Young tableaus(applications of the heap properties)

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Please check the integrated code in young.py

1.Draw 4×4 tableau containing the elements {9,16,3,2,4,8,5,14,12}

2	3	5	14
4	8	9	16
12	8	8	8
∞	∞	∞	∞

2. Argue that an m×n Young tableau Y is empty if $Y[1,1]=\infty$. Argue that Y is full (contains mn elements) if $Y[m,n]<\infty$.

If $Y[1,1] = \infty$, it means all elements in the first row and column are ∞ . Y[1,1] is the smallest element in Y. The rest element of Y will be larger than Y[1,1]. If $Y[1,1] = \infty$, other elements in Y will also equal to ∞ to satisfy the rules of Young tableau. Hence, the Young tableau Y is empty.

If $Y[m, n] < \infty$, it means the last "children" element(the bottom right element) is not ∞ . And this bottom right element is the largest element in the Y. Each element is smaller than Y[m, n]. If Y[m, n] is smaller than ∞ , each element in Y will smaller than ∞ . Hence Y is not empty.

3. Give an algorithm to implement EXTRACT-MIN.

Time complexity:

$$T(p) = T(p-1) + O(1)$$

$$= T(p-2) + O(1) + O(1)$$

$$= T(p - (p - 1)) + (p - 1) * O(1)$$

$$= T(1) + (p - 1) * O(1)$$

$$= p * O(1)$$

$$= O(p)$$

$$= O(m + n)$$

Result:

Code:

def Young(Y, i, j, m, n):

```
x = i

y = j

if i < m \text{ and } Y[i][j] > Y[i + 1][j]:

x = i + 1

y = j

if j < n \text{ and } Y[x][y] > Y[i][j + 1]:

x = i

y = j + 1

if x != i \text{ or } y != j:

Y[i][j], Y[x][y] = Y[x][y], Y[i][j]

Young(Y, x, y, m, n)

def \text{ ExtractMin}(Y, m, n):

tmp = Y[0][0]

Y[0][0] = \text{sys.maxsize}

young(Y, 0, 0, m, n)

return tmp
```

4. Show how to insert a new element into a nonfull m×n Young tableau in O(m+n) time

Time Complexity:

The similar algorithm as 3. But an inversing process.

$$T(p) = T(p-1) + O(1)$$

$$= T(p-2) + O(1) + O(1)$$

$$= T(p - (p - 1)) + (p - 1) * O(1)$$

$$= T(1) + (p - 1) * O(1)$$

$$= p * O(1)$$

$$= O(p)$$

$$= O(m + n)$$

Result:

Code:

$$def Inverse(Y, m, n):$$

$$i = m$$

$$j = n$$

$$if m>0 \ and \ Y[m][n] < Y[m-1][n]:$$

```
i = m - 1
j = n

if n > 0 \text{ and } Y[i][j] < Y[m][n-1]:
i = m
j = n - 1

if i! = m \text{ or } j! = n:
Y[m][n], Y[i][j] = Y[i][j], Y[m][n]
Inverse(Y, i, j)

def Insert(Y, m, n, value):
if m < 0 \text{ and } n < 0:
return
Y[m][n] = value
Inverse(Y, m, n)
```

5. Using no other sorting method as a subroutine, show how to use an $n \times n$ Young tableau to sort n2 numbers in O(n3) time.

Inserting is actually an in-order operation. It will maintain the order in its algorithm. For inserting one element, we need to spend O(m + n) = O(n + n) = O(n). For inserting n^2 elements, we need to spend $n^2 O(n) = O(n^3)$.

6. Give an O(m+n)-time algorithm to determine whether a given number is stored in a given m×n Young tableau.

Time complexity:

$$T(p) = T(p-1) + O(1)$$

$$= T(p-2) + O(1) + O(1)$$

$$= T(p - (p - 1)) + (p - 1) * O(1)$$

$$= T(1) + (p - 1) * O(1)$$

$$= p * O(1)$$

$$= O(p)$$

$$= O(m + n)$$

Result:

```
Sorting
previous list is:
[9, 16, 3, 2, 4, 8, 5, 14, 12]
Sorted list is:
[2, 3, 4, 5, 8, 9, 12, 14, 16]

Searching
array is:
[9, 16, 3, 2, 4, 8, 5, 14, 12]
Search number: 3 search result: True
Search number: 50 search result: False
```

Code:

```
def \ Young Search(Y, i, j, m, n, value):
if \ Y[i][j] == value:
return \ True
x = i
y = j
if \ i > 0 \ and \ Y[i][j] > value:
x = i - 1
y = j
if \ j < n \ and \ Y[i][j] < value:
x = i
y = j + 1
if \ x == i \ and \ y == j:
return \ False
```

return YoungSearch(Y,x,y,m,n,value)
def search(Y, m, n, value):
 return YoungSearch(Y, m, 0, m, n, value)