

A, B size n

\Downarrow

C size $2(n-1)+1 = 2n-2+1 = 2n-1$

$$P = \{x_0, x_1, \dots, x_{2(n-1)+1}\}$$

$$\{C(x_0), C(x_1), \dots, C(x_{2(n-1)+1})\} \Rightarrow \text{determines uniquely } C.$$

$$C(x_i) = A(x_i) B(x_i)$$

FFT-Based Multiplication

A, B in coef. form $\dots \rightarrow C = A \cdot B$ in coef. form

\uparrow
DFT⁻¹

\downarrow

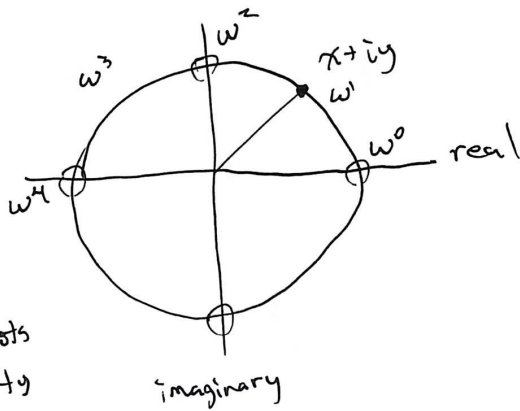
A, B in pt.-val form

\xrightarrow{DFT}

C in pt.-val form

	$A(x)$	$B(x)$	C
x_0	$A(x_0)$	$B(x_0)$	$C(x_0) = A(x_0)B(x_0)$
\vdots			
x_{2n-2}	$A(x_{2n-2})$		$C(x_{2n-2}) = A(x_{2n-2})B(x_{2n-2})$

DFT "Discrete Fourier Transform"



$$\omega^2 = \omega^1 \cdot \omega^1$$
$$\omega^3 = \omega^1 \cdot \omega^2$$

In general, " n^{th} roots of unity"
are the solutions to

$$\omega^n = 1 \quad \omega = 1 \text{ is a solution}$$

$$\omega_n^k = e^{i \frac{2\pi}{n} \cdot k}$$

Properties

$$1) \omega_{dn}^{dk} = \omega_n^k$$

$$\left(\begin{smallmatrix} n \text{ is a} \\ \text{power} \\ \text{of } 2 \end{smallmatrix} \right) \quad 2) \omega_n^{n/2} = \omega_2^1 = -1$$

3) Halving Lemma if $n > 0$ (even)

then the squares of the n^{th} roots of unity are the $(n/2)^{\text{th}}$ roots of unity.

$$0 \leq k < n/2 \quad (\omega_n^k)^2 = \omega_n^k \cdot \omega_n^k = \omega_n^{2k} = (\omega_n^2)^k = \omega_{n/2}^k$$

$$(\omega_n^{k+n/2})^2 = \omega_n^{(2k+n)} = \omega_n^{2k} \cdot \underbrace{\omega_n^n}_{=1} = \omega_n^{2k} = \omega_{n/2}^k$$

4) Summation Lemma

Assume $n \geq 1$ $k = \text{non-zero integer}$
not divisible by n .

$$\sum_{j=0}^{n-1} (\omega_n^k)^j = 0$$

pf

$$\sum_{j=0}^{n-1} (\omega_n^k)^j = \frac{(\omega_n^k)^n - 1}{\omega_n^k - 1} = 0$$

$\neq 0$

DFT: We want to evaluate $A(x)$ $x = \omega_n^0, \omega_n^1, \omega_n^2, \dots, \omega_n^{n-1}$

$$\text{Let } y = A(\omega_n^k).$$

$$\vec{y} = \text{DFT}_n(\vec{a})$$

$$\vec{a} = (a_0, a_1, \dots, a_{n-1})$$

$$\vec{y} = (y_0, y_1, \dots, y_{n-1})$$

$$A(x) = a_0 + a_1x + a_2x^2 + \dots + a_{n-1}x^{n-1}$$

$$= a_0 + a_2x^2 + a_4x^4 + \dots$$

$$+ a_1x + a_3x^3 + \dots$$

$$= A^{[0]}(x^2) + x \cdot A^{[1]}(x^2)$$

where $A^{[0]} = a_0 + a_2x + a_4x^2 + \dots$

$$A^{[1]} = a_1 + a_3x + a_5x^2 + \dots$$

A, B

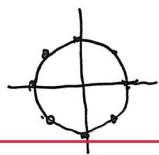
$(a_0, a_1, a_2, \dots, a_{n-1})$

$(b_0, b_1, b_2, \dots, b_{n-1})$

$$C = A \cdot B$$

$(c_0, c_1, \dots, c_{2(n-1)})$

FFT



pt. value form
(2n points)

pointwise
multiplication

C

pt. val
form

$$\omega_n^k = k^{\text{th}} \text{ } n^{\text{th}} \text{ root of unity} = e^{i \frac{2\pi}{n} k}$$

$$\omega_n^0 = 1$$

Summation Lemma

$n \geq 1$, k = non-zero integer not divisible by n .

$$\sum_{j=0}^{n-1} (\omega_n^k)^j = 0$$

DFT "discrete Fourier transform"

$$y_k = A(\omega_n^k)$$

$$\vec{y} = \text{DFT}_n(\vec{a})$$

$$A(x) = A^{[0]}(x^2) + x A^{[1]}(x^2)$$

RECURSIVE-DFT(a)

Let $n = \text{length}(a)$

if $(n_i = 1)$ return a

$$\omega_n = e^{i \frac{2\pi}{n}}$$

$\omega = 1$ (current root)

$$a^{[0]} = (a_0, a_2, a_4, \dots, a_{n-2})$$

$$a^{[1]} = (a_1, a_3, a_5, \dots, a_{n-1})$$

$$y^{[0]} = \text{RECURSIVE-DFT}(a^{[0]})$$

$$y^{[1]} = \text{RECURSIVE-DFT}(a^{[1]})$$

for $k = 0$ to $\frac{n}{2} - 1$ {

$$y_k = y_k^{[0]} + \omega \cdot y_k^{[1]}$$

$$y_{k+\frac{n}{2}} = y_k^{[0]} - \omega \cdot y_k^{[1]}$$

}

return y

running time

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$= \Theta(n \lg(n)) = \tilde{\Theta}(n)$$

"Inverse DFT"

We have

$$y_k = \sum_{j=0}^{n-1} a_j \cdot (\omega_n^k)^j$$

$$\vec{y} = V_n \cdot \vec{a}$$

$$V_n = \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \omega_n & \omega_n^2 & \dots & \omega_n^{n-1} \\ 1 & \omega_n^2 & \omega_n^4 & \dots & (\omega_n^2)^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega_n^{n-1} & (\omega_n^{n-1})^2 & \dots & (\omega_n^{n-1})^{n-1} \end{bmatrix}$$

$$\begin{pmatrix} y_0 \end{pmatrix}^T = \begin{pmatrix} \text{---} \end{pmatrix} \begin{pmatrix} | \end{pmatrix}$$

$$V_n^{-1} \vec{y} = V_n^{-1} V_n \cdot \vec{a}$$

$$V_n^{-1} \vec{y} = \vec{a}$$

V_n is invertible

$$[V_n^{-1}]_{jk} = \frac{1}{n} \omega_n^{-kj}$$

check $V_n^{-1} V_n = I_n$

$$[V_n^{-1} \cdot V_n]_{ij} = \sum_{k=0}^{n-1} \frac{\omega_n^{-ik}}{n} \omega_n^{kj}$$

$$= \frac{1}{n} \sum_{k=0}^{n-1} \omega_n^{k(j-i)}$$

$$= \begin{cases} 1 & \text{if } i=j \\ 0 & \text{if } i \neq j \end{cases}$$

✓

$$a_j = \frac{1}{n} \sum_{k=0}^{n-1} y_k \cdot (w_n^{-1})^{kj}$$

Modify RECURSIVE-DFT as follows:

- switch $a_k \leftrightarrow y_k$ variables
- replace w_n with w_n^{-1}
- divide a_k 's by n .

⇓
DFT⁻¹_n

Selection Problem

input: + A array (unsorted, distinct values) length = n.

+ k ($0 \leq k \leq n-1$)

output: $A_{\text{sorted}}[k] = ?$