## Connect 4 Coursework Report

#### Introduction

The purpose of this coursework was to create a working implementation of the game Connect Four in any programming language, this could only be done through a command line console, i.e. GUI could not be used. For this task, I have chosen to implement the game in C# as my work environment makes use of C#. Moreover, I found that it would be a good opportunity to further my knowledge of the language. The implemented game is 5 columns by 4 rows, similar to the traditional 7x6 game by HASBRO¹, and consists of the option to play against another human player; human player against an "AI"; or an "AI" against "AI". The player has the option to undo a move that has just been entered, in case a better move is observed. Once the move is entered, the computer selects randomly a feasible move and the control of the game returns to the player until either the game is won or tied.

### Game Design

To implement the game there are some basic structures and functions that need to be considered:

- 1) The board is a 5x4 2D-array which needs to be updated at each stage. In its simplest form, the board states are stored in a *Queue* and the moves are stored in a *Stack*.
- 2) An initial screen output is needed to describe the game and gives users the options to play against the computer or another player.
- 3) To show the current state of play, the 2D-array has to be displayed on the console. A PrintBoard() method was implemented for this task.
- 4) A series of methods are then needed to identify which player is making a move; if the player's move is allowed; and check if the player has won, or if the board is full to declare a tie.

In what follows the main methods of the code are described in more detail.

PrintBoard() – this method prints the game board by iterating through a nested *for* loop. To display the state of play, 0 identifies an empty slot; 1 corresponds to a tile of player 1; and 2 corresponds to a tile of player 2. The main method initialises the array to all 0 values. The outer loop iterates through the rows and the inner loop iterates through the columns of the 2D-array.

<sup>&</sup>lt;sup>1</sup> https://shop.hasbro.com/en-us/product/connect-4-game:80FB5BCA-5056-9047-F5F4-5EB5DF88DAF4

To help the player identify the move, a column with the row number is added to the left of the array, and a row with the column number is included both above and below the board. An example is shown in Figure A1 in the appendix.

GetPosition() – this method sets the corresponding row to the user's column input, and only returns the row value if the input is within the board boundaries. This is achieved through an *if* condition, otherwise the methods returns a 0 value. There are two supplementary checks in *place: a try-catch that validates that the input is numerical, and an if that validates that* GetPosition() is greater than 0.

PvP() – PvPC() – PCvPC() methods are called in the main program when either "h", "c" or "a" (respectively) is entered by the user at the start of the game. These functions pool together all the other methods in order to create the game that is then played. All the methods make use of a *while* loop that is only exited once a game has either been won or tied. In order to switch between players, a Boolean is used, which is initiated as player = true and then flips to false when it is the next player's turn, and vice versa.

GameWon() – this is implemented through an iterative algorithm, consisting of nested *for* loops with *if* conditionals. The section with this winning *if* condition is shown in Figure A2 in the appendix. Figure 1 below shows the possible winning outcomes in the 4x4 sub-array. These are horizontal, vertical, and ascending and descending diagonal chains of 4 tiles.

1				X	1	X				1			X		1				
2			X		2		X			2			X		2	X	X	X	X
3		X			3			X		3			X		3				
4	X				4				X	4			X		4				
	1	2	3	4		1	2	3	4		1	2	3	4		1	2	3	4

Figure 1. Four possible winning chains identified by GameWon().

TotalMoves() – this was originally implemented through an iterative algorithm, consisting of nested *for* loops with an *if* conditional that checks whether there was an empty space i.e. a 0. This has been reimplemented as to counting the total number of moves using the Replay. *Count* value and comparing it to the size of the board.

PCPlayer() – This method uses a *while* loop to generate a random integer between 1 and 5 and checks for a valid move using GetPosition(). The feasible column number generated is used by PvPC() and PCvPC() to create a very simplistic AI player.

MoveUndo() – this method is implemented using an undo *Stack* data structure as with the last-in first-out feature of *Stack* structures it is relatively easy to go back one or more moves. To be able to replay the same sequence as the moves are removed using *Pop*(), a replay *Stack* is used to store these moves. One could also opt to use *pointer* variables of the entire board at each step and implement the undo method replacing the entire 2D-array by previous values. The use of stacks is more efficient given that only one array element is changed.

GameReplay()—this method allows the user to review the moves of the game that has just been completed by making use of a *Queue* data structure. Each time a move is made the new board state is enqueued into the Replay *Queue* and once the game has been finished the user is given the option to review the game, this is done with a *for* loop that until Replay. *Count* is 0 will *Dequeue* and print the board.

GameMoves{} – this class is a *Struct* data structure, which when used in conjunction with *Push*() for the *Stack* and *Enqueue*() for the *Queue*, retrieves and stores the most recent move made by the player into the Undo *Stack* and the board state into the Replay *Queue* via getters and setters, storing three values: the player's column input and corresponding row value, and which player made the move.

## **Evaluation**

Implementing the game board via the use of a multi-dimensional array data structure, a 2D-array for this task, made it straightforward to write code that is easier to read compared to a one dimensional array, which would need a conversion of rows/columns into a position [(RowN-1)\*NCol+ColN] at each occurrence. The 2D-array provides a one to one correspondence with the board and is easy to visualise. This helps also when considering how to implement the win algorithm, which was the first major task to resolve.

The win algorithm that was used for this task presents both advantages and disadvantages. An advantage is that the logic implemented will work for the winning four in a row regardless of the game board size and player position, due to the *if* conditional that checks for the 4 possible winning moves shown Figure 1. However, as it is a nested *for* loop which runs in quadratic time, as the board size gets bigger, it will take longer to eventually return true for a winning conditional being met (as shown in Table A1 in the appendix). This became apparent at the evaluation stage as the execution time on the 5x4 array is very fast. A more efficient option would be to keep track of possible chains originating from the current move, but rather than run on the entire 2D-array, it would only check a sub-array  $\pm 3$  positions, with respect to the

current move position, in both rows and columns limiting the maximum cycle to a 7x7 array for larger boards.

A similar issue was found in case of a game that has been tied, with the execution time being  $O(n^2)$ . Eventually, it would become near impossible to play the game due to the ever-increasing time taken for the methods to finish looping, and potentially could take up too many resources - as even a board of only 1000x1000 averaged at ~15% CPU usage. Another thing to take into consideration is that as the board size becomes bigger, the chance of a tied game becomes smaller, meaning that looping through GameTied() would have been a significant waste of resources and this is the reason why this was changed to simply tracking the total number of moves.

The simple way the "AI" player was implemented works relatively well with a smaller board and keeps the game interesting for both the player vs AI and AI vs AI games, and it is the reason why the board size 5x4 was chosen. Nevertheless, this could have been executed in a smarter way, for example through the use of a Tree data structure, which could map out all the possible moves and would then select the most efficient path towards a winning conditional. Another possible defensive approach could have been implemented to check if a move would allow to break the opponent's chains of 2 tiles and then block 3 tiles in a row.

Due to the way that the undo feature was implemented – only allowing the player to undo the move they have just made, meaning only one undo move per player – meant that it made adding a redo function unnecessary as once a move is undone, the player is asked for a new input. In order to add a useful redo function, the undo method would have to be implemented in a way that allows the player to undo not only their own moves, but also the moves of the other player, bringing the game to a standstill. This execution however seems less fair in comparison to the way that it has been implemented in the code.

The table shown in the appendix is the result of tests done using the C# System. Diagnostics Stopwatch feature which gave a reliable execution time result compared to attempting to manually time results. An example test: setting the board size to 10,000x10,000 resulted in the executable taking 1hr-1-mins-24s (3684404ms) to print just one instance of the board, whereas with the normal 5x4 it only takes 2ms.

# **Appendix**

https://github.com/CatB1794/Brandani\_Caterina\_ADS/tree/master/WinGifs

```
if (Board[x, y] == move && Board[x + 1, y - 1] == move && Board[x + 2, y - 2] == move && Board[x + 3, y - 3] == move)
```

Figure A2 - Ascending if condition

52, 60, 3 51, 59, 3 694, 44.729, <b>Board si</b> 23, 18, 2 23, 17, 2 6.591, 5.145 <b>Board si</b> Mo 18, 11, 2	ize: 50x50 oves 27, 26, 20 27, 26, 20	V	Vin:Lose:Tie ratio 4:1:0 1:4:0  Vin:Lose:Tie ratio 1:4:0 4:1:0  Vin:Lose:Tie ratio 2:3:0			
51, 59, 3 594, 44.729, <b>Board si</b> Mo 23, 18, 2 23, 17, 2 6.591, 5.145 <b>Board si</b> Mo 18, 11, 2	38, 52, 44 38.613 ize: 50x50 oves 27, 26, 20 27, 26, 20 5 ize: 25x25 oves 20, 16, 27		1:4:0  Vin:Lose:Tie ratio  1:4:0  4:1:0  Vin:Lose:Tie ratio  2:3:0			
Board si  23, 18, 2 23, 17, 2 6.591, 5.143  Board si  Mo  18, 11, 2	38.613 ize: 50x50 oves 27, 26, 20 27, 26, 20 5 ize: 25x25 oves 20, 16, 27		Vin:Lose:Tie ratio 1:4:0 4:1:0  Vin:Lose:Tie ratio 2:3:0			
Board si  Mo  23, 18, 2  23, 17, 2  6.591, 5.145  Board si  Mo  18, 11, 2	ize: 50x50  oves  27, 26, 20  27, 26, 20  5  ize: 25x25  oves  20, 16, 27		1:4:0 4:1:0 Vin:Lose:Tie ratio 2:3:0			
23, 18, 2 23, 17, 2 6.591, 5.145 <b>Board si</b> Mo 18, 11, 2	27, 26, 20 27, 26, 20 5 ize: 25x25 oves 20, 16, 27		1:4:0 4:1:0 Vin:Lose:Tie ratio 2:3:0			
23, 18, 2 23, 17, 2 6.591, 5.145 <b>Board si</b> Mo 18, 11, 2	27, 26, 20 27, 26, 20 5 ize: 25x25 oves 20, 16, 27		1:4:0 4:1:0 Vin:Lose:Tie ratio 2:3:0			
23, 17, 2 6.591, 5.145 <b>Board si</b> Mo 18, 11, 2	27, 26, 20 5 ize: 25x25 oves 20, 16, 27	V	4:1:0  Vin:Lose:Tie ratio 2:3:0			
6.591, 5.145 <b>Board si</b> Mo  18, 11, 2	5 ize: 25x25 oves 20, 16, 27	V	Vin:Lose:Tie ratio 2:3:0			
<b>Board si</b> Mo 18, 11, 2	ize: 25x25 oves 20, 16, 27	V	2:3:0			
Mo 18, 11, 2	oves 20, 16, 27	V	2:3:0			
18, 11, 2	20, 16, 27	V	2:3:0			
18, 11, 2	20 15, 26					
	•	3:2:0				
1.075, 1.844	4					
Board	size: 5x5					
Mo	oves	V	Vin:Lose:Tie ratio			
9, 12, 11, 8, 13			2:2:1			
9, 11, 10, 8, 12			2:2:1			
0.068, 0.100	6					
Moves (1)	Average Mov	ves (2)	Average time (s)			
.6	48.8		42.5268			
.8	22.6		5.7862			
.4	18		1.2642			
.6	10		0.0872			
	9, 12, 1 9, 11, 1 0.068, 0.10 <b>Moves (1)</b> 6 8	9, 12, 11, 8, 13 9, 11, 10, 8, 12 0.068, 0.106  Moves (1) Average Mov 6 48.8 8 22.6 4 18 6 10	9, 12, 11, 8, 13 9, 11, 10, 8, 12  0.068, 0.106  Moves (1)			

Table of results