

COMP9311 16S2 Assignment 3

Question 1

i.

a.

Candidate key: ACEF, BCEF

b.

Not BCNF, e.g.

- $AD \rightarrow B$, AD doesn't contain a key
- $C \rightarrow D$, C doesn't contain a key
- $BC \rightarrow A$, BC doesn't contain a key
- $B \rightarrow D$, B doesn't contain a key

c.

- The FD $C \rightarrow D$ violates BCNF
- To fix, we need to decompose into tables: CD and $ABCEF$.
- FDs for CD is $\{ C \rightarrow D \}$, therefore key is C , therefore BCNF.
- FDs for $ABCEF$ is $\{ BC \rightarrow A \}$, Keys are $BCEF$, and $BC \rightarrow A$ violates BCNF
- We then decompose BCA : FDs $\{ BC \rightarrow A \}$, the key is BC
- There is no FDs for $BCEF$, so it is BCNF

Final schema (with keys boldened): **CD**, **BCA**, **BCEF**

ii.

a.

Candidate key: AF, CF

b.

Not BCNF, e.g.

- $BC \rightarrow E$, BC doesn't contain a key
- $C \rightarrow AB$, C doesn't contain a key

c.

- The FD $C \rightarrow AB$ violates BCNF
- To fix, we need to decompose into tables: ABC and $CDEF$.
- FDs for ABC is $\{ C \rightarrow AB \}$, therefore key is C , therefore BCNF.
- There is no FDs for $CDEF$, so it is BCNF

Final schema (with keys boldened): **CAB**, **CDEF**

iii.

a.

Candidate key: ABCF, BCDF

b.

Not BCNF, e.g.

- $ABF \rightarrow D$, ABF doesn't contain a key
- $CD \rightarrow E$, CD doesn't contain a key
- $BD \rightarrow A$, BD doesn't contain a key

c.

- The FD $CD \rightarrow E$ violates BCNF
- To fix, we need to decompose into tables: CDE and $ABCDF$.
- FDs for CDE is $\{ CD \rightarrow E \}$, therefore key is CD , therefore BCNF.
- FDs for $ABCDF$ are $\{ ABF \rightarrow D, BD \rightarrow A \}$, Keys are $BCDF$, and $BD \rightarrow A$ violates BCNF
- We then decompose BDA : FDs $\{ BD \rightarrow A \}$, the key is BD
- There is no FDs for $BCDF$, so it is BCNF

Final schema (with keys boldened): **CDE**, **BDA**, **BCDF**

iv.

a.

Candidate key: AB

b.

Not BCNF, e.g.

- $BCD \rightarrow EF$, BCD doesn't contain a key
- $B \rightarrow C$, B doesn't contain a key

c.

- The FD $B \rightarrow C$ violates BCNF
- To fix, we need to decompose into tables: BC and $ABDEF$.
- FDs for BC is $\{ B \rightarrow C \}$, therefore key is B , therefore BCNF.
- FDs for $ABDEF$ is $\{ AB \rightarrow D \}$, Keys are $ABEF$, and $AB \rightarrow D$ violates BCNF
- We then decompose ABD : FDs $\{ AB \rightarrow D \}$, the key is AB
- There is no FDs for $ABEF$, so it is BCNF

Final schema (with keys boldened): **BC**, **ABD**, **ABEF**

Question 2

- i. $\text{TechCode} = \text{Proj}[\text{Code}](\text{Sel}[\text{Sector} = \text{"Technology"}](\text{Category}))$
 $\text{Answer} = \text{Proj}[\text{Name}](\text{TechCode Join Company})$

- ii. $\text{CodeGroup} = \text{GroupBy}[\text{Code}, \text{Count}[\text{Code}]](\text{Executive})$
 $\text{CodeGroup_2} = \text{Rename}[1 \rightarrow \text{Code}, 2 \rightarrow \text{Count}](\text{CodeGroup})$
 $\text{Answer} = \text{Proj}[\text{Code}](\text{Sel}[\text{Count} \geq 6](\text{CodeGroup_2}))$

- iii. $\text{NameGroup} = \text{GroupBy}[\text{Person}, \text{Count}[\text{Person}]](\text{Executive})$
 $\text{NameGroup_2} = \text{Rename}[1 \rightarrow \text{Person}, 2 \rightarrow \text{Count}](\text{NameGroup})$
 $\text{Answer} = \text{Proj}[\text{Person}](\text{Sel}[\text{Count} \geq 2](\text{NameGroup_2}))$

- iv. $\text{CatGroup} = \text{GroupBy}[\text{Industry}, \text{Count}[\text{Industry}]](\text{Category})$
 $\text{CatGroup_2} = \text{Rename}[1 \rightarrow \text{Industry}, 2 \rightarrow \text{Count}](\text{CatGroup})$
 $\text{Answer} = \text{Proj}[\text{Code}, \text{Industry}](\text{Sel}[\text{Count} = 1](\text{CatGroup_2 Join Category}))$

Question 3

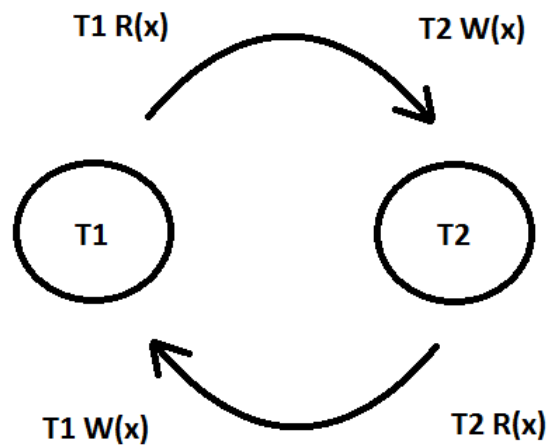
- i. Min: r , when $S \cap T$ is zero
 Max: $r + t$, when $T \subseteq S$, or $r + s$, when $S \subseteq T$

- ii. Min: 0, when no tuple meets the condition, c
 Max: $r * s$, when all tuples meet the condition, c

- iii. Min: 0, when $R \Join S = R$
 Max: r , when $R \Join S = \emptyset$

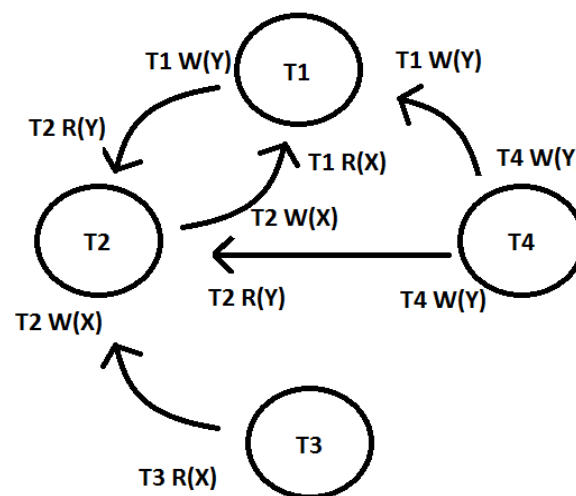
Question 4

i.



The graph has a cycle due to the edge from T1 to T2 and the edge from T2 to T1, hence it is not serializable.

ii.



There are conflicts that cause the cycle between:

T3 R(X) and T2 W(X),

T2 W(Y) and T1 W(Y),

T4 W(Y) and T2 R(Y),

T1 W(Y) and T2 R(Y),

T2 W(X) and T1 R(X)

Hence, the schedule is not serializable.