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**Computer Aided Decision Support
International Research Master 2
2013 - 2014**

Operational research project draft

**Ant colony optimization algorithm for a
combined routing and scheduling problem**

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Formalization of the problem

The problem can be defined from three different subproblems :

- The two machines permutation flowshop. We have two machines M_1 and M_2 organized as a flowshop. A set of n jobs have to be scheduled into a sequence. Each job j has a value a_j and b_j that correspond to the time required to complete the job j on M_1 and M_2 respectively.
- The grouping of jobs to be delivered. When the jobs are completed for the two machines permutation flowshop we must form groups of jobs, where a group must contain at least one job. These groups are then used for the third subproblem.
- A traveling salesman problem. We consider one truck taking the group of jobs previously formed where each job has a destination k_j . The idea is to find the sequence of destinations in order to deliver every job from the group while minimizing the traveling distance. We have a matrix K ($m \times m$) where m is the number of destinations and K_i^j is the distance between the destination i and j .

The objective function is to minimize the sum of lateness of the jobs $\sum_{j=1}^{j \leq n} L_j$. The lateness L_j of a job j is the difference between the completion time of the job c_j and its due date d_j : $L_j = c_j - d_j$. The completion time of a job is not equal to the completion time of the job for the flowshop problem but for the traveling salesman problem. **c_j is equal to the date at which the truck is coming back to the factory after its deliveries. We consider that the truck has no capacity limit, it can carry as many jobs as possible.**

The input data of our problem are the following :

- The number n of jobs,
- The values a_j , b_j and d_j of job j ($\forall j \in [1, n]$),
- The number of m destinations,
- The destination k_j of job j ($\forall j \in [1, n]$),
- The matrix K ($m \times m$) of distances between each destination.

During the resolution of the problem we must find the sequence of jobs on the flowshop, the groups of jobs to be delivered and the sequence of delivery for each group.

We proposed a first way of encoding a solution : a table of size n containing the sequence of job for the flowshop subproblem, a table of size $2n$ containing the groups and sequences for the traveling salesman subproblem. The first cell of the second table contains the number of jobs for the first group, the following cells contain the number of jobs in the delivery sequence and so on. The figure 1.1 shows a solution for a simple instance of the problem with its encoding.

The sequence here is $\{1, 2, 3\}$ the first group G_1 of jobs is $G_1 = \{1, 2\}$ and 1 is delivered first, the second group G_2 only contains one job $G_2 = \{3\}$.

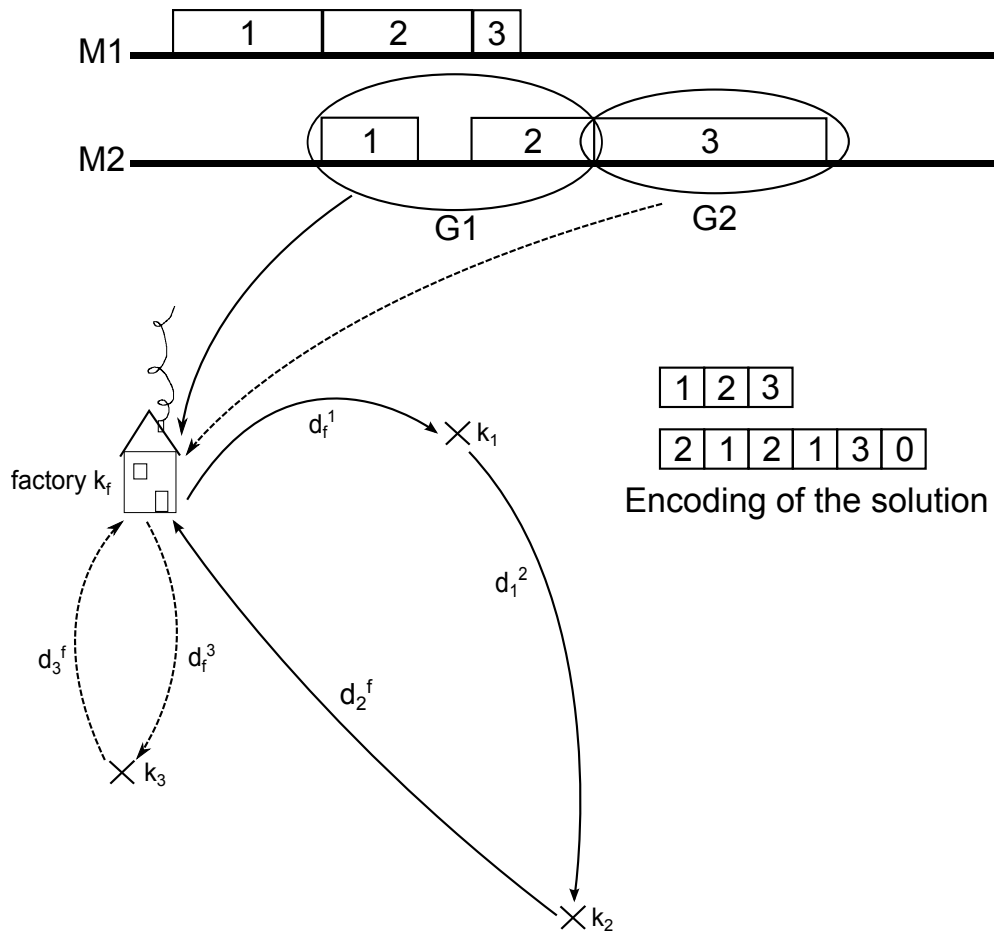


Figure 1.1: A solution of the problem for a simple instance

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