## Monadic design patterns for the Web

Putting data, query and analytic in the hands of the people who use them

L.G. Meredith

## Overview

- Confessions of a mad (computer) scientist
- An architecture for semantic search
- A path to this architecture
- Informed by the monadic design pattern
- Demo

## Confessions of a mad (computer) scientist

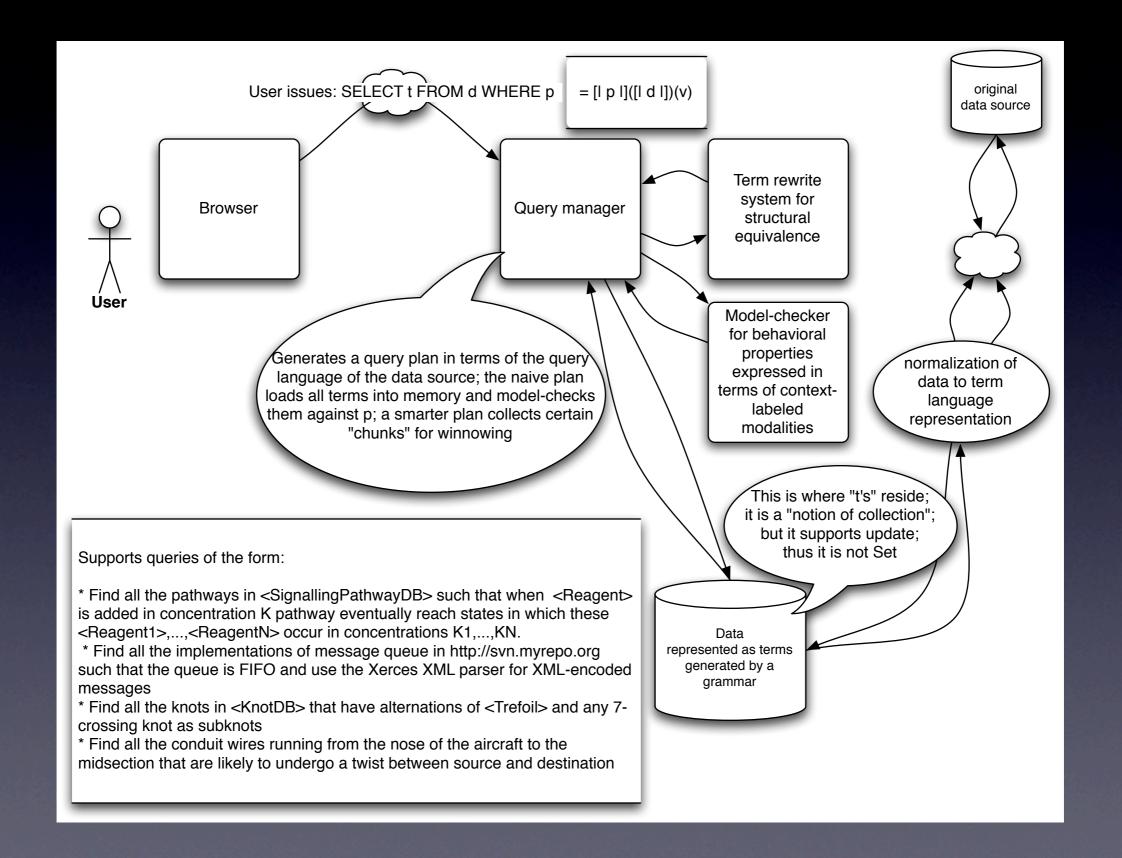
- What am i passionate about?
  - A revolutionary approach to search not based on "semantic web" techniques and yet results in genuine semantics-based search capability.
  - The technical details of this idea have taken 5 years to work out

- (Behavioral and structural queries in physical systems) Find all the pathways, p, in <SignallingPathwayDB> such that when <Reagent> is added in concentration K p eventually reaches states in which we see these <Reagent I>,..., <Reagent N> in concentrations K I,..., KN.
- (Behavioral and structural queries in logical systems) Find all the implementations of message queue in <a href="http://svn.myrepo.org">http://svn.myrepo.org</a> such that the queue is FIFO and use the Xerces XML parser for XML-encoded messages
- (Structural queries in physical systems) Find all the knots in <KnotDB> that have alternations of <Trefoil> and any 7-crossing knot as subknots
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## Confessions of a mad (computer) scientist

- How am i going to build this?
  - Get an active open source community excited by the ideas and let the magic happen!
  - Everybody wins if this can be realized at the right scope and scale

#### Architecture



- What is a collection?
- What is an item?
- What is an analytic?
- A machine for mashing things up
- Questions?

 Back in undergraduate maths we knew what a collection was: it was anything you could write like this

```
{ x in D | C1(x), C2(x), ..., Cn(x) }
```

 Later, when we got jobs in computing it was anything you could write like this

```
select x from D where C1(x) and ... and Cn(x)
```

- Then along came Haskell (and Python, and C#, and F#, and Scala, and XQuery, and ...)
   and the difference between these two ways of thinking about collections began to blur
- [x for x in S if x % 2 == 1] (Python)
- [x | x <- S, odd x] (Haskell)

• ...



This key language feature came about because some very clever people in computing and maths saw some very deep patterns that unified a lot of different phenomena





And a few large scale companies and institutions knew a good thing when they saw it



• The fact is these are just two syntaxes for <u>one</u> underlying language for denoting collections

```
{ x \text{ in } D \mid C1(x), C2(x), \ldots, Cn(x) } select x \text{ from } D \text{ where } C1(x) \text{ and } \ldots \text{ and } Cn(x) In scala
```

for (  $x \leftarrow D$  if  $C1(x) \&\& \dots \&\& Cn(x)$  ) yield x

 That can be backed in memory, or storage of many different shapes and kinds

- And that language has already proven itself to be an excellent way to represent data sets
- ... for the last 30 years of enterprise computing
- ... and the last 100 years of mathematics

 Back in undergraduate maths we knew what an item was: it was anything you could put between

```
'{' and '}'
```

 Later, when we got jobs in computing it was anything you could stick in a table and retrieve with statements like

```
select x from D where C1(x) and ... and Cn(x)
```

- Then along came the Internet and it became important that individuals and companies that didn't have common access to a single store be able to exchange items
- So.... items were things you could put
   between <tag> and </tag>

- But tag soup is pretty unsavory... so several proposals for what an item is have been in play
  - DTD, XSD, RelaxNG, ...
- All of these line up with the notion of algebraic data type as witnessed in Haskell and F# and Scala and OCaml and Erlang and ...

Some 80/20 rules to keep in mind:

- Mutable data types don't make very good message types
- And they don't make very good record types (for storage media)
- As the Telco's figured out a long time ago the intersection between algebraic data types and what's described by an EBNF grammar is more than adequate

- Things are looking pretty rosy
  - A collection is anything denoted by an existing store or a select-from-where construct
  - An item is anything described by an EBNF grammar
- There's always some little detail

- Sometimes items have parts
  - 'Menus' have 'items' and 'submenus'
  - 'Graphs' have 'vertices' and 'edges' and 'subgraphs'
- ... and sometimes they have lots of parts

We never write out all the integers, we just say Int

```
lazy val intStrm =

List( 0 ).toStream +++ ( intStrm map { _ => _ + 1 } )
```

 ... but more importantly the web 2.0 is teaching us that what initially looks like a small thing

```
type TwitterMsg = Char [140] // not legal Scala
```

becomes pretty valuable when it is iterated into a stream

```
abstract class TwitterStream

case class Twitterful( msg : TwitterMsg, strm : TwitterStream )
extends TwitterStream
```

And the <u>streams are bundled</u>

```
case class Tweeter( myStrm : TwitterStream, followers :
[TwitterStream] )
```

And then we mine them for common patterns

```
def hottopic( tweeters : List[Tweeter] ) =
  tweeters match {
  case ( out, follows ) :: ts =>
    for ( x <- merge( follows )
        if ( follows precedes x out ) ) yield x
  ... }</pre>
```

- Notice hottopic is also a stream
- Many interesting items have parts (components) that are programmatically specified -- we refer to the parts programmatically either because
  - there are too many of them or
  - they arrive on an as-needed basis
  - because we gain a significant information compression advantage in our exchange of these items

We are familiar with this idea from simple arithmetic.

- When there are a reasonable number of numbers to add we write out the addition, like so:  $a_1 + a_2 + ... + a_n$
- When there are more than a reasonable number of numbers to explicitly write out the addition, we write it out in indexed form, like so:  $\sum_{i \in I} a_i$

In this example '+' provides the composition and numbers are both the container and the component (this is the operational definition of what it means for a data type to be compositional: instances can be in the role of container and in the role of component)

- Likewise, when an item is compositional in nature (like numbers or menus or graphs) then we need to add indexed forms of composition to their <u>term</u> language
- Example:

```
case class Menu[A]( items : List[(A,Menu[A])] )

really needs to be

abstract class Menu[A]

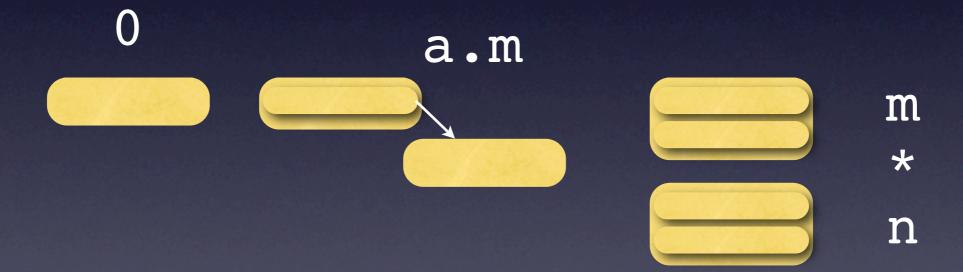
case class MenuExplicit( items : List[(A , Menu[A])] ) extends Menu[A]

case class MenuProgrammatic(

   mptn : MPtn[A], msrc : MGen[A], mcond : List[MenuCondition[A]]

) extends Menu[A]
```

A more compact term language for menus



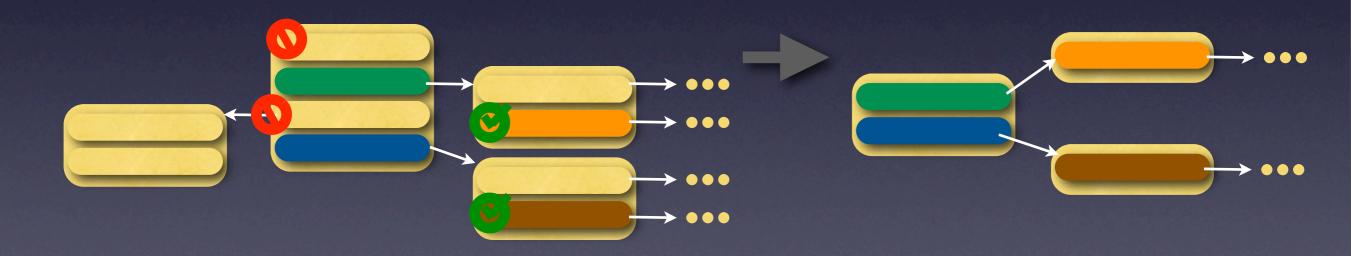
This data structure

```
M a ::= 0 | a . M a | M a * M a
```

Becomes

```
{ x . r : x . l*r <- m, r in true.~0, l in true.~0 }
```

represents a kind of "projection" that picks out submenus satisfying a more global property



# A machine for mashing things up

- (QM a) is <u>calculated</u> from (M a)
  - Plus a 'notion of collection' -- first line of MCnd grammar says this is something that is a boolean lattice -- hence essentially the powerset of the set of all (M a)'s
  - A grammar giving formulae of multiplicative LL would have implied the notion of collection is powerset of sequences of (M a)'s -- another gives streams
  - This means the Q process is applicable to any type with compositional structure -- hence mashable

• This data structure

```
M a := 0 a . M a M a * M a
```

Data type + Collection Type

Becomes

```
QM a ::= 0 | a . QM a | QM a * QM a
```

Data type w/ built in indexing

```
MPtn a ::= M (LitOrVar a)

MSat a ::= MPtn a in MCnd a
```

Pattern language

MCnd a ::= true | ~ MCnd a | MCnd a or MCnd a

Logic

0 | a . MCnd a | MCnd a \* MCnd a

{ MPtn a : (MPtn <- Ma)\*, (MSat a)\* }

## Summary

- A collection is anything we can write with a comprehension
  - Effectively given by a monad C
- An item is anything we can represent with EBNF
  - Effectively given by a monad G
- An analytic is a machine for mashing things up
  - Effectively given by a distributive law D : CG -> GC

## Summary

- D shows how to express collections of terms as terms of collections
- The existence of such distributive law for reasonable notions of collections and items is the generalization of Codd's theorem

# A machine for mashing things up

- Capability set is strictly richer than what can be said with LINQ
  - Menu example already causes exponential blow-up for MS LINQ
  - "Q" procedure extends to <u>behavioral</u> predicates these are not expressible in LINQ
- Efficient approach compiles this to a <u>lazy</u> query plan

# A machine for mashing things up

- (QM a) semantics is independent of the means by which the constraints are solved
- Oleg Kiselyov's LogicT monad transformer is a good example
   -- he shows that for a certain 'notion of collection' one can give a 'natural' backtracking solution
- Nothing prevents one from using a special purpose solver
  - This is one of the many reasons why having a specification of the logic that is <u>as tight as possible</u> is a good thing -- the solver can often be chosen from the constraint language

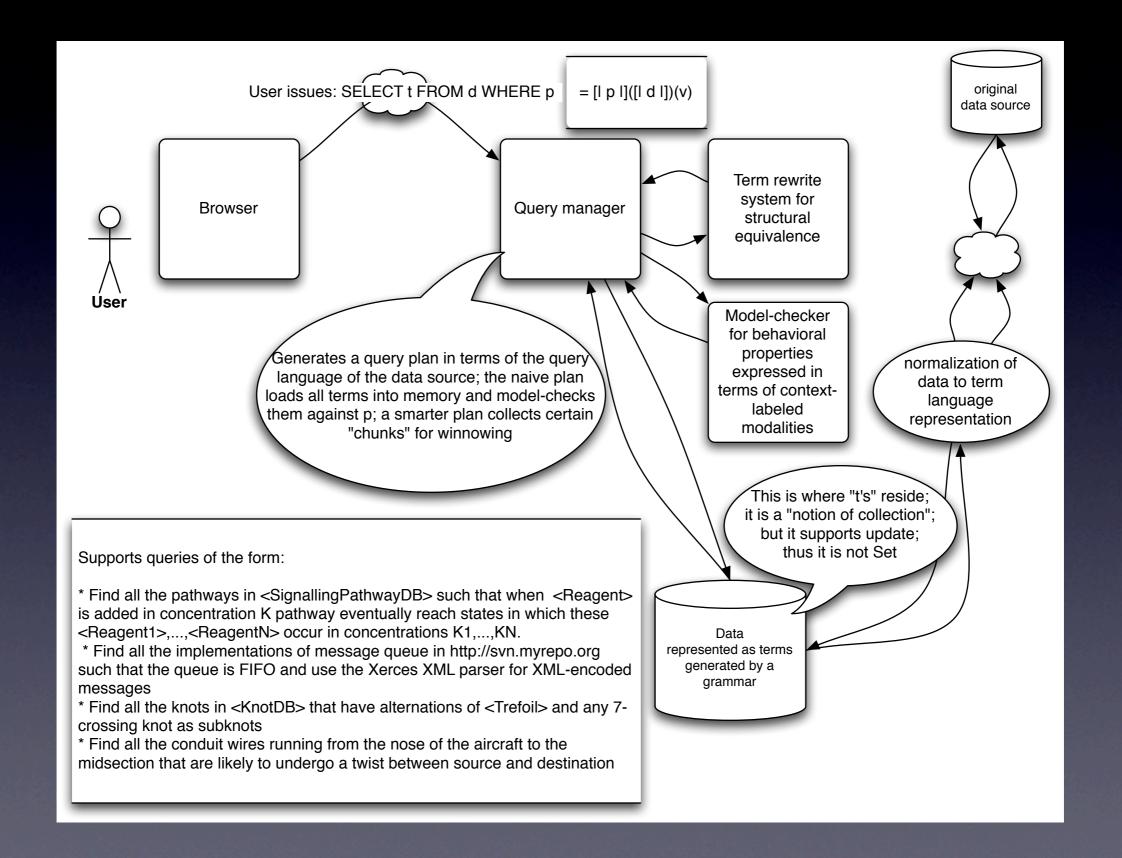
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  - Following Priami/Regev/Shapiro, et al, encode biological processes as processes in Milner's π-calculus
  - Term language monad -- Milner's π-calculus
  - Collection monad -- Nominal sets
  - Distributive law -- semantics of Caires' logic of spatial and behavioral observations
  - SELECT process FROM xmldb WHERE formula-in-logic-of-spatial-andbehavioral observations

- (Structural queries in physical systems) Find all the knots in <KnotDB> that have alternations of <Trefoil> and any 7-crossing knot as subknots
  - Following Meredith and Snyder knots as processes in Milner's  $\pi$ -calculus
  - Term language monad -- Milner's π-calculus
    - Alternate term language -- Conway's knotation
  - Collection monad -- Nominal sets
  - Distributive law -- semantics of Caires' logic of spatial and behavioral observations
  - SELECT process FROM xmldb WHERE formula-in-logic-of-spatial-andbehavioral observations

- (Behavioral and structural queries in logical systems) Find all the implementations of message queue in <a href="http://svn.myrepo.org">http://svn.myrepo.org</a> such that the queue is FIFO and use the Xerces XML parser for XML-encoded messages
  - Following Berger, Honda and Yoshida encode Java programs as strategies in games
  - Term language monad -- Games
  - Collection monad -- Mutable sets
  - Distributive law -- semantics of Berger, Honda, Yoshida Hoare-logic
  - SELECT strategy FROM xmldb WHERE formula-in-BHY-logic

- (Geometric queries in physical systems) Find all the conduit wires running from the nose of the aircraft to the midsection that are likely to undergo a twist between source and destination
  - Encode geometric information in Clifford algebra
  - Term language monad -- Clifford algebra
  - Collection monad -- Mutable sets
  - Distributive law -- exhibited by Meredith and Beckman
  - SELECT ensemble FROM xmldb WHERE formula-in-Geometric-logic

#### Architecture



## References

- Comprehending monads, Phil Wadler
- Backtracking, interleaving and terminating monad transformers, Oleg Kiselyov, Chung-chie Shan, Daniel P Friedman
- Behavioral and spatial observations in a logic for the π-calculus, Luis Caires
- Domain theory in logical form, Samson Abramsky

## Haskell/Scala cheat sheet

Haskell	Scala
do var1 <- e1 var2 <- e2 e	<pre>for( var1 &lt;- e1;      var2 &lt;- e2;      rslt &lt;- e ) yield rslt</pre>
do var1 <- e1 var2 <- e2 return e	for( var1 <- e1; var2 <- e2 ) yield e
do var1 <- e1 >> e2 return e	for( var1 <- e1; var2 <- e2 ) yield e

#### Haskell/Scala cheat sheet

#### Haskell Scala $do \{e; stmts\} =$ for( var <- e ) yield rslt = e >> do {stmts} $e map { x => rslt }$ do {p <- e; stmts} = for( var <- e if c ) yield rslt =</pre> let ok p = do {stmts} ok = fail "..." e filter { $x \Rightarrow c$ } map { $x \Rightarrow rslt$ } in e >>= ok for ( v1 <- e1; v2 <- e2 ) yield rslt = do {let decls; stmts} = let decls in do {stmts} el flatMap { $x \Rightarrow e2 map \{ x \Rightarrow rslt \} \}$

#### Haskell/Scala cheat sheet

```
class Monad m where
   -- chain
   (>>=) :: m a -> (a -> m b) -> m b
   -- inject
   return :: a -> m a
trait Monad {
   type T
   def map(f: Monad => T): T
   def flatMap(f: Monad => T): T
   def filter(f: Monad => Boolean): Monad
}
```

#### Category theory

```
M: C -> C
unit: Id -> M
mult: M<sup>2</sup> -> M
```

## Questions?