# Blocking Point-to-Point Communications

### Types of Communication

- In MPI there are two main types of communication
- Point-to-point communication:
- A communication involving a pair of processes
- Good if a process only needs to communicate with a limited number of other processes

#### Collective communication:

- A communication involving all the processes
- May seem like point-to-point communication is all that is required as could use a set of them to communicate with everyone
  - Very inefficient Sending data from one process to every other process involves p-1 communications from the originating node
  - Collective communications can do things like relaying data Original process sends to a subset of the nodes, which, in term send the data on to other nodes that have not yet received it
  - Communication time scales as approximately log(p)

### Blocking point-to-point communication

- We will initially do blocking point-to-point communication
  - Easier to code and shows the basic idea
  - Later on we will look at non-blocking communications

#### Two main functions involved:

- MPI\_Send Sends data to another process
- MPI\_Recv Receives data from another process
- Because these are blocking, you must be careful of the send and receive order
  - If a process tries to send data without all the previous receives being handled or vice versa there will be a block in communication and the code will hang or crash
  - We will later look at using probe as a way around this problem so that messages can be received from different processes in any order (amongst other uses)

### Simple point-to-point communication

This program sends data from process zero to all other processes (a different random number sent to each

process):

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
#include <time.h>
using namespace std;
int id, p;
int main(int argc, char *argv[])
     MPI Init(&argc, &argv);
     MPI Comm rank(MPI COMM WORLD, &id);
     MPI Comm size(MPI COMM WORLD, &p);
     srand(time(NULL)+id*10);
     int tag num = 1;
```

```
if (id == 0)
     for (int i = 1; i < p; i++)
          int send data = rand();
           MPI Send(&send_data, 1, MPI_INT, i, tag_num, MPI_COMM_WORLD);
          cout << send data << " sent to processor " << i << endl;</pre>
          cout.flush();
else
     int recv data;
     MPI Recv(&recv data, 1, MPI INT, 0, tag num, MPI COMM WORLD,
          MPI STATUS IGNORE);
     cout << recv data << " received on processor " << id << endl;</pre>
     cout.flush();
MPI Finalize();
```

### Random Numbers

- In a lot of the examples I use random numbers to give me something to send
  - rand generates an integer random number between 0 and RAND\_MAX
  - srand seeds the random number generator
  - Both of these functions are found in <cstdlib>
- To seed the random number I use srand(time(NULL)+id\*10);
  - time(NULL) gives a time as the number of seconds since the beginning of 1970
  - As all the processes will start at the same time, I try and have them generate different random number by adding their id times 10 – These are still only pseudo-random and so are likely to be related to one another

### MPI\_Send

```
int MPI_Send(
   void* data,
   int count,
   MPI_Datatype datatype,
   int destination,
   int tag,
   MPI_Comm communicator)
```

- data is a pointer to the data to be sent
- count is the number of items to be sent
  - Note that this is the number of variables of type datatype to be sent, not the number of bytes
- datatype is the MPI data type to be sent
- destination is the id of process that is to receive the data
- tag is an identifier for the communication
- For both MPI\_Send and MPI\_Recv the return value indicates error or success

### MPI\_Recv

```
int MPI_ Recv(
   void* data,
   int count,
   MPI_Datatype datatype,
   int source,
   int tag,
   MPI_Comm communicator,
   MPI_Status* status)
```

- data is a pointer to the data to be received
- count is the number of items to be received
- datatype is the MPI data type to be received
- source is the id of process from which the data is to be received
- tag is an identifier for the communication
- status is a pointer to a structure that contains information about the communication
  - Stores information such as the processes involved and the size of the data being sent, as well as any communication error information
  - To ignore use MPI\_STATUS\_IGNORE

# MPI Types and their C/C++ equivalents

MPI_SHORT	short int
MPI_INT	int
MPI_LONG	long int
MPI_LONG_LONG	long long int
MPI_UNSIGNED_CHAR	unsigned char
MPI_UNSIGNED_SHORT	unsigned short int
MPI_UNSIGNED	unsigned int
MPI_UNSIGNED_LONG	unsigned long int
MPI_UNSIGNED_LONG_LONG	unsigned long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_BYTE	char

- These are the simple MPI\_Datatypes available
  - Later we will look at creating our own MPI data types

### More on MPI data types

- Note that MPI\_BYTE should be used for send generic data of a known size in bytes
  - E.g. if Object is a complex object of fixed size (watch out for objects, for instance, containing pointers!):

```
MPI_Send(&Object, sizeof(Object), MPI_BYTE, 0, 0, MPI_COMM_WORLD);
```

- MPI\_BYTE should also be used to send variables of type bool
  - bool may only represent true or false (0 or 1), but it is stored in a variable of one byte in size as this is the smallest addressable memory unit

### More on tags

- Data can potentially be received out of the expected order
  - Especially true in non-blocking communications
  - Need a method to ensure that the correct communication is being dealt with
  - A "race condition" refers to the situation where communications are undesirably handled in the wrong order
- Each communication is given a tag
- The tag should identify the specific type and/or order of communication
  - Either increment the tag for each communication (what I usually do)
  - ...or have a different tag for each type of communication (though this might run the risk of being out by a whole cycle of communications)
  - ...or do a combination by, for instance, having a unique identifier for each communication with a value between 0 and 100 and then add 100 times the number of communication cycles to this value to produce the tag
- Note that I don't change the tags in most of the examples because I am showing a single communication
- Tags can be any non-negative value below MPI\_TAG\_UB

### MPI return codes

 Most MPI communications give a return value that indicates either success or what went wrong. These are some those codes:

#### MPI SUCCESS

Communication completed successfully

#### MPI\_ERR\_COMM

Invalid communicator

#### MPI\_ERR\_COUNT

Invalid count argument (usually a negative count – a zero is usually valid)

#### MPI\_ERR\_TYPE

• Invalid datatype argument (Either not a valid MPI type or, for types you created yourself, an type that has not been committed – more on this in a later lecture)

#### MPI ERR TAG

Invalid tag argument

#### MPI\_ERR\_RANK

Invalid source or destination rank. Ranks must be between zero and the size of the communicator minus one

### Using a Probe

- A probe allows information about a communication to be read before the communication is completed
  - i.e. before MPI\_Recv is called in the context of blocking communications
- Two very useful pieces of information can be obtained by doing this
  - Which process has sent the data
  - ...and how much data is being sent
    - You need to be able to assign enough memory to the data pointer that is receiving the data
      - You could assign more than the maximum that that will ever be sent, but this is inefficient and you need to know what is the maximum that could be send
    - If you wish to assign the correct amount of memory for the amount of data being sent you can either send a communication saying how much data to expect or you can use a probe to find out

### MPI\_Probe

• MPI\_Probe is the function that probes the communication:

```
int MPI_Probe(
   int source,
   int tag,
   MPI_Comm comm,
   MPI_Status *status)
```

- source is the source process
- tag for the communication
- comm is the communicator
- status is a pointer to the structure containing the status information
- Having to specify the source and the tag does seem to remove some of the utility of the function, but MPI\_ANY\_SOURCE and/or MPI\_ANY\_TAG can be set

### Using the information from MPI\_Status

```
typedef struct {
    int count;
    int cancelled;
    int MPI_SOURCE;
    int MPI_TAG;
    int MPI_ERROR;
} MPI_Status;
```

- count is the number of received entries
  - Use MPI\_Get\_count in conjunction with the status rather than using this variable directly – avoids potential issues with mismatches in variable type or size
- cancelled is true or false depending on whether the corresponding communication was cancelled
- MPI\_SOURCE is the source process for the communication
- MPI\_TAG is the tag associated with the communication
- MPI\_ERROR is an error code (there are a lot of potential errors and you can look them up if you want)
  - Has a value of MPI\_SUCCESS if there is no error

### MPI\_Get\_count

 MPI\_Get\_count is used in conjunction with the status to get the number of items (not bytes) sent:

```
int MPI_Get_count(
    const MPI_Status *status,
    MPI_Datatype datatype,
    int *count)
```

- status is a pointer to the status structure
- datatype of the data being sent
- count is a pointer to an integer that will store the number of items being sent
- Note that the return value is again an error code and not the number of items (which is stored in the variable pointed to by count)

### Using MPI\_Probe to get the size of the data

• Similar to the previous example, with processor zero sending data to each of the other processes – Difference is that it randomly sends between 1 and 5 items

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
#include <time.h>
using namespace std;
int id, p;
int main(int argc, char *argv[])
     MPI Init(&argc, &argv);
     MPI Comm rank(MPI COMM WORLD, &id);
     MPI Comm size(MPI COMM WORLD, &p);
     srand(time(NULL)+id*10);
     int tag num = 1;
     if (id == 0)
           for (int i = 1; i < p; i++)
                int num send = 1 + rand() \% 5;
                int *send data = new int[num send];
                for (int j = 0; j < num send; j++) send data[j] = rand();
                MPI Send(send data, num send, MPI INT, i, tag num, MPI COMM WORLD);
```

```
for (int j = 0; j < num send; j++) cout << send data[j] << "\t";
           cout << " sent to processor " << i << endl;
           cout.flush();
           delete[] send data;
else
     int *recv data;
     int num recv;
     MPI Status status;
     MPI Probe(0, tag num, MPI COMM WORLD, &status);
     MPI Get count(&status, MPI INT, &num recv);
     recv data = new int[num recv];
     MPI Recv(recv data, num recv, MPI INT, 0, tag num, MPI COMM WORLD,
           MPI STATUS IGNORE);
     for (int j = 0; j < num recv; j++) cout << recv data[j] << "\t";
     cout << " received on processor " << id << endl;
     cout.flush();
MPI Finalize();
```

### Using MPI\_Probe to get the source of the data

- MPI\_Probe can be used to both wait for data to be sent and to find out where it has come from
  - Note that MPI\_Probe can time out (you can check for this as it will give an error) and, if you wish to, probe again to continue waiting
  - Note that MPI\_Recv could directly wait for a message from any process, but it is often useful to prepare to receive data from a specific process
- The following example plays multi-processor ping-pong, with a processor receiving a communication and then passing it on to another random process
- I exit when one process has received 10 messages
  - As other processes don't know how many messages have been received by another process I exit messily using MPI\_Abort (likely to result in error messages from other processes).
- Processor zero starts the ping-pong
- You will notice that some of the outputs will be out of the logical order
  - The process sends it cout back to originating process and so may arrive back in a different order to which the processes generated them – Watch out for this problem when trying to debug parallel code

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
#include <time.h>
using namespace std;
int id, p;
int tag num = 1;
void Send Random Data(void)
     int to proc;
     while ((to proc = rand() \% p) == id); //Stop code sending to itself
     int send data = id;
     MPI Send(&send data, 1, MPI INT, to proc, tag num, MPI COMM WORLD);
     cout << "Processor " << id << " sent data to processor " << to proc << endl;
     cout.flush();
```

### MPI\_Probe example continued

```
void Recv_Data(void)
{
    int from_proc;
    int recv_data;
    MPI_Status status;

while (MPI_Probe(MPI_ANY_SOURCE, tag_num, MPI_COMM_WORLD, &status)
!= MPI_SUCCESS)
    cout << "Communication timed out or failed! " << id << endl;

from_proc = status.MPI_SOURCE;

MPI_Recv(&recv_data, 1, MPI_INT, from_proc, tag_num, MPI_COMM_WORLD, &status);

cout << "Processor" << id << " received data for processor" << from_proc << endl;
cout.flush();
}</pre>
```

```
int main(int argc, char *argv[])
     int recv cnt = 0;
      MPI Init(&argc, &argv);
      MPI_Comm_rank(MPI_COMM_WORLD, &id);
      MPI Comm size(MPI COMM WORLD, &p);
      srand(time(NULL)+id*10);
      if (id == 0) Send Random Data();
                                         //Starts the ping-pong
      while (true)
            Recv Data();
                                         //Waits for and receives data
            recv cnt++;
            if (recv cnt<10)
                                         //Sends it on to a new processor
                  Send Random Data();
            else break;
                                          //Exits if 10 communications are received
      cout << endl << "Process " << id << " has received 10 communications and is
            aborting (the exit will be messy!)" << endl;
      cout.flush();
      MPI Abort(MPI COMM WORLD, 0);
```

## What is wrong with blocking communications?

- The main problem with blocking communications is that it is hard to do them efficiently if a lot of communications is required
  - Either you can let each process communicate in turn
    - Lots of idle time waiting for other processes to finish their communications
  - ...or you can try and do clever ordering of the communications so that all processes get to do communications at the same time
    - Easy to get wrong Will block and crash if, for instance, two process are trying to send data to
      one another or receive data from one another at the same time
    - Still likely to be inefficient as even efficient ordering will rely on, for instance, an assumption that all communications take about the same time
- Things can be much more efficient if communications don't block one another
  - MPI provides the tools to do this
  - More complex to implement, but potentially much more efficient code