

Part I) From these simple codes, determine their asymptotic running time

1) nums is a list of size n
 nums.append(1)

	cost	No. of Iteration	
	1	1	total = $1(1) = 1$

 $O(1)$

2) nums is a list of size n
 nums.insert(0,2)

	cost	No. of Iteration	
	1	1	total = $1(1) = 1$

 $O(1)$

3) seq is a list with n elements
 s = 0
 for x in seq:
 s += x

	cost	No. of Iteration	
s = 0	1	1	
for x in seq:	2	$n+1$	total = $1(1) + 2(n+1) + 2(n) = 4n + 3$
s += x	2	n	

 $O(n)$

4) seq is a list with n elements
 squares = [x**2 for x in seq]

	cost	No. of Iteration	
for x in seq:	2	$n+1$	
square = x**2	2	n	total = $2(n+1) + 2n = 4n + 2$

 $O(n)$

5) seq is a list with n elements
 s = 0
 for x in seq:
 for y in seq:
 s += x*y

	cost	No. of Iteration	
s = 0	1	1	
for x in seq:	2	$n+1$	total = $1(1) + 2(n+1) + n(n+1) + (n \times n)$
for y in seq:	2	$n \times (n+1)$	$= 2n^2 + 3n + 1$
s += x*y	3	$n \times n$	

 $O(n^2)$

6) seq is a list with n elements
 s = 0
 for x in seq:
 for y in seq:
 s += x*y
 for z in seq:
 for w in seq:
 s += x-w

	cost	No. of Iteration	
s = 0	1	1	
for x in seq:	2	$n+1$	total = $1(1) + 2(n+1) + n(n+1) +$
for y in seq:	2	$n(n+1)$	$(h \times n) + n(n+1) + n(n)(n+1)$
s += x*y	3	$n \times n$	$+ (n \times n \times n)$
for z in seq:	2	$n(n+1)$	$= 2n^3 + 4n^2 + 4n + 1$
for w in seq:	2	$n(n)(n+1)$	
s += x-w	3	$n \times n \times n$	

 $O(n^3)$

7) seq1 contains n elements and seq2 contains m elements
 s = 0
 for x in seq1:
 for y in seq2:
 s += x*y

	cost	No. of Iteration	
s = 0	1	1	total = $1(1) + 2(n+1) + n(m+1) + (n \times m)$
for x in seq1:	2	$n+1$	$= 2n^2 + 3n + 3$
for y in seq2:	2	$n(m+1)$	
s += x*y	3	$n \times m$	

 $O(n^2)$

8) seq1 = [[0,1],[2],[3,4,5]] cost No. of Iteration
 s = 0 1 1 total = $1(1) + 2(n+1) + 2(n)(n+1) + (n \times n)$
 for seq2 in seq1: 2 $n+1$
 for x in seq2: 2 $n(n+1)$
 s += x 2 $n \times n$
 $= 3n^2 + 4n + 3$

 $O(n^2)$

9) seq is a list with n elements cost No. of Iteration
 s = 0 1 1 total = $1(1) + 1(1) + 2n + 2(n)(n-1)$
 n = len(seq) n 1
 for i in range(n-1): i, n 2 n
 for j in range(i+1, n): 2 $n(n-1)$
 s += seq[i] * seq[j] 3 $(n-1)(n-1)$
 $= 3n^2 - 2n + 3$

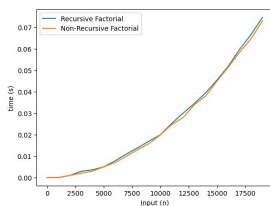
 $O(n^2)$

10) seq is a list with n elements Best case Worst case
 def sort_w_check(seq): cost No. of Iteration cost No. of Iteration
 n = len(seq) 1 1 1 1
 for i in range(n-1): 1 n 2 n
 if seq[i] > seq[i+1]: 1 n 2 n
 break 1 1
 else: total = $1(1) + 1(1) + 1(1)$ total = $1(1) + 2n + 2n + 1(1)$
 return = 3 = $4n + 2$
 ...

>>> sort_w_check(seq)

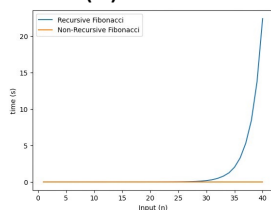
Best case : $O(1)$, Worst case : $O(n)$ **Part II) Write code and do some experiments for these problem**

1) Run a timing experiment for the recursive and non-recursive algorithm for computing $n!$ and discuss with the theoretical estimation



จากการกราฟจะเห็นว่าแบบ Recursive นาน Non-recursive Factorial
 ใช้เวลาใกล้เคียงกัน ได้ time complexity คือ $O(n)$

2) Run a timing experiment for the recursive and non-recursive algorithm for computing Fibonacci(n) and discuss with the theoretical estimation



จากการกราฟจะเห็นว่าแบบ Recursive นาน Non-recursive Fibonacci
 ใช้เวลาต่างกัน โดยแบบ Recursive จะใช้เวลาเพิ่มขึ้นเมื่อ $n > 30$ อย่าง
 เหนือขีด โดยได้ time complexity คือ $O(2^n)$ ในขณะที่แบบ Non-Recursive
 ใช้เวลาเท่าไร 0 วินาที จึงได้ time complexity คือ $O(1)$ ซึ่งดีกว่า

∴ การใช้ Algorithms ที่ต่างกันอาจทำให้ค่า Big O ออกมาต่างกัน
 จึงควรเลือก Algorithms ที่ได้ time complexity ดีที่สุด