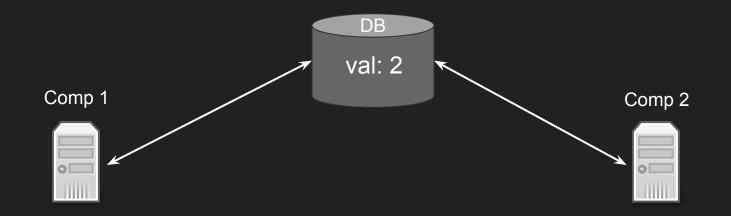
#### **Prelim**

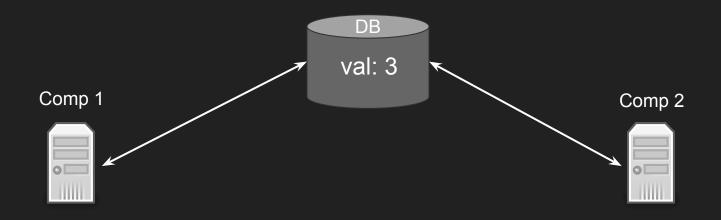
- Grades have been released on Gradescope
  - Median: 79
  - Mean: 76
  - o High: 100
- Significant error in the writing of 6.1 so it was not graded
  - We want to go over this question to help reinforce DP; still working on plan for that but will be before the final
- Use Gradescope to request regrades if you believe your submission was graded incorrectly
  - Original grader will re-grade first, can appeal to Ramin if you're still unsatisfied

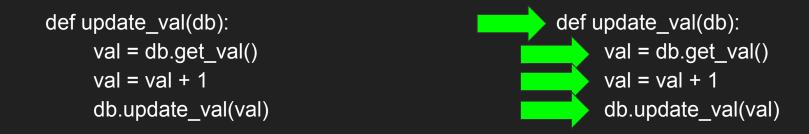
#### Administrivia

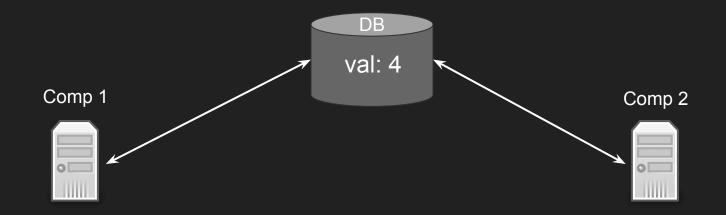
- Drop deadline blanket extended until Nov. 1
  - Additional extension is possible if necessary; please reach out to Ramin/Angy
- Ramin is away at a conference this week
  - Will hold extended office hours next week
  - Available for teleconference for urgent matters
- Deadline for finding HW2 partners is tonight









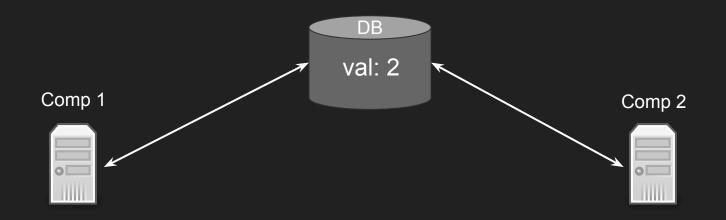


```
      def update_val(db):
      def update_val(db):

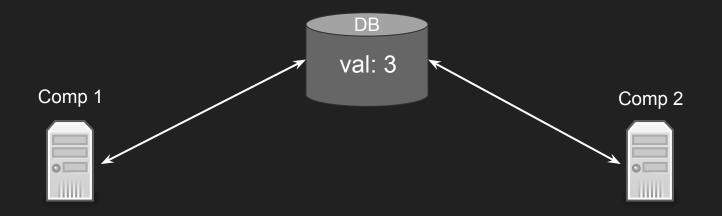
      val = db.get_val()
      val = db.get_val()

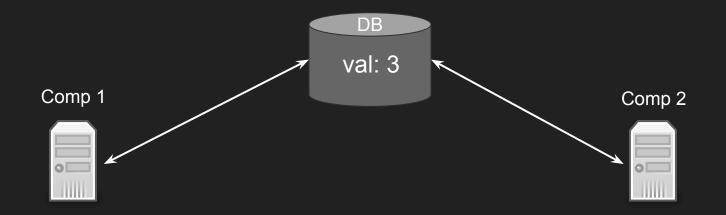
      val = val + 1
      val = val + 1

      db.update_val(val)
      db.update_val(val)
```









```
      def update_val(db):
      def update_val(db):

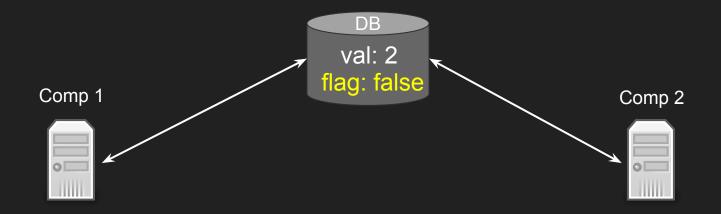
      val = db.get_val()
      val = db.get_val()

      val = val + 1
      val = val + 1

      db.update_val(val)
      db.update_val(val)
```

#### Race Condition

- When a system's behavior depends on timing, sequencing, or other uncontrollable events
- Bugs occurring from race conditions are often hard to pin down
- Nondeterministic "Heisenbug"
- Multiple threads accessing shared state is a classic place where race conditions happen.
- The "critical section" is the area where the shared state is accessed/updated

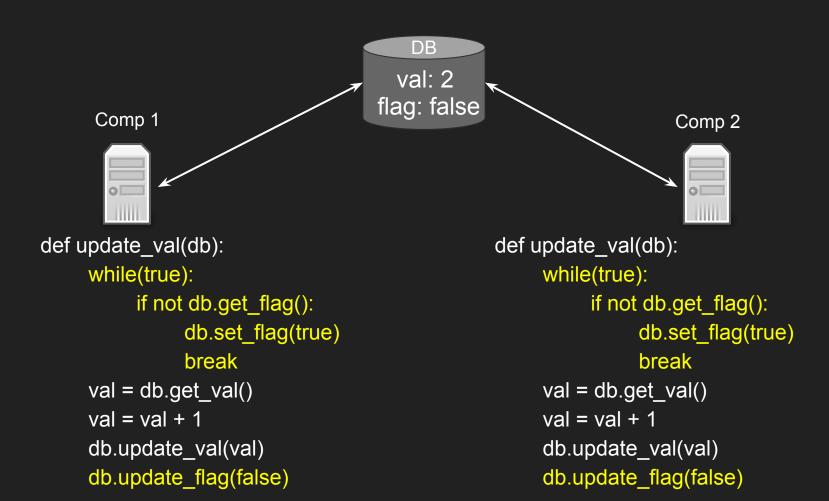


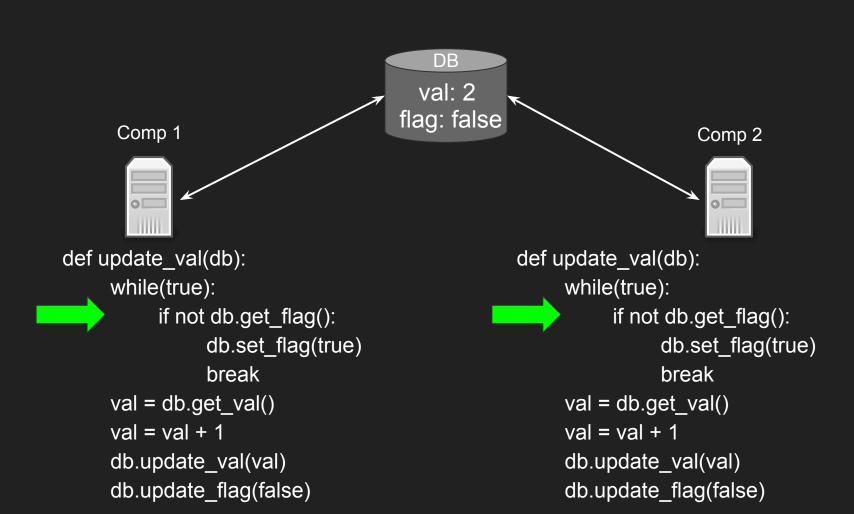
```
      def update_val(db):
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      val = db.get_val()
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      val = val + 1
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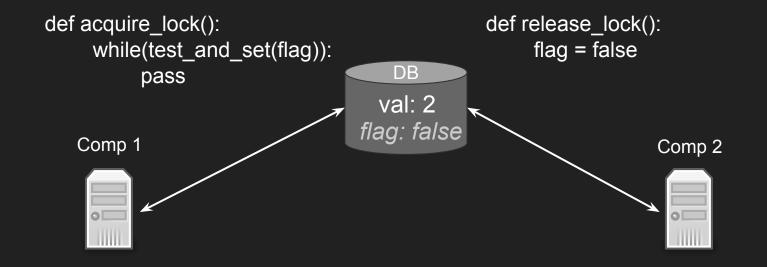
      db.update_val(val)
      db.update_val(val)
```





- There always seems to be a way to "interrupt" the flag checking/setting
- Need something "atomic"
  - Atomic == indivisible; guaranteed to all happen as one single operation
- test\_and\_set
  - Sets a boolean memory value to true and returns what the value WAS as one atomic operation
  - Can be used to create a spin lock

- Generally abstracted into another data structure called a "mutex"
  - Stands for "mutual exclusion"

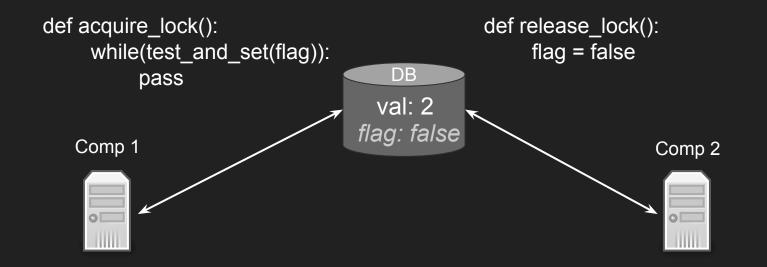


```
      def update_val(db):
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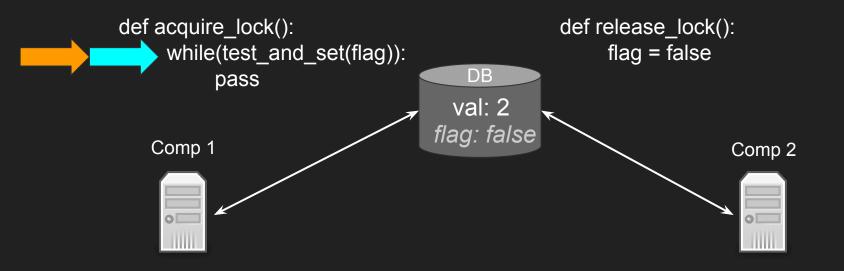
      val = db.get_val()
      val = db.get_val()

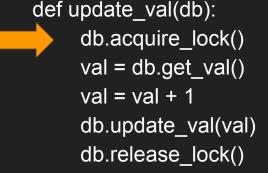
      val = val + 1
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      db.update_val(val)
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```

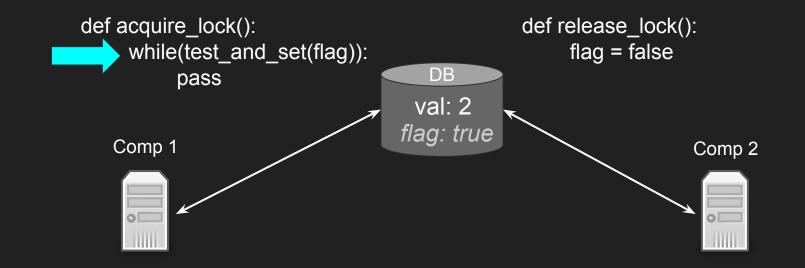


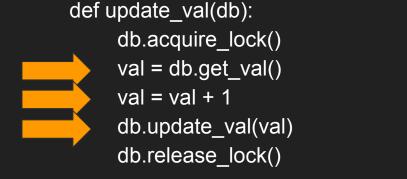
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def update_val(db):def update_val(db):db.acquire_lock()db.acquire_lock()val = db.get_val()val = db.get_val()val = val + 1val = val + 1db.update_val(val)db.update_val(val)db.release_lock()db.release_lock()
```



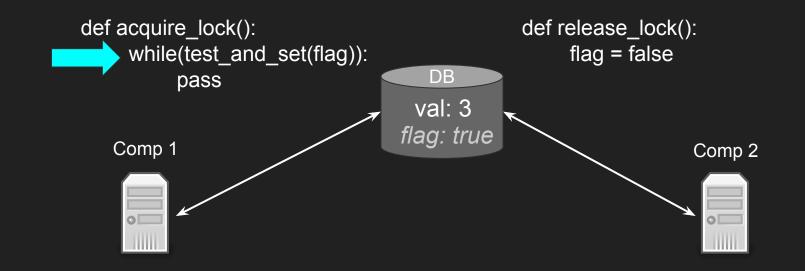


def update\_val(db):
 db.acquire\_lock()
 val = db.get\_val()
 val = val + 1
 db.update\_val(val)
 db.release\_lock()



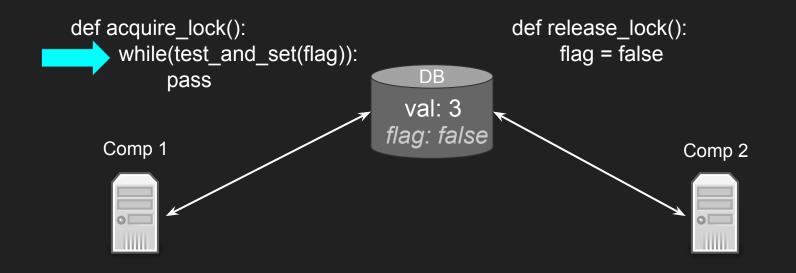


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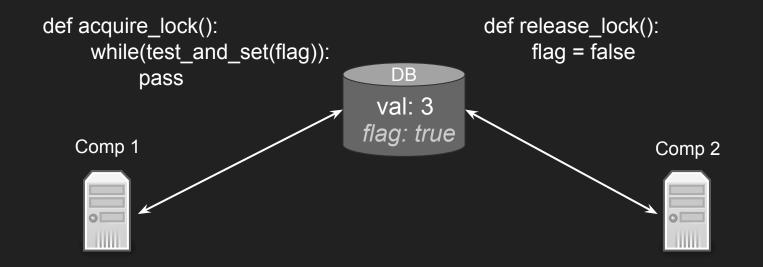


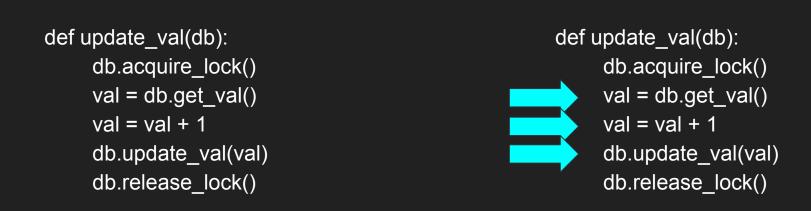
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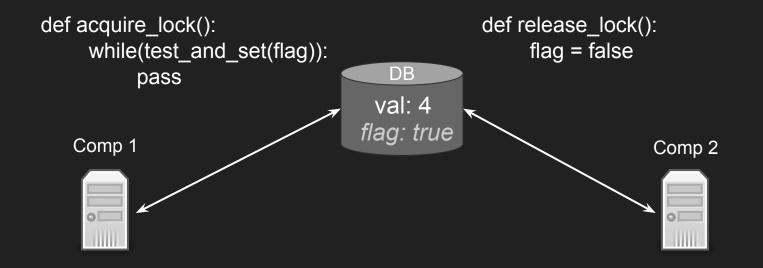
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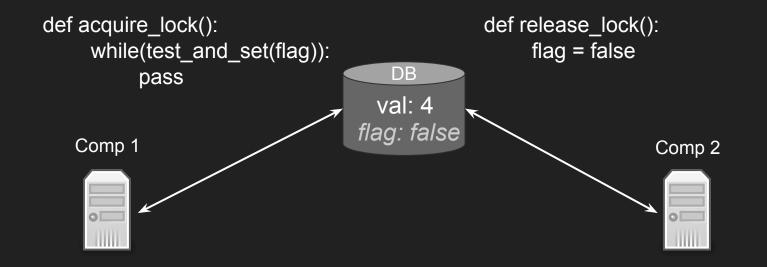
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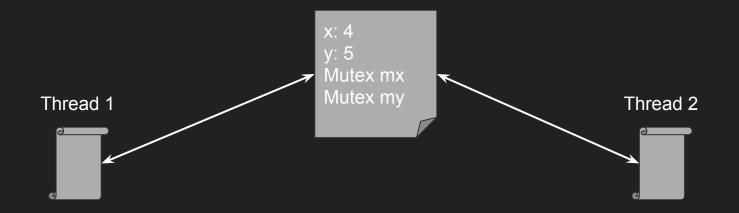




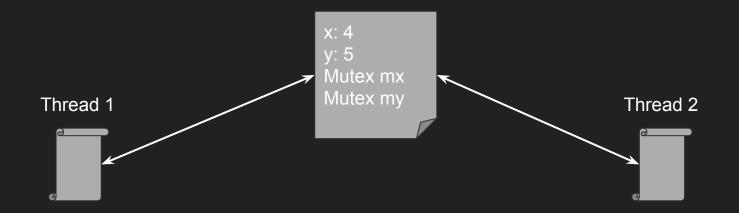




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```



```
def update_x():def update_x():mx.acquire_lock()my.acquire_lock()my.acquire_lock()mx.acquire_lock()x = x + yx = x + ymy.release_lock()mx.release_lock()mx.release_lock()my.release_lock()
```



```
def update_x():

mx.acquire_lock()

my.acquire_lock()

my.acquire_lock()

x = x + y

my.release_lock()

mx.release_lock()

mx.release_lock()

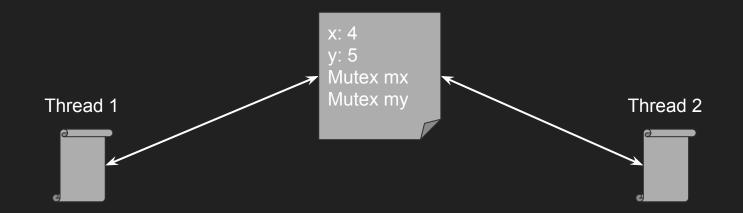
my.acquire_lock()

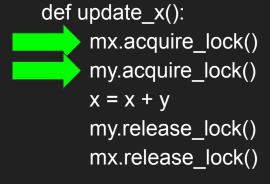
mx.acquire_lock()

x = x + y

my.release_lock()

mx.release_lock()
```





def update\_x():

my.acquire\_lock()

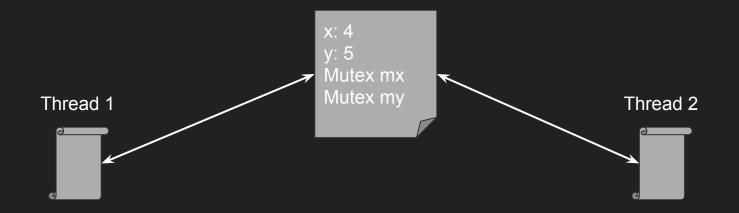
mx.acquire\_lock()

x = x + y

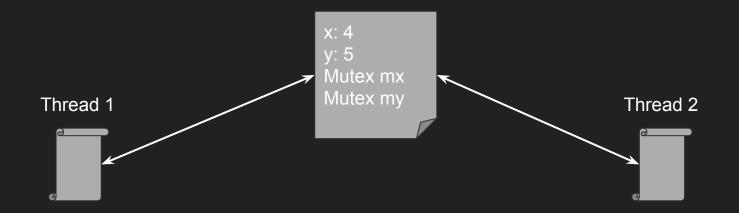
mx.release\_lock()

my.release\_lock()

- Mutual Exclusion can be the cause of <u>deadlocks</u>
  - Deadlock is when threads are not progressing because they're circularly waiting on each other
- Deadlock requires four things:
  - Mutual exclusion (locks around a critical section)
  - Resource holding (already having lock for one resource, and holding it while requesting another)
  - No pre-emption (no way of breaking locks early)
  - Circular wait (thread 1 is waiting on thread 2, 2 on 3, 3 on 4, etc.... n on 1)
- Fixing deadlock requires mitigating one of those four problems
  - In practice, the most straight-forward is avoiding circular wait by ordering the locks.
  - "Ostrich Algorithm" just hope it doesn't happen



```
def update_x():def update_x():mx.acquire_lock()my.acquire_lock()my.acquire_lock()mx.acquire_lock()x = x + yx = x + ymy.release_lock()mx.release_lock()mx.release_lock()my.release_lock()
```



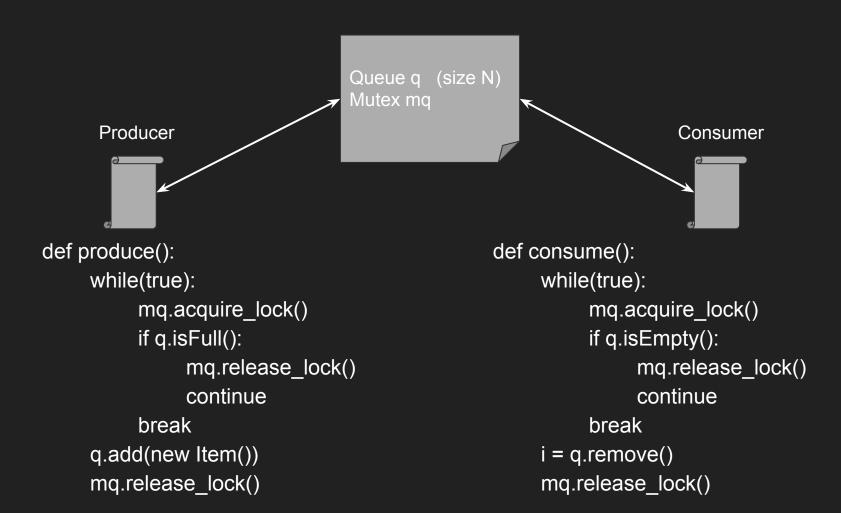
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def update_x():def update_x():mx.acquire_lock()mx.acquire_lock()my.acquire_lock()my.acquire_lock()x = x + yx = x + ymy.release_lock()my.release_lock()mx.release_lock()mx.release_lock()
```

#### Resource Starvation

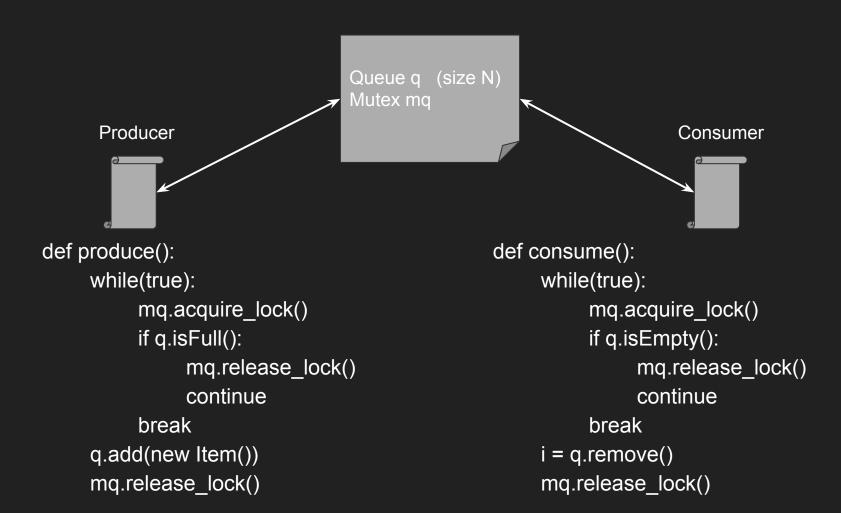
- Many threads may be waiting on a lock, which gets the lock next?
- If unlucky, some threads may never run at all
- Scheduling, age prioritization, etc. can help fix this

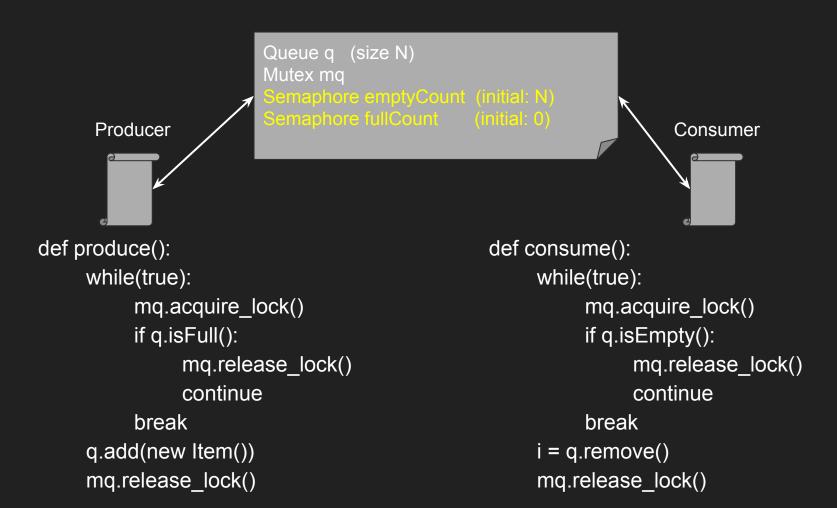
#### Livelock

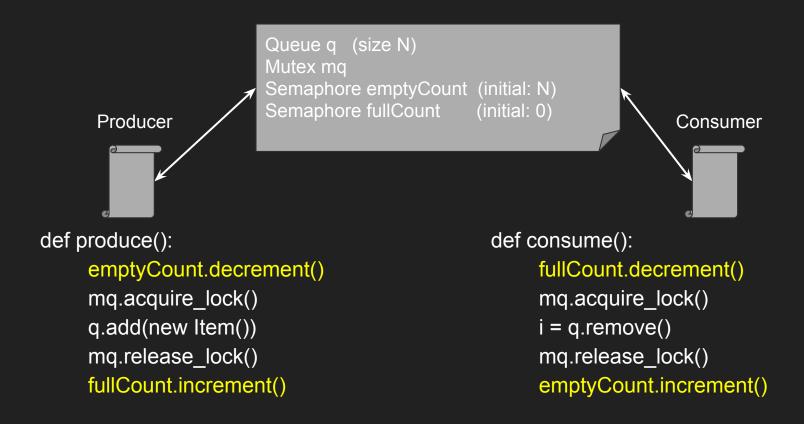
- Elements of the system are changing, but none are making any practical progress.
- Can be a risk for deadlock-detection systems



- Threads could spend a lot of time spin-waiting for the queue to be in the right state
  - Waste of CPU resources!
- Solution: semaphores
  - Semaphores are essentially a "count" of how much resource is available
  - Using a resource decrements the count, releasing it increments the count
  - If no resources are available, a decrement just waits
  - A mutex is essentially a semaphore with 1 resource count







- Semaphores are just a counter!
  - If implemented naively, could end up like val in the first example
  - Semaphore increment/decrement also needs to happen atomically
- Alternate solution: Condition Variables
  - Condition variables have two methods:
    - wait: release mutex and sleep until awoken
    - notify: wake up a sleeping thread and allow them to automatically re-take the mutex
  - The combination of a condition variable and a mutex is called a *monitor*

