

Design document:

Group 2D

Registers

\$0: A register that always contains 0. Reg #0

\$acc: The accumulator register. Reg #1

\$a0-3: argument temporary registers. Reg #2-5

\$v0-1: return temporary registers. Reg #6-7

\$t0-1: temporary registers. Reg #8-9

\$s0-1: safe registers. Reg #10-11

\$ra: A register which contains the return address. Reg #12

\$sp: A register which contains the stack pointer. Reg #13

\$at0-1: Assembler temporaries. Reg #14-15

Instructions

Types

Instruction Fields:

Type I:

op	reg	imm/unused
4 bits	4 bits	8 bits

op: Basic operation of the instruction.

des: The register/accumulator destinator operand

imm: immediate value that represents a constant value or an address in memory.

Type II:

op	imm/unused	imm
4 bits	4 bits	8 bits

op: Basic operation of the instruction

imm: immediate value.

Type III: For beq and bne

op	reg	des	unused
4 bits	4 bits	4 bits	4 bits

op: Basic operation of the instruction

List of instructions

add: Type I instruction with opcode 0 0

Adds value at register into accumulator value and stores at accumulator

sub: Type I instruction with opcode 1 1

Subtracts value at register with accumulator value and stores at accumulator(decide order using unused bits. If no imm in instruction, assembler assumes 0. 1 means $acc = acc - A$, 0 means $acc = A - acc$)

lmem: Type I instruction with opcode 10 2

Takes the value at a register, adds the immediate to create a memory address, then gets the value from memory and puts it in the accumulator

smem: Type I instruction with opcode 11 3

Takes the value at a register, adds the immediate to create a memory address, then sets the memory at the location to the value in the accumulator

beq: Type III instruction with opcode 100 4

Conditionally branch if $reg == acc$, then go to the mem location contained in dest.

bne: Type III instruction with opcode 101 5

Conditionally branch if $reg != acc$, then go to the mem location contained in dest.

sll: Type II instruction with opcode 110 6

Left shift the value in accumulator by constant value

slt: Type I instruction with opcode 111 7

Decide order using unused bits. If no imm in instruction, assembler assumes 0. 0 means $acc < reg$, 1 means $reg < acc$. If inequality is true, acc gets 1. else acc gets 0.

or: Type I instruction with opcode 1000 8

Bitwise OR operation for a reg and acc

and: Type I instruction with opcode 1001 9

Bitwise AND operation for a reg and acc

lui: Type II instruction with opcode 1010 A

Load the upper half of an immediate into accumulator

jr: Type I instruction with opcode 1011 B

Jump to an address within a register

jal: Type I instruction with opcode 1100 C

Jump to an address within a register and save ra in the register

ori: Type II instruction with opcode 1101 D

Bitwise or zero extended immediate into the accumulator

lacc: Type I instruction with opcode 1110 E

Load value in specified register into the acc

sacc: Type I instruction with opcode 1111 F

Store value in acc to given reg

Pattern

Procedure calls

During procedure calls, we will still back up all the registers we need including \$ra into stack, and restore them when finished. We use \$a0-\$a3 to store arguments, and \$v0 and \$v1 to store return values. If we need more arguments or return values, we will store them in the stack. Backing up registers looks like the following: If acc is needed, back it up first. Afterwards, move values from registers into the acc, then onto the stack.

Translate to machine code

For logical instructions, it will zero extend immediates. For arithmetic instructions, it will be sign extend.

For op codes and registers, it will zero extend to 4 bits.

For beq, bne, jr, and jal we use direct addressing, meaning it will go to the address in the register specified.

Example and test cases

Example Assembly Code for Relative Prime:

Hex Code:

```
relPrime:      #$a0=n
    lui    0                A000
    ori    2                D020
    sacc   $s0 #s0=m=2      FA00

loop:
    lacc   $s0 #for loop    EA00
    sacc   $a1 #a1=m        F200
    lacc   $sp #stack       ED00
    lui    -12              AF40
    ori    -12              DFF0
    add    $sp              0D00
    sacc   $sp #$sp-=12     FD00

    #Code back up $ra, $a0, $s0
    lacc   $ra              EC00
    smem   $sp, 0           3D00
    lacc   $a0              E200
    smem   $sp, 4           3D04
    lacc   $s0              EA00
    smem   $sp, 8           3D08
    lui    gcd              A address of gcd
    ori    gcd              D address of gcd
    jal    $acc             C100

    #Code to restore $ra, $a0, $s0
    lmem   $sp, 0           2D00
    sacc   $ra              FC00
    lmem   $sp, 4           2D04
    sacc   $a0              F200
    lmem   $sp, 8           2D08
    sacc   $s0              FA00
```

lacc	\$sp	#stack	ED00
lui	0		A000
ori	12		D0C0
add	\$sp		0D00
sacc	\$sp	#\$sp+=12	FD00
lui	output		A address of output
ori	output		D address of output
sacc	\$t0		F800
lui	0		A000
ori	1		D010
beq	\$v0, \$t0		4680
add	\$s0		0A00
sacc	\$s0	#m=m+1	FA00
lui	loop		A upper of addr of loop
ori	loop		D upper of addr of loop
jr	\$acc		B100

output:

lacc	\$s0	#\$acc = m	EA00
sacc	\$v0	#\$v0 = m	F600
jr	\$ra		BC00

gcd:

lui	bcase		A address of bcase
ori	bcase		D address of bcase
sacc	\$t0		F800
lui	0		A000
beq	\$a0, \$t0		4810

loop:

lui	acase		A address of acase
ori	acase		D address of acase

	sacc	\$t0	F800
	lui	0	A000
	beq	\$a1, \$t0	4380
	lacc	\$a0 #acc=a	E200
	slt	\$a1 #acc=(b<a?1:0)	7300
	sacc	\$t0	F800
	lui	asub	A address of asub
	ori	asub	D address of asub
	sacc	\$t1	F900
	lui	0	A000
	ori	1	D010
	beq	\$t0, \$t1	4890
	lacc	\$a1	E300
	sub	\$a0 #a<b	1200
	sacc	\$a1	F300
loop	lui	loop	A upper address of
loop	ori	loop	D lower address of
	jr	\$acc	B100
bcase:			
	lacc	\$a1	E300
	sacc	\$v0	F600
	jr	\$ra	BC00
asub:			
	lacc	\$a0	E200
	sub	\$a1	1300
	sacc	\$a0	F200
loop	lui	loop	A upper address of
loop	ori	loop	D lower address of
	jr	\$acc	B100

```

acase:
    lacc    $a0            E200
    sacc    $v0            F600
    jr      $ra            BC00

```

Use psuedo instrutions (we use to run the relprime):

```

relPrime:
li $s0 2
loop:
lacc $s0
sacc $a1
addi $sp $sp 3
lacc $ra
smem $sp 0
lacc $a0
smem $sp -1
lacc $s0
smem $sp -2
jalp gcd
lmem $sp 0
sacc $ra
lmem $sp -1
sacc $a0
lmem $sp -2
sacc $s0
addi $sp $sp -3
li $acc 1
beqp $v0 $acc output
add $s0
sacc $s0
j loop
output:
lacc $s0

```



```

sacc $v0
jr $ra
gcd:
beq $a0 $0 bcase
loop:
beq $a1 $0 acase
lacc $a0
slt $a1
sacc $t0
li $acc 1
beq $t0 $acc asub
lacc $a1
sub $a0 1
sacc $a1
j loop
bcase:
lacc $a1
sacc $v0
jr $ra
asub:
lacc $a0
sub $a1 1
sacc $a0
j loop
acase:
lacc $a0
sacc $v0
jr $ra

```

Loading an address into accumulator:

```

    lui    addr
addr
    ori    addr
addr

```

A upper address of

D lower address of

Iteration:

```
loop:
    lui    0                A000
    ori    1                D010
    add    $a0              0200
    sacc   $a0              F200
    lui    loop             A upper address of
loop
    ori    loop             D lower address of
loop
    jr     $acc             B100
```

Conditional Statement:

```
top:
    lui    0                A000
    slt    $a0              7200
    sacc   $t0              F800
    lui    0                A000
    lui    case             A address of case
    ori    case             D address of case
    beq    $t0, $acc        4810

case:
    lui    top              A upper address of top
    ori    top              D lower address of top
    jr     $acc             B100
```

RTL

See below

Logical/Arithmetic/Slr	sacc/lacc	lmem/smem	sll	beq/bne	lui	jr/jal
IR = Mem[PC] PC = PC + 1						
C = Reg[1] A = Reg[IR[11-8]] B = Reg[IR[7-4]] ori: order = IR[8] otherwise: order = IR[0]						
ori: Use ZE[IR[7-0]] instead of A if(order==0) ALUOut=A op C else ALUOut=C op A	sacc: Reg[IR[11-8]] = C lacc: Reg[1] = A	ALUOut=A+ SE(IR[7-0])	Reg[1] = C<<IR[11-0]	beq: if(A==C) PC=B bne: if(A!=C) PC=B	Reg[1] = ZE(IR[7-0])<<8	jal: Reg[12]=PC jr and jal: PC = A
Reg[1] = ALUOut		lmem: MDR = Mem[ALUOut] smem: Mem[ALUOut] = C				
		lmem: Reg[1] = MDR				

Component description

Component	Description	Inputs	Outputs	RTL Interactions
Register file	<p>Used for register operations including the Accumulator. Will take in IR[11-8] and IR[7-4] and output into A and B respectively. The register file will also have a write address input that will mainly be the accumulator register, and a write data which will be the data written into the register. This data will be coming from ALUout and the memory. Lastly, it will output the value in \$acc to C.</p> <p>Other register files are A, B, C, ALUOut, IR, MDR, and PC</p>	IR[11-8] - ra1 - 4b IR[7-4] - ra2 - 4b write address - wa - 4b write data - wd - 16b	rd1 - A - 16b rd2 - B - 16b rd3 - C - 16b	IR PC C A B Order
Memory	<p>Will mainly be used to find the instruction, but also will be used to store values. Takes in an input from the PC if the memory read signal is on, and outputs the correct instruction from memory. Also can take in a signal and write it to memory if the Memory write signal is on.</p>	Address in - addin - 16b Data in - din - 16b	Data out - do - 16b	PC IR ALUout MDR
ALU	<p>ALU will take in two inputs and perform an operation depending on the ALU control. This can be +, -, &, , or <.</p>	input 1 - A - Output of ALU source A mux - 16b input 2 - B - Output of ALU source B mux - 16b	ALUout - 16b	PC A C ZE(IR[11-4])

		ALU control -3b input 3 - order bits that decide which order the inputs should be operated on		
Control	Decides the value of the various muxes, write and read codes to perform a certain operation based on the Instruction Registers output.	instruction from memory - 16b	Various mux signals, write and read signals, and ALU control - bits depend on amount of mux inputs	Controls what RTL runs
Shift left	shifts the value coming in left filling the end with zeroes and then outputting it.	Value coming in from ALU or instruction - 16b Number to shift by - 12b	Value shifted	IR C Register file
Comparator	Compares two signals to see if they are equal	C and A	BranchDecide	A and C

Control Signals

Control Signals	Description
Instruction or data IorD - 1b	Controls whether the PC goes into memory to find instruction or data to write to memory
ALU source B ALUSrcB - 2b	Controls which value goes into the second input of the ALU
ALU source A ALUSrcA - 2b	Controls which value goes into the first input of the ALU
PC source PCSrc - 2b	Controls whether PC + 1 goes into the PC or a jump/branch target
Shift Amount ShiftAmt - 1b	Decides the value that goes into the shift input to choose how much to shift by.

Shift Select ShiftSrc - 1b	Decides the value that will be shifted. Sign extended order or Accumulator.
Register Destination RegDst - 2b	Chooses which register to write to.
Register Data RegData - 3b	Chooses which value gets put in the register.
Order Control OrderControl - 1b	Chooses what bit of the instruction actually constitutes the order bits
Register File write: regwrite - 1b	Chooses Whether or not to write to the register file.
Memory write: memw - 1b	Chooses whether or not to write to Memory.
Memory read: memr - 1b	Chooses whether or not to read from Memory.
PC Write PCWrite 1b	Chooses whether or not to write to the PC
PC Write Conditional PCWriteCond 1b	Chooses whether a conditional will allow the PC to write
Instruction Register Write IRwrite 1b	Chooses whether or not to write to the IR.
ALU control 3b	Decides the operation of the ALU
EorNE	Chooses whether the comparator should compare equals or not equals for the two inputs to the comparator

RTL Test

This is code and RTL written by Anderson and Yuankai. Anderson was a little bit involved with creating the original RTL, but most of the RTL design was done by Caleb and Brian. We made this test code for them to think about what the RTL did and to find any errors we may have made when designing the RTL.

```
RA=0x6666
```

```
PC=0x0000
```

```
SP=0x7FFF
```

```
test:
```

```
0x0000      lui    0
             ori    0
             sacc   $t0
             lui    0
             ori    1
             add    $t0
             sacc   $a0
             lui    (upper)-4
             ori    (lower)-4
             add    $sp
             sacc   $sp
             lacc   $ra
             smem   $sp, 0
             lui    (upper)branch
             ori    (lower)branch
             jal    $acc
             lmem   $sp, 0
             sacc   $ra
             lui    0
             ori    4
             add    $sp
             sacc   $sp
             jr     $ra
```

```
branch:
```

```
0x0100      lui    0
             sacc   $t0
             sub    $a0
             slt    $t0, 1
```

```

        sacc $t1
        lui  0
        beq  $t1, $acc
        jr   $ra

```

start:

```

        IR=0x0000
        PC=0x0000

```

lui 0:

```

        0xA000
        PC=0x0004
        Reg[1]=0x0000=>(sll 8) 0x0000

```

ori 2:

```

        IR=0xD020
        PC=0x0008
        A=>0x02=>(ZE) 0x0002
        ALUOut=A or Reg[1]=0x0002
        Reg[1]=ALUOut=0x0002

```

sacc \$t0:

```

        IR=0xF800
        PC=0x000C
        Reg[8]=Reg[1]=0x0002

```

lui 0:

```

        IR=0xA000
        PC=0x0010
        Reg[1]=0x0000=>(sll 8) 0x0000

```

Ori 1:

```

        IR=0x0010
        PC=0x0014
        A=>0x01=>(ZE) 0x0001
        ALUOut=A or Reg[1]=0x0001
        Reg[1]=ALUOut=0x0001

```

add \$t0:

```

        IR=0x0800
        PC=0x0018
        A=Reg[8]
        ALUOut=Reg[1]+Reg[8]=0x0001+0x0002=0x0003
        Reg[1]=ALUOut=0x0003

```

sacc \$a0:

```

        IR=0xF200
        PC=0x001C
        Reg[2]=Reg[1]=0x0003

```



```

lui (upper)-4:
    IR=0xAFF0
    PC=0x0020
    Reg[1]=0x00FF=>(sll 8) 0xFF00
ori (lower)-4:
    IR=0xDFC0
    PC=0x0024
    A=>0xFC=>(ZE) 0x00FC
    ALUOut=A or Reg[1]=0xFFFC
    Reg[1]=ALUOut=0xFFFC
add $sp:
    IR=0x0000
    PC=0x0028
    A=Reg[13]
    ord=0
    ALUOut=Reg[1]+A=0xFFFC+0x7FFF=0x7FFB
sacc $sp:
    IR=0xFD00
    PC=0x002C
    Reg[13]=Reg[1]=0x7FFB

lacc $ra:
    IR=0xEC00
    PC=0x0030
    A=Reg[12]
    Reg[1]=Reg[13]=0x6666
smem $sp, 0:
    IR=0x3D00
    PC=0x0034
    A=Reg[13]=0x7FFB
    ALUOut=A+(SE) 0=0x7FFB
    MEM[0x7FFB]=Reg[1]=0x6666
lui (upper)branch:
    IR=0xA010
    PC=0x0038
    Reg[1]=0x0001=>(sll 8) 0x0100
ori (lower) branch?
    IR=0xD000
    PC=0x003C
    A=0x00
    ALUOut=Reg[1] or ZE(A)=0x????
    Reg[1]=0x0100

```

```

jal $acc:
    IR=0xC100
    PC=0x0040
    Reg[12]=0x0040
    PC=Reg[1]=0x0100
lmem $sp, 0:
    IR=0x2D00
    PC=0x0044
    A=Reg[13]
    Reg[1]=A=
sacc $ra:
    IR=0xFC00
    PC=0x0048
    A=0x0C=12
    Reg[12]=Reg[1]=0x6666
lui 0:
    IR=0xA000
    PC=0x004C
    Reg[1]=0x00=>(sll 8) 0x0000
ori 4:
    IR=0xD400
    PC=0x0050
    A=0x04
    Reg[1]=Reg[1] or ZE(A)=0x0004
add $sp:
    IR=0x0D00
    PC=0x0054
    Reg[1]=Reg[1]+Reg[13]=0x0004+0x7FFB=0x7FFF
sacc $sp:
    IR=0xFD00
    PC=0x0058
    A=0x0D=13
    Reg[13]=Reg[1]=0x7FFF
jr $ra:
    IR=0xBC00
    PC=0x005C
    PC=Reg[1]=0x6666

branch:
lui 0:
    IR=0xA000
    PC=0x0100

```

```

        Reg[1]=0x0000=>sll 8 0x0000
sacc $t0:
        IR=0xF800
        PC=0x0104
        Reg[8]=reg[1]=0x0000
sub $a0, 1:
        IR=0x2201
        PC=0x0108
        A=Reg[2]=0x0003
        ord=1
        ALUOut=Reg[2]-Reg[1]=0x0003-0x0000=0x0003
        Reg[1]=ALUOut=0x0003
slt $t0, 1:
        IR=0x7801
        PC=0x010C
        ord=1
        ALUOut=Reg[8]<Reg[1]=0x0000<0x0000=0
        Reg[1]=ALUOut=0x0000
sacc $t1:
        IR=0xF900
        PC=0x0110
        Reg[9]=Reg[1]=0x0000
lui 0:
        IR=0xA000
        PC=0x0114
        Reg[1]=0x0000=>(sll 8) 0x0000
lui (upper) test:
        IR=0xA000
        PC=0x0118
        Reg[1]=0x0000=>(sll 8) 0x0000
ori (lower) test:
        IR=0xD000
        PC=0x001C
        A=0x00
        Reg[1]=Reg[1] or ZE(A)=0x0000
beq $t1, $acc:
        IR=0x4910
        PC=0x0120
        PC=Reg[1]==>0x0000
jr $ra:
        IR=0xBC00
        PC=0x0124

```

`PC=Reg[1]=0x0000`

Comments for RTL design: When Anderson and I (Kevin) wrote the test for rtl, we feel the instructions are easy to understand and write. We can easily understand the way it works. The RTL combines many similar instructions, so it is easy to know the pattern. The test is mainly show every small part works as expected. It does not purposefully calculate anything.

Component Test Plan

Step1: Individual test

1. ALU Test

Input	Output
ALUcontrol:000(and) input1[7:0]:from 0000 0000 to 1111 1111 input2[7:0]:from 0000 0000 to 1111 1111	[7:0]:input1 and input2
ALUcontrol:001(or) input1[7:0]:from 0000 0000 to 1111 1111 input2[7:0]:from 0000 0000 to 1111 1111	[7:0]:input1 or input2
ALUcontrol:010(add) input1[7:0]:from 0000 0000 to 1111 1111 input2[7:0]:from 0000 0000 to 1111 1111	[7:0]:input1+input2 (may overflow)
ALUcontrol:110(sub) input1[7:0]:from 0000 0000 to 1111 1111 input2[7:0]:from 0000 0000 to 1111 1111	[7:0]:input1-input2 (may overflow)
ALUcontrol:111(slt) input1[7:0]:from 0000 0000 to 1111 1111 input2[7:0]:from 0000 0000 to 1111 1111	[7:0]: if (input1<input2) 1 else 0

2. Mem Test

Input	Output
save things into the memory (once at a time)	load from memory to see if it the same
save things into the memory (several data at a time)	load from memory to see if it the same

3. Register file Test

Input	Output
save things into register (once at a time)	load from register to see if it the same
save things into register (several data at a time)	load from register to see if it the same

4. SE Test

Input	Output
input[x:0]:from 0000 0000 to 1111 1111	[15:x+1]=input[x] [x:0]=input

5. ZE Test

Input	Output
input[x:0]:from 0000 0000 to 1111 1111	[15:x+1]=0 [x:0]=input

6. Left Shift Test

Input	Output
number to shift[7:0]:from 0000 0000 to 1111 1111 bits to shift[1:0]: from 0 to 8	[7:0]:number to shift<<bits to shift

7. == Test

Input	Output
input1[7:0]:from 0000 0000 to 1111 1111 input2[7:0]:from 0000 0000 to 1111 1111	input1==input2 1 for true, 0 for false

8. mux Test

Input	Output
mux control:from 000 to 101 input0, input1, input2, input3, input4, input5	mux control=000, input0 mux control=001, input1 mux control=010, input2 mux control=011, input3 mux control=100, input4 mux control=101, input5

Step2: Integration test

1. PC+1 Test

Components	Input	Output
PC, ALU	PC (include edge cases like 0000 0000, FFFF FFFF)	PC=PC+1

2. IR=Mem[PC] Test

Components	Input	Output
Mem, IR, PC	PC (include edge cases like 0000 0000, FFFF FFFF)	IR=Mem[PC]

3. sacc Test

Components	Input	Output
Reg, IR, C	C IR	Reg[IR[11-8]]=C Check if the register has stored the value

4. lacc Test

Components	Input	Output
Reg, A	A	Reg[1]=A load the value in accumulator out to see the value matches

5. Reg and Other calculation Test

Components	Input	Output
IR, Reg, A, B, C, ==, SE, ZE, <<	A=Reg[IR[11-8]] B=Reg[IR[7-4]] C=Reg[1]	Get the data from instruction and register file, and run them in the == or SE or ZE or << to do some calculation

6. RegToMem Test

Components	Input	Output
------------	-------	--------

ALU, ALUOut, A, Mem, C, IR	IR, A, C	ALUOut=A+SE(IR[7-0]) Mem[ALUOut]=C Check if the memory has stored the value at required address
----------------------------	----------	---

7. MemToReg Test

Components	Input	Output
ALU, ALUOut, MDR, A, IR, Reg, Mem	ALU, A	ALUOut=A+SE(IR[7-0]) MDR=Mem[ALUOut] Reg[1]=MDR Check if the \$acc has stored the changed value

8. PCReg Test

Components	Input	Output
PC, IR, ALU, ALUOut, Reg	IR, PC	Write the value of the PC into Register Load the value stored in Register into PC

9. beq Test

Components	Input	Output
A, C, PC, B	A, C, PC, B	if(A==C) PC=B Check if PC is equal to B when A is equal to C

10. bne Test

Components	Input	Output
A, C, PC, B	A, C, PC, B	if(A!=C) PC=B Check if PC is equal to B when A is not equal to C

11. lui Test

Components	Input	Output
Left Shift Unit(LSU), SE Unit, Reg, IR	IR	Reg[1]=IR[7-0]<<8

		Check if \$acc has the correct value in register
--	--	--

12. jr Test

Components	Input	Output
PC, Reg, IR	PC, IR	PC=Reg[IR[11-8]] Check if PC jumps to the correct value

13. jal Test

Components	Input	Output
PC, Reg, IR	PC, IR	Reg[12]=PC PC=Reg[IR[11-8]] Check if PC jumps to the correct value Check if \$ra is equal to the place it should jump back

14. PCMemReg Test

Components	Input	Output
PC, Mem, IR, Reg	PC, IR	Store PC into Register ,and then store this value into Memory Load from Memory into Register, then set its value to PC

15. PCMemRegALU Test

Components	Input	Output
PC, Mem, IR, Reg, ALU	PC, IR	Store PC into Register, do some calculation, and then store this value into Memory Load from Memory into Register, do some calculation, and then set its value to PC

16. Reg ALU Test

Components	Input	Output
Reg, A, B, C, ALU	A=Reg[IR[11-8]] B=Reg[IR[7-4]] C=Reg[1]	Get the data from register file and run them in the ALU to do some calculation

Step3: Full datapath test

When we test it, we will need to save all the instructions to memory first, edit any possible inputs in the test bench so that the input will be set to \$a0.

1. A small piece of test code to just combine some instructions together.

Column1	Cout	t0	t1	s0	instructions in hex
lui 0	0	X	X	X	a000
ori 12	12	X	X	X	d00c
sacc \$t2	12	X	X	12	fa00
lui 0	0	X	X	12	a000
ori 4	4	X	X	12	d004
sacc \$t1	4	X	4	12	f900
lui 0	0	X	4	12	a000
ori 1	1	X	4	12	d001
sacc \$t0	1	1	4	12	f800
lacc \$t0	1	1	4	12	e800
add \$t0	2	1	4	12	800
sacc \$t0	2	2	4	12	f800
lacc \$t1	4	2	4	12	e900
bne \$t0 \$t2	4	2	4	12	58a0
lacc \$t0	2	2	4	12	e800
add \$t0	4	2	4	12	800
sacc \$t0	4	4	4	12	f800
lacc \$t1	4	4	4	12	e900
bne \$t0 \$t2	4	4	4	12	58a0
lui 0	0	4	4	12	a000
ori 3	3	4	4	12	d003

jr \$acc	3	4	4	12	b100
----------	---	---	---	----	------

We look into the waves shown in the simulation and it all works as expected. We manually select the address we need to branch or jump to in this test code.

2. The Eculid's algorithm

We mainly test it first with an input of 4 to see how it works on a small number. After it works, we try large numbers. We mainly debug by seeing the position of our PC is in the right place or not. If it is in the right place, then we have a large chance that it did the correct calculation, and therefore branch or jump to the right place.

PC	Lable	Pseudo Instructions	Assembly Code	Hex	Binary
0		lui b00000000	lui b00000000	a000	1010000000000000
1		ori b01110100	ori b01110100	d074	1101000001110100
2		sacc \$sp	sacc \$sp	fd00	1111110100000000
3		lacc \$0	lacc \$0	e000	1110000000000000
4	relPri me:	li \$s0 2	sacc \$at0	fe00	1111111000000000
5			lui b00000000	a000	1010000000000000
6			ori b00000010	d002	1101000000000010
7			sacc \$s0	fa00	1111101000000000
8			lacc \$at0	ee00	1110111000000000
9	loop:	lacc \$s0	lacc \$s0	ea00	1110101000000000
10		sacc \$a1	sacc \$a1	f300	1111001100000000
11		addi \$sp \$sp3	sacc \$at0	fe00	1111111000000000
12			lui b00000000	a000	1010000000000000
13			ori b00000011	d003	1101000000000011
14			add \$sp	0d00	0000110100000000

15		sacc \$sp	fd00	1111110100000000
16		lacc \$at0	ee00	1110111000000000
17	lacc \$ra	lacc \$ra	ec00	1110110000000000
18	smem \$sp0	smem \$sp0	3d00	0011110100000000
19	lacc \$a0	lacc \$a0	e200	1110001000000000
20	smem \$sp-1	smem \$sp-1	3dff	0011110111111111
21	lacc \$s0	lacc \$s0	ea00	1110101000000000
22	smem \$sp-2	smem \$sp-2	3dfe	0011110111111110
23	jal pgcd	sacc \$at0	fe00	1111111000000000
24		lui b00000000	a000	1010000000000000
25		ori b00111100	d03c	1101000000111100
26		sacc \$at1	ff00	1111111100000000
27		lacc \$at0	ee00	1110111000000000
28		jal \$at1	cf00	1100111100000000
63	gcd: addi \$sp \$sp3	sacc \$at0	fe00	1111111000000000
64		lui b00000000	a000	1010000000000000
65		ori b01011111	d05f	1101000001011111
66		sacc \$at1	ff00	1111111100000000
67		lacc \$a0	e200	1110001000000000
68		beq \$0 \$at1	40f0	0100000011110000
69		lacc \$at0	ee00	1110111000000000
70	loop: beqp \$a1 \$0 acase	sacc \$at0	fe00	1111111000000000
71		lui b00000000	a000	1010000000000000
72		ori b01101001	d069	1101000001101001
73		sacc \$at1	ff00	1111111100000000
74		lacc \$a1	e300	1110001100000000
75		beq \$0 \$at1	40f0	0100000011110000

76		lacc \$at0	ee00	1110111000000000
77	lacc \$a0	lacc \$a0	e200	1110001000000000
78	slt \$a1	slt \$a1	7300	0111001100000000
79	sacc \$t0	sacc \$t0	f800	1111100000000000
80	li \$acc1	sacc \$at0	fe00	1111111000000000
81	beqp \$t0 \$acc asub	lui b00000000	a000	1010000000000000
82		ori b00000001	d001	1101000000000001
83		sacc \$acc	f100	1111000100000000
84		sacc \$at0	fe00	1111111000000000
85		lui b00000000	a000	1010000000000000
86		ori b01100001	d061	1101000001100001
87		sacc \$at1	ff00	1111111100000000
88		lacc \$t0	e800	1110100000000000
89		beq \$at0 \$at1	4ef0	0100111011110000
90		lacc \$at0	ee00	1110111000000000
103	asub: lacc \$a0	lacc \$a0	e200	1110001000000000
104	sub \$a1 1	sub \$a1	1301	0001001100000001
105	sacc \$a0	sacc \$a0	f200	1111001000000000
106	j loop	sacc \$at0	fe00	1111111000000000
107		lui b00000000	a000	1010000000000000
108		ori b01000010	d042	1101000001000010
109		sacc \$at1	ff00	1111111100000000
110		lacc \$at0	ee00	1110111000000000
111		jr \$at1	bf00	1011111100000000
70	loop: beqp \$a1 \$0 acase	sacc \$at0	fe00	1111111000000000
71		lui b00000000	a000	1010000000000000
72		ori b01101001	d069	1101000001101001

73		sacc \$at1	ff00	1111111100000000
74		lacc \$a1	e300	1110001100000000
75		beq \$0 \$at1	40f0	0100000011110000
76		lacc \$at0	ee00	1110111100000000
77	lacc \$a0	lacc \$a0	e200	1110001000000000
78	slt \$a1	slt \$a1	7300	0111001100000000
79	sacc \$t0	sacc \$t0	f800	1111100000000000
80	li \$acc1	sacc \$at0	fe00	1111111100000000
81	beqp \$t0 \$acc asub	lui b00000000	a000	1010000000000000
82		ori b00000001	d001	1101000000000001
83		sacc \$acc	f100	1111000100000000
84		sacc \$at0	fe00	1111111100000000
85		lui b00000000	a000	1010000000000000
86		ori b01100001	d061	1101000001100001
87		sacc \$at1	ff00	1111111110000000
88		lacc \$t0	e800	1110100000000000
89		beq \$at0 \$at1	4ef0	0100111011110000
90		lacc \$at0	ee00	1110111100000000
91	lacc \$a1	lacc \$a1	e300	1110001100000000
92	sub \$a0 1	sub \$a0	1201	0001001000000001
93	sacc \$a1	sacc \$a1	f300	1111001100000000
94	j loop	sacc \$at0	fe00	1111111100000000
95		lui loop		
96		ori loop		
97		sacc \$at1		
98		lacc \$at0		
99		jr \$at1		

70	loop:	beqp \$a1 \$0 acase	sacc \$at0	fe00	1111111000000000
71			lui b00000000	a000	1010000000000000
72			ori b01101001	d069	1101000001101001
73			sacc \$at1	ff00	1111111100000000
74			lacc \$a1	e300	1110001100000000
75			beq \$0 \$at1	40f0	0100000011110000
76			lacc \$at0	ee00	1110111000000000
112	acase	lacc \$a0	lacc \$a0	e200	1110001000000000
	:				
113		sacc \$v0	sacc \$v0	f600	1111011000000000
114		jr \$ra	jr \$ra	bc00	1011110000000000
29		lmem \$sp0	lmem \$sp0	2d00	0010110100000000
30		sacc \$ra	sacc \$ra	fc00	1111110000000000
31		lmem \$sp-1	lmem \$sp-1	2dff	0010110111111111
32		sacc \$a0	sacc \$a0	f200	1111001000000000
33		lmem \$sp-2	lmem \$sp-2	2dfe	0010110111111110
34		sacc \$s0	sacc \$s0	fa00	1111101000000000
35		addi \$sp \$sp-3	sacc \$at0	fe00	1111111000000000
36			lui b11111111	a0ff	1010000011111111
37			ori b11111111	d0ff	1101000011111111
38			add \$sp	0d00	0000110100000000
39			sacc \$sp	fd00	1111110100000000
40			lacc \$at0	ee00	1110111000000000
41		li \$acc 1	sacc \$at0	fe00	1111111000000000
42			lui b00000000	a000	1010000000000000
43			ori b00000001	d001	1101000000000001
44			sacc \$acc	f100	1111000100000000

45	beq \$v0 \$acc output	sacc \$at0	fe00	1111111000000000
46		lui b00000000	a000	1010000000000000
47		ori b00111010	d03a	1101000000111010
48		sacc \$at1	ff00	1111111100000000
49		lacc \$v0	e600	1110011000000000
50		beq \$at0 \$at1	4ef0	0100111011110000
51		lacc \$at0	ee00	1110111000000000
52	add \$s0	add \$s0	0a00	0000101000000000
53	sacc \$s0	sacc \$s0	fa00	1111101000000000
54	j loop	sacc \$at0	fe00	1111111000000000
55		lui b00000000	a000	1010000000000000
56		ori b00001000	d008	1101000000001000
57		sacc \$at1	ff00	1111111100000000
58		lacc \$at0	ee00	1110111000000000
59		jr \$at1	bf00	1011111100000000
9	loop: lacc \$s0	lacc \$s0	ea00	1110101000000000
10	sacc \$a1	sacc \$a1	f300	1111001100000000
11	addi \$sp \$sp3	sacc \$at0	fe00	1111111000000000
12		lui b00000000	a000	1010000000000000
13		ori b00000011	d003	1101000000000011
14		add \$sp	0d00	0000110100000000
15		sacc \$sp	fd00	1111110100000000
16		lacc \$at0	ee00	1110111000000000
17	lacc \$ra	lacc \$ra	ec00	1110110000000000
18	smem \$sp0	smem \$sp0	3d00	0011110100000000
19	lacc \$a0	lacc \$a0	e200	1110001000000000

20		smem \$sp-1	smem \$sp-1	3dff	0011110111111111
21		lacc \$s0	lacc \$s0	ea00	1110101000000000
22		smem \$sp-2	smem \$sp-2	3dfe	0011110111111110
23		jalp gcd	sacc \$at0	fe00	1111111000000000
24			lui b00000000	a000	1010000000000000
25			ori b00111100	d03c	1101000000111100
26			sacc \$at1	ff00	1111111100000000
27			lacc \$at0	ee00	1110111000000000
28			jal \$at1	cf00	1100111100000000
63	gcd:	beqp \$a0 \$0 bcase	sacc \$at0	fe00	1111111000000000
64			lui b00000000	a000	1010000000000000
65			ori b01011111	d05f	1101000001011111
66			sacc \$at1	ff00	1111111100000000
67			lacc \$a0	e200	1110001000000000
68			beq \$0 \$at1	40f0	0100000011110000
69			lacc \$at0	ee00	1110111000000000
70	loop:	beqp \$a1 \$0 acase	sacc \$at0	fe00	1111111000000000
71			lui b00000000	a000	1010000000000000
72			ori b01101001	d069	1101000001101001
73			sacc \$at1	ff00	1111111100000000
74			lacc \$a1	e300	1110001100000000
75			beq \$0 \$at1	40f0	0100000011110000
76			lacc \$at0	ee00	1110111000000000
77		lacc \$a0	lacc \$a0	e200	1110001000000000
78		slt \$a1	slt \$a1	7300	0111001100000000
79		sacc \$t0	sacc \$t0	f800	1111100000000000
80		li \$acc1	sacc \$at0	fe00	1111111000000000

81	beq \$t0 \$acc	asub	lui b00000000	a000	1010000000000000	
82			ori b00000001	d001	1101000000000001	
83			sacc \$acc	f100	1111000100000000	
84			sacc \$at0	fe00	1111111100000000	
85			lui b00000000	a000	1010000000000000	
86			ori b01100001	d061	1101000001100001	
87			sacc \$at1	ff00	1111111100000000	
88			lacc \$t0	e800	1110100000000000	
89			beq \$at0 \$at1	4ef0	0100111011110000	
90			lacc \$at0	ee00	1110111100000000	
103	asub:	lacc \$a0	lacc \$a0	e200	1110001000000000	
104		sub \$a1 1	sub \$a1	1301	0001001100000001	
105		sacc \$a0	sacc \$a0	f200	1111001000000000	
106		j loop	sacc \$at0	fe00	1111111100000000	
107			lui b00000000	a000	1010000000000000	
108			ori b01000010	d042	1101000001000010	
109			sacc \$at1	ff00	1111111100000000	
110			lacc \$at0	ee00	1110111100000000	
111			jr \$at1	bf00	1011111100000000	
70	loop:	beq \$a1 \$0	acase	sacc \$at0	fe00	1111111100000000
71			lui b00000000	a000	1010000000000000	
72			ori b01101001	d069	1101000001101001	
73			sacc \$at1	ff00	1111111100000000	
74			lacc \$a1	e300	1110001100000000	
75			beq \$0 \$at1	40f0	0100000011110000	
76			lacc \$at0	ee00	1110111100000000	
77		lacc \$a0	lacc \$a0	e200	1110001000000000	

78	slt \$a1	slt \$a1	7300	0111001100000000
79	sacc \$t0	sacc \$t0	f800	1111100000000000
80	li \$acc1	sacc \$at0	fe00	1111111000000000
81	beqp \$t0 \$acc asub	lui b00000000	a000	1010000000000000
82		ori b00000001	d001	1101000000000001
83		sacc \$acc	f100	1111000100000000
84		sacc \$at0	fe00	1111111000000000
85		lui b00000000	a000	1010000000000000
86		ori b01100001	d061	1101000001100001
87		sacc \$at1	ff00	1111111100000000
88		lacc \$t0	e800	1110100000000000
89		beq \$at0 \$at1	4ef0	0100111011110000
90		lacc \$at0	ee00	1110111000000000
91	lacc \$a1	lacc \$a1	e300	1110001100000000
92	sub \$a0 1	sub \$a0	1201	0001001000000001
93	sacc \$a1	sacc \$a1	f300	1111001100000000
94	j loop	sacc \$at0	fe00	1111111000000000
95		lui b00000000	a000	1010000000000000
96		ori b01000010	d042	1101000001000010
97		sacc \$at1	ff00	1111111100000000
98		lacc \$at0	ee00	1110111000000000
99		jr \$at1	bf00	1011111100000000
70	loop: beqp \$a1 \$0 acase	sacc \$at0	fe00	1111111000000000
71		lui b00000000	a000	1010000000000000
72		ori b01101001	d069	1101000001101001
73		sacc \$at1	ff00	1111111100000000
74		lacc \$a1	e300	1110001100000000

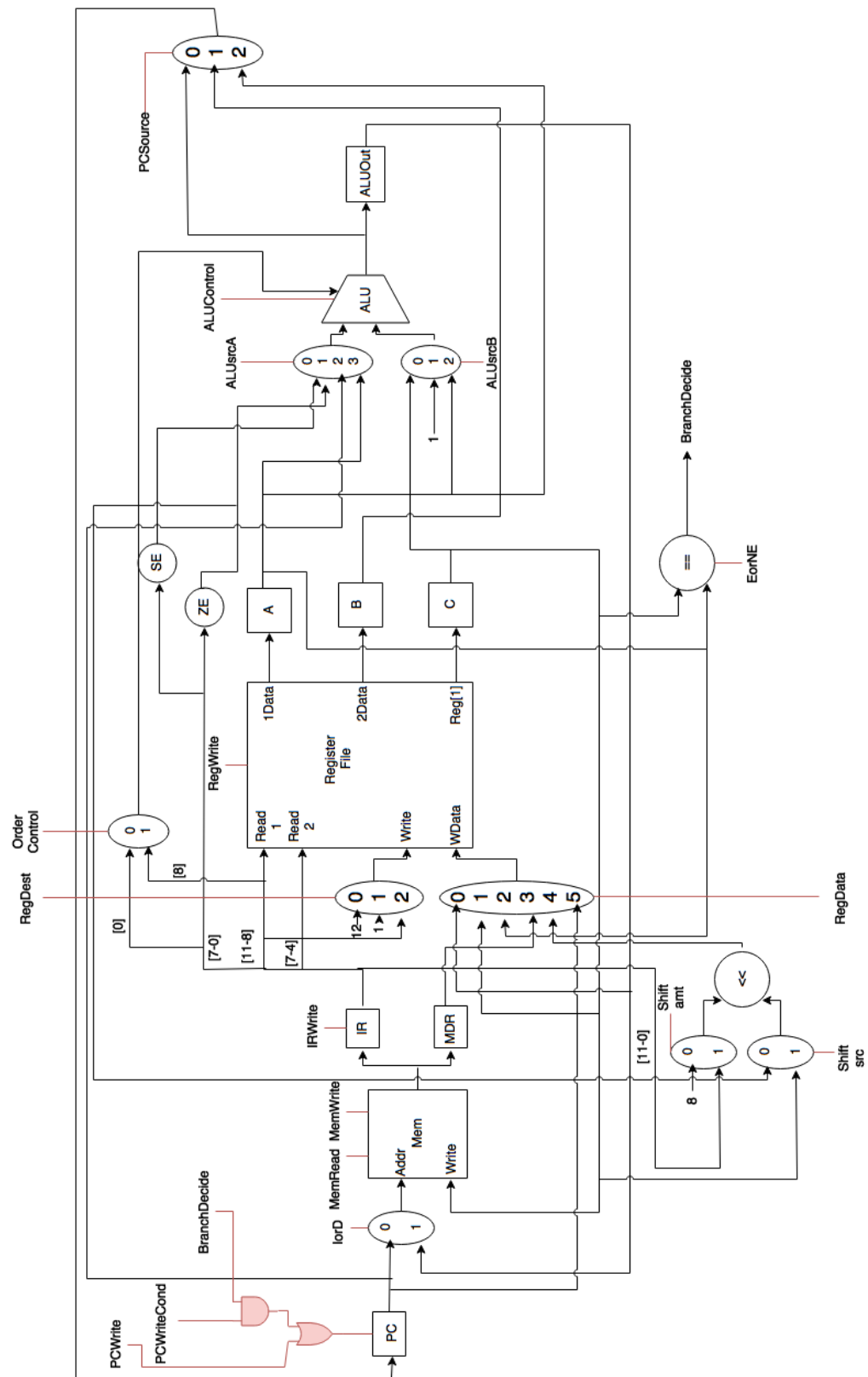
75		beq \$0 \$at1	40f0	0100000011110000	
76		lacc \$at0	ee00	1110111000000000	
77	lacc \$a0	lacc \$a0	e200	1110001000000000	
78	slt \$a1	slt \$a1	7300	0111001100000000	
79	sacc \$t0	sacc \$t0	f800	1111100000000000	
80	li \$acc1	sacc \$at0	fe00	1111111000000000	
81	beqp \$t0 \$acc asub	lui b00000000	a000	1010000000000000	
82		ori b00000001	d001	1101000000000001	
83		sacc \$acc	f100	1111000100000000	
84		sacc \$at0	fe00	1111111000000000	
85		lui b00000000	a000	1010000000000000	
86		ori b01100001	d061	1101000001100001	
87		sacc \$at1	ff00	1111111100000000	
88		lacc \$t0	e800	1110100000000000	
89		beq \$at0 \$at1	4ef0	0100111011110000	
90		lacc \$at0	ee00	1110111000000000	
91	lacc \$a1	lacc \$a1	e300	1110001100000000	
92	sub \$a0 1	sub \$a0	1201	0001001000000001	
93	sacc \$a1	sacc \$a1	f300	1111001100000000	
94	j loop	sacc \$at0	fe00	1111111000000000	
95		lui b00000000	a000	1010000000000000	
96		ori b01000010	d042	1101000001000010	
97		sacc \$at1	ff00	1111111100000000	
98		lacc \$at0	ee00	1110111000000000	
99		jr \$at1	bf00	1011111100000000	
70	loop:	beqp \$a1 \$0 acase	sacc \$at0	fe00	1111111000000000
71		lui b00000000	a000	1010000000000000	

72		ori b01101001	d069	1101000001101001
73		sacc \$at1	ff00	1111111100000000
74		lacc \$a1	e300	1110001100000000
75		beq \$0 \$at1	40f0	0100000011110000
76		lacc \$at0	ee00	1110111000000000
77	lacc \$a0	lacc \$a0	e200	1110001000000000
78	slt \$a1	slt \$a1	7300	0111001100000000
79	sacc \$t0	sacc \$t0	f800	1111100000000000
80	li \$acc1	sacc \$at0	fe00	1111111000000000
81	beqp \$t0 \$acc asub	lui b00000000	a000	1010000000000000
82		ori b00000001	d001	1101000000000001
83		sacc \$acc	f100	1111000100000000
84		sacc \$at0	fe00	1111111000000000
85		lui b00000000	a000	1010000000000000
86		ori b01100001	d061	1101000001100001
87		sacc \$at1	ff00	1111111100000000
88		lacc \$t0	e800	1110100000000000
89		beq \$at0 \$at1	4ef0	0100111011110000
90		lacc \$at0	ee00	1110111000000000
91	lacc \$a1	lacc \$a1	e300	1110001100000000
92	sub \$a0 1	sub \$a0	1201	0001001000000001
93	sacc \$a1	sacc \$a1	f300	1111001100000000
94	j loop	sacc \$at0	fe00	1111111000000000
95		lui b00000000	a000	1010000000000000
96		ori b01000010	d042	1101000001000010
97		sacc \$at1	ff00	1111111100000000
98		lacc \$at0	ee00	1110111000000000

99			jr \$at1	bf00	1011111100000000
70	loop:	beqp \$a1 \$0 acase	sacc \$at0	fe00	1111111100000000
71			lui b00000000	a000	1010000000000000
72			ori b01101001	d069	1101000001101001
73			sacc \$at1	ff00	1111111100000000
74			lacc \$a1	e300	1110001100000000
75			beq \$0 \$at1	40f0	0100000011110000
76			lacc \$at0	ee00	1110111000000000
112	acase	lacc \$a0	lacc \$a0	e200	1110001000000000
	:				
113		sacc \$v0	sacc \$v0	f600	1111011000000000
114		jr \$ra	jr \$ra	bc00	1011110000000000
29		lmem \$sp0	lmem \$sp0	2d00	0010110100000000
30		sacc \$ra	sacc \$ra	fc00	1111110000000000
31		lmem \$sp-1	lmem \$sp-1	2dff	0010110111111111
32		sacc \$a0	sacc \$a0	f200	1111001000000000
33		lmem \$sp-2	lmem \$sp-2	2dfe	0010110111111110
34		sacc \$s0	sacc \$s0	fa00	1111101000000000
35		addi \$sp \$sp-3	sacc \$at0	fe00	1111111000000000
36			lui b11111111	a0ff	1010000011111111
37			ori b11111111	d0ff	1101000011111111
38			add \$sp	0d00	0000110100000000
39			sacc \$sp	fd00	1111110100000000
40			lacc \$at0	ee00	1110111000000000
41		li \$acc1	sacc \$at0	fe00	1111111000000000
42			lui b00000000	a000	1010000000000000
43			ori b00000001	d001	1101000000000001

44		sacc \$acc	f100	1111000100000000
45	beqp \$v0 \$acc output	sacc \$at0	fe00	1111111100000000
46		lui b00000000	a000	1010000000000000
47		ori b00111010	d03a	1101000000111010
48		sacc \$at1	ff00	1111111110000000
49		lacc \$v0	e600	1110011000000000
50		beq \$at0 \$at1	4ef0	0100111011110000
51		lacc \$at0	ee00	1110111000000000
60	output : lacc \$s0	lacc \$s0	ea00	1110101000000000
61	sacc \$v0	sacc \$v0	f600	1111011000000000
62	jr \$ra	jr \$ra	bc00	1011110000000000

Datapath



Control

lorD (1 bit)

Value	Effect
0	The memory address is set to PC
1	The memory address is set to ALUOut

MemRead

Value	Effect
0	None
1	Data memory contents designated by address input are put on Read data output.

MemWrite

Value	Effect
0	None
1	Data memory contents designated by the address input are replaced by the value on the Write data input.

IRWrite

Value	Effect
0	None
1	Value in IR is replaced by the value readed from memory (which is current instruction).

RegDest

Value	Effect
00	The register destination number for the write register is set to 12.
01	The register destination number for the write register is set to 1.
10	The register destination number for the write register comes from the IR field (Bits [11:8]).
11	The register destination number for the write register is set to 3

RegData

Value	Effect
000	The data for writing into register is set to ALUOut.
001	The data for writing into register is set to value in C (\$Accumulator).
010	The data for writing into register is set to value in A.
011	The data for writing into register is set to value in MDR.
100	The data for writing into register is set to value after shift
101	The data for writing into register is set to value PC
110	The data for writing into register is set to value in Input

Shiftamt

Value	Effect
0	Shift offset is set to 8
1	Shift offset is set to value in IR (bits [11:0])

Shiftsrc

Value	Effect
0	Value for shiftment is set to value in IR (Zero extended bits[7:0]).
1	Value for shiftment is set to value in C.

RegWrite

Value	Effect
0	None
1	The register on the Write register input is written with the value on the Write data input.

ALUSrcA

Value	Effect
00	Value inputting into first source of ALU is set to value in IR (Sign extended bits[7:0]).
01	Value inputting into first source of ALU is ZE(IR[7-0])
10	Value inputting into first source of ALU is set to value in PC
11	Value inputting into first source of ALU is set to the value in A

ALUSrcB

Value	Effect
00	Value inputting into second source of ALU is set to value in C
01	Value inputting into second source of ALU is set to 1
10	Value inputting into second source of ALU is set to value in A

ALUControl

Value	Effect
000	The operation performed in ALU will be and The output of the ALU will be ALUSrcA and ALUSrcB.
001	The operation performed in ALU will be or.

	The output of the ALU will be ALUSrcA or ALUSrcB.
010	The operation performed in ALU will be add The output of the ALU will be ALUSrcA add ALUSrcB.
110	The operation performed in ALU will be sub The output of the ALU will be ALUSrcA sub ALUSrcB.
111	The operation performed in ALU will be slt The output of the ALU will be ALUSrcA slt by ALUSrcB bits.

PCSource

Value	Effect
00	Output value of mux is set to value come from ALU
01	Output value of mux is set to value in B
10	Output value of mux is set to value in A

PCWriteCond

Value	Effect
0	PC will only update if PCWrite is 1.
1	If ALUOut is zero (or say, zero is 1), PC's value is updated

PCWrite

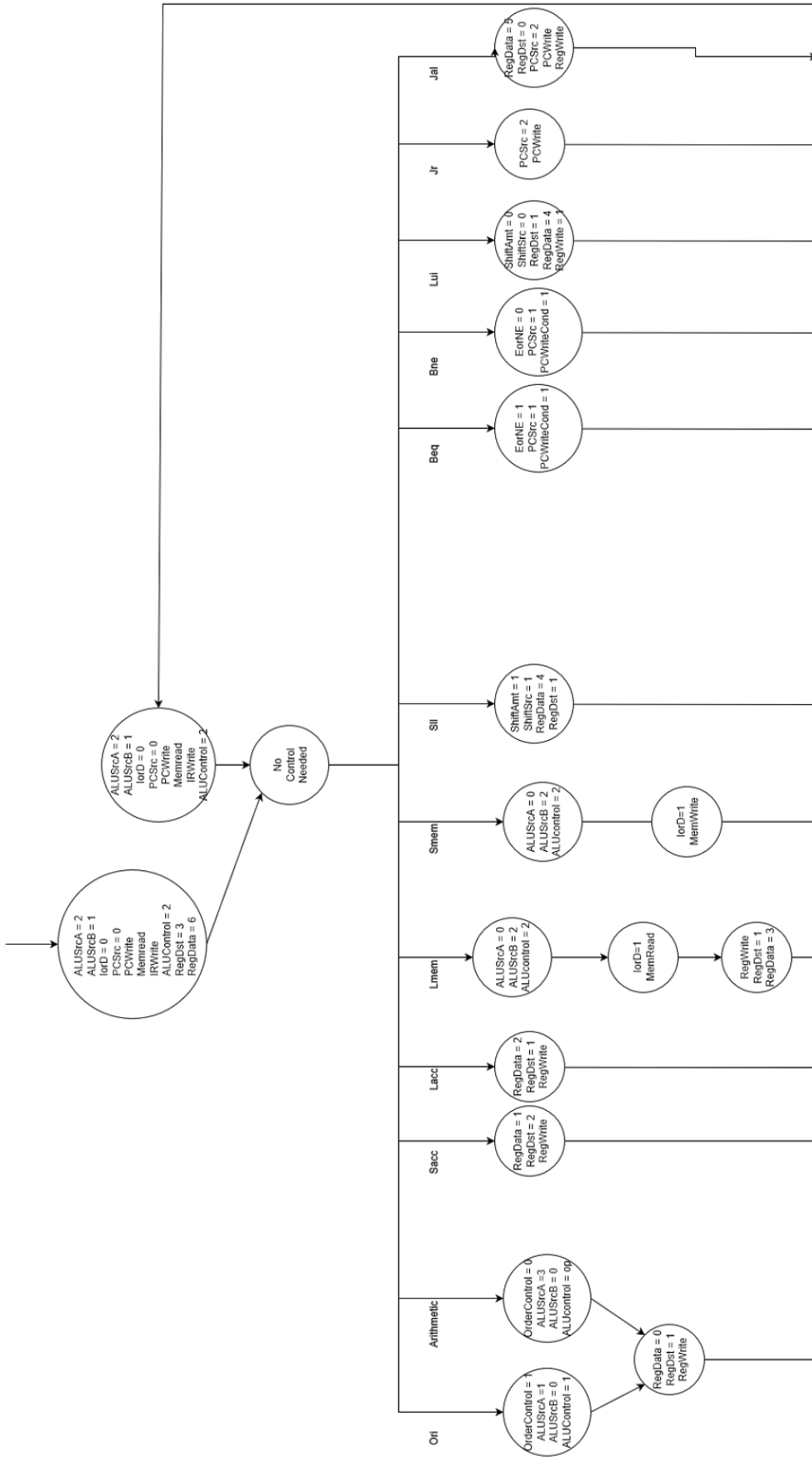
Value	Effect
0	PC will only update if the value of PCWriteCond and zero are both 1.
1	PC's value is updated

EorNE

Value	Effect
0	Value of BranchDecide is 1 if two inputs to the comparator are not equal. If two inputs are equal, value of BranchDecide is set to 0.
1	Value of BranchDecide is 0 if two inputs to the comparator are not equal. If two inputs are equal, value of BranchDecide is set to 1.

OrderControl

Value	Effect
0	Order bits for ALU are IR[7-0]
1	Order bits for ALU are ZE(IR[11-8])



Control Test Plan

Cycle One

Control Signal	What it ends up doing
ALUSrcA = 2	Puts PC into the ALU
ALUSrcB = 1	Puts 1 into the ALU
PCWrite = 1	Allows a value to be written into PC
MemRead = 1	Allows a value to be read from memory
IRWrite = 1	Allows an instruction to be written to IR
PCSrc = 0	Sets the value going to PC the output of the ALU
ALUControl = 2	Allows 1 to be added to the PC in the ALU

In combination, these signals add one to PC and get the instruction at PC from memory. This cycle's control works.

Cycle Two

No control signals needed for cycle two.

Cycle Three

Ori

Control Signal	What it ends up doing
OrderControl = 1	Makes the order bits from the correct part of the instruction
ALUSrcA = 1	Sets the top input to the ALU the output of AorO
ALUSrcB = 0	Sets the bottom input to the ALU to C

This combination of these control signals gets C and the immediate from the instruction and ors them using the correct order(not that it really matters).

Arithmetic/Slt

Control Signal	What it actually does
OrderControl = 0	Sets the order bits to the last 8 bits of the instruction, which is where the order bits are for every arithmetic instruction.
ALUSrcA = 1	Sets the top input of the ALU to A
ALUSrcB = 0	Sets the bottom input of the ALU to C

This combination of control signals is very similar to the ones for ori, but it chooses A and a different set of order bits instead of an immediate and a different set of order bits.

Sacc

Control Signals	What they actually do
RegData = 1	Sets the value to put into the register file the value in C
RegDst = 2	Sets the register to write to the register in bits 11-8 of the instruction
RegWrite = 1	Allows for the register file to be written to

This combination of control signals sort of bypasses the ALU and writes directly to another register. It chooses to put C into whatever register is specified by the instructions

Lacc

Control Signals	What they actually do
RegData = 2	Sets the write data from the register file to the data in A
RegDst = 1	Sets the register that will be written to to the accumulator
RegWrite = 1	Allows for the register file to be written to

Exactly the same as sacc but register gets written to acc instead of the other way around. The changes in the control bits are what allow for this.

Lmem

Control Signal	What it actually does
ALUSrcA = 0	Same as above
ALUSrcB = 2	Puts A as the bottom input to the ALU

This set of control signals adds the immediate to the register specified.

Smem

Control Signal	What it actually does
AorO = 1	Puts the imm as the top input to the ALU
ALUSrcA = 0	Same as above
ALUSrcB = 2	Puts A as the bottom input to the ALU
IorD = 1	Makes memory location the value in ALUOut
MemWrite = 1	Allows for memory to be written to

This set of control signals does the exact same thing as Imem, but instead of reading from memory, it writes the accumulator value to memory at the calculated address.

Sll

Control Signal	What it actually does
ShiftAmt = 1	Sets the shift amount to the shifter to whatever is in the 11-0 imm
ShiftSrc = 1	Sets the thing to shift to C
RegData = 4	Sets the reg write data to the output of the shifter
RegDst = 1	Sets the register to write to to the accumulator

Operates with the shifter to shift the accumulator a certain amount

Beq

Control Signal	What it actually does
EorNE = 1	Sets the comparator to perform equals
PCSrc = 1	Sets the PC to value in B
PCWriteCond = 1	Makes sure the PC can be written to only if the condition is right

This puts the correct inputs into the comparator and jumps if correct

Bne

Control Signal	What it actually does
EorNE = 0	Sets the comparator to perform not equals
PCSrc = 1	Sets the PC to value in B
PCWriteCond = 1	Makes sure the PC can be written to only if the condition is right

This puts the correct inputs into the comparator and jumps if correct

Lui

Control Signal	What it actually does
ShiftAmt = 0	Sets the shift amount to the shifter to 8
ShiftSrc = 1	Sets the thing to shift to ZE(order)
RegData = 4	Sets the reg write data to the output of the shifter
RegDst = 1	Sets the register to write to to the accumulator

Very similar to sll, except it shifts an imm left 8 every time.

Jr

Control Signal	What it actually does
PCSrc = 2	Sets the value to put in PC to A
PCWrite = 1	Allows the PC to be written to

This basically just puts A into PC

Jal

Control Signal	What it actually does
RegData = 5	Sets the info to put in the register file to PC
RegDst = 0	Sets the register to write to to \$ra
RegWrite = 1	Allows \$ra to be written to

PCWrite = 1	Allows PC to be written to
PCSrc = 2	Sets PC to the jump target

This writes the PC into \$ra. Basically just sets the controls for the Reg input stuff

Cycle 4

Ori/Arithmetic/Slt

RegData = 0	Sets the data that you are writing to the register to ALUOut
RegDst = 1	Sets the register you are writing to the accumulator
RegWrite = 1	Allows you to write to the register file

Writes the value in ALUOut to the accumulator

Lmem

Control Signal	What is actually does
lorD = 1	Makes memory location the value in ALUOut
MemRead = 1	Allows the memory to be read from

This writes to MDR

Smem

lorD = 1	Makes memory location the value in ALUOut
MemWrite = 1	Allows for memory to be read from

Cycle 5

Imem

Control Signal	What it actually does
RegWrite	Allows the register to be written to
RegDst = 1	Set the accumulator as the register to be written to.
RegData = 3	Writes MDR to the selected register.

Performance

While we test our datapath with 0x13B0 (5040 in decimal) as input, we get these results.

Total used memory: 274bytes

Total instructions: 316509

Total cycles: 949662

Average cycles per instruction: 3.0004265

Cycle time: 14.794ns

Total execution time: 15.194744ms

Gate count: 1800

Device utilization summary:

Selected Device : 3sd1800afg676-4

Number of Slices:	398	out of	16640	2%
Number of Slice Flip Flops:	374	out of	33280	1%
Number of 4 input LUTs:	796	out of	33280	2%
Number of IOs:	147			
Number of bonded IOBs:	147	out of	519	28%
Number of BRAMs:	1	out of	84	1%
Number of GCLKs:	1	out of	24	4%