

# Thesis

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## 1 Introduction

This is an introduction, welcome!

Introduce notation which will be used in text. A list of notation and description would be nice, so that the reader might scroll back up if something is unclear.

## 2 Triangulated Categories

Probably introduce this section, what is happening and what will be done etc. I can maybe say something about algebraic triangulated categories and topological triangulated categories, and explaining the name cone, fiber and cofiber.

### 2.1 Definition and First Properties

In this section  $\mathcal{T}$  denotes an additive category and  $T : \mathcal{T} \rightarrow \mathcal{T}$  is an additive autoequivalence of  $\mathcal{T}$ , which is often called translation or suspension functor.

**Definition 2.1.** A candidate triangle is a collection  $(A, B, C, a, b, c)$  of objects  $A, B, C \in \mathcal{T}$  and morphisms  $a : A \rightarrow B$ ,  $b : B \rightarrow C$ ,  $c : C \rightarrow TA$ . These candidate triangles can be drawn as diagrams in the following way:

$$A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA$$

A morphism between candidate triangles is a triple of morphism  $(\alpha, \beta, \gamma)$ , where  $\alpha : A \rightarrow A'$ ,  $\beta : B \rightarrow B'$  and  $\gamma : C \rightarrow C'$  such that the following diagram commutes:

$$\begin{array}{ccccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\ \downarrow \alpha & & \downarrow \beta & & \downarrow \gamma & & \downarrow T\alpha \\ A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C & \xrightarrow{c'} & TA' \end{array}$$

The naming convention of the candidate triangles isn't standardized, some literatures calls the candidate triangles for triangles instead; see [1]. This name arises from an alternate description of the diagrams given above. To remove confusion about the domain or codomain of the arrows, one arrow of the triangle is decorated with " $T$ ". This decorator means that the functor  $T$  has to be applied to the corresponding edge of the arrow. Thus the  $c$  arrow points to  $TA$ , not  $A$ .

$$\begin{array}{ccc} A & \xrightarrow{a} & B \\ T \uparrow & & \downarrow c \\ C & \xleftarrow{b} & A \end{array} \qquad \begin{array}{ccccccc} A & \xrightarrow{\phi_a} & A' & & & & \\ T \uparrow & \searrow a & & \swarrow a' & & T \uparrow & \\ & B & \xrightarrow{\phi_b} & B' & & & \\ & \swarrow b & & \searrow b' & & & \\ C & \xleftarrow{\phi_c} & C' & & & & \end{array}$$

A triangulated category is an additive category together with a translation functor  $T$  and a triangulation  $\Delta$  consisting of candidate triangles. When a candidate triangle is an element

of  $\Delta$  it is usually called a triangle, an exact triangles or just a triangle. Note that if the candidate triangles are referred to as triangles it is common to either call the elements of  $\Delta$  for triangles or exact triangles. As this is not the case for this thesis these objects will be referred to as triangles.

**Definition 2.2.** A triangulation of an additive category  $\mathcal{T}$  with translation  $T$  is a collection  $\Delta$  of triangles consisting of candidate triangles in  $\mathcal{T}$  satisfying the following axioms:

1. (TR1) Formation axiom

- (a) A candidate triangle isomorphic to a triangle is a triangle.
- (b) Every morphism  $a : A \rightarrow B$  can be embedded into a triangle  $(A, B, C, a, b, c)$ :

$$A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA$$

- (c) For every object  $A$  there is a triangle  $(A, A, 0, id_A, 0, 0)$ :

$$A \xrightarrow{id_A} A \xrightarrow{0} 0 \xrightarrow{0} TA$$

2. (TR2) Rotation axiom

For every triangle  $(A, B, C, a, b, c)$  there is a triangle  $(B, C, TA, b, c, Ta)$

$$A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA \implies B \xrightarrow{b} C \xrightarrow{c} TA \xrightarrow{-Ta} TB$$

3. (TR3) Morphism axiom

Given the two triangles  $(A, B, C, a, b, c)$  (1) and  $(A', B', C', a', b', c')$  (2)

$$(1) \quad A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA \quad (2) \quad A' \xrightarrow{a'} B' \xrightarrow{b'} C' \xrightarrow{c'} TA'$$

and morphisms  $\phi_A : A \rightarrow A'$  and  $\phi_B : B \rightarrow B'$  such that the square (1) commutes, then there is a morphism  $\phi_C : C \rightarrow C'$  (not necessarily unique) such that  $(\phi_A, \phi_B, \phi_C)$  is a morphism of triangles (2).

$$(1) \quad \begin{array}{ccc} A & \xrightarrow{a} & B \\ \downarrow \phi_A & & \downarrow \phi_B \\ A' & \xrightarrow{a'} & B' \end{array} \quad (2) \quad \begin{array}{ccccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\ \downarrow \phi_A & & \downarrow \phi_B & & \downarrow \phi_C & & \downarrow T\phi_A \\ A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C' & \xrightarrow{c'} & TA' \end{array}$$

4. (TR4) Octahedron axiom

Given the triangles  $(A, B, C', a, x, x')$  (1),  $(B, C, A', b, y, y')$  (2) and  $(A, C, B', b \circ a, z, z')$  (3)

$$(1) \quad A \xrightarrow{a} B \xrightarrow{x} C' \xrightarrow{x'} TA$$

$$(2) \quad B \xrightarrow{b} C \xrightarrow{y} A' \xrightarrow{y'} TB$$

$$(3) \quad A \xrightarrow{b \circ a} C \xrightarrow{z} B' \xrightarrow{z'} TA$$

then there exist morphisms  $f : C' \rightarrow B'$  and  $g : B' \rightarrow A'$ , the following diagram commutes and the third row is a triangle:

$$\begin{array}{ccccccc}
T^{-1}B' & \xrightarrow{T^{-1}z'} & A & \xrightarrow{id_A} & A & & \\
\downarrow T^{-1}g & & \downarrow a & & \downarrow b \circ a & & \\
T^{-1}A' & \xrightarrow{T^{-1}y'} & B & \xrightarrow{b} & C & \xrightarrow{y} & A' \xrightarrow{y'} TB \\
& & \downarrow x & & \downarrow z & & \parallel id_{A'} \downarrow Tx' \\
& & C' & \xrightarrow{f} & B' & \xrightarrow{g} & A' \xrightarrow{Ti \circ y'} TC' \\
& & \downarrow x' & & \downarrow z' & & \\
& & TA & \xrightarrow{id_{TA}} & TA & & 
\end{array}$$

A triangulated category is denoted as  $(\mathcal{T}, T, \Delta)$ , where  $\mathcal{T}$  is the additive category,  $T$  is the translation and  $\Delta$  is the triangulation. When  $\mathcal{T}$  is called a triangulated category, it should be understood as the triple given above.

*Remark.* The third object in a triangle is usually called cone, fiber, cofiber. These names are more historic than describing what these objects are, as the prototypes for these categories used these names for the third object. In this thesis it will be referred to as cone.

*Remark.* The rotation axiom has a dual, and it can be thought of as a rotation in the opposite direction. This dual can be proved by the other axioms, so it is here omitted as an axiom. The dual rotation axiom goes as:

Given a triangle  $A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA$ , there is a triangle  $T^{-1}C \xrightarrow{T^{-1}c} A \xrightarrow{a} B \xrightarrow{b} C$ . To be able to prove this, some more lemmata are needed.

*Remark.* The final axiom is referred to as the octahedron axiom. By using the alternative description of the triangle diagram, it is possible to rewrite the diagram as an octahedron. The axiom can be restated as the following:

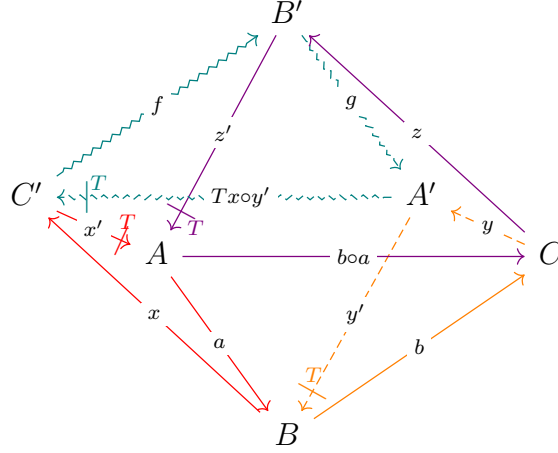
Given the triangles  $(A, B, C', a, x, x')$  (1),  $(B, C, A', b, y, y')$  (2) and  $(A, C, B', b \circ a, z, z')$  (3)

$$(1) \quad \begin{array}{ccc} & A & \\ \uparrow T & \searrow a & \\ x' & & B \\ \downarrow x' & \swarrow x & \\ & C' & \end{array}$$

$$(2) \quad \begin{array}{ccc} & B & \\ \uparrow T & \searrow b & \\ y' & & C \\ \downarrow y' & \swarrow y & \\ & A' & \end{array}$$

$$(3) \quad \begin{array}{ccc} & A & \\ \uparrow T & \searrow b \circ a & \\ z' & & C \\ \downarrow z' & \swarrow z & \\ & B' & \end{array}$$

then there exists morphisms  $f : C' \rightarrow B'$  and  $g : B' \rightarrow A'$ , the following diagram commutes and the squiggly teal back face is a triangle.

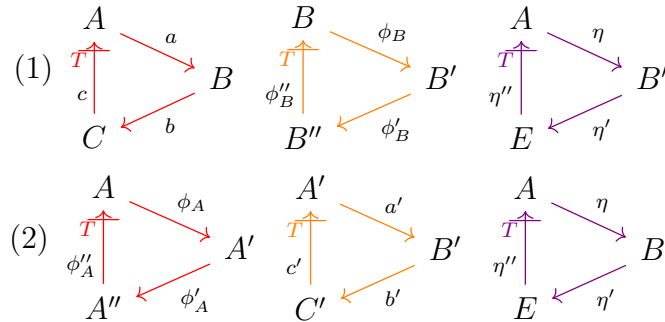


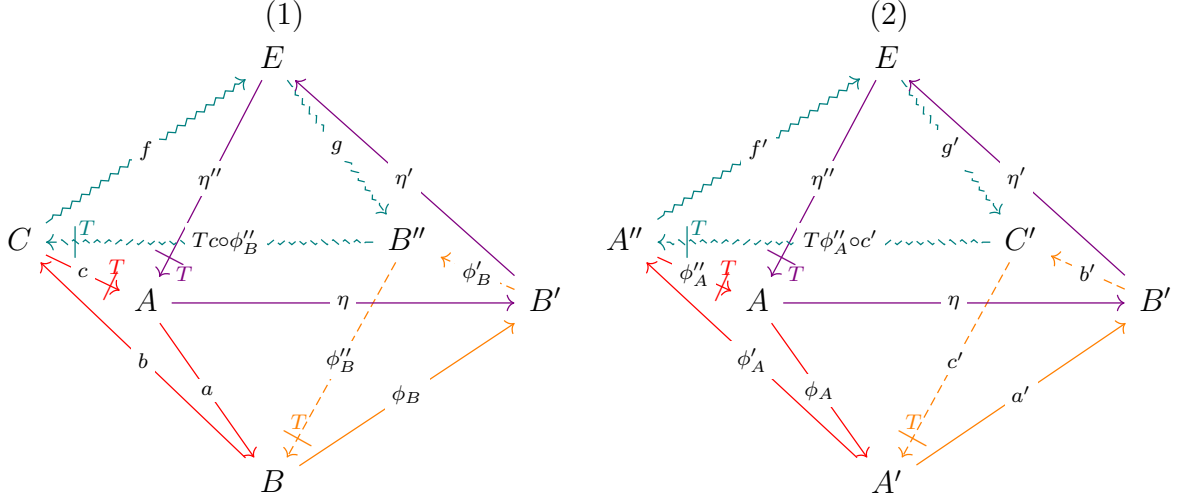
**Proposition 2.1.** *The axiom TR3 can be proven from TR1 and TR4.*

*Proof.* Suppose that there are two triangles and a commutative square as follows.

$$\begin{array}{ccc}
 A & \xrightarrow{a} & B \\
 \downarrow \phi_A & \searrow \eta & \downarrow \phi_B \\
 A' & \xrightarrow{a'} & B'
 \end{array}
 \qquad
 \begin{array}{ccccccc}
 A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\
 \downarrow \phi_A & & \downarrow \phi_B & & & & \downarrow T\phi_A \\
 A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C' & \xrightarrow{c'} & TA'
 \end{array}$$

Now we are able to use the Octahedron axiom twice after the square has been completed to triangles in the following sense. The upper commutative simplex and the lower commutative simplex form the following diagrams, and thus the following octahedra:





following the blue squiggly lines there is a morphism  $g'f : C \rightarrow C'$ , it remains to see that this fits into the triangle morphism. By chasing the different paths in the octahedra and combining diagrams we get that the following diagrams are commutative, which is exactly the requirement for the collection  $(\phi_A, \phi_B, g'f, T\phi_A)$  to be a morphism of triangles.

$$\begin{array}{ccccc}
 B & \xrightarrow{b} & C & & \\
 \downarrow \phi'_B & & \downarrow f & \searrow c & \\
 B' & \xrightarrow{\eta'} & E & \xrightarrow{\eta''} & TA \\
 & \searrow b' & \downarrow g' & & \downarrow Ta \\
 & & C' & \xrightarrow{c'} & TA'
 \end{array}$$

□

**Lemma 2.2.** *Let  $(A, B, C, a, b, c)$  be a triangle, then  $b \circ a = 0$*

*Proof.* By TR2 the triangle  $(A, B, C, a, b, c)$  can be rotated to  $(B, C, TA, b, c, Ta)$ .

$$\begin{array}{ccc}
 A & \xrightarrow{a} & B \\
 \uparrow T & & \uparrow T \\
 C & \xleftarrow{b} & B
 \end{array}
 \implies
 \begin{array}{ccc}
 B & \xrightarrow{b} & C \\
 \uparrow T & & \uparrow T \\
 TA & \xleftarrow{c} & C
 \end{array}$$

The triangle exists  $(C, C, 0, id_C, 0, 0)$  by TR1 and TR3 says there exists a morphism from  $TA$  to  $0$  making the diagram below commute.

$$\begin{array}{ccccccc}
 B & \xrightarrow{b} & C & \xrightarrow{c} & TA & \xrightarrow{-Ta} & TB \\
 \downarrow b & & \downarrow id_C & & \downarrow 0 & & \downarrow Tb \\
 C & \xrightarrow{id_C} & C & \xrightarrow{0} & 0 & \xrightarrow{0} & TC
 \end{array}$$

Thus  $0 = Tb \circ -Ta = T(-ba) \implies b \circ a = 0$  as  $T$  is a translation.  $\square$

**Definition 2.3.** An additive functor between triangulated categories  $F : (\mathcal{T}, T, \Delta) \rightarrow (\mathcal{R}, R, \Gamma)$  is called exact or triangulated if there exist a natural isomorphisms  $\alpha : FT \rightarrow RF$  such that  $F(\Delta) \subseteq \Gamma$ .

A functor  $F : \mathcal{T} \rightarrow \mathcal{R}$  is called a triangle-equivalence if it is triangulated and an equivalence of categories. In this case  $\mathcal{T}$  and  $\mathcal{R}$  are called triangle-equivalent.

**Definition 2.4.** Let  $\mathcal{T}$  be a triangulated category and  $\mathcal{A}$  be an abelian category. A covariant functor  $H : \mathcal{T} \rightarrow \mathcal{A}$  is called a homological functor if  $\forall(A, B, C, a, b, c) : \Delta$  there is a long exact sequence in  $\mathcal{A}$ .

$$\begin{array}{c} A \\ \uparrow T \\ C \end{array} \begin{array}{c} \nearrow a \\ \searrow b \\ \downarrow c \end{array} B \implies \begin{array}{c} \dots \longrightarrow H(T^i A) \xrightarrow{H(T^i a)} H(T^i B) \xrightarrow{H(T^i b)} H(T^i C) \\ \hspace{10em} \searrow H(T^i c) \\ \hspace{10em} \longrightarrow H(T^{i+1} A) \xrightarrow{H(T^{i+1} a)} H(T^{i+1} B) \xrightarrow{H(T^{i+1} b)} H(T^{i+1} C) \longrightarrow \dots \end{array}$$

Dually, a contravariant functor  $H : \mathcal{T} \rightarrow \mathcal{A}$  is called cohomological if  $\forall(A, B, C, a, b, c) : \Delta$  there is a long exact sequence in  $\mathcal{A}$ .

$$\begin{array}{c} A \\ \uparrow T \\ C \end{array} \begin{array}{c} \nearrow a \\ \searrow b \\ \downarrow c \end{array} B \implies \begin{array}{c} \dots \longleftarrow H(T^{i-1} A) \xleftarrow{H(T^{i-1} a)} H(T^{i-1} B) \xleftarrow{H(T^{i-1} b)} H(T^{i-1} C) \\ \hspace{10em} \searrow H(T^{i-1} c) \\ \hspace{10em} \longleftarrow H(T^i A) \xleftarrow{H(T^i a)} H(T^i B) \xleftarrow{H(T^i b)} H(T^i C) \longleftarrow \dots \end{array}$$

**Lemma 2.3.** Let  $M : \mathcal{T}$  be any object of  $\mathcal{T}$ , then the represented functors  $\mathcal{T}(M, -)$  is homological and  $\mathcal{T}(-, M)$  is cohomological.

*Proof.* Only the covariant case needs to be proved, as the contravariant case is dual. For  $\mathcal{T}(M, -)$  to be homological, it has to create long exact sequences for every triangle in  $\Delta$ . Let  $(A, B, C, a, b, c) : \Delta$  be a triangle, then there can be extracted sequences in  $\text{Ab}$  for any  $i : \mathbb{N}$ .

$$\begin{array}{c} A \\ \uparrow T \\ C \end{array} \begin{array}{c} \nearrow a \\ \searrow b \\ \downarrow c \end{array} B \implies \mathcal{T}(M, T^i A) \xrightarrow{T^i a_*} \mathcal{T}(M, T^i B) \xrightarrow{T^i b_*} \mathcal{T}(M, T^i C)$$

Observe that it is enough to prove that these types of diagrams are exact, as the other diagrams can be obtained by the rotation axiom, thus reducing it to same case.

The goal is then to prove that  $\text{Im}(T^i a_*) = \text{Ker}(T^i b_*)$ . Since  $ba = 0$  it follows that  $\text{Im}(T^i a_*) \subseteq \text{Ker}(T^i b_*)$ . Assume that  $f : \text{Ker}(T^i b_*)$ , that is  $f : M \rightarrow T^i B$  such that



$b_*(f) = 0$ . The current goal is to show that  $f$  factors through  $T^i A$ , as this means that  $\text{Ker}(T^i b_*) \subseteq \text{Im}(T^i a_*)$ . Note that since  $T$  is a translation, it is necessarily a right adjoint to the inverse translation; thus  $\mathcal{T}(M, T^i B) \simeq \mathcal{T}(T^{-i} M, B)$ , and by this assertion it suffices to assume that  $f : T^{-i} M \rightarrow B$  such that  $b \circ f = 0$ . By TR1 and TR2 there exists triangles  $(T^{-i} M, 0, T^{-i+1} M, 0, 0, -T^{-i+1} \text{id})$  and  $(B, C, TA, b, c, -Ta)$ .

$$\begin{array}{ccccccc} T^{-i} M & \xrightarrow{0} & 0 & \xrightarrow{0} & T^{-i+1} M & \xrightarrow{T^{-i+1} \text{id}} & T^{-i+1} M \\ \downarrow f & & \downarrow 0 & & \downarrow g & & \downarrow Tf \\ B & \xrightarrow{b} & C & \xrightarrow{c} & TA & \xrightarrow{-Ta} & TB \end{array}$$

The left square commutes by the assumption, thus the morphism  $g$  exist by TR3, such that  $-Ta \circ h = -Tf \circ T^{-i+1} \text{id} = -Tf \implies Ta \circ h = Tf$ , thus  $f = a \circ T^{-1} h$  asserting that  $f$  factors through  $A$ .  $\square$

**Lemma 2.4.** *Let  $(\phi_A, \phi_B, \phi_C) : (A, B, C, a, b, c) \rightarrow (A', B', C', a', b', c')$  be a morphism of triangles. If 2 of the maps are isomorphisms, then the last one is an isomorphism as well.*

$$\begin{array}{ccccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\ \wr \downarrow \phi_A & & \wr \downarrow \phi_B & & \wr \downarrow \phi_C & & \wr \downarrow T\phi_A \\ A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C' & \xrightarrow{c'} & TA' \end{array}$$

*Proof.* Without loss of generality, assume that  $\phi_A$  and  $\phi_B$  are the isomorphisms. This can be done as the rotation axiom reduce the other cases to this case. Then we have the following diagram:

$$\begin{array}{ccccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\ \wr \downarrow \phi_A & & \wr \downarrow \phi_B & & \downarrow \phi_C & & \wr \downarrow T\phi_A \\ A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C' & \xrightarrow{c'} & TA' \end{array}$$

By applying the functor  $\mathcal{T}(C', -)$  we get the following diagram in  $\text{Ab}$ :

$$\begin{array}{ccccccccc} \mathcal{T}(C', A) & \xrightarrow{a_*} & \mathcal{T}(C', B) & \xrightarrow{b_*} & \mathcal{T}(C', C) & \xrightarrow{c_*} & \mathcal{T}(C', TA) & \xrightarrow{Ta_*} & \mathcal{T}(C', TB) \\ \wr \downarrow (\phi_A)_* & & \wr \downarrow (\phi_B)_* & & \downarrow (\phi_C)_* & & \wr \downarrow (\phi_{TA})_* & & \wr \downarrow (T\phi_B)_* \\ \mathcal{T}(C', A') & \xrightarrow{a'_*} & \mathcal{T}(C', B') & \xrightarrow{b'_*} & \mathcal{T}(C', C') & \xrightarrow{c'_*} & \mathcal{T}(C', TA') & \xrightarrow{Ta'_*} & \mathcal{T}(C', TB) \end{array}$$

By the 5-lemma, we get that  $(\phi_C)_*$  is an isomorphisms, i.e.  $(\phi_C)_*$  is both mono and epi. Thus for some unique  $s$  in  $\mathcal{T}(C', C)$ ,  $\phi_{C*}(s) = \text{id}_{C'}$ .

By applying the functor  $\mathcal{T}(-, C)$  we get the diagram:

$$\begin{array}{ccccccccc}
\mathcal{T}(A, C) & \xleftarrow{a^*} & \mathcal{T}(B, C) & \xleftarrow{b^*} & \mathcal{T}(C, C) & \xleftarrow{c^*} & \mathcal{T}(TA, C) & \xleftarrow{Ta^*} & \mathcal{T}(TB, C) \\
(\phi_A)^* \uparrow \wr & & (\phi_B)^* \uparrow \wr & & (\phi_C)^* \uparrow & & (\phi_{TA})^* \uparrow \wr & & (\phi_{TB})^* \uparrow \wr \\
\mathcal{T}(A', C) & \xleftarrow{a'^*} & \mathcal{T}(B, C) & \xleftarrow{b'^*} & \mathcal{T}(C', C) & \xleftarrow{c'^*} & \mathcal{T}(TA', C) & \xleftarrow{Ta'^*} & \mathcal{T}(TB', C)
\end{array}$$

By the 5-lemma, we get that  $(\phi_C)^*$  is an isomorphisms. By the same argument  $id_C = s' \circ \phi_C$  for some unique  $s'$ .  $\phi_C$  is both split mono and split epi, which means it is an isomorphism.  $\square$

**Corollary 2.4.1.**  $(A, B, 0, a, 0, 0)$  is a triangle if and only if  $a$  is an isomorphism.

*Proof.* Assume that  $a$  is an isomorphism. Then it is seen that  $(a, id_B, 0)$  is an isomorphism of triangles.

$$\begin{array}{ccccccc}
A & \xrightarrow{a} & B & \xrightarrow{0} & 0 & \xrightarrow{0} & TA \\
\wr \downarrow a & & \wr \downarrow id_B & & \wr \downarrow 0 & & \wr \downarrow Ta \\
B & \xrightarrow{id_B} & B & \xrightarrow{0} & 0 & \xrightarrow{0} & TB
\end{array}$$

Converesly, assume that  $(A, B, 0, a, 0, 0)$  is a triangle. Then the same diagram as above can be constructed, and by the 2 out of 3 property,  $a$  has to be an isomorphism.  $\square$

**Lemma 2.5.** For a triangle  $(A, B, C, a, b, c)$  the following are equivalent:

$$\begin{array}{ccc}
A & & \\
\uparrow T & \searrow a & \\
C & \swarrow b & B \\
& c &
\end{array}$$

- $a$  is split mono
- $b$  is split epi
- $c = 0$

*Proof.* The proof has two parts. First assume that  $a$  is split mono, and prove that  $b$  is split epi, and  $c = 0$ . By duality, it is then known that  $b$  being split epi implies that  $a$  is split mono and  $c = 0$ . The final part is to assume that  $c = 0$ , and prove either  $a$  is split mono or  $b$  is split epi.

Assume that  $a$  is split mono, then there exist an  $a^{-1}$  such that  $id_A = a^{-1}a$ . Let  $M : \mathcal{T}$  be any object, then we can make a long exact sequence:

$$\begin{array}{ccccccc}
\mathcal{T}(M, T^{-1}C) & \xrightarrow{T^{-1}c_*} & \mathcal{T}(M, A) & \xrightarrow{a_*} & \mathcal{T}(M, B) & \xrightarrow{b_*} & \mathcal{T}(M, C) \xrightarrow{c_*} \mathcal{T}(M, TA) \\
& & & \curvearrowright & & & \\
& & & a_*^{-1} & & &
\end{array}$$

By assumption  $a_*$  is split mono, thus  $T^{-1}c_* = 0$  and in particular  $c = 0$ . This implies that  $b_*$  is epi, making a split short exact sequence.

$$\begin{array}{ccccccc}
0 & \xrightarrow{0} & \mathcal{T}(M, A) & \xrightarrow{a_*} & \mathcal{T}(M, B) & \xrightarrow{b_*} & \mathcal{T}(M, C) \xrightarrow{0} 0 \\
& & & \nwarrow \text{dashed} & \nwarrow \text{dashed} & & \\
& & & a_*^{-1} & b_*^{-1} & & 
\end{array}$$

This gives that  $b$  is split epi, completing the first part of the proof.

For the next part, assume that  $c = 0$ ; then we can construct the following triangles.

$$\begin{array}{ccc}
(1) \quad \begin{array}{ccc} A & \xrightarrow{a} & B \\ T \uparrow & & \\ 0 \downarrow & & \\ C & \xleftarrow{b} & \end{array} & \Rightarrow & \begin{array}{ccc} C & \xrightarrow{0} & TA \\ T \uparrow & & \\ -Tb \downarrow & & \\ TB & \xleftarrow{-Ta} & \end{array} \\
(2) \quad \begin{array}{ccc} A & \xrightarrow{id_A} & A \\ T \uparrow & & \\ 0 \downarrow & & \\ 0 & \xleftarrow{0} & \end{array} & \Rightarrow & \begin{array}{ccc} 0 & \xrightarrow{0} & TA \\ T \uparrow & & \\ 0 \downarrow & & \\ TA & \xleftarrow{-id_{TA}} & \end{array}
\end{array}$$

(1) is constructed by applying TR2 twice, while (2) is constructed with TR1 and TR2 twice. Observe that there is a commutative square between the triangles, allowing for TR3 to make a morphism of triangles.

$$\begin{array}{ccccccc}
C & \xrightarrow{0} & TA & \xrightarrow{-Ta} & TB & \xrightarrow{-Tb} & TC \\
\downarrow 0 & & \parallel id_{TA} & & \downarrow Ta^{-1} & & \downarrow 0 \\
0 & \xrightarrow{0} & TA & \xrightarrow{-id_{TA}} & TA & \xrightarrow{0} & 0
\end{array}$$

Thus  $T(a^{-1}a) = id_{TA} = T(id_A) \Rightarrow id_A = a^{-1}a$ , making a split mono.  $\square$

**Lemma 2.6.** *Given two triangles  $(A, B, C, a, b, c)$  and  $(A', B', C', a', b', c')$  the following are equivalent:*

$$\begin{array}{ccc}
A & \xrightarrow{a} & B \xrightarrow{b} C \xrightarrow{c} TA \\
\downarrow f & & \downarrow g \quad \downarrow h \quad \downarrow Tf \\
A' & \xrightarrow{a'} & B' \xrightarrow{b'} C' \xrightarrow{c'} TA'
\end{array}
\quad \begin{array}{l}
1. (f, g, h) \text{ is a morphism of triangles} \\
2. \exists g : B \rightarrow B' \text{ such that } b'ga = 0
\end{array}$$

Moreover, if  $\mathcal{T}(A, T^{-1}C') \simeq 0$ , then  $f$  and  $h$  are unique.

*Proof.* 1.  $\Rightarrow$  2. as the composition  $ba = 0 = b'a'$ , so assume 2. The existence of  $f$  and  $h$  is evident from the long exact sequence of the bottom triangle at the functor represented by  $A$ .

$$\mathcal{T}(A, T^{-1}C') \xrightarrow{T^{-1}c'_*} \mathcal{T}(A, A') \xrightarrow{a'_*} \mathcal{T}(A, B') \xrightarrow{b'_*} \mathcal{T}(A, C')$$

The morphism  $ga : \mathcal{T}(A, B')$  such that  $b'_*ga = b'_*(ga) = 0$ , thus  $ga : \text{Ker}(b'_*)$ . By exactness  $\exists f : \mathcal{T}(A, A')$  such that  $a'_*f = ga$ , and by TR3  $\exists h : C \rightarrow C'$  such that  $(f, g, h)$  is a morphism of triangles. Now assume that  $\mathcal{T}(A, T^{-1}C') \simeq 0$ . Exactness determines that  $a'_*$  is a monomorphism, and  $f$  is then unique. Since  $T$  is a translation, we have that  $\mathcal{T}(A, T^{-1}C') \simeq \mathcal{T}(TA, C')$ . By using the functor  $\mathcal{T}(-, C')$  at the top triangle, we get that  $b^*$  is a monomorphism, and thus  $h$  is chosen uniquely.  $\square$

**Lemma 2.7.** *If  $(A, B, C, a, b, c)$  is a triangle, then  $(T^{-1}C, A, B, T^{-1}c, a, b)$  is a triangle.*

*Proof.* By TR2 we can construct a triangle.

$$\begin{array}{ccc} A & \xrightarrow{a} & B \\ \uparrow T & & \uparrow T \\ C & \xleftarrow{b} & TB \end{array} \implies \begin{array}{ccc} C & \xrightarrow{c} & TA \\ \uparrow T & & \uparrow T \\ TB & \xleftarrow{Ta} & TC \end{array}$$

By TR1 we can create a triangle  $(T^{-1}C, A, B, T^{-1}c, a', b')$ , and then use TR3 to find a morphism.

$$\begin{array}{ccccccc} C & \xrightarrow{c} & TA & \xrightarrow{Ta'} & TB' & \xrightarrow{Tb'} & TC \\ \parallel id_C & & \parallel id_{TA} & & \downarrow h & & \parallel id_{TC} \\ C & \xrightarrow{c} & TA & \xrightarrow{Ta} & TB & \xrightarrow{Tb} & TC \end{array}$$

By the 2 out of 3 property it is seen that  $h$  is an isomorphism, so the triple  $(id_{T^{-1}C}, id_A, T^{-1}h)$  is an isomorphism of candidate triangles, and by TR1, is an isomorphism of triangles, asserting that  $(T^{-1}C, A, B, T^{-1}c, a, b)$  is in fact a triangle.  $\square$

**Lemma 2.8.** *Let  $(A, B, C, a, b, c)$  and  $(A', B', C', a', b', c')$  be two triangles, then the direct sum of these triangles is a triangle.*

*Proof.* Observe that for any functor  $\mathcal{T}(K, -)$  there is still a long exact sequence of Hom(ology) since  $\mathcal{T}(K, A \oplus A') \simeq \mathcal{T}(K, A) \oplus \mathcal{T}(K, A')$ . Thus for the direct sum of the triangles we have the following.

$$\begin{array}{ccccc} A \oplus A' & \xrightarrow{\begin{pmatrix} a & 0 \\ 0 & a' \end{pmatrix}} & B \oplus B' & \xrightarrow{\begin{pmatrix} b & 0 \\ 0 & b' \end{pmatrix}} & C \oplus C' & \xrightarrow{\begin{pmatrix} c & 0 \\ 0 & c' \end{pmatrix}} & TA \oplus TC \\ & & & \Downarrow & & & \end{array}$$

$$\begin{array}{c}
\cdots \longrightarrow \mathcal{T}(K, A) \oplus \mathcal{T}(K, A') \longrightarrow \mathcal{T}(K, B) \oplus \mathcal{T}(K, B') \\
\downarrow \\
\mathcal{T}(K, C) \oplus \mathcal{T}(K, C') \longrightarrow \mathcal{T}(K, TA) \oplus \mathcal{T}(K, TA') \longrightarrow \cdots
\end{array}$$

Thus the 2 out of 3 property holds for the direct sum. By TR1 there is a triangle

$$A \oplus A' \longrightarrow B \oplus B' \longrightarrow D \longrightarrow TA \oplus TA'$$

By TR3 there are morphisms from this triangle to the direct summands. Adding these maps together there is a map from this triangle to direct sum, and by using the 2 out of 3 property this is an isomorphism of candidate triangles. Thus the direct sum is a triangle.

$$\begin{array}{ccccccc}
A \oplus A' & \longrightarrow & B \oplus B' & \longrightarrow & D & \longrightarrow & TA \oplus TA' \\
\downarrow & & \downarrow & & \downarrow & & \downarrow \\
A & \longrightarrow & B & \longrightarrow & C & \longrightarrow & TA \\
& & & & \& & \\
A \oplus A' & \longrightarrow & B \oplus B' & \longrightarrow & D & \longrightarrow & TA \oplus TA' \\
\downarrow & & \downarrow & & \downarrow & & \downarrow \\
A' & \longrightarrow & B' & \longrightarrow & C' & \longrightarrow & TA' \\
& & & & \Downarrow & & \\
A \oplus A' & \longrightarrow & B \oplus B' & \longrightarrow & D & \longrightarrow & TA \oplus TA' \\
\parallel & & \parallel & & \downarrow \wr & & \parallel \\
A \oplus A' & \longrightarrow & B \oplus B' & \longrightarrow & A'' \oplus B'' & \longrightarrow & TA \oplus TA'
\end{array}$$

□

**Lemma 2.9.** *The direct summands of a triangle is a triangle.*

*Proof.* The proof can be found in [2]

□

## 2.2 Mapping Cones, Homotopies and Contractibility

*Remark.* The observant reader might have seen that the Octahedron axiom have not yet been used once! A lot of the theory proven for triangulated categories work without this axiom, and this motivates the definition of a pre-triangulated category. This odd axiom might seem to be pointless, but it is one of the most fundamental axioms, as TR3 can be proven with TR1, TR2 and TR4; see [need sources here].

**Definition 2.5.** A pre-triangulation of an additive category  $\mathcal{T}$  with translation  $T$  is a collection  $\Delta'$  of triangles consisting of candidate triangles in  $\mathcal{T}$  satisfying the following axioms:

1. (TR1) Formation axiom

- (a) A candidate triangle isomorphic to a triangle is a triangle.
- (b) Every morphism  $a : A \rightarrow B$  can be embedded into a triangle  $(A, B, C, a, b, c)$ :

$$A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA$$

- (c) For every object  $A$  there is a triangle  $(A, A, 0, id_A, 0, 0)$ :

$$A \xrightarrow{id_A} A \xrightarrow{0} 0 \xrightarrow{0} TA$$

2. (TR2) Rotation axiom

For every triangle  $(A, B, C, a, b, c)$  there is a triangle  $(B, C, TA, b, c, Ta)$

$$A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA \implies B \xrightarrow{b} C \xrightarrow{c} TA \xrightarrow{-Ta} TB$$

3. (TR3) Morphism axiom

Given the two triangles  $(A, B, C, a, b, c)$  (1) and  $(A', B', C', a', b', c')$  (2)

$$(1) \quad A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} TA \quad (2) \quad A' \xrightarrow{a'} B' \xrightarrow{b'} C' \xrightarrow{c'} TA'$$

and morphisms  $\phi_A : A \rightarrow A'$  and  $\phi_B : B \rightarrow B'$  such that the square (1) commutes, then there is a morphism  $\phi_C : C \rightarrow C'$  (not necessarily unique) such that  $(\phi_A, \phi_B, \phi_C)$  is a morphism of triangles (2).

$$(1) \quad \begin{array}{ccc} A & \xrightarrow{a} & B \\ \downarrow \phi_A & & \downarrow \phi_B \\ A' & \xrightarrow{a'} & B' \end{array} \quad (2) \quad \begin{array}{ccccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\ \downarrow \phi_A & & \downarrow \phi_B & & \downarrow \phi_C & & \downarrow T\phi_A \\ A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C' & \xrightarrow{c'} & TA' \end{array}$$

The category  $\mathcal{T}$  with the pre-triangulation  $\Delta'$  is called a pre-triangulated category, and the candidate triangles in  $\Delta'$  are called triangles. We will only use this notion of triangles in this subsection.

The main goal of this subsection is to see how we can find triangles, and when these are triangles. We will also look at triangulated functors and triangulated subcategories. For the rest of this subsection it is assumed that we work in a pre-triangulated category  $\mathcal{T}$ .

**Definition 2.6.** Suppose there is a morphism of candidate triangles  $\phi : (A, B, C, a, b, c) \rightarrow (A', B', C', a', b', c')$ .

$$\begin{array}{ccccccc}
A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\
\downarrow \phi_A & & \downarrow \phi_B & & \downarrow \phi_C & & \downarrow T\phi_A \\
A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C & \xrightarrow{c'} & TA'
\end{array}$$

We define the mapping cone to be the candidate triangle:

$$B \oplus A' \xrightarrow{\begin{pmatrix} -b & 0 \\ \phi_B & a' \end{pmatrix}} C \oplus B' \xrightarrow{\begin{pmatrix} -c & 0 \\ \phi_C & b' \end{pmatrix}} TA \oplus C' \xrightarrow{\begin{pmatrix} -Ta & 0 \\ T\phi_A & c' \end{pmatrix}} TB \oplus TA'$$

**Definition 2.7.** A morphism  $\alpha : A, B, C, a, b, c \rightarrow (A', B', C', a', b', c')$  between candidate triangles is called null-homotopic if it factors through a homotopy. That is, there exists maps  $\Theta, \Phi, \Psi$  in the following diagram:

$$\begin{array}{ccccccc}
A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\
\downarrow \alpha_A & \swarrow \Theta & \downarrow \alpha_B & \swarrow \Phi & \downarrow \alpha_C & \swarrow \Psi & \downarrow T\alpha_A \\
A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C & \xrightarrow{c'} & TA'
\end{array}$$

such that  $\alpha_A = \Theta a + T^{-1}(c' \Psi)$ ,  $\alpha_B = \Phi b + a' \Theta$  and  $\alpha_C = \Psi c + b' \Phi$ . Two maps are called homotopic if their difference is null-homotopic

**Lemma 2.10.** *The mapping cone only depends on morphisms up to homotopy. I.e. if two maps are homotopic, their mapping cones are isomorphic.*

*Proof.* Suppose that  $(f, g, h)$  and  $(f', g', h')$  are two homotopic morphisms of triangles:

$$\begin{array}{ccccccc}
A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\
\downarrow & & \downarrow & & \downarrow & & \downarrow \\
A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C' & \xrightarrow{c'} & TA'
\end{array}$$

Let  $(\Theta, \Phi, \Psi)$  be the homotopy between the triangle morphisms. Then there is an isomorphism of triangles.

$$\begin{array}{ccccccc}
B \oplus A' & \xrightarrow{\begin{pmatrix} -b & 0 \\ g & a' \end{pmatrix}} & C \oplus B' & \xrightarrow{\begin{pmatrix} -c & 0 \\ h & b' \end{pmatrix}} & TA \oplus C' & \xrightarrow{\begin{pmatrix} -Ta & 0 \\ Tf & c' \end{pmatrix}} & TB \oplus TA' \\
\downarrow \begin{pmatrix} 1 & 0 \\ \Theta & 1 \end{pmatrix} & & \downarrow \begin{pmatrix} 1 & 0 \\ \Phi & 1 \end{pmatrix} & & \downarrow \begin{pmatrix} 1 & 0 \\ \Psi & 1 \end{pmatrix} & & \downarrow \begin{pmatrix} 1 & 0 \\ T\Theta & 1 \end{pmatrix} \\
B \oplus A' & \xrightarrow{\begin{pmatrix} -b & 0 \\ g' & a' \end{pmatrix}} & C \oplus B' & \xrightarrow{\begin{pmatrix} -c & 0 \\ h' & b' \end{pmatrix}} & TA \oplus C' & \xrightarrow{\begin{pmatrix} -Ta & 0 \\ Tf' & c' \end{pmatrix}} & TB \oplus TA'
\end{array}$$

□

**Lemma 2.11.** *Let  $A$  denote the triangle  $(A, A', A'')$  and  $B$  denote  $(B, B', B'')$ . Suppose  $\alpha, \beta : A \rightarrow B$  are two homotopic morphisms of candidate triangles. Then for any map  $\gamma : \tilde{A} \rightarrow A$  and any map  $\delta : B \rightarrow \tilde{B}$  the maps  $\delta\alpha\gamma$  and  $\delta\beta\gamma$  are homotopic as well.*

*Proof.* To prove this statement it is enough to prove that  $\alpha\gamma$  is homotopic to  $\beta\gamma$  due to the symmetry of the statement. The goal is then to show that  $(\Theta\gamma', \Phi\gamma'', \Psi T\gamma)$  is the homotopy. This can be seen as

$$\alpha'\gamma' - \beta'\gamma' = (\alpha' - \beta')\gamma' = (b\Theta + \Phi a')\gamma' = b\Theta\gamma' + \Phi a'\gamma' = b(\Theta\gamma') + (\Phi\gamma'')\tilde{a}'$$

□

**Definition 2.8.** A candidate triangle  $A$  is called a contractible triangle if  $id_A$  is null-homotopic.

*Remark.* If  $A$  is a contractible triangle and  $F : \mathcal{T} \rightarrow \mathcal{A}$  is an additive functor to an abelian category, then the identity of the cochain

$$\dots \longrightarrow F(A) \longrightarrow F(A') \longrightarrow F(A'') \longrightarrow F(TA) \longrightarrow \dots$$

is null-homotopic as well. Thus the homology of this sequence is 0, asserting that it is an exact sequence. The exactness of such sequences allow us to use the 2 out of 3 property on morphisms between contractible triangles.

**Corollary 2.11.1.** *If  $A$  is a contractible triangle, then any map in  $\mathcal{T}(A, \_)$  or  $\mathcal{T}(\_, A)$  is null-homotopic.*

*Proof.* Being null-homotopic is the same as that there is a homotopy between the map and the zero map. If  $id_A \sim 0 \implies f \circ id_A = f \sim f \circ 0 = 0$ . So any map  $f$  is null-homotopic. □

**Lemma 2.12.** *A contractible triangle is a triangle.*

*Proof.* Let  $A$  be the triangle  $(A, A', A'')$  such that  $id_A \sim 0$ . This means that there is a homotopy

$$\begin{array}{ccccccc} A & \xrightarrow{a} & A' & \xrightarrow{a'} & A'' & \xrightarrow{a''} & TA \\ \downarrow id_A & \swarrow \Theta & \downarrow id_{A'} & \swarrow \Phi & \downarrow id_{A''} & \swarrow \Psi & \downarrow id_{TA} \\ A & \xrightarrow{a} & A' & \xrightarrow{a'} & A'' & \xrightarrow{a''} & TA \end{array}$$

By using TR1 there is a triangle, and a long exact sequence.



$$\begin{array}{ccccccc}
A & \xrightarrow{a} & A' & \xrightarrow{e} & E & \xrightarrow{e'} & TA \\
& & & \Downarrow & & & \\
\cdots & \longrightarrow & \mathcal{T}(TA, A) & \longrightarrow & \mathcal{T}(TA, A') & \longrightarrow & \mathcal{T}(TA, E) \xrightarrow{e'_*} \mathcal{T}(TA, TA) \xrightarrow{Ta_*} \cdots
\end{array}$$

Since the map  $Ta \circ a''\Psi = 0$ , the kernel  $Ker Ta_* = Im e'_* \neq 0$ . This means that there is a map  $\Psi' : \mathcal{T}(TA, E)$  such that  $e'\Psi' = a''\Psi$ . Then the map  $(id_A, id_{A'}, e\Theta + \Psi'a'')$  is a well defined map of candidate triangles. By the remark, we can use the 2 out of 3 properties to assert that the map found is an isomorphism, giving an isomorphism of triangles, asserting that the contractible triangle is a triangle.  $\square$

**Corollary 2.12.1.** *The mapping cone of the zero map between triangles is a triangle.*

*Proof.* The mapping cone of the zero map can be seen to be the direct sum of two triangles. Thus it is a triangle.  $\square$

**Corollary 2.12.2.** *The mapping cone of a null-homotopic map between triangles is a triangle.*

*Remark.* Suppose we have a morphism of triangles where one of the triangles are contractible, then the mapping cone is a triangle as well.

**Definition 2.9.** A morphism of triangles will be called good if the mapping cone of the morphism is a triangle.

**Theorem 2.13.** *A pre-triangulated category  $\mathcal{T}$  is triangulated if given two triangles  $(A, B, C, a, b, c)$  and  $(A', B', C', a', b', c')$  and diagram (1) commutes, then diagram (1) can be completed to diagram (2) such that  $\phi$  is good. That is the mapping cone of  $\phi$  is a triangle.*

$$\begin{array}{ccc}
(1) \quad \begin{array}{ccc} A & \xrightarrow{a} & B \\ \downarrow \phi_A & & \downarrow \phi_B \\ A' & \xrightarrow{a'} & B' \end{array} & (2) \quad \begin{array}{ccccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C & \xrightarrow{c} & TA \\ \downarrow \phi_A & & \downarrow \phi_B & & \downarrow \phi_C & & \downarrow T\phi_A \\ A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C' & \xrightarrow{c'} & TA' \end{array}
\end{array}$$

*Remark.* This condition is equivalent to the Octahedron axiom.

**Definition 2.10.** A commutative square (1) is called homotopy cartesian if it arises from a triangle. That is, (2) is a triangle.

$$(1) \quad \begin{array}{ccc} D & \longrightarrow & A \\ \downarrow & \lrcorner HO \lrcorner & \downarrow \\ B & \longrightarrow & C \end{array} \implies (2) \quad \begin{array}{ccc} D & & \\ \uparrow T & \searrow & \\ A \oplus B & & \\ \downarrow & \swarrow & \\ C & & \end{array}$$

*Remark.* A way to construct homotopy cartesian squares is with homotopy pullbacks. That is to us TR1 on the following map to get a triangle.

$$\begin{array}{c}
A \\
\downarrow a \\
B \xrightarrow{b} C
\end{array}
\Rightarrow A \oplus B \xrightarrow{\begin{pmatrix} a & b \end{pmatrix}} C$$

$$\Rightarrow A \oplus B \xrightarrow{\begin{pmatrix} a & b \end{pmatrix}} C \longrightarrow TD \longrightarrow TA \oplus TB \Rightarrow
\begin{array}{ccc}
D & \longrightarrow & A \\
\downarrow & \lrcorner HO \lrcorner & \downarrow \\
B & \longrightarrow & C
\end{array}$$

Dually, one can use homotopy push-outs to get homotopy cartesian squares.

*Remark.* A remark about good maps and homotopy cartesian squares???

**Lemma 2.14.** *Suppose that there is a homotopy cartesian square (1). Then there are triangles and a triangle isomorphism as in (2).*

$$(1) \quad \begin{array}{ccc} D & \xrightarrow{g'} & A \\ \downarrow f' \lrcorner HO \lrcorner \downarrow f & & \\ B & \xrightarrow{g} & C \end{array} \quad (2) \quad \begin{array}{ccccccc} D & \longrightarrow & A & \longrightarrow & E & \longrightarrow & TD \\ \downarrow f' & & \downarrow f & & \parallel & & \downarrow Tf' \\ B & \longrightarrow & C & \longrightarrow & E & \longrightarrow & TB \end{array}$$

*Proof.* There is a commutative square which satisfies the requirements of the octahedron axiom.

$$\begin{array}{ccccc}
D \xrightarrow{\begin{pmatrix} g' \\ f' \end{pmatrix}} A \oplus B & & D \xrightarrow{\begin{pmatrix} g' \\ f' \end{pmatrix}} A \oplus B & & A \oplus B \xrightarrow{\begin{pmatrix} 0 & 1 \end{pmatrix}} A \\
\parallel & & \uparrow T \quad \downarrow 0 & & \uparrow T \quad \downarrow 0 \\
D \xrightarrow{g'} A & & C & & TB \xrightarrow{\begin{pmatrix} 0 \\ 1 \end{pmatrix}} A \\
& & \downarrow (f \ g) & & \downarrow 0 \\
& & & & E
\end{array}$$

This setup shows that there is a triangle with a commutative square:

$$\begin{array}{ccccccc}
C & \longrightarrow & E & \longrightarrow & TB & \longrightarrow & TC \\
& & & & & & \downarrow (f \ g) \\
& & & & & & C \longrightarrow E
\end{array}
\quad
\begin{array}{ccc}
A \oplus B & \xrightarrow{\begin{pmatrix} 1 & 0 \end{pmatrix}} & A \\
\downarrow (f \ g) & & \downarrow \\
C & \longrightarrow & E
\end{array}$$

. Since  $\begin{pmatrix} 1 & 0 \end{pmatrix}$  is a splitmono, we have that the following is a morphism of triangles.

$$\begin{array}{ccccccc}
D & \longrightarrow & A & \longrightarrow & E & \longrightarrow & TD \\
\downarrow & & \downarrow & & \parallel & & \downarrow \\
B & \longrightarrow & C & \longrightarrow & E & \longrightarrow & TB
\end{array}$$

□

## 2.3 Calculus of Fractions and the Verdier Quotient

Localization is a method for adding formal inverses to a category. It is most notably known in commutative algebra where we can invert elements with respect to some ideal of the ring. The rational numbers can be shown to be a localization of the integers at every number except 0. The category gained from localizing at some set  $S$  of morphisms is the universal category where these morphisms are isomorphisms.

**Definition 2.11.** Let  $S$  be a collection of morphisms in the category  $\mathcal{C}$ . We say that the Localization of  $\mathcal{C}$  on  $S$  is the category  $\mathcal{C}[S^{-1}]$  together with a functor  $q : \mathcal{C} \rightarrow \mathcal{C}[S^{-1}]$  such that:

- $\forall s : S | q(s)$  is an isomorphism
- For any functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  such that  $\forall s : S$  such that  $F(s)$  is an isomorphism, then  $F$  factors through  $q$ . That is to say that there is a natural isomorphism  $\eta : F \rightarrow F' \circ q$  so that  $\mathcal{C}[S^{-1}]$  is the universal category where morphisms in  $S$  are isomorphisms.

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{D} \\ & \searrow q & \nearrow F' \\ & S^{-1}\mathcal{C} & \end{array}$$

*Remark.* Even though if we know that  $\mathcal{C}$  is locally small, then we cannot be sure that the category  $\mathcal{C}[S^{-1}]$  is again locally small.

These categories are in general pretty hard to describe. When the set of morphisms is what we call a multiplicative system, we get the same calculus of fractions description of these localization the same style as for localizations of rings.

**Definition 2.12.** A set  $S$  of morphisms in a category  $\mathcal{C}$  is called right multiplicative if it satisfies the following conditions:

- $S$  is closed under composition, i.e. if  $f, g : S$  are composable then  $gf : S$ . Every identity morphism in  $\mathcal{C}$  is in  $S$ .
- (Right Ore condition) If  $t : X \rightarrow Y$  is a morphism in  $S$ , then  $\forall g : Z \rightarrow Y$  there is a commutative square (1) such that  $f : W \rightarrow X$  and  $s : W \rightarrow Z$  exists and  $s : S$  as well.

$$(1) \quad \begin{array}{ccc} W & \xrightarrow{f} & X \\ \downarrow s & & \downarrow t \\ Z & \xrightarrow{g} & Y \end{array}$$

- (Left cancellation) Suppose  $f, g : X \rightarrow Y$  are parallel morphisms in  $\mathcal{C}$ , then 1.  $\implies$  2.:
  1.  $sf = sg$  for som  $s : S$  starting at  $Y$
  2.  $ft = gt$  for som  $t : S$  ending at  $X$

*Remark.* The previous definition has a dual statement. We say that a set  $S$  of morphisms is a left multiplicative system if it satisfies:

- $S$  is closed under composition, i.e. if  $f, g : S$  are composable then  $gf : S$ . Every identity morphism in  $\mathcal{C}$  is in  $S$ .
- (Left Ore condition) If  $s : Y \rightarrow Z$  is a morphism in  $S$ , then  $\forall f : Y \rightarrow X$  there is a commutative square (1) such that  $g : Z \rightarrow W$  and  $t : X \rightarrow W$  exists and  $t : S$  as well.

$$(1) \quad \begin{array}{ccc} Y & \xrightarrow{f} & X \\ \downarrow s & & \downarrow t \\ Z & \xrightarrow{g} & W \end{array}$$

- (Right cancellation) Suppose  $f, g : X \rightarrow Y$  are parallel morphisms in  $\mathcal{C}$ , then 1.  $\implies$  2.:
  1.  $ft = gt$  for som  $t : S$  ending at  $X$
  2.  $sf = sg$  for som  $s : S$  starting at  $Y$

If  $S$  is both right multiplicative and left multiplicative then we just say that it is multiplicative.

*Prototype.* Let  $R$  be a commutative integral domain, ... (look at Bacharaya and how they define the field of fractions, or ask Andreas if he have any good literature on this topic)

As with the definition of localization of rings, localization of a category  $\mathcal{C}$  at a multiplicative system will be defined with fractions. That is the morphisms will be "fractions" of morphisms. These morphisms will be described as diagrams over spans for right multiplicative systems (or dually cospans for left multiplicative systems), together with an equivalence relation.

**Definition 2.13.** A span is a diagram of the form:

$$\cdot \longleftarrow \cdot \longrightarrow \cdot$$

**Definition 2.14.** Let  $S$  be a right multiplicative system of morphisms in a category  $\mathcal{C}$ . Given a morphism  $s : Y \rightarrow X$  in  $S$  and a morphism  $t : Y \rightarrow Z$  we define the right fraction of  $s$  and  $t$  to be the span of the morphisms. That is  $s$  and  $t$  fit in the diagram:

$$X \xleftarrow{s} Y \xrightarrow{t} Z$$

We denote the right fraction as  $ts^{-1}$ . Let  $\sim$  be the equivalence relation of right fractions given by the diagram (1) such that  $ts^{-1} \sim t's'^{-1}$  if and only if  $\exists w, w' : \mathcal{C}$  making the diagram commute and that the middle row is a right fraction.

$$\begin{array}{ccccc} & & Y & & \\ & \swarrow s & \uparrow w & \searrow t & \\ X & \xleftarrow{\quad} & W & \xrightarrow{\quad} & Z \\ & \swarrow s' & \downarrow w' & \searrow t' & \\ & & Y' & & \end{array}$$

Dually, we define left fractions as diagrams over cospans such that if  $t : S$  we get a left fraction  $t^{-1}s$  as the diagram:

$$X \xrightarrow{s} Y \xleftarrow{t} Z$$

The equivalence relation  $\sim$  is given by the diagram in the same manner as above.

$$\begin{array}{ccccc} & & Y & & \\ & \swarrow s & \downarrow w & \swarrow t & \\ X & \xrightarrow{\quad} & W & \xleftarrow{\quad} & Z \\ & \swarrow s' & \uparrow w' & \swarrow t' & \\ & & Y' & & \end{array}$$

**Proposition 2.15.** *Suppose that  $S$  is a right multiplicative system, then the relation stated above is in fact an equivalence relation.*

*Proof.* We will need to prove that  $\sim$  is reflexive, symmetric and transitive.

- (Reflexive) Let  $fs^{-1}$  be a right fraction. Then  $fs^{-1} \sim fs^{-1}$  by the diagram:

$$\begin{array}{ccccc} & & W & & \\ & \swarrow s & \parallel & \searrow f & \\ X & \xleftarrow{s} & W & \xrightarrow{f} & Y \\ & \swarrow s & \parallel & \swarrow f & \\ & & W & & \end{array}$$

- (Symmetric) Let  $fs^{-1}$  and  $gt^{-1}$  be two right fractions such that  $fs^{-1} \sim gt^{-1}$ , that is the diagram commute. Due to inherent symmetric nature of the diagram it follows that  $gt^{-1} \sim fs^{-1}$ .

$$\begin{array}{ccc}
 & W & \\
 & \uparrow w & \searrow f \\
 X & \xleftarrow{s} \widetilde{W} \xrightarrow{g} & Y \\
 & \downarrow w' & \nearrow g \\
 & W' &
 \end{array}
 \Rightarrow
 \begin{array}{ccc}
 & W' & \\
 & \uparrow w' & \searrow g \\
 X & \xleftarrow{t} \widetilde{W} \xrightarrow{f} & Y \\
 & \downarrow w & \nearrow f \\
 & W &
 \end{array}$$

- (Transitive) Suppose that there are three right fractions  $fs^{-1}$ ,  $gt^{-1}$  and  $hu^{-1}$  such that  $fs^{-1} \sim gt^{-1}$  and  $gt^{-1} \sim hu^{-1}$ . That is, diagrammatically speaking:

$$\begin{array}{ccc}
 & W' & \\
 & \uparrow w' & \searrow f \\
 X & \xleftarrow{s} \widetilde{W} \xrightarrow{g} & Y \\
 & \downarrow \widetilde{w}' & \nearrow g \\
 & W &
 \end{array}
 \quad
 \begin{array}{ccc}
 & W & \\
 & \uparrow \widetilde{w}'' & \searrow g \\
 X & \xleftarrow{t} \widetilde{W} \xrightarrow{h} & Y \\
 & \downarrow w'' & \nearrow h \\
 & W'' &
 \end{array}$$

By the Ore condition we can create two new maps from  $\widetilde{w}'$  and  $\widetilde{w}''$ . Since both morphisms are assumed to be in  $S$ , then we get that both  $\widetilde{w}'$  and  $\widetilde{w}''$  are in  $S$ . The following diagram then shows that  $fs^{-1} \sim hu^{-1}$ . A simple diagram chase shows that it is commutative.

$$\begin{array}{ccc}
 \widetilde{\widetilde{W}} & \xrightarrow{\widetilde{w}''} & \widetilde{\widetilde{W}} \\
 \downarrow \widetilde{w}' & & \downarrow \widetilde{w}'' \\
 \widetilde{W} & \xrightarrow{\widetilde{w}'} & W
 \end{array}
 \quad
 \begin{array}{ccc}
 & W' & \\
 & \uparrow \widetilde{w}' & \searrow f \\
 & \widetilde{\widetilde{W}} & \\
 & \uparrow \widetilde{w}'' & \searrow g \\
 X & \xleftarrow{s} \widetilde{W} \xrightarrow{h} & Y \\
 & \downarrow \widetilde{w}'' & \nearrow h \\
 & \widetilde{\widetilde{W}} & \\
 & \downarrow \widetilde{w}' & \nearrow f \\
 & W'' &
 \end{array}$$

□

**Definition 2.15.** Let  $S$  be a multiplicate system in a category  $\mathcal{C}$ . Given two right fractions  $fs^{-1}$  and  $gt^{-1}$  in the diagrams:

$$X \xleftarrow{s} W \xrightarrow{f} Y \quad \& \quad Y \xleftarrow{t} W' \xrightarrow{g} Z$$

we can define the composite of these fractions  $gt^{-1} \circ fs^{-1}$  by the Ore condition:

$$\begin{array}{ccccc} \widetilde{W} & \xrightarrow{h} & W' & \xrightarrow{g} & Z \\ \downarrow u & & \downarrow t & & \\ X \xleftarrow{s} W & \xrightarrow{f} & Y & & \end{array}$$

The composite is then the right fraction  $gt^{-1} \circ fs^{-1} = gh(su)^{-1}$ .

**Proposition 2.16.** *The composition of right fractions is well-defined up to equivalence.*

*Proof.* In order to prove that the composite is well-defined we need to prove that the composite is independent from the different choices of the right Ore condition and that it is independent from choice of right fraction. There will only be presented a proof for that the choice of Ore maps is independent, as the other two cases are analogous.

Suppose we have two right fractions  $fs^{-1}$  and  $gt^{-1}$  as indicated from the diagrams.

$$X \xleftarrow{s} W_1 \xrightarrow{f} Y \quad \& \quad Y \xleftarrow{t} W_2 \xrightarrow{g} Z$$

Further suppose that there are at least two different choices for the maps gained by the right Ore condition. That is for example  $(\widetilde{W}, \widetilde{s}, \widetilde{f})$  and  $(\widehat{W}, \widehat{s}, \widehat{g})$ . We can draw the two compositions as:

$$\begin{array}{ccc} \widetilde{W} & \xrightarrow{\widetilde{g}} & W_2 \xrightarrow{g} Z \\ \downarrow \widetilde{s} & & \downarrow t \\ X \xleftarrow{s} W_1 & \xrightarrow{f} & Y \end{array} \qquad \begin{array}{ccc} \widehat{W} & \xrightarrow{\widehat{f}} & W_2 \xrightarrow{g} Z \\ \downarrow \widehat{s} & & \downarrow t \\ X \xleftarrow{s} W_1 & \xrightarrow{f} & Y \end{array}$$

By combining the diagrams at  $W_1$  by using the right Ore condition again we can find  $(W, \widetilde{w}, \widehat{w})$  as in the diagram below.

$$\begin{array}{ccccccc}
& & \bar{W} & & & & \\
& & \searrow \xi & & & & \\
& & W & \xrightarrow{\hat{w}} & \widehat{W} & \xrightarrow{\hat{f}} & W_2 \xrightarrow{g} Z \\
& & \downarrow \tilde{w} & & \downarrow \hat{f} & & \downarrow t \\
& & \widetilde{W} & \xrightarrow{\tilde{s}} & W_1 & \xrightarrow{f} & Y \\
& & & & \downarrow s & & \\
& & & & X & & 
\end{array}$$

We see that the three squares commute, as by the definition of right Ore condition. Thus we have that  $s\tilde{s}\tilde{w} = s\hat{s}\hat{w}$ . As the three squares commute we get that  $t\hat{f}\hat{w} = t\tilde{g}\tilde{w}$ . As  $t : S$  we can use right cancellation to find a  $\xi : \bar{W} \rightarrow W$  such that  $\hat{f}\hat{w}\xi = \tilde{g}\tilde{w}\xi \implies g\hat{f}\hat{w}\xi = g\tilde{g}\tilde{w}\xi$ . Thus the equivalence relation diagram commutes:

$$\begin{array}{ccccc}
& & \widehat{W} & & \\
& \swarrow s\hat{s} & \uparrow \hat{w}\xi & \searrow g\hat{f} & \\
X & \longleftarrow & \bar{W} & \longrightarrow & Z \\
& \swarrow s\tilde{s} & \downarrow \tilde{w}\xi & \searrow g\tilde{g} & \\
& & \widetilde{W} & & 
\end{array}$$

□

**Proposition 2.17.** *The composition of right fractions is associative.*

*Proof.* Let  $fs^{-1}$ ,  $gt^{-1}$  and  $hu^{-1}$  be right fractions as in the diagrams below:

$$A \xleftarrow{s} X \xrightarrow{f} B \quad , \quad B \xleftarrow{t} Y \xrightarrow{g} C \quad \& \quad C \xleftarrow{u} Z \xrightarrow{h} D$$

There are two different ways of calculating the composition. Every morphism in  $S$  will be marked blue.



$$\begin{array}{ccc}
\frac{hu^{-1} \circ (gt^{-1} \circ fs^{-1})}{W \longrightarrow Z \longrightarrow D} & & \frac{(hu^{-1} \circ gt^{-1}) \circ fs^{-1}}{V' \longrightarrow W' \longrightarrow Z \longrightarrow D} \\
\downarrow & & \downarrow \\
V \longrightarrow Y \longrightarrow C & & Y \longrightarrow C \\
\downarrow & & \downarrow \\
X \longrightarrow B & & X \longrightarrow B \\
\downarrow & & \downarrow \\
A & & A
\end{array}$$

To be able to find a relation between these diagrams we create another diagram with the right Ore condition.

$$\begin{array}{ccc}
T & \dashrightarrow & V' \\
\downarrow & & \downarrow \\
W & \longrightarrow & X
\end{array}$$

To finish the proof, one would need to show that the maps to  $A$  and  $D$  commute. The maps to  $A$  commute right out of the bat, by the right Ore condition. To prove that the maps to  $D$  commute, first apply right cancellation on the maps to  $B$ , then on the maps to  $C$ .

□

**Definition 2.16.** Let  $S$  be a right multiplicative system in a category  $\mathcal{C}$ . We define a category  $\mathfrak{r}S^{-1}\mathcal{C}$  to have objects  $\mathfrak{Ob}\mathfrak{r}S^{-1}\mathcal{C} = \mathfrak{Ob}\mathcal{C}$  and morphisms  $\mathfrak{Mor}\mathfrak{r}S^{-1}\mathcal{C} = \{\text{right fractions of } S\} / \sim$ . This means that the morphisms  $\mathfrak{r}S^{-1}\mathcal{C}(X, Y)$  are spans in  $\mathcal{C}$  where one of the maps are in  $S$  up to equivalence.

$$X \longleftarrow A \longrightarrow Y$$

This is well-defined by the previous results and the identity morphisms are the right fractions of the form:

$$X \longequal{\quad} X \longequal{\quad} X$$

*Remark.* Dually there is a category  $\mathfrak{l}S^{-1}\mathcal{C}$  for a left multiplicative system  $S$  in a category  $\mathcal{C}$ . It is defined in the same manner as  $\mathfrak{r}S^{-1}\mathcal{C}$ , but with left fractions instead.

*Remark.* Given that  $S$  is right multiplicative, A right fraction from the object  $A$  to the object  $B$  can be described with a special kind of diagram. Let  $A \downarrow S$  be the comma category of arrows from  $S$  ending in  $A$  and let  $\delta : A \downarrow S \rightarrow \mathcal{C}$  be the forgetful functor, sending each arrow to its domain. A morphism in  $A \downarrow S$  from the objects  $(b, B')$  to  $(c, C')$  is a morphism  $t : B' \rightarrow C'$  such that  $b = ct$ . We see that there is a correspondance between right fractions and elements in components of diagrams over  $A \downarrow S$  such as  $\mathcal{C}(\delta b, B) = \mathcal{C}(B', B)$

and right fractions. That is, let  $f : \delta \downarrow, \mathcal{B}$ , then  $f$  can be regarded as  $fb^{-1}$ . A morphism from  $\mathcal{C}(\delta c, B)$  to  $\mathcal{C}(\delta b, B)$  is a morphism induced by a morphism from  $(b, B')$  to  $(c, C')$  in  $A \downarrow S$ . By the equivalence relation above we want to fractions  $fb^{-1}$  and  $gc^{-1}$  to be identified if there exists morphisms from  $\mathcal{C}(\delta d, B)$  with maps  $b' : (d, D') \rightarrow (b, B')$  and  $c' : (d, D') \rightarrow (c, C')$  in  $A \downarrow S$  such that  $b' * f = c' * g$ . This would be the same as saying that the right fractions are the coequalizer of the diagram  $\mathcal{C}(\delta b, B) \coprod \mathcal{C}(\delta c, B) \rightrightarrows \mathcal{C}(\delta d, B)$ . This observation motivates that the right fractions from  $A$  to  $B$  is described as the colimit of the functor  $\mathcal{C}(\delta -, B) : A \downarrow S \rightarrow \mathbf{SET}$ . Dually, if  $S$  is left multiplicative we get that the left fractions from  $A$  to  $B$  can be described as the colimit of the functor  $\mathcal{C}(A, \rho_-) : S \downarrow B \rightarrow \mathbf{SET}$ . More details can be found in [3] and [4].

To ensure us that these categories  $\mathbf{r}S^{-1}\mathcal{C}$  does indeed exist there are many different criteria which we can place upon our assumptions. A natural restriction is to ensure that the colimits above exists as sets

**Definition 2.17.** A multiplicative system  $S$  in a locally small category  $\mathcal{C}$  is called locally small on the right if for every object  $X : \mathcal{C}$  there is a set  $S_X$  of morphisms from  $S$  such that for every morphism  $f : X_1 \rightarrow X$  in  $S$  there is a morphism  $f' : X' \rightarrow X$  in  $S_X$  factoring through  $f$ .

The dual of this definition will be called a locally small multiplicative system on the left. If it is both locally small on the left and the right, we will simply call it locally small.

*Remark.* If  $S$  is a left multiplicative system, then  $S \downarrow A$  is a filtered category for every object  $A$ . Dually, if  $S$  is right multiplicative then  $A \downarrow S$  is cofiltered.

*Remark.* Equipped with this notion we are now able to prove that the localizations exists as locally small categories. Locally small right multiplicative systems allows us to prove that the classes  $\mathbf{r}S^{-1}\mathcal{C}(X, Y)$  are sets. This can be seen as we can regard  $S_X$  as a small category. By using the right Ore condition and left cancellation we can extend  $S_X$  such that it is again cofiltered and admits the same colimit, i.e.

$$\varinjlim \mathcal{C}(\delta -, B) : A \downarrow S \rightarrow \mathbf{Set} \simeq \varinjlim \mathcal{C}(\delta -, B) : S_A \rightarrow \mathbf{Set}.$$

**Theorem 2.18.** (Gabriel-Zisman) Let  $S$  be a locally small right multiplicative system of morphisms in a category  $\mathcal{C}$ . Then the category  $\mathbf{r}S^{-1}\mathcal{C}$  exists and is the localization of  $\mathcal{C}$  on  $S$ . That is there is an equivalence of categories  $\mathcal{C}[S^{-1}] \simeq \mathbf{r}S^{-1}\mathcal{C}$  together with a functor  $q : \mathcal{C} \rightarrow \mathbf{r}S^{-1}\mathcal{C}$  sending a morphism  $f : X \rightarrow Y$  to the right fraction  $fid_X^{-1}$ .

*Proof.* To prove the theorem we have to show that  $q$  is a functor, and that it is universal. Suppose that  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$  are morphisms in  $\mathcal{C}$ . Then  $q(gf) = (gf)id_X^{-1}$  and  $q(g)q(f) = (gid_Y^{-1}) \circ (fid_X^{-1})$ . We can then choose the composition to be defined by the diagram:

$$\begin{array}{ccccc}
X & \xrightarrow{f} & Y & \xrightarrow{g} & Z \\
\parallel & & \parallel & & \\
X & \xrightarrow{f} & Y & & \\
\parallel & & & & \\
X & & & & 
\end{array}$$

Thus we can see that  $(gid_Y^{-1}) \circ (fid_X^{-1}) = (gf)id_X^{-1}$  asserting the functoriality of  $q$ .

To see that  $q$  is universal let  $\mathcal{D}$  be a category where every morphism of  $S$  is an isomorphism, and suppose there is a functor  $F : \mathcal{C} \rightarrow \mathcal{D}$ . We can define a functor  $\mathfrak{r}S^{-1}F : \mathfrak{r}S^{-1}\mathcal{C} \rightarrow \mathcal{D}$  by  $\mathfrak{r}S^{-1}F(fs^{-1}) = F(f)F(s)^{-1}$ . One can see that  $F = \mathfrak{r}S^{-1}F \circ q$ , it remains to show that it is well-defined. Suppose  $fs^{-1} = gt^{-1}$ , that means there is a diagram in  $\mathcal{C}$  with the blue arrows in  $S$ :

$$\begin{array}{ccccc}
& & W' & & \\
& \swarrow s & \uparrow w' & \searrow f & \\
X & & W & & Y \\
& \swarrow t & \downarrow w'' & \searrow g & \\
& & W'' & & 
\end{array}$$

Thus in  $\mathcal{D}$  we have that  $F(t) = F(sw')F(w'')^{-1}$  and  $F(g) = F(fw')F(w'')^{-1}$ , and this again gives us that

$$\begin{aligned}
\mathfrak{r}S^{-1}F(gt^{-1}) &= F(g)F(t)^{-1} \\
&= F(fw')F(w'')^{-1}(F(fw')F(w'')^{-1})^{-1} = F(fw')F(w'')^{-1}F(w'')F(sw')^{-1} \\
&= F(f)F(w')F(w')^{-1}F(s)^{-1} = F(f)F(s)^{-1} = \mathfrak{r}S^{-1}F(fs^{-1})
\end{aligned}$$

Thus  $\mathfrak{r}S^{-1}F$  is well-defined and is unique by construction.  $\square$

**Corollary 2.18.1.** *If  $S$  is a locally small left multiplicative system instead then  $\mathfrak{l}S^{-1}\mathcal{C}$  is the localization of  $\mathcal{C}$  on  $S$ .*

*If moreover  $S$  is a locally small multiplicative system, then there is an equivalence of categories  $\mathfrak{r}S^{-1}\mathcal{C} \simeq \mathfrak{l}S^{-1}\mathcal{C}$ .*

*Proof.* The first statement is dual to the theorem.

To see the other statement, note that both  $\mathfrak{r}S^{-1}\mathcal{C}$  and  $\mathfrak{l}S^{-1}\mathcal{C}$  are the universal categories where the morphisms of  $S$  are isomorphisms. Thus it follows that these categories have to be equivalent.  $\square$

*Remark.* Since righthandedness of lefthandedness of the multiplicative system  $S$  doesn't affect the localization we can simply call the localization of a (left/right) multiplicative system for  $S^{-1}\mathcal{C}$ .

*Remark.* A morphisms  $f : \mathcal{C}(X, Y)$  will be invertible in the localized category if it is in the same equivalence class as the identity, both  $id_X$  and  $id_Y$ . This forces a morphism  $f$  to be invertible in  $S^{-1}\mathcal{C}$  if and only if there is  $g, h : S$  such that  $fg, hf : S$ .

**Proposition 2.19.** *Let  $\mathcal{C}$  be a category, and  $S$  a right multiplicative set of morphisms. The canonical functor  $q : \mathcal{C} \rightarrow S^{-1}\mathcal{C}$  commutes with finite limits.*

*Proof.* Let  $T : \mathcal{D} \rightarrow \mathcal{C}$  be a diagram over a finite category  $\mathcal{D}$ . Then for any object  $A : S^{-1}\mathcal{C}$  we have the following equation.

$$\begin{aligned} S^{-1}\mathcal{C}(qA, q(\varprojlim T_-)) &\simeq \varinjlim \mathcal{C}(\delta_-, \varprojlim T_-) \\ &\simeq \varinjlim \varprojlim \mathcal{C}(\delta_-, T_-) \simeq \varprojlim \varinjlim \mathcal{C}(\delta_-, T_-) \simeq \varprojlim S^{-1}\mathcal{C}(qA, q(T_-)) \end{aligned}$$

The first isomorphism is by the remark, the second is by the representative nature of limits and the third isomorphism is from the fact that filtered colimits commute with finite limits. The colimits are filtered by the remark that  $S_A$  is cofiltered and that the functor  $\mathcal{C}(, A)$  is contravariant.  $\square$

*Remark.* I am not quite sure yet how this argument proves the statement I want to prove, but that can be figure out later. I also don't know the proof for why filtered colimits commute with finite colimits, but it is in Riehl's book.

**Proposition 2.20.** *Let  $\mathcal{C}$  be a category with a zero. That is an object which is both initial and terminal. Suppose that  $S$  is a right multiplicative system, then  $q0$  is a zero object in  $S^{-1}\mathcal{C}$ .*

*Proof.* The claim that  $q0$  is initial follows from that initial is a limit of a diagram over the empty category. To see that  $q0$  is terminal we only need to prove that every right fraction is equivalent with  $0id_A^{-1}$ . The following diagram proves this:

$$\begin{array}{ccccc} & & X & & \\ & \swarrow f & \downarrow f & \searrow 0 & \\ A & \xlongequal{\quad} & A & \xrightarrow{0} & 0 \end{array}$$

$\square$

**Proposition 2.21.** *If  $\mathcal{A}$  is an additive category and  $S$  is a right multiplicative system, then  $S^{-1}\mathcal{A}$  is additive as well.*



$$\begin{array}{ccc}
A & \xrightarrow{a} & B \\
\uparrow T & & \\
c \downarrow & & \\
C & \xleftarrow{b} & 
\end{array}
\quad \Longrightarrow \quad
\begin{array}{ccc}
F(A) & \xrightarrow{F(a)} & F(B) \\
\uparrow T & & \\
F(c) \downarrow & & \\
F(C) & \xleftarrow{F(b)} & 
\end{array}$$

**Definition 2.19.** A triangulated subcategory  $\mathcal{S}$  of a triangulated category  $\mathcal{T}$  is a full additive subcategory closed under isomorphisms such that the inclusion functor is triangulated.

**Definition 2.20.** Let  $F : \mathcal{S} \rightarrow \mathcal{T}$  be a triangulated functor. The kernel of  $F$  is defined to be the full subcategory  $Ker(F)$  of  $\mathcal{S}$  such that every object in  $Ker(F)$  gets mapped to 0 by  $F$ . That is,  $Ker(F)$  is the class of objects  $\{K : \mathcal{S} | F(K) \simeq 0\}$ .

**Lemma 2.22.** *The kernel of a triangulated functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is a triangulated subcategory.*

*Proof.* Let  $X : Ker F$ , since  $F$  is a triangulated functor  $CX : Ker F$  as  $F(CX) = D(FX) = D0 = 0$ . As  $F$  is triangulated, we have that every triangle maps to a triangle. Let  $X, Y : Ker F$ , then:

$$\begin{array}{ccc}
X & \xrightarrow{\quad} & Y \\
\uparrow T & & \\
Z & \xleftarrow{\quad} & 
\end{array}
\quad \Longrightarrow \quad
\begin{array}{ccc}
0 & \xrightarrow{\quad} & 0 \\
\uparrow T & & \\
F(Z) & \xleftarrow{\quad} & 
\end{array}$$

By TR3 and the 2 out of 3 property  $F(Z) \simeq 0 \implies Z : Ker F$ . Thus  $Ker F$  is a triangulated subcategory of  $\mathcal{C}$ .  $\square$

**Definition 2.21.** A subcategory  $\mathcal{S}$  of a triangulated category  $\mathcal{T}$  is called thick if it contains all the direct summands of its objects.

**Lemma 2.23.** *The kernel of a triangulated functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is thick.*

*Proof.* Let  $X \oplus Y : Ker F$ , since  $F$  is additive we have that  $0 \simeq F(X \oplus Y) \simeq F(X) \oplus F(Y)$ , but then there is a splitmono  $F(X) \rightarrow 0 \implies F(X) \simeq 0 \simeq F(Y)$ .  $\square$

**Lemma 2.24.** *Let  $F : \mathcal{C} \rightarrow \mathcal{D}$  be a triangulated functor. Suppose that  $f : X \rightarrow Y$  is a morphism such that  $F(f)$  is an isomorphism. Then the cone of  $f$  is in  $Ker F$ .*

*Proof.* There is an isomorphism of triangles in  $\mathcal{D}$ , showing that the cone of  $f$  is in  $Ker F$ .

$$\begin{array}{ccccccc}
FX & \xrightarrow{Ff} & FY & \longrightarrow & F(\text{cone}(f)) & \longrightarrow & FTX \\
\parallel & & \parallel & & \downarrow \cong & & \parallel \\
FX & \xrightarrow{Ff} & FY & \longrightarrow & 0 & \longrightarrow & FTX
\end{array}$$

□

The goal for the rest of this section is to prove that there is a localization at any triangulated subcategory  $\mathcal{S} \subseteq \mathcal{C}$ . This localization will yield a functor  $q : \mathcal{C} \rightarrow \mathcal{C}/\mathcal{S}$  such that  $\mathcal{S} \subseteq \text{Ker} q$ . We will define a set of morphism  $\text{Mor}_{\mathcal{S}}$  related to  $\mathcal{S}$  such that this set is multiplicative.

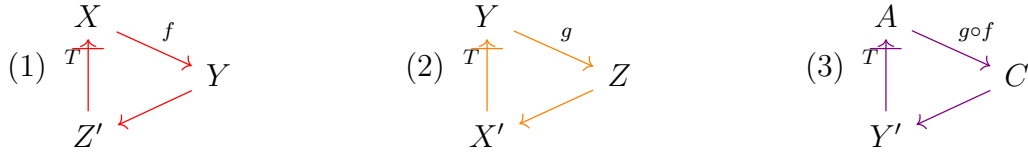
**Definition 2.22.** Let  $\mathcal{C}$  be a triangulated category and  $\mathcal{S} \subseteq \mathcal{C}$  be a triangulated subcategory. Define a collection  $\text{Mor}_{\mathcal{S}}$  to be a collection of morphisms in  $\mathcal{C}$  such that for any  $f : \text{Mor}_{\mathcal{S}}$  there is a triangle with  $C : \mathcal{S}$ .

$$A \xrightarrow{f} B \longrightarrow C \longrightarrow TA$$

*Remark.* Every isomorphism is in  $\text{Mor}_{\mathcal{S}}$ . That is because isomorphisms are found in triangles  $(A, B, 0, f, 0, 0)$  and  $0 : \mathcal{S}$  for any triangulated subcategory.

**Lemma 2.25.** Let  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$  be two morphisms. If any two of the morphisms  $f$ ,  $g$  and  $gf$  are in  $\text{Mor}_{\mathcal{S}}$  then so is the third.

*Proof.* We are able to find three triangles in  $\mathcal{C}$ .



By the Octahedron axiom there exist another triangle in  $\mathcal{C}$ :

$$Z' \longrightarrow X' \longrightarrow Y' \longrightarrow TZ'$$

Note that  $f$  is in  $\text{Mor}_{\mathcal{S}}$  if and only if  $Z' : \mathcal{S}$ . WLOG assume that  $f$  and  $g$  is in  $\text{Mor}_{\mathcal{S}}$ , this can be done due to the rotation axiom. Thus we can find the triangle  $(Z', X', Y'')$  in  $\mathcal{S}$  proving that  $Y'$  is in  $\mathcal{S}$  by the following diagram:

$$\begin{array}{ccccccc} Z' & \longrightarrow & X' & \longrightarrow & Y' & \longrightarrow & TZ' \\ \parallel & & \parallel & & \downarrow \cong & & \parallel \\ Z' & \longrightarrow & X' & \longrightarrow & Y'' & \longrightarrow & TZ' \end{array}$$

□

**Proposition 2.26.** Let  $\mathcal{S} \subseteq \mathcal{C}$  be a triangulated subcategory, then  $\text{Mor}_{\mathcal{S}}$  satisfies the Ore condition.

*Proof.* To prove that a system satisfies the Ore condition there has to be a proof for both right and left condition. Luckily, the arguments presented here can be dualized to give a proof for the other condition. Thus there will only be presented a proof for the right Ore condition. Let  $f : A \rightarrow C$  be in  $Mor_{\mathcal{S}}$  and  $g : B \rightarrow C$  in  $\mathcal{C}$ . Then we can form a homotopy pullback forming a homotopy cartesian square as follows:

$$\begin{array}{ccc} & A & \\ & \downarrow f & \\ B & \xrightarrow{g} & C \end{array} \implies \begin{array}{ccccc} D & \xrightarrow{g'} & A & & \\ \downarrow f' & \lrcorner & HO & \lrcorner & \downarrow f \\ B & \xrightarrow{g} & C & & \end{array}$$

By Lemma 2.12 there are triangles along this homotopy square identifying the cones. Since the cone of  $f$  is assumed to be in  $\mathcal{S}$ , the cone of  $f'$  is also in  $\mathcal{S}$ . This proves that  $f' : Mor_{\mathcal{S}}$ .  $\square$

**Proposition 2.27.** *For any parallell morphism  $f, g : X \rightarrow Y$  in  $\mathcal{C}$  the following are equivalent:*

1.  $sf = sg$  for some  $s : Mor_{\mathcal{S}}$  starting at  $Y$ .
2.  $ft = gt$  for some  $t : Mor_{\mathcal{S}}$  ending at  $X$ .
3.  $f - g$  factors through an object  $C : \mathcal{S}$ .

*Proof.* (1.  $\iff$  3.): Suppose that there exists an  $s : Y \rightarrow Z$  such that  $s(f - g) = 0$ . By TR1 there is a triangle  $Y \xrightarrow{s} Z \xrightarrow{Ts'} TC \longrightarrow TY$  and a long exact sequence of hoom(ology).

$$\mathcal{C}(X, C) \xrightarrow{s'_*} \mathcal{C}(X, Y) \xrightarrow{s_*} \mathcal{T}(X, Z)$$

$$p \xrightarrow{s'_*} f - g \xrightarrow{s_*} 0$$

Since  $s(f - g) = 0$  there exists a  $p : \mathcal{C}(X, C)$  such that  $f - g = s'_*p$ . By definition,  $s : Mor_{\mathcal{S}} \iff C : \mathcal{S}$ , but  $s : Mor_{\mathcal{S}} \implies f - g$  factors through  $C$ , and vice versa. (2.  $\iff$  3.): This argument is dual.  $\square$

This has shown that  $Mor_{\mathcal{S}}$  is a multiplicative system, and Theorem 2.17 say that the localization exists given that  $Mor_{\mathcal{S}}$  is locally small. The category  $Mor_{\mathcal{S}}^{-1}$  will be called for  $\mathcal{C}/\mathcal{S}$  and it is called the Verdier quotient. As  $\mathcal{C}$  is additive, it is known that  $\mathcal{C}/\mathcal{S}$  is additive as well by Proposition 2.20. The remaining part is to show that  $\mathcal{C}/\mathcal{S}$  is triangulated and that localization functor  $q : \mathcal{C} \rightarrow \mathcal{C}/\mathcal{S}$  is a triangulated functor.



**Theorem 2.28.** *Let  $\mathcal{S} \subseteq \mathcal{C}$  be triangulated categories. Then the Verdier quotient  $\mathcal{C}/\mathcal{S}$  and the functor  $q : \mathcal{C} \rightarrow \mathcal{C}/\mathcal{S}$  is the universal triangulated category where morphisms in  $\text{Mor}_{\mathcal{S}}$  are isomorphisms.*

*Proof.* The triangulation on  $\mathcal{C}/\mathcal{S}$  is defined as the following. Let  $C/S : \mathcal{C}/\mathcal{S} \rightarrow \mathcal{C}/\mathcal{S}$  be the additive autoequivalence defined by  $C/S(A) = C(A)$ . Since  $q : \mathcal{C} \rightarrow \mathcal{C}/\mathcal{S}$  maps every object to itself it follows that  $q(C(A)) \simeq C/S(A) = C/S(q(A))$ , and define  $\Delta_{\mathcal{C}/\mathcal{S}} \supseteq q(\Delta_{\mathcal{C}})$  such that  $\Delta_{\mathcal{C}/\mathcal{S}}$  has every isomorphism class of  $q(\Delta_{\mathcal{C}})$ .

$$\begin{array}{ccc} qX & & X \\ \uparrow T & \searrow & \uparrow T \\ & qY & Y \\ \downarrow & \swarrow & \downarrow \\ qZ & & Z \end{array} \iff$$

Then by definition  $q$  is triangulated if the category  $\mathcal{C}/\mathcal{S}$  is triangulated. By definition, the triangles are closed under isomorphisms,  $(X, X, 0, id_X, 0, 0)$  is a triangle, and TR2 holds. Thus it remains to show TR1 and TR4 (TR3 is implied by the other axioms). To prove TR1, let  $f s^{-1} : \mathcal{C}/\mathcal{S}(qW, qY)$ . Then expand  $f : \mathcal{C}(X, Y)$  to a triangle in  $\mathcal{C}$  with TR1, it will induce a triangle in  $\mathcal{C}/\mathcal{S}$ .

$$qX \xrightarrow{fid_X^{-1}} qY \xrightarrow{gid_Y^{-1}} qZ \xrightarrow{hid_Z^{-1}} qTX$$

Then there is an isomorphism to the following candidate triangle from the induced triangle, proving TR1.

$$\begin{array}{ccccccc} qX & \xrightarrow{fid_X^{-1}} & qY & \xrightarrow{gid_Y^{-1}} & qZ & \xrightarrow{hid_Z^{-1}} & qTX \\ \wr \downarrow \text{\scriptsize } sid_X^{-1} & & \parallel & & \parallel & & \wr \downarrow \text{\scriptsize } (Ts)id_{Tx}^{-1} \\ qW & \xrightarrow{fs^{-1}} & qY & \xrightarrow{gid_Y^{-1}} & qZ & \xrightarrow{(Ts)hid_Z^{-1}} & qTW \end{array}$$

To show the Octahedron axiom, suppose that there are three triangles in  $\mathcal{C}/\mathcal{S}$ . By construction, these triangles can be chosen such that only the first map is a fraction up to isomorphism of triangles.

$$(1) \quad \begin{array}{ccc} Z & \xrightarrow{t'} & X \\ \swarrow s & & \downarrow a \\ A & \xrightarrow{as^{-1}} & B \\ \uparrow x' & \searrow x & \\ C' & & \end{array} \quad (2) \quad \begin{array}{ccc} Y & & \\ \swarrow t & & \downarrow b \\ B & \xrightarrow{bt^{-1}} & C \\ \uparrow y' & \searrow y & \\ A' & & \end{array} \quad (3) \quad \begin{array}{ccc} Z & & \\ \swarrow st' & & \downarrow ba' \\ A & \xrightarrow{boa} & C \\ \uparrow z' & \searrow z & \\ B' & & \end{array}$$

By composing the fractions from  $A$  to  $B$  and  $B$  to  $C$  we can find an object  $Z$  as in the diagram with the Ore condition. To illustrate with triangle (1), there is a correspondance of triangles in  $\mathcal{C}/\mathcal{S}$  and  $\mathcal{C}$ . That is by the following isomorphism:

$$\begin{array}{ccccccc}
Z & \xrightarrow{at'} & B & \longrightarrow & Z' & \longrightarrow & TZ \\
\downarrow \wr & & \parallel & & \downarrow \wr & & \parallel \\
X & \xrightarrow{a} & B & \longrightarrow & C' & \longrightarrow & TX \\
\downarrow s & & & & & & \downarrow Ts \\
A & & & & & & TA
\end{array}$$

The result of the octahedron axiom follows as we can instead consider the triangles found by the composition of morphisms as below.

$$\begin{array}{ccc}
Z & & \\
\downarrow f' & \searrow bf' & \\
Y & \xrightarrow{b} & C
\end{array}$$

□

**Proposition 2.29.** *Let  $\mathcal{S} \subseteq \mathcal{C}$  be triangulated categories. If  $0 : X \rightarrow 0$  is an isomorphism in  $\mathcal{C}/\mathcal{S}$ , then there is an object  $Y$  such that  $X \oplus Y : \mathcal{S}$ .*

*Proof.* If  $0 : X \rightarrow 0$  is invertible, then there exist a map  $0 : 0 \rightarrow Y$ , such that  $0 : X \rightarrow Y$  is in  $Mor_{\mathcal{S}}$ . By definition  $X \oplus Y$  is in  $\mathcal{S}$ . □

This proposition shows us that the kernel of  $q : \mathcal{C} \rightarrow \mathcal{C}/\mathcal{S}$  is the smallest thick subcategory of  $\mathcal{C}$  such that  $\mathcal{C}/Kerq$  is the universal category where every morphism in  $Mor_{\mathcal{S}}$  is an isomorphism. For this reason  $\widehat{\mathcal{S}} = Kerq$  is called the thick closure of  $\mathcal{S}$ .

## 2.4 Universal Homological Embedding

Do Yoneda embedding into functor categories.

### 3 Exact Categories

I can maybe write some of the history of the development of the idea of exact categories.

#### 3.1 Definitions and First Properties

In this section we will focus on defining what an exact category is and the first elementary properties. We will prove the axiom dubbed as "the obscure axiom" and motivate that it is not as obscure as its name suggest. Some "short" variants of some homological diagram lemmas will also be proved.

To start with the exact categories we will first take a look towards the abelian ones first. Short exact sequences are of great interest, and they can be characterized with two morphisms  $p : A \rightarrow B$  and  $q : B \rightarrow C$  such that  $p$  is the kernel of  $q$  and  $q$  is the cokernel of  $p$ . This leads to the first definition.

**Definition 3.1.** A kernel-cokernel pair is a pair of maps  $(p, q)$  such that  $p$  is the kernel of  $q$  and  $q$  is the cokernel of  $p$ . A morphism of kernel-cokernel pairs  $(p, q)$  and  $(p', q')$  is a triple  $(f, g, h)$  such that the following diagram commutes. An isomorphisms is a triple in which each morphism is an isomorphism.

$$\begin{array}{ccccc} A & \xrightarrow{p} & B & \xrightarrow{q} & C \\ \downarrow f & & \downarrow g & & \downarrow h \\ A' & \xrightarrow{p'} & B' & \xrightarrow{q'} & C' \end{array}$$

**Lemma 3.1.** Let  $(p, q)$  be a kernel-cokernel pair, then the image and coimage of  $p$  exists and are isomorphic. I.e. this diagram exists, such that the left square is a push-out and the right square is a pull-back:

$$\begin{array}{ccccccc} 0 & \xrightarrow{0} & A & \xrightarrow{p} & B & \xrightarrow{q} & C \\ & \searrow 0 & \downarrow & & \uparrow & \nearrow 0 & \\ & & \text{Coim}(p) & \xrightarrow{\text{iso}} & \text{Im}(p) & & \end{array}$$

*Proof.* Since  $(p, q)$  is a kernel-cokernel pair we have that the first square is bicartesian and the second square is a push-out.

$$\begin{array}{ccc} A & \xrightarrow{p} & B \\ & \searrow 0 & \downarrow q \\ & & C \end{array} \quad \begin{array}{ccc} 0 & \xrightarrow{0} & A \\ & \searrow 0 & \parallel \\ & & A \end{array}$$

Thus  $Im(p) = Coim(p) = A$ , asserting the isomorphism as the identity in the diagram.

$$\begin{array}{ccccccc}
0 & \xrightarrow{0} & A & \xrightarrow{p} & B & \xrightarrow{q} & C \\
& \searrow 0 & \parallel & & \uparrow p & \nearrow 0 & \\
& & A & = & A & & 
\end{array}$$

□

**Corollary 3.1.1.** *Suppose that  $(p, q)$  is a kernel-cokernel pair. If  $p$  is an epimorphism, then  $p$  is an isomorphism.*

**Definition 3.2.** An exact structure for an additive category  $\mathcal{A}$  is a class  $\mathcal{E}$  of kernel-cokernel pairs which are closed under isomorphisms. A pair  $(p, q) : \mathcal{E}$  is called a conflation, here  $p$  is called an inflation and  $q$  is called a deflation.  $(\mathcal{A}, \mathcal{E})$  is called exact when the following axioms holds:

- (QE0)  $\forall A : \mathcal{A} \text{ } id_A$  is both an inflation and a deflation.
- (QE1) Both inflations and deflations are closed under composition.
- (QE2) The push-out of an inflation is an inflation.
- (QE2<sup>op</sup>) The pull-back of a deflation is a deflation.

An exact category is the additive category  $\mathcal{A}$  together with an exact structure  $\mathcal{E}$ .

*Remark.* When writing diagrams we use decorated arrows to indicate that a morphism is either an inflation or a deflation. A tail with a circle means inflation:  $A \rightarrowtail \circ \rightarrow B$ . Double heads with a circle means deflation:  $A \rightarrowtail \circ \twoheadrightarrow B$ . We can now rewrite the (QE2\*) axioms as:

$$\begin{array}{ccc}
A \rightarrowtail \circ \rightarrow B & A \rightarrowtail \circ \twoheadrightarrow B & \\
\downarrow & \lrcorner & \downarrow \\
C \rightarrowtail \circ \rightarrow D & C \rightarrowtail \circ \twoheadrightarrow D & 
\end{array}$$

*Remark.* Inflations are also called admissible monomorphisms, deflations are also called admissible epimorphisms and conflations are also called short exact sequences.

*Remark.* The axioms for an exact category is made in such a way that  $\mathcal{E}$  is an exact structure for  $\mathcal{A}$  if and only if  $\mathcal{E}^{op}$  is an exact structure for  $\mathcal{A}^{op}$ .

**Lemma 3.2.** *The map  $0 : 0 \rightarrow A$  is an inflation. Dually, the map  $0 : A \rightarrow 0$  is a deflation.*

*Proof.* Consider the diagram  $0 \rightarrowtail \xrightarrow{0} A \xrightarrow{id_A} A$ . The left morphism is the kernel of the right morphism making a kernel-cokernel pair  $(0, id_A)$ . The identity  $id_A$  is assumed to be a deflation, implying that the pair is a conflation. □

*Remark.* It can be seen that isomorphisms are deflations. Let  $f : A \rightarrow B$  be an isomorphism, then there are two kernel-cokernel pairs:  $(0, id_A)$  and  $(0, f)$ . Between these there is an isomorphism which is the triple  $(0, id_A, f^{-1})$ . As the conflationations are closed under isomorphism,  $(0, f)$  is a conflation, making  $f$  into a deflation. By dualizing this argument  $f$  is also an inflation.

$$\begin{array}{ccccc} 0 & \xrightarrow{0} & A & \xrightarrow{f} & B \\ \downarrow 0 & & \downarrow id_A & & \downarrow f^{-1} \\ 0 & \xrightarrow{0} & A & \xrightarrow{id_A} & A \end{array}$$

**Corollary 3.2.1.** *A kernel-cokernel pair  $(i, p)$  found as a split short-exact sequence (1) is a conflation.*

$$(1) \quad A \xrightarrow{i} A \oplus B \xrightarrow{p} B$$

*Proof.* In a category with an initial object the coproduct can be thought of as the push-out with the initial in the upper right corner. This can be assembled into push-out (1). By the lemma the zero morphisms are inflations, asserting that  $i$  and  $i'$  are inflations by (QE2). Thus there are conflationations  $(i, p)$  and  $(i', p')$ .

$$(1) \quad \begin{array}{ccc} 0 & \xrightarrow{0} & A \\ \downarrow 0 & & \downarrow i \\ B & \xrightarrow{i'} & A \oplus B \end{array}$$

□

**Corollary 3.2.2.** *The direct sum of conflationations is a conflation. I.e. there is a diagram:*

$$\begin{array}{ccccc} A \xrightarrow{i} B & \xrightarrow{p} & C & \mathcal{E} & A' \xrightarrow{i'} B' \xrightarrow{p'} C' \\ & & \downarrow & & \\ A \oplus A' & \xrightarrow{i \oplus i'} & B \oplus B' & \xrightarrow{p \oplus p'} & C \oplus C' \end{array}$$

*Proof.* We start with only considering the conflation  $(i, p)$ .  $\forall D$  there is a conflation  $(i \oplus id_D, p \oplus 0)$ , drawn as the diagram.

$$A \oplus D \xrightarrow{\begin{pmatrix} i & 0 \\ 0 & 1 \end{pmatrix}} B \oplus D \xrightarrow{\begin{pmatrix} p & 0 \end{pmatrix}} C$$

As kernels and cokernels are preserved by direct sums, this pair is in fact a kernel-cokernel pair. The epimorphism is a deflation as it can be factored by the deflations:

$$B \oplus D \xrightarrow{\begin{pmatrix} 1 & 0 \\ \circ & \circ \end{pmatrix}} B \xrightarrow{p} C$$

Thus it is seen that  $(i \oplus id_D, p \oplus 0)$  is a conflation, and dually  $(i \oplus 0, p \oplus id_D)$  is also a conflation. To finish off the proof it is seen that the morphism  $i \oplus i'$  factors as  $i \oplus id_{A'} \circ id_A \oplus i'$  asserting that it is an inflation by (QE1). By the dual argument we then get that the direct sum of conflations is a conflation.  $\square$

**Definition 3.3.** A square is bicartesian if it is both a pull-back and a push-out.

$$\begin{array}{ccc} A & \longrightarrow & B \\ \downarrow & \lrcorner & \downarrow \\ C & \longrightarrow & D \end{array}$$

**Proposition 3.3.** *The following statements are equivalent:*

1. *The square (1) is a push-out.*
2. *The sequence (2) is a conflation.*
3. *The square (1) is bicartesian.*
4. *The square (1) is a part of the commutative diagram (3)*

$$(1) \quad \begin{array}{ccc} A & \xrightarrow{i} & B \\ \downarrow f & & \downarrow g \\ C & \xrightarrow{j} & D \end{array} \quad (2) \quad A \xrightarrow{\begin{pmatrix} i \\ -f \end{pmatrix}} B \oplus C \xrightarrow{\begin{pmatrix} g & j \end{pmatrix}} D \quad (3) \quad \begin{array}{ccccc} A & \xrightarrow{i} & B & \xrightarrow{p} & E \\ \downarrow f & & \downarrow g & & \parallel \\ C & \xrightarrow{j} & D & \xrightarrow{q} & E \end{array}$$

Before the proof for this proposition will be presented a useful lemma will be proved first.

**Lemma 3.4.** *Assume that there is a commutative square (1) and an associated sequence (2). (1) is a push-out square if and only if  $(p \ q)$  is the cokernel of the morphism  $\begin{pmatrix} i \\ -j \end{pmatrix}$*

$$(1) \quad \begin{array}{ccc} A & \xrightarrow{i} & B \\ \downarrow j & & \downarrow p \\ C & \xrightarrow{q} & D \end{array} \quad (2) \quad A \xrightarrow{\begin{pmatrix} i \\ -j \end{pmatrix}} B \oplus C \xrightarrow{\begin{pmatrix} p & q \end{pmatrix}} D$$

*Proof.* For any test object  $T$  and two maps  $t_1 : B \rightarrow T$  and  $t_2 : C \rightarrow T$ , we can construct the diagrams for the universal properties of both the cokernel and the push-out. It is seen that these diagrams are equivalent, proving the lemma.

$$\begin{array}{ccc}
A \xrightarrow{\begin{pmatrix} i \\ -f \end{pmatrix}} B \oplus C & \xrightarrow{(t_1 \ t_2)} & A \xrightarrow{i} B \\
\searrow 0 & \downarrow (g \ j) & \downarrow f \quad \downarrow g \\
D & \xrightarrow{t'} T & C \xrightarrow{j} D \\
& & \searrow t_2 \quad \searrow t' \\
& & T
\end{array} \Leftrightarrow$$

□

**Corollary 3.4.1.** *For the same diagrams (1) and (2) as above the dual statement is also true. (1) is a pull-back square if and only if  $\begin{pmatrix} i \\ -j \end{pmatrix}$  is kernel of the morphism  $(p \ q)$ . Thus we have that (1) is bicartesian (i.e. both a pull-back and a push-out) if and only if the morphisms make a kernel-cokernel pair.*

*Proof. of Proposition 3.3* 1.  $\Rightarrow$  2.: By the previous lemma we know that  $(g \ j)$  is the cokernel of  $\begin{pmatrix} i \\ -j \end{pmatrix}$ . Thus proving that  $\begin{pmatrix} i \\ -j \end{pmatrix}$  is an inflation, will prove that the pair is a conflation.

Observe that the morphism  $\begin{pmatrix} i \\ -f \end{pmatrix}$  can be factored through the sequence.

$$A \xrightarrow{\begin{pmatrix} 1 \\ 0 \end{pmatrix}} A \oplus C \xrightarrow{\begin{pmatrix} 1 & 0 \\ -f & 1 \end{pmatrix}} A \oplus C \xrightarrow{\begin{pmatrix} i & 0 \\ 0 & 1 \end{pmatrix}} B \oplus C$$

By corollary 3.2.1 the first map is an inflation, as the second map is an isomorphism it is also an inflation and the last map is the direct sum of two inflations. Thus the composite of all these maps surely is an inflation by (QE1), proving the first implication.

2.  $\Rightarrow$  3.: This follows from corollary 3.4.1.

3  $\Rightarrow$  1.: This is by definition.

1.  $\Rightarrow$  4.: Let  $p$  be the cokernel of  $i$ , then we can form the diagram below.

$$\begin{array}{ccc}
A \xrightarrow{i} B & & \\
\downarrow f & \lrcorner & \downarrow g \\
C \xrightarrow{j} D & & \\
\searrow 0 & \searrow p' & \searrow p \\
& & T
\end{array}$$

$p'$  is an epimorphism as  $p = p'g$  is epi. To prove that  $p'$  is the cokernel of  $j$  we let  $T'$  be another test object with a map  $t' : D \rightarrow T'$  such that  $0 = t'j$ . By doing some diagram chases we have that  $0 = t'jg = t'gi$ , thus by the universal property of  $p$  the morphism  $t'g$  factors through  $T$  such that  $t'g = tp$  for some unique  $t$ . By rearranging we have that  $t'g = tp'g = tp$ , and  $t'j = tp'j = 0$ , thus since  $t'$  is the unique morphism satisfying this equation we demand that  $t' = tp'$ .  $t$  is also unique, for if there exist another map  $h$  such that  $tp' = hp'$ , then  $t = h$  as  $p'$  is epic. The unique existence proves the universal property, and  $p'$  is the cokernel of  $j$ .

$$\begin{array}{ccccc}
A & \xrightarrow{i} & B & \xrightarrow{p} & T \\
\downarrow f & & \downarrow g & & \parallel \\
C & \xrightarrow{j} & D & \xrightarrow{p'} & T \\
& \searrow 0 & \searrow t' & \downarrow t & \\
& & & & T'
\end{array}$$

4.  $\Rightarrow$  2.: We start by taking the pullback of  $p$  and  $p'$  using  $(QE2^{op})$ , and determines the diagram with the dual statement of the last implication.

$$\begin{array}{ccc}
\begin{array}{ccccc}
A & \xlongequal{\quad} & A \\
\downarrow \phi_{i'} & & \downarrow \phi_i \\
C & \xrightarrow{j'} & E & \xrightarrow{q'} & B \\
\parallel & & \downarrow \phi_q & \swarrow g & \downarrow \phi_p \\
C & \xrightarrow{j} & D & \xrightarrow{p'} & T
\end{array} & \Rightarrow & \begin{array}{ccccc}
B & & & & \\
\downarrow k & & & & \\
E & \xrightarrow{q'} & B \\
\downarrow \phi_q & & \downarrow \phi_p \\
D & \xrightarrow{p'} & T
\end{array}
\end{array}$$

From these diagrams we can deduce that  $q'$  is a split-epimorphism. The composite  $q'(id_E - kq') = q' - q'kq' = q' - q' = 0$  as  $q'$  is split-epi, so  $(id_E - kq')$  factors over  $j'$  as in the following diagram.

$$\begin{array}{ccc}
E & \xrightarrow{id_E - kq'} & E \\
\downarrow l & \searrow & \downarrow \phi_{q'} \\
C & \xrightarrow{j'} & E \\
& \searrow 0 & \downarrow \phi_{q'} \\
& & B
\end{array}$$

From these diagrams we can extract three different equations:

- $0 = k - k = k - kq'k = (id_E - kq')k = j'lk \Rightarrow lk = 0$  as  $j'$  is monic
- $j'lj' = (id_E - kq')j' = j' \Rightarrow lj' = id_C$  as  $j'$  is monic



- $jli' = (qj')li' = q(id_E - kq')i' = -(qk)(q'i') = -gi = -jf \implies li' = -f$  as  $j$  is monic

The morphisms  $(k \ j')$  and  $\begin{pmatrix} q' \\ l \end{pmatrix}$  are inverses:

- $(k \ j') \begin{pmatrix} q' \\ l \end{pmatrix} = kq' + j'l = kq' + id_E - kq' = id_E$
- $\begin{pmatrix} q' \\ l \end{pmatrix} (k \ j') = \begin{pmatrix} q'k & q'j' \\ lk & lj' \end{pmatrix} = \begin{pmatrix} id_B & 0 \\ 0 & id_C \end{pmatrix}$

Thus we have an isomorphism of kernel-cokernel pairs  $(id_A, \begin{pmatrix} q' \\ l \end{pmatrix} (k \ j'))$ ,

from  $(\begin{pmatrix} i \\ -f \end{pmatrix}, (f' \ i'))$  to  $(i', q)$ . This proves 2.  $\square$

**Corollary 3.4.2.** *The pull-back of an inflation along a deflation is an inflation.*

$$\begin{array}{ccc} A & \xrightarrow{i'} & B \\ \downarrow \phi_{e'} & \lrcorner & \downarrow \phi_e \\ C & \xrightarrow{i} & D \end{array}$$

*Proof.* By (QE2) this pullback exists, as there is a deflation in the pullback. Extend the diagram by adding the deflation of the inflation in the following manner.

$$\begin{array}{ccccc} & & T & & \\ & \swarrow t' & \downarrow t & \searrow 0 & \\ A & \xrightarrow{i'} & B & \xrightarrow{pe} & C \\ \downarrow \phi_{e'} & \lrcorner & \downarrow \phi_e & & \parallel \\ C & \xrightarrow{i} & D & \xrightarrow{p} & C \end{array}$$

$pe$  is a deflation by (QE1), and  $i'$  is a mono as a limit of a mono is a mono. Our goal is to prove that  $i'$  is the kernel of  $pe$ . Let  $T$  be a test object such that  $pet = 0$ . Then we have that  $te$  factorizes over  $i$ , such that we can apply the universal property of the pullback to factorize  $te$  over  $i'$ . Uniqueness of  $t'$  is achieved with  $i'$  being monic. This proves that  $(i', pe)$  is a conflation.  $\square$

**Theorem 3.5. The obscure axiom.** *Assume that  $i : A \rightarrow B$  is a morphism with a cokernel. If there is a morphism  $j : B \rightarrow C$  such that  $ji$  is an inflation, then  $i$  is an inflation.*

*Proof.*

□

*Proof.* Let  $p : B \rightarrow D$  be the cokernel of  $i$ . We start the proof with forming the push-out of  $i$  and  $ji$ .

$$\begin{array}{ccc} A & \xrightarrow{j \circ i} & C \\ \downarrow i & \lrcorner & \downarrow \\ B & \xrightarrow{\quad} & E \end{array}$$

By proposition 3.3 we get that  $\begin{pmatrix} i \\ ji \end{pmatrix}$  is an inflation.  $\begin{pmatrix} i \\ 0 \end{pmatrix} = \begin{pmatrix} id_B & 0 \\ -j & id_C \end{pmatrix} \begin{pmatrix} i \\ ji \end{pmatrix}$ , this is an inflation by (QE1) as the 2x2 matrix is an isomorphism. Observer that the cokernel of this map is  $\begin{pmatrix} k & 0 \\ 0 & id_C \end{pmatrix}$ . Our final trick will be to show that there is a pullback square, and then use (QE2) to say that  $k$  is a deflation.

$$\begin{array}{ccccc} T & & & & \\ & \searrow t & & \searrow t_1 & \\ & B & \xrightarrow{k} & D & \\ & \downarrow \begin{pmatrix} 1 \\ 0 \end{pmatrix} & & \downarrow \begin{pmatrix} 1 \\ 0 \end{pmatrix} & \\ \begin{pmatrix} t_2 \\ 0 \end{pmatrix} \searrow & B \oplus C & \xrightarrow{\begin{pmatrix} k & 0 \\ 0 & id_C \end{pmatrix}} & D \oplus C & \end{array}$$

Note that setting  $t = t_2$  we get the universal property. This is well defined as  $kt_2 = t_1$  by assumption, thus  $kt = t_1$ . This is what we need to prove that the square is a pullback, proving the obscure axiom. □

*Remark.* Write a bit about the dual of the obscure axiom.

**Lemma 3.6.** *Let  $(p, q)$  and  $(p', q')$  be the conflations:*

- $(p, q): A \xrightarrow{p} B \xrightarrow{q} \twoheadrightarrow C$
- $(p', q'): A' \xrightarrow{p'} B' \xrightarrow{q'} \twoheadrightarrow C'$

*A morphism of the conflations  $(f, g, h) : (p, q) \rightarrow (p', q')$  factors through the conflation  $A \xrightarrow{\quad} D \twoheadrightarrow C'$  such that we have the following diagram where  $g = g_2 g_1$ .*

$$\begin{array}{ccccc}
A & \xrightarrow{p} & B & \xrightarrow{q} & C \\
\downarrow f & \lrcorner & \downarrow g_1 & & \parallel \\
A' & \xrightarrow{\quad} & D & \xrightarrow{\quad} & C \\
\parallel & & \downarrow g_2 & \lrcorner & \downarrow h \\
A' & \xrightarrow{p'} & B' & \xrightarrow{q'} & C'
\end{array}$$

*Proof.* Observe that the upper part of the diagram is made by taking a push-out of  $p$  and  $f$ , where the right part is gained from proposition 3.3. Next we will combine the upper part with the lower part using the push-out property.

$$\begin{array}{ccc}
\begin{array}{ccccc}
A & \xrightarrow{p} & B & & \\
\downarrow f & \lrcorner & \downarrow g_1 & & \\
A' & \xrightarrow{\quad} & D & & \\
& & \searrow g_2 & & \\
& & & B' & 
\end{array} & \Rightarrow & \begin{array}{ccccc}
A & \xrightarrow{p} & B & \xrightarrow{q} & C \\
\downarrow f & \lrcorner & \downarrow g_1 & & \parallel \\
A' & \xrightarrow{\quad} & D & \xrightarrow{\quad} & C \\
\parallel & & \downarrow g_2 & & \downarrow h \\
A' & \xrightarrow{p'} & B' & \xrightarrow{q'} & C'
\end{array}
\end{array}$$

It remains to show that the lower right square is commutative, then we can apply the dual of proposition 3.3 to see that the square is bicartesian. Note that  $q = cg_1$  by prop 3.3 thus  $q'g_2g_1 = q'g = hq = hcg_1$ . By uniqueness of the push-out property we have that  $hc = q'g_2$ .  $\square$

**Corollary 3.6.1. The short five lemma.** Suppose that there is a morphism of conflations  $(f, g, h)$  as above. If  $f$  and  $h$  are isomorphisms, then  $g$  is an isomorphism.

*Proof.* Since  $f$  is an isomorphism it is at least an inflation, thus  $g_1$  is an inflation by (QE2). As colimits preserve epis,  $g_1$  is also an epimorphism. by corollary 3.1.1 we know that  $g_1$  is an iso, and dually that  $g_2$  is an iso. Since isomorphisms are closed under composition we have that  $g$  is an isomorphism.

$$\begin{array}{ccccc}
A & \xrightarrow{p} & B & \xrightarrow{q} & C \\
\downarrow f & \lrcorner & \downarrow g_1 & & \parallel \\
A' & \xrightarrow{\quad} & D & \xrightarrow{\quad} & C \\
\parallel & & \downarrow g_2 & \lrcorner & \downarrow h \\
A' & \xrightarrow{p'} & B' & \xrightarrow{q'} & C'
\end{array}$$

$\square$

**Lemma 3.7. Noethers isomorphism lemma.** Suppose there is a diagram with rows as conflations and the first column as a conflation. Then the final column is also a conflation.

$$\begin{array}{ccccc}
A & \xrightarrow{\circ} & B & \xrightarrow{\circ} & X \\
\parallel & & \downarrow \phi & \lrcorner & \downarrow \phi \\
A & \xrightarrow{\circ} & C & \xrightarrow{\circ} & Y \\
& & \downarrow \phi & & \downarrow \phi \\
& & Z & \xlongequal{\quad} & Z
\end{array}$$

*Proof.* Assume that we only have the solid part of the diagram above. By the universal property of cokernels, the upper dashed map exists, and by the dual of proposition 3.3 the square is bicartesian. We can state that the upper dashed map is an inflation, and since the square is a push-out we get that the lower dashed map exists such that the final column is a conflation by proposition 3.3.  $\square$

## 3.2 The Stable Frobenis Category

**Definition 3.4.** Let  $(\mathcal{A}, \mathcal{E})$  and  $(\mathcal{A}', \mathcal{E}')$  be two exact categories. A functor  $F : (\mathcal{A}, \mathcal{E}) \rightarrow (\mathcal{A}', \mathcal{E}')$  is called exact if it is additive and  $F(\mathcal{E}) \subseteq \mathcal{E}'$ . That is to say that conflations get mapped to conflations. One speaks of a reflective exact functor for whenever the pair  $(Fp, Fq)$  is a conflation, then  $(p, q)$  is a conflation.

**Definition 3.5.** Let  $(\mathcal{A}, \mathcal{E})$  be an exact category. An object  $P : \mathcal{A}$  is called projective if  $\mathcal{A}(P, -) : (\mathcal{A}, \mathcal{E}) \rightarrow \mathbf{Ab}$  is an exact functor. Objects  $I : \mathcal{A}$  are called injective whenever  $\mathcal{A}(-, I) : \mathcal{A}^{op} \rightarrow \mathbf{Ab}$  is an exact functor.

*Remark.* In the case of exact functors  $F : (\mathcal{A}, \mathcal{E}) \rightarrow \mathbf{Ab}$ , one generally speaks of a functor which maps conflations to short exact sequences.

*Remark.* The hom-functor is something that is called left-exact. That means that conflations gets mapped to sequences which is only exact in the two first terms.

**Proposition 3.8.** Let  $(\mathcal{A}, \mathcal{E})$  be an exact category.  $P : \mathcal{A}$  is projective if and only if for every deflation  $q : A \rightarrow B$  and morphism  $f : P \rightarrow B$  there is a morphism  $f' : P \rightarrow A$  rendering the diagram below commutative.

$$\begin{array}{ccc}
& & P \\
& \swarrow f' & \downarrow f \\
A & \xrightarrow{q} & B
\end{array}$$

*Proof.* Suppose that  $P$  is projective, then  $\mathcal{A}(P, -)$  is an exact functor. Let  $(p : A \rightarrow B, q : B \rightarrow C)$  be a conflation, then there is a short exact sequence.

$$0 \xrightarrow{0} \mathcal{A}(P, A) \xrightarrow{p_*} \mathcal{A}(P, B) \xrightarrow{q_*} \mathcal{A}(P, C) \xrightarrow{0} 0$$

Pick  $f : \mathcal{A}(P, C)$ , since  $q_*$  is a surjection there exists an  $f' : \mathcal{A}(P, B)$  such that  $pf' = f$ . Now, suppose that  $P$  has the diagram property and that  $(p : A \rightarrow B, q : B \rightarrow C)$  is a conflation. Then there is an exact sequence in  $\mathbf{Ab}$  by  $\mathcal{A}(P, -)$ .

$$0 \xrightarrow{0} \mathcal{A}(P, A) \xrightarrow{p_*} \mathcal{A}(P, B) \xrightarrow{q_*} \mathcal{A}(P, C)$$

To see that  $q_*$  is a surjection and, let  $f : P \rightarrow C$ . As  $q$  is a deflation there exists an  $f' : P \rightarrow B$  such that  $q_*(f') = f$ . Thus the sequence above is short exact and  $P$  is projective.  $\square$

**Corollary 3.8.1.** *Let  $P$  be projective, then if  $q : A \rightarrow P$  is a deflation, it is split-epi.*

**Corollary 3.8.2.** *Two objects  $P$  and  $Q$  are projective if and only if  $P \oplus Q$  is projective.*

**Definition 3.6.** A category  $(\mathcal{A}, \mathcal{E})$  has enough projective objects if for any object  $A : \mathcal{A}$  there is a projective object  $P$  along with a deflation  $q : P \rightarrow A$ .

**Definition 3.7.** An exact category is called a Frobenius category if it has enough projective and injective objects and the class of projective objects coincide with the injective objects.

From the Frobenius categories, one can construct the stable Frobenius category where every morphism factoring through an injective will be identified. These stable categories will give rise to a class of categories called the algebraic triangulated categories. In order to construct the stable Frobenius category we will use the construction of the quotient category.

**Definition 3.8.** A congruence relation  $\sim$  on a category  $\mathcal{C}$  is a relation on the homsets, such that:

1.  $\forall A, B : \mathcal{C}$  the relation  $\sim_{A,B}$  is an equivalence relation.
2. Given that  $f, f' : A \rightarrow B$  is related ( $f \sim f'$ ) and morphisms  $g : A' \rightarrow A$  and  $h : B \rightarrow B'$ , then  $hgf \sim hf'g$ .

**Proposition 3.9.** *Let  $\mathcal{C}$  be a category and  $\sim$  be a congruence relation. Then there is a universal category  $\mathcal{C}/\sim$  together with a functor  $q : \mathcal{C} \rightarrow \mathcal{C}/\sim$  such that morphisms  $f, g : A \rightarrow B$  are identified if  $f \sim g$ . Universality means that if there is a functor  $H : \mathcal{C} \rightarrow \mathcal{D}$  such that  $Hf = Hg$  for any  $f, g$  such that  $f \sim g$ , then  $H$  factors uniquely through  $\mathcal{C}/\sim$ .*

*Proof.* Define the category  $\mathcal{C}/\sim$  to have the same objects as  $\mathcal{C}$ , and define  $\mathcal{C}/\sim(A, B) = \mathcal{C}(A, B)/\sim_{a,b}$ . This definition is well defined as  $\sim$  is a congruence relation. A sketch of this proof can be found in [5].  $\square$

*Remark.* Any functor gives rise to a congruence relation. That is, if  $F : \mathcal{C} \rightarrow \mathcal{D}$  is a functor, then there is a congruence relation  $\sim$  defined as follows:  $\forall A, B : \mathcal{C}$  and  $f, g : A \rightarrow B$ , we define  $f \sim_{A,B} g \iff Ff = Fg$ . This is a congruence as equality within  $\mathcal{D}$  gives rise to an equivalence relation, and functoriality gives the congruence.

*Remark.* For any relation  $\sim$  the universal category  $\mathcal{C}/\sim$  exists. As in the case for the Verdier quotient,  $\mathcal{C}/\sim$  is the same as the quotient category of the smallest congruence relation having the same relations as  $\sim$ .

If  $\mathcal{A}$  is an additive category, the quotient categories which respects the additive structures are of interest. That is to say that the functor  $q : \mathcal{A} \rightarrow \mathcal{C}/\sim$  is additive and the equivalence relation  $\sim$  should respect the additive structure. Then a quotient category is additive if  $f \sim f'$  and  $g \sim g'$ , then  $f + g \sim f' + g'$ . This leads to the following definition.

**Definition 3.9.** Let  $\mathcal{A}$  be an additive category.  $\mathcal{I}$  is an ideal of  $\mathcal{A}$  if:

1. (subgroup) for every abelian group  $\mathcal{A}(A, B)$  there is a subgroup  $\mathcal{I}(A, B) \subseteq \mathcal{A}(A, B)$ .
2. (absorption) For every  $g : A' \rightarrow A$ ,  $h : B \rightarrow B'$  and  $f : \mathcal{I}(A, B)$  we have that  $hfg : \mathcal{I}(A', B')$

This is equivalent of saying that the equivalence relation gained from  $\mathcal{I}(A, B)$  ( $f \sim g \iff f - g : \mathcal{I}(A, B)$ ) is a congruence relation.

**Corollary 3.9.1.** Let  $\mathcal{A}$  be an additive category and  $\mathcal{I}$  be an ideal of  $\mathcal{A}$ , then  $\mathcal{A}/\mathcal{I}$  is an additive category.

Let  $\mathcal{A}$  be a Frobenius category. Define the ideal  $\mathcal{I}$  as the subgroups of every morphism factoring through injective objects.

**Proposition 3.10.** For any Frobenius category  $\mathcal{A}$  the ideal  $\mathcal{I}$  is well defined and  $\underline{\mathcal{A}} = \mathcal{A}/\mathcal{I}$  is the stable Frobenius category.

*Proof.* To prove this we have to show that  $\mathcal{I}(A, B)$  is a subgroup for any  $A, B : \mathcal{A}$ , and that it is absorptive. First observe that  $0 : \mathcal{I}(A, B)$ . Let  $f, g : \mathcal{I}(A, B)$ . Since  $\mathcal{A}$  has enough injectives there exists an injective object with an inflation from  $A$ .

$$\begin{array}{ccccc}
 & & J_1 & & \\
 & f_1 \nearrow & \uparrow f'_1 & \nwarrow f_2 & \\
 A & \xrightarrow{i} & I & & B \\
 & g_1 \searrow & \downarrow g'_1 & \nearrow g_2 & \\
 & & J_2 & & 
 \end{array}$$

$f - g = f_2 \circ f_1 - g_2 \circ g_1 = (f_2 \circ f'_1 - g_2 \circ g'_1) \circ i$ . Thus  $f - g$  factors through an injective, and  $\mathcal{I}(A, B)$  is a subgroup. To see that it is absorptive is to see that if  $f$  factors through an injective, then  $gf$  factors through an injective as well.  $\square$

Objects in the stable Frobenius category is denoted as  $\underline{X}$  and morphisms are denoted as  $\underline{f}$ . That is the functor  $q : \mathcal{A} \rightarrow \underline{\mathcal{A}}$  is defined as  $q(X) = \underline{X}$  and  $q(f) = \underline{f}$ . One important property is that in the stable Frobenius category, taking syzygies or cosyzygies is a functor.

**Lemma 3.11.** *Let  $\mathcal{A}$  be a Frobenius category and suppose that there are two conflations with injectives as below. Then  $\underline{X}' \simeq \underline{X}''$ .*

$$X \rightharpoonup \circlearrowleft \xrightarrow{i} I \rightharpoonup \circlearrowleft \xrightarrow{i'} X'$$

$$X \rightharpoonup \circlearrowleft \xrightarrow{j} J \rightharpoonup \circlearrowleft \xrightarrow{j'} X''$$

*Proof.* Observe that with the diagram we can find morphisms as  $I$  and  $J$  are injective.

$$\begin{array}{ccc} X & \xrightarrow{i} & I \\ \parallel & \searrow j & \downarrow f \\ X & \xrightarrow{j} & J \end{array}$$

By iterating this we get the following commutative diagram in  $\mathcal{A}$ .

$$\begin{array}{ccccc} X & \xrightarrow{i} & I & \xrightarrow{i'} & X' \\ \parallel & & \downarrow f & & \downarrow g \\ X & \xrightarrow{j} & J & \xrightarrow{j'} & X'' \\ \parallel & & \downarrow f' & & \downarrow g' \\ X & \xrightarrow{i} & I & \xrightarrow{i'} & X' \end{array}$$

By chasing the diagram we get that  $i - f'fi = (id_I - f'f) \circ i = 0$ , this means that  $(f'f - id_I)$  factors through  $X'$ , i.e. there exists  $h : X' \rightarrow I$  and  $f'f = hi' + id_I$ . We also have that  $g'gi' = i'f'f = i'(hi' + id_I) = i'hi' + i' = (i'h + id_{X'})i'$ . As  $i'$  is an epi we get that  $g'g = i'h + id_{X'} \implies \underline{g'g} = id_{\underline{X'}}$  as  $i'h$  factors through  $I$ .  $\underline{gg'} = id_{\underline{X''}}$  is dual.  $\square$

**Corollary 3.11.1.** *Cosyzygy is a well defined functor  $\mathcal{U} : \underline{\mathcal{A}} \rightarrow \underline{\mathcal{A}}$*

*Proof.* Let  $f : X \rightarrow Y$  be a morphism in  $\mathcal{A}$ . Then we have the following diagrams representing the different choices of syzygies.

$$\begin{array}{ccccc} X & \xrightarrow{i} & I & \xrightarrow{p} & \mathcal{U}X \\ \downarrow f & & \downarrow & & \downarrow \mathcal{U}f \\ Y & \xrightarrow{j} & J & \xrightarrow{q} & \mathcal{U}Y \end{array} \quad \begin{array}{ccccc} X & \xrightarrow{i'} & I' & \xrightarrow{p'} & \mathcal{U}'X \\ \downarrow f & & \downarrow & & \downarrow \mathcal{U}'f \\ Y & \xrightarrow{j'} & J' & \xrightarrow{q'} & \mathcal{U}'Y \end{array}$$

By the previous proof we know there are maps between the diagrams making an almost commutative diagram where all the outer squares commute.

To see that the definition of the cosyzygy is well defined is to show that the inner squares is commutative in the quotient category.

Observe that the inner square commutes by definition, and that the central inner square commutes in the quotient as every morphism gets related to 0. Thus it remains to show that  $\gamma(Y) \circ \bar{U}f = \bar{U}'f \circ \gamma(X)$ , which is the same as to say that  $\gamma(Y) \circ \bar{U}f - \bar{U}'f \circ \gamma(X)$  factors over an injective.

By doing a diagram chase in the left cube one may find the following equation  $(I'f \circ \beta(X) - \beta(Y) \circ If)i = 0$ . This means that the map  $I'f \circ \beta(X) - \beta(Y) \circ If$  factors through the cokernel of  $i$  as  $\chi p$ . By chasing the right cube we get the following equation  $q' \chi p = (\gamma(Y) \circ \bar{U}f - \bar{U}'f \circ \gamma(X))p$ , thus  $q' \chi = \gamma(Y) \circ \bar{U}f - \bar{U}'f \circ \gamma(X)$ .  $\square$

**Corollary 3.11.2.** *Cosyzygy is an autoequivalence with syzygy as inverse.*

### 3.3 Self-injective Algebras

### 3.4 The Homotopy Category

### 3.5 Admissable Morphisms, Homology and Long Exact Sequences

### 3.6 The Derived Category



## 4 Auslander-Reiten Triangles

### 4.1 Krull-Schmidt Categories

### 4.2 Definition and First Properties

### 4.3 Description of Derived Categories

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