

$$\begin{aligned} \max \quad & c^T x \\ \text{s.t.} \quad & Ax = b \\ & x \geq 0 \end{aligned}$$

Let x and y be optimal solutions to the problem. Prove or disprove the following statements:

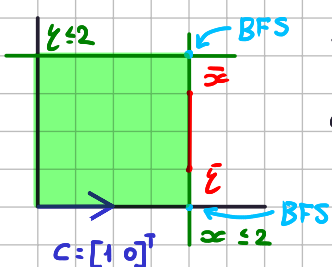
i. Any convex combination of x and y is also an optimal solution of the linear program. Recall that a convex combination of two points x and y is any point $z = \alpha x + (1 - \alpha)y$ for some $\alpha \in [0, 1]$

ii. There exists an optimal solution z on the line through x and y which is also basic feasible solution.

i) Sia γ l'ottimo $\Rightarrow c^T x = c^T y = \gamma$. Sia $z = \alpha x + (1 - \alpha)y$

$$c^T(\alpha x + (1 - \alpha)y) = \alpha c^T x + (1 - \alpha)c^T y = \alpha \gamma + (1 - \alpha)\gamma = \gamma \Rightarrow c^T z = \gamma \Rightarrow z \text{ e' soluzione ottimale.}$$

ii) Non e' necessariamente vero, se x oppure y sono una BFS, allora il segmento di linea comprende uno dei 2, se invece non sono BFS, sul segmento di linea non ci sara' una BFS (vertice). Il controesempio e':



sul segmento in rosso non c'e' una BFS.

$$LP: \begin{cases} \max & x \\ \text{s.t.} & x \leq 2 \\ & y \geq 2 \\ & x, y \geq 0 \end{cases}$$

$$\bar{x} = [2, \frac{3}{2}]^T \quad BFS = \left\{ \begin{bmatrix} 2 \\ 0 \end{bmatrix}, \begin{bmatrix} 2 \\ 2 \end{bmatrix} \right\}$$

$$\bar{y} = [2, \frac{1}{2}]^T$$

$$L = \alpha \bar{x} + (1 - \alpha) \bar{y} = \left\{ \begin{bmatrix} x \\ y \end{bmatrix} : x = 2 \wedge y \in [\frac{1}{2}, \frac{3}{2}] \right\}$$

$$\Rightarrow L \cap BFS = \emptyset.$$

Exercise 2: consider the primal problem (P)

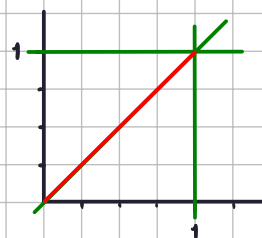
$$\begin{aligned} \max \quad & -p_1 x + p_2 y \\ \text{s.t.} \quad & x - y = 0 \\ & 0 \leq x \leq 1 \\ & 0 \leq y \leq 1 \end{aligned}$$

State the dual problem (D) and discuss the range of optimal feasible solutions to (P) and (D) in each of these cases:

- * $p_1 < p_2$
- * $p_1 = p_2$
- * $p_2 < p_1$

Il problema duale e':

$$(D): \begin{cases} \min & \gamma_2 + \gamma_3 \\ & \gamma_1 + \gamma_2 \geq -p_1 \\ & -\gamma_1 + \gamma_3 \geq p_2 \\ & \gamma_1 \in \mathbb{R} \\ & \gamma_2, \gamma_3 \geq 0 \end{cases}$$



e' la retta in rosso.

CASO $p_1 = p_2$ Il poliedro del problema (P) e' il seguente:

Se $p_1 = p_2$, il vettore $c = [-p, p]^T$ e' parallelo alla retta delle sol. ammissibili, e tutte le sol. sono ottimali con ottimo = 0, per il teo della dualita'

Forte, anche l'ottimo di (D) e' 0 $\Rightarrow \gamma_2, \gamma_3 = 0 \Rightarrow \begin{cases} \gamma_1 \geq -p_1 = -p_2 \\ \gamma_1 \leq p_2 \end{cases} \Rightarrow \gamma_1 = -p_1 \Rightarrow$ l'unica soluzione di (D) e' $\gamma^* = [p_1 \ 0 \ 0]^T$.

CASO $p_1 > p_2$ in tal caso, la retta parallela a $c = [-p_1, p_2]^T$ interseca la retta delle soluzioni solo se passante per l'origine, per ogni $(x, y) \neq (0, 0) \Rightarrow c^T \begin{bmatrix} x \\ y \end{bmatrix} < 0 \Rightarrow$ l'ottimo e' in $[0, 0]^T$. Riguardo (D), valgono considerazioni simili:

$$\text{l'ottimo di (D) e' } 0 \Rightarrow \gamma_2, \gamma_3 = 0 \Rightarrow \begin{cases} \gamma_1 \geq -p_1 \\ -\gamma_1 \geq p_2 \end{cases} \Rightarrow \begin{cases} \gamma_1 \geq -p_1 \\ \gamma_1 \leq -p_2 \end{cases} \Rightarrow \gamma_1 \in [-p_2, -p_1] \Rightarrow$$

$$\text{l'insieme delle sol. ottimali e' } \left\{ \begin{bmatrix} \kappa \\ 0 \\ 0 \end{bmatrix} \in \mathbb{R}^3 : \kappa \in [-p_2, -p_1] \right\}$$

CASO $p_1 < p_2$ in tal caso, per (P) la sol ottimale e' $[1, 1]^T$ perche'

$$C^T x = -p_1 + p_2 = (p_2 - p_1) > 0 \Rightarrow \text{l'ottimo di (D) e' } p_2 - p_1 \Rightarrow \zeta_2 + \zeta_3 = p_2 - p_1$$

\Rightarrow L'insieme delle sol. ottimali e' dato da $\zeta_1, \zeta_2, \zeta_3$ t.c.

$$\begin{cases} \zeta_1 + \zeta_2 \geq -p_1 \\ -\zeta_1 + \zeta_3 \geq p_2 \\ \zeta_2 + \zeta_3 = p_2 - p_1 \end{cases} \Rightarrow \begin{cases} \zeta_1 \geq \zeta_3 - p_2 \\ \zeta_3 \geq \zeta_1 + p_1 \\ \zeta_2 = p_2 - p_1 - \zeta_3 \end{cases}$$

Exercise 3: give an example of a linear program with a degenerate pivot. Include a justification of why the pivot is degenerate.

Si consideri

$$\begin{cases} \max x_2 \\ x_1 + x_2 \leq 2 \\ x_1 \leq 1 \\ x_2 \leq 1 \\ x \geq 0 \end{cases} \equiv \begin{cases} \max x_2 \\ x_1 + x_2 + x_3 = 2 \\ x_1 + x_4 = 1 \\ x_2 + x_5 = 1 \\ x \geq 0 \end{cases} \Rightarrow \text{la sol. } x = [0 \ 1 \ 0 \ 0 \ 0]^T \text{ e' degenera perche' ha} \\ \text{delle variabili di base uguali a zero.}$$

Exercise 4: solve the following linear program with the simplex method

$$\begin{aligned} \max & 2x_1 + x_2 - x_3 \\ & x_1 + 2x_2 + x_3 \leq 8 \\ & -x_1 + x_2 - 2x_3 \leq 4 \\ & x_1, x_2, x_3 \geq 0 \end{aligned}$$

Aggiungo le slack variable e comincio con $B = \{4, 5\}$

$B = \{1, 5\}$ e' ottimale

$$\begin{cases} \zeta_4 = -x_1 - 2x_2 - x_3 + 8 \\ \zeta_5 = x_1 - x_2 + 2x_3 + 4 \\ z = 2x_1 + x_2 - x_3 \end{cases} \Rightarrow \begin{cases} x_1 = -\zeta_4 - 2x_2 - x_3 + 8 \\ \zeta_5 = -\zeta_4 - 3x_2 + x_3 + 12 \\ z = -2\zeta_4 - 3x_2 - 3x_3 + 16 \end{cases} \Rightarrow \text{la sol.} \\ \text{ottimale} \\ \text{al problema} \\ \text{originale e' } x^* = \begin{bmatrix} 8 \\ 0 \\ 0 \end{bmatrix}$$

Exercise 5: prove or disprove the following statement

If a linear programming problem has a unique solution, then the dual also has a unique optimal solution.

E' falso, si consideri

$$P: \begin{cases} \max x_1 \\ x_1 \leq 1 \\ x_2 \leq 1 \\ x \geq 0 \end{cases} \quad D: \begin{cases} \min \zeta_1 + \zeta_2 \\ \zeta_1 \geq 1 \\ \zeta \geq 0 \end{cases}$$

\uparrow infinite soluzioni

\uparrow unica sol. in $\zeta = [1 \ 0]^T$