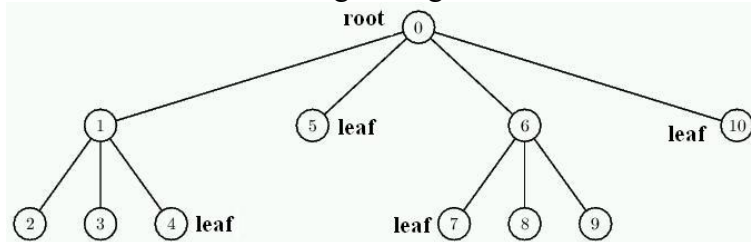


CSCI 3110

Tree, Binary tree, Binary search tree

A tree is an undirected simple graph G that is connected and has no simple cycles.

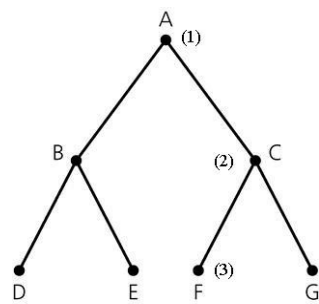
- A graph is connected if there is at least one path between any two vertices
- It is a hierarchical structure for organizing data



Tree terminologies:

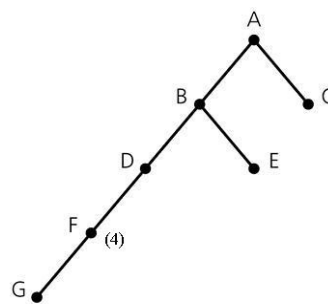
- Parent of node n
 - The node directly above node n in the tree
- Child of node n
 - a node directly below node n in the tree
- Root
 - The only node in the tree with no parent
 - When referring to a tree, we need to specify its root. A graph can be viewed as different trees by choosing different vertex as the root
- Leaf
 - A node with no children
- Siblings
 - Nodes with a common parent
- Ancestor of node n
 - A node on the path from the root to n
- Descendant of node n
 - A node on the path from n to a leaf
- Degree of a node
 - Number of child nodes for the node
- Subtree : any node in the tree together with all its descendants
- Subtree of a node n : subtree rooted at a child of n
- Height of a tree
 - The number of nodes on the longest path from the root to a leaf
- Level of a node n
 - If n is the root, it is at level 1
 - Otherwise, its level is 1 greater than the level of its parent.

Examples of these concepts in the following tree?



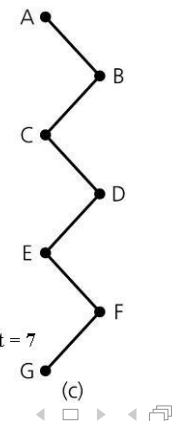
height = 3

(a)



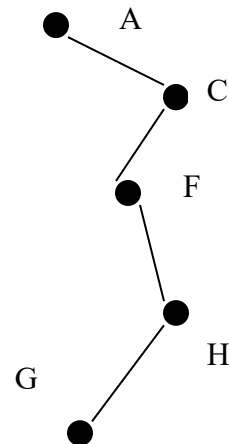
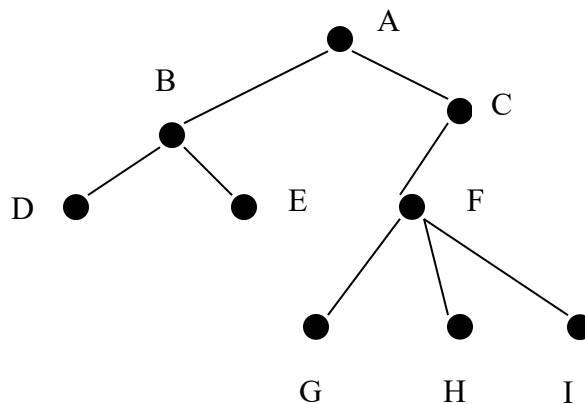
height = 5

(b)



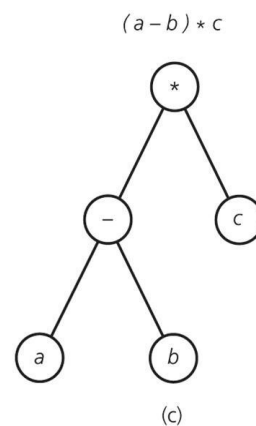
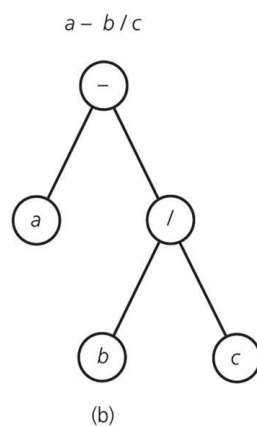
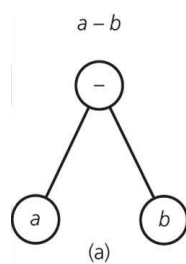
height = 7

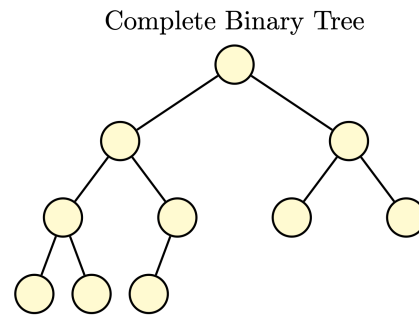
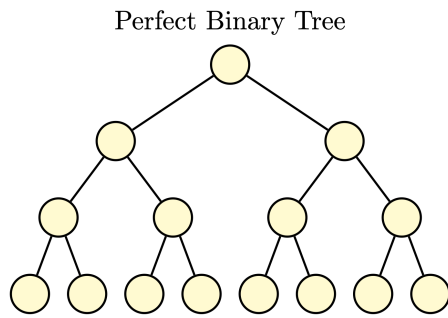
(c)



Binary tree – tree where each node has at most 2 child nodes, called left child (if any) and right child (if any).

- Equivalently, each node in a binary tree has at most two subtrees, called left subtree T_L (if any) and right subtree T_R (if any).



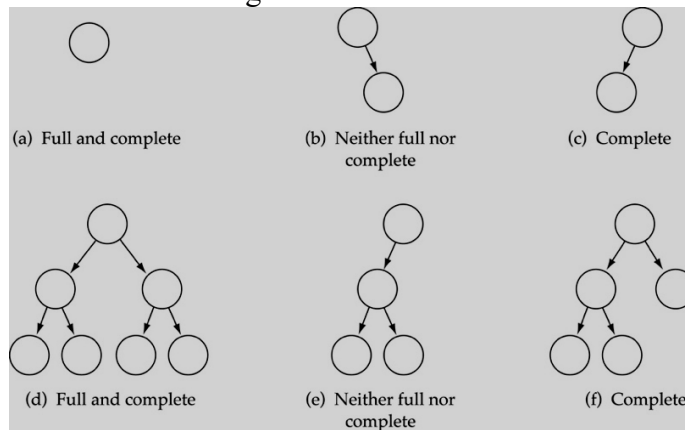
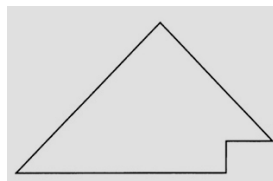


Perfect binary tree - a binary tree of height h , where all nodes at levels less than h have 2 children nodes

- How many nodes are in a perfect binary tree of height h ?
 $2^0 + 2^1 + 2^2 + 2^3 \dots 2^{(h-1)} = 2^h - 1$
- Given a perfect binary tree with N nodes, what is the height of the tree?
 $2^h - 1 = N \rightarrow h = \log_2(N+1)$

Complete binary tree:

- a binary tree that is either full or full through the next to last level
- The last level is filled in from the left to right.

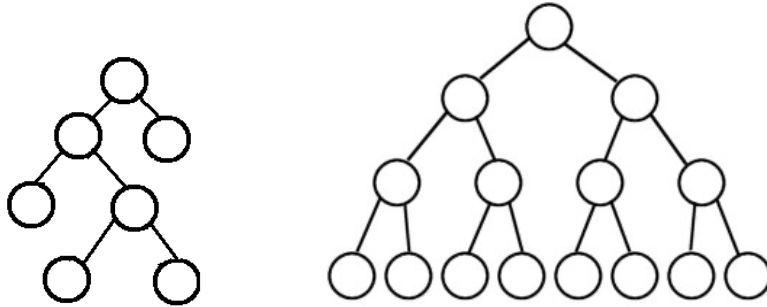


Complete binary tree : binary tree full down the level $(h-1)$, with level h filled in from left to right

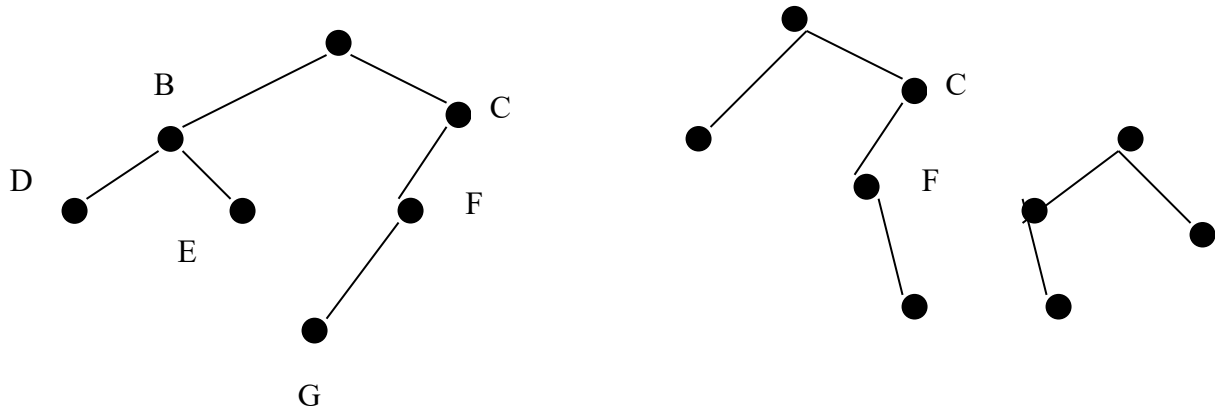
- All nodes at level $\leq h-2$ have 2 children
- When a node at level $h-1$ has children, all nodes to its left at the same level have 2 children each
- When a node at level $h-1$ has one child, it is a left child

Full binary tree:

- a binary tree in which **all of the nodes have either 0 or 2 offspring**, and
- a binary tree in which all nodes, except the leaf nodes, have two offspring.



Balanced binary tree: binary tree is balanced, if the height of any node's left subtree differs from the height of any node's right subtree by no more than 1



If a binary tree is perfect, is it always balanced?

If a binary tree is complete, is it always balanced?

If a binary tree is full, is it always balanced?

If a binary tree is balanced, is it always full?

If a binary tree is balanced, is it always complete?

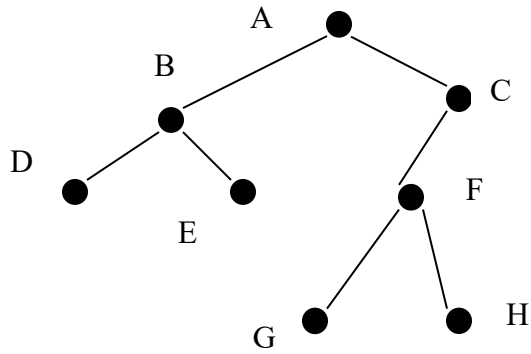
Operations of a binary tree

- **Create** an empty binary tree `BinaryTree();` // default constructor
Create a one-node(the root) binary tree `BinaryTree(const TreeItemType root);`
- Create a binary tree given the root, and its subtrees
`BinaryTree(const TreeItemType root, BinaryTree left, BinaryTree right);`
- **Destroy** a binary tree
`~BinaryTree();`
- **Determine** whether a binary tree is empty
`isEmpty();`
- **Determine or change the data in the binary tree's root**
`TreeItemType getRootData();`
`void setRootData(TreeItemType data);`
- **Attach a left or right node to the binary tree's root**
`void attachLeft(TreeItemType item) ;`
`void attachRight(TreeItemType item);`
- **Detach** a left or right subtree of the binary tree's root
`void detachLeftSubtree(BinaryTree& leftTree);`
`void detachRightSubTree(BinaryTree& rightTree);`
- **Return** a copy of the left or right subtree of the binary tree's root
`BinaryTree leftSubtree(void);`
`BinaryTree rightSubTree(void);`
- **Traverse** the nodes in a binary tree
preorder: `void preorderTraverse(FunctionType visit);`
inorder: `void inorderTraverse(FunctionType visit);`
postorder: `void postorderTraverse(FunctionType visit);`

Traversal of a binary tree

A traversal algorithm for a binary tree visits each node in the tree

- Preorder traversal : visit **root first**, then left child, than right child
(ABDECFGHI)
- Inorder traversal : visit left child first, then **visit root**, then visit right child
- Postorder traversal : visit left node, then visit right child node, **then visit root**



Preorder traversal (used to copy a tree)

```

preorder ( BinaryTree tree, FunctionType visit ) {
    if ( tree is not empty ) {
        visit(the root of tree);
        preorder(Left subtree of tree's root, visit);
        preorder(Right subtree of tree's root, visit);
    }
}
  
```

Inorder traversal (used to visit nodes in ascending order of keys, to extract inorder expression)

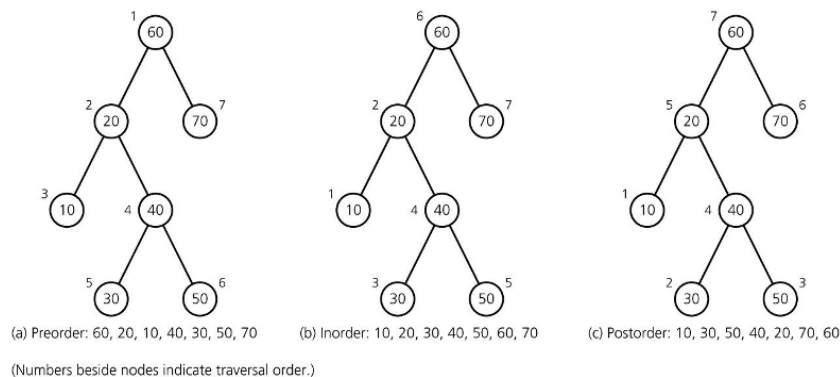
```

inorder ( BinaryTree tree, FunctionType visit ) {
    if ( tree is not empty ) {
        inorder(Left subtree of tree's root, visit);
        visit(the root of tree);
        inorder(Right subtree of tree's root, visit);
    }
}
  
```

Postorder traversal (used to extract postorder expression from a BST)

```

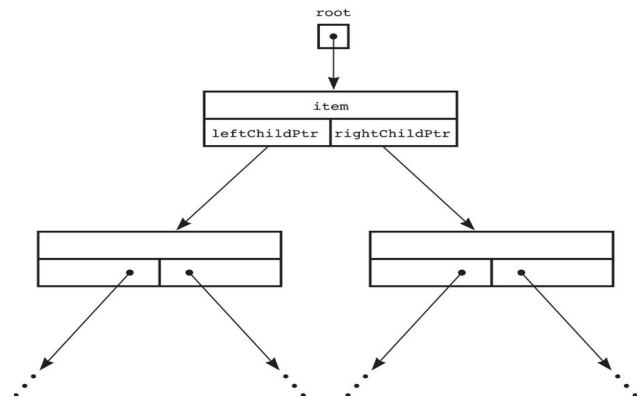
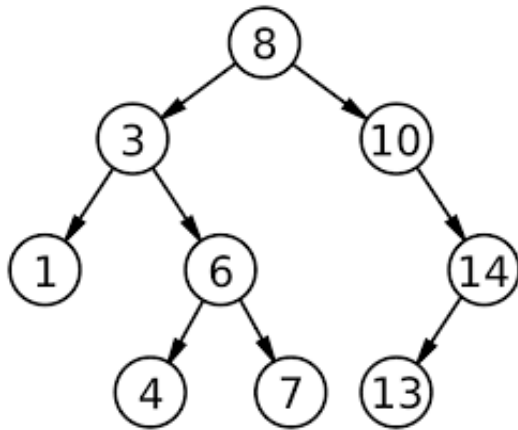
postorder ( BinaryTree tree, FunctionType visit ) {
    if ( tree is not empty ) {
        postorder(Left subtree of tree's root, visit);
        postorder(Right subtree of tree's root, visit);
        visit(the root of tree);
    }
}
  
```



The function visit is provided and passed by clients as the computation performed for each node.

A Binary Search Tree (BST)

- A binary tree that has the following properties for each node n
 - n 's value is greater than all values in its left subtree T_L
 - n 's value is less than all values in its right subtree T_R
 - Both T_L and T_R are binary search trees
- A deficiency of the ADT binary tree which is corrected by the ADT binary search tree
 - Searching for a particular item
- A data item in a binary search tree has a specially designated search key
 - A search key is the part of a record that identifies it within a collection of records
- *KeyedItem* class
 - Contains the search key as a data field and a function for accessing the search key



Defining a Binary Search Tree :

```

struct treeNode;
typedef treeNode * ptrType;

struct treeNode
{
    treeItemType Item;
    ptrType    LChildPtr, RChildPtr;
    treeNode(const treeItemType& NodeItem, ptrType L, ptrType R);
}; // end struct

treeNode::treeNode(const treeItemType& NodeItem, ptrType L, ptrType R):
Item(NodeItem), LChildPtr(L), RChildPtr(R)
{
}
  
```

Binary search tree traversal examples:

1. Binary Search Tree Operation: Search

- A recursive search function
 - Searches the binary search tree *bst* for the item whose search key is *key*.
 - Searches recursively in the right or left subtree (depending on the value of the item) until *key* matches the search key of the node's item

```
Search ( BinarySearchTree bst, KeyType key)
{
    if ( bst is empty )
        The desired record is not found, throw an exception
    else if ( key == search key of root's item )
        The desired record is found
    else if ( key < search key of root's item )
        Search(left subtree of bst, key)
    else // key > search key of root's item
        Search(right subtree of bst, key)
}
```

2. Binary Search Tree Operation: Insert

Inserts *newItem* into the binary search tree to which *treePtr* points

```
InsertItem( TreeNodePtr treePtr, TreeItemType newItem )
// Question: Should treePtr be passed by value or by reference???
{
    if ( treePtr == nullptr )
        Create a new Node, Initialize it appropriately, and let treePtr point
        to it.
    else if ( newItem.getKey() < treePtr->item.getKey() )
        InsertItem( treePtr->leftChildPtr, newItem );
    else
        InsertItem( treePtr->rightChildPtr, newItem );
}
```

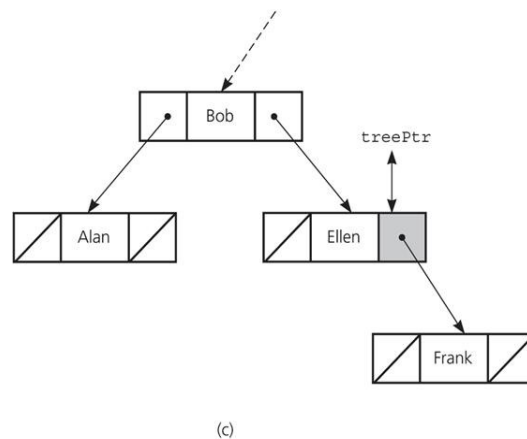
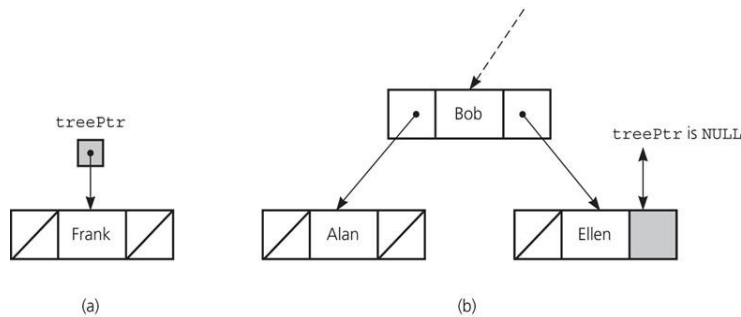
How are these functions called?

If these functions are member functions of a class, how should they be called?

Why is it necessary to have a wrapper function in these cases?

?? How to write this function as a member function of BST class ??

Ans: a wrapper function is needed to deal with recursive member function that require the access of the private data, the root, of the class



(a) Insertion into an empty tree; (b) search terminates at a leaf; (c) insertion at a leaf

Practice Question:

How to build a binary search tree by inserting the following key values one by one into an empty tree ? 45, 23, 100, 80, 56, 87, 5, 25

3. Binary Search Tree Operation: **Delete**

- **Three possible cases** for deleting a node N
 1. N is a leaf
 - Set the pointer in N's parent to nullptr (or NULL)
 2. Delete a node with only left child
 3. Delete a node with only right child
 4. N has two children
 - Locate another node M that is the leftmost node in N's right subtree
M's search key is called the in-order successor of N's search key
 - Copy the item that is in node M to node N
 - Update the pointer to M to point to M's right subtree.
 - Remove the node M from the tree

4. How to determine the level of each node in a binary (search) tree? (use pre-order traversal)

function call: `AssignLevel(root, 1);` // assuming root at level 1.

function definition:

```
void AssignLevel(ptrType nodePtr, int level)
{
    if (nodePtr != nullptr)
    {
        cout << "Node with key" << nodePtr->Item.Key()
              << "is at level " << level << endl;
        nodePtr->PutLevel(level);
        AssignLevel(nodePtr->LChildPtr, level+1);
        AssignLevel(nodePtr->RChildPtr, level+1);
    }
}
```

When defined in the BST class as a method:

define the method:

```
void BST::AssignLevel(ptrType nodePtr, int level)
{
    if (nodePtr != nullptr)
    {
        cout << "Node with key" << nodePtr->Item.Key()
              << "is at level " << level << endl;
        nodePtr->Level = level;
        AssignLevel(nodePtr->LChildPtr, level+1);
        AssignLevel(nodePtr->RChildPtr, level+1);
    }
}

void BST::DetermineLevel()
{
    AssignLevel(root, 1);
}
```

call the method: `OneBSTree.DetermineLevel();`

5. How to determine the height of a binary (search) tree? (use post-order traversal)

function call: `height = FindHeight(root);`

function definition:

```
int FindHeight(ptrType nodePtr)
{
    int leftHeight, rightHeight;
    if (nodePtr != NULL)
    {
        leftHeight = FindHeight(nodePtr->LChildPtr)+1;
        rightHeight = FindHeight(nodePtr->RChildPtr)+1;
        if (leftHeight < rightHeight)
            return rightHeight;
        else
            return leftHeight;
    }
    else
        return 0;
}
```

6. Saving and restoring binary search tree

- **Restoring the binary search tree to its original shape**
(use **pre-order traversal** to save the tree to a file, then rebuild the tree by inserting the items one by one into an empty tree)
- **Saving the binary search tree in order and Restoring the binary search tree to the minimum height**
 - Uses **inorder traversal** to save the tree to a file
 - Can be used if the data is sorted and the number of nodes in the tree is known

```
// save the records in a binary search tree in order
void ReadSave(ofstream& outFile, ptrType nodePtr)
{
    if (nodePtr != nullptr)
    {
        ReadSave(outFile, nodePtr->LChildPtr);

        // Write information in nodePtr to outFile (overloaded << operator)
        outFile << nodePtr->Item << endl;

        ReadSave(outFile, nodePtr->RChildPtr);
    }
}

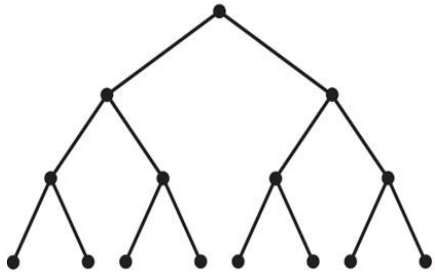
// restore the binary search tree to minimum height
void Restore(ifstream & inFile, ptrType & nodePtr, int numOfRecords)
{
    if (numOfRecords > 0)
    {
        nodePtr = new treeNode;

        Restore (inFile, nodePtr->LChildPtr, numOfRecords/2);

        inFile >> newItem; // overloaded >> operator
        nodePtr->Item = newItem // overloaded = operator

        Restore(inFile, nodePtr->RChildPtr, (numOfRecords-1)/2);
    }
    else
        nodePtr = nullptr;
}
```

The efficiency of Binary Search Tree operations

	Level	Number of nodes at this level	Number of nodes at this and previous levels
	1	$1 = 2^0$	$1 = 2^1 - 1$
	2	$2 = 2^1$	$3 = 2^2 - 1$
	3	$4 = 2^2$	$7 = 2^3 - 1$
	4	$8 = 2^3$	$15 = 2^4 - 1$
• • • •	• •		•
• • • •	• •		•
• • • •	• •		•
	h	2^{h-1}	$2^h - 1$

Uses the ADT binary search tree to sort an array of records into search-key order

- Average case: $O(n * \log n)$
- Worst case: $O(n^2)$