lecture 14(a): Milk!!! Parallel Computer for Data Science (10/10/2022)

Alice

- Arrive home
- Look in fridge for milk
- Leave for grocery
- Buy milk
- Arrive home with purchased milk



Bob

- Arrive home
- Look in fridge for milk
- Leave for grocery
- Buy milk
- Arrive home with purchased milk





Why is synchronization hard?

- Design an algorithm for purchasing milk between two roommates
- Steps:
 - Look in fridge for milk
 - Leave for grocery
 - Buy milk
 - Arrive home with purchased milk
- Too much milk!
- Problem is impossible without communication between parties



Let's Try Using Notes

• Algorithm #1: If you find that there is no milk in fridge, leave a note on the door, go to store and purchase milk, on return home remove note

```
if (no note) then
  if (no milk) then
  leave note
  buy milk
  remove note
  fi
fi
```



They can't see each other!!

Alice

```
if (no note) then
  if (no milk) then
   leave note
  buy milk
  remove note
  fi
fi
```

Bob

```
if (no note) then
  if (no milk) then
  leave note
  buy milk
  remove note
  fi
fi
```





Let's Try Using Notes

• Algorithm #2: Based on leaving a note (with one's name) before checking fridge

```
leave note A
if (no note B) then
  if (no milk) then
    buy milk
  fi
fi
remove note
```



Alice

```
leave note A
if (no note B) then
  if (no milk) then
    buy milk
  fi
fi
remove note
```

Bob

```
leave note B
if (no note A) then
  if (no milk) then
   buy milk
  fi
fi
remove note
```



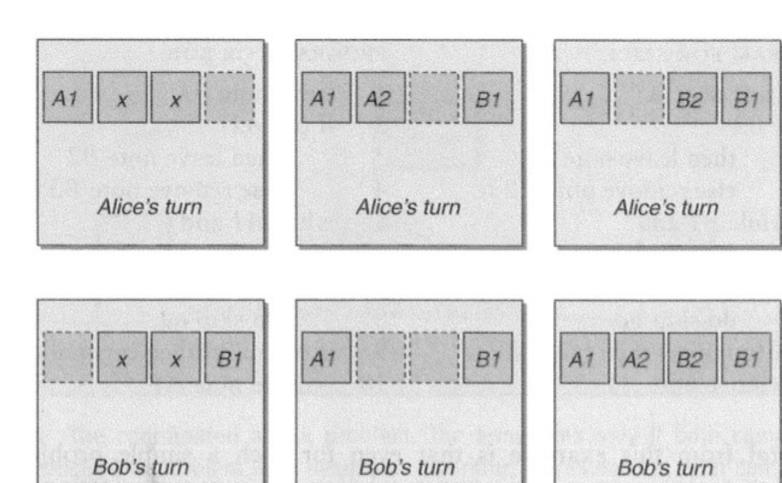


A Correct Algorithm

```
leave note A1
                            leave note B1
if (B2)
                            if (no A2)
  then leave note A2
                              then leave note B2
                              else remove note B2 fi
  else remove note A2 fi
while B1 and
                            while A1 and
  ((A2 and B2) or
                              ((A2 and no B2) or
                               (no A2 and B2))
   (no A2 and no B2))
  do skip do
                              do skip do
if (no milk)
                            if (no milk)
  then buy milk fi
                              then buy milk fi
                            remove note B1
remove note A1
```



Configurations





Two Notes

- First one to identify contention
 - Are two parties vying for this resource
- Second one to break ties during contention
 - Essentially even and odd configurations
- These notes are the analogies of atomic shared registers in computing
 - Essentially a volatile variable of basic type



Some Properties

- Correct
- Asynchronous: doesn't depend on timing
- Symmetric: equal chance of A/B buying milk
 - Notably steps aren't symmetric
- Two parties

• Even simple synchronization is hard and subtle