${\bf IUT\ Brogrammers}, {\bf IUT}$

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1 Al Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,sse3,sse4,popcnt,abm
    ,mmx,avx,tune=native")
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
    //find by order(k) --> returns iterator to the kth
        largest element counting from 0
    //order_of_key(val) --> returns the number of items
        in a set that are strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type, less<DT>,
    rb_tree_tag,tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef SFT
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\'' << x << '\'';}</pre>
void show(const string &x) {cerr << '\"' << x << '\"';}</pre>
void show(bool x) {cerr << (x ? "true" : "false");}</pre>
template<typename T, typename V>
void show(pair<T, V> x) { cerr << ','; show(x.first);</pre>
    cerr << ", "; show(x.second); cerr << '}'; }</pre>
template<typename T>
void show(T x) {int f = 0; cerr << "{"; for (auto &i: x)</pre>
    cerr << (f++ ? ", " : ""), show(i); cerr << "}";}</pre>
void debug_out(string s) {
    cerr << '\n';
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
    s.erase(remove(s.begin(), s.end(), ''), s.end());
    cerr << "
                   "; // 8 spaces
    cerr << s.substr(0, s.find(','));</pre>
    s = s.substr(s.find(',') + 1);
    cerr << " = ";
    show(t):
    cerr << endl;</pre>
    if(sizeof...(v)) debug_out(s, v...);
}
```

```
debug_out(#x, x); cerr << endl;</pre>
#else
#define debug(x...)
#endif
const int RANDOM = chrono::high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) \hat{x};
unsigned hash_combine(unsigned a, unsigned b) { return a
    * 31 + b: }
struct chash {
 int operator()(int x) const { return hash_f(x); }
typedef gp_hash_table<int, int, chash> gp;
gp table;
```

Data Structure

2.1 Segment Tree

```
const int N = 1000006;
using DT = LL;
using LT = LL;
constexpr DT I = 0;
constexpr LT None = 0;
DT val[4 * N]:
LT lazy[4 * N];
int L, R;
void pull(int s, int e, int node) {
  val[node] = val[node << 1] + val[node << 1 | 1];</pre>
void apply(const LT &U, int s, int e, int node) {
 val[node] += (e - s + 1) * U:
 lazv[node] += U;
void reset(int node) { lazy[node] = None; }
DT merge(const DT &a, const DT &b) { return a + b; }
DT get(int s, int e, int node) { return val[node]; }
void push(int s, int e, int node) {
 if (s == e) return;
  apply(lazy[node], s, s + e \gg 1, node \ll 1);
  apply(lazy[node], s + e + 2 >> 1, e, node << 1 | 1);
 reset(node):
```

```
#define debug(x...) cerr << "LINE: " << LINE << endl: void build(int s. int e. vector<DT> &v. int node = 1) {
                                                            int m = s + e >> 1:
                                                            if (s == e) {
                                                              val[node] = v[s];
                                                              return;
                                                            build(s, m, v, node * 2);
                                                            build(m + 1, e, v, node * 2 + 1);
                                                            pull(s, e, node);
                                                           void update(int S, int E, LT uval, int s = L, int e = R,
                                                               int node = 1) {
                                                            if (S > E) return:
                                                            if (S == s \text{ and } E == e) {
                                                              apply(uval, s, e, node);
                                                              return:
                                                            }
                                                            push(s, e, node);
                                                            int m = s + e >> 1:
                                                            update(S, min(m, E), uval, s, m, node * 2);
                                                            update(max(S, m + 1), E, uval, m + 1, e, node * 2 + 1);
                                                            pull(s, e, node);
                                                          DT query(int S, int E, int s = L, int e = R, int node =
                                                              1) {
                                                            if (S > E) return I;
                                                            if (s == S and e == E) return get(s, e, node);
                                                            push(s, e, node);
                                                            int m = s + e >> 1:
                                                            DT L = query(S, min(m, E), s, m, node * 2);
                                                            DT R = query(\max(S, m + 1), E, m + 1, e, node * 2 + 1);
                                                            return merge(L, R);
                                                           void init(int _L, int _R, vector<DT> &v) {
                                                            L = _L, R = _R;
                                                            build(L, R, v);
```

2.2 Persistent Segment Tree

```
struct Node {
Node *1, *r;
 int sum;
 Node(int val) : l(nullptr), r(nullptr), sum(val) {}
 Node(Node* 1, Node* r) : 1(1), r(r), sum(0) {
  if (1) sum += 1->sum;
  if (r) sum += r->sum:
}
```

```
int a[MAXN]:
Node* root[MAXN];
Node* Build(int bg, int ed) {
 if (bg == ed) return new Node(a[bg]);
 int mid = (bg + ed) / 2;
 return new Node(Build(bg, mid), Build(mid + 1, ed));
int Query(Node* v, int bg, int ed, int 1, int r) {
 if (1 > ed || r < bg) return 0;</pre>
 if (1 <= bg && ed <= r) return v->sum;
 int mid = (bg + ed) / 2;
 return Query(v->1, bg, mid, 1, r) + Query(v->r, mid +
      1, ed, l, r);
Node* Update(Node* v, int bg, int ed, int pos, int
    new val) {
  if (bg == ed) return new Node(v->sum + new_val);
 int mid = (bg + ed) / 2;
  if (pos <= mid)</pre>
   return new Node(Update(v->1, bg, mid, pos, new_val),
        v->r):
  else
   return new Node(v->1, Update(v->r, mid + 1, ed, pos,
        new_val));
```

2.3 Implicit Segment Tree

```
struct node {
  int val;
  node *lft, *rt;
 node() {}
  node(int val = 0) : val(val), lft(NULL), rt(NULL) {}
};
struct implicit_segtree {
  node *root:
  implicit_segtree() {}
  implicit_segtree(int n) {
   root = new node(n);
  void update(node *now, int L, int R, int idx, int val)
      {
   if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2:
```

```
if (now->lft == NULL) now->lft = new node(mid - L +
        1);
   if (now->rt == NULL) now->rt = new node(R - mid);
   if (idx <= mid) update(now->lft, L, mid, idx, val);
   else update(now->rt, mid + 1, R, idx, val);
   now->val = (now->lft)->val + (now->rt)->val;
 int query(node *now, int L, int R, int k) {
   if (L == R) return L;
   int mid = L + (R - L) / 2:
   if (now->lft == NULL) now->lft = new node(mid - L +
   if (now->rt == NULL) now->rt = new node(R - mid):
   if (k <= (now->lft)->val) return query(now->lft, L,
   else return query(now->rt, mid + 1, R, k - (now->lft) |}
        ->val);
 }
};
```

2.4 DSU With Rollbacks

```
struct Rollback DSU {
 int n;
 vector<int> par, sz;
 vector<pair<int, int>> op;
 Rollback_DSU(int n) : par(n), sz(n, 1) {
   iota(par.begin(), par.end(), 0);
   op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par[node])
     ; // no path compression
   return node;
 void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)]) swap(x, y);
   op.emplace_back(x, y);
   par[y] = x;
   sz[x] += sz[y];
 void Undo(int t) {
   for (; op.size() > t; op.pop_back()) {
     par[op.back().second] = op.back().second;
     sz[op.back().first] -= sz[op.back().second];
   }
 }
```

2.5 BIT-2D

```
const int N = 1008:
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
 for (; x < N; x += -x & x)
   for (int j = y; j < N; j += -j & j) bit[x][j] += val;
int get(int x, int y) {
 int ans = 0;
 for (; x; x -= x & -x)
   for (int j = y; j; j -= j & -j) ans += bit[x][j];
 return ans;
int get(int x1, int y1, int x2, int y2) {
 return get(x2, y2) - get(x1 - 1, y2) - get(x2, y1 - 1)
      + get(x1 - 1, y1 - 1);
```

2.6 Merge Sort Tree

if (bg == ed) {

Tree[node].push_back(arr[bg]);

```
vector<LL> Tree[4 * MAXN];
LL arr[MAXN];
vector<LL> merge(vector<LL> v1, vector<LL> v2) {
 LL i = 0, j = 0;
 vector<LL> ret;
 while (i < v1.size() || j < v2.size()) {</pre>
   if (i == v1.size()) {
     ret.push_back(v2[j]);
     j++;
   } else if (j == v2.size()) {
     ret.push_back(v1[i]);
     i++;
   } else {
     if (v1[i] < v2[i]) {</pre>
       ret.push_back(v1[i]);
       i++;
     } else {
       ret.push_back(v2[j]);
       j++;
   }
 }
 return ret;
void Build(LL node, LL bg, LL ed) {
```

```
3
```

```
return:
 LL leftNode = 2 * node, rightNode = 2 * node + 1;
 LL mid = (bg + ed) / 2;
  Build(leftNode, bg, mid);
  Build(rightNode, mid + 1, ed);
 Tree[node] = merge(Tree[leftNode], Tree[rightNode]);
LL query(LL node, LL bg, LL ed, LL 1, LL r, LL k) {
 if (ed < 1 || bg > r) return 0;
 if (1 <= bg && ed <= r)</pre>
   return upper_bound(Tree[node].begin(), Tree[node].end |}
        (), k) -
         Tree[node].begin();
 LL leftNode = 2 * node, rightNode = 2 * node + 1;
 LL mid = (bg + ed) / 2;
 return query(leftNode, bg, mid, l, r, k) +
        query(rightNode, mid + 1, ed, l, r, k);
```

2.7 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs = 25;
int arr[N], freq[2 * N], cnt[2 * N], id[N], ans[N];
struct query {
  int 1, r, t, L, R;
  query(int l = 1, int r = 0, int t = 1, int id = -1)
     : 1(1), r(r), t(t), L(1 / sz), R(r / sz) {}
  bool operator<(const query &rhs) const {</pre>
    return (L < rhs.L) or (L == rhs.L and R < rhs.R) or
          (L == rhs.L and R == rhs.R and t < rhs.t);
 }
} Q[N]:
struct update {
 int idx, val, last;
} Up[N];
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]++;
  cnt[freq[arr[idx]]]++;
```

```
void remove(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]--;
 cnt[freq[arr[idx]]]++;
void apply(int t) {
 const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].val;
 if (f) add(Up[t].idx);
void undo(int t) {
 const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].last;
 if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
 assert(0):
int main() {
 int n, q;
 cin >> n >> q;
 int counter = 0;
 map<int, int> M;
 for (int i = 1; i <= n; i++) {</pre>
   cin >> arr[i]:
   if (!M[arr[i]]) M[arr[i]] = ++counter;
   arr[i] = M[arr[i]];
 iota(id, id + N, 0);
 while (q--) {
   int tp, x, y;
   cin >> tp >> x >> y;
   if (tp == 1)
     Q[++qi] = query(x, y, ui);
   else {
     if (!M[y]) M[y] = ++counter;
     v = M[v];
     Up[++ui] = \{x, y, arr[x]\};
     arr[x] = v;
   }
 t = ui;
 cnt[0] = 3 * n;
 sort(id + 1, id + qi + 1, [&](int x, int y) { return Q[
      x] < Q[v]; \});
 for (int i = 1; i <= qi; i++) {
```

```
int x = id[i];
while (Q[x].t > t) apply(++t);
while (Q[x].t < t) undo(t--);
while (Q[x].1 < 1) add(--1);
while (Q[x].r > r) add(++r);
while (Q[x].1 > 1) remove(1++);
while (Q[x].r < r) remove(r--);
ans[x] = mex();
}
for (int i = 1; i <= qi; i++) cout << ans[i] << '\n';
}</pre>
```

2.8 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9:
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN][MAXN][LOGN][LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
   for (int j = 1; j <= N; j++) {</pre>
     M[i][i][0][0] = A[i][i];
   for (int q = 1; (1 << q) <= N; q++) {
     int add = 1 << (q - 1);</pre>
     for (int j = 1; j + add \le N; j++) {
       M[i][j][0][q] = max(M[i][j][0][q - 1], M[i][j +
            add][0][q - 1]);
     }
   }
 }
 for (int p = 1; (1 << p) <= N; p++) {
   int add = 1 << (p - 1);
   for (int i = 1; i + add <= N; i++) {</pre>
     for (int q = 0; (1 << q) <= N; q++) {
       for (int j = 1; j <= N; j++) {</pre>
         M[i][j][p][q] = max(M[i][j][p - 1][q], M[i + add
             ][i][p - 1][q]);
 }
```

```
// returns max of all A[i][j], where x1<=i<=x2 and y1<=j | ll dp[MAX];
int Query(int x1, int y1, int x2, int y2) {
 int kX = log2(x2 - x1 + 1);
 int kY = log2(y2 - y1 + 1);
 int addX = 1 << kX;
 int addY = 1 \ll kY:
 int ret1 = max(M[x1][y1][kX][kY], M[x1][y2 - addY + 1][
      kX][kY]);
 int ret2 = max(M[x2 - addX + 1][v1][kX][kY],
               M[x2 - addX + 1][y2 - addY + 1][kX][kY]);
 return max(ret1, ret2);
```

3 DP

3.1 Convex Hull Trick

```
struct line {
 11 m, c;
 line() {}
  line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
  vector<line> lines;
  int ptr = 0;
  convex_hull_trick() {}
  bool bad(line a, line b, line c) {
   return 1.0 * (c.c - a.c) * (a.m - b.m) < 1.0 * (b.c - #include <bits/stdc++.h>
         a.c) * (a.m - c.m);
  void add(line L) {
    int sz = lines.size();
    while (sz >= 2 && bad(lines[sz - 2], lines[sz - 1], L
       )) {
     lines.pop_back();
     sz--;
   lines.pb(L);
  ll get(int idx, int x) { return (1ll * lines[idx].m * x
       + lines[idx].c); }
  11 query(int x) {
    if (lines.empty()) return 0;
    if (ptr >= lines.size()) ptr = lines.size() - 1;
    while (ptr < lines.size() - 1 && get(ptr, x) > get(
        ptr + 1, x)) ptr++;
    return get(ptr, x);
};
```

```
ll sum[MAX]:
int arr[MAX];
int main() {
 fastio:
 int t:
 cin >> t;
 while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >> sum[i];
   for (int i = 1; i <= n; i++) dp[i] = 0, sum[i] += sum</pre>
        [i - 1]:
   convex hull trick cht:
   cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 111 * a * sqr(sum[
         pos]) - c:
     cht.add(line(211 * a * sum[pos], dp[pos] - a * sqr(
          sum[pos])));
   }
   ll ans = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
```

3.2 Dynamic CHT

```
using namespace std;
typedef long long 11;
const ll IS_QUERY = -(1LL << 62);
struct line {
 ll m, b;
  mutable function <const line*()> succ:
  bool operator < (const line &rhs) const {</pre>
    if (rhs.b != IS_QUERY) return m < rhs.m;</pre>
    const line *s = succ();
   if (!s) return 0:
   11 x = rhs.m;
   return b - s -> b < (s -> m - m) * x;
 }
};
struct HullDynamic : public multiset <line> {
```

```
bool bad (iterator y) {
   auto z = next(y);
   if (y == begin()) {
    if (z == end()) return 0;
     return y -> m == z -> m && y -> b <= z -> b;
   auto x = prev(y);
   if (z == end()) return y -> m == x -> m && y -> b <=
        x \rightarrow b:
   return 1.0 * (x \rightarrow b - y \rightarrow b) * (z \rightarrow m - y \rightarrow m) >=
         1.0 * (y \rightarrow b - z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
 void insert line (ll m. ll b) {
   auto y = insert({m, b});
   y \rightarrow succ = [=] {return next(y) == end() ? 0 : &*next}
        (y);};
   if (bad(y)) {erase(y); return;}
   while (next(y) != end() && bad(next(y))) erase(next(y))
   while (y != begin() && bad(prev(y))) erase(prev(y));
 ll eval (ll x) {
   auto 1 = *lower_bound((line) {x, IS_QUERY});
   return 1.m * x + 1.b;
int main() {
 HullDynamic hull;
 hull.insert_line(1, 1);
 hull.insert_line(-1, 1);
 cout << hull.eval(69) << endl;</pre>
 cout << hull.eval(420) << endl;</pre>
 return 0;
```

3.3 Li Chao Tree

```
struct line {
 LL m, c;
 line(LL m = 0, LL c = 0) : m(m), c(c) {}
LL calc(line L, LL x) { return 1LL * L.m * x + L.c; }
struct node {
 LL m, c;
 line L:
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft = NULL, node *rt =
     NULL)
```

```
: L(line(m, c)), lft(lft), rt(rt) {}
};
struct LiChao {
  node *root:
  LiChao() { root = new node(); }
  void update(node *now, int L, int R, line newline) {
    int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi:
     return;
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new node();
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo:
     if (now->lft == NULL) now->lft = new node();
     update(now->lft, L, mid, hi);
   }
  LL query(node *now, int L, int R, LL x) {
   if (now == NULL) return -inf;
   int mid = L + (R - L) / 2;
   if (x \le mid)
     return max(calc(now->L, x), query(now->lft, L, mid,
           x)):
    else
     return max(calc(now->L, x), query(now->rt, mid + 1,
           R, x));
 }
};
```

4 Geometry

4.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) \{ return abs(x) < EPS ? 0 : (x < 0 ? -1 : 
    1); }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
 Point operator+(const Point& u) const { return Point(x
      + u.x, y + u.y); }
```

```
Point operator-(const Point& u) const { return Point(x | Point normal(Point a) {
      -u.x, y - u.y); }
  Point operator*(const LL u) const { return Point(x * u,
       v * u); }
  Point operator*(const Tf u) const { return Point(x * u, // returns 1 if c is left of ab, 0 if on ab && -1 if
       y / u); }
  bool operator==(const Point& u) const {
   return dcmp(x - u.x) == 0 && dcmp(y - u.y) == 0;
  bool operator!=(const Point& u) const { return !(*this
      == u): }
  bool operator<(const Point& u) const {</pre>
   return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x) == 0 &&
        dcmp(y - u.y) < 0);
 }
Ti dot(Point a, Point b) { return a.x * b.x + a.y * b.y;
Ti cross(Point a, Point b) { return a.x * b.y - a.y * b.x
Tf length(Point a) { return sqrt(dot(a, a)); }
Ti sqLength(Point a) { return dot(a, a); }
Tf distance(Point a, Point b) { return length(a - b); }
Tf angle(Point u) { return atan2(u.y, u.x); }
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Point a, Point b) {
 Tf ans = angle(b) - angle(a);
  return ans <= -PI ? ans + 2 * PI : (ans > PI ? ans - 2
      * PI : ans);
// Rotate a ccw by rad radians, Tf Ti same
Point rotate(Point a, Tf rad) {
 return Point(a.x * cos(rad) - a.y * sin(rad),
             a.x * sin(rad) + a.y * cos(rad));
// rotate a ccw by angle th with cos(th) = co && sin(th)
    = si, tf ti same
Point rotatePrecise(Point a, Tf co, Tf si) {
 return Point(a.x * co - a.y * si, a.y * co + a.x * si);
Point rotate90(Point a) { return Point(-a.y, a.x); }
// scales vector a by s such that length of a becomes s,
    Tf Ti same
Point scale(Point a, Tf s) { return a / length(a) * s; }
// returns an unit vector perpendicular to vector a, Tf
    Ti same
```

```
Tf l = length(a);
                                                          return Point(-a.y / 1, a.x / 1);
                                                             right of ab
Point operator/(const Tf u) const { return Point(x / u, | int orient(Point a, Point b, Point c) { return dcmp(cross
                                                             (b - a, c - a)); }
                                                         /// Use as sort(v.begin(), v.end(), polarComp(0, dir))
                                                         /// Polar comparator around O starting at direction dir
                                                         struct polarComp {
                                                          Point O, dir;
                                                           polarComp(Point 0 = Point(0, 0), Point dir = Point(1,
                                                               0)) : 0(0), dir(dir) {}
                                                           bool half(Point p) {
                                                            return dcmp(cross(dir, p)) < 0 ||</pre>
                                                                   (dcmp(cross(dir, p)) == 0 && dcmp(dot(dir, p))
                                                                       > 0);
                                                           bool operator()(Point p, Point q) {
                                                            return make_tuple(half(p), 0) < make_tuple(half(q),</pre>
                                                                 cross(p, q));
                                                         struct Segment {
                                                          Point a, b;
                                                           Segment(Point aa, Point bb) : a(aa), b(bb) {}
                                                         typedef Segment Line;
                                                         struct Circle {
                                                          Point o;
                                                           Tf r:
                                                           Circle(Point o = Point(0, 0), Tf r = 0) : o(o), r(r) {}
                                                           // returns true if point p is in || on the circle
                                                           bool contains(Point p) { return dcmp(sqLength(p - o) -
                                                               r * r) <= 0; }
                                                          // returns a point on the circle rad radians away from
                                                               +X CCW
                                                           Point point(Tf rad) {
                                                             static_assert(is_same<Tf, Ti>::value);
                                                            return Point(o.x + cos(rad) * r, o.y + sin(rad) * r);
                                                          // area of a circular sector with central angle rad
                                                           Tf area(Tf rad = PI + PI) { return rad * r * r / 2; }
                                                           // area of the circular sector cut by a chord with
                                                               central angle alpha
                                                           Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - sin )
                                                               (alpha)); }
```

4.2 Linear

```
// **** LINE LINE INTERSECTION START ****
// returns true if point p is on segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p)) == 0 && dcmp(dot(s
      .a - p, s.b - p)) <= 0;
// returns true if segment p && q touch or intersect
bool segmentsIntersect(Segment p, Segment q) {
 if (onSegment(p.a, q) || onSegment(p.b, q)) return true
  if (onSegment(q.a, p) || onSegment(q.b, p)) return true
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) <
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.a)) == 0;
// lines are represented as a ray from a point: (point,
// returns false if two lines (p, v) && (q, w) are
    parallel or collinear
// true otherwise, intersection point is stored at o via
    reference, Tf Ti Same
bool lineLineIntersection(Point p, Point v, Point q,
    Point w, Point& o) {
 if (dcmp(cross(v, w)) == 0) return false;
  Point u = p - q;
  o = p + v * (cross(w, u) / cross(v, w));
 return true;
// returns false if two lines p && q are parallel or
// true otherwise, intersection point is stored at o via
    reference
bool lineLineIntersection(Line p, Line q, Point& o) {
 return lineLineIntersection(p.a, p.b - p.a, q.a, q.b -
      q.a, o);
// returns the distance from point a to line 1
// **** LINE LINE INTERSECTION FINISH ****
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(1.b - 1.a, p - 1.a) / length(1.b - 1.a | Tf circleTriangleIntersectionArea(Circle c, Segment s) {
     ));
```

```
// returns the shortest distance from point a to segment
Tf distancePointSegment(Point p, Segment s) {
 if (s.a == s.b) return length(p - s.a);
 Point v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
 if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
  else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
 else
   return abs(cross(v1, v2) / length(v1));
// returns the shortest distance from segment p to
    segment q
Tf distanceSegmentSegment(Segment p, Segment q) {
 if (segmentsIntersect(p, q)) return 0;
 Tf ans = distancePointSegment(p.a, q);
  ans = min(ans, distancePointSegment(p.b, q));
  ans = min(ans, distancePointSegment(q.a, p));
  ans = min(ans, distancePointSegment(q.b, p));
 return ans;
// returns the projection of point p on line 1, Tf Ti
Point projectPointLine(Point p, Line 1) {
 Point v = 1.b - 1.a:
 return 1.a + v * ((Tf)dot(v, p - 1.a) / dot(v, v));
```

4.3 Circular

```
// Extremely inaccurate for finding near touches
// compute intersection of line 1 with circle c
// The intersections are given in order of the ray (1.a,
    1.b), Tf Ti same
vector<Point> circleLineIntersection(Circle c, Line 1) {
 vector<Point> ret;
 Point b = 1.b - 1.a, a = 1.a - c.o;
 Tf A = dot(b, b), B = dot(a, b):
 Tf C = dot(a, a) - c.r * c.r, D = B * B - A * C;
 if (D < -EPS) return ret;</pre>
 ret.push_back(1.a + b * (-B - sqrt(D + EPS)) / A);
 if (D > EPS) ret.push_back(1.a + b * (-B + sqrt(D)) / A
      );
 return ret;
// signed area of intersection of circle(c.o, c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)/2]
 using Linear::distancePointSegment;
```

```
Tf OA = length(c.o - s.a);
 Tf OB = length(c.o - s.b);
 // sector
 if (dcmp(distancePointSegment(c.o, s) - c.r) >= 0)
   return angleBetween(s.a - c.o, s.b - c.o) * (c.r * c.
       r) / 2.0:
 // triangle
 if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r) \le 0)
   return cross(c.o - s.b, s.a - s.b) / 2.0;
 // three part: (A, a) (a, b) (b, B)
 vector<Point> Sect = circleLineIntersection(c, s);
 return circleTriangleIntersectionArea(c, Segment(s.a,
      Sect[0])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [0], Sect[1])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [1], s.b));
// area of intersecion of circle(c.o, c.r) && simple
    polyson(p[])
Tf circlePolyIntersectionArea(Circle c, Polygon p) {
 Tf res = 0;
 int n = p.size();
 for (int i = 0; i < n; ++i)
   res += circleTriangleIntersectionArea(c, Segment(p[i
       ], p[(i + 1) % n]));
 return abs(res):
// locates circle c2 relative to c1
// interior
                     (d < R - r)
                                        ---> -2
// interior tangents (d = R - r)
// concentric
                  (0 = b)
// secants
                     (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r)
// exterior
                     (d > R + r)
int circleCirclePosition(Circle c1, Circle c2) {
 Tf d = length(c1.o - c2.o):
 int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (c1.r))
     r + c2.r):
 return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex >
     0 ? 2 : 0;
// compute the intersection points between two circles c1
     && c2, Tf Ti same
vector<Point> circleCircleIntersection(Circle c1, Circle
    c2) {
 vector<Point> ret;
 Tf d = length(c1.o - c2.o);
 if (dcmp(d) == 0) return ret;
 if (dcmp(c1.r + c2.r - d) < 0) return ret;
```

```
if (dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
 Point v = c2.o - c1.o;
 Tf co = (c1.r * c1.r + sqLength(v) - c2.r * c2.r) / (2)
      * c1.r * length(v));
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co, -si), c1.r) + c1.
 Point p2 = scale(rotatePrecise(v, co, si), c1.r) + c1.o
 ret.push_back(p1);
 if (p1 != p2) ret.push_back(p2);
 return ret;
// intersection area between two circles c1, c2
Tf circleCircleIntersectionArea(Circle c1, Circle c2) {
 Point AB = c2.0 - c1.0;
 Tf d = length(AB);
 if (d \ge c1.r + c2.r) return 0;
 if (d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
 if (d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
 Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r) /
      (2.0 * c1.r * d));
 Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r) /
      (2.0 * c2.r * d));
 return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
// returns tangents from a point p to circle c, Tf Ti
vector<Point> pointCircleTangents(Point p, Circle c) {
 vector<Point> ret;
 Point u = c.o - p;
 Tf d = length(u);
 if (d < c.r)
  ;
 else if (dcmp(d - c.r) == 0) {
   ret = {rotate(u, PI / 2)};
 } else {
   Tf ang = asin(c.r / d);
   ret = {rotate(u, -ang), rotate(u, ang)};
 return ret;
// returns the points on tangents that touches the circle
    , Tf Ti Same
vector<Point> pointCircleTangencyPoints(Point p, Circle c
    ) {
 Point u = p - c.o;
```

```
Tf d = length(u);
 if (d < c.r)
   return {}:
  else if (dcmp(d - c.r) == 0)
   return {c.o + u};
  else {
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
   return {c.o + rotate(u, -ang), c.o + rotate(u, ang)};
// for two circles c1 && c2, returns two list of points a }
// such that a[i] is on c1 && b[i] is c2 && for every i
// Line(a[i], b[i]) is a tangent to both circles
// CAUTION: a[i] = b[i] in case they touch | -1 for c1 =
int circleCircleTangencyPoints(Circle c1, Circle c2,
    vector<Point> &a,
                            vector<Point> &b) {
  a.clear(), b.clear();
  int cnt = 0;
  if (dcmp(c1.r - c2.r) < 0) {
   swap(c1, c2);
   swap(a, b);
 Tf d2 = sqLength(c1.o - c2.o);
 Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r;
 if (dcmp(d2 - rdif * rdif) < 0) return 0;</pre>
  if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r) == 0) return -1;
  Tf base = angle(c2.o - c1.o);
  if (dcmp(d2 - rdif * rdif) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(base));
   cnt++;
   return cnt;
  Tf ang = acos((c1.r - c2.r) / sqrt(d2));
  a.push_back(c1.point(base + ang));
  b.push_back(c2.point(base + ang));
  cnt++;
  a.push_back(c1.point(base - ang));
  b.push_back(c2.point(base - ang));
  cnt++;
  if (dcmp(d2 - rsum * rsum) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(PI + base));
```

```
cnt++;
} else if (dcmp(d2 - rsum * rsum) > 0) {
   Tf ang = acos((c1.r + c2.r) / sqrt(d2));
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(PI + base + ang));
   cnt++;
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(PI + base - ang));
   cnt++;
}
return cnt;
}
```

4.4 Convex

```
/// minkowski sum of two polygons in O(n)
Polygon minkowskiSum(Polygon A, Polygon B) {
 int n = A.size(), m = B.size();
 rotate(A.begin(), min_element(A.begin(), A.end()), A.
      end());
 rotate(B.begin(), min_element(B.begin(), B.end()), B.
      end());
 A.push_back(A[0]);
 B.push_back(B[0]);
 for (int i = 0; i < n; i++) A[i] = A[i + 1] - A[i];
 for (int i = 0; i < m; i++) B[i] = B[i + 1] - B[i];
 Polygon C(n + m + 1);
 C[0] = A.back() + B.back();
 merge(A.begin(), A.end() - 1, B.begin(), B.end() - 1, C
      .begin() + 1,
       polarComp(Point(0, 0), Point(0, -1)));
 for (int i = 1; i < C.size(); i++) C[i] = C[i] + C[i -</pre>
      1];
 C.pop_back();
 return C;
// finds the rectangle with minimum area enclosing a
    convex polygon and
// the rectangle with minimum perimeter enclosing a
    convex polygon
// Tf Ti Same
pair<Tf, Tf> rotatingCalipersBoundingBox(const Polygon &p
 using Linear::distancePointLine;
 int n = p.size();
 int l = 1, r = 1, j = 1;
 Tf area = 1e100;
 Tf perimeter = 1e100;
 for (int i = 0; i < n; i++) {</pre>
```

```
Point v = (p[(i + 1) \% n] - p[i]) / length(p[(i + 1)
       % n] - p[i]);
   while (dcmp(dot(v, p[r % n] - p[i]) - dot(v, p[(r +
       1) \% n] - p[i])) < 0)
     r++;
   while (j < r \mid | dcmp(cross(v, p[j % n] - p[i]) -
                       cross(v, p[(j + 1) % n] - p[i])) <
                            0)
     j++;
   while (1 < j ||
          dcmp(dot(v, p[1 \% n] - p[i]) - dot(v, p[(1 + 1)
               % n] - p[i])) > 0)
     1++:
   Tf w = dot(v, p[r \% n] - p[i]) - dot(v, p[1 \% n] - p[
        i]);
   Tf h = distancePointLine(p[j % n], Line(p[i], p[(i +
       1) % n]));
   area = min(area, w * h);
   perimeter = min(perimeter, 2 * w + 2 * h);
 return make_pair(area, perimeter);
// returns the left side of polygon u after cutting it by
     ray a->b
Polygon cutPolygon(Polygon u, Point a, Point b) {
  using Linear::lineLineIntersection;
  using Linear::onSegment;
  Polygon ret;
  int n = u.size();
  for (int i = 0; i < n; i++) {
   Point c = u[i], d = u[(i + 1) \% n];
   if (dcmp(cross(b - a, c - a)) >= 0) ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0) {
     Point t;
     lineLineIntersection(a, b - a, c, d - c, t);
     if (onSegment(t, Segment(c, d))) ret.push_back(t);
   }
 }
 return ret;
// returns true if point p is in or on triangle abc
bool pointInTriangle(Point a, Point b, Point c, Point p)
 return dcmp(cross(b - a, p - a)) >= 0 && dcmp(cross(c -
      b, p - b)) >= 0 &&
        dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no three collinear points
// returns inside = -1, on = 0, outside = 1
```

```
int pointInConvexPolygon(const Polygon &pt, Point p) {
 int n = pt.size();
 assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p - pt[0])) > 0)
     lo = mid;
   else
     hi = mid:
 }
 bool in = pointInTriangle(pt[0], pt[lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p - pt[lo - 1])) == |}
       0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p - pt[lo])) == 0)
 if (dcmp(cross(pt[hi] - pt[(hi + 1) % n], p - pt[(hi +
      1) % n (1) == 0)
   return 0;
 return -1;
// Extreme Point for a direction is the farthest point in
     that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly, Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 + n) % n], u)) >=
           0) {
     return c;
   bool a_up = dcmp(dot(poly[(a + 1) \% n] - poly[a], u))
   bool c_{up} = dcmp(dot(poly[(c + 1) % n] - poly[c], u))
   bool a_above_c = dcmp(dot(poly[a] - poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c;
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
```

```
if (a_above_c)
       b = c;
     else
       a = c;
   } else {
     if (!a_above_c)
       b = c;
     else
       a = c;
 }
 if (dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 &&
     dcmp(dot(poly[a] - poly[(a - 1 + n) \% n], u)) > 0)
   return a;
 return b % n:
// For a convex polygon p and a line l, returns a list of
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i], p[(i + 1) modulo
// #1 If a segment is collinear with the line, only that
    is returned
// #2 Else if 1 goes through i'th point, the i'th segment
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const Polygon &p,
    Line 1) {
 assert((int)p.size() >= 3);
 assert(1.a != 1.b);
 int n = p.size();
 vector<int> ret;
 Point v = 1.b - 1.a;
 int lf = extremePoint(p, rotate90(v));
 int rt = extremePoint(p, rotate90(v) * Ti(-1));
 int olf = orient(l.a, l.b, p[lf]);
 int ort = orient(l.a, l.b, p[rt]);
 if (!olf || !ort) {
   int idx = (!olf ? lf : rt);
   if (orient(1.a, 1.b, p[(idx - 1 + n) \% n]) == 0)
     ret.push_back((idx - 1 + n) \% n);
     ret.push_back(idx);
   return ret;
 if (olf == ort) return ret;
```

```
for (int i = 0; i < 2; ++i) {</pre>
   int lo = i ? rt : lf;
   int hi = i ? lf : rt;
   int olo = i ? ort : olf;
   while (true) {
     int gap = (hi - lo + n) \% n;
     if (gap < 2) break;</pre>
     int mid = (lo + gap / 2) % n;
     int omid = orient(l.a, l.b, p[mid]);
     if (!omid) {
       lo = mid:
       break:
     if (omid == olo)
       lo = mid;
     else
       hi = mid;
   }
   ret.push_back(lo);
 return ret;
// Calculate [ACW, CW] tangent pair from an external
    point
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) == dir ? u : v;
Point pointPolyTangent(const Polygon &pt, Point Q, int
    dir, int lo, int hi) {
  while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
   bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
   if (pvs && nxt) return pt[mid];
   if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q, dir, mid + 1, hi
         );
     Point p2 = pointPolyTangent(pt, Q, dir, lo, mid -
         1);
     return better(p1, p2, Q, dir);
```

```
if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
       hi = mid - 1;
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
      hi = mid - 1;
     else
       lo = mid + 1;
   }
   if (!nxt) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
       lo = mid + 1;
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
       hi = mid - 1:
     else
       lo = mid + 1;
   }
 }
 Point ret = pt[lo];
 for (int i = lo + 1; i <= hi; i++) ret = better(ret, pt</pre>
      [i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(const Polygon &pt,
    Point Q) {
 int n = pt.size();
 Point acw_tan = pointPolyTangent(pt, Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q, CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
4.5 Polygon
```

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &poly) {
 Polygon ret;
 int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = poly[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0 && (ret.empty() ||
         b != ret.back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p of n vertices
```

```
Tf signedPolygonArea(const Polygon &p) {
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() - 1; i++)</pre>
  ret += cross(p[i] - p[0], p[i + 1] - p[0]);
 return ret / 2;
// given a polygon p of n vertices, generates the convex
   hull in in CCW
// Tested on https://acm.timus.ru/problem.aspx?space=1&
// Caution: when all points are colinear AND
    removeRedundant == false
// output will be contain duplicate points (from upper
   hull) at back
Polygon convexHull(Polygon p, bool removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
 p.erase(unique(p.begin(), p.end()), p.end());
 int n = p.size();
 Polygon ch(n + n);
 int m = 0; // preparing lower hull
 for (int i = 0; i < n; i++) {</pre>
   while (m > 1 &&
         dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m-
               1])) <= check)
     m--:
   ch[m++] = p[i];
 int k = m; // preparing upper hull
 for (int i = n - 2; i \ge 0; i--) {
   while (m > k \&\&
         dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m-
               21)) <= check)
     m--;
   ch[m++] = p[i];
 }
 if (n > 1) m--;
 ch.resize(m);
 return ch;
// returns inside = -1, on = 0, outside = 1
int pointInPolygon(const Polygon &p, Point o) {
 using Linear::onSegment;
 int wn = 0, n = p.size();
 for (int i = 0; i < n; i++) {</pre>
   int j = (i + 1) \% n;
   if (onSegment(o, Segment(p[i], p[j])) || o == p[i])
   int k = dcmp(cross(p[j] - p[i], o - p[i]));
```

```
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```

```
int d1 = dcmp(p[i].y - o.y);
   int d2 = dcmp(p[j].y - o.y);
   if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
   if (k < 0 && d2 <= 0 && d1 > 0) wn--;
 return wn ? -1 : 1;
// Given a simple polygon p, and a line l, returns (x, y)
// x = longest segment of 1 in p, y = total length of 1
pair<Tf, Tf> linePolygonIntersection(Line 1, const
    Polygon &p) {
  using Linear::lineLineIntersection;
  int n = p.size();
 vector<pair<Tf, int>> ev;
 for (int i = 0; i < n; ++i) {
   Point a = p[i], b = p[(i + 1) \% n], z = p[(i - 1 + n)
        % n];
   int ora = orient(l.a, l.b, a), orb = orient(l.a, l.b,
       orz = orient(1.a, 1.b, z);
   if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(d, 0);
       // else // Point Touch
     } else if (orz)
       ev.emplace_back(d, orz);
     else if (orb)
       ev.emplace_back(d, orb);
   } else if (ora == -orb) {
     Point ins:
     lineLineIntersection(l, Line(a, b), ins);
     ev.emplace_back(dot(ins - 1.a, 1.b - 1.a), 0);
   }
  sort(ev.begin(), ev.end());
 Tf ans = 0, len = 0, last = 0, tot = 0;
  bool active = false;
  int sign = 0;
  for (auto &qq : ev) {
   int tp = qq.second;
   Tf d = qq.first; /// current Segment is (last, d)
   if (sign) {
                 /// On Border
    len += d - last;
     tot += d - last:
     ans = max(ans, len);
     if (tp != sign) active = !active;
     sign = 0;
```

```
} else {
     if (active) { /// Strictly Inside
       len += d - last;
       tot += d - last;
       ans = max(ans, len);
     if (tp == 0)
       active = !active;
       sign = tp;
   last = d;
   if (!active) len = 0;
  ans /= length(l.b - l.a);
 tot /= length(l.b - l.a);
 return {ans, tot};
    Graph
|5.1 \quad LCA, ETT, VT|
struct lca_table {
 tree &T:
 int n, LOG = 20;
  vector<vector<int>> anc;
  vector<int> level;
 void setupLifting(int node, int par) {
   for (int v : T[node])
     if (v != par) {
       anc[v][0] = node, level[v] = level[node] + 1;
       for (int k = 1; k < LOG; k++) anc[v][k] = anc[anc
           [v][k-1]][k-1];
       setupLifting(v, node);
     }
 }
 lca_table(tree \&T, int root = 0) : T(T), n(T.n) {
   LOG = 33 - __builtin_clz(n);
   anc.assign(n, vector<int>(LOG, root));
   level.resize(n);
   setupLifting(root, root);
 }
 int lca(int u, int v) {
   if (level[u] > level[v]) swap(u, v);
   for (int k = LOG - 1; "k; k--)
     if (level[u] + (1 << k) <= level[v]) v = anc[v][k];
   if (u == v) return u:
```

for (int k = LOG - 1; ~k; k--)

v][k];

```
return anc[u][0]:
                                                       int getAncestor(int node, int ht) {
                                                         for (int k = 0; k < LOG; k++)
                                                           if (ht & (1 << k)) node = anc[node][k];</pre>
                                                         return node:
                                                       }
                                                       int distance(int u, int v) {
                                                         int g = lca(u, v);
                                                         return level[u] + level[v] - 2 * level[g];
                                                      struct euler_tour {
                                                       int time = 0:
                                                       tree &T;
                                                       int n:
                                                        vector<int> start, finish, level, par;
                                                        euler_tour(tree &T, int root = 0)
                                                           : T(T), n(T.n), start(n), finish(n), level(n), par(
                                                         time = 0;
                                                         call(root);
                                                        void call(int node, int p = -1) {
                                                         if (p != -1) level[node] = level[p] + 1;
                                                         start[node] = time++;
                                                         for (int e : T[node])
                                                           if (e != p) call(e, node);
                                                         par[node] = p;
                                                         finish[node] = time++;
                                                       bool isAncestor(int node, int par) {
                                                         return start[par] <= start[node] and finish[par] >=
                                                              finish[node];
                                                       int subtreeSize(int node) { return finish[node] - start
                                                            [node] + 1 >> 1: }
                                                      tree virtual_tree(vector<int> &nodes, lca_table &table,
                                                          euler tour &tour) {
                                                        sort(nodes.begin(), nodes.end(),
                                                            [&](int x, int y) { return tour.start[x] < tour.</pre>
                                                                 start[v]; });
                                                        int n = nodes.size();
                                                       for (int i = 0; i + 1 < n; i++)</pre>
                                                         nodes.push_back(table.lca(nodes[i], nodes[i + 1]));
                                                        sort(nodes.begin(), nodes.end());
                                                        nodes.erase(unique(nodes.begin(), nodes.end()), nodes.
                                                            end());
if (anc[u][k] != anc[v][k]) u = anc[u][k], v = anc[
                                                        sort(nodes.begin(), nodes.end(),
```

```
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```

```
[&](int x, int y) { return tour.start[x] < tour.</pre>
          start[v]; });
 n = nodes.size();
 stack<int> st;
 st.push(0);
 tree ans(n);
 for (int i = 1; i < n; i++) {
   while (!tour.isAncestor(nodes[i], nodes[st.top()]))
        st.pop();
   ans.addEdge(st.top(), i);
   st.push(i);
 return ans;
set<int> getCenters(tree &T) {
 int n = T.n:
 vector<int> deg(n), q;
 set<int> s;
 for (int i = 0; i < n; i++) {</pre>
   deg[i] = T[i].size();
   if (deg[i] == 1) q.push_back(i);
   s.insert(i);
 for (vector<int> t; s.size() > 2; q = t) {
   for (auto x : q) {
    for (auto e : T[x])
       if (--deg[e] == 1) t.push_back(e);
     s.erase(x);
   }
 }
 return s;
```

5.2 SCC

```
int main() {
  for (int i = 1; i <= n; i++) {
    if (!vis[i]) dfs(i);
  }
  memset(vis, 0, sizeof vis);
  reverse(order.begin(), order.end());
  for (int i : order) {
    if (!vis[i]) {
        // one component is found
        dfs(i, 1), component.clear();
    }
  }
}
</pre>
```

5.3 Euler Tour on Edge

```
// for simplicity, G[idx] contains the adjacency list of
    a node
// while G(e) is a reference to the e-th edge.
const int N = 2e5 + 5;
int in[N], out[N], fwd[N], bck[N];
int t = 0;
void dfs(graph &G, int node, int par) {
 out[node] = t:
 for (int e : G[node]) {
   int v = G(e).to(node);
   if (v == par) continue;
   fwd[e] = t++;
   dfs(G, v, node);
   bck[e] = t++;
 in[node] = t - 1;
void init(graph &G, int node) {
 t = 0;
 dfs(G, node, node);
```

5.4 LCA In O(1)

```
/* LCA in O(1)
 * depth calculates weighted distance
 * level calculates distance by number of edges
 * Preprocessing in NlongN */
LL depth[N];
int level[N];
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
vector<PII> adj[N];
```

```
int n, root, Time, cur;
void init(int nodes, int root_) {
 n = nodes, root = root_, LOG[0] = LOG[1] = 0;
 for (int i = 2; i <= n; i++) LOG[i] = LOG[i >> 1] + 1;
 for (int i = 0; i <= n; i++) adj[i].clear();</pre>
void addEdge(int u, int v, int w) {
 adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
 if (en[u] > en[v]) swap(u, v);
 if (st[v] <= st[u] && en[u] <= en[v]) return v;</pre>
 int 1 = LOG[id[v] - id[u] + 1];
 int p1 = id[u], p2 = id[v] - (1 << 1) + 1;
 int d1 = level[table[1][p1]], d2 = level[table[1][p2]];
 if (d1 < d2)
   return par[table[1][p1]];
 else
   return par[table[1][p2]];
LL dist(int u, int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[1] * 2));
/* Euler tour */
void dfs(int u, int p) {
 st[u] = ++Time, par[u] = p;
 for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
 en[u] = ++Time;
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
 cur = Time = 0, dfs(root, root);
```

```
for (int i = 1: i \le n: i++) table[0][i] = a[i]:
for (int 1 = 0; 1 < L - 1; 1++) {
 for (int i = 1; i <= n; i++) {</pre>
   table[1 + 1][i] = table[1][i];
   bool C1 = (1 << 1) + i <= n;
   bool C2 = level[table[1][i + (1 << 1)]] < level[</pre>
        table[1][i]];
   if (C1 && C2) table[1 + 1][i] = table[1][i + (1 <<</pre>
        1)];
 }
}
```

5.5 HLD

```
const int N = 1e6 + 7:
template <typename DT>
struct Segtree {
 // write lazy segtree here
Segtree<int> tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
 depth[u] = depth[p] + 1;
  for (auto v : adj[u]) {
   if (v == p) continue;
   par[v] = u;
   int subsz = dfs(v, u);
   if (subsz > mxsz) heavy[u] = v, mxsz = subsz;
   SZ += subsz;
 return SZ:
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
 if (heavy[u] != -1) decompose(heavy[u], h);
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
 }
int query(int a, int b) {
```

```
int ret = 0:
for (; head[a] != head[b]; b = par[head[b]]) {
 if (depth[head[a]] > depth[head[b]]) swap(a, b);
 ret += tree.query(1, 0, cnt, pos[head[b]], pos[b]);
if (depth[a] > depth[b]) swap(a, b);
ret += tree.query(1, 0, cnt, pos[a], pos[b]);
return ret;
```

5.6 Dinic Max Flow

```
/// flow with demand(lower bound) only for DAG
// create new src and sink
// add_edge(new src, u, sum(in_demand[u]))
// add_edge(u, new sink, sum(out_demand[u]))
// add edge(old sink, old src, inf)
// if (sum of lower bound == flow) then demand satisfied
// flow in every edge i = demand[i] + e.flow
using Ti = long long;
const Ti INF = 1LL << 60;</pre>
struct edge {
 int v, u;
 Ti cap, flow = 0;
 edge(int v, int u, Ti cap) : v(v), u(u), cap(cap) {}
const int N = 1e5 + 50:
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
bool bfs(int s, int t) {
 for (q.push(s), level[s] = 0; !q.empty(); q.pop()) {
   for (int id : adj[q.front()]) {
     auto &ed = edges[id];
     if (ed.cap - ed.flow > 0 and level[ed.u] == -1)
       level[ed.u] = level[ed.v] + 1, q.push(ed.u);
   }
 return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
 if (pushed == 0) return 0;
 if (v == t) return pushed;
 for (int &cid = ptr[v]; cid < adj[v].size(); cid++) {</pre>
   int id = adj[v][cid];
   auto &ed = edges[id];
```

```
if (level[v] + 1 != level[ed.u] || ed.cap - ed.flow <</pre>
        1) continue;
   Ti tr = dfs(ed.u, min(pushed, ed.cap - ed.flow), t);
   if (tr == 0) continue;
   ed.flow += tr;
   edges[id ^ 1].flow -= tr;
   return tr;
 }
 return 0;
roid init(int nodes) {
 m = 0, n = nodes;
 for (int i = 0; i < n; i++) level[i] = -1, ptr[i] = 0,
     adj[i].clear();
void addEdge(int v, int u, Ti cap) {
 edges.emplace_back(v, u, cap), adj[v].push_back(m++);
 edges.emplace_back(u, v, 0), adj[u].push_back(m++);
Ti maxFlow(int s, int t) {
 Ti f = 0:
 for (auto &ed : edges) ed.flow = 0;
 for (; bfs(s, t); memset(level, -1, n * 4)) {
   for (memset(ptr, 0, n * 4); Ti pushed = dfs(s, INF, t
       ); f += pushed)
 return f;
```

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5.7 Min Cost Max Flow

```
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count()):
const LL inf = 1e9;
struct edge {
 int v, rev;
 LL cap, cost, flow;
 edge() {}
 edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(cost), flow(0) {}
struct mcmf {
 int src, sink, n;
 vector<int> par, idx, Q;
 vector<bool> ing;
 vector<LL> dis;
 vector<vector<edge>> g;
 mcmf() {}
 mcmf(int src, int sink, int n)
     : src(src),
```

```
sink(sink).
     n(n),
     par(n),
     idx(n),
     inq(n),
     dis(n),
     g(n),
     \mathbb{Q}(10000005) {} // use \mathbb{Q}(n) if not using random
void add_edge(int u, int v, LL cap, LL cost, bool
    directed = true) {
 edge _u = edge(v, g[v].size(), cap, cost);
 edge _v = edge(u, g[u].size(), 0, -cost);
 g[u].pb(_u);
 g[v].pb(_v);
 if (!directed) add_edge(v, u, cap, cost, true);
bool spfa() {
 for (int i = 0; i < n; i++) {
   dis[i] = inf, inq[i] = false;
 int f = 0, 1 = 0;
 dis[src] = 0, par[src] = -1, Q[1++] = src, inq[src] =
 while (f < 1) {
   int u = Q[f++];
   for (int i = 0; i < g[u].size(); i++) {</pre>
     edge &e = g[u][i];
     if (e.cap <= e.flow) continue;</pre>
     if (dis[e.v] > dis[u] + e.cost) {
       dis[e.v] = dis[u] + e.cost;
       par[e.v] = u, idx[e.v] = i;
       if (!inq[e.v]) inq[e.v] = true, Q[1++] = e.v;
      // if (!inq[e.v]) {
       // inq[e.v] = true;
       // if (f \&\& rnd() \& 7) Q[--f] = e.v;
       // else Q[1++] = e.v;
       // }
     }
   }
   inq[u] = false;
 return (dis[sink] != inf);
pair<LL, LL> solve() {
 LL mincost = 0, maxflow = 0;
 while (spfa()) {
  LL bottleneck = inf;
   for (int u = par[sink], v = idx[sink]; u != -1; v =
        idx[u], u = par[u]) {
     edge &e = g[u][v];
```

```
bottleneck = min(bottleneck, e.cap - e.flow);
     for (int u = par[sink], v = idx[sink]; u != -1; v = |graph &create_tree() {
          idx[u], u = par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck;
     mincost += bottleneck * dis[sink], maxflow +=
         bottleneck;
   }
   return make_pair(mincost, maxflow);
// want to minimize cost and don't care about flow
// add edge from sink to dummy sink (cap = inf, cost = 0)
// add edge from source to sink (cap = inf, cost = 0)
// run mcmf, cost returned is the minimum cost
```

5.8 Bridge Tree

```
vector<vector<int>> components;
vector<int> depth, low;
stack<int> st;
vector<int> id;
vector<edge> bridges;
graph tree;
void find_bridges(int node, graph &G, int par = -1, int d
 low[node] = depth[node] = d;
 st.push(node);
 for (int id : G[node]) {
   int to = G(id).to(node);
   if (par != to) {
     if (depth[to] == -1) {
       find_bridges(to, G, node, d + 1);
       if (low[to] > depth[node]) {
         bridges.emplace_back(node, to);
         components.push_back({});
         for (int x = -1; x != to; x = st.top(), st.pop()
           components.back().push_back(st.top());
      }
     low[node] = min(low[node], low[to]);
   }
 if (par == -1) {
   components.push_back({});
   while (!st.empty()) components.back().push_back(st.
       top()), st.pop();
```

```
for (auto &comp : components) {
  int idx = tree.addNode();
  for (auto &e : comp) id[e] = idx;
for (auto &[l, r] : bridges) tree.addEdge(id[l], id[r])
 return tree;
void init(graph &G) {
 int n = G.n;
 depth.assign(n, -1), id.assign(n, -1), low.resize(n);
 for (int i = 0; i < n; i++)</pre>
   if (depth[i] == -1) find_bridges(i, G);
```

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5.9 Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
 int cnt = 0;
 vector<int> vs;
 for (auto v : g1[u]) {
   if (v != p) {
     int got = dfs(v, u);
     vs.pb(got);
     cnt++;
   }
 }
 if (!cnt) return 1;
 sort(vs.begin(), vs.end());
 string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
 s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] = ++ind;
 int ret = mp[s];
 return ret;
```

Math

6.1 Linear Sieve

```
const int N = 1e7;
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5], cnt[N + 5], POW[N
     + 51:
bool prime[N + 5];
```

```
int SOD[N + 5];
void init() {
 fill(prime + 2, prime + N + 1, 1);
 SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
 for (LL i = 2; i <= N; i++) {</pre>
   if (prime[i]) {
     primes.push_back(i), spf[i] = i;
     phi[i] = i - 1;
    NOD[i] = 2, cnt[i] = 1;
     SOD[i] = i + 1, POW[i] = i;
   for (auto p : primes) {
     if (p * i > N or p > spf[i]) break;
     prime[p * i] = false, spf[p * i] = p;
     if (i % p == 0) {
      phi[p * i] = p * phi[i];
       NOD[p * i] = NOD[i] / (cnt[i] + 1) * (cnt[i] + 2)
              cnt[p * i] = cnt[i] + 1;
       SOD[p * i] = SOD[i] / SOD[POW[i]] * (SOD[POW[i]]
           + p * POW[i]),
              POW[p * i] = p * POW[i];
       break:
     } else {
       phi[p * i] = phi[p] * phi[i];
      NOD[p * i] = NOD[p] * NOD[i], cnt[p * i] = 1;
       SOD[p * i] = SOD[p] * SOD[i], POW[p * i] = p;
    }
  }
```

6.2 Pollard Rho

```
for (LL x : a) {
   LL p = bigmod(x \% n, d, n), i = s;
   for (; p != 1 and p != n - 1 and x % n and i--; p =
       mul(p, p, n))
   if (p != n - 1 and i != s) return false;
 return true;
LL get_factor(LL n) {
 auto f = [\&](LL x) \{ return mul(x, x, n) + 1; \};
 LL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
 for (; t++ % 40 or gcd(prod, n) == 1; x = f(x), y = f(f)
      (v))) {
   (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) - min(x, y), n)) ? q
        : prod;
 return gcd(prod, n);
map<LL, int> factorize(LL n) {
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = \{2, 3, 5, 7, 11, 13, 17, 19, 23,
      29, 31, 37, 41,
                     43, 47, 53, 59, 61, 67, 71, 73, 79,
                         83, 89, 97};
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
 auto _factor = [&](LL n, auto &_factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++;
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
 };
 _factor(n, _factor);
 return res;
```

6.3 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
  if (b == 0)
```

```
return {1, 0, a};
else {
    auto [x, y, g] = EGCD(b, a % b);
    return {y, x - a / b * y, g};
    }
}
// given modulo equations, will apply CRT
PLL CRT(vector<PLL> &v) {
    LL V = 0, M = 1;
    for (auto &[v, m] : v) { // value % mod
        auto [x, y, g] = EGCD(M, m);
        if ((v - V) % g != 0) return {-1, 0};
        V += x * (v - V) / g % (m / g) * M, M *= m / g;
        V = (V % M + M) % M;
    }
    return make_pair(V, M);
}
```

6.4 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
  memset(mob, -1, sizeof mob);
  mob[1] = 1;
  for (int i = 2; i < N; i++)
    if (mob[i]) {
      for (int j = i + i; j < N; j += i) mob[j] -= mob[i ];
    }
}</pre>
```

6.5 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
 assert((n & (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {</pre>
   for (int i = 0; i < k; i++) {
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
 }
 wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
```

```
wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N)
       ));
   wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
       N));
 }
void fft(vector<CD> &v, bool invert = false) {
 if (v.size() != perm.size()) precalculate(v.size());
 for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
 for (int len = 2; len <= N; len *= 2) {</pre>
   for (int i = 0, d = N / len; i < N; i += len) {</pre>
     for (int j = 0, idx = 0; j < len / 2; j++, idx += d
         ) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
     }
   }
 }
 if (invert) {
   for (int i = 0; i < N; i++) v[i] /= N;</pre>
 }
void pairfft(vector<CD> &a, vector<CD> &b, bool invert =
    false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0,</pre>
      1);
 fft(p, invert);
 p.push_back(p[0]);
 for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
   }
 }
vector<LL> multiply(const vector<LL> &a, const vector<LL>
     &b) {
 int n = 1;
  while (n < a.size() + b.size()) n <<= 1;</pre>
 vector<CD> fa(a.begin(), a.end()), fb(b.begin(), b.end
      ());
 fa.resize(n):
                                                            const LL N = 1 << 18;
```

```
fb.resize(n):
           fft(fa); fft(fb);
 pairfft(fa, fb);
 for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];</pre>
 fft(fa, true);
 vector<LL> ans(n):
 for (int i = 0; i < n; i++) ans[i] = round(fa[i].real()</pre>
      ):
 return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const vector<LL> &
 int n = 1:
 while (n < a.size() + b.size()) n <<= 1;</pre>
 vector<CD> al(n), ar(n), bl(n), br(n);
 for (int i = 0; i < a.size(); i++) al[i] = a[i] % M / B</pre>
      , ar[i] = a[i] % M % B;
 for (int i = 0; i < b.size(); i++) bl[i] = b[i] % M / B</pre>
      , br[i] = b[i] % M % B;
 pairfft(al, ar);
 pairfft(bl, br);
          fft(al); fft(ar); fft(bl); fft(br);
 for (int i = 0; i < n; i++) {</pre>
   CD 11 = (al[i] * bl[i]), lr = (al[i] * br[i]);
   CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
   al[i] = 11:
   ar[i] = lr;
   bl[i] = rl:
   br[i] = rr;
 pairfft(al, ar, true);
 pairfft(bl, br, true);
           fft(al, true); fft(ar, true); fft(bl, true);
      fft(br, true);
 vector<LL> ans(n);
 for (int i = 0: i < n: i++) {</pre>
   LL right = round(br[i].real()), left = round(al[i].
        real());
   LL mid = round(round(bl[i].real()) + round(ar[i].real
   ans[i] = ((left \% M) * B * B + (mid \% M) * B + right)
         % M;
 return ans;
6.6 NTT
```

```
const LL MOD = 786433:
vector<LL> P[N];
LL rev[N], w[N | 1], a[N], b[N], inv_n, g;
LL Pow(LL b, LL p) {
 LL ret = 1:
 while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) \% MOD;
   p >>= 1;
 return ret;
LL primitive_root(LL p) {
 vector<LL> factor;
 LL phi = p - 1, n = phi;
 for (LL i = 2; i * i <= n; i++) {
   if (n % i) continue;
   factor.emplace_back(i);
   while (n \% i == 0) n /= i;
 }
 if (n > 1) factor.emplace_back(n);
 for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true;
   for (LL i = 0; i < factor.size() && ok; i++)</pre>
     ok &= Pow(res, phi / factor[i]) != 1;
   if (ok) return res;
 }
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - __builtin_clz(n));
 LL r = Pow(g, (MOD - 1) / n);
 inv_n = Pow(n, MOD - 2);
 w[0] = w[n] = 1;
 for (LL i = 1; i < n; i++) w[i] = (w[i - 1] * r) % MOD;</pre>
 for (LL i = 1: i < n: i++)
   rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
for (LL i = 1; i < n - 1; i++)
   if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (LL m = 2; m <= n; m <<= 1) {
   for (LL i = 0; i < n; i += m) {
     for (LL j = 0; j < (m >> 1); j++) {
       LL &u = a[i + j], &v = a[i + j + (m >> 1)];
       LL t = v * w[dir ? n - n / m * j : n / m * j] %
       v = u - t < 0 ? u - t + MOD : u - t;
       u = u + t >= MOD ? u + t - MOD : u + t:
```

```
16
```

```
}
  if (dir)
   for (LL i = 0; i < n; i++) a[i] = (inv_n * a[i]) %</pre>
vector<LL> mul(vector<LL> p, vector<LL> q) {
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
  while (sz < t) sz <<= 1;
  prepare(sz);
  for (LL i = 0; i < n; i++) a[i] = p[i];</pre>
  for (LL i = 0; i < m; i++) b[i] = q[i];
  for (LL i = n: i < sz: i++) a[i] = 0:
  for (LL i = m; i < sz; i++) b[i] = 0;</pre>
  NTT(a, sz);
  NTT(b, sz);
  for (LL i = 0; i < sz; i++) a[i] = (a[i] * b[i]) % MOD;</pre>
  NTT(a, sz, 1);
  vector<LL> c(a, a + sz);
  while (c.size() && c.back() == 0) c.pop_back();
  return c;
```

7 String

7.1 Aho Corasick

```
const int sg = 26, N = 1e3 + 9;
struct aho_corasick {
 struct node {
   node *link, *out, *par;
   bool leaf:
   LL val;
   int cnt, last, len;
   char p_ch;
   array<node*, sg> to;
   node(node* par = NULL, char p_ch = '$', int len = 0)
       : par(par), p_ch(p_ch), len(len) {
     val = leaf = cnt = last = 0;
    link = out = NULL:
  }
 vector<node> trie;
 node* root;
 aho_corasick() {
   trie.reserve(N), trie.emplace_back();
```

```
root = &trie[0]:
   root->link = root->out = root;
 inline int f(char c) { return c - 'a'; }
  inline node* add_node(node* par = NULL, char p_ch = '$' PLL inv_base = inverse(base);
      , int len = 0) {
   trie.emplace_back(par, p_ch, len);
   return &trie.back();
  void add_str(const string& s, LL val = 1) {
   node* now = root:
   for (char c : s) {
     int i = f(c):
     if (!now->to[i]) now->to[i] = add_node(now, c, now
         ->len + 1);
     now = now - > to[i]:
   }
   now->leaf = true, now->val++;
  void push_links() {
   queue<node*> q;
   for (q.push(root); q.empty(); q.pop()) {
     node *cur = q.front(), *link = cur->link;
     cur->out = link->leaf ? link : link->out;
     int idx = 0;
     for (auto& next : cur->to) {
       if (next != NULL) {
         next->link = cur != root ? link->to[idx++] :
             root:
         q.push(next);
       } else
         next = link->to[idx++]:
     }
   }
   cur->val += link->val;
};
```

```
7.2 Double hash
// define +, -, * for (PLL, LL) and (PLL, PLL), % for (
    PLL. PLL):
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
PLL power(PLL a, LL p) {
 PLL ans = PLL(1, 1);
 for(; p; p >>= 1, a = a * a % mod) {
     if(p \& 1) ans = ans * a % mod;
 }
 return ans;
```

```
PLL inverse(PLL a) { return power(a, (mod.ff - 1) * (mod.
    ss - 1) - 1); }
PLL val:
vector<PLL> P;
void hash_init(int n) {
 P.resize(n + 1);
 P[0] = PLL(1, 1);
 for (int i = 1; i \le n; i++) P[i] = (P[i-1] * base) %
PLL append(PLL cur, char c) { return (cur * base + c) %
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) { return (P[k] * c +
    cur) % mod: }
/// replaces the i-th (0-indexed) character from right
    from a to b:
PLL replace(PLL cur, int i, char a, char b) {
 cur = (cur + P[i] * (b - a)) \% mod;
 return (cur + mod) % mod;
/// Erases c from the back of the string
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv_base) % mod + mod) % mod;
/// Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod + mod) % mod;
/// concatenates two strings where length of the right is
PLL concat(PLL left, PLL right, int k) { return (left * P
    [k] + right) % mod; }
/// Calculates hash of string with size len repeated cnt
/// This is O(\log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
 PLL ret = (hash * mul) % mod;
 if (P[len].ff == 1) ret.ff = hash.ff * cnt;
 if (P[len].ss == 1) ret.ss = hash.ss * cnt;
 return ret:
LL get(PLL hash) { return ((hash.ff << 32) ^ hash.ss); }
struct hashlist {
```

```
int len;
 vector<PLL> H, R;
  hashlist() {}
  hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL
        (0, 0));
   for (int i = 1; i <= len; i++) H[i] = append(H[i -</pre>
        1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i] = append(R[i +
        1], s[i - 1]);
 /// 1-indexed
  PLL range_hash(int 1, int r) {
   return ((H[r] - H[l - 1] * P[r - l + 1]) \% mod + mod)
        % mod:
  PLL reverse_hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 + 1]) \% mod + mod)
        % mod:
 PLL concat_range_hash(int 11, int r1, int 12, int r2) {
   return concat(range_hash(l1, r1), range_hash(l2, r2),
        r2 - 12 + 1);
 PLL concat_reverse_hash(int 11, int r1, int 12, int r2)
   return concat(reverse_hash(12, r2), reverse_hash(11,
       r1), r1 - 11 + 1);
 }
};
```

7.3 KMP

```
vector<int> prefix_function(string s) {
  int n = (int)s.length();
  vector<int> pi(n);
  for (int i = 1; i < n; i++) {
    int j = pi[i-1];
    while (j > 0 && s[i] != s[j])
        j = pi[j-1];
    if (s[i] == s[j])
        j++;
    pi[i] = j;
  }
  return pi;
}
```

7.4 Manacher's

```
vector<int> d1(n);
```

7.5 String Match FFT

push_back(i - m + 1);

```
//find occurrences of t in s where '?'s are automatically
     matched with any character
//res[i + m - 1] = sum_j = 0 to m - 1_{s[i + j]} * t[j] * (s]
    [i + i] - t[i]
vector<int> string_matching(string &s, string &t) {
 int n = s.size(), m = t.size();
 vector<int> s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[i] == '?' ? 0 : s[</pre>
     i] - 'a' + 1; //assign any non zero number for non
      171s
 for(int i = 0; i < n; i++) s2[i] = s1[i] * s1[i];</pre>
 for(int i = 0; i < n; i++) s3[i] = s1[i] * s2[i];
 vector<int> t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[i] == '?' ? 0 : t[]
     i] - 'a' + 1:
 for(int i = 0; i < m; i++) t2[i] = t1[i] * t1[i];</pre>
 for(int i = 0; i < m; i++) t3[i] = t1[i] * t2[i];</pre>
 reverse(t1.begin(), t1.end());
 reverse(t2.begin(), t2.end());
 reverse(t3.begin(), t3.end());
 vector<int> s1t3 = multiply(s1, t3);
 vector<int> s2t2 = multiply(s2, t2);
 vector<int> s3t1 = multiply(s3, t1);
 vector<int> res(n);
 for(int i = 0; i < n; i++) res[i] = s1t3[i] - s2t2[i] *</pre>
       2 + s3t1[i]:
 vector<int> oc;
 for(int i = m - 1; i < n; i++) if(res[i] == 0) oc.
```

```
return oc:
#include <bits/stdc++.h>
using namespace std;
 O(|S| + |alphabet|) Suffix Array
LIM := max{s[i]} + 2
void inducedSort (const vector <int> &vec, int val_range,
     vector <int> &SA, const vector <int> &sl, const
    vector <int> &lms idx) {
 vector <int> l(val_range, 0), r(val_range, 0);
 for (int c : vec) {
   ++r[c]; if (c + 1 < val_range) ++l[c + 1];
 partial_sum(1.begin(), 1.end(), 1.begin());
 partial_sum(r.begin(), r.end(), r.begin());
 fill(SA.begin(), SA.end(), -1);
 for (int i = lms_idx.size() - 1; i \ge 0; --i) SA[--r[
     vec[lms_idx[i]]] = lms_idx[i];
 for (int i : SA) if (i > 0 and sl[i - 1]) SA[l[vec[i -
     1]]++] = i - 1;
 fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c]:
 partial_sum(r.begin(), r.end(), r.begin());
 for (int k = SA.size() - 1, i = SA[k]; k; --k, i = SA[k]
     ]) {
   if (i and !sl[i-1]) SA[--r[vec[i-1]]] = i-1;
vector <int> suffixArray (const vector <int> &vec, int
    val_range) {
 const int n = vec.size():
 vector <int> sl(n), SA(n), lms_idx;
 for (int i = n - 2; i >= 0; --i) {
   sl[i] = vec[i] > vec[i + 1] or (vec[i] == vec[i + 1]
       and sl[i + 1]);
   if (sl[i] and !sl[i + 1]) lms_idx.emplace_back(i + 1)
 reverse(lms_idx.begin(), lms_idx.end());
 inducedSort(vec, val_range, SA, sl, lms_idx);
 vector <int> new_lms_idx(lms_idx.size()), lms_vec(
```

lms idx.size()):

```
for (int i = 0, k = 0; i < n; ++i) {
   if (SA[i] > 0 and !sl[SA[i]] and sl[SA[i] - 1])
        new_lms_idx[k++] = SA[i];
  }
  int cur = 0; SA[n - 1] = 0;
  for (int k = 1; k < new_lms_idx.size(); ++k) {</pre>
    int i = new_lms_idx[k - 1], j = new_lms_idx[k];
   if (vec[i] ^ vec[j]) {
     SA[j] = ++cur; continue;
    bool flag = 0;
    for (int a = i + 1, b = j + 1; ; ++a, ++b) {
     if (vec[a] ^ vec[b]) {
       flag = 1; break;
     if ((!sl[a] and sl[a - 1]) or (!sl[b] and sl[b -
          1])) {
       flag = !(!sl[a] and sl[a - 1] and !sl[b] and sl[b | 7.7 Trie
             - 1]); break;
     }
   }
    SA[j] = flag ? ++cur : cur;
  for (int i = 0; i < lms_idx.size(); ++i) lms_vec[i] =</pre>
      SA[lms_idx[i]];
  if (cur + 1 < lms_idx.size()) {</pre>
    auto lms_SA = suffixArray(lms_vec, cur + 1);
   for (int i = 0; i < lms_idx.size(); ++i) new_lms_idx[</pre>
        i] = lms_idx[lms_SA[i]];
  inducedSort(vec, val_range, SA, sl, new_lms_idx);
      return SA:
vector <int> getSuffixArray (const string &s, const int
    LIM = 128) {
  vector <int> vec(s.size() + 1):
  copy(begin(s), end(s), begin(vec)); vec.back() = '$';
  auto ret = suffixArray(vec, LIM);
  ret.erase(ret.begin()); return ret;
}
// build RMQ on it to get LCP of any two suffix
vector <int> getLCParray (const string &s, const vector <</pre>
    int> &SA) {
  int n = s.size(), k = 0;
  vector <int> lcp(n), rank(n);
  for (int i = 0; i < n; ++i) rank[SA[i]] = i;</pre>
  for (int i = 0; i < n; ++i, k ? --k : 0) {
   if (rank[i] == n - 1) {
```

```
k = 0: continue:
   int j = SA[rank[i] + 1];
   while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j + k]
       kl) ++k:
   lcp[rank[i]] = k;
 lcp[n-1] = 0; return lcp;
int main() {
 string s; cin >> s;
 for (const int i : getSuffixArray(s)) printf("%d ", i);
 puts(""):
 return 0;
template<int sz>
struct Trie {
 Trie() : id(1) {
   memset(endMark, 0, sizeof endMark);
   for_each(all(trie), [](vector<int> &v) { v.assign(sz,
         0); });
 }
 void insert(const string &s) {
   int cur = 0:
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       trie[cur][val] = id++;
     cur = trie[cur][val];
   endMark[cur] = true;
 bool search(const string &s) {
   int cur = 0:
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       return false;
     cur = trie[cur][val];
   }
   return endMark[cur];
 }
private:
 int id, endMark[100005];
 vector<int> trie[100005];
```

```
7.8 Z Algo

vector<int> calcz(string s) {
  int n = s.size();
  vector<int> z(n);
  int l = 0, r = 0;
  for (int i = 1; i < n; i++) {
    if (i > r) {
      l = r = i;
      while (r < n && s[r] == s[r - 1]) r++;
      z[i] = r - 1, r--;
  } else {
    int k = i - 1;
    if (z[k] < r - i + 1) z[i] = z[k];
```

while (r < n && s[r] == s[r - 1]) r++;

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8 Extra

return z;

}

}

8.1 Make File

else {

1 = i:

z[i] = r - 1, r--;

```
.PHONY: clean
D ?= 0
ifeq ($(D), 1)
CXXFLAGS=-std=c++2a -g -DMUSHFIQ_DEBUG -
     DDBG_MACRO_NO_WARNING -Wall -Wextra -Wpedantic -
     Wshadow -Wformat=2 -Wfloat-equal -Wconversion -
     Wlogical-op -Wshift-overflow=2 -Wduplicated-cond -
     Wcast-qual -Wcast-align -Wno-variadic-macros -
     D_GLIBCXX_DEBUG -D_GLIBCXX_DEBUG_PEDANTIC -
     fsanitize=address -fsanitize=undefined -fno-
     sanitize-recover -fstack-protector -fsanitize-
     address-use-after-scope
else
CXXFLAGS=-02 -std=c++2a
endif
all:
@echo 'Error: No file given.'
clean:
@echo 'Cleaning executables...'
Ofind . ! -name stress.sh -type f -executable -delete
```

```
make $2
make $3
# $1 is actual code
# $2 is good code
# $3 is generator
for((i=1;i<1001;i++)); do</pre>
  echo 'Test #'$i
  timeout 5s ./$3 $RANDOM > in
  timeout 5s ./$1 < in> out
  timeout 5s./$2 < in> ans
 diff out ans | | break
done
    Sublime Build
  "shell_cmd": "make $file_base_name && timeout 5s ./
      $file_base_name < in > out",
  "working_dir": "$file_path",
  "selector": "source.cpp"
8.4 vimrc
" Auto import & Compile
:autocmd BufNewFile *.cpp Or ~/template.cpp
nnoremap <F4> :!xclip -o -sel clip > ~/cp/in <CR><CR>
inoremap <F4> <ESC>:!xclip -o -sel clip > ~/cp/in <CR><CR
nnoremap <F6> :!xclip -sel clip % <CR><CR>
inoremap <F6> <ESC>:!xclip -sel clip % <CR><CR>
autocmd filetype cpp nnoremap <F9> :wa \| !make %:r &&
    timeout 5s ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F9> <ESC>:wa \| !make %:r
    && timeout 5s ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp nnoremap <F10> :wa \| !make clean
    && make %:r D=1 \&\& ./%:r < ^/cp/in> ^/cp/out<CR>
autocmd filetype cpp inoremap <F10> <ESC>:wa \| !make
    clean && make %:r D=1 && ./%:r < ~/cp/in > ~/cp/out<</pre>
autocmd filetype python nnoremap <F9> :wa \| !python % <
    ~/cp/in> ~/cp/out<CR>
```

@echo 'Executables cleaned.'

8.2 Stress Tester

make \$1

```
autocmd filetype python inoremap <F9> <ESC>:wa \| !python set expandtab
     % < ~/cp/in> ~/cp/out<CR>
  Auto Completion
inoremap ( ()<left>
inoremap <expr> ) strpart(getline('.'), col('.')-1, 1) == set cindent
     ")" ? "\<Right>" : ")"
inoremap { {}<left>
inoremap <expr> } strpart(getline('.'), col('.')-1, 1) == set guifont=*
     "}" ? "\<Right>" : "}"
inoremap [ [] < left >
inoremap <expr> ] strpart(getline('.'), col('.')-1, 1) == set incsearch
     "]" ? "\<Right>" : "]"
inoremap <expr> " strpart(getline('.'), col('.')-1, 1) == set hlsearch
     "\"" ? "\<Right>" : "\"\"\<left>"
inoremap <expr> ' strpart(getline('.'), col('.')-1, 1) == " bubt site
     "\'" ? "\<Right>" : "\'\'\<left>"
inoremap <expr> <CR> <sid>insert_newline()
function s:insert_newline() abort
 let pair = strpart(getline('.'), col('.')-2, 2)
 return stridx('{}[]', pair) % 2 == 0 && strlen(pair) == set showcmd
       2 ? "\<CR>\<ESC>\O" : "\<CR>"
endfunction
inoremap <expr> <space> <sid>insert_space()
function s:insert_space() abort
 let pair = strpart(getline('.'), col('.')-2, 2)
 return stridx('(){}[]', pair) % 2 == 0 && strlen(pair)
      == 2 ? "\<space>\<left>" : "\<space>"
endfunction
inoremap <expr> <bs> <sid>rm_pair()
function s:rm_pair() abort
let pair = strpart(getline('.'), col('.')-2, 2)
return stridx('(){}[]''', pair) % 2 == 0 && strlen(
     pair) == 2 ? "\<del>\<bs>" : "\<bs>"
endfunction
set nocompatible
                           " be iMproved, required
filetype on
                          " required
filetype plugin on
filetype plugin indent on
syntax on
set mouse=a
set number
set relativenumber
set tabstop=2
set shiftwidth=2
```

```
set softtabstop=2
set smartindent
set smarttab
set autoindent
set noerrorbells
set ruler
set backspace=indent,eol,start
set ignorecase
set nowrap
  set termguicolors
set foldmethod=indent
set nofoldenable
set cursorline
set laststatus=2
set wildmenu
if !has('nvim')
  set clipboard=unnamedplus
endif
if !has('nvim')
  set ttymouse=xterm2
endif
nnoremap \langle S-j \rangle : m .+1 \langle CR \rangle ==
nnoremap \langle S-k \rangle : m .-2 \langle CR \rangle ==
vnoremap <S-j> :m '>+1<CR>gv==gv
vnoremap \langle S-k \rangle : m \langle -2 \langle CR \rangle gv == gv
```

Equations and Formulas

Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1 = 1 \text{ and } C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

The number of ways to completely parenthesize n+1 factors. The number of triangulations of a convex polygon with n+2sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (ver- $S^d(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$ tices are not numbered). A rooted binary tree is full if every 9.4 Other Combinatorial Identities vertex has either two children or no children.

Number of permutations of $1, \ldots, n$ that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321

9.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) = \frac{n}{2}$$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$\begin{bmatrix} n & n \\ n & -k \end{bmatrix} = \sum_{0 \le i_1 \le i_2 \le i_k \le n} i_1 i_2 \dots i_k.$$

9.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

$$S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$$
, where $S(0,0) = 1$, $S(n,0) = S(0,n) = 0$ $S(n,2) = 2^{n-1} - 1$ $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k \text{ such that if } m \text{ is any integer, then } \gcd(a+m\cdot b,b) = \gcd(a,b)$ The gcd is a multiplicative function in the follow

An r-associated Stirling number of the second kind is the num- if a_1 and a_2 are relatively prime, then $gcd(a_1 \cdot a_2, b) =$ ber of ways to partition a set of n objects into k subsets, with $\gcd(a_1,b) \cdot \gcd(a_2,b)$.

each subset containing at least r elements. It is denoted by $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c))$. $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) = |\operatorname{lcm}(a,\operatorname{gcd}(b,c))| = \operatorname{gcd}(\operatorname{lcm}(a,b),\operatorname{lcm}(a,c)).$ $\left| kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \right|$

Denote the n objects to partition by the integers $1, 2, \dots, n$. $\gcd(a, b) = \sum \phi(k)$ Define the reduced Stirling numbers of the second kind, denoted $S^d(n,k)$, to be the number of ways to partition the integers $1, 2, \ldots, n$ into k nonempty subsets such that all ele- $\sum [\gcd(i, n) = k] = \phi(\frac{n}{k})$ ments in each subset have pairwise distance at least d. That i=1The number of ways to connect the 2n points on a circle to is, for any integers i and j in a given subset, it is required that $|i-j| \geq d$. It has been shown that these numbers satisfy,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

9.5 Different Math Formulas

Picks Theorem : A = i + b/2 - 1

Deragements: $d(i) = (i-1) \times (d(i-1) + d(i-2))$

$$\frac{n}{ab}$$
 - $\left\{\frac{b'n}{a}\right\}$ - $\left\{\frac{a'n}{b}\right\}$ +

The gcd is a multiplicative function in the following sense:

For non-negative integers a and b, where a and b are not both zero, $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$ k|a and k|b $\sum_{k=1}^{n} \gcd(k, n) = \sum_{d \mid n} d \cdot \phi\left(\frac{n}{d}\right)$ $\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{d \mid n} x^d \cdot \phi\left(\frac{n}{d}\right)$ $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{n=1}^{\infty} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{n=1}^{\infty} d \cdot \phi(d)$ $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$ $\left| \sum_{k=1}^{n} \frac{n}{\gcd(k,n)} \right| = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$ $\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{i=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$ $\sum_{i=1}^n \sum_{j=1}^n \gcd(i,j) = \sum_{d=1}^n \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^2$ $\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{j=1}^{n} \phi(i)i^{2}$ $F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left(\frac{\left(1 + \lfloor \frac{n}{l} \rfloor\right) \left(\lfloor \frac{n}{l} \rfloor\right)}{2} \right)^{2} \sum_{i=1}^{n} \mu(d) l d$