# Solving a coin-flip 12 minimax estimation problem

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#### Abstract

We set up and solve a minimax under 12 loss problem derived from a coin flipping problem. The solution is interesting as it involves use of cancellation to solve the minimax problem.

### 1 Introduction

Wald [Wald, 1949] set up statistical estimation as a game played against nature where the researcher picks a (possibly probabilistic) decision function and nature picks an adversarial distribution. Nature's distribution plays the role of Bayesian priors, but is not considered to be the an objective true distribution or a subjective estimate. It is instead a worst-possible distribution so that any inference bounds proven in this formalism hold in general. This game theoretic form of probability is fascinating and leads quickly to interesting questions and procedures.

# 2 The Problem

Take as our problem the task of estimating the unknown win-rate (or heads-rate) p of a random process or coin. We assume the process is memory-less and stationary (p is not changing and does not depend on earlier flips). We observe a sequence of n flips showing h wins/heads, and then are asked to return an estimate  $\phi_n(h)$  for p. This problem was discussed and given context in [Mount, 2014a], [Mount, 2014c], [Mount, 2014b], and [Bauer et al., 2014].

Fix  $n \in \mathbb{N}$ ,  $n \ge 1$ . Let  $p \in [0,1]$  and  $\phi = (\phi(0), \dots, \phi(n))$  be a (n+1)-dimensional real vector in  $[0,1]^{n+1}$ , and define:

$$L_n(p,\phi) := \sum_{h=0}^n \binom{n}{h} p^h (1-p)^{n-h} (\phi(h) - p)^2.$$
 (1)

 $L_n(p,\phi)$  represents the expected square-error of encountered when using  $\phi$  to estimate the win-rate of a coin with (unknown) true win-rate p by observing n flips/outcomes. The estimate is: use  $\phi(h)$  when you see h wins/heads. This is related to Wald's game-theoretic formalism, but we are insisting on pure strategies for both the estimate (a single deterministic  $\phi$ ) and a single unknown true probability p. We are going to assume that nature picks p in an adversarial manner with full knowledge of  $\phi$ .

Define:

$$f_n(\phi) = \max_{p \in [0,1]} L_n(p,\phi).$$
 (2)

What we are looking for is  $\operatorname{argmin}_{\phi \in [0,1]^{n+1}} f_n(\phi)$ . The issue is: the definition of  $f_n()$  has two quantifiers so it seems like it will be difficult to derive or even check solutions.

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## 3 A Solution

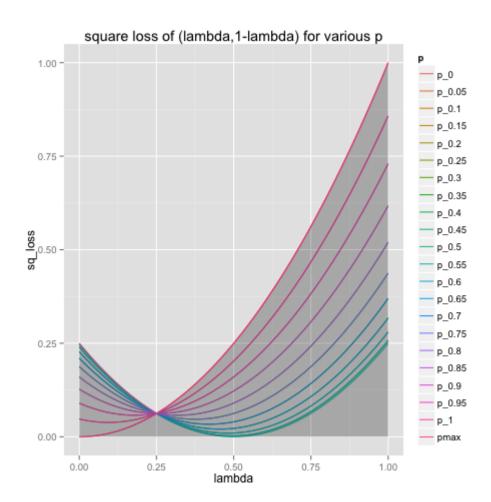


Figure 1:  $L_1(p,(\lambda,1-\lambda))$ 

**Lemma 1.** Suppose  $\phi$  is in the interior of  $[0,1]^{n+1}$  and is such that  $L_n(p,\phi)-\phi(0)^2=0$  simultaneously for all p. Then:  $\phi$  is the unique global minimizer of  $f_n()$ .

*Proof.* Suppose  $\phi$  is as stated. We will confirm  $\phi$  is an isolated local minimum by checking partial derivatives. Look at  $\frac{\partial}{\partial \phi(h)} f_n(\phi)$  and  $\frac{\partial}{\partial -\phi(h)} f_n(\phi)$ . If we can show these are always both positive for all h we are done.

Because  $L_n(p,\phi)$  is a constant independent of p we know  $\frac{\partial}{\partial \phi(h)} f_n(\phi) \geq \frac{\partial}{\partial \phi(h)} L_n(p,\phi)$  for any  $p \in [0,1]$  (i.e. all p curves are active or on the boundary boundary, figure 1 shows an example). So

$$\frac{\partial}{\partial \phi(h)} f_n(\phi) \ge \max_p \frac{\partial}{\partial \phi(h)} L_n(p, \phi)$$

$$= \max_p \binom{n}{h} p^h (1 - p)^{n-h} 2(\phi(h) - p)$$

$$\ge \binom{n}{h} p^h (1 - p)^{n-h} 2(\phi(h) - p) \Big|_{p = \phi(h)/2}$$

$$> 0$$

Similarly we know  $\frac{\partial}{\partial -\phi(h)} f_n(\phi) \geq \frac{\partial}{\partial -\phi(h)} L_n(p,\phi)$  for any  $p \in [0,1]$ . So

$$\frac{\partial}{\partial - \phi(h)} f_n(\phi) \ge \max_p \frac{\partial}{\partial - \phi(h)} L_n(p, \phi)$$

$$= \max_p \binom{n}{h} p^h (1 - p)^{n-h} 2(p - \phi(h))$$

$$\ge \binom{n}{h} p^h (1 - p)^{n-h} 2(p - \phi(h)) \Big|_{p = (1 + \phi(h))/2}$$

$$> 0$$

So we know  $\phi$  is an isolated local minimum of  $f_n()$  But  $L_n(p,\phi)$  is convex in  $\phi$  for any fixed n,p  $(n \geq 1, p \in \mathbb{R})$ , so  $f_n(\phi)$  is also convex in  $\phi$ . So an isolated local minimum  $\phi$  is also the unique global minimum.

Define:  $\phi_n$  as the vector in  $\mathbb{R}^{n+1}$  such that

$$\phi_n(h) := (\frac{1}{2}\sqrt{n} + h)/(\frac{1}{2}\sqrt{n} + n). \tag{3}$$

**Lemma 2.**  $\phi_n$  from equation 3 is in the interior of  $[0,1]^{n+1}$  has  $L_n(p,\phi) - \phi(0)^2 = 0$  simultaneously for all p.

*Proof.* It is obvious is in the interior of  $[0,1]^{n+1}$ . So it is just a matter of checking  $L_n(p,\phi_n) - \phi_n(0)^2 = 0$  using arguments from [Bauer et al., 2014] or by checking  $\phi_n$  obeys the recurrences in [Mount, 2014b].

**Theorem 1.**  $\phi_n$  from equation 3 is the unique minimizer of  $f_n(\phi)$  and the only  $\phi$  in the interior of  $[0,1]^{n+1}$  such that  $L_n(p,\phi) - \phi(0)^2 = 0$ .

Proof. By lemma 2 we know  $\phi_n$  meets the conditions of lemma 1. Therefore  $\phi_n$  is the unique global minimizer of  $f_n()$ . This follows by combining lemma 1 and lemma 2. The uniqueness of the minimizer of  $f_n()$  means there can be no other solutions of  $L_n(p,\phi) - \phi(0)^2 = 0$  that meet the pre-conditions of lemma 1, so  $\phi_n$  must be the only solution  $L_n(p,\phi) - \phi(0)^2 = 0$  in the interior of  $[0,1]^{n+1}$ . Note  $L_n(p,\phi) - \phi(0)^2 = 0$  can have solutions outside of  $[0,1]^{n+1}$  (for example it is known to have non-real solutions).

It is kind of neat we get that these is no more than one solution of  $L_n(p,\phi) - \phi(0)^2 = 0$  in the interior of  $[0,1]^{n+1}$  from the convexity of the related optimization problem.

# References

[Bauer et al., 2014] Bauer, A., Mount, J., and Dotsenko, V. (2014). Mathoverflow: Existence of solutions of a polynomial system. http://mathoverflow.net/questions/177574/existence-of-solutions-of-a-polynomial-system. Accessed: 2014-08-08.

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