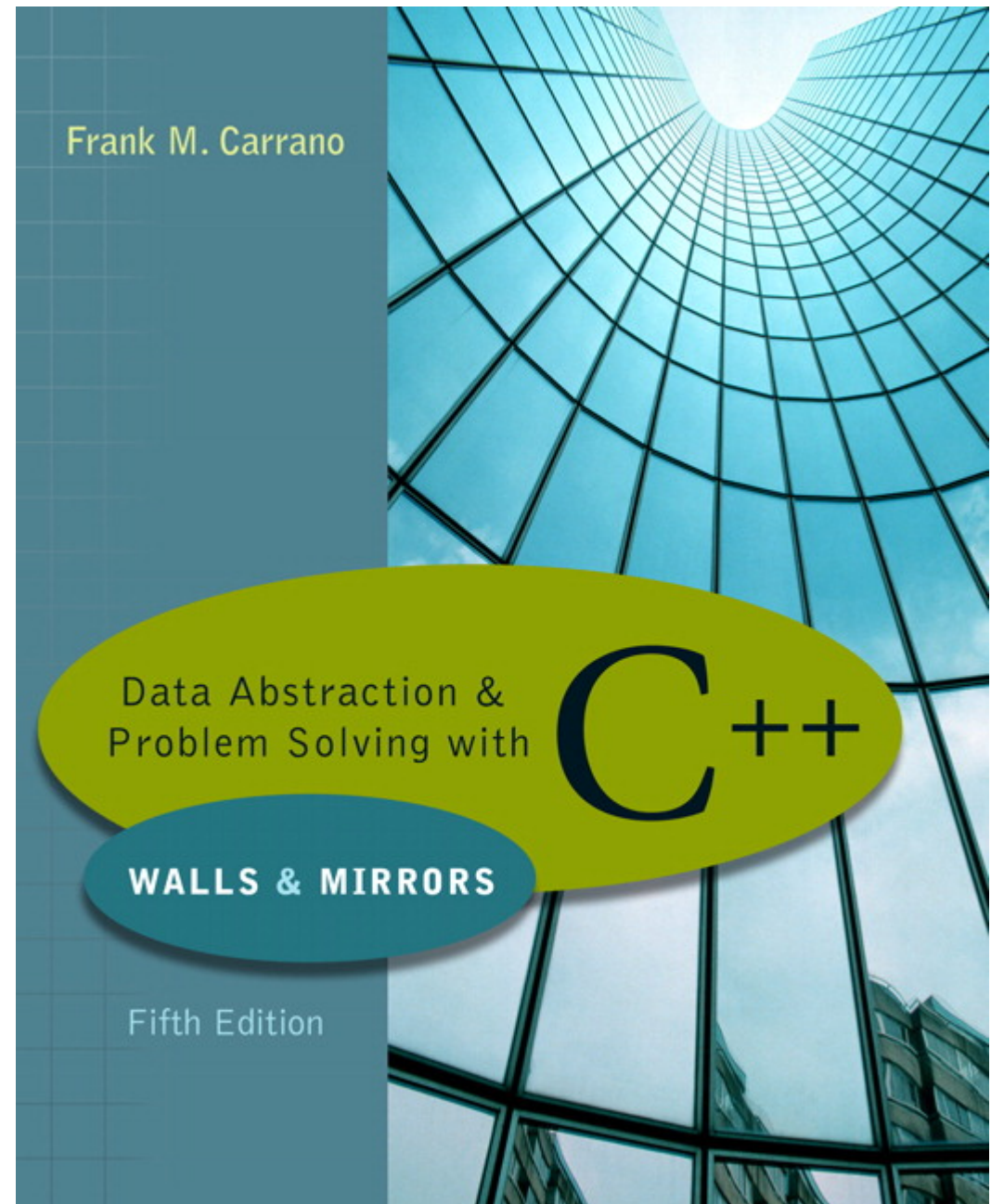


CHAPTER 10

Trees





What have we learned so far?

- ADT list, stack, queue
- General data management operations
 - Insert
 - Delete
 - Query: is empty? Length? retrieve.
- Operations
 - Position-oriented
 - Value-oriented

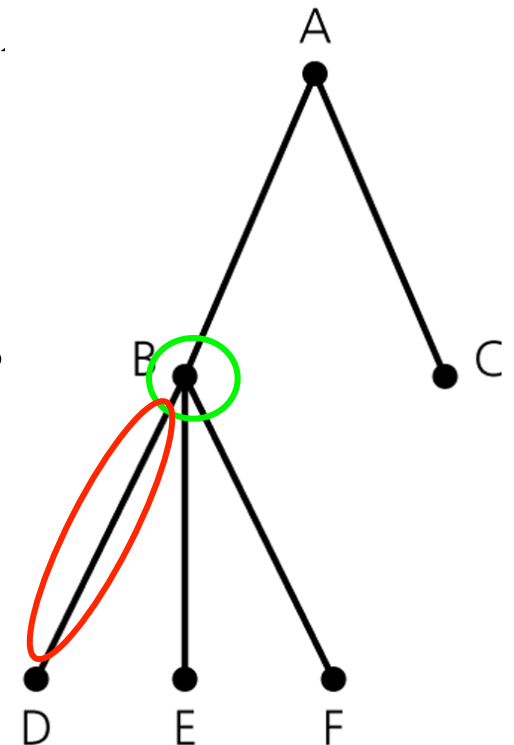
What is coming

- ❑ Trees
- ❑ Terminologies
- ❑ Binary search tree
- ❑ Operations



Terminology

- Tree: A connected, undirected graph without cycles
 - The process of arranging the vertices into a topological order
 - graph: a set of vertices and a set of edges
 - vertices
- Vertex: a node in a graph
- Edge: line (connection) between nodes



Terminology

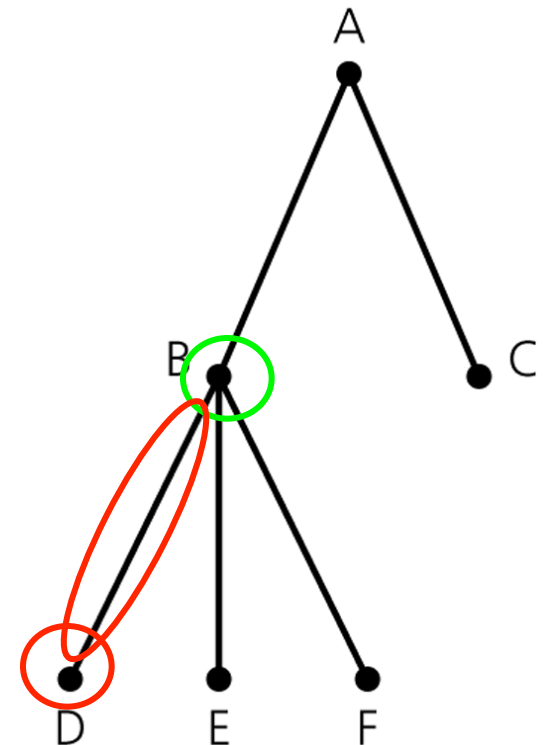
- Tree: all trees are hierarchical in nature
 - Hierarchical: parent/child relationship between nodes in the tree

- Parent / child

If an edge is between two nodes

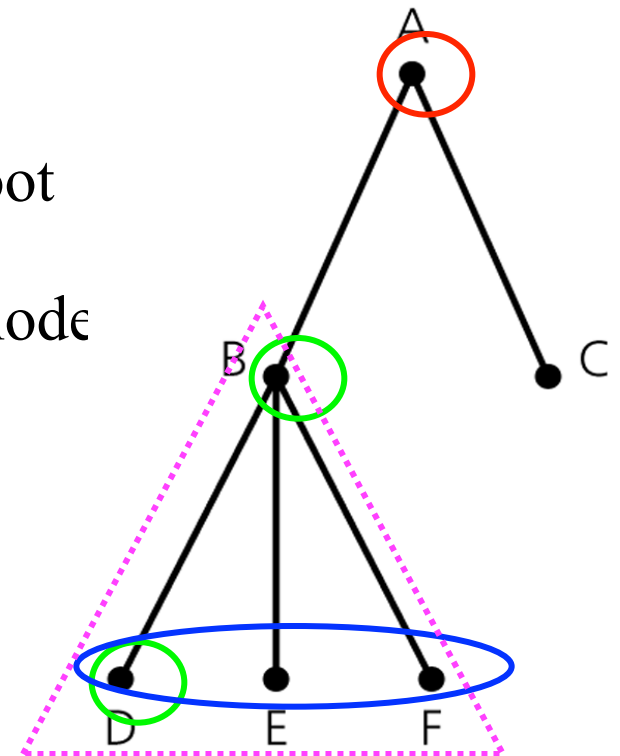
B and D, B is above D in the tree

B is the parent of D, D is a child of B



Terminology continued ...

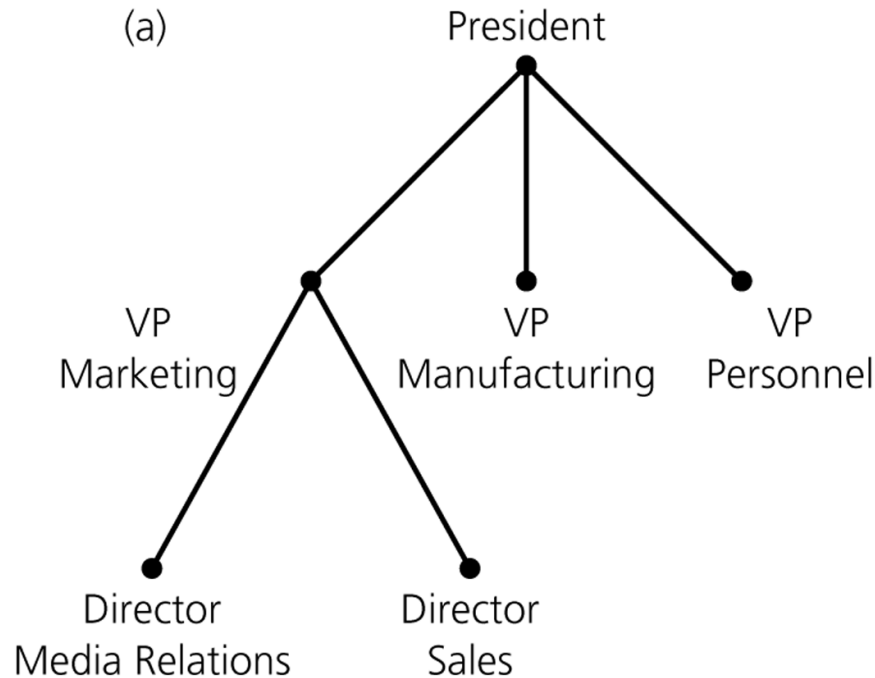
- ❑ Sibling: children of the same parent (DEF)
- ❑ Root: node with no parent (one node in a tree)
- ❑ ancestor / descendant
 - ancestor: a node on the path from the root of a tree to the node
 - descendant: a node on a path from the node to a leaf of a tree
- ❑ Subtree : any node in a tree, together with all of the node's descendants



■ Represent information that is hierarchical in nature

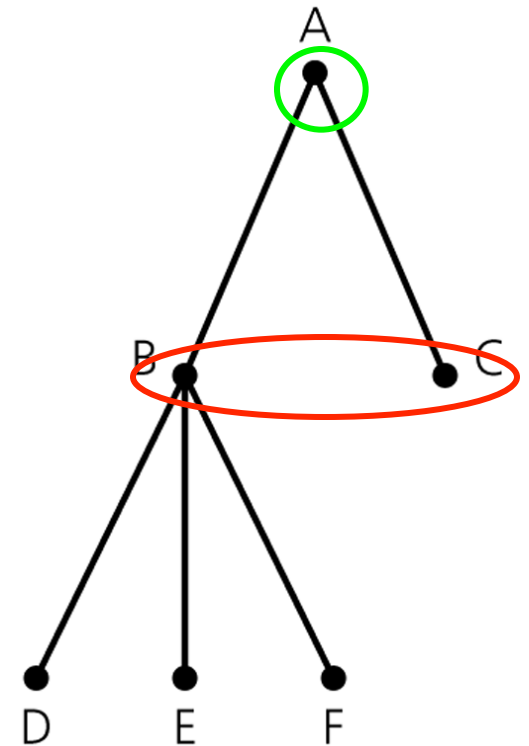
(a) An organization chart;

(b) a family tree



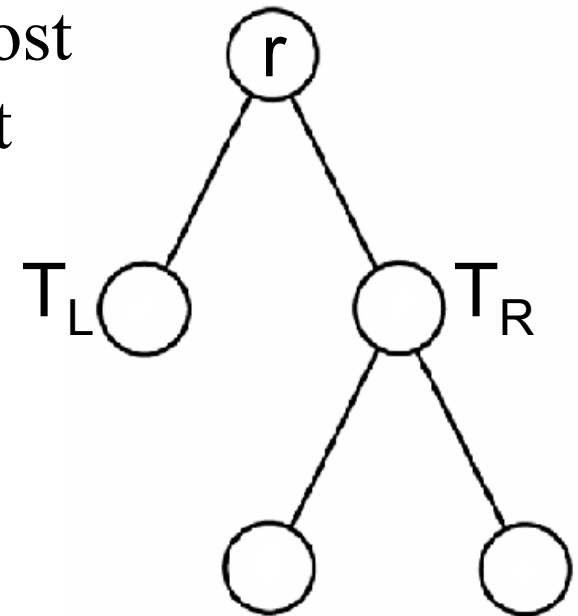
A general tree (T)

- A set of one or more nodes such that T is partitioned into disjoint subsets
 - A single node r, the root
 - Sets that are general trees, called subtrees



A binary tree

- A set of zero or more nodes, partitioned into a root node and two possibly empty sets that are binary trees.
- Each node in a binary tree has at most two children, the left child and right child
- T_L : left subtree
- T_R : right subtree



An application of binary tree:

Binary trees that represent algebraic expressions

$$a - b$$

$$a - b / c$$

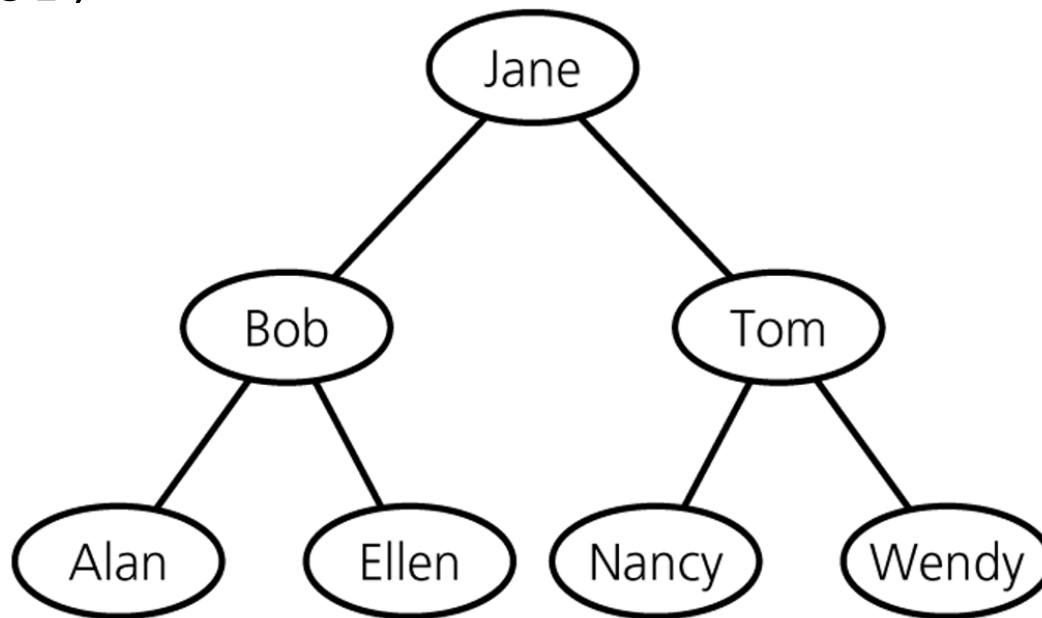
$$(a - b) * c$$

Leaves of these trees contain the expressions operands.
Other tree nodes contain the operators.

An application of binary tree:

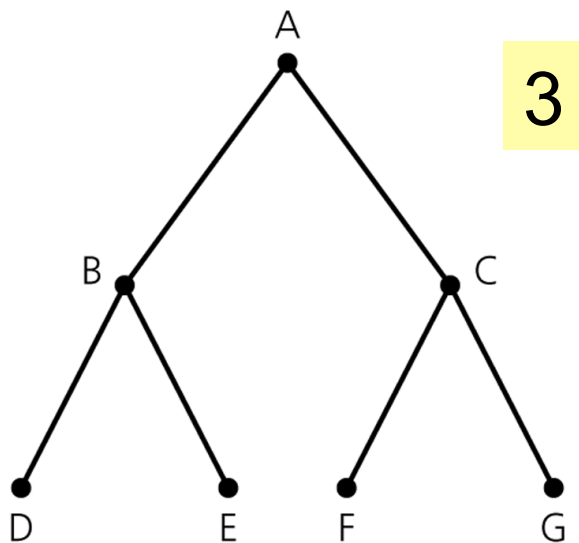
A binary search tree (BST)

- A binary tree where the search key in any node n is
 - greater than the search key in its left subtree (BST),
 - but less than the search key in any node in its right subtree (BST)

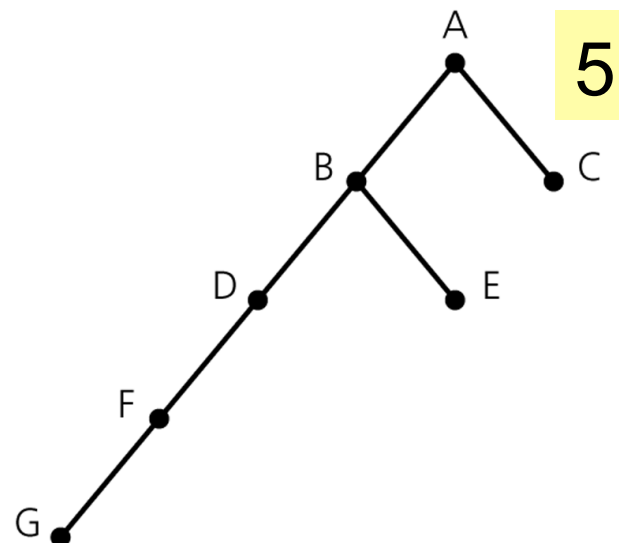


The height of trees

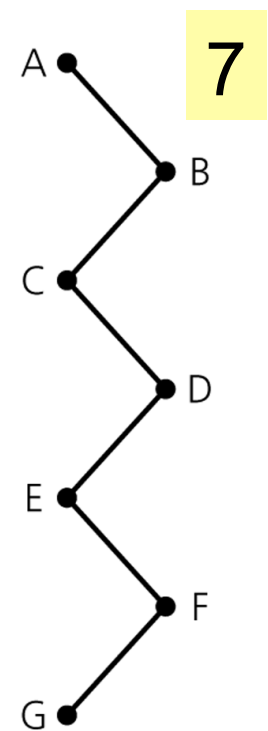
- The number of nodes on the longest path from the root to a leaf.



(a)



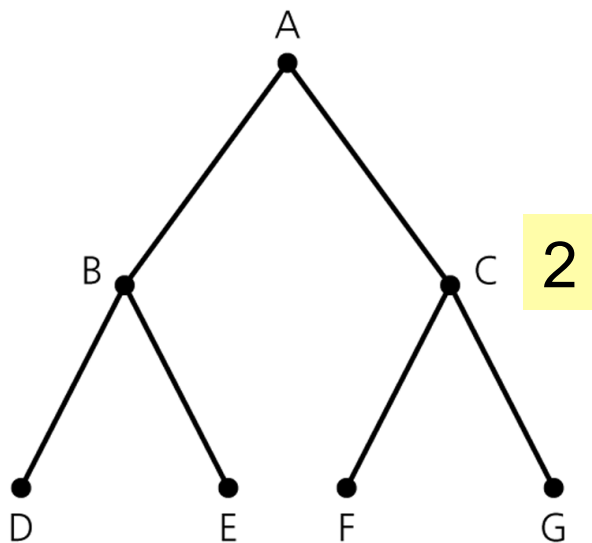
(b)



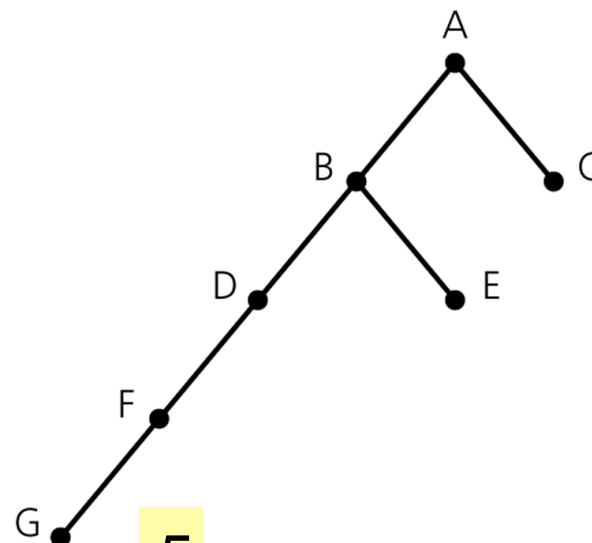
(c)

The level of a node n

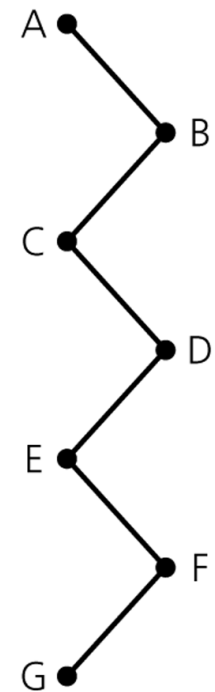
- If n is the root of T , it is at level 1
- If n is not the root of T , its level is 1 greater than the level of its parent



(a)



(b)



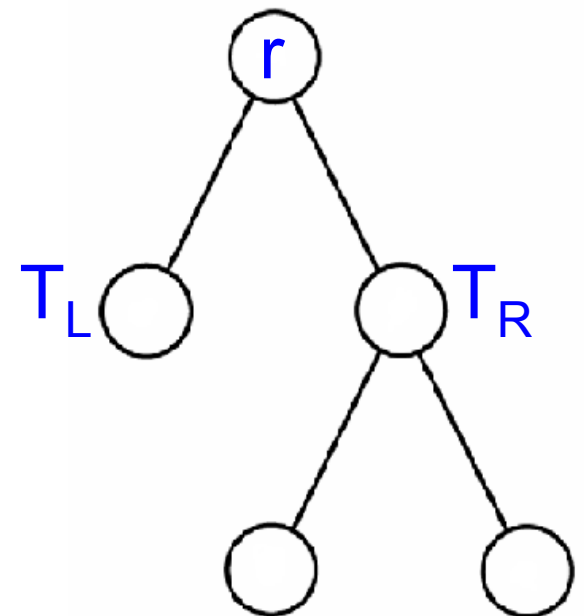
(c)

Height in a binary tree

- If T is empty, its height is 0
- If T is a nonempty binary tree,

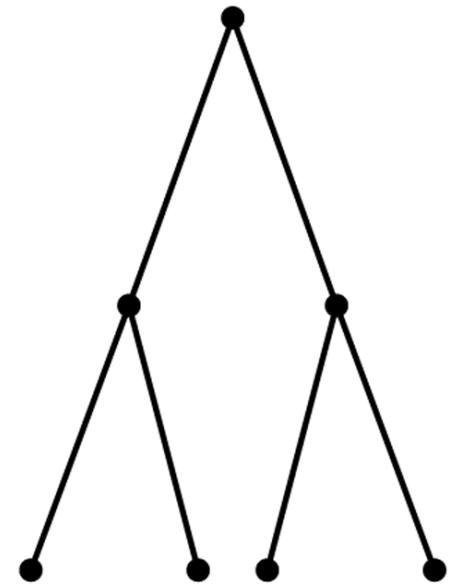
$$\text{Height}(T) = 1 + \max(\text{height}(T_L), \text{height}(T_R))$$

- T_L : left subtree
- T_R : right subtree



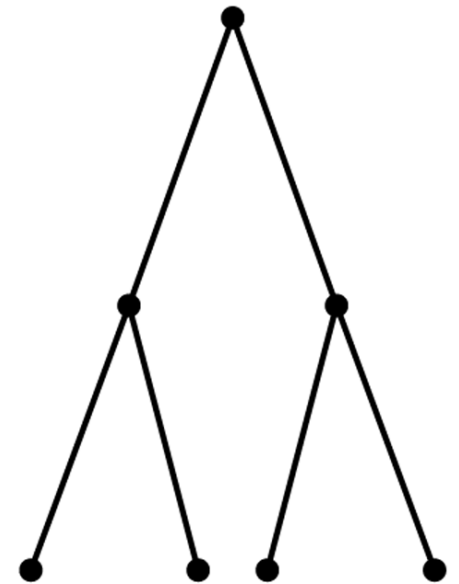
Full binary tree

- Full binary tree: all nodes that are at a level less than the height, h , has two children each
 - If T is empty, T is a full binary tree of height 0
 - If T is not empty and has height h , T is a full binary tree if its root's subtrees are both full binary tree of height $h-1$



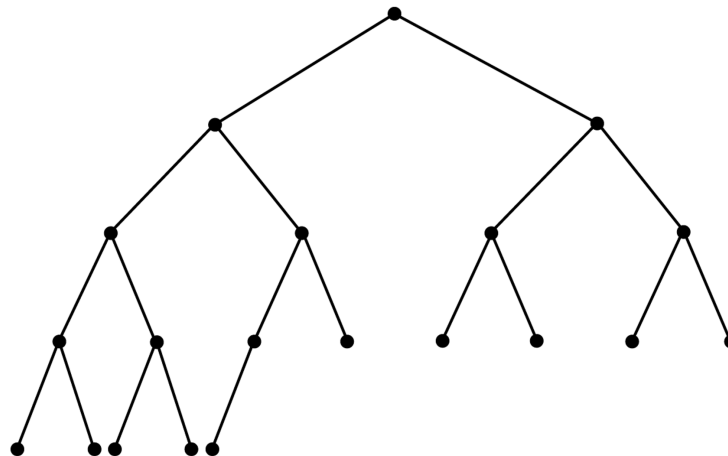
Full binary tree

- A full binary tree of height h has
 - $2^h - 1$ nodes
 - 2^{h-1} terminal nodes (leaves)
- A full binary tree of n nodes
 - $\lceil \log_2 n \rceil$ height



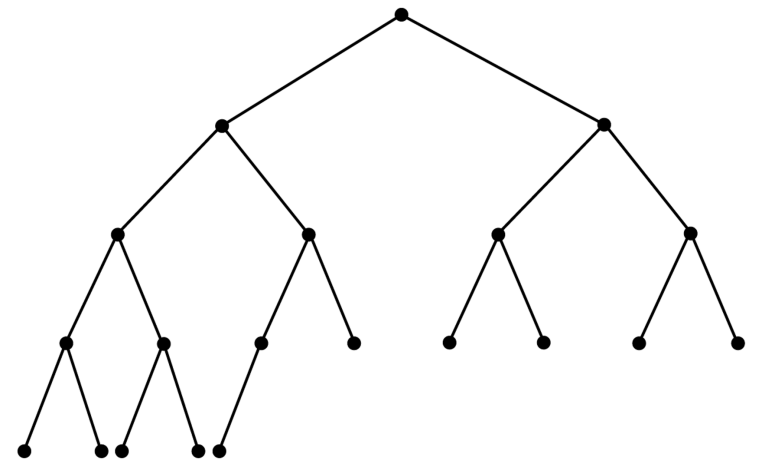
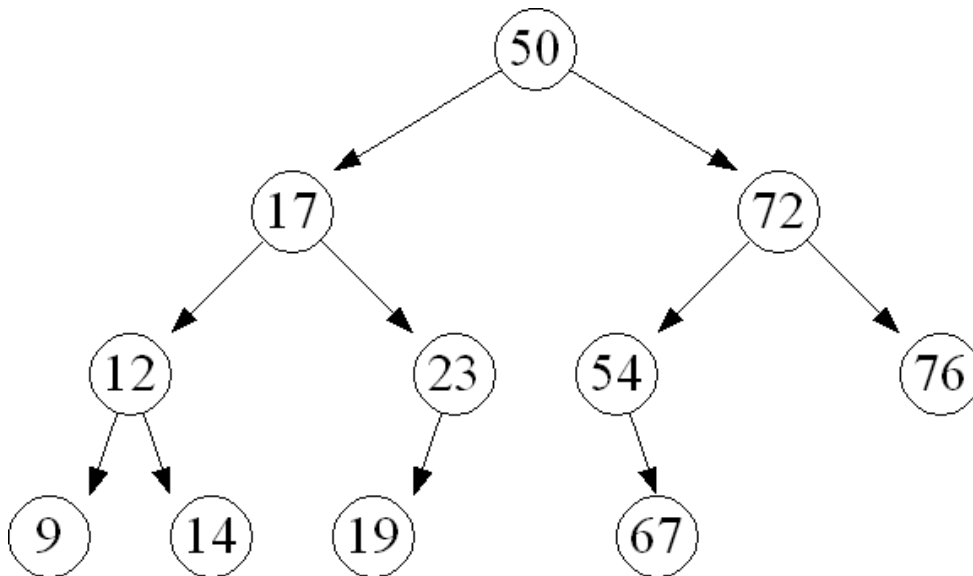
Complete binary tree

- Complete binary tree: full down to level $h-1$, with level h filled in from left to right
 - All nodes at level $h-2$ and above has two children each,
 - When a node at level $h-1$ has children, all nodes to its left at the same level have two children each
 - When a node at level $h-1$ has one child, it is left child

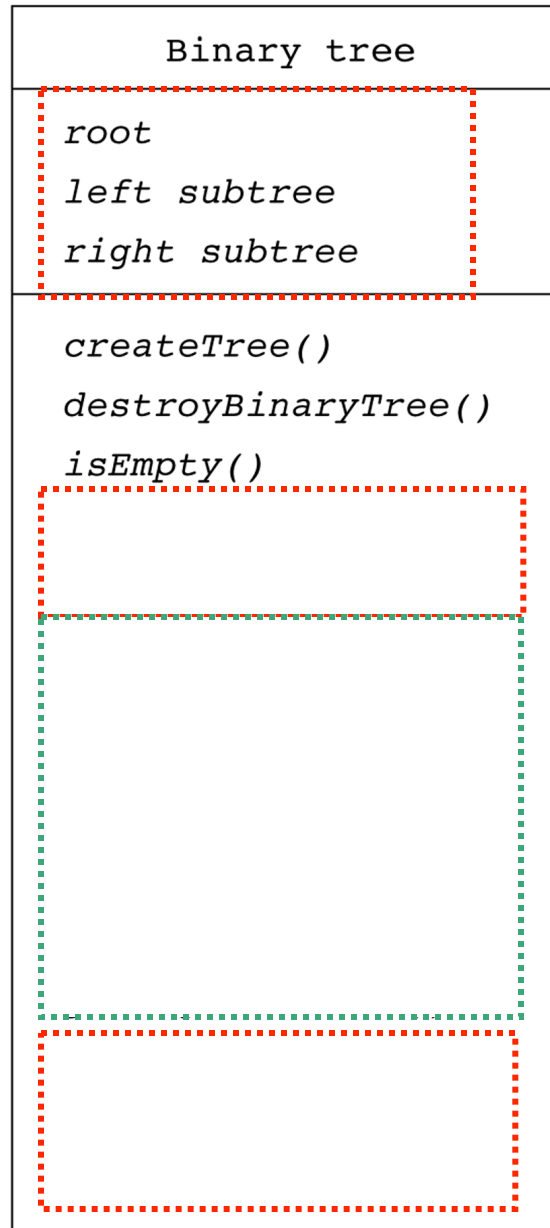


■ AVL tree: balanced binary search tree

- A binary tree is height balanced:
 - If the height of any node's right subtree differs from the height of the node's left subtree by no more than 1
- The AVL tree is named after its two inventors, G.M. Adelson-Velsky and E.M. Landis



UML diagram for the class *BinaryTree*



ADT binary tree operations

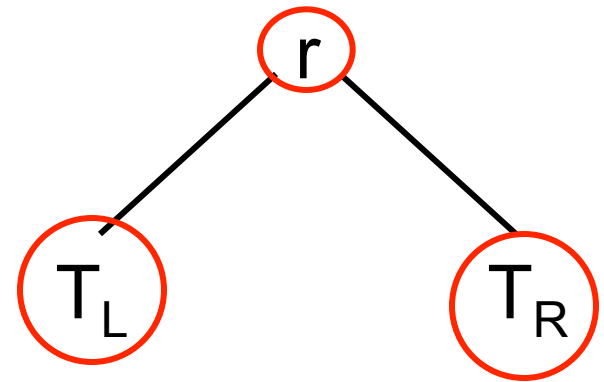
- Create an empty binary tree
- Create a one-node binary tree given an item
- Create a binary tree given an item for its root and two binary tree for the root's subtrees
- Destroy a binary tree
- Determine whether a binary tree is empty

ADT binary tree operations

- Determine or change the data in the binary tree's root
- Attach a left or right child to the binary tree's root
- Attach a left or right subtree to the binary tree's root
- Detach the left or right subtree of the binary tree's root
- Return a copy of the left or right subtree of binary tree's root
- Traverse the nodes in a binary tree in preorder, inorder, or postorder

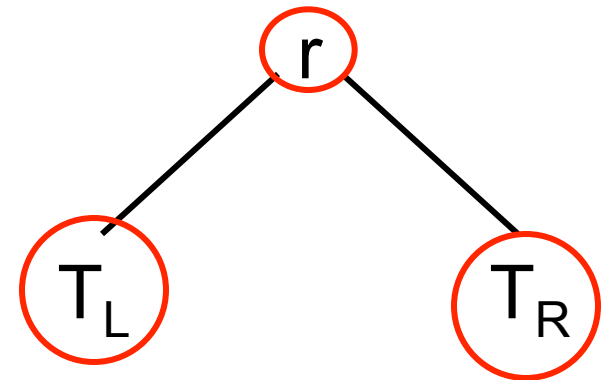
Binary tree traversal

□ Preorder : $r \rightarrow T_L \rightarrow T_R$



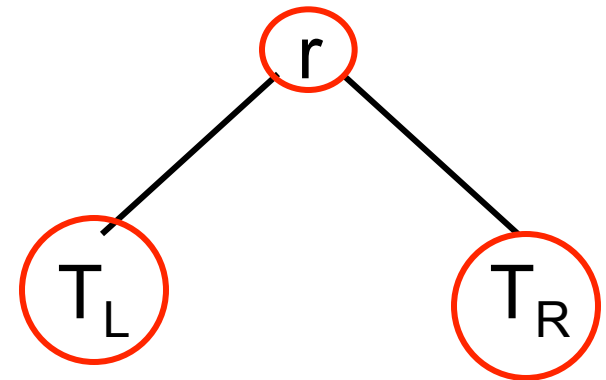
Binary tree traversal

- Preorder : $r \rightarrow T_L \rightarrow T_R$
- inorder : $T_L \rightarrow r \rightarrow T_R$



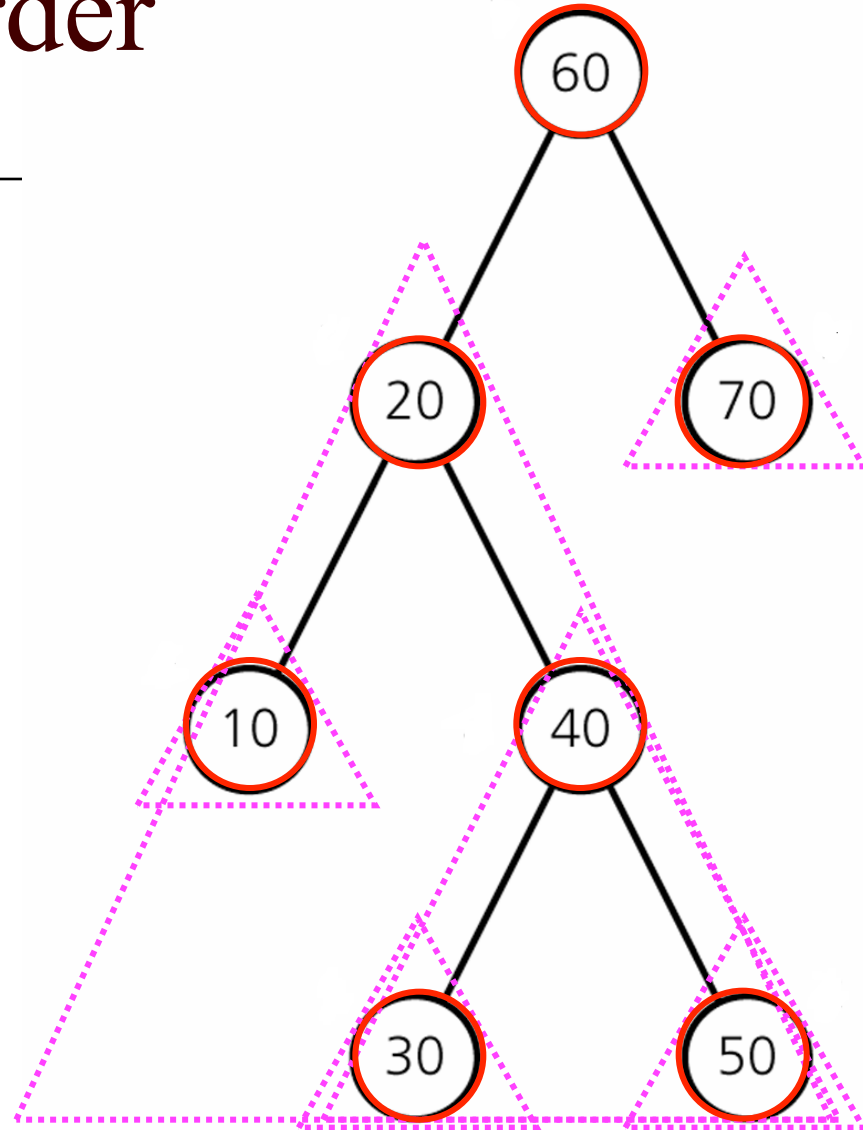
Binary tree traversal

- Preorder : $r \rightarrow T_L \rightarrow T_R$
- inorder : $T_L \rightarrow r \rightarrow T_R$
- Postorder : $T_L \rightarrow T_R \rightarrow r$



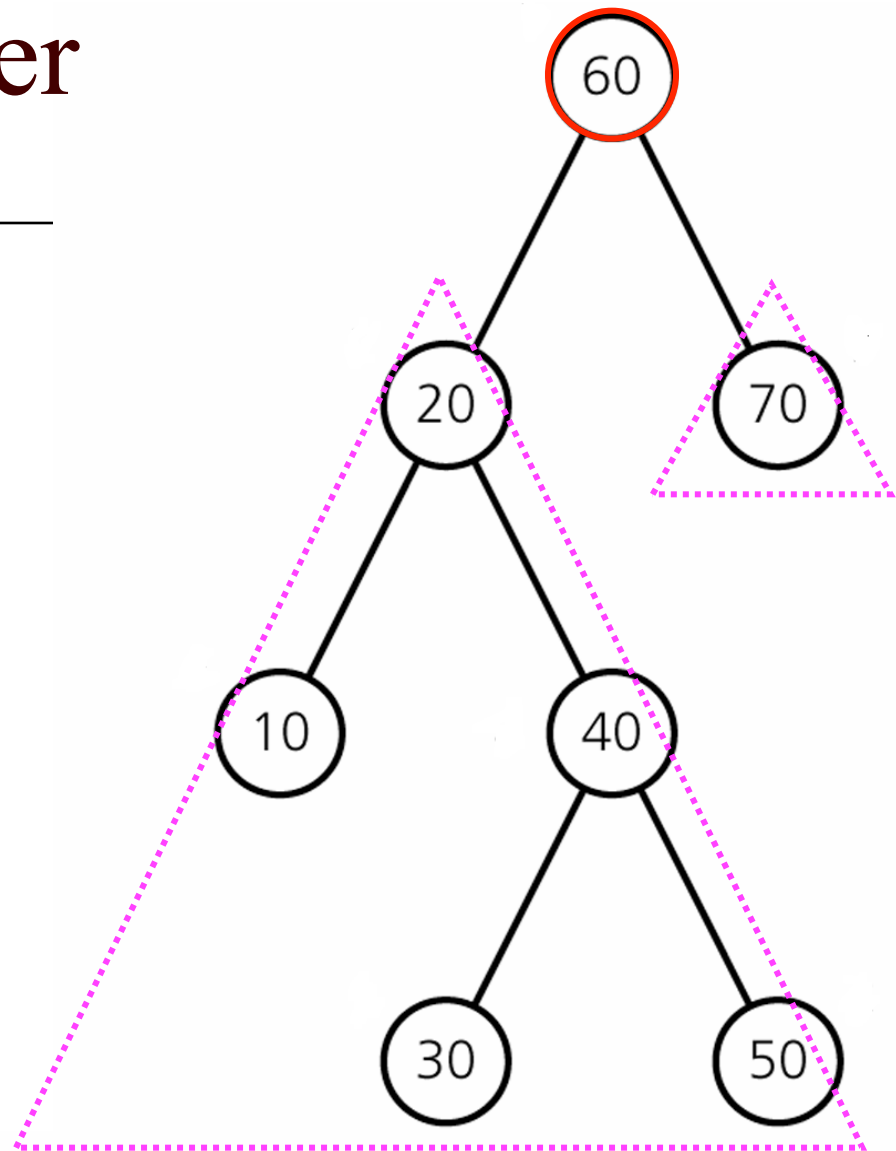
Traversal: preorder

1. Root
 2. Left child
 3. Right child
- Recursively!!



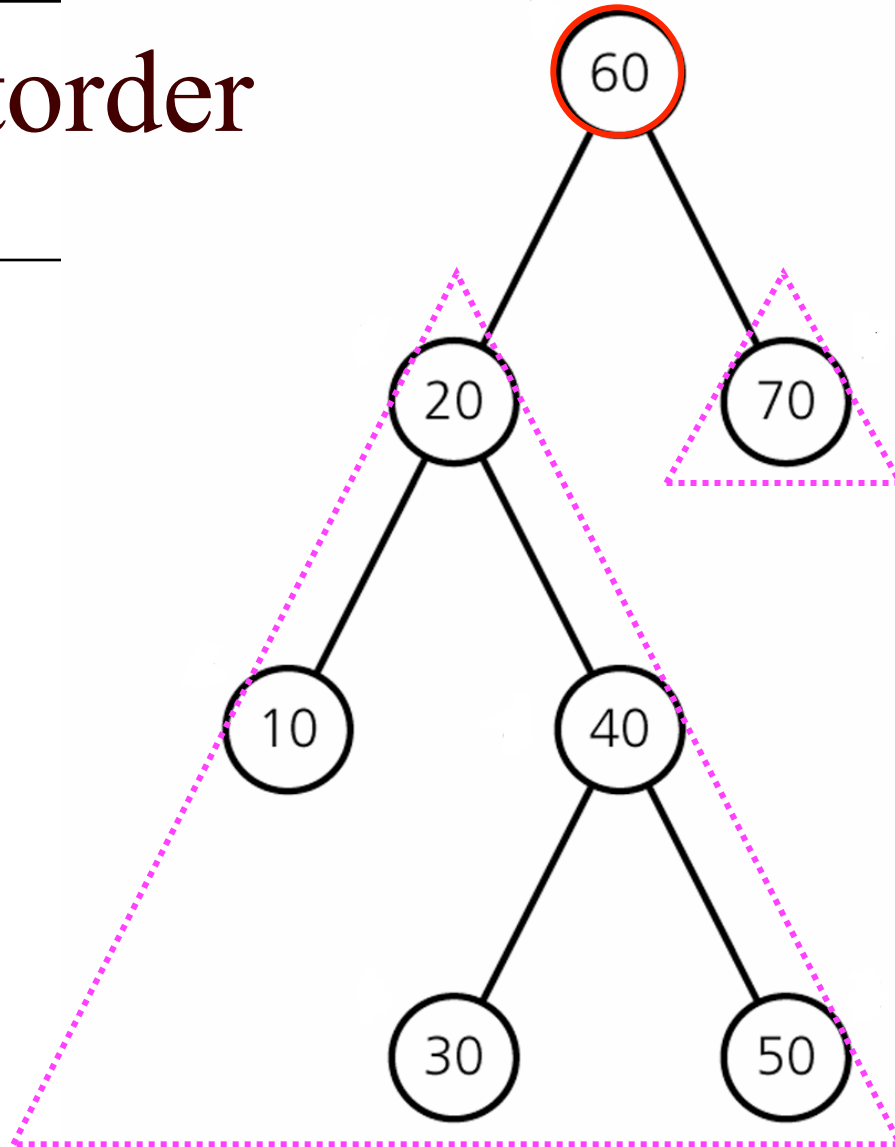
Traversal: inorder

1. Left child
 2. Root
 3. Right child
- Recursively!!



Traversal: postorder

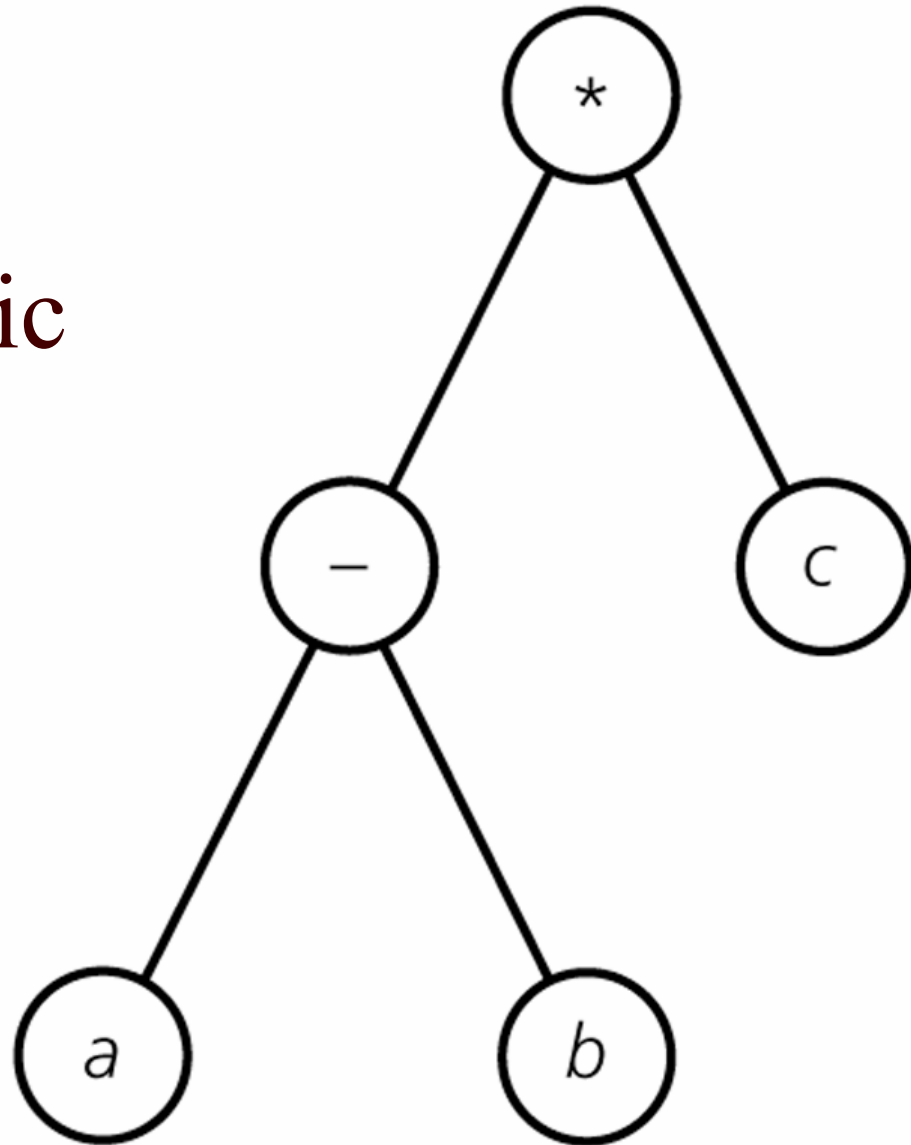
1. Left child
 2. Right child
 3. Root
- Recursively!!



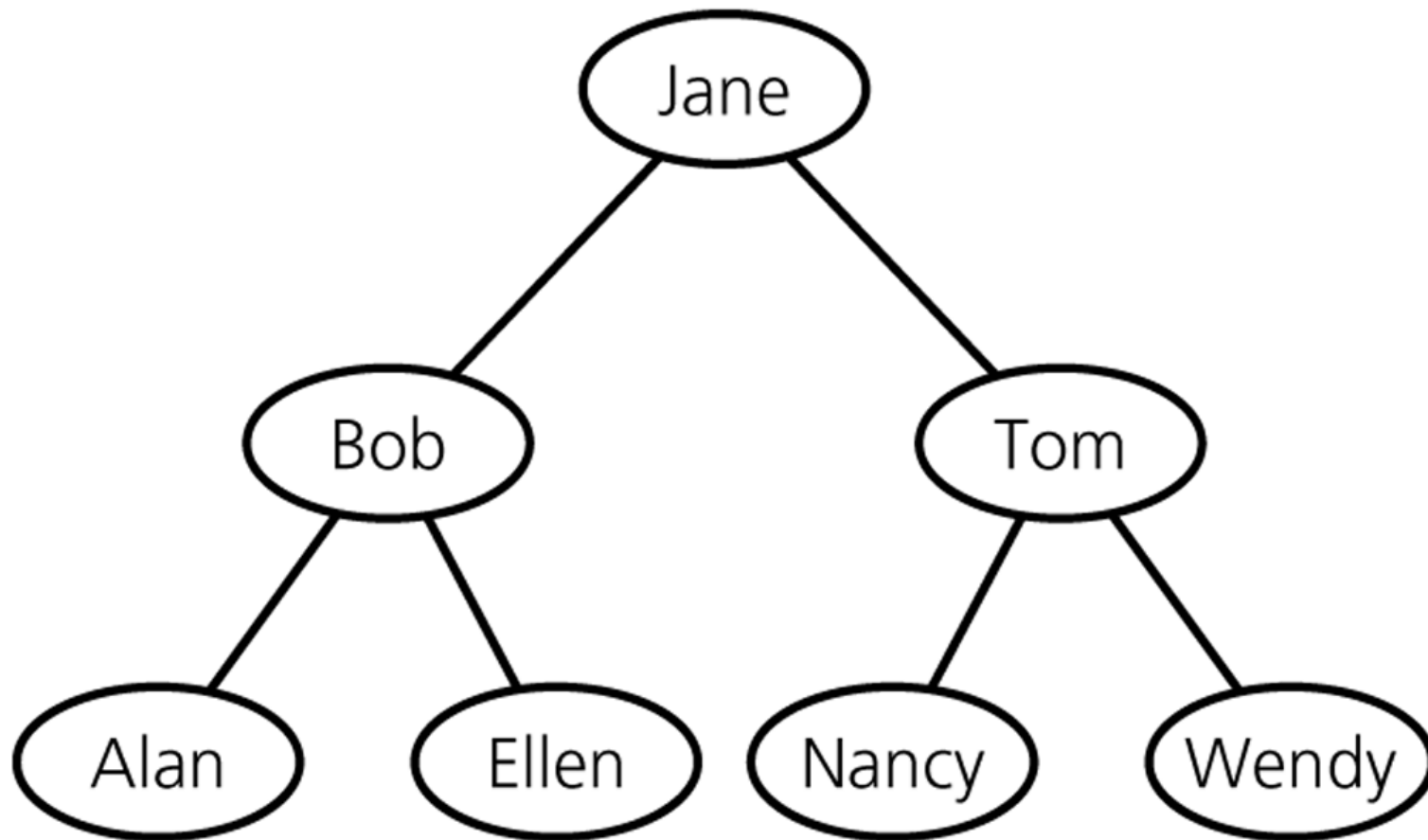
Traverse

a binary tree that
represent algebraic
expressions

- Preorder
- Inorder
- Postorder



Traverse a binary search tree



What is coming

- Trees
- Operations





Preorder traversal in a binary tree T

Preorder (T)

```
{  
    if (T is not empty) {  
        visit root;  
        Preorder ( $T_L$ );  
        Preorder( $T_R$ );  
    }
```



Inorder traversal in a binary tree T

Inorder (T)

```
{  
    if (T is not empty) {  
        Inorder ( $T_L$ );  
        visit root;  
        Inorder( $T_R$ );  
    }
```




Postorder traversal in a binary tree T

Postorder (T)

```
{  
    if (T is not empty) {  
        Postorder ( $T_L$ );  
        Postorder( $T_R$ );  
        visit root;  
    }
```



Representations of a binary tree

- Pointer-based representation

A pointer-based implementation of a binary tree

