CS2040S Data Structures and Algorithms

Augmented Trees!

Where we are...

Dictionaries

Binary search trees

Tries

Balanced search trees

- AVL trees
- Scapegoat Trees
- B-trees

Today: Dynamic Data Structures

1. Maintain a set of items

2. Modify the set of items

3. Answer queries.

Big picture idea:

Trees are a good way to store, summarize, and search dynamic data.

Dynamic Data Structures

- Operations that create a data structure
 - build (preprocess)

- Operations that modify the structure
 - insert
 - delete

- Query operations
 - search, select, etc.

Augmented Data Structures

Many problems require storing additional data in a standard data structure.

Augment more frequently than invent...

Plan

Three examples of augmenting balanced BSTs

1. Order Statistics

2. Interval Queries

3. Orthogonal Range Searching

Basic methodology:

1. Choose underlying data structure (tree, hash table, linked list, stack, etc.)

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(subject to insert/delete/etc.)

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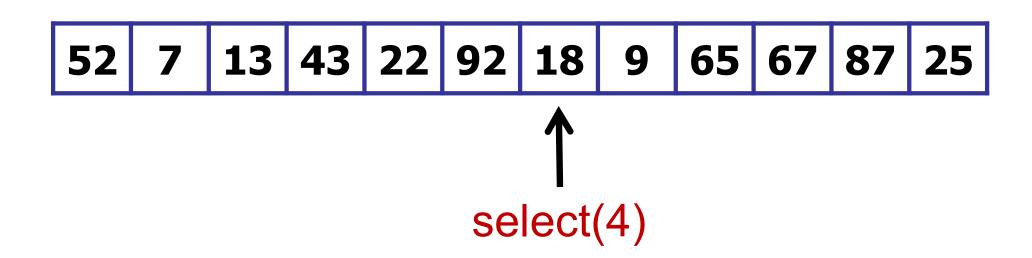
4. Develop new operations.

Input

A set of integers.

Output: select(k)

The kth item in the set.

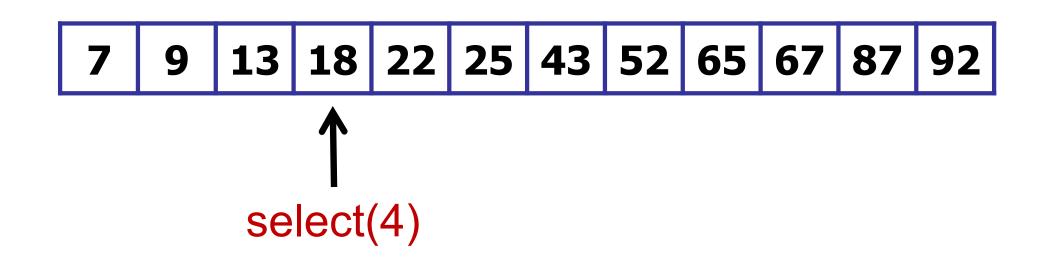


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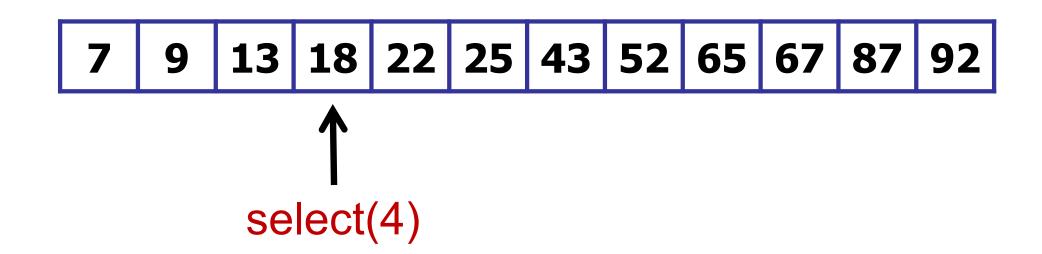


Input

A set of integers.

Output: $select(k) \longrightarrow Sort: O(n log n)$

The kth item in the set.

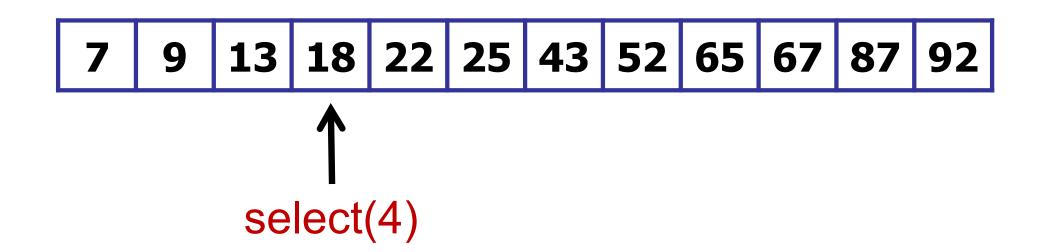


Input

A set of integers.

Output: select(k) ———— QuickSelect: O(n)

The kth item in the set.

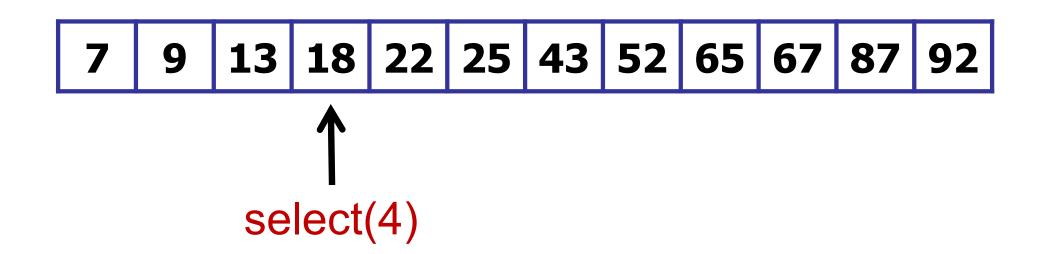


Solution 1:

Sort: O(n log n)

Solution 2:

QuickSelect: O(n)



More than one query

Solution 1:

Preprocess: sort --- O(n log n)

Select: O(1)

Solution 2:

Preprocess: nothing --- O(1)

QuickSelect: O(n)

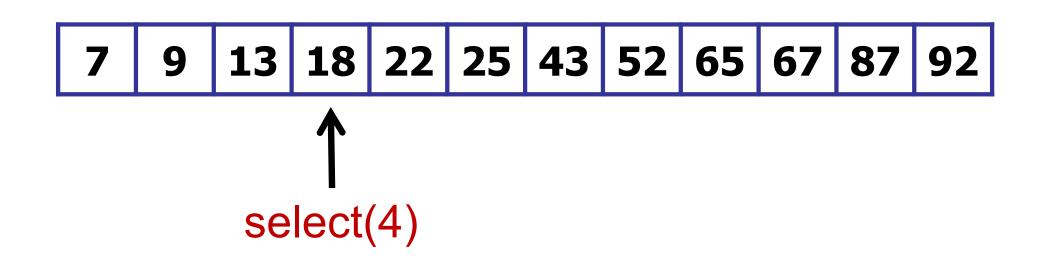
Trade-off: how many items to select?

Implement a data structure that supports:

- insert(int key)
- delete(int key)

and also:

select(int k)



Solution 1:

Basic structure: sorted array A.

insert(int item): add item to sorted array A.

select(int k): return A[k]

7 9 13 18 22 25 43 52 65 67 87 92

Solution 2:

Basic structure: unsorted array A.

insert(int item): add item to end of array A.

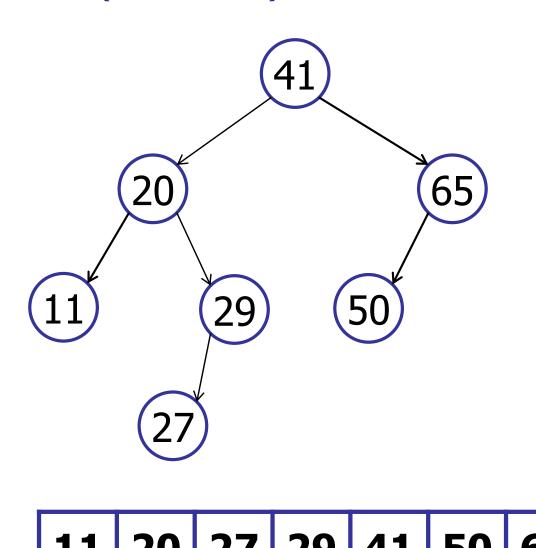
select(int k): run QuickSelect(k)

7 9 13 18 22 25 43 52 65 67 87 92

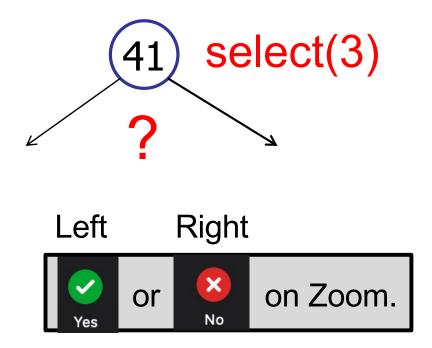
	Insert	Select
Solution 1: Sorted Array	O(n)	O(1)
Solution 2: Unsorted Array	O(1)	O(n)



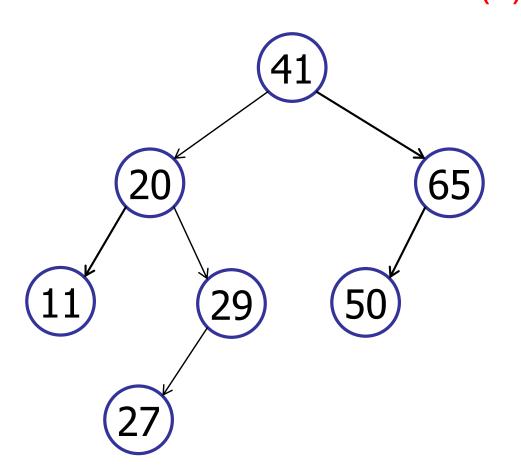
Today: use a (balanced) tree



How to find the right item?



Simple solution: traversal select(k): O(k)

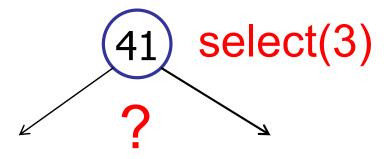


in-order traversal

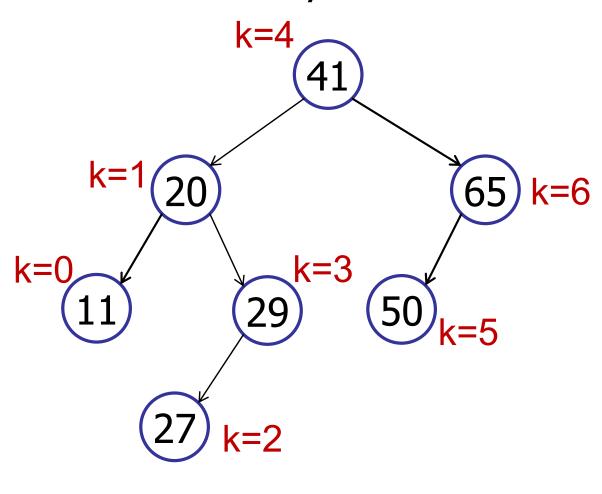
 11
 20
 27
 29
 41
 50
 65

Augment!

What extra information would help?

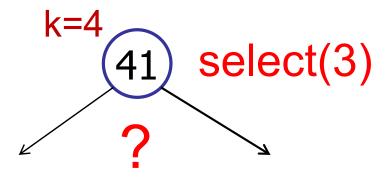


Idea: store rank in every node



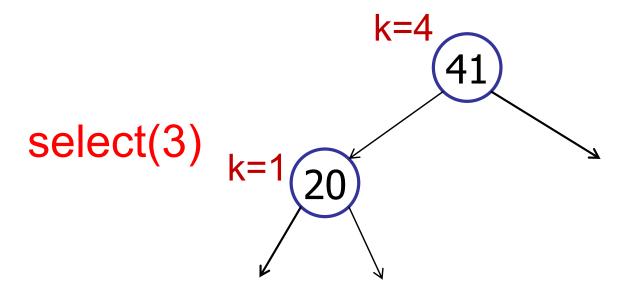
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Idea: store rank in every node



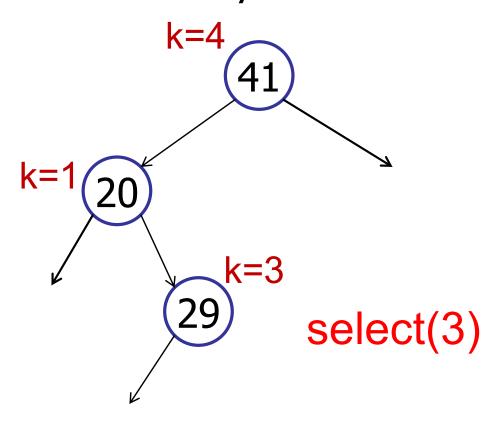
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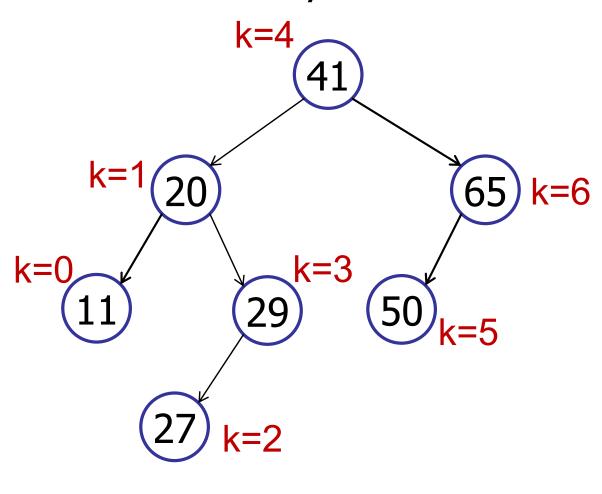
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Idea: store rank in every node



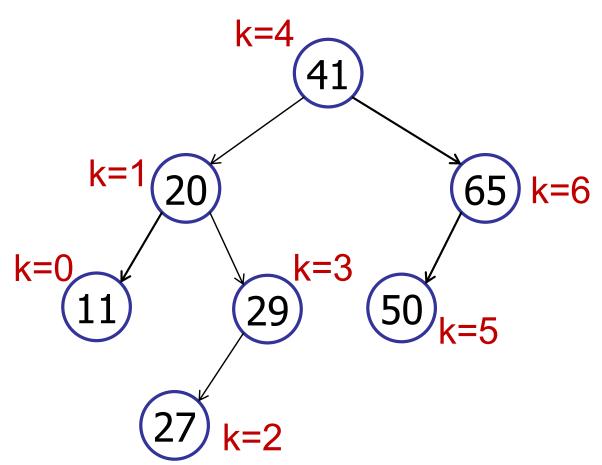
11 20 27 29 41 50 65

Idea: store rank in every node



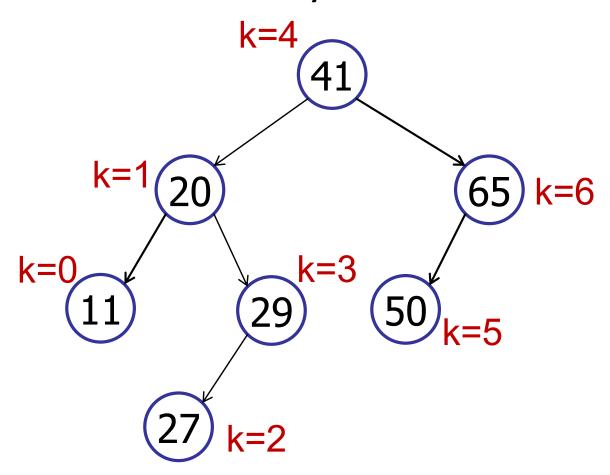
 11
 20
 27
 29
 41
 50
 65

Question: What goes wrong if you store ranks on every node??



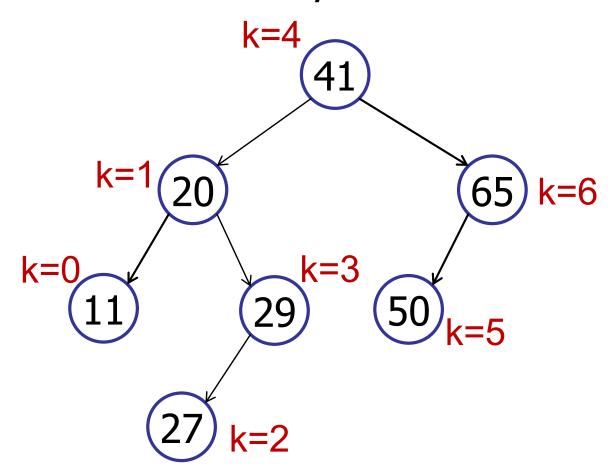


Idea: store rank in every node



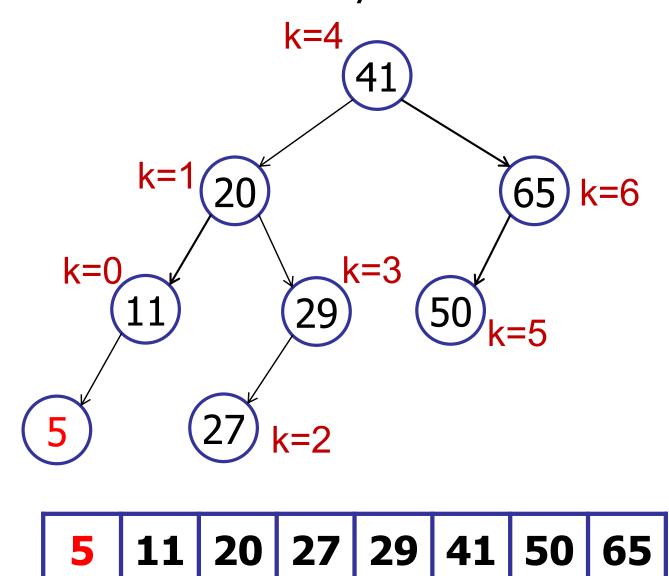
Problem: insert(5)

Idea: store rank in every node

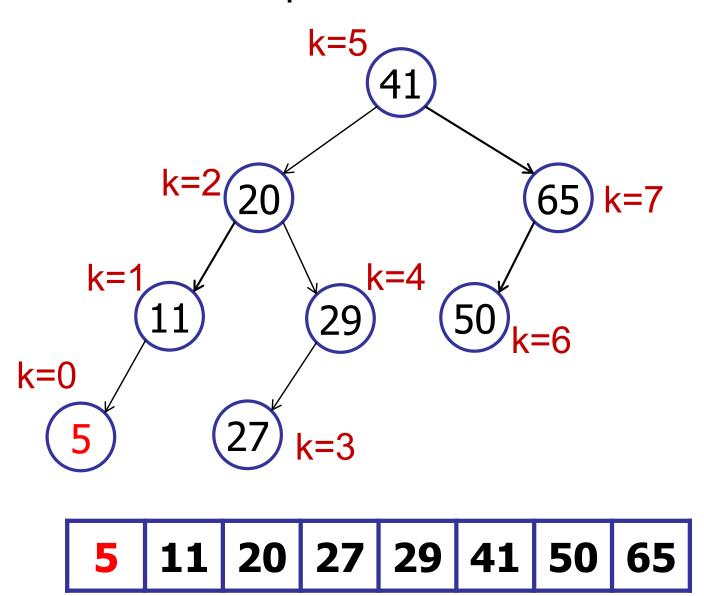


Problem: insert(5) requires updating all the ranks!

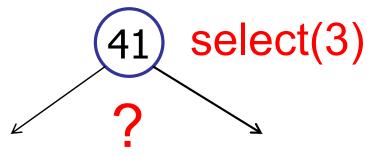
Idea: store rank in every node



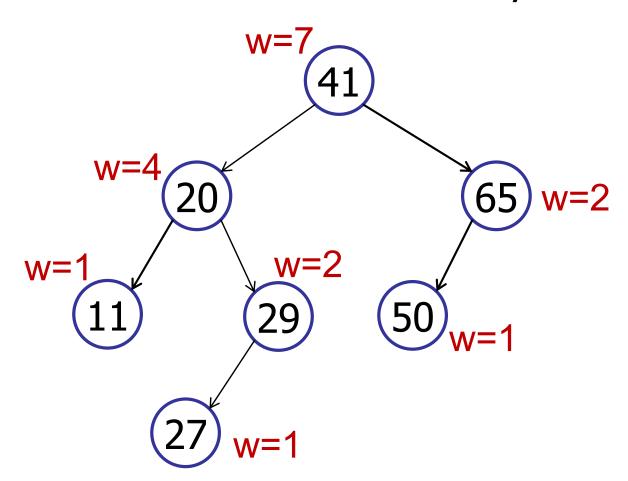
Conclusion: too expensive to store rank in every node!



What should we store in each node?



Idea: store *size* of sub-tree in every node



Idea: store size of sub-tree in every node

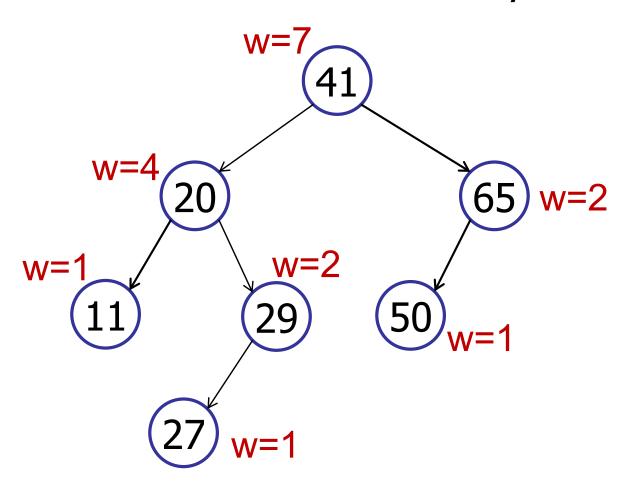
The <u>weight</u> of a node is the size of the tree rooted at that node.

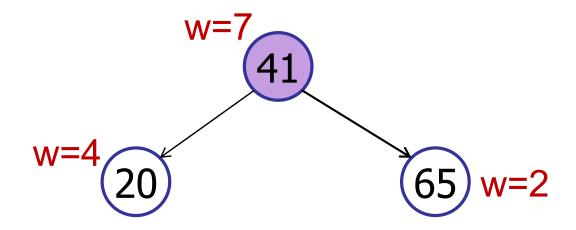
Define weight:

```
w(leaf) = 1

w(v) = w(v.left) + w(v.right) + 1
```

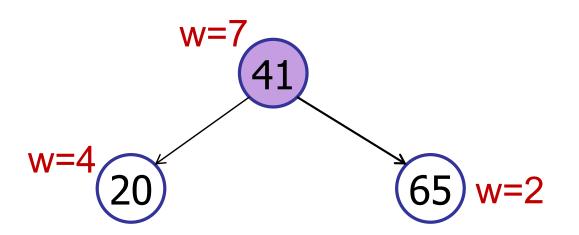
Idea: store *size* of sub-tree in every node





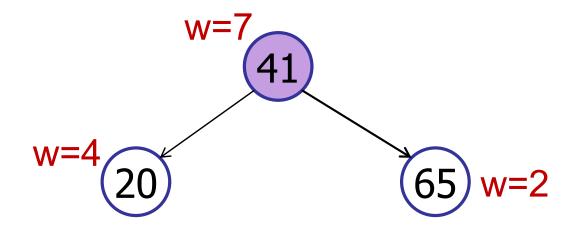
What is the rank of 41?

- 1. 1
- 2. 3
- **✓**3. 5
 - 4. 7
 - 5. 9
 - 6. Can't tell.

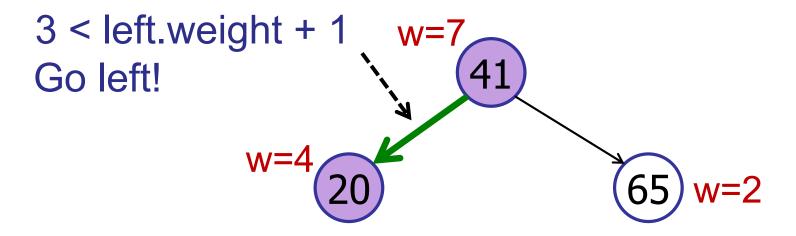


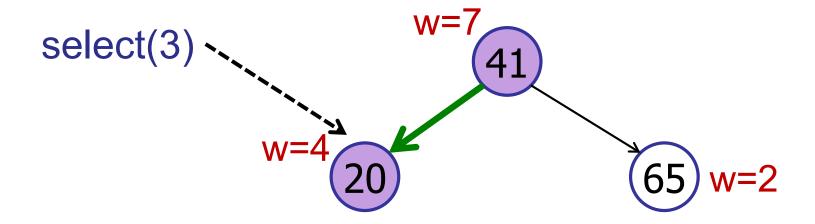


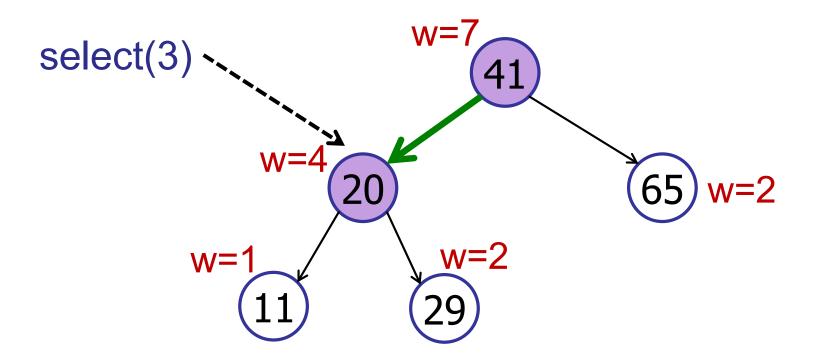
Example: select(3)

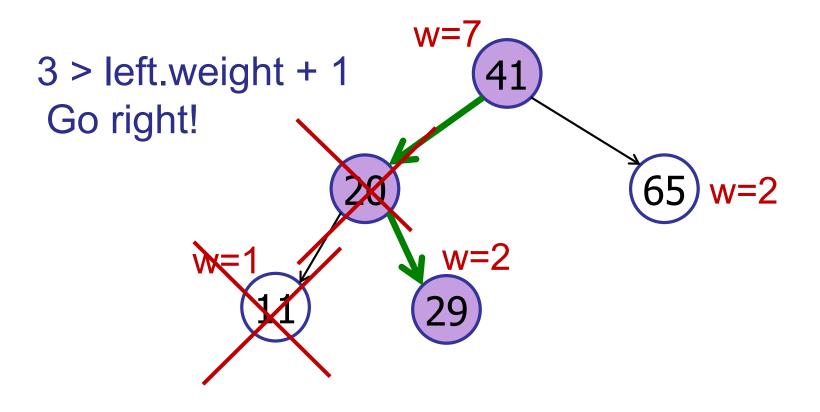


"rank in subtree" = left.weight + 1

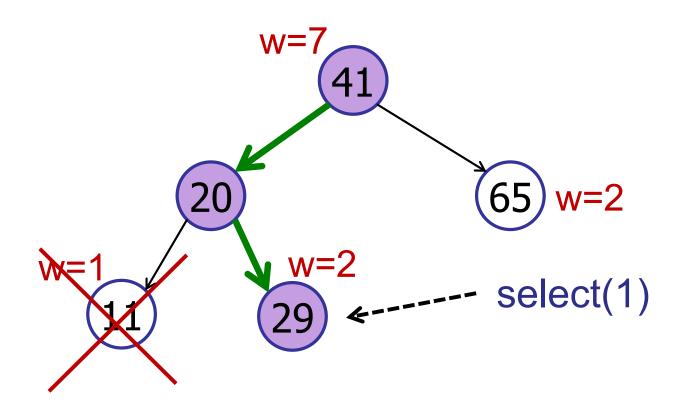








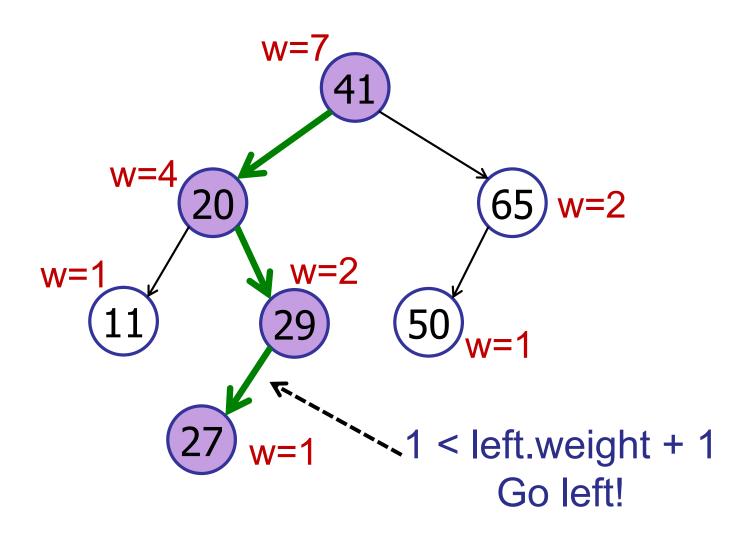
Example: select(3)

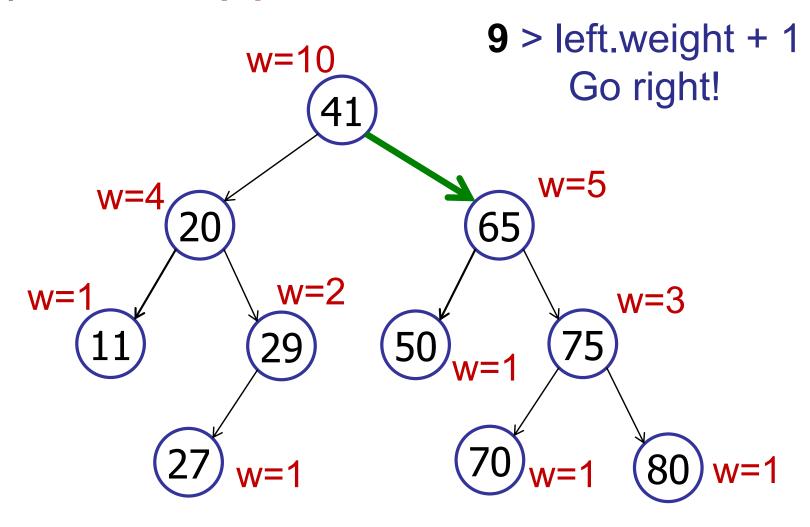


Item to select:

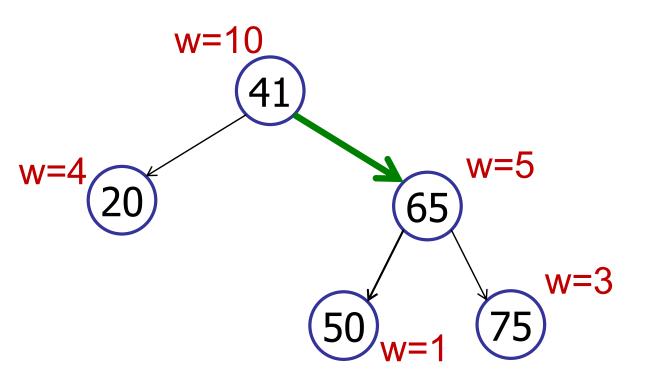
$$3 - (left.weight + 1) =$$

$$3 - (1 + 1) = 1$$

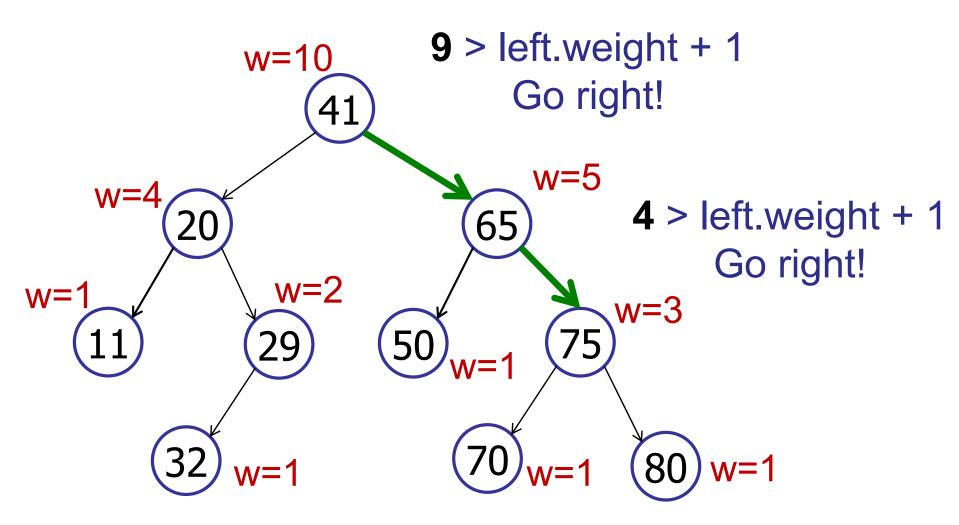




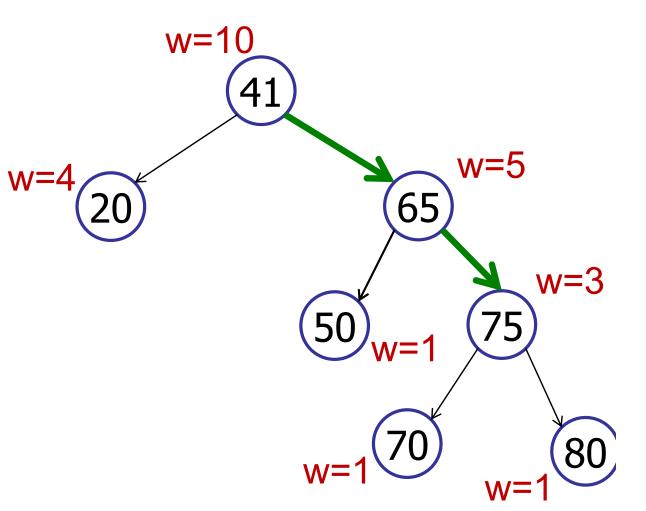
- 1. Go left at 65
- ✓2. Go right at 65
 - 3. Stop at 65
 - 4. I'm confused



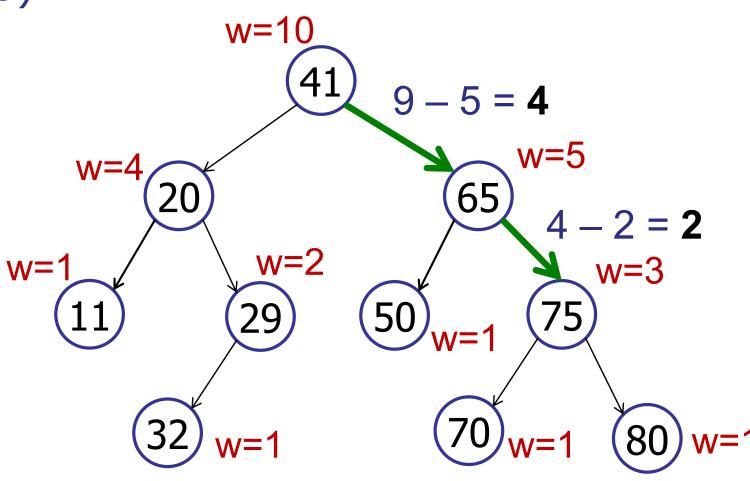




- 1. Go left at 75
- 2. Go right at 75
- **✓** 3. Stop at 75
 - 4. I'm confused







select(k)

```
rank = m left.weight + 1;
if (k == rank) then
    return v;
else if (k < rank) then
    return m left.select(k);
else if (k > rank) then
    return m right.select(k-rank);
```

select(k): finds the node with rank k

Example: find the 10th tallest student in the class.

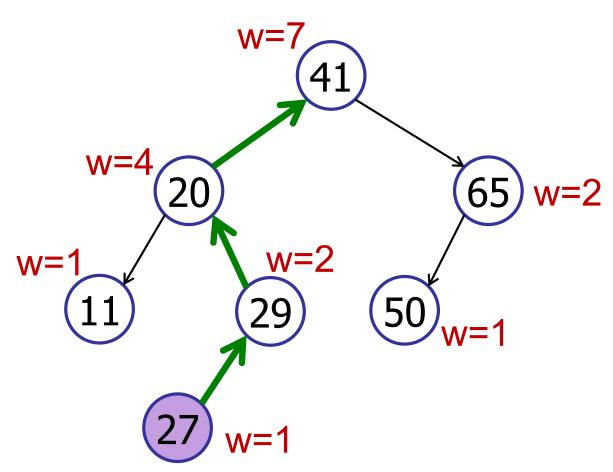
select(k): finds the node with rank k

Example: find the 10th tallest student in the class.

rank(v): computes the rank of a node v

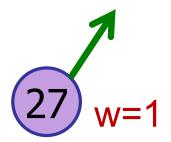
Example: determine the percentile of Johnny's height. Is Johnny in the 10th percentile or the 90th percentile?

Example: rank(27)



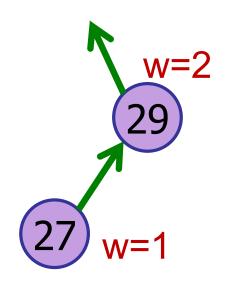
rank = 2

Example: rank(27)



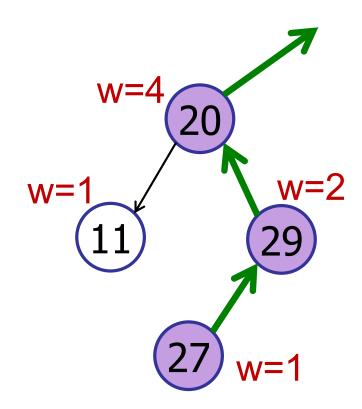
Initially: rank = 1

Example: rank(27)



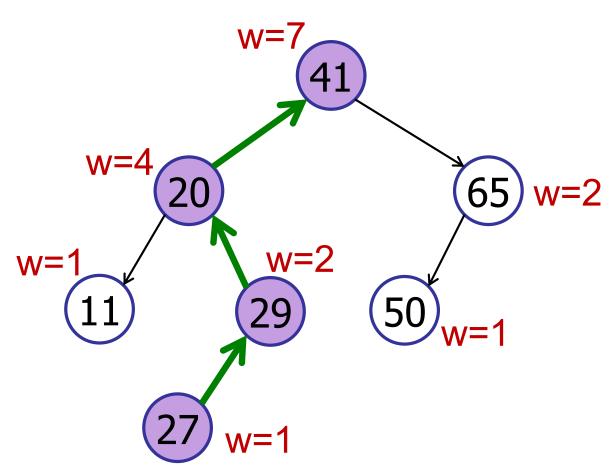
rank = 1

Example: rank(27)



$$rank = 1 + 2$$

Example: rank(27)

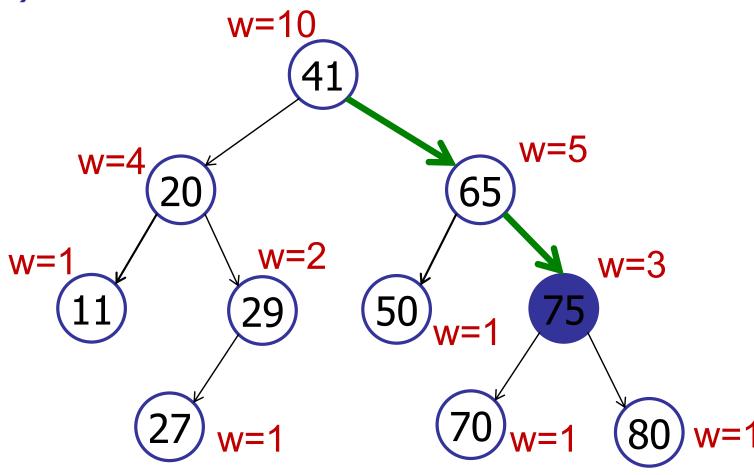


$$rank = 1 + 2 = 3$$

Rank(v): computes the rank of a node v

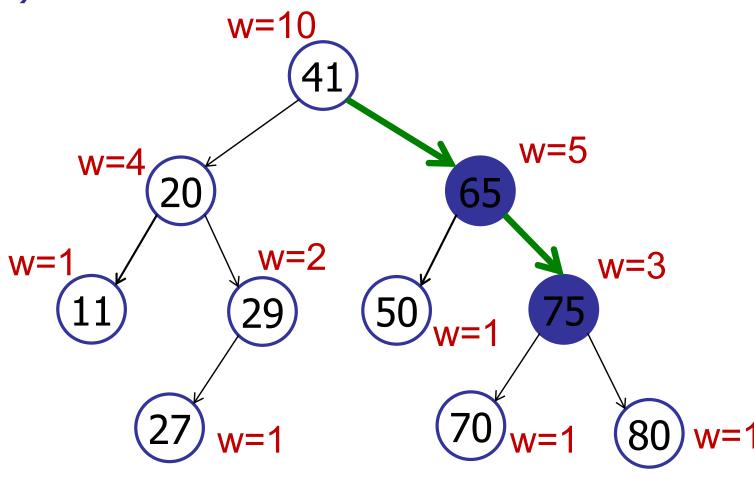
```
rank(node)
     rank = node.left.weight + 1;
     while (node != null) do
           if node is left child then
                 do nothing
           else if node is right child then
                 rank += node.parent.left.weight + 1;
           node = node.parent;
     return rank;
```

rank(75)



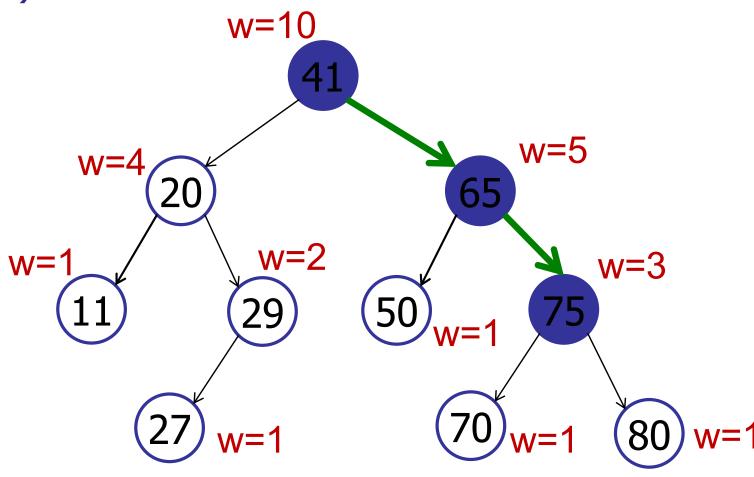
rank = 2

rank(75)



$$rank = 2 + 2$$

rank(75)



$$rank = 2 + 2 + 5 = 9$$

Rank(v): computes the rank of a node v

```
rank(node)
     rank = node.left.weight + 1;
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```

Augmenting data structures

Basic methodology:

1. Choose underlying data structure:

AVL tree

2. Determine additional info needed: Weight of each node

3. Maintained info as data structure is modified.

Update weights as needed

4. Develop new operations using the new info.

Select and Rank

Augmenting data structures

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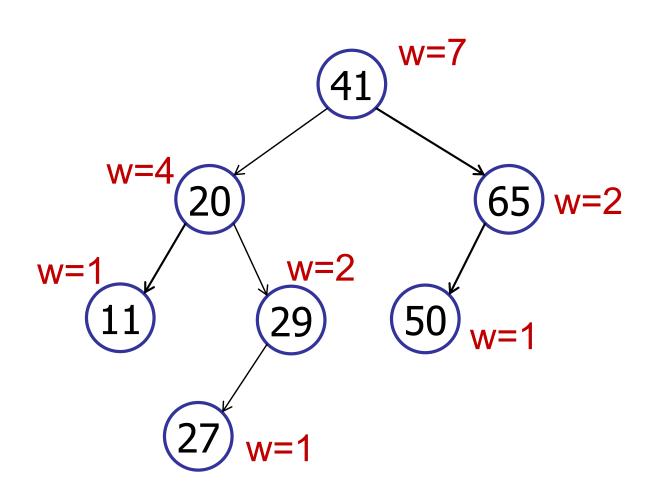
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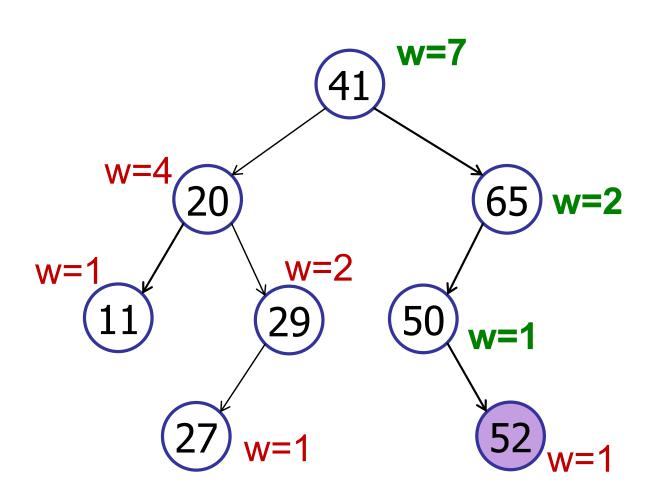
AVL tree

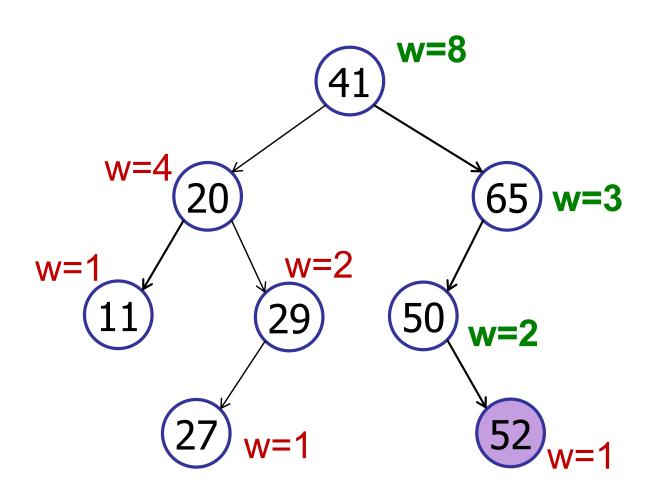
- 2. Determine additional info needed: Weight of each node
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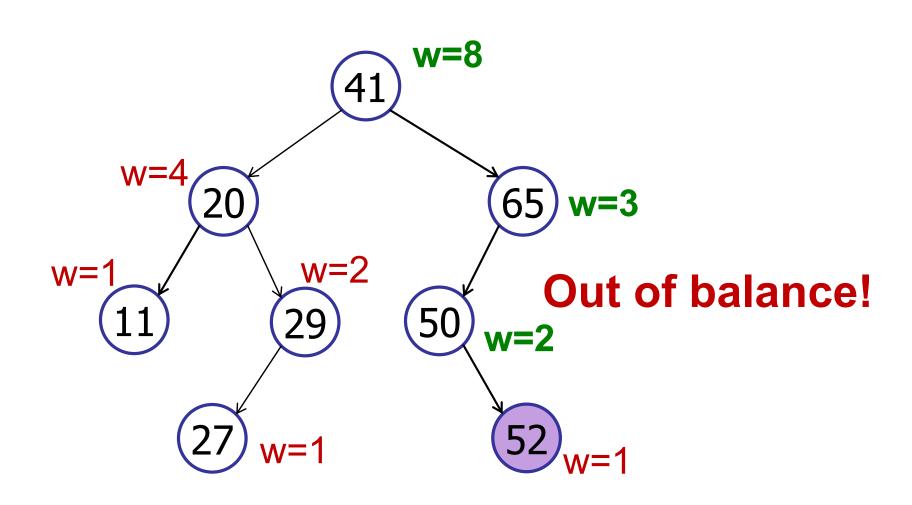
 Update weights as needed
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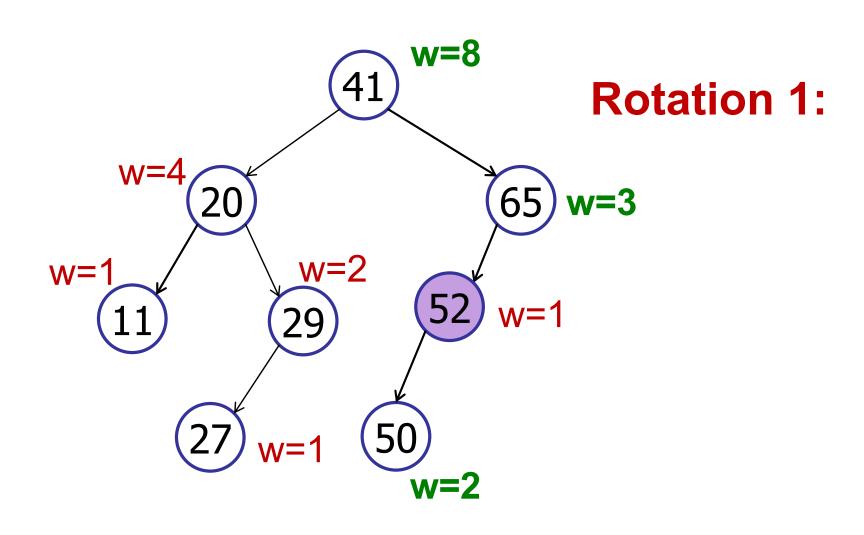
 Select and Rank



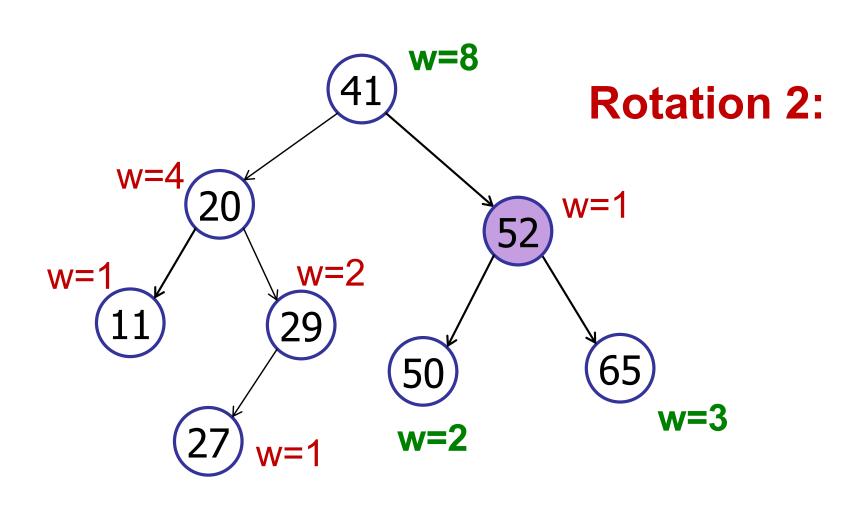




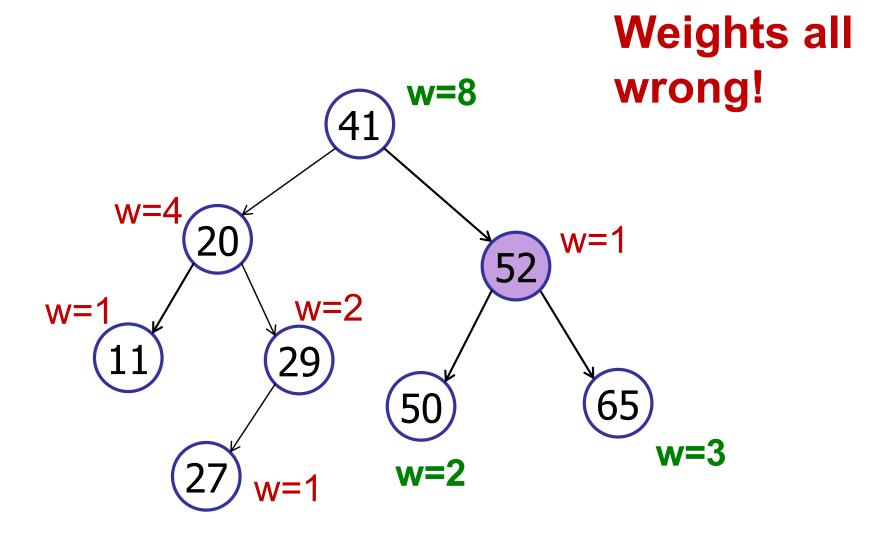


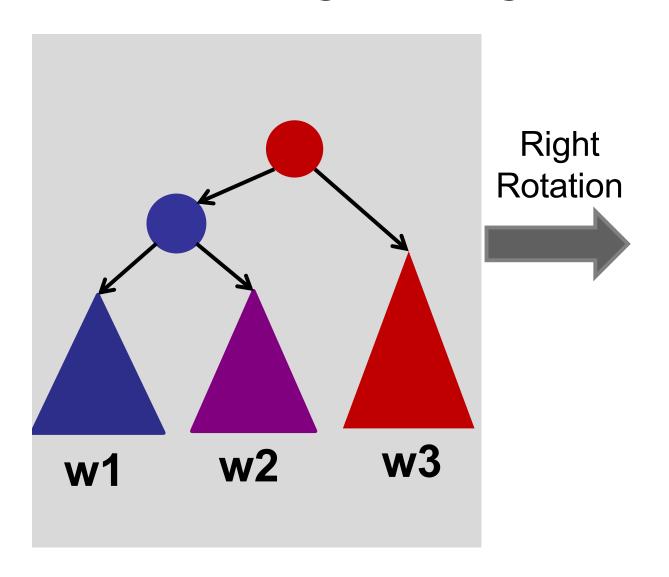


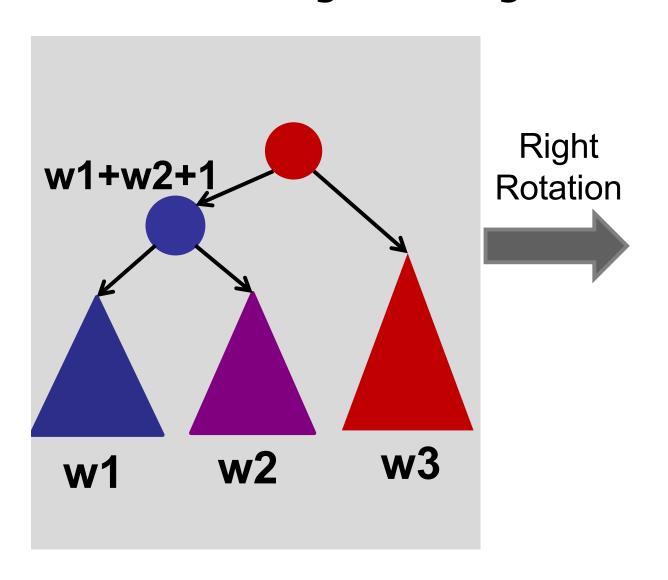
Maintain weight during insertions:

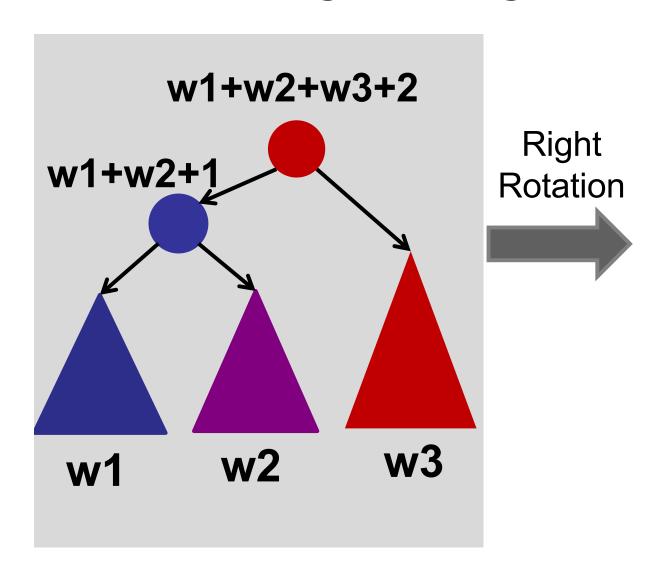


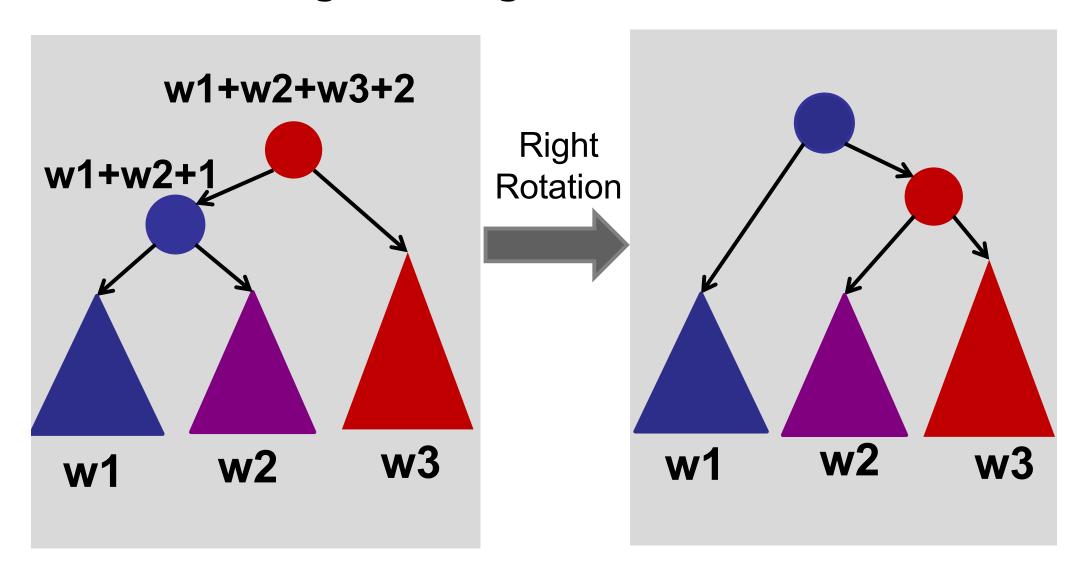
How to update weights on rotation?

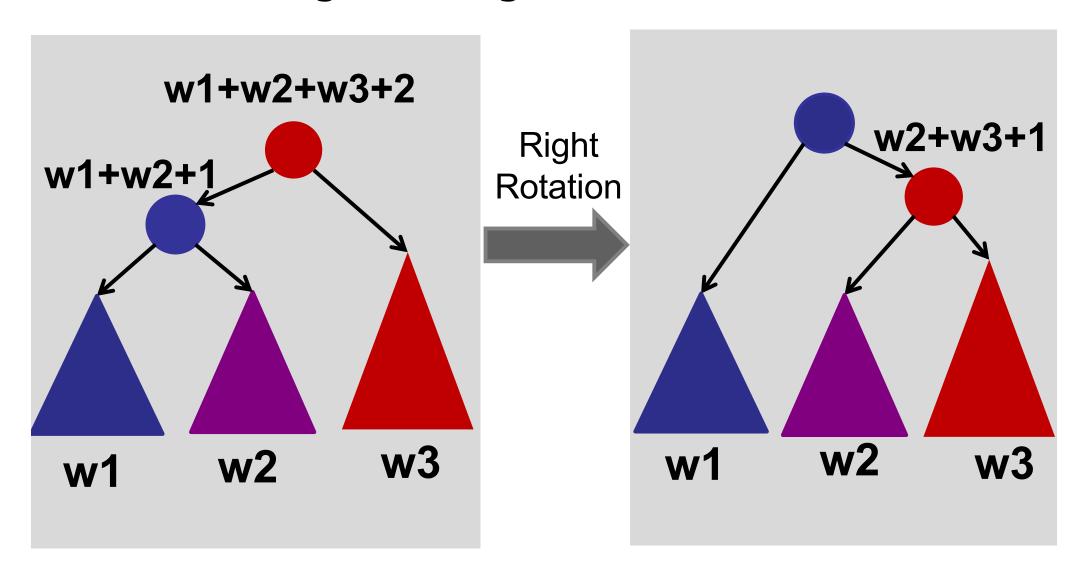


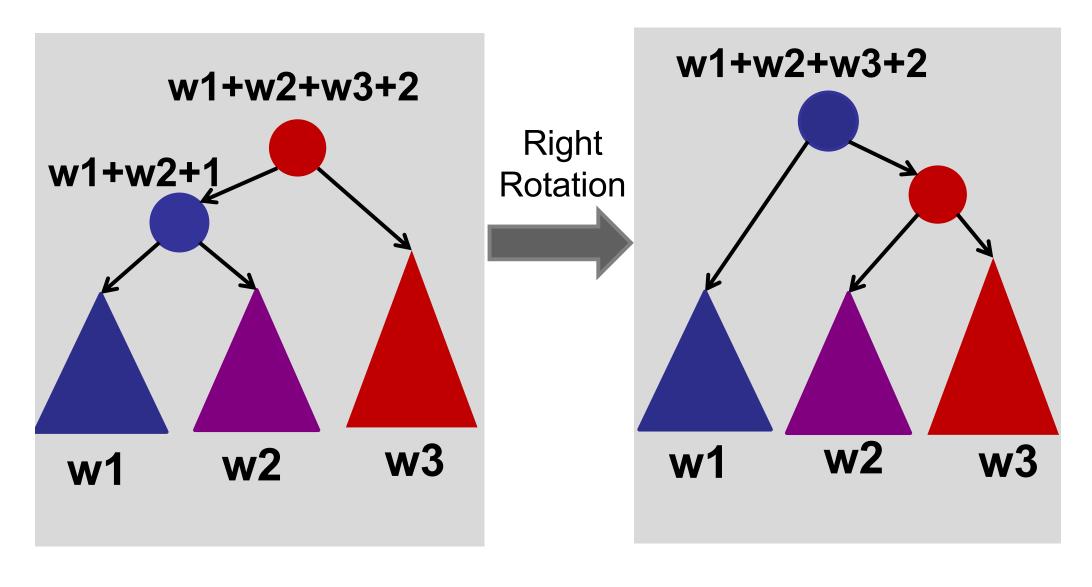


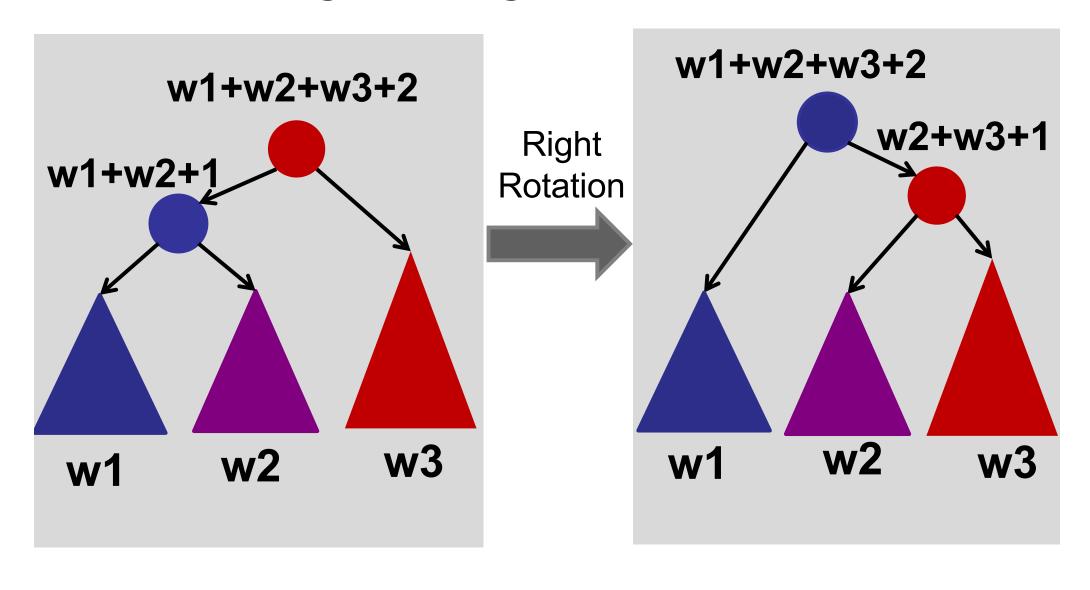








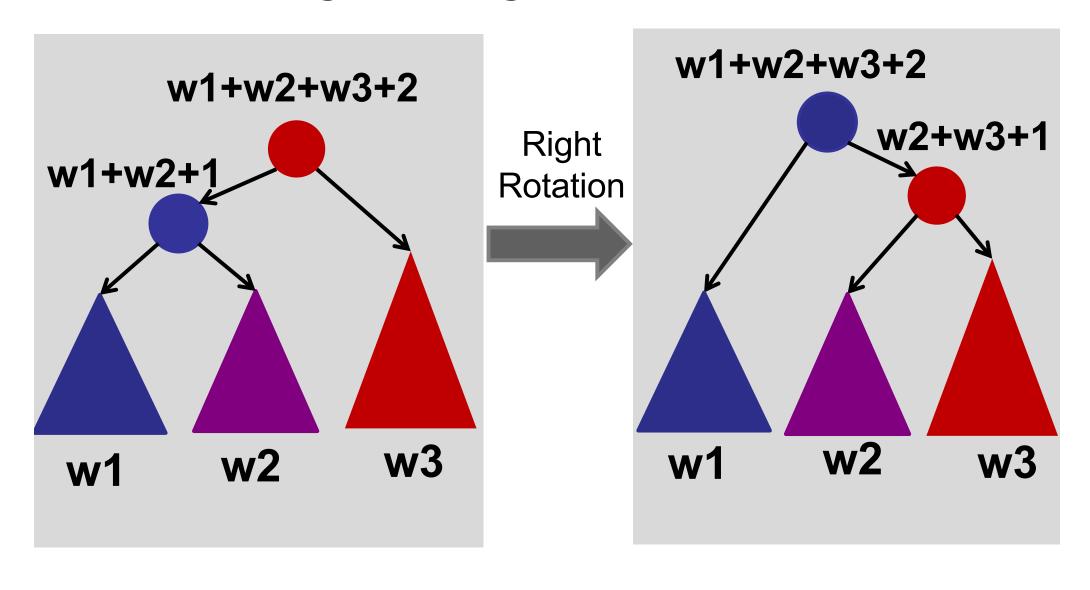




How long does it take to update the weights during a rotation?

- 1. O(1)
- 2. O(log n)
- 3. O(n)
- 4. $O(n^2)$
- 5. What is a rotation?





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(subject to insert/delete/etc.)

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Today

Three examples of augmenting balanced BSTs

1. Order Statistics

2. Intervals

3. Orthogonal Range Searching