

Network Programming

Introduction to Computer Systems
23rd Lecture, Dec. 9, 2024

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The Notion of an internet Protocol

Human protocols:

- “what’s the time?”
- “I have a question”
- introductions

Network protocols:

- computers (devices) rather than humans
- all communication activity in Internet governed by protocols

■ How is it possible to send bits across incompatible LANs and WANs?

■ Solution: *protocol* software running on each host and router

- Protocol is a set of rules that governs how hosts and routers should cooperate when they transfer data from network to network.
- Smooths out the differences between the different networks

What Does an internet Protocol Do?

■ Provides a *naming scheme*

- An internet protocol defines a uniform format for *host addresses*
- Each host (and router) is assigned at least one of these internet addresses that uniquely identifies it

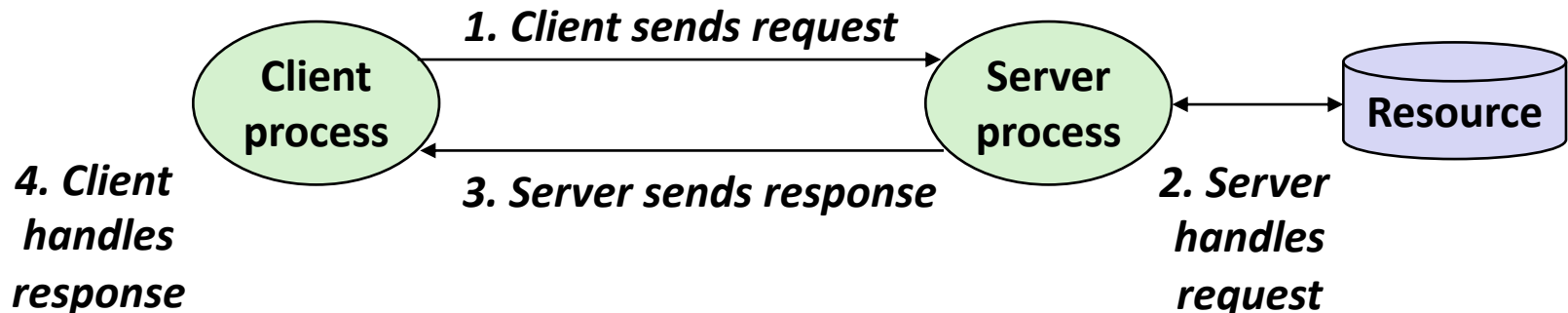
■ Provides a *delivery mechanism*

- An internet protocol defines a standard transfer unit (*packet*)
- Packet consists of *header* and *payload*
 - Header: contains info such as packet size, source and destination addresses
 - Payload: contains data bits sent from source host

A Client-Server Transaction

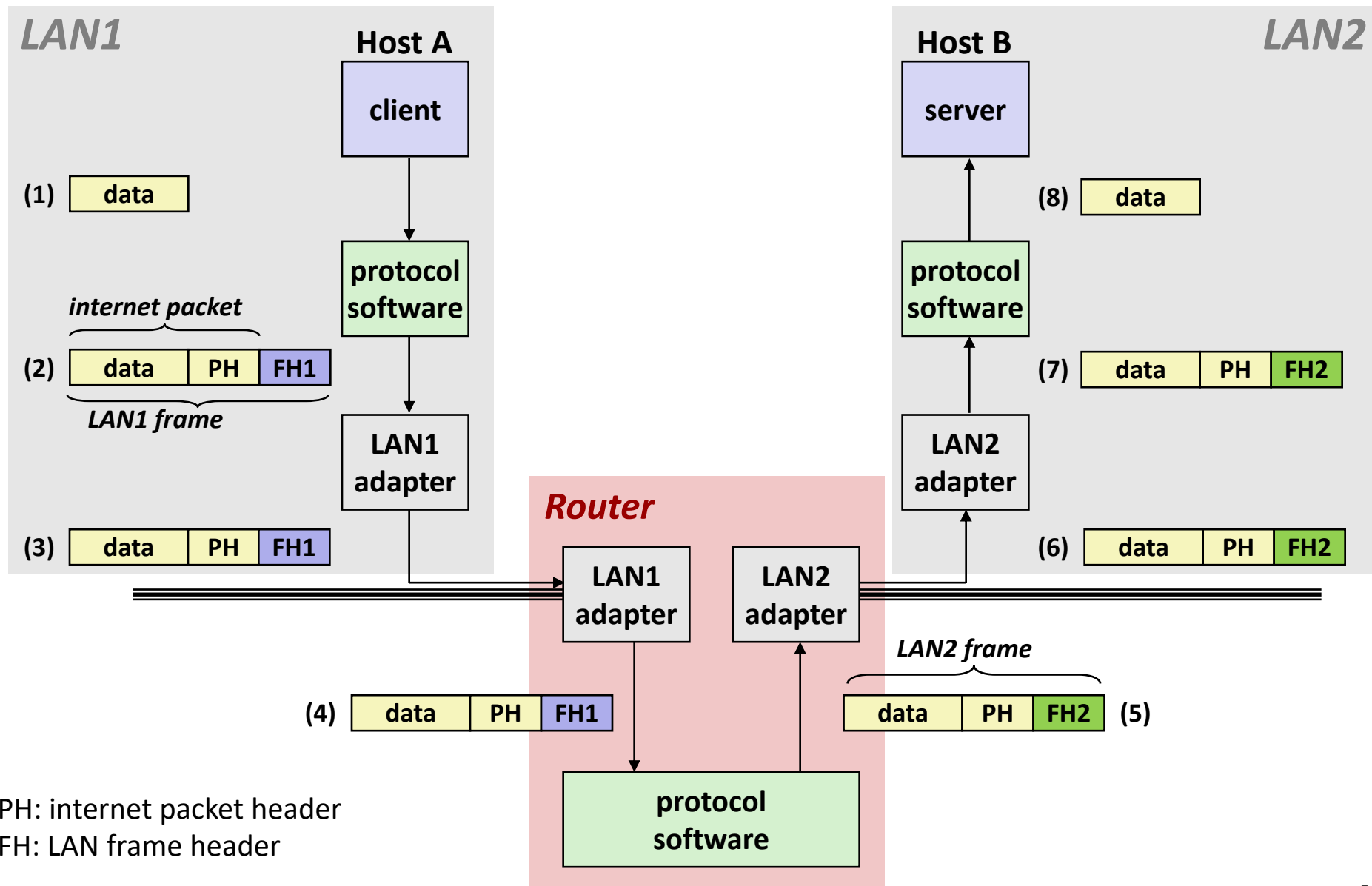
■ Most network applications are based on the client-server model:

- A **server** process and one or more **client** processes
- Server manages some **resource**
- Server provides **service** by manipulating resource for clients
- Server activated by request from client (vending machine analogy)



Note: clients and servers are processes running on hosts (can be the same or different hosts)

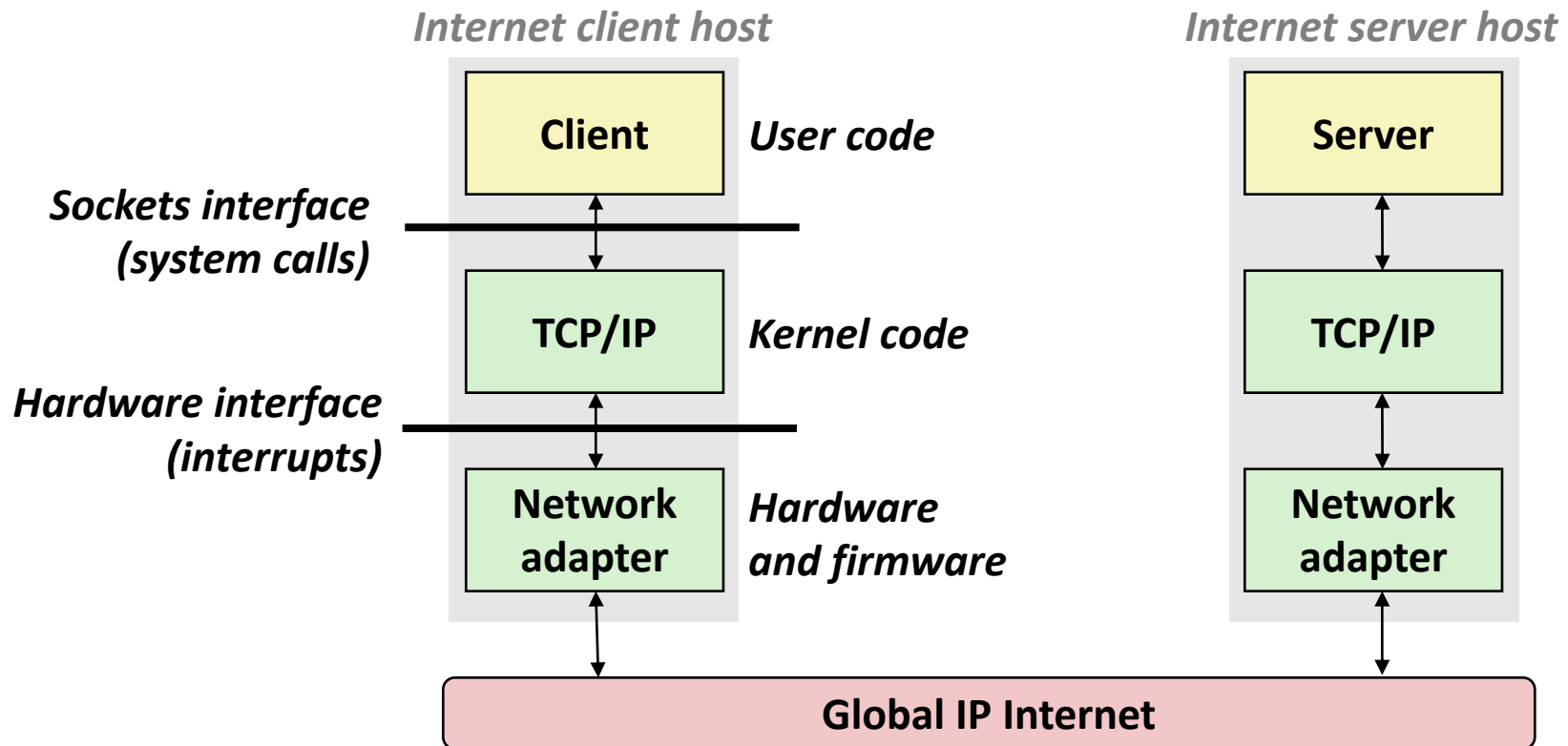
Transferring internet Data Via Encapsulation



Global IP Internet (upper case)

- Most famous example of an internet
- Based on the TCP/IP protocol family
 - IP (Internet Protocol)
 - Provides *basic naming scheme* and unreliable *delivery capability* of packets (datagrams) from *host-to-host*
 - UDP (Unreliable Datagram Protocol)
 - Uses IP to provide *unreliable* datagram delivery from *process-to-process*
 - TCP (Transmission Control Protocol)
 - Uses IP to provide *reliable* byte streams from *process-to-process* over *connections*
- Accessed via a mix of Unix file I/O and functions from the *sockets interface*

Hardware and Software Organization of an Internet Application



A Programmer's View of the Internet

1. Hosts are mapped to a set of 32-bit *IP addresses*

- 128.2.203.179
- 127.0.0.1 (always *localhost*)

2. As a convenience for humans, the Domain Name System maps a set of identifiers called Internet *domain names* to IP addresses:

- 128.2.217.3 is mapped to www.cs.cmu.edu
- www.cs.cmu.edu “resolves to” 128.2.217.3

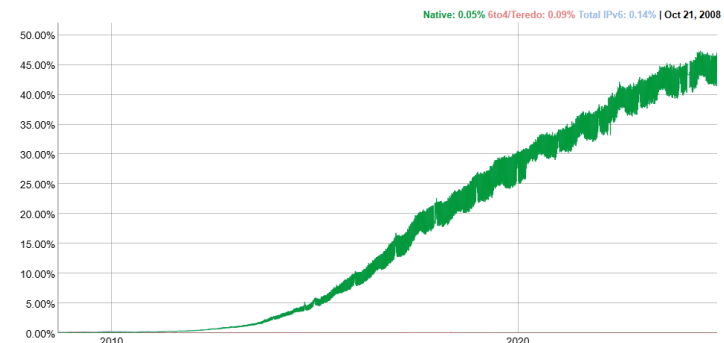
3. A process on one Internet host can communicate with a process on another Internet host over a *connection*

Aside: IPv4 and IPv6

- **IPv4** (*Internet Protocol Version 4*) specified 1981
 - 32-bit host addresses
 - Known to not be enough for everyone since ~1990
 - Majority of Internet traffic still carried by IPv4
- **IPv6** (*Internet Protocol Version 6*) specified 1996
 - 128-bit addresses (2001:0db8:0:0:0:0:cafe:1a7e)
 - Intended to replace IPv4
 - Very slow adoption due to need to replace routers
- **Application programmers mostly don't have to care**
 - Sockets API makes it easy to write code that seamlessly uses either, as necessary

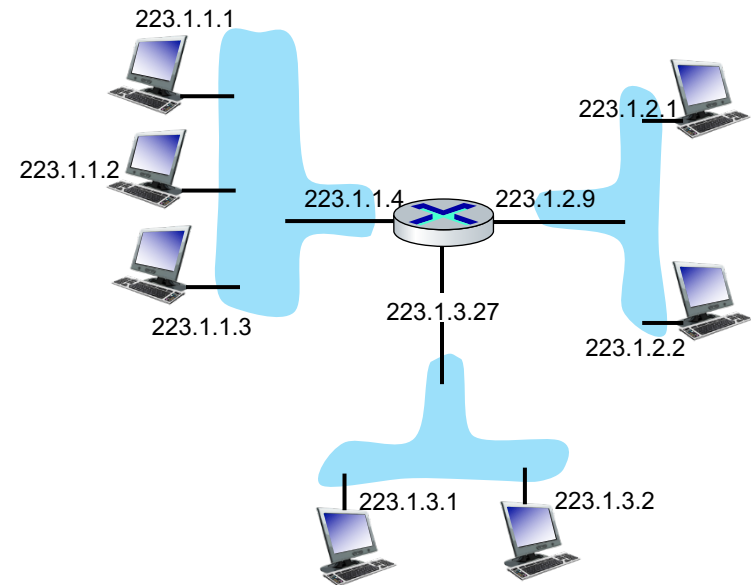
IPv6 traffic to Google

<https://www.google.com/intl/en/ipv6/statistics.html>



(1) IP Addresses

- **IP address:** 32-bit identifier associated with each host or router *interface*
- **interface:** connection between host/router and physical link
 - router's typically have multiple interfaces
 - host typically has one or two interfaces (e.g., wired Ethernet, wireless 802.11)



dotted-decimal IP address notation:

223.1.1.1 = 11011111 00000001 00000001 00000001

223
1
1
1

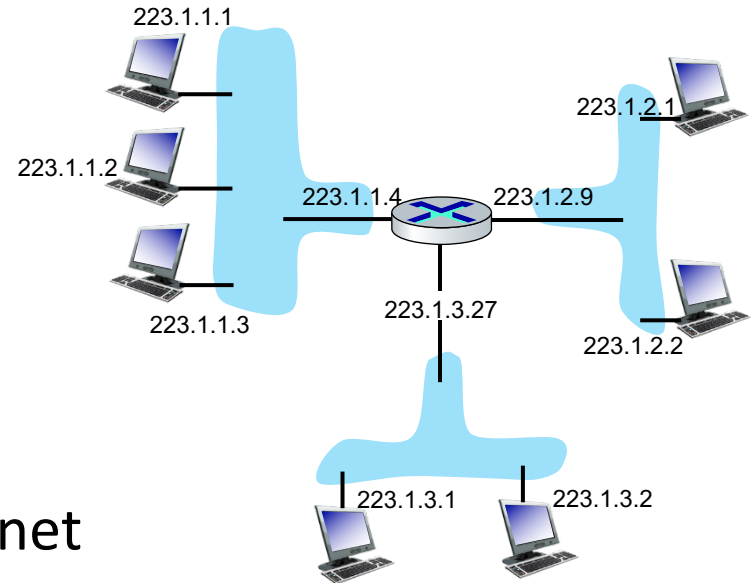
Subnets

■ *What's a subnet ?*

- device interfaces that can physically reach each other **without passing through an intervening router**

■ IP addresses have structure:

- **subnet part**: devices in same subnet have common high order bits
- **host part**: **remaining** low order bits

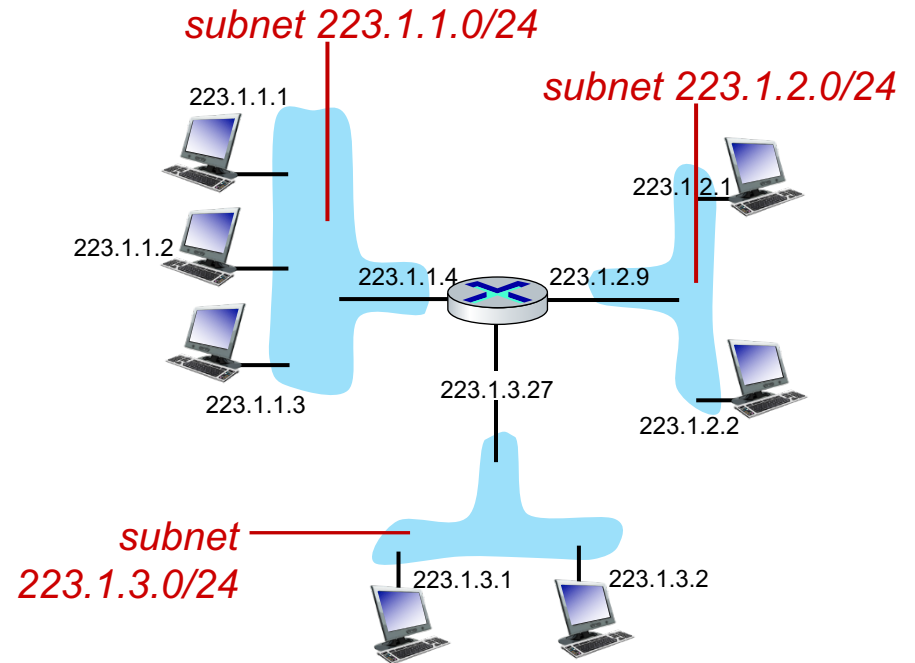


network consisting of 3 subnets

Subnets

Recipe for defining subnets:

- detach each interface from its host or router, creating “islands” of isolated networks
- each isolated network is called a *subnet*

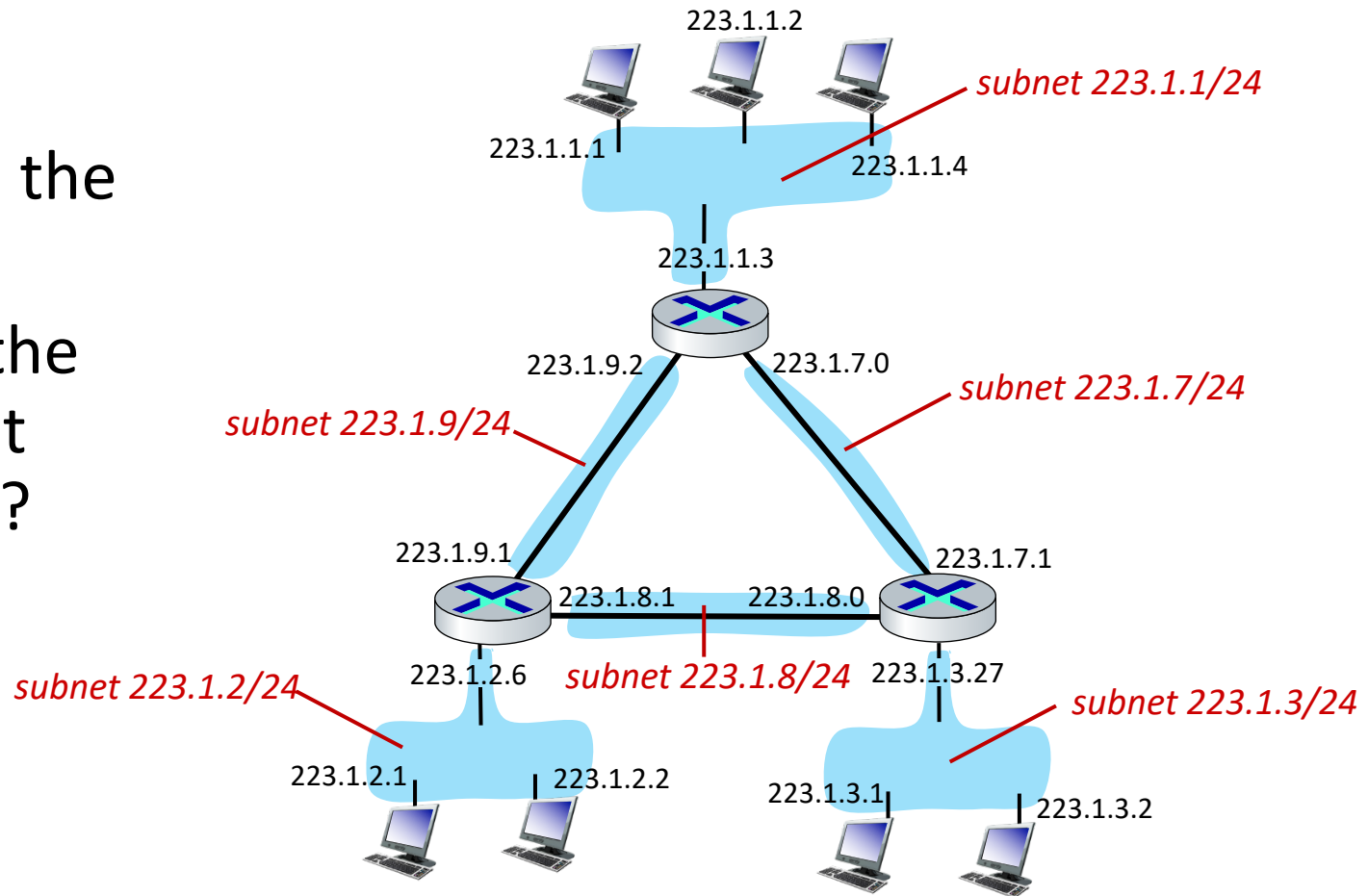


subnet mask: /24

(high-order 24 bits: subnet part of IP address)

Subnets

- where are the subnets?
- what are the /24 subnet addresses?



IP Address Structure

■ IP (V4) Address space divided into classes:

	0	1	2	3	8	16	24	31		
Class A	0	Net ID				Host ID				
Class B	1	0	Net ID				Host ID			
Class C	1	1	0	Net ID				Host ID		
Class D	1	1	1	0	Multicast address					
Class E	1	1	1	1	Reserved for experiments					

■ Network ID Written in form w.x.y.z/n

- n = number of bits in host address
- E.g., CMU written as 128.2.0.0/16
 - Class B address

■ Unrouted (private) IP addresses:

10.0.0.0/8 172.16.0.0/12 192.168.0.0/16

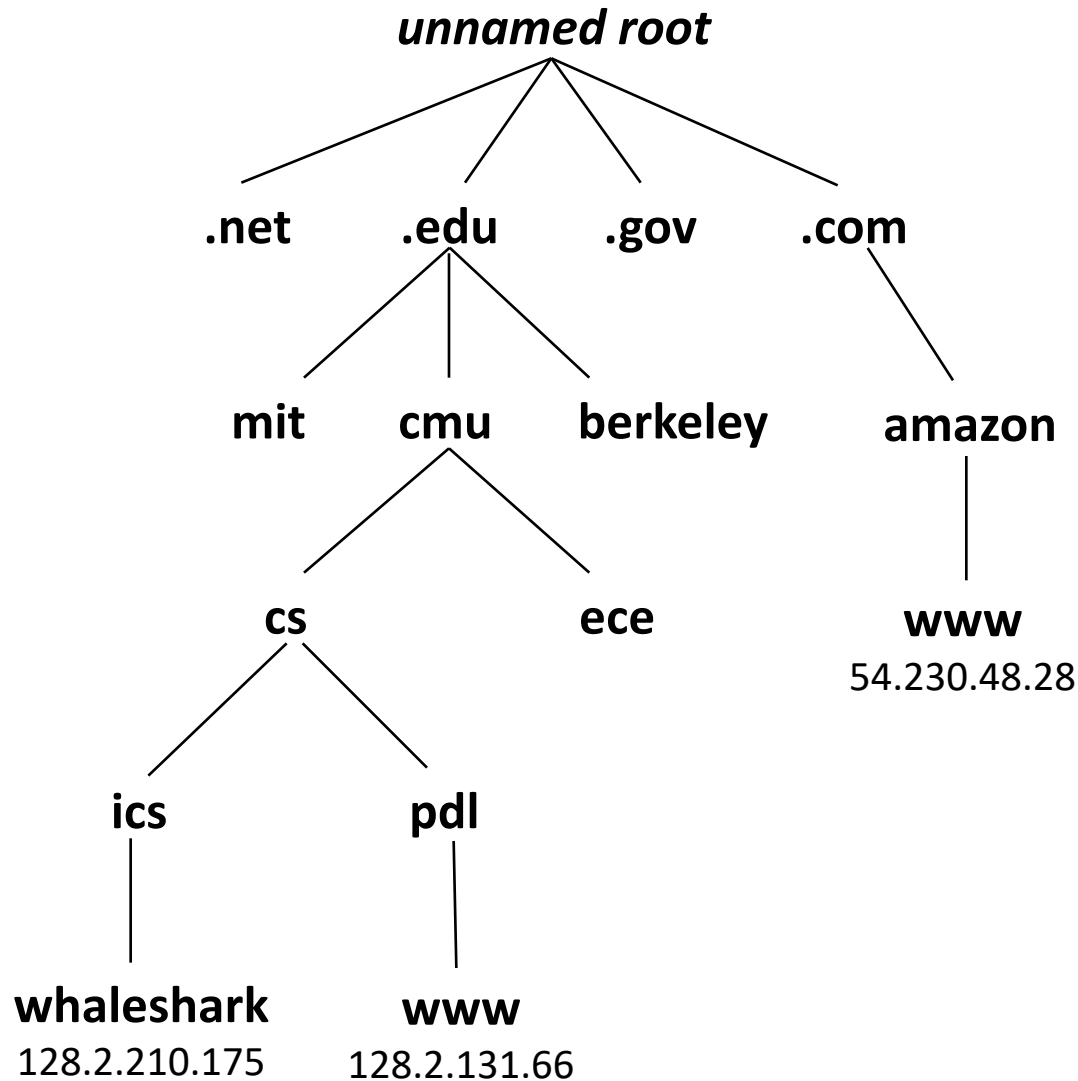
IP Addresses

- **32-bit IP addresses are stored in an *IP address struct***
 - IP addresses are always stored in memory in *network byte order* (big-endian byte order)
 - True in general for any integer transferred in a packet header from one machine to another.
 - E.g., the port number used to identify an Internet connection.

```
/* Internet address structure */
struct in_addr {
    uint32_t    s_addr; /* network byte order (big-endian) */
};
```

- **By convention, each byte in a 32-bit IP address is represented by its decimal value and separated by a period**
 - IP address: `0x8002C2F2` = `128.2.194.242`

(2) Internet Domain Names



Domain Naming System (DNS)

- The Internet maintains a mapping between IP addresses and domain names in a huge worldwide distributed database called **DNS**
- Conceptually, programmers can view the DNS database as a collection of millions of *host entries*.
 - Each host entry defines the mapping between a set of domain names and IP addresses.
 - In a mathematical sense, a host entry is an equivalence class of domain names and IP addresses.

Properties of DNS Mappings

- Can explore properties of DNS mappings using `nslookup`
 - (Output edited for brevity)
- Each host has a locally defined domain name `localhost` which always maps to the *loopback address* `127.0.0.1`
- Use `hostname` to determine real domain name of local host:

```
linux> nslookup localhost  
Address: 127.0.0.1
```

```
linux> hostname  
whaleshark.ics.cs.cmu.edu
```

Properties of DNS Mappings (cont)

- Simple case: one-to-one mapping between domain name and IP address:

```
linux> nslookup whaleshark.ics.cs.cmu.edu  
Address: 128.2.210.175
```

- Multiple domain names mapped to the same IP address:

```
linux> nslookup cs.mit.edu  
Address: 18.25.0.23  
linux> nslookup eecs.mit.edu  
Address: 18.25.0.23
```

- And backwards:

```
linux> nslookup 18.25.0.23  
23.0.25.18.in-addr.apra    name = eecs.mit.edu.
```

Properties of DNS Mappings (cont)

- Multiple domain names mapped to multiple IP addresses:

```
linux> nslookup www.tencent.com
Address: 116.131.60.102
Address: 115.56.90.84
Address: 116.196.155.54
Address: 115.56.90.198
Address: 116.196.145.223
```

```
linux> nslookup www.tencent.com
Address: 116.196.155.54
Address: 116.196.145.223
Address: 115.56.90.84
Address: 116.131.60.102
Address: 115.56.90.198
```

- Some valid domain names don't map to any IP address:

```
linux> nslookup ics.cs.cmu.edu
(No Address given)
```

(3) Internet Connections

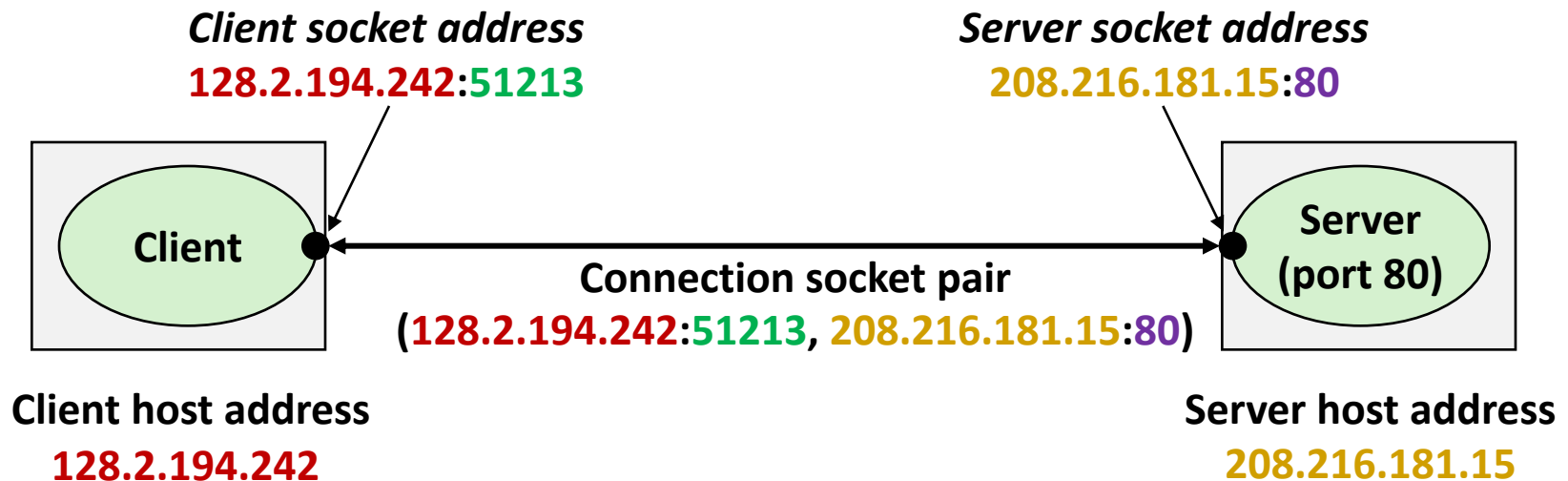
- Clients and servers most often communicate by sending streams of bytes over TCP *connections*. Each connection is:
 - *Point-to-point*: connects a pair of processes.
 - *Full-duplex*: data can flow in both directions at the same time,
 - *Reliable*: stream of bytes sent by the source is eventually received by the destination in the same order it was sent.
- A *socket* is an endpoint of a connection
 - *Socket address* is an `IPAddress:port` pair
- A *port* is a 16-bit integer that identifies a process:
 - *Ephemeral port*: Assigned automatically by client kernel when client makes a connection request.
 - *Well-known port*: Associated with some *service* provided by a server (e.g., port 80 is associated with Web servers)

Well-known Service Names and Ports

- Popular services have permanently assigned *well-known ports* and corresponding *well-known service names*:
 - echo servers: echo 7
 - ftp servers: ftp 21
 - ssh servers: ssh 22
 - email servers: smtp 25
 - Unencrypted Web servers: http 80
 - SSL/TLS encrypted Web: https 443
- Mappings between well-known ports and service names is contained in the file `/etc/services` on each Linux machine.

Anatomy of a Connection

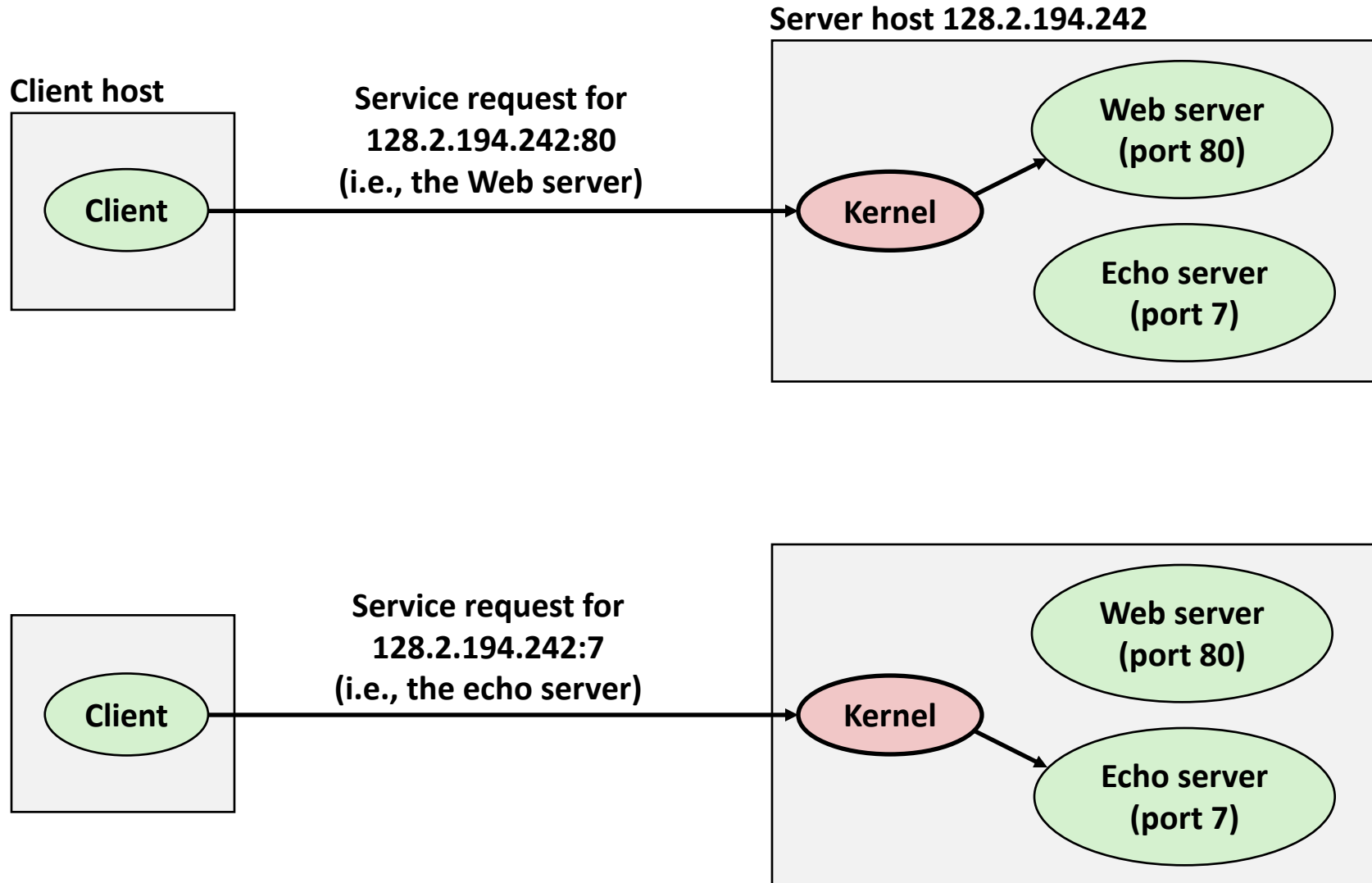
- A connection is uniquely identified by the socket addresses of its endpoints (*socket pair*)
 - (cliaddr:cliport, servaddr:servport)



51213 is an ephemeral port allocated by the kernel

80 is a well-known port associated with Web servers

Using Ports to Identify Services



Sockets Interface

- Set of system-level functions used in conjunction with Unix I/O to build network applications.
- Created in the early 80's as part of the original Berkeley distribution of Unix that contained an early version of the Internet protocols.
- Available on all modern systems
 - Unix variants, Windows, OS X, IOS, Android, ARM

Sockets

■ What is a socket?

- To the kernel, a socket is an endpoint of communication
- To an application, a socket is a file descriptor that lets the application read/write from/to the network
- Using the FD abstraction lets you reuse code & interfaces
 - **Remember:** All Unix I/O devices, including networks, are modeled as files

■ Clients and servers communicate with each other by reading from and writing to socket descriptors



- The main distinction between regular file I/O and socket I/O is how the application “opens” the socket descriptors

Socket Address Structures

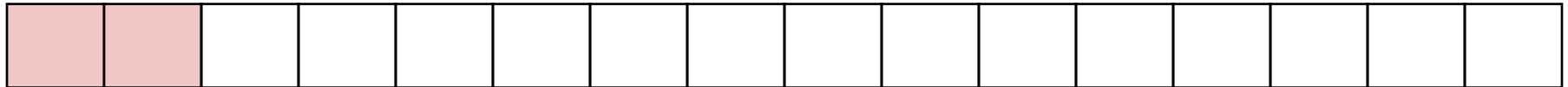
■ Generic socket address:

- For address arguments to **connect**, **bind**, and **accept** (*next lecture*)
- Necessary only because C did not have generic (**void ***) pointers when the sockets interface was designed
- For casting convenience, we adopt the Stevens convention:

typedef struct sockaddr SA;

```
struct sockaddr {  
    uint16_t  sa_family;    /* Protocol family */  
    char      sa_data[14]; /* Address data */  
};
```

sa_family



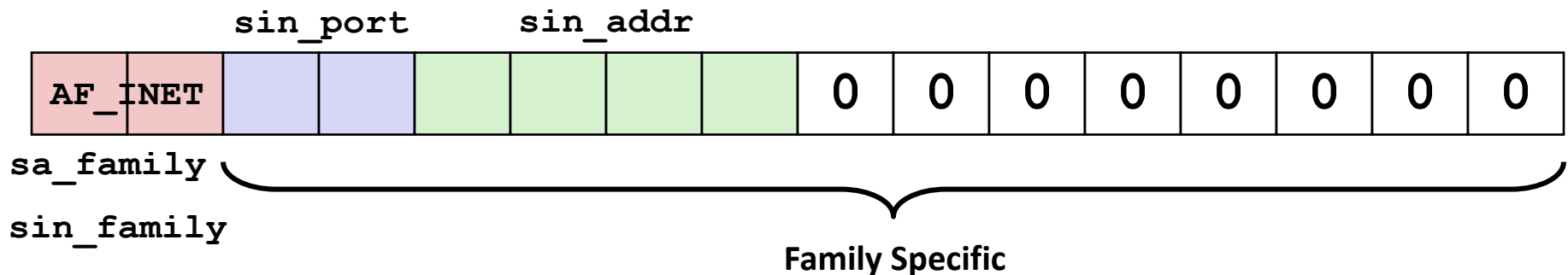
Family Specific

Socket Address Structures

■ Internet (IPv4) specific socket address:

- Must cast (`struct sockaddr_in *`) to (`struct sockaddr *`) for functions that take socket address arguments.

```
struct sockaddr_in {
    uint16_t      sin_family; /* Protocol family (always AF_INET) */
    uint16_t      sin_port;   /* Port num in network byte order */
    struct in_addr sin_addr;   /* IP addr in network byte order */
    unsigned char sin_zero[8]; /* Pad to sizeof(struct sockaddr) */
};
```



Host and Service Conversion: `getaddrinfo`

- `getaddrinfo` is the modern way to convert string representations of hostnames, host addresses, ports, and service names to socket address structures.
 - Replaces obsolete `gethostbyname` and `getservbyname` funcs.
- **Advantages:**
 - Reentrant (can be safely used by threaded programs).
 - Allows us to write portable protocol-independent code
 - Works with both IPv4 and IPv6
- **Disadvantages**
 - Somewhat complex
 - Fortunately, a small number of usage patterns suffice in most cases.

Host and Service Conversion: `getaddrinfo`

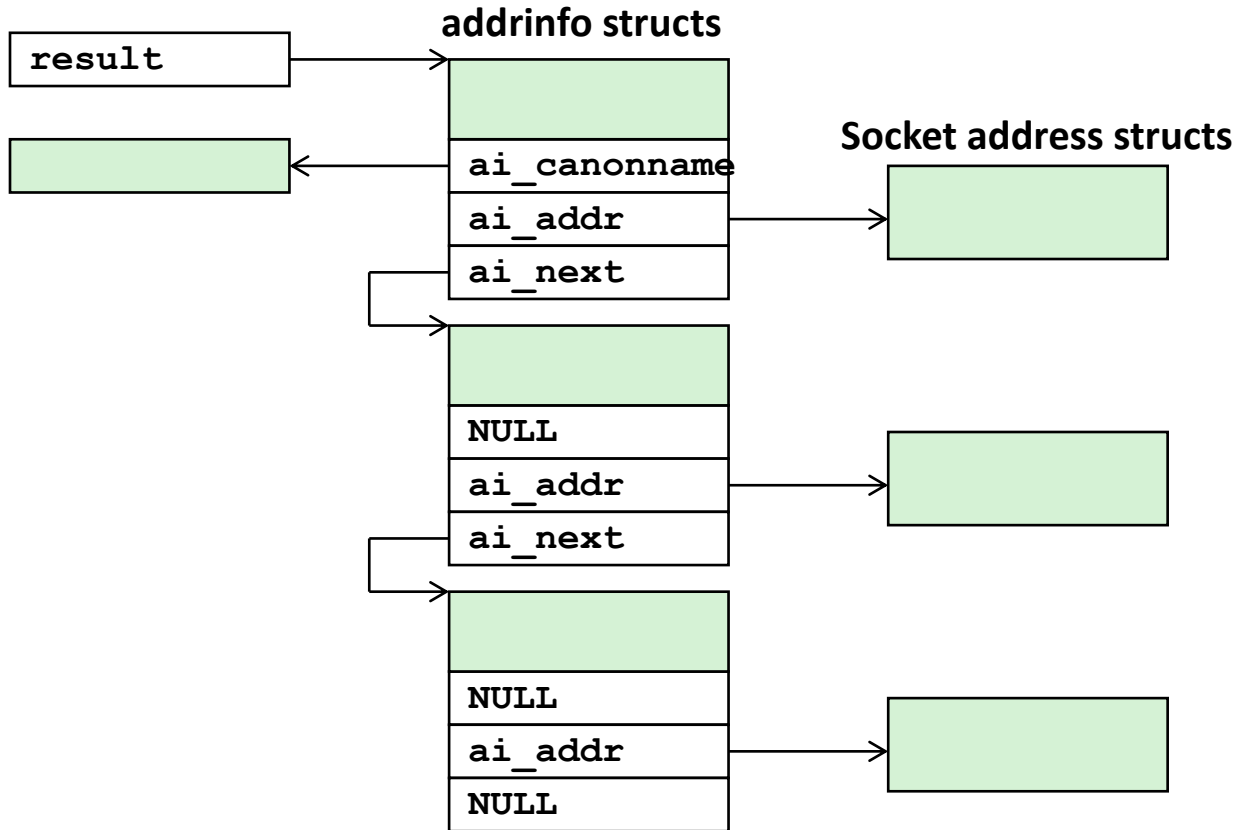
```
int getaddrinfo(const char *host,           /* Hostname or address */
               const char *service,       /* Port or service name */
               const struct addrinfo *hints, /* Input parameters */
               struct addrinfo **result);  /* Output linked list */

void freeaddrinfo(struct addrinfo *result); /* Free linked list */

const char *gai_strerror(int errcode);     /* Return error msg */
```

- Given `host` and `service`, `getaddrinfo` returns `result` that points to a linked list of **`addrinfo`** structs, each of which points to a corresponding socket address struct, and which contains arguments for the sockets interface functions.
- **Helper functions:**
 - `freeaddrinfo` frees the entire linked list.
 - `gai_strerror` converts error code to an error message.

Linked List Returned by getaddrinfo



addrinfo Struct

```
struct addrinfo {  
    int          ai_flags;      /* Hints argument flags */  
    int          ai_family;     /* First arg to socket function */  
    int          ai_socktype;   /* Second arg to socket function */  
    int          ai_protocol;   /* Third arg to socket function */  
    char         *ai_canonname; /* Canonical host name */  
    size_t       ai_addrlen;    /* Size of ai_addr struct */  
    struct sockaddr *ai_addr;    /* Ptr to socket address structure */  
    struct addrinfo *ai_next;    /* Ptr to next item in linked list */  
};
```

- Each `addrinfo` struct returned by `getaddrinfo` contains arguments that can be passed directly to `socket` function.
- Also points to a socket address struct that can be passed directly to `connect` and `bind` functions.

Host and Service Conversion: `getnameinfo`

- `getnameinfo` is the inverse of `getaddrinfo`, converting a socket address to the corresponding host and service.
 - Replaces obsolete `gethostbyaddr` and `getservbyport` funcs.
 - Reentrant and protocol independent.

```
int getnameinfo(const SA *sa, socklen_t salen, /* In: socket addr */
               char *host, size_t hostlen, /* Out: host */
               char *serv, size_t servlen, /* Out: service */
               int flags); /* optional flags */
```

Conversion Example

```
#include "csapp.h"

int main(int argc, char **argv)
{
    struct addrinfo *p, *listp, hints;
    char buf[MAXLINE];
    int rc, flags;

    /* Get a list of addrinfo records */
    memset(&hints, 0, sizeof(struct addrinfo));
    // hints.ai_family = AF_INET;          /* IPv4 only */
    hints.ai_socktype = SOCK_STREAM; /* Connections only */
    if ((rc = getaddrinfo(argv[1], NULL, &hints, &listp)) != 0) {
        fprintf(stderr, "getaddrinfo error: %s\n", gai_strerror(rc));
        exit(1);
    }
}
```

hostinfo.c

Conversion Example (cont)

```
/* Walk the list and display each IP address */
flags = NI_NUMERICHOST; /* Display address instead of name */
for (p = listp; p; p = p->ai_next) {
    Getnameinfo(p->ai_addr, p->ai_addrlen,
                buf, MAXLINE, NULL, 0, flags);
    printf("%s\n", buf);
}

/* Clean up */
Freeaddrinfo(listp);

exit(0);
}
```

hostinfo.c

Running hostinfo

```
whaleshark> ./hostinfo localhost  
127.0.0.1
```

```
whaleshark> ./hostinfo whaleshark.ics.cs.cmu.edu  
128.2.210.175
```

```
whaleshark> ./hostinfo twitter.com  
199.16.156.230  
199.16.156.38  
199.16.156.102  
199.16.156.198
```

```
whaleshark> ./hostinfo google.com  
172.217.15.110  
2607:f8b0:4004:802::200e
```

Socket Programming Example

- **Echo server and client**
- **Server**
 - Accepts connection request
 - Repeats back lines as they are typed
- **Client**
 - Requests connection to server
 - Repeatedly:
 - Read line from terminal
 - Send to server
 - Read reply from server
 - Print line to terminal

Echo Server/Client Session Example

Client

```

bambooshark: ./echoclient whaleshark.ics.cs.cmu.edu 6616      (A)
This line is being echoed                                     (B)
This line is being echoed
This one is, too                                             (C)
This one is, too
^D
bambooshark: ./echoclient whaleshark.ics.cs.cmu.edu 6616      (D)
This one is a new connection                                 (E)
This one is a new connection
^D

```

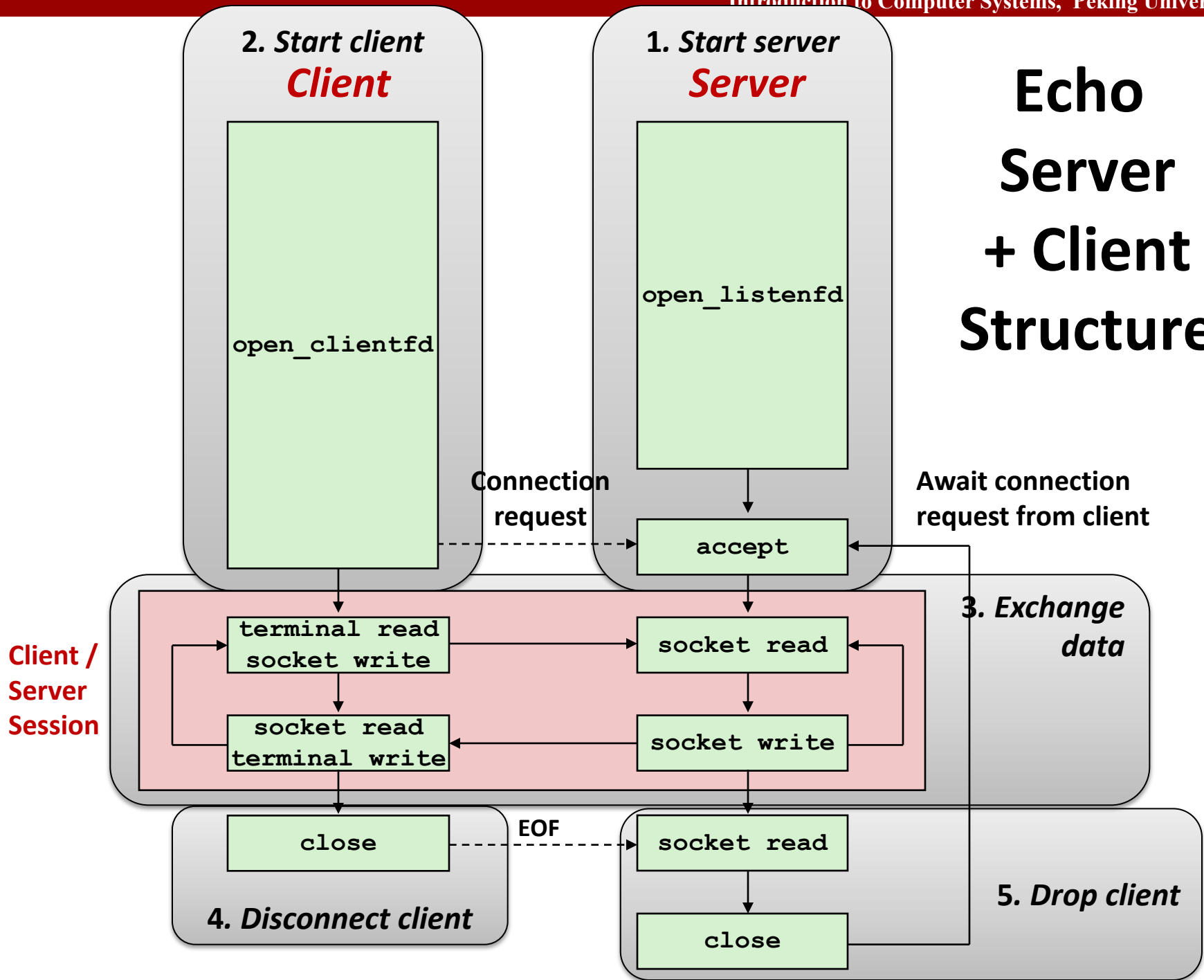
Server

```

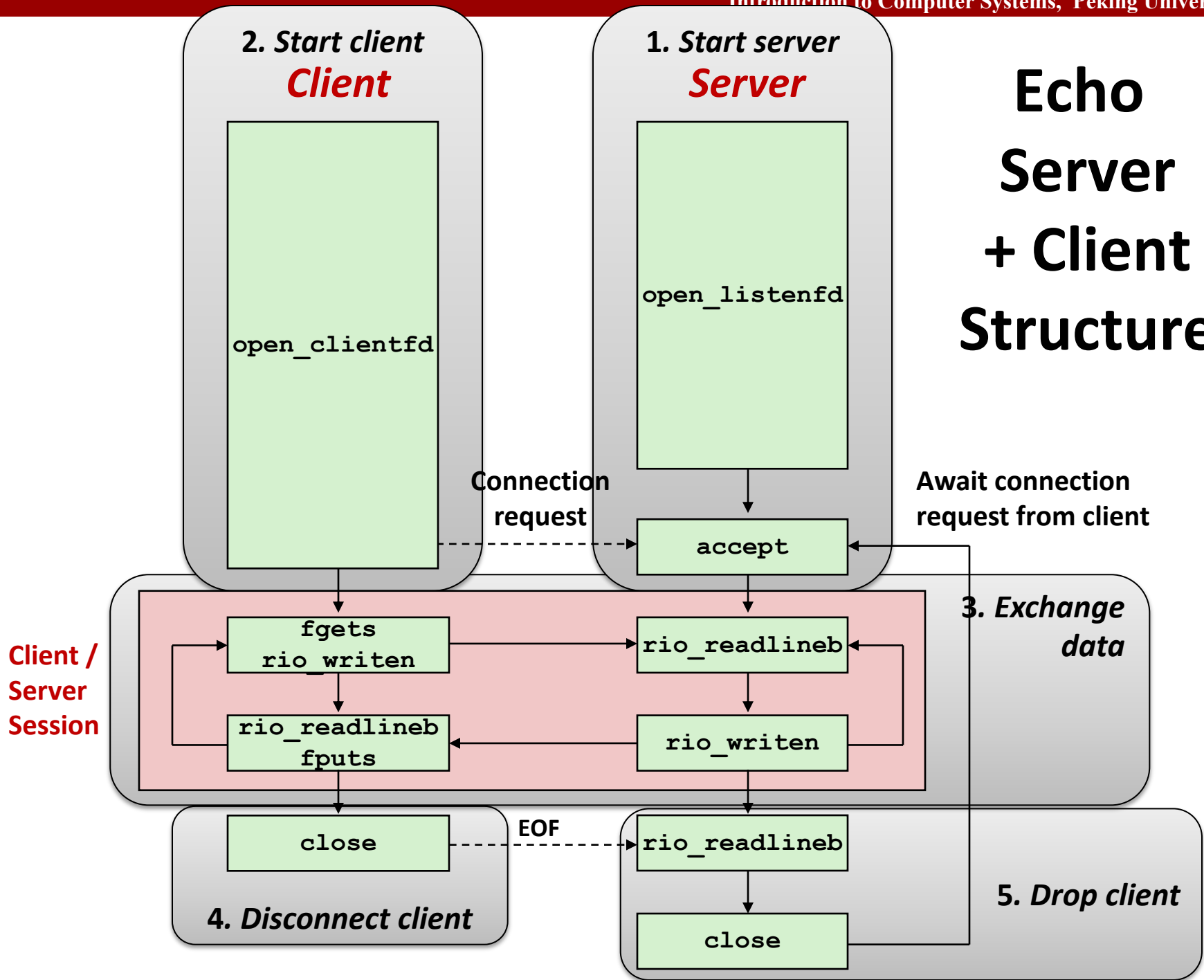
whaleshark: ./echoserveri 6616
Connected to (BAMBOOSHARK.ICS.CS.CMU.EDU, 33707)            (A)
server received 26 bytes                                     (B)
server received 17 bytes                                     (C)
Connected to (BAMBOOSHARK.ICS.CS.CMU.EDU, 33708)            (D)
server received 29 bytes                                     (E)

```

Echo Server + Client Structure



Echo Server + Client Structure



Recall: Unbuffered RIO Input/Output

- Same interface as Unix `read` and `write`
- Especially useful for transferring data on network sockets

```
#include "csapp.h"
```

```
ssize_t rio_readn(int fd, void *usrbuf, size_t n);  
ssize_t rio_writen(int fd, void *usrbuf, size_t n);
```

Return: num. bytes transferred if OK, 0 on EOF (`rio_readn` only), -1 on error

- `rio_readn` returns short count only if it encounters EOF
 - Only use it when you know how many bytes to read
- `rio_writen` never returns a short count
- Calls to `rio_readn` and `rio_writen` can be interleaved arbitrarily on the same descriptor

Recall: Buffered RIO Input Functions

- Efficiently read text lines and binary data from a file partially cached in an internal memory buffer

```
#include "csapp.h"

void rio_readinitb(rio_t *rp, int fd);

ssize_t rio_readlineb(rio_t *rp, void *usrbuf, size_t maxlen);
ssize_t rio_readnb(rio_t *rp, void *usrbuf, size_t n);
```

Return: num. bytes read if OK, 0 on EOF, -1 on error

- **rio_readlineb** reads a *text line* of up to **maxlen** bytes from file **fd** and stores the line in **usrbuf**
 - Especially useful for reading text lines from network sockets
- Stopping conditions
 - **maxlen** bytes read
 - EOF encountered
 - Newline ('\n') encountered

Echo Client: Main Routine

```
#include "csapp.h"

int main(int argc, char **argv)
{
    int clientfd;
    char *host, *port, buf[MAXLINE];
    rio_t rio;

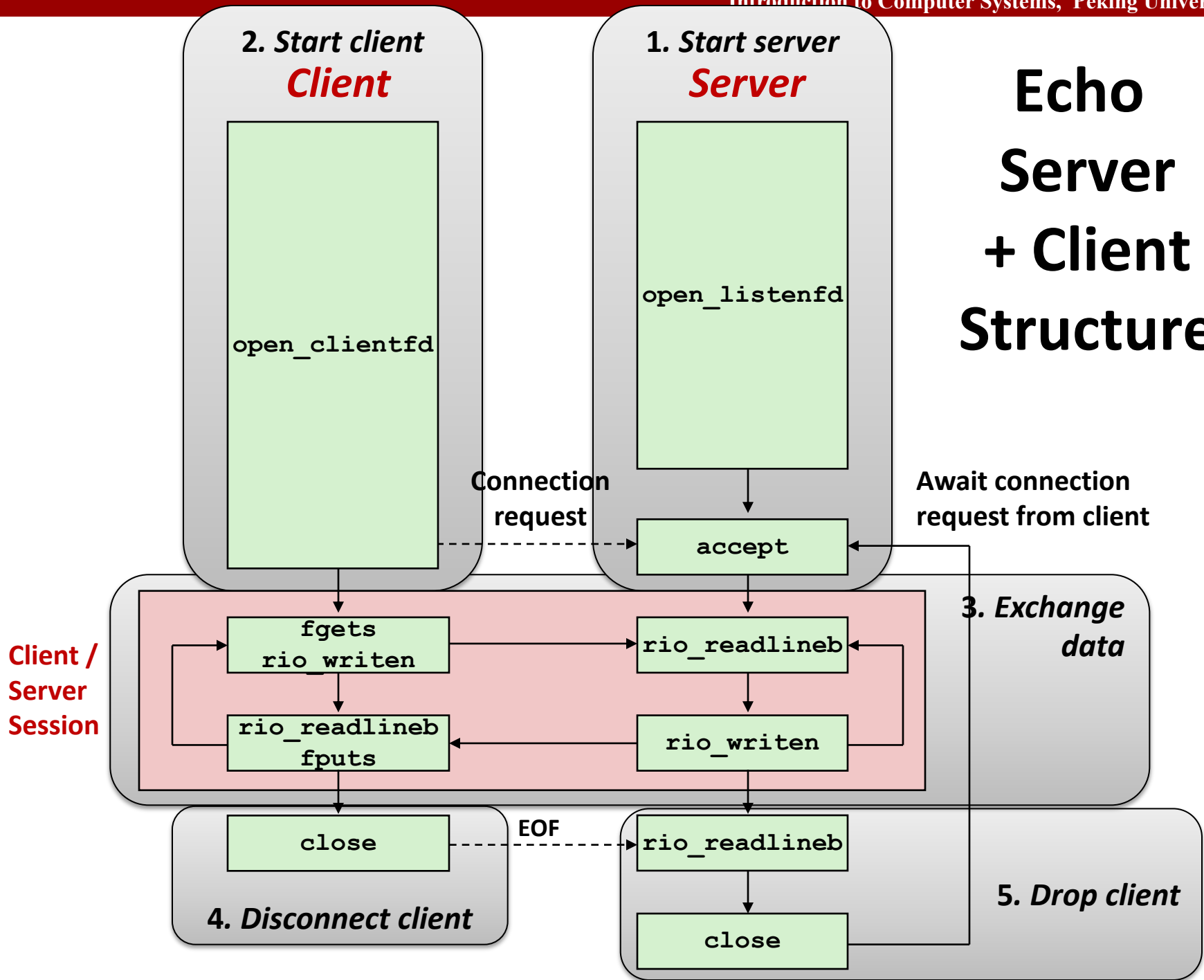
    host = argv[1];
    port = argv[2];

    clientfd = Open_clientfd(host, port);
    Rio_readinitb(&rio, clientfd);

    while (Fgets(buf, MAXLINE, stdin) != NULL) {
        Rio_writen(clientfd, buf, strlen(buf));
        Rio_readlineb(&rio, buf, MAXLINE);
        Fputs(buf, stdout);
    }
    Close(clientfd);
    exit(0);
}
```

echoclient.c

Echo Server + Client Structure



Iterative Echo Server: Main Routine

```
#include "csapp.h"
void echo(int connfd);

int main(int argc, char **argv)
{
    int listenfd, connfd;
    socklen_t clientlen;
    struct sockaddr_storage clientaddr; /* Enough room for any addr */
    char client_hostname[MAXLINE], client_port[MAXLINE];

    listenfd = Open_listenfd(argv[1]);
    while (1) {
        clientlen = sizeof(struct sockaddr_storage); /* Important! */
        connfd = Accept(listenfd, (SA *)&clientaddr, &clientlen);
        Getnameinfo((SA *)&clientaddr, clientlen,
                    client_hostname, MAXLINE, client_port, MAXLINE, 0);
        printf("Connected to (%s, %s)\n", client_hostname, client_port);
        echo(connfd);
        Close(connfd);
    }
    exit(0);
}
```

echoserveri.c

Echo Server: echo function

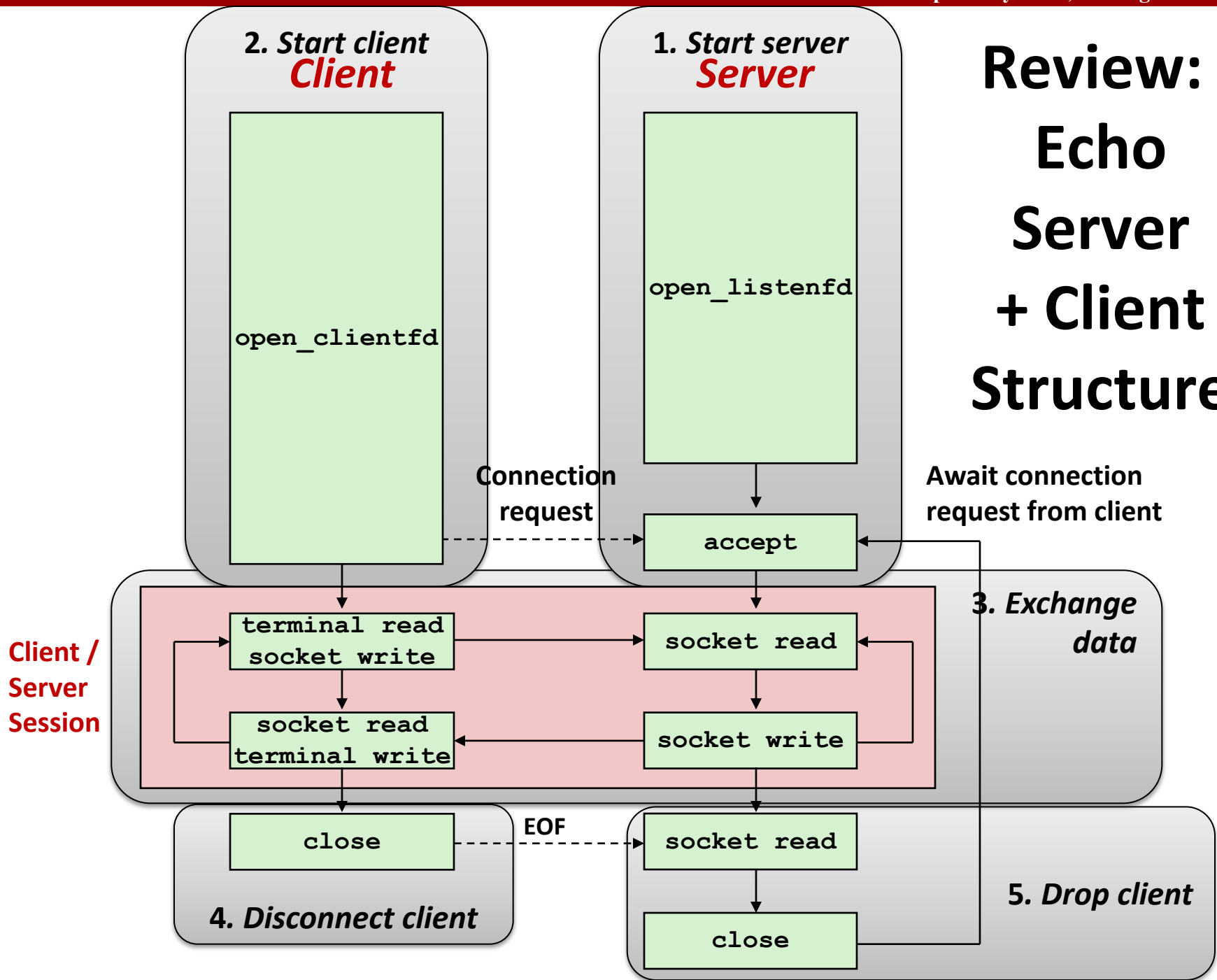
- The server uses RIO to read and echo text lines until EOF (end-of-file) condition is encountered.
 - EOF condition caused by client calling `close(clientfd)`

```
void echo(int connfd)
{
    size_t n;
    char buf[MAXLINE];
    rio_t rio;

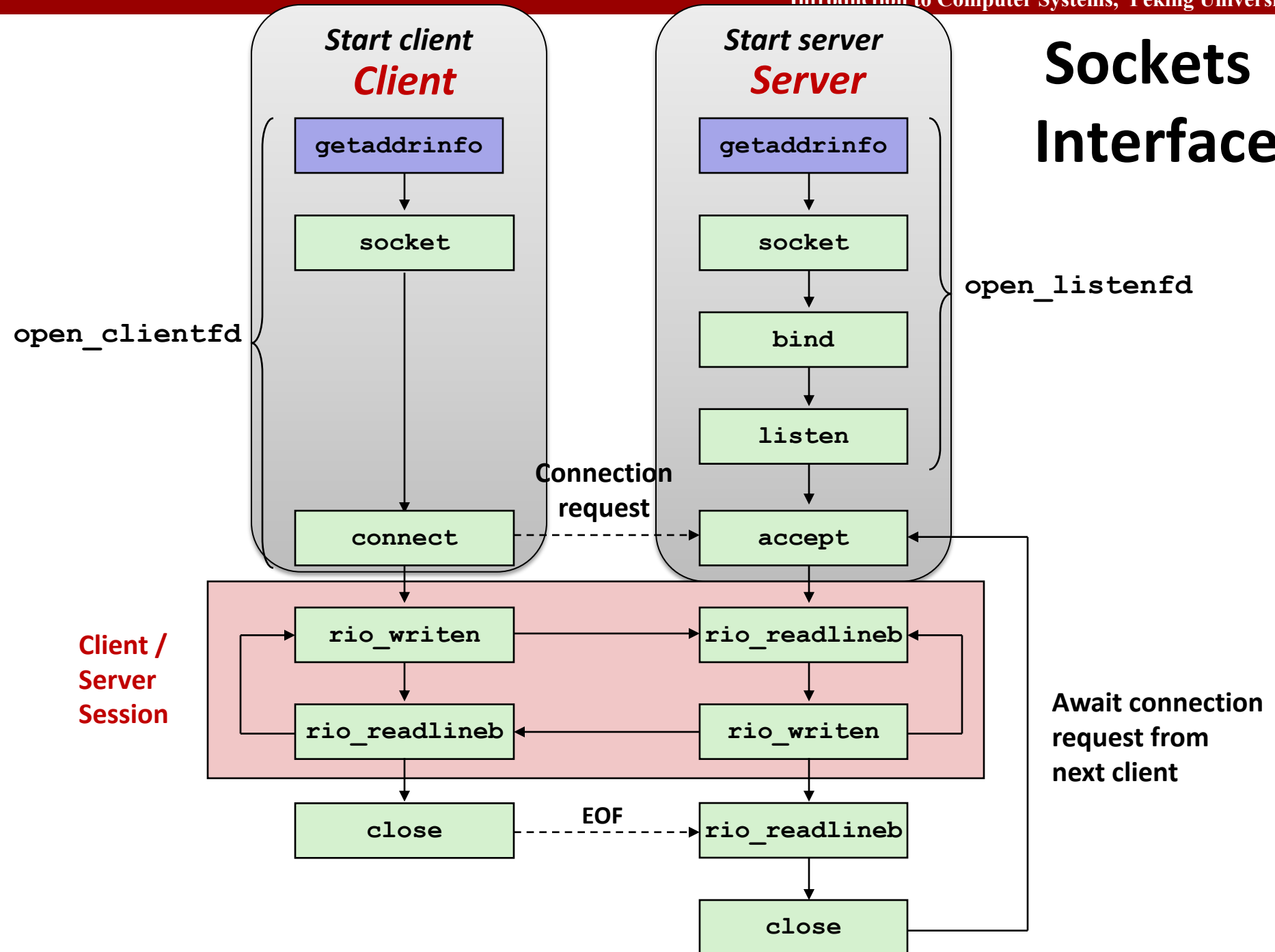
    Rio_readinitb(&rio, connfd);
    while((n = Rio_readlineb(&rio, buf, MAXLINE)) != 0) {
        printf("server received %d bytes\n", (int)n);
        Rio_writen(connfd, buf, n);
    }
}
```

echo.c

Review: Echo Server + Client Structure

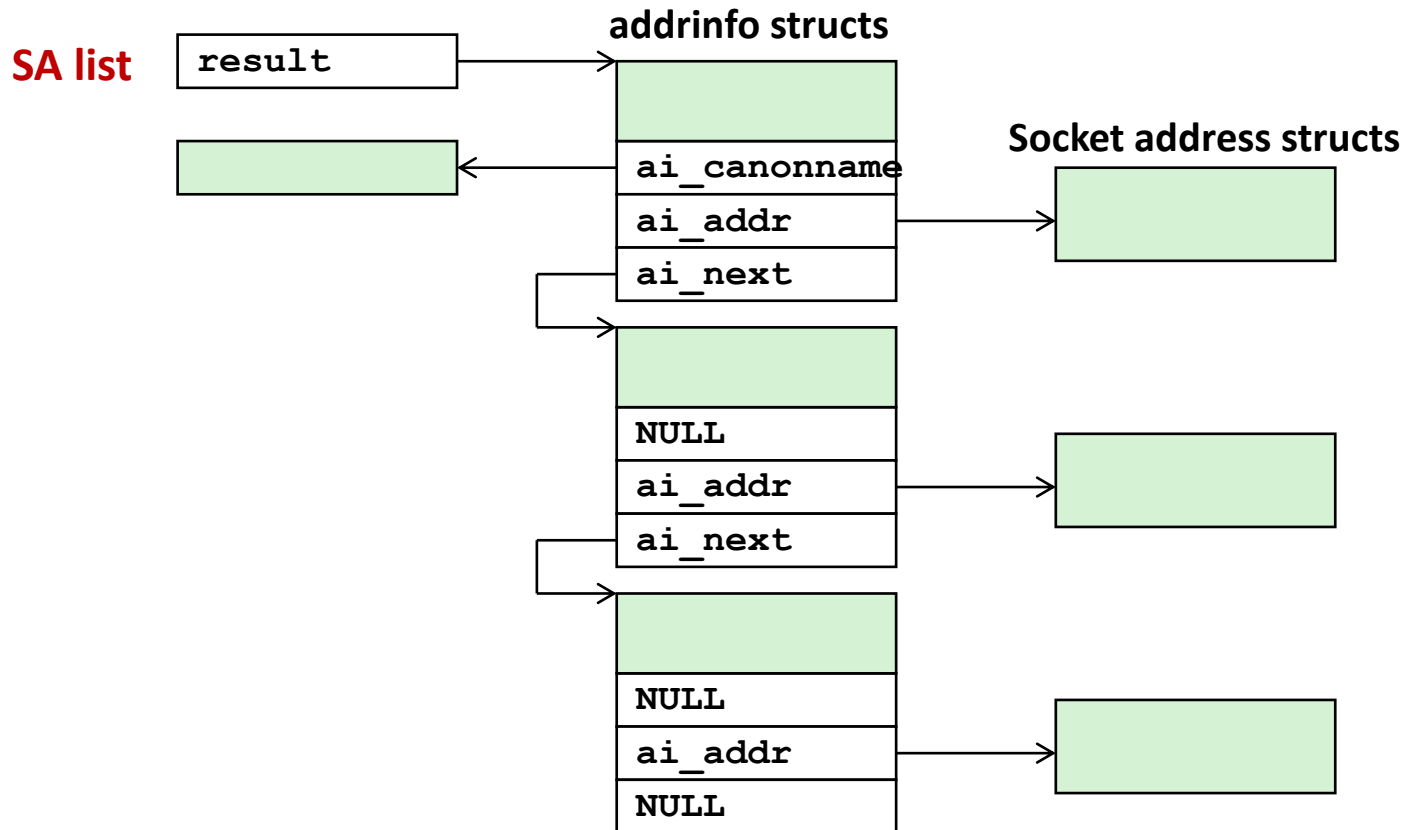


Sockets Interface

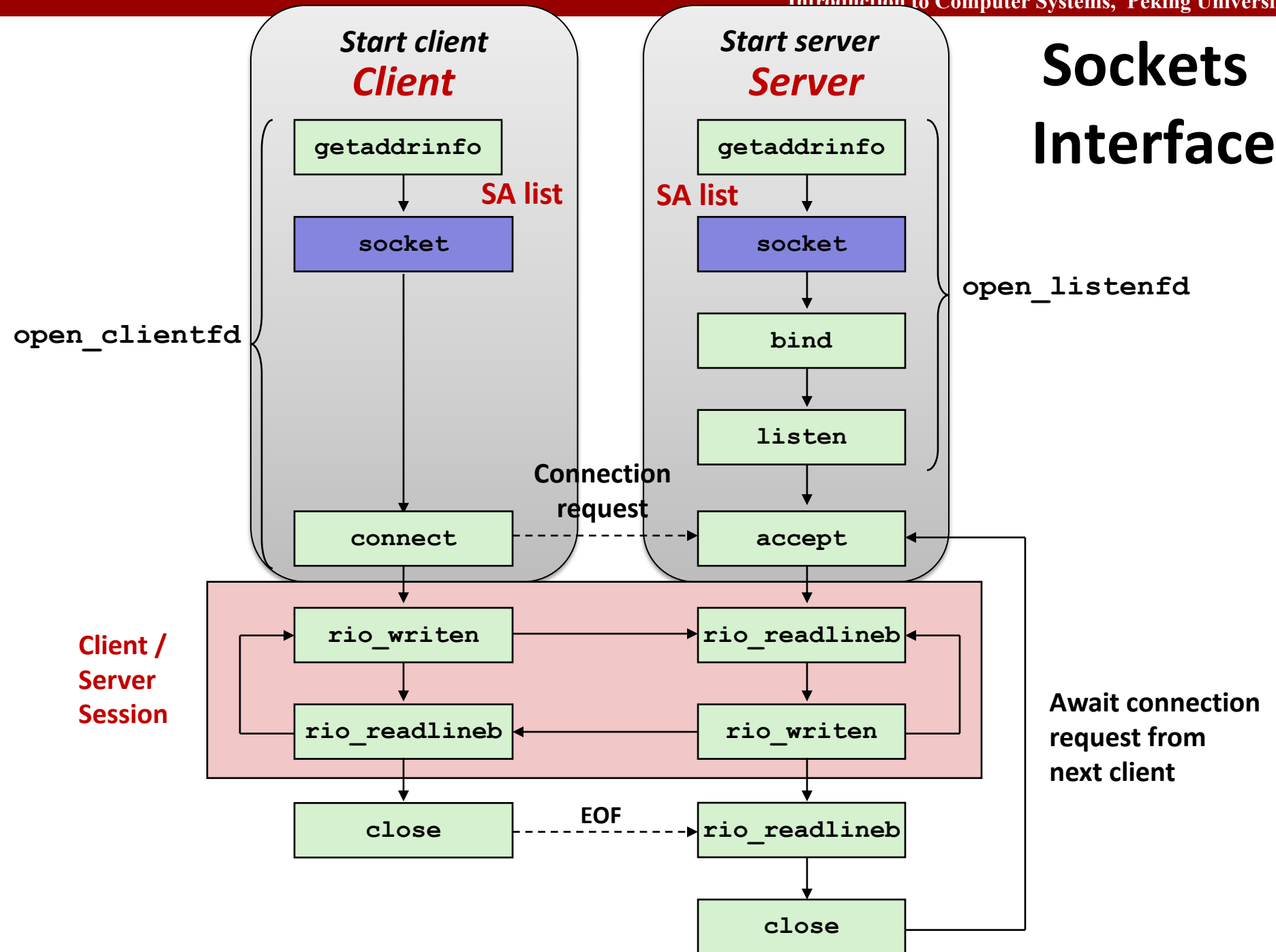


Review: getaddrinfo

- `getaddrinfo` converts string representations of hostnames, host addresses, ports, service names to socket address structures



Sockets Interface



Sockets Interface: `socket`

- Clients and servers use the `socket` function to create a *socket descriptor*:

```
int socket(int domain, int type, int protocol)
```

- Example:

```
int clientfd = socket(AF_INET, SOCK_STREAM, 0);
```

Protocol specific!

Indicates that we are using
32-bit IPV4 addresses

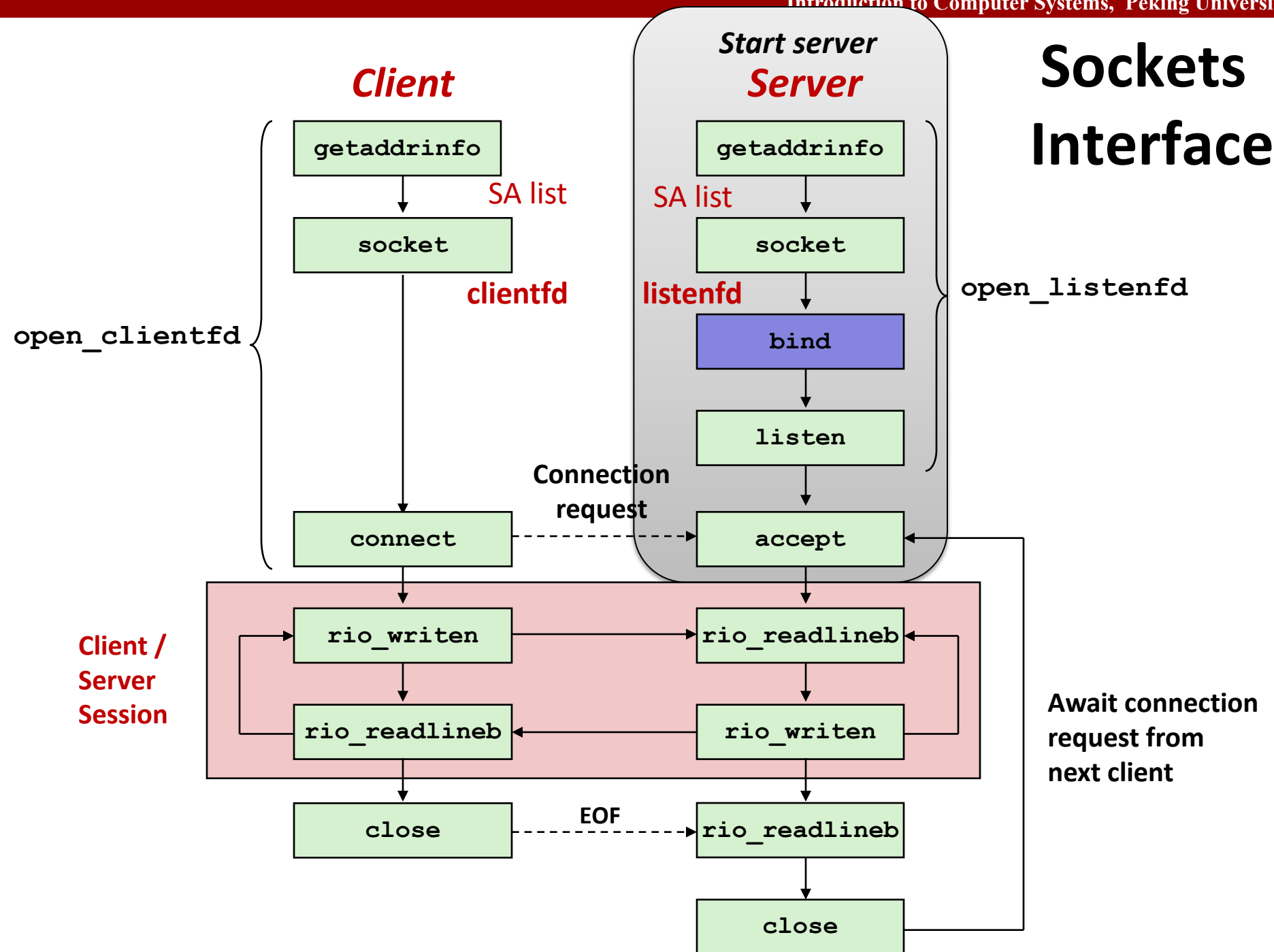
Indicates that the socket
will be the end point of a
reliable (TCP) connection

- Example:

```
int clientfd = socket(ai->ai_family, ai->ai_socktype,  
                    ai->ai_protocol);
```

*Use `getaddrinfo` to generate the parameters automatically,
so that code is protocol independent.*

Sockets Interface



Sockets Interface: `bind`

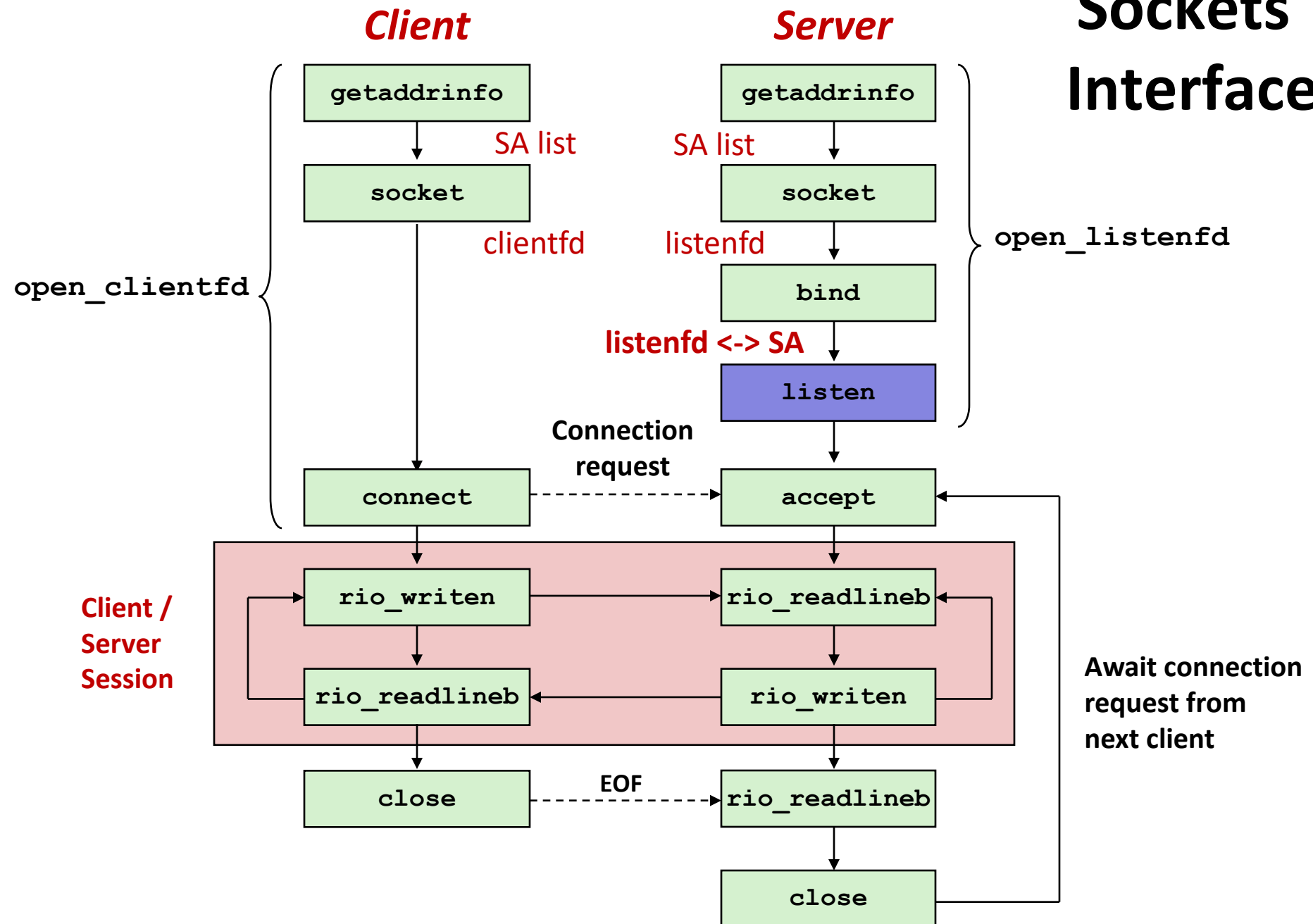
- A server uses `bind` to ask the kernel to associate the server's socket address with a socket descriptor:

```
int bind(int sockfd, SA *addr, socklen_t addrlen);
```

Our convention: `typedef struct sockaddr SA;`

- Process can read bytes that arrive on the connection whose endpoint is `addr` by reading from descriptor `sockfd`
- Similarly, writes to `sockfd` are transferred along connection whose endpoint is `addr`
- Best practice is to use `getaddrinfo` to supply the arguments `addr` and `addrlen`.

Sockets Interface



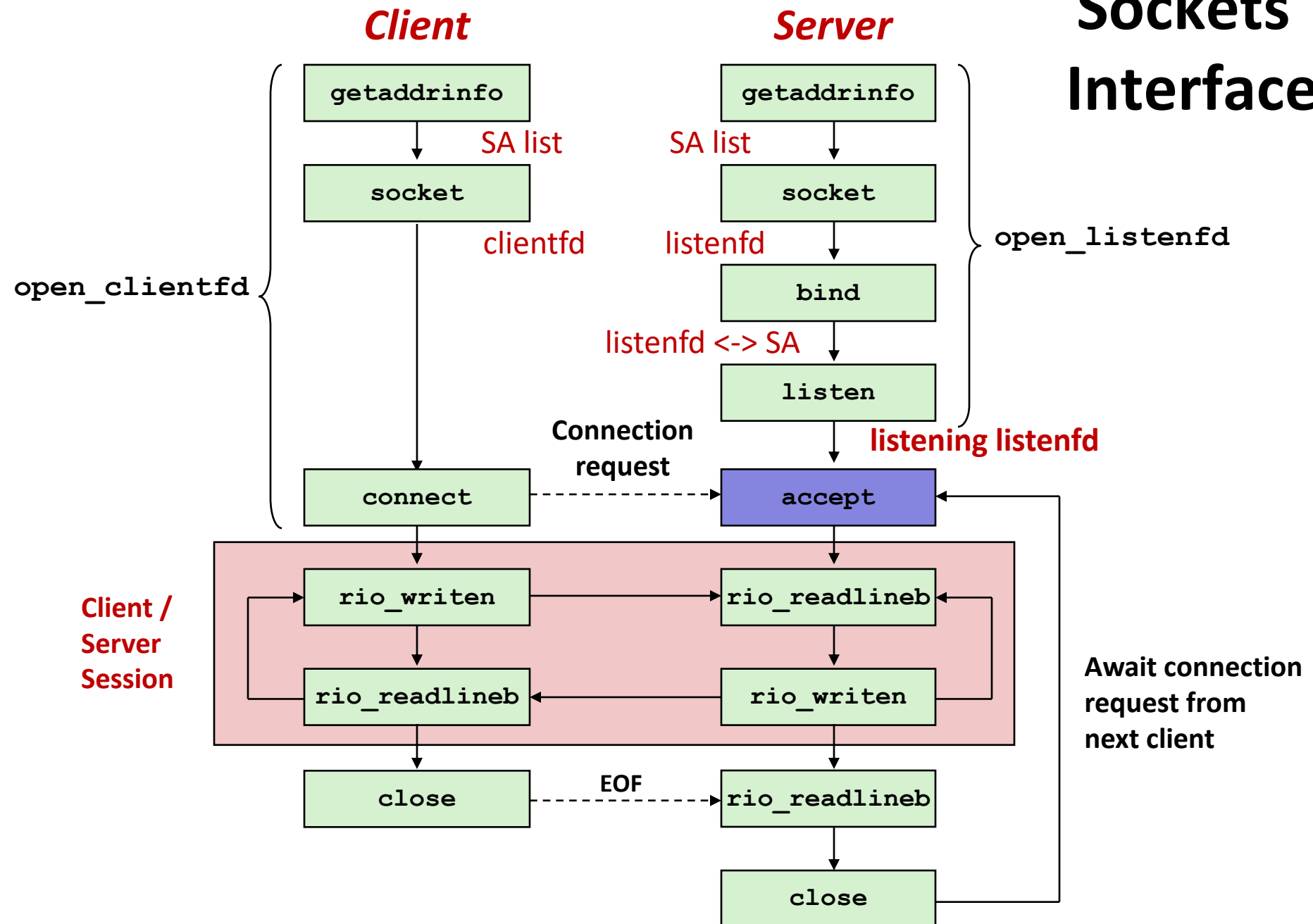
Sockets Interface: `listen`

- Kernel assumes that descriptor from `socket` function is an *active socket* that will be on the client end
- A server calls the `listen` function to tell the kernel that a descriptor will be used by a server rather than a client:

```
int listen(int sockfd, int backlog);
```

- Converts `sockfd` from an active socket to a *listening socket* that can accept connection requests from clients.
- `backlog` is a hint about the number of outstanding connection requests that the kernel should queue up before starting to refuse requests (128-ish by default)

Sockets Interface



Sockets Interface: `accept`

- Servers wait for connection requests from clients by calling `accept`:

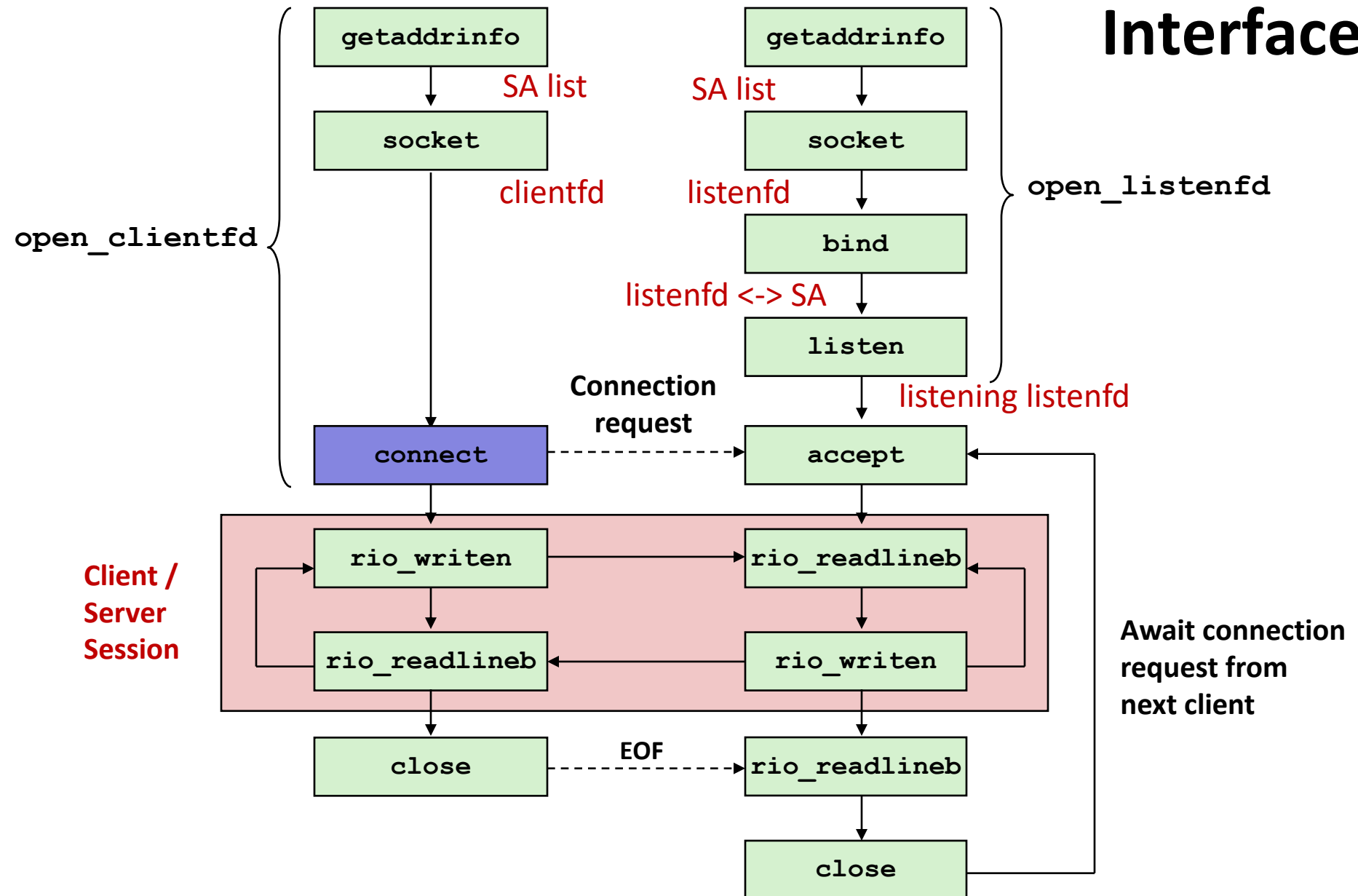
```
int accept(int listenfd, SA *addr, int *addrlen);
```

- Waits for connection request to arrive on the connection bound to `listenfd`, then fills in client's socket address in `addr` and size of the socket address in `addrlen`.
- Returns a *connected descriptor* `connfd` that can be used to communicate with the client via Unix I/O routines.

Sockets Interface

Client

Server



Sockets Interface: connect

- A client establishes a connection with a server by calling `connect`:

```
int connect(int clientfd, SA *addr, socklen_t addrlen);
```

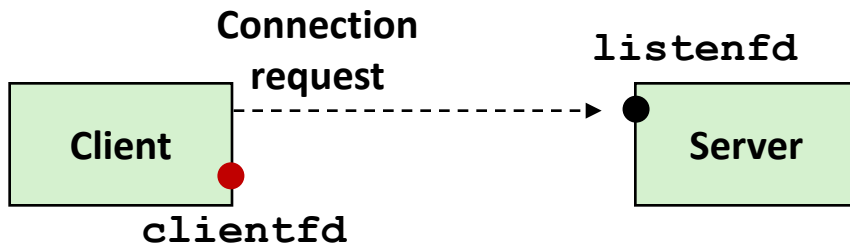
- Attempts to establish a connection with server at socket address `addr`
 - If successful, then `clientfd` is now ready for reading and writing.
 - Resulting connection is characterized by socket pair
(`x:y`, `addr.sin_addr:addr.sin_port`)
 - `x` is client address
 - `y` is ephemeral port that uniquely identifies client process on client host

Best practice is to use `getaddrinfo` to supply the arguments `addr` and `addrlen`.

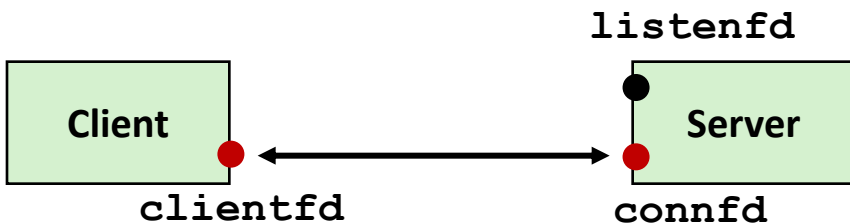
connect/accept Illustrated



1. Server blocks in `accept`, waiting for connection request on listening descriptor `listenfd`



2. Client makes connection request by calling and blocking in `connect`



3. Server returns `connfd` from `accept`. Client returns from `connect`. Connection is now established between `clientfd` and `connfd`

Connected vs. Listening Descriptors

■ Listening descriptor

- End point for client connection requests
- Created once and exists for lifetime of the server

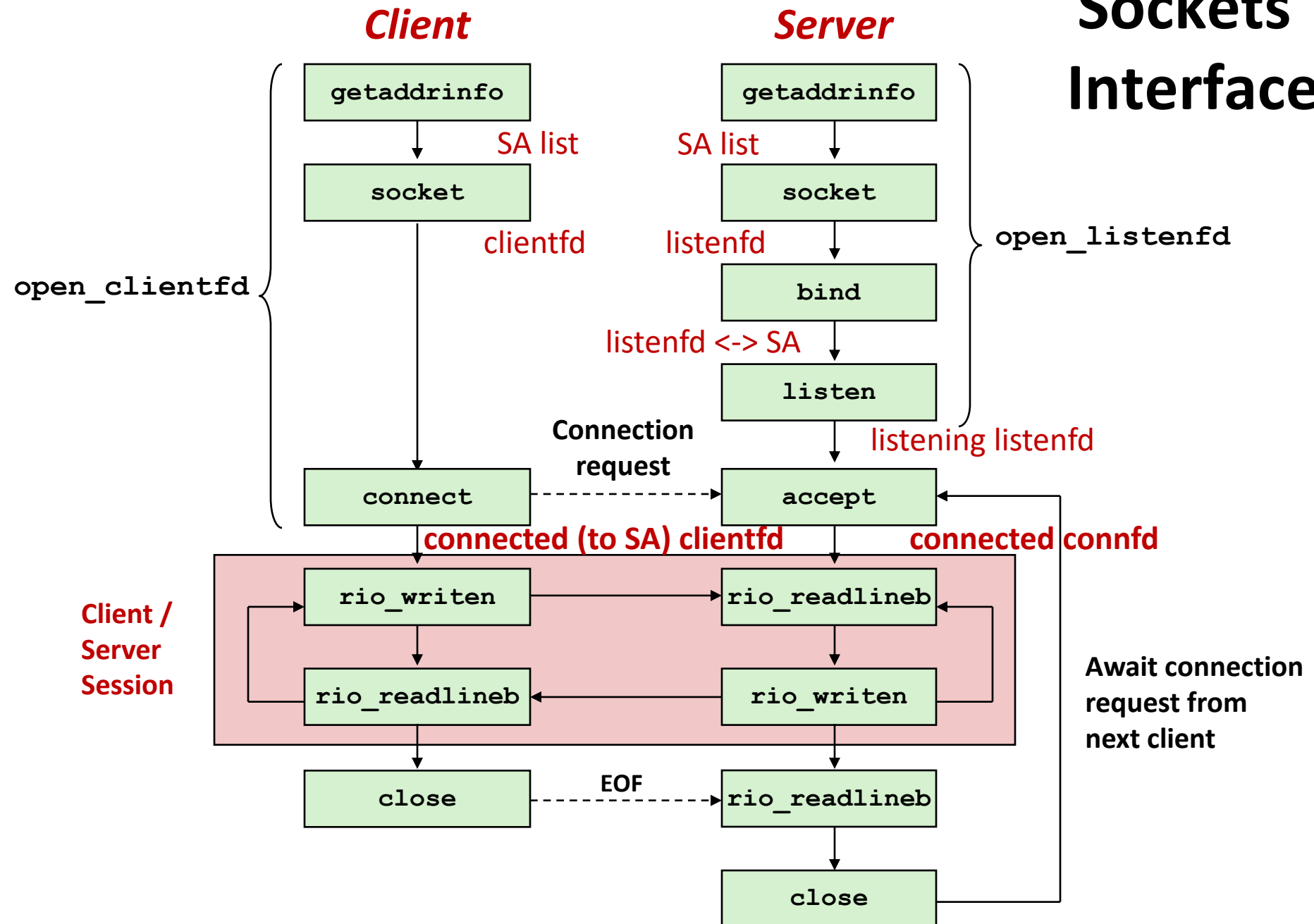
■ Connected descriptor

- End point of the connection between client and server
- A new descriptor is created each time the server accepts a connection request from a client
- Exists only as long as it takes to service client

■ Why the distinction?

- Allows for concurrent servers that can communicate over many client connections simultaneously
 - E.g., Each time we receive a new request, we fork a child to handle the request

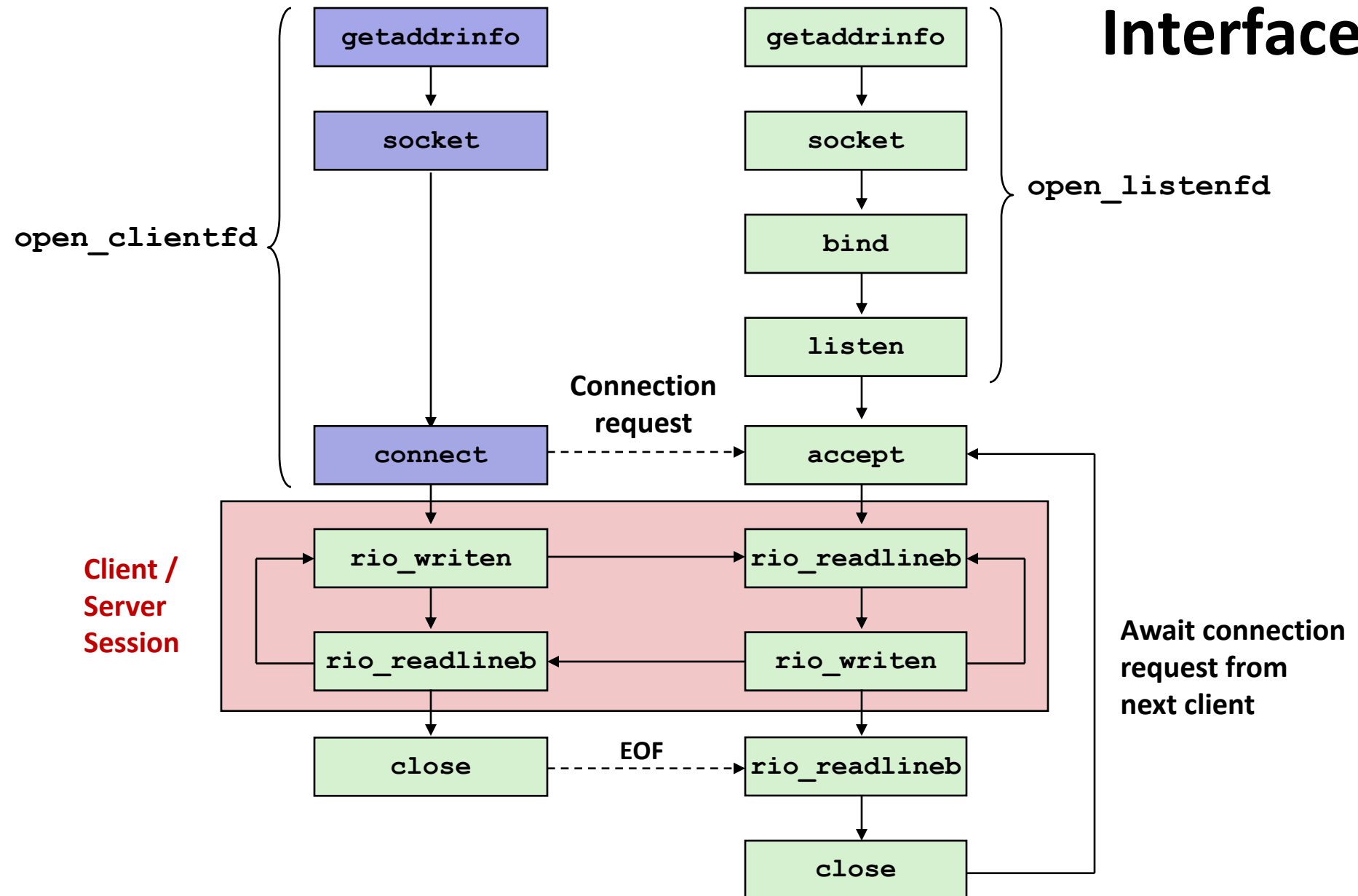
Sockets Interface



Sockets Interface

Client

Server



Sockets Helper: `open_clientfd`

■ Establish a connection with a server

```
int open_clientfd(char *hostname, char *port) {
    int clientfd;
    struct addrinfo hints, *listp, *p;

    /* Get a list of potential server addresses */
    memset(&hints, 0, sizeof(struct addrinfo));
    hints.ai_socktype = SOCK_STREAM; /* Open a connection */
    hints.ai_flags = AI_NUMERICSERV; /* ...using numeric port arg. */
    hints.ai_flags |= AI_ADDRCONFIG; /* Recommended for connections */
    Getaddrinfo(hostname, port, &hints, &listp);
```

csapp.c

AI_ADDRCONFIG means “use whichever of IPv4 and IPv6 works on this computer”. Good practice for clients, not for servers.

Sockets Helper: open_clientfd (cont)

```
/* Walk the list for one that we can successfully connect to */
for (p = listp; p; p = p->ai_next) {
    /* Create a socket descriptor */
    if ((clientfd = socket(p->ai_family, p->ai_socktype,
                          p->ai_protocol)) < 0)
        continue; /* Socket failed, try the next */

    /* Connect to the server */
    if (connect(clientfd, p->ai_addr, p->ai_addrlen) != -1)
        break; /* Success */
    Close(clientfd); /* Connect failed, try another */
}

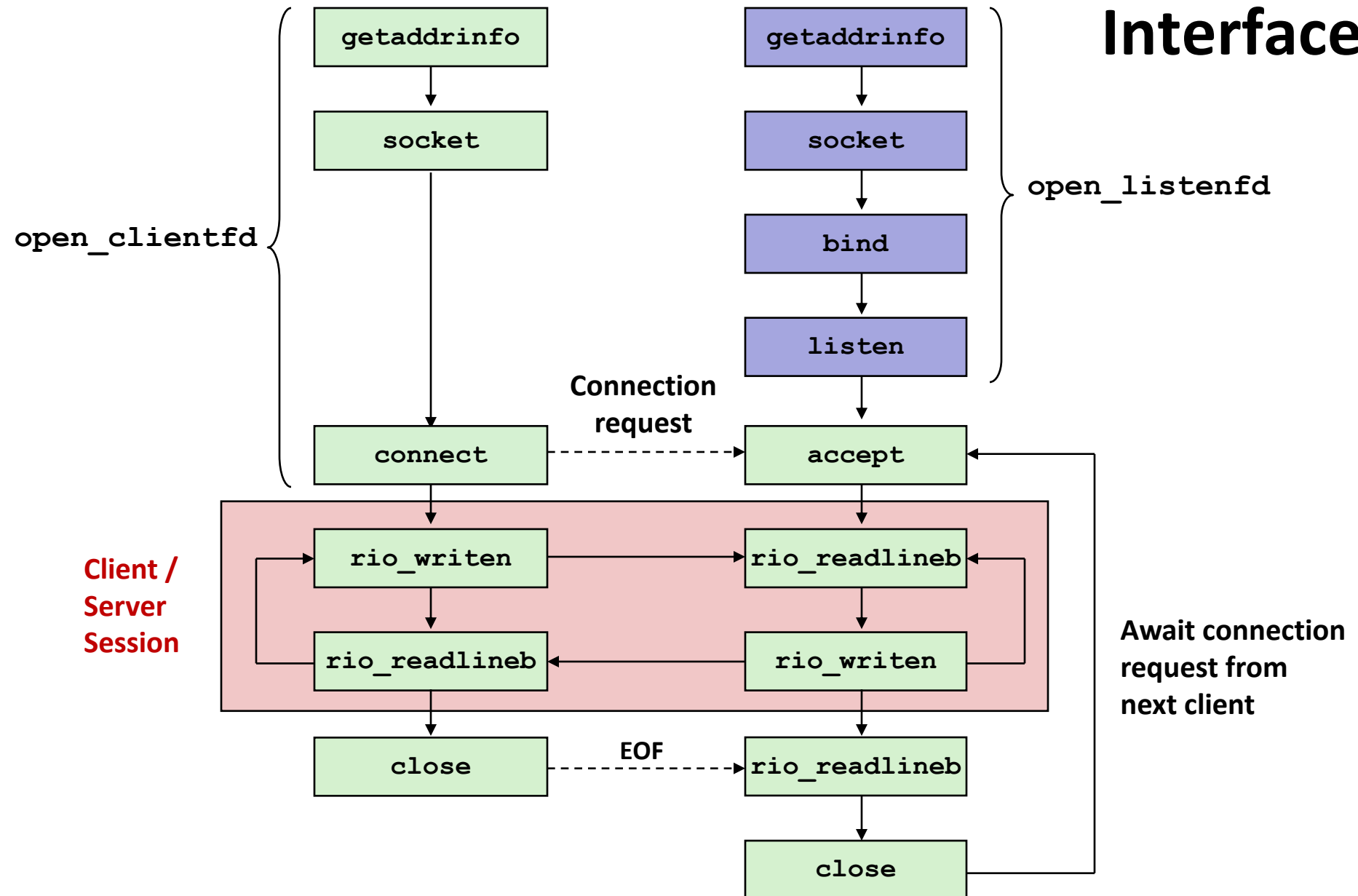
/* Clean up */
Freeaddrinfo(listp);
if (!p) /* All connects failed */
    return -1;
else /* The last connect succeeded */
    return clientfd;
}
```

csapp.c

Sockets Interface

Client

Server



Sockets Helper: `open_listenfd`

- Create a listening descriptor that can be used to accept connection requests from clients.

```
int open_listenfd(char *port)
{
    struct addrinfo hints, *listp, *p;
    int listenfd, optval=1;

    /* Get a list of potential server addresses */
    memset(&hints, 0, sizeof(struct addrinfo));
    hints.ai_socktype = SOCK_STREAM;                /* Accept connect. */
    hints.ai_flags = AI_PASSIVE | AI_ADDRCONFIG;    /* ...on any IP addr */
    hints.ai_flags |= AI_NUMERICSERV;              /* ...using port no. */
    Getaddrinfo(NULL, port, &hints, &listp);
}
```

csapp.c

`AI_PASSIVE` means “I plan to listen on this socket.”

`AI_ADDRCONFIG` normally not used for servers, but we use it for convenience

Sockets Helper: open_listenfd (cont)

```
/* Walk the list for one that we can bind to */
for (p = listp; p; p = p->ai_next) {
    /* Create a socket descriptor */
    if ((listenfd = socket(p->ai_family, p->ai_socktype,
                          p->ai_protocol)) < 0)
        continue; /* Socket failed, try the next */

    /* Eliminates "Address already in use" error from bind */
    Setsockopt(listenfd, SOL_SOCKET, SO_REUSEADDR,
               (const void *)&optval , sizeof(int));

    /* Bind the descriptor to the address */
    if (bind(listenfd, p->ai_addr, p->ai_addrlen) == 0)
        break; /* Success */
    Close(listenfd); /* Bind failed, try the next */
}
```

csapp.c

A production server would not break out of the loop on the first success. We do that for simplicity only.

Sockets Helper: `open_listenfd` (cont)

```
/* Clean up */
Freeaddrinfo(listp);
if (!p) /* No address worked */
    return -1;

/* Make it a listening socket ready to accept conn. requests */
if (listen(listenfd, LISTENQ) < 0) {
    Close(listenfd);
    return -1;
}
return listenfd;
}
```

csapp.c

- **Key point:** `open_clientfd` and `open_listenfd` are both independent of any particular version of IP.

Testing Servers Using `telnet`

- The `telnet` program is invaluable for testing servers that transmit ASCII strings over Internet connections
 - Our simple echo server
 - Web servers
 - Mail servers
- Usage:
 - `linux> telnet <host> <portnumber>`
 - Creates a connection with a server running on `<host>` and listening on port `<portnumber>`

Testing the Echo Server With telnet

```
whaleshark> ./echoserveri 15213
Connected to (MAKOSHARK.ICS.CS.CMU.EDU, 50280)
server received 11 bytes
server received 8 bytes

makoshark> telnet whaleshark.ics.cs.cmu.edu 15213
Trying 128.2.210.175...
Connected to whaleshark.ics.cs.cmu.edu (128.2.210.175) .
Escape character is '^]'.
Hi there!
Hi there!
Howdy!
Howdy!
^]
telnet> quit
Connection closed.
makoshark>
```

For More Information

- **W. Richard Stevens et. al. “Unix Network Programming: The Sockets Networking API”, Volume 1, Third Edition, Prentice Hall, 2003**
 - THE network programming bible.
- **Michael Kerrisk, “The Linux Programming Interface”, No Starch Press, 2010**
 - THE Linux programming bible.
- **Complete versions of all code in this lecture is available from the 213 schedule page.**
 - `http://www.cs.cmu.edu/~213/schedule.html`
 - `csapp.{c,h}`, `hostinfo.c`, `echoclient.c`, `echoserveri.c`, `tiny.c`, `adder.c`
 - You can use any of this code in your assignments.

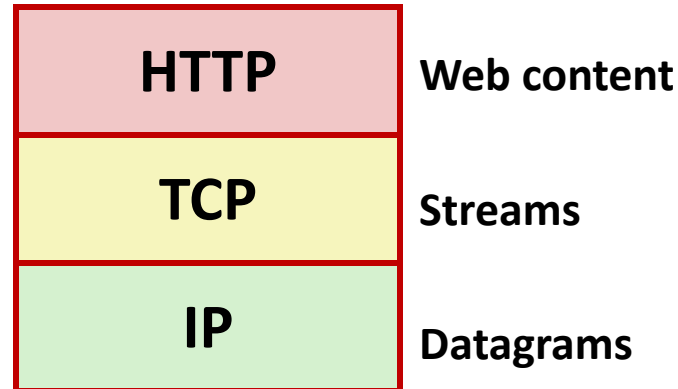
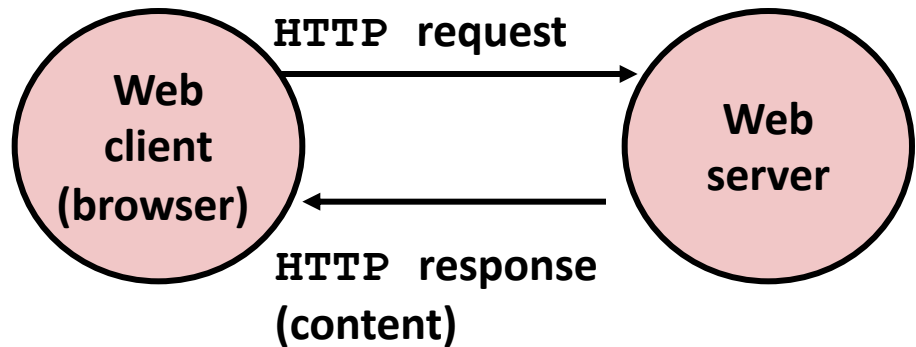
Additional slides

Key Layers of the Internet

early milestones	Key Layers of the Internet	milestones
email@-1971 Ray Tomlinson	CONTENT	1987-HyperCard Bill Atkinson
Archie-1990 Emtage & Deutsch	SEARCH ENGINE*	1998-Google Brin & Page
DOS Houdini-1986 Neil Larson	BROWSERS	1993-Mosaic Marc Andreessen
(Vannevar Bush, Ted Nelson, Douglas Engelbart)	WORLD WIDE WEB	1990-http:// Tim Berners-Lee
ARPANET-1969 J.C.R. Licklider	INTERNET	1975-TCP/IP Cerf & Kahn
SAGE-1956 George Valley	NETWORKS	1973-Ethernet Robert Metcalfe
Z3-1941 Konrad Zuse	COMPUTERS	1976-Apple Jobs & Wozniak

Web Server Basics

- **Clients and servers communicate using the HyperText Transfer Protocol (HTTP)**
 - Client and server establish TCP connection
 - Client requests content
 - Server responds with requested content
 - Client and server close connection (eventually)
- **Current version is HTTP/1.1**
 - RFC 2616, June, 1999.



`http://www.w3.org/Protocols/rfc2616/rfc2616.html`

Web Content

■ Web servers return *content* to clients

- *content*: a sequence of bytes with an associated MIME (Multipurpose Internet Mail Extensions) type

■ Example MIME types

- | | |
|---------------------------|-------------------------------------|
| ■ <code>text/html</code> | HTML document |
| ■ <code>text/plain</code> | Unformatted text |
| ■ <code>image/gif</code> | Binary image encoded in GIF format |
| ■ <code>image/png</code> | Binary image encoded in PNG format |
| ■ <code>image/jpeg</code> | Binary image encoded in JPEG format |

You can find the complete list of MIME types at:

<http://www.iana.org/assignments/media-types/media-types.xhtml>

Static and Dynamic Content

- The content returned in HTTP responses can be either *static* or *dynamic*
 - *Static content*: content stored in files and retrieved in response to an HTTP request
 - Examples: HTML files, images, audio clips, Javascript programs
 - Request identifies which content file
 - *Dynamic content*: content produced on-the-fly in response to an HTTP request
 - Example: content produced by a program executed by the server on behalf of the client
 - Request identifies file containing executable code
- ***Web content associated with a file that is managed by the server***

URLs and how clients and servers use them

- Unique name for a file: URL (Universal Resource Locator)
- Example URL: `http://www.cmu.edu:80/index.html`
- Clients use *prefix* (`http://www.cmu.edu:80`) to infer:
 - What kind (protocol) of server to contact (HTTP)
 - Where the server is (`www.cmu.edu`)
 - What port it is listening on (80)
- Servers use *suffix* (`/index.html`) to:
 - Determine if request is for static or dynamic content.
 - No hard and fast rules for this
 - One convention: executables reside in `cgi-bin` directory
 - Find file on file system
 - Initial “/” in suffix denotes home directory for requested content.
 - Minimal suffix is “/”, which server expands to configured default filename (usually, `index.html`)

HTTP Requests

- HTTP request is a *request line*, followed by zero or more *request headers*
- Request line: `<method> <uri> <version>`
 - `<method>` is one of GET, POST, OPTIONS, HEAD, PUT, DELETE, or TRACE
 - `<uri>` is typically URL for proxies, URL suffix for servers
 - A URL is a type of URI (Uniform Resource Identifier)
 - See <http://www.ietf.org/rfc/rfc2396.txt>
 - `<version>` is HTTP version of request (HTTP/1.0 or HTTP/1.1)
- Request headers: `<header name>: <header data>`
 - Provide additional information to the server

HTTP Responses

- HTTP response is a *response line* followed by zero or more *response headers*, possibly followed by *content*, with blank line (“\r\n”) separating headers from content.
- Response line:
 - `<version> <status code> <status msg>`
 - `<version>` is HTTP version of the response
 - `<status code>` is numeric status
 - `<status msg>` is corresponding English text
 - 200 OK Request was handled without error
 - 301 Moved Provide alternate URL
 - 404 Not found Server couldn't find the file
- Response headers: `<header name>: <header data>`
 - Provide additional information about response
 - **Content-Type**: MIME type of content in response body
 - **Content-Length**: Length of content in response body

Example HTTP Transaction

whaleshark> telnet www.cmu.edu 80	Client: open connection to server
Trying 128.2.42.52...	Telnet prints 3 lines to terminal
Connected to WWW-CMU-PROD-VIP.ANDREW.cmu.edu.	
Escape character is '^['.	
GET / HTTP/1.1	Client: request line
Host: www.cmu.edu	Client: required HTTP/1.1 header
	Client: blank line terminates headers
HTTP/1.1 301 Moved Permanently	Server: response line
Date: Wed, 05 Nov 2014 17:05:11 GMT	Server: followed by 5 response headers
Server: Apache/1.3.42 (Unix)	Server: this is an Apache server
Location: http://www.cmu.edu/index.shtml	Server: page has moved here
Transfer-Encoding: chunked	Server: response body will be chunked
Content-Type: text/html; charset=...	Server: expect HTML in response body
	Server: empty line terminates headers
15c	Server: first line in response body
<HTML><HEAD>	Server: start of HTML content
...	
</BODY></HTML>	Server: end of HTML content
0	Server: last line in response body
Connection closed by foreign host.	Server: closes connection

- HTTP standard requires that each text line end with “\r\n”
- Blank line (“\r\n”) terminates request and response headers

Example HTTP Transaction, Take 2

```

whaleshark> telnet www.cmu.edu 80
Trying 128.2.42.52...
Connected to WWW-CMU-PROD-VIP.ANDREW.cmu.edu.
Escape character is '^]'.
GET /index.shtml HTTP/1.1
Host: www.cmu.edu

HTTP/1.1 200 OK
Date: Wed, 05 Nov 2014 17:37:26 GMT
Server: Apache/1.3.42 (Unix)
Transfer-Encoding: chunked
Content-Type: text/html; charset=...

1000
<html ..>
...
</html>
0
Connection closed by foreign host.

```

Client: open connection to server
 Telnet prints 3 lines to terminal

Client: request line
 Client: required HTTP/1.1 header
 Client: blank line terminates headers
 Server: response line
 Server: followed by 4 response headers

Server: empty line terminates headers
 Server: begin response body
 Server: first line of HTML content

Server: end response body
 Server: close connection

Example HTTP(S) Transaction, Take 3

```
whaleshark> openssl s_client www.cs.cmu.edu:443
CONNECTED(00000005)
...
Certificate chain
...
-
Server certificate
-----BEGIN CERTIFICATE-----
MIIGDjCCBPagAwIBAgIRAMiF7LBPDoYSilnNoU+mp+gwDQYJKoZIhvcNAQELBQAw
djELMAkGA1UEBhMCVVMxCzAJBgNVBAGTAk1JMRIwEAYDVQQHEwlBbm4gQXJib3Ix
EjAQBgNVBAoTCUluudGVybmV0MjERMA8GA1UECzMISW5Db21tb24xHzAdBgNVBAMT
wkWkvDVBBCwKXrShVxQNsJ6J
...
-----END CERTIFICATE-----
subject=/C=US/postalCode=15213/ST=PA/L=Pittsburgh/street=5000 Forbes
Ave/O=Carnegie Mellon University/OU=School of Computer
Science/CN=www.cs.cmu.edu issuer=/C=US/ST=MI/L=Ann
Arbor/O=Internet2/OU=InCommon/CN=InCommon RSA Server CA
SSL handshake has read 6274 bytes and written 483 bytes
...
>GET / HTTP/1.0

HTTP/1.1 200 OK
Date: Tue, 12 Nov 2019 04:22:15 GMT
Server: Apache/2.4.10 (Ubuntu)
Set-Cookie: SHIBLOCATION=scsweb; path=/; domain=.cs.cmu.edu
... HTML Content Continues Below ...
```

Tiny Web Server

■ Tiny Web server described in text

- Tiny is a sequential Web server
- Serves static and dynamic content to real browsers
 - text files, HTML files, GIF, PNG, and JPEG images
- 239 lines of commented C code
- Not as complete or robust as a real Web server
 - You can break it with poorly-formed HTTP requests (e.g., terminate lines with “\n” instead of “\r\n”)

Tiny Operation

- **Accept connection from client**
- **Read request from client (via connected socket)**
- **Split into <method> <uri> <version>**
 - If method not GET, then return error
- **If URI contains “cgi-bin” then serve dynamic content**
 - (Would do wrong thing if had file “abcgi-bingo.html”)
 - Fork process to execute program
- **Otherwise serve static content**
 - Copy file to output

Tiny Serving Static Content

```
void serve_static(int fd, char *filename, int filesize)
{
    int srcfd;
    char *srcp, filetype[MAXLINE], buf[MAXBUF];

    /* Send response headers to client */
    get_filetype(filename, filetype);
    sprintf(buf, "HTTP/1.0 200 OK\r\n");
    sprintf(buf, "%sServer: Tiny Web Server\r\n", buf);
    sprintf(buf, "%sConnection: close\r\n", buf);
    sprintf(buf, "%sContent-length: %d\r\n", buf, filesize);
    sprintf(buf, "%sContent-type: %s\r\n\r\n", buf, filetype);
    Rio_writen(fd, buf, strlen(buf));

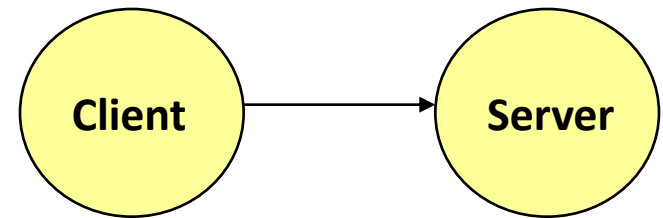
    /* Send response body to client */
    srcfd = Open(filename, O_RDONLY, 0);
    srcp = Mmap(0, filesize, PROT_READ, MAP_PRIVATE, srcfd, 0);
    Close(srcfd);
    Rio_writen(fd, srcp, filesize);
    Munmap(srcp, filesize);
}
```

tiny.c

Serving Dynamic Content

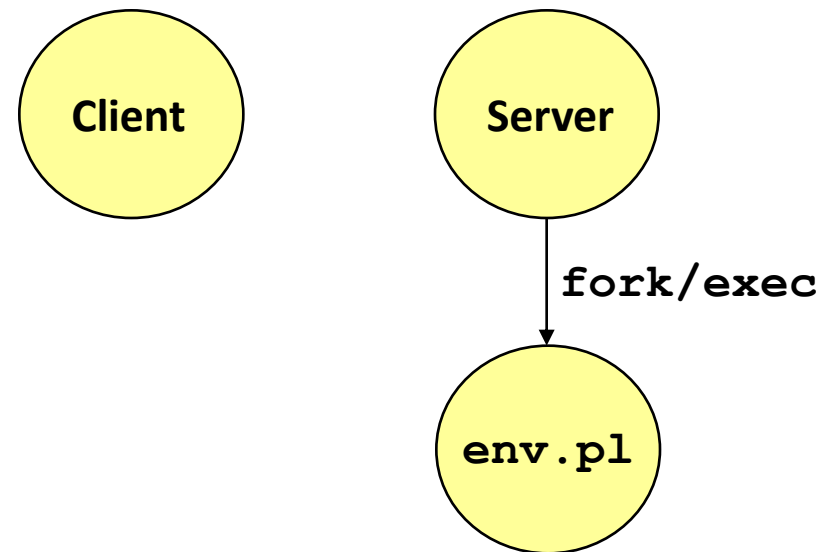
- Client sends request to server
- If request URI contains the string `/cgi-bin`, the Tiny server assumes that the request is for dynamic content

`GET /cgi-bin/env.pl HTTP/1.1`



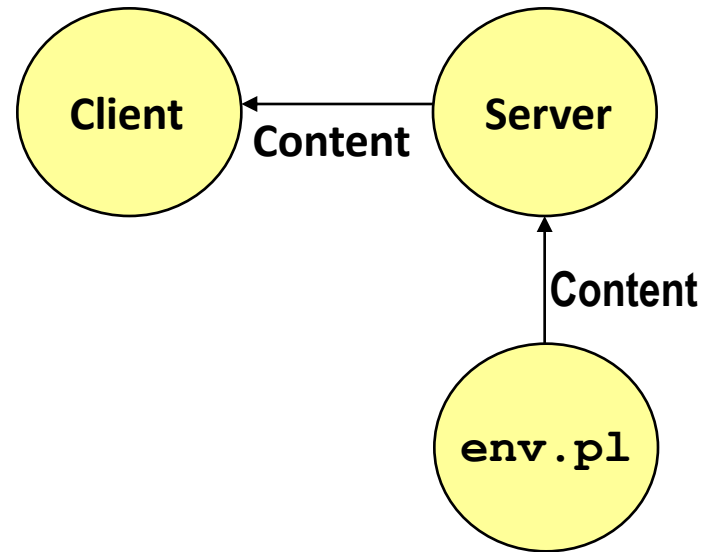
Serving Dynamic Content (cont)

- The server creates a child process and runs the program identified by the URI in that process



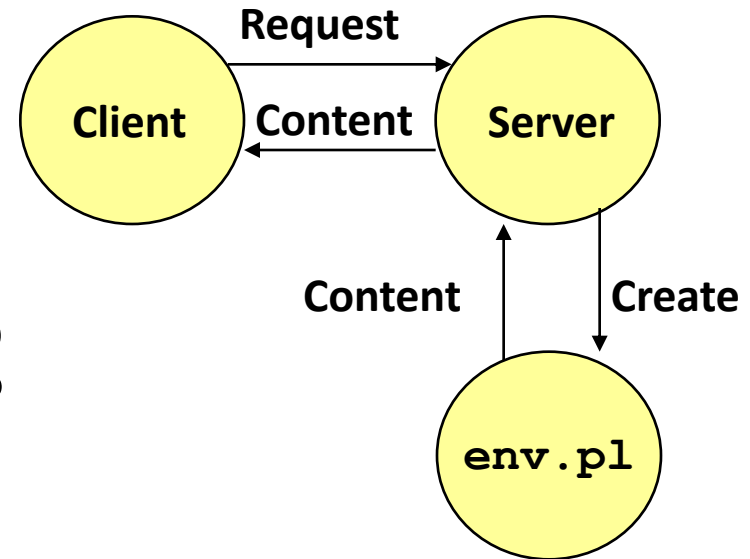
Serving Dynamic Content (cont)

- The child runs and generates the dynamic content
- The server captures the content of the child and forwards it without modification to the client



Issues in Serving Dynamic Content

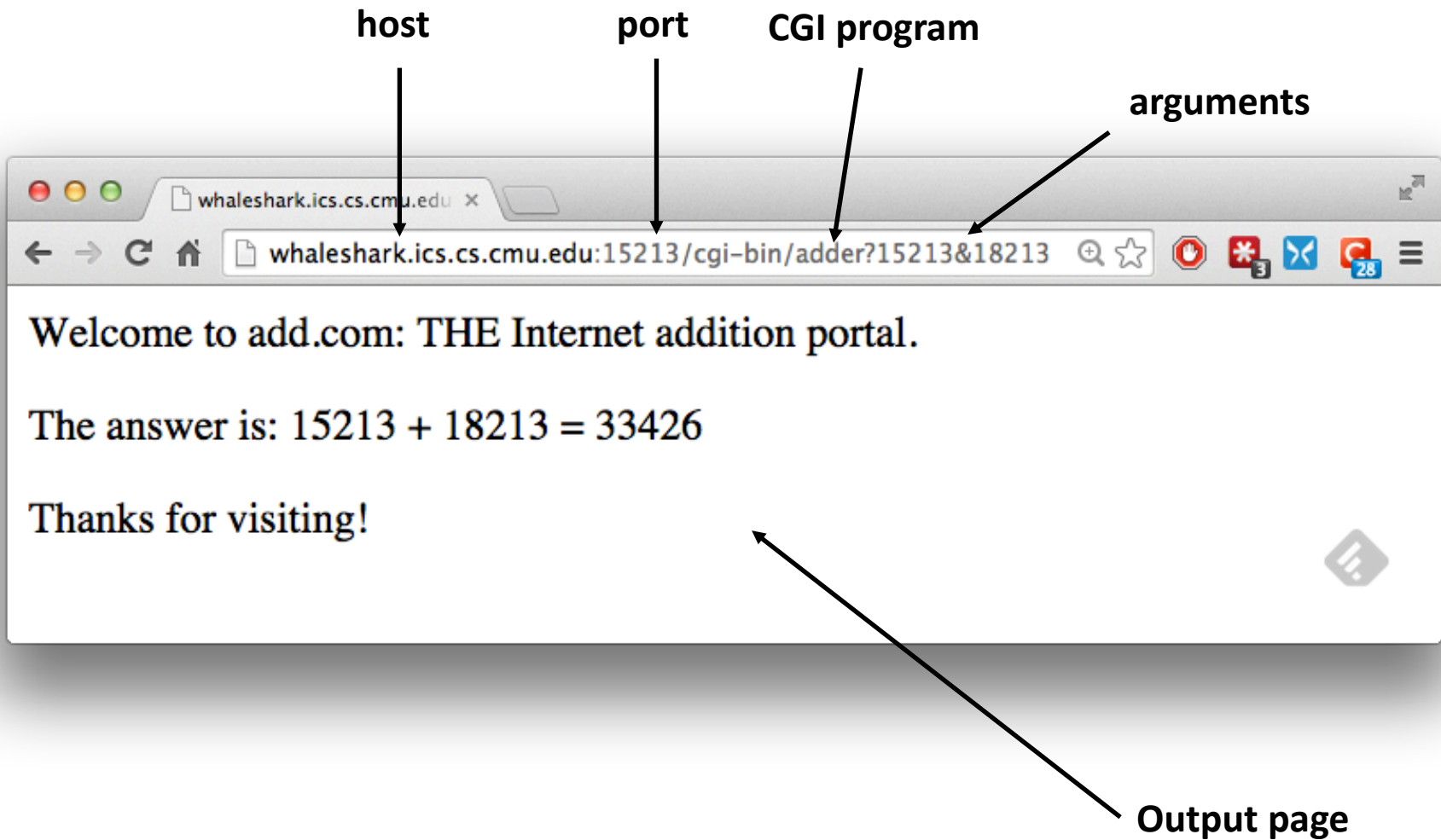
- How does the client pass program arguments to the server?
- How does the server pass these arguments to the child?
- How does the server pass other info relevant to the request to the child?
- How does the server capture the content produced by the child?
- These issues are addressed by the **Common Gateway Interface (CGI)** specification.



CGI

- Because the children are written according to the CGI spec, they are often called *CGI programs*.
- However, CGI really defines a simple standard for transferring information between the client (browser), the server, and the child process.
- CGI is the original standard for generating dynamic content. Has been largely replaced by other, faster techniques:
 - E.g., fastCGI, Apache modules, Java servlets, Rails controllers
 - Avoid having to create process on the fly (expensive and slow).

The add.com Experience



Serving Dynamic Content With GET

- Question: How does the client pass arguments to the server?
- Answer: The arguments are appended to the URI
- Can be encoded directly in a URL typed to a browser or a URL in an HTML link
 - `http://add.com/cgi-bin/adder?15213&18213`
 - **adder** is the CGI program on the server that will do the addition.
 - argument list starts with “?”
 - arguments separated by “&”
 - spaces represented by “+” or “%20”

Serving Dynamic Content With GET

- URL suffix:
 - `cgi-bin/adder?15213&18213`

- Result displayed on browser:

```
Welcome to add.com: THE Internet  
addition portal.
```

```
The answer is: 15213 + 18213 = 33426
```

```
Thanks for visiting!
```

Serving Dynamic Content With GET

- **Question**: How does the server pass these arguments to the child?
- **Answer**: In environment variable `QUERY_STRING`
 - A single string containing everything after the “?”
 - For add: `QUERY_STRING = “15213&18213”`

```
/* Extract the two arguments */  
if ((buf = getenv("QUERY_STRING")) != NULL) {  
    p = strchr(buf, '&');  
    *p = '\0';  
    strcpy(arg1, buf);  
    strcpy(arg2, p+1);  
    n1 = atoi(arg1);  
    n2 = atoi(arg2);  
}
```

adder.c

Serving Dynamic Content with GET

- Question: How does the server capture the content produced by the child?
- Answer: The child generates its output on `stdout`. Server uses `dup2` to redirect `stdout` to its connected socket.

```
void serve_dynamic(int fd, char *filename, char *cgiargs)
{
    char buf[MAXLINE], *emptylist[] = { NULL };

    /* Return first part of HTTP response */
    sprintf(buf, "HTTP/1.0 200 OK\r\n");
    Rio_writen(fd, buf, strlen(buf));
    sprintf(buf, "Server: Tiny Web Server\r\n");
    Rio_writen(fd, buf, strlen(buf));

    if (Fork() == 0) { /* Child */
        /* Real server would set all CGI vars here */
        setenv("QUERY_STRING", cgiargs, 1);
        Dup2(fd, STDOUT_FILENO); /* Redirect stdout to client */
        Execve(filename, emptylist, environ); /* Run CGI program */
    }
    Wait(NULL); /* Parent waits for and reaps child */
}
```


Serving Dynamic Content with GET

- Notice that only the CGI child process knows the content type and length, so it must generate those headers.

```
/* Make the response body */
sprintf(content, "Welcome to add.com: ");
sprintf(content, "%sTHE Internet addition portal.\r\n<p>", content);
sprintf(content, "%sThe answer is: %d + %d = %d\r\n<p>",
        content, n1, n2, n1 + n2);
sprintf(content, "%sThanks for visiting!\r\n", content);

/* Generate the HTTP response */
printf("Content-length: %d\r\n", (int)strlen(content));
printf("Content-type: text/html\r\n\r\n");
printf("%s", content);
fflush(stdout);

exit(0);
```

adder.c

Serving Dynamic Content With GET

```
bash:makoshark> telnet whaleshark.ics.cs.cmu.edu 15213
Trying 128.2.210.175...
Connected to whaleshark.ics.cs.cmu.edu (128.2.210.175) .
Escape character is '^]'.
GET /cgi-bin/adder?15213&18213 HTTP/1.0
```

HTTP request sent by client

```
HTTP/1.0 200 OK
Server: Tiny Web Server
Connection: close
Content-length: 117
Content-type: text/html
```

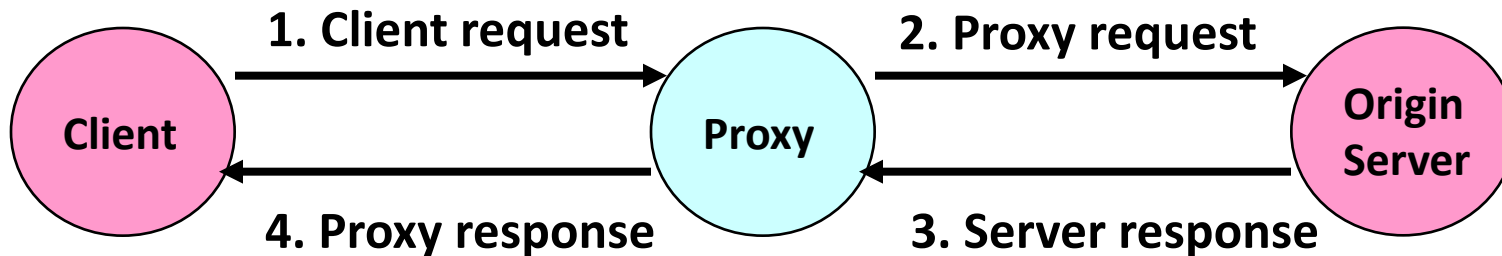
*HTTP response generated
by the server*

```
Welcome to add.com: THE Internet addition portal.
<p>The answer is: 15213 + 18213 = 33426
<p>Thanks for visiting!
Connection closed by foreign host.
bash:makoshark>
```

*HTTP response generated
by the CGI program*

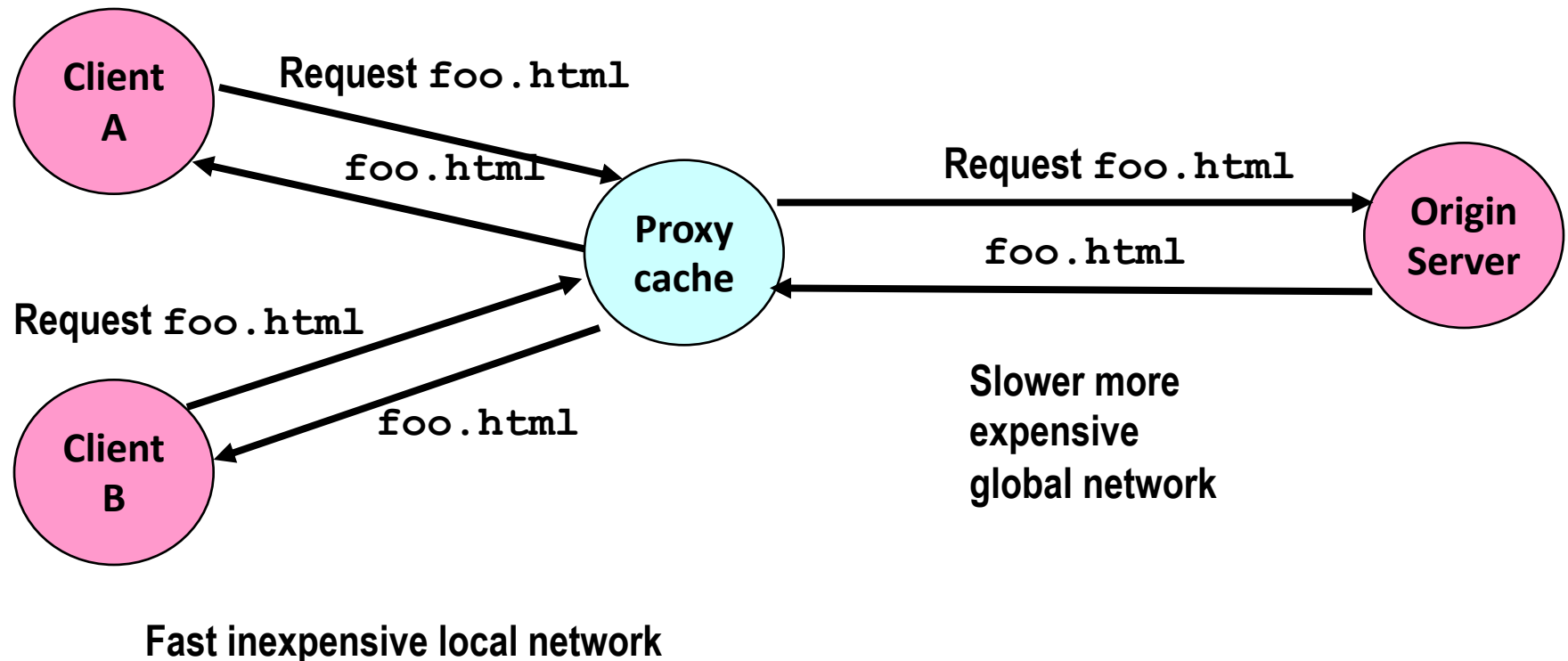
Proxies

- A **proxy** is an intermediary between a client and an **origin server**
 - To the client, the proxy acts like a server
 - To the server, the proxy acts like a client



Why Proxies?

- Can perform useful functions as requests and responses pass by
 - Examples: Caching, logging, anonymization, filtering, transcoding



Evolution of Internet

■ Original Idea

- Every node on Internet would have unique IP address
 - Everyone would be able to talk directly to everyone
- No secrecy or authentication
 - Messages visible to routers and hosts on same LAN
 - Possible to forge source field in packet header

■ Shortcomings

- There aren't enough IP addresses available
- Don't want everyone to have access or knowledge of all other hosts
- Security issues mandate secrecy & authentication

Evolution of Internet: Naming

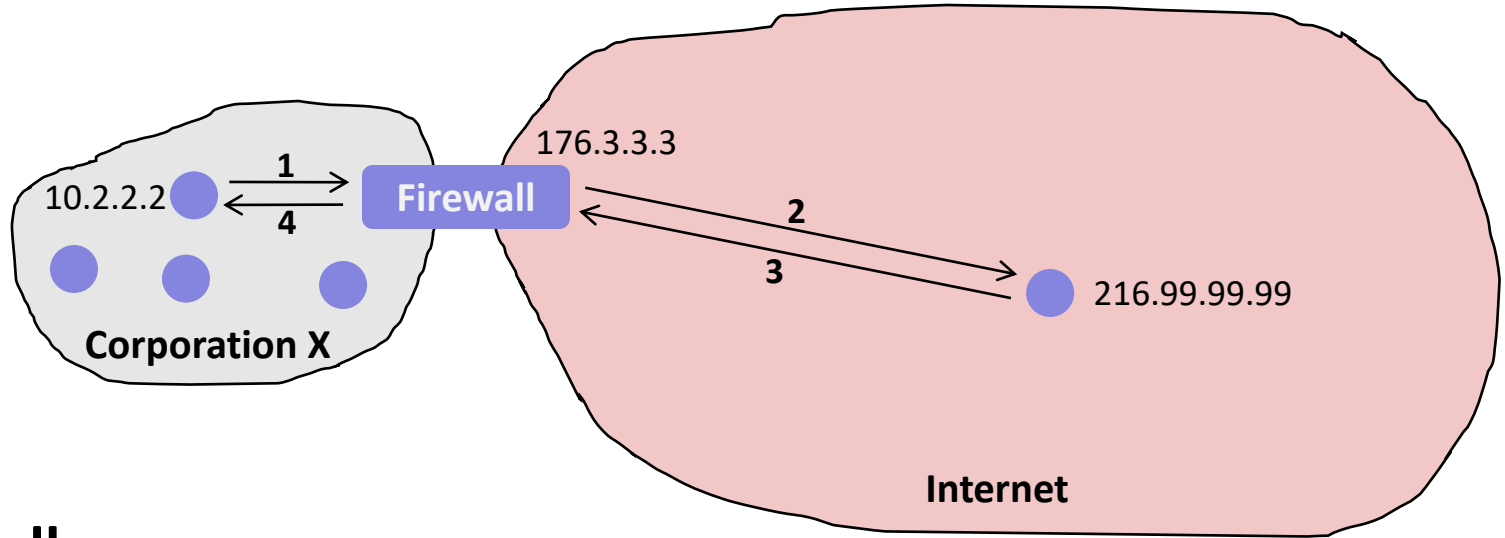
■ Dynamic address assignment

- Most hosts don't need to have known address
 - Only those functioning as servers
- DHCP (Dynamic Host Configuration Protocol)
 - Local ISP assigns address for temporary use

■ Example:

- Laptop at CMU (wired connection)
 - IP address 128.2.213.29 (**bryant-tp4.cs.cmu.edu**)
 - Assigned statically
- Laptop at home
 - IP address 192.168.1.5
 - Only valid within home network

Evolution of Internet: Firewalls



■ Firewalls

- Hides organizations nodes from rest of Internet
- Use local IP addresses within organization
- For external service, provides proxy service
 1. Client request: src=10.2.2.2, dest=216.99.99.99
 2. Firewall forwards: src=176.3.3.3, dest=216.99.99.99
 3. Server responds: src=216.99.99.99, dest=176.3.3.3
 4. Firewall forwards response: src=216.99.99.99, dest=10.2.2.2