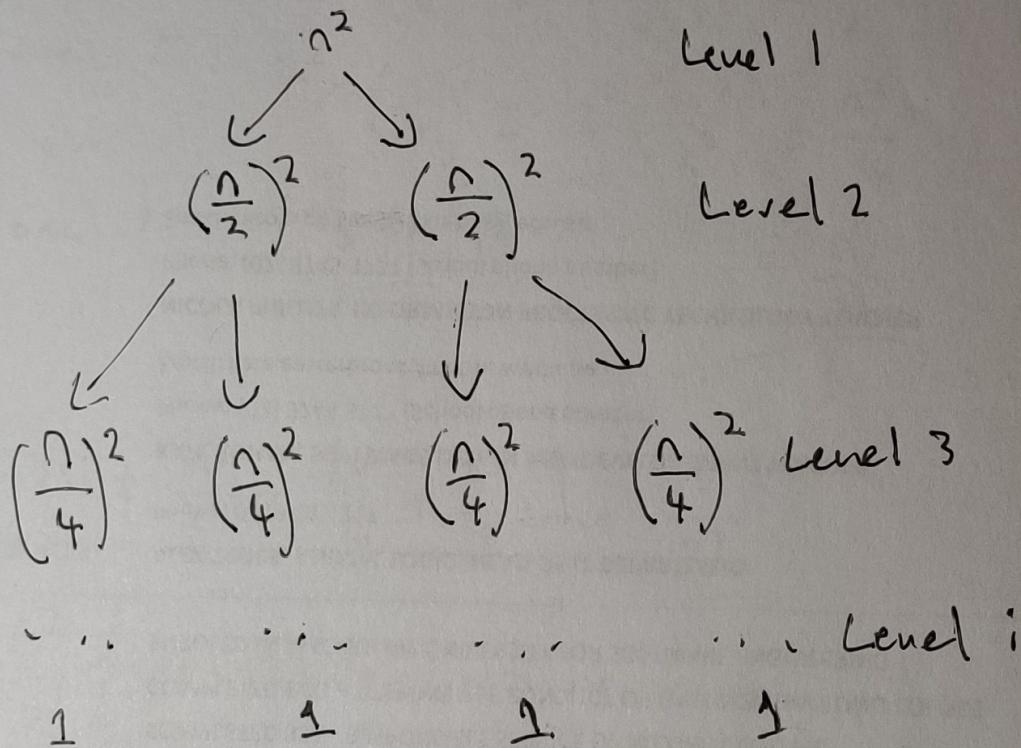


Q1

$$T(n) = 4T\left(\frac{n}{2}\right) + n^2$$

Recursion Tree:



Therefore this is reduced to  $T(?)$  at level ;

$$\frac{n}{2^i} = 1$$

$n = 2^i$  (take log on both sides)

$$\log_2 n = \log_2 2^i$$

$$\therefore i = \log_2 n$$

$$\therefore i = \log_2 n$$

# Q 1 Continued . -

constant at each level

$$\text{level 0} = n^2$$

$$\text{level 1} = 2 \times \left(\frac{n}{2}\right)^2 = 2^1 \left(\frac{n}{2^1}\right)^2 = \frac{n^2}{2}$$

$$\text{level 2} = 4 \times \left(\frac{n}{4}\right)^2 = 2^2 \left(\frac{n}{2^2}\right)^2 = \frac{n^2}{4}$$

$$\text{level } i = 2^i \times \left(\frac{n}{2^i}\right)^2 = \underbrace{\frac{n^2}{2^i}}$$

$$T(n) = n^2 + \frac{n^2}{2} + \frac{n^2}{4} + \frac{n^2}{8} + \dots + \frac{n^2}{2^i}$$

$$\frac{n^2}{2^i} \left( 1 + 2 + 2^2 + 2^3 + 2^4 + 2^5 + \dots + 2^i \right)$$

$$\frac{n^2}{2^i} (2^i - 1)$$

$$\frac{n^2}{2^{\log_2 n}} (2^{\log_2 n} - 1)$$

$$\frac{n^2}{n^2} (n - 1)$$

$$n(n-1)$$

$$n^2 - n$$

$$\therefore T(n) = \underline{\underline{\Theta(n^2)}}$$

Q2

• Base Case

When  $n$  is equal to 1.

• induction hypothesis

Assume that  $\exp(x, n)$   
is equal to  $x^n$

y

Inductive proof

Q 3

a) 4 transactions:

T1: \$80, \$90, \$100

T2: \$50, \$60, \$70      Total discount: \$150

T3: \$20, \$30, \$40

T4 \$10

b)

1) Sort items from largest to smallest

2) loop through every 3rd item add every 3rd item to integer (total)

3. return total

c) The greedy choice in this algorithm is to only visit every 3rd item. starting from index 2 This then becomes an optimal solution

Q 4

a)

Bottom up table;

	4	1	3	6	7	8
1	0	1	2	3	4	5
2	1	0	1	2	3	4
3	2	1	1	2	3	4
4	3	2	2	2	3	4
5	4	3	3	3	3	4
6	5	4	4	4	3	5
7	6	5	5	4	4	5
8	7	6	5	5	5	5
9	8	7	6	6	6	5
						6

= 00123345

b) let 45426317 = SID

let 413678 = NUM

SID: replace 5 with 1 at pos 2

SID: replace 4 with 3 at pos 3

SID: delete 2 at pos 4

SID: replace 3 with 7 at pos 6

NUM: replace 1 with 8 at pos 7

NUM: delete 7 at pos 8