

HOMEWORK 3

Question 1:

+ bowling balls are sporting equipment:

$$\forall x \text{ ball}(x) \wedge \text{bowling}(x) \rightarrow \text{equipment}(x, \text{sporting})$$

+ All domesticated horses have an owner:

$$\forall x (\text{horse}(x) \wedge \text{domesticated}(x) \rightarrow \exists y \text{Owner}(x, y))$$

+ the rider of a horse can be different than the owner.

$$\exists x \text{rider}(\text{horse}, x) \wedge \neg \text{owner}(\text{horse}, x).$$

+ horses move faster than frogs

$$\forall x, y \text{horse}(x) \wedge \text{frog}(y) \rightarrow \text{isfaster}(x, y)$$

+ a finger is any digit on a hand other than the thumb.

$$\forall x \text{finger}(x) \leftrightarrow \text{digit}(x) \wedge \neg \text{isthumb}(x)$$

+ an isosceles triangle is defined as a polygon w 3 edges, which are connected at 3 vertices, where 2 (but not 3) edges have the same length.

$$\forall a \text{triangle}(a, \text{isosceles}) \rightarrow \exists x, y, z \text{edge}(x) \wedge \text{edge}(y) \wedge \text{edge}(z) \wedge \text{connected}(x, y) \wedge \text{connected}(y, z) \wedge \text{connected}(x, z) \wedge \text{equal}(x, y) \wedge \neg \text{equal}(x, z) \wedge \neg \text{equal}(y, z)$$

Question 2:

$$\forall x \text{person}(x) \wedge [\exists z \text{petOf}(x, z) \wedge \forall y \text{petOf}(x, y) \rightarrow \text{dog}(y)] \rightarrow \text{doglover}(x).$$

$$1 [\forall x \text{person}(x) \wedge [\exists z \text{petOf}(x, z) \wedge \forall y \text{petOf}(x, y) \rightarrow \text{dog}(y)]] \vee \text{doglover}(x)$$

$$2 \exists x \forall person(x) \vee [\forall z \forall petOf(x, z) \vee \exists y \forall (petOf(x, y) \rightarrow dog(y))] \vee \text{doglover}(x)$$

$$3 \exists x \exists z \forall y [(\neg \text{person}(x) \vee \neg \text{petOf}(x, z) \vee \neg (\text{petOf}(x, y) \rightarrow \text{dog}(y))) \vee \text{doglover}(x)]$$

$$4 \exists x \exists z \forall y [(\neg \text{person}(x) \vee \neg \text{petOf}(x, z) \vee \neg \text{petOf}(x, y) \vee \text{dog}(y)) \vee \text{doglover}(x)]$$

$$5 \exists x \forall y [(\neg \text{person}(x, y) \vee \neg \text{petOf}(x, y) \vee \neg \text{petOf}(y, x) \vee \text{dog}(y) \vee \text{doglover}(x, y)]$$

Question 3:

. ones(owner(X), citibank, cost(X)) ones(owner(ferrari), Z, cost(Y))

unifiable ✓ $u = \{ X/Y, X/ferrari, Z/citibank \}$.

. gives(bill, jerry, book1) . gives(X, brother(X), Z)

not unifiable ✓ because you can't replace constant with function.

. opened(X, result(open(X), SD))) and opened(toolbox, Z)

unifiable ✓ $u = \{ X/toolbox, result(open(X), SD)/Z \}$

Question 4:

(a) Marcus is a Pompeian: Pompeian(Marcus) (1)

All pompeians are Romans: $\forall x \text{ Pompeian}(x) \rightarrow \text{Romans}(x)$ (2)

Cesar is a ruler : Ruler(Caesar). (3)

All Romans are either loyal to Caesar or hate Caesar (not both)

$\forall x \text{ Romans}(x) \rightarrow \text{Loyal}(x, \text{Caesar}) \wedge \neg \text{H}(x, \text{Caesar}) \wedge \forall y (\text{L}(x, y) \vee \text{H}(x, y)) \wedge y \neq \text{Caesar}))$. (4)

Everyone is loyal to someone :

$\forall x \text{ person}(x) \rightarrow \exists y \text{ loyal}(x, y)$. (5)

People only try to assassinate rulers they are not loyal to :

$\forall x, y \text{ person}(x) \wedge \text{ruler}(y) \wedge \text{assassinate}(x, y) \wedge \neg \text{loyal}(x, y)$ (6)

Marcus tries to assassinate Caesar:

assassinate(Marcus, Caesar). (7)

(b) Prove that Marcus hates Caesar by Natural Deduction.

8. Romans(Marcus) (MP, 1, 2) $\Theta = \{x / \text{Marcus}\}$

9. Loyal(Marcus, Caesar) \vee Hates(Marcus, Caesar) (MP, 4, 8) $\Theta = \{x / \text{Marcus}\}$

10. $\exists y \text{ Loyal}(\text{Marcus}, y)$ (5, existential) $\Theta = \{x / \text{Marcus}\}$

11. $\neg \text{Loyal}(\text{Marcus}, \text{Caesar})$ (MP, 6, 7) $\Theta = \{x / \text{Marcus}, y / \text{Caesar}\}$

12. Hates(Marcus, Caesar) (MP, 9, 11) $\Theta = \{x / \text{Marcus}\}$

(c) Convert all the sentences into CNF. {need to do Skolem}

* Convert to conjunctive form:

. $P(\text{marcus}) : \neg \text{Marcus} \vee P(x)$ (1)

. $\forall x P(x) \rightarrow R(x) : \neg P(x) \vee R(x)$ (2)

. Ruler(Caesar) : $\neg \text{Caesar} \vee \text{Ruler}(x)$ (3)

. $\neg \text{H}(x) \vee \text{L}(x, \text{Caesar}) \vee \text{H}(x, \text{caesar})$ (4)

. $\forall x \exists y \text{ L}(x, y) / \text{L}(x, s(x))$ (5) \rightarrow Skolem

. $\forall x \forall y (\neg A(x, y) \vee \neg \text{L}(x, y)) \wedge (\neg \text{A}(x, y) \vee R(y)) \wedge (\neg A(x, y) \vee R(x))$ (6)

. A(Marcus, Caesar) (7)

(d) Prove that Marcus hates Caesar using Resolution Refutation.

Negation: $\neg \text{Hate}(\text{Marcus}, \text{Caesar})$

(8) $\neg \text{Roman}(\text{Marcus}) \vee \text{Loyal}(\text{Marcus}, \text{Caesar})$

(9) $\text{Loyal}(\text{Marcus}, \text{Caesar}) \vee \neg \text{Pompeian}(\text{Marcus})$

(10) $\text{Loyal}(\text{Marcus}, \text{Caesar})$

(11) $\neg \text{Ruler}(\text{Caesar}) \vee \neg \text{Assassinate}(\text{Marcus}, \text{Caesar})$

(12) $\neg \text{Ruler}(\text{Caesar}) \quad \rightarrow \text{Contradiction} \rightarrow \text{Marcus hates Caesar.}$

Question 5:

a. Map-Coloring.

- Assume that there are five states: WA, NT, SA, Q, NSW.
- Assume these neighboring states: WA & NT, WA & SA, NT & SA, NT & Q, SA & Q, SA & NSW, Q & NSW

KB = $\{ \forall c \text{ color}(c), \forall s \text{ state}(s), \forall s_1 s_2 \text{ adjacent}(s_1, s_2) \rightarrow (\text{color}(s_1, c_1) \wedge \text{color}(s_2, c_2)) \rightarrow (c_1 = c_2) \}$

$$\forall s_1 \forall s_2 (\text{adjacent}(s_1, s_2) \rightarrow (\forall c_1 \forall c_2 (\text{color}(s_1, c_1) \wedge \text{color}(s_2, c_2)) \rightarrow (c_1 = c_2)))$$

b. Sammy's Sport Shop

Facts: $\{ \text{obs}(1, Y), \text{obs}(2, W), \text{obs}(3, Y), \text{label}(1, W), \text{label}(2, Y), \text{label}(3, B) \}$

- $\forall x \text{ Contains}(x, W) \vee \text{Contains}(x, Y) \vee \text{Contains}(x, B)$

1. box1-3 represents 3 boxes

2. W,Y,B: white, yellow, both

Predicates: $\text{Contains}(x, c) \rightarrow x \text{ contains balls color } c$.

$\text{Obs}(x, c)$: observe box have balls color c.

$\text{Label}(x, c)$: label box have balls color c.

each box contains either y>w, both.

Axioms: $\{ \text{b1}, \text{b2}, \text{b3} \mid (\text{contains}(\text{b1}, W) \vee \text{contains}(\text{b2}, Y) \vee \text{contains}(\text{b3}, B)) \}$

$\forall b (\text{contains}(b, B) \rightarrow (\neg(\text{contains}(b, W) \wedge \text{contains}(b, Y)) \wedge \neg(\text{contains}(b, W) \vee \text{contains}(b, Y) \vee \text{contains}(b, B)))$

unique color for each box

label wrong

$\forall x, c \text{ label}(x, c) \rightarrow \neg \text{contains}(x, c)$,

$\{ \text{obs}(1, Y), \text{obs}(2, W), \text{obs}(3, Y), \text{label}(1, W), \text{label}(2, Y), \text{label}(3, B) \}$.

c. Wumpus World:

. Room(x,y): room @ location (x,y)

. Pit(x,y): a pit @ (x,y)

. Wumpus(x,y): wumpus @ (x,y)

. Gold(x,y): gold @ (x,y)

. Adjacent(x₁,y₁,x₂,y₂): rooms are adjacent.

. Safe(x,y): safe to explore.

KB: $\{ \forall x, y (\text{Pit}(x, y) \rightarrow \exists x_2, y_2 (\text{Adjacent}(x, y, x_2, y_2) \wedge \text{Breeze}(x_2, y_2))) \}$

$\forall x, y (\text{Wumpus}(x, y) \rightarrow \exists x_2, y_2 (\text{Adjacent}(x, y, x_2, y_2) \wedge \text{Stench}(x_2, y_2)))$,

$\forall x, y (\text{Gold}(x, y) \rightarrow \text{Glitter}(x, y))$, $\forall x, y (\text{Safe}(x, y) \rightarrow \neg \text{Pit}(x, y) \wedge \neg \text{Wumpus}(x, y))$,

$\forall x, y, z (\text{Pit}(x, y) \wedge \text{Wumpus}(x, y)) \rightarrow \forall x, y, z (\text{Adjacent}(x, y, z, y) \rightarrow \neg \text{Pit}(z, y) \rightarrow \text{Breeze}(z, y))$,

$\forall x, y, z (\text{Adjacent}(x, y, z, y) \rightarrow \neg \text{Wumpus}(x, y) \rightarrow \text{Stench}(z, y))$,

$\forall x, y \text{ Safe}(x, y) \rightarrow (\exists z \text{ Adjacent}(x, y, z, y) \rightarrow \text{Safe}(z, y)) \wedge (\exists z \text{ Adjacent}(x, y, x, z) \rightarrow \text{Safe}(x, z)) \}$.

(a) A queen can only be in 1 row + column :

$$\forall r \forall c \forall c' ((\text{queen}(r,c) \wedge \text{queen}(r,c') \wedge c \neq c') \rightarrow \text{False})$$

$$\forall r \forall r' \forall c ((\text{queen}(r,c) \wedge \text{queen}(r',c) \wedge r \neq r') \rightarrow \text{false})$$

2. No two queens can be on the same diagonal:

$$\forall r \forall r' \forall c \forall c' (((\text{queen}(r,c) \wedge \text{queen}(r',c')) \wedge |\text{abs}(r-r')| = |\text{abs}(c-c')|) \rightarrow \text{false})$$

3. There must be exactly 4 queens in total:

$$\exists r_1 \exists r_2 \exists r_3 \exists r_4 (\text{queen}(r_1, \text{col}(1)) \wedge \text{queen}(r_2, \text{col}(2)) \wedge \text{queen}(r_3, \text{col}(3)) \wedge (\text{queen}(r_4, \text{col}(4)))$$