

## Week2

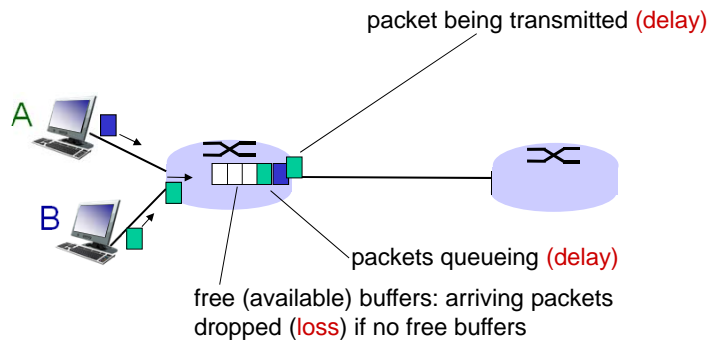
### *Delay, loss, and throughput in networks*

#### *Addresses*

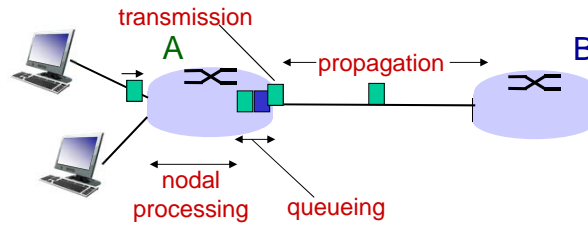
## How do loss and delay occur?

packets *queue* in router buffers

- ❖ packet arrival rate to link (temporarily) exceeds output link capacity
- ❖ packets queue, wait for turn



## Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

### $d_{\text{proc}}$ : nodal processing

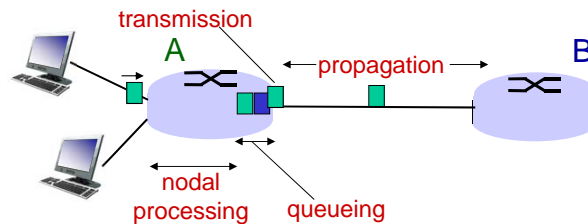
- check bit errors
- determine output link
- typically < msec

### $d_{\text{queue}}$ : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

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## Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

### $d_{\text{trans}}$ : transmission delay:

- $L$ : packet length (bits)
- $R$ : link bandwidth (bps)
- $d_{\text{trans}} = L/R$

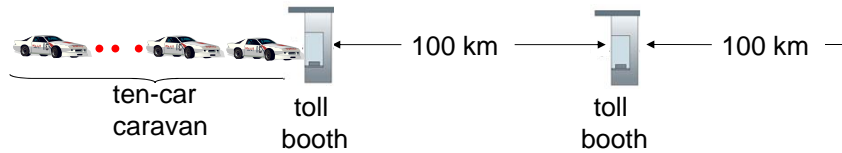
### $d_{\text{prop}}$ : propagation delay:

- $d$ : length of physical link
- $s$ : propagation speed in medium ( $\sim 2 \times 10^8$  m/sec)
- $d_{\text{prop}} = d/s$

$d_{\text{trans}}$  and  $d_{\text{prop}}$   
very different

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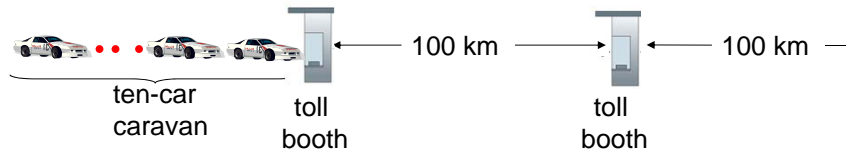
## Caravan analogy



- ❖ cars “propagate” at 100 km/hr
- ❖ toll booth takes 12 sec to service car (bit transmission time)
- ❖ car ~ bit; caravan ~ packet
- ❖ **Q: How long until caravan is lined up before 2nd toll booth?**
  - time to “push” entire caravan through toll booth onto highway =  $12 \times 10 = 120$  sec
  - time for last car to propagate from 1st to 2nd toll booth:  $100 \text{ km} / (100 \text{ km/hr}) = 1$  hr
  - **A: 62 minutes**

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## Caravan analogy (more)



- ❖ suppose cars now “propagate” at 1000 km/hr
- ❖ and suppose toll booth now takes one min to service a car
- ❖ **Q: Will cars arrive to 2nd booth before all cars serviced at first booth?**
  - **A: Yes!** after 7 min, 1st car arrives at second booth; three cars still at 1st booth.

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## Exercise

Consider two hosts, A and B, are separated by 5,000 kilometers and are directly connected by a link with transmission rate  $R=10\text{Mbps}$ . The propagation speed over the link is  $2.5 \times 10^8$  meters/sec.

- What is the propagation delay of the link?

- Solution:

propagation delay

= link length / propagation speed

=  $5,000,000 \text{ m} / 2.5 \times 10^8 \text{ meters/sec}$

= 0.02 sec

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## Exercise

Consider two hosts, A and B, are separated by 5,000 kilometers and are directly connected by a link with transmission rate  $R=10\text{Mbps}$ . The propagation speed over the link is  $2.5 \times 10^8$  meters/sec.

- Consider only one file of 50,000 bits from Host A to Host B. What is the transmission delay?

- Solution:

transmission delay

= file size / transmission rate

=  $50,000 \text{ bits} / 10 \times 10^6 \text{ bits/s}$

= 0.005 sec

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## Exercise

Consider two hosts, A and B, are separated by 5,000 kilometers and are directly connected by a link with transmission rate  $R=10\text{Mbps}$ . The propagation speed over the link is  $2.5 \times 10^8$  meters/sec.

- How to interpret the bandwidth –propagation delay product?

- Solution:

bandwidth-propagation delay product

$$= 10 \times 10^6 \text{ bits/s} \times 0.02 \text{ sec}$$

$$= 2 \times 10^5 \text{ bits}$$

The amount of data that could be in transit in the link.

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## Exercise

Consider two hosts, A and B, are separated by 5,000 kilometers and are directly connected by a link with transmission rate  $R=10\text{Mbps}$ . The propagation speed over the link is  $2.5 \times 10^8$  meters/sec.

- What is the width (in meters) of a bit in the link?

- Solution:

width = propagation speed \* transmission time for one bit

$$= \text{propagation speed} / \text{transmission rate}$$

$$= 2.5 \times 10^8 \text{ meters/sec} / 10 \times 10^6 \text{ bits/sec}$$

$$= 25 \text{ meters/bit}$$

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## Exercise

Consider two hosts, A and B, are separated by 5,000 kilometers and are directly connected by a link with transmission rate  $R=10\text{Mbps}$ . The propagation speed over the link is  $2.5 \times 10^8$  meters/sec.

- Suppose a 5000bits file is sent continuously as one big message. What is the maximum number of bits that will be in the link at any given time?

- Solution:

link distance / bit width

$$= 5,000,000 \text{ m} / 25 \text{ meters/bit}$$

$$= 2 \times 10^5 \text{ bits} > 5000 \text{ bits, so the solution is 5000 bits.}$$

propagation delay \* transmission rate

$$= 0.02 \text{ sec} * 10 \times 10^6 \text{ bits/sec}$$

$$= 2 \times 10^5 \text{ bits} > 5000 \text{ bits, so the solution is 5000 bits.}$$

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## Exercise

Consider two hosts, A and B, are separated by 5,000 kilometers and are directly connected by a link with transmission rate  $R=10\text{Mbps}$ . The propagation speed over the link is  $2.5 \times 10^8$  meters/sec.

- Let  $x$  denote the size of the file, what is the minimum value of  $x$  for the link to be continuously transmitting?

- solution

link distance / bit width

$$= 5,000,000 \text{ m} / 25 \text{ meters/bit}$$

$$= 2 \times 10^5 \text{ bits}$$

propagation delay \* transmission rate

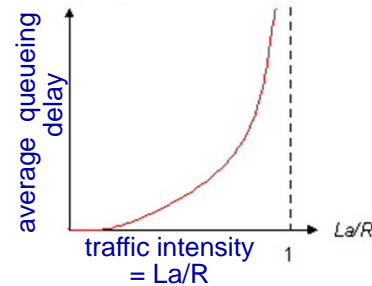
$$= 0.02 \text{ sec} * 10 \times 10^6 \text{ bits/sec}$$

$$= 2 \times 10^5 \text{ bits}$$

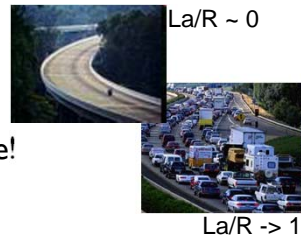
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## Queuing delay

- ❖  $R$ : link bandwidth (bps)
- ❖  $L$ : packet length (bits)
- ❖  $a$ : average packet arrival rate



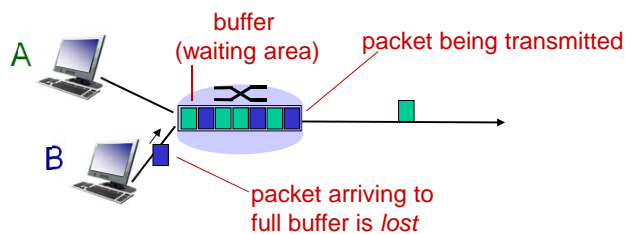
- ❖  $La/R \sim 0$ : avg. queueing delay small
- ❖  $La/R \rightarrow 1$ : avg. queueing delay large
- ❖  $La/R > 1$ : more “work” arriving than can be serviced, average delay infinite!



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## Packet loss

- ❖ queue (aka buffer) preceding link in buffer has finite capacity
- ❖ packet arriving to full queue dropped (aka lost)
- ❖ lost packet may be retransmitted by previous node, by source end system, or not at all



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## End-to-End Delay

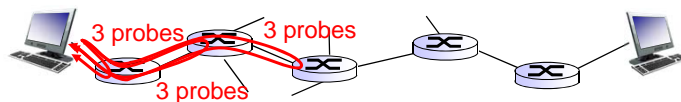
- ❖ Suppose there are  $N-1$  routers between the source and the destination

$$D_{\text{end-end}} = N (d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}})$$

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## “Real” Internet delays and routes

- ❖ what do “real” Internet delay & loss look like?
- ❖ **traceroute** program: provides delay measurement from source to router along end-end Internet path towards destination. For all  $i$ :
  - sends three packets that will reach router  $i$  on path towards destination
  - router  $i$  will return packets to sender
  - sender times interval between transmission and reply.



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## “Real” Internet delays, routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

3 delay measurements from  
gaia.cs.umass.edu to cs-gw.cs.umass.edu

```

1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 ***
18 ***
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

```

trans-oceanic link

\* means no response (probe lost, router not replying)

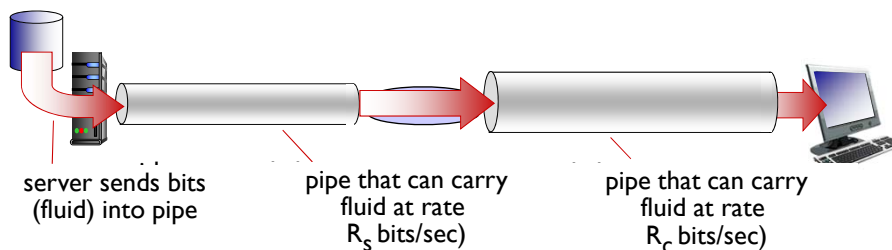
\* Do some traceroutes from exotic countries at [www.traceroute.org](http://www.traceroute.org)

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## Throughput

❖ **throughput**: rate (bits/time unit) at which bits transferred between sender/receiver

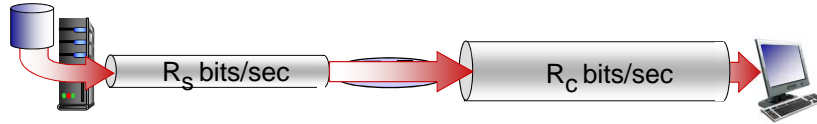
- **instantaneous**: rate at given point in time
- **average**: rate over longer period of time



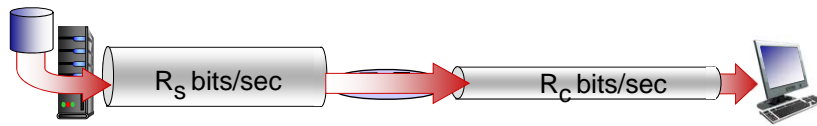
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## Throughput (more)

- ❖  $R_s < R_c$  What is average end-end throughput?



- ❖  $R_s > R_c$  What is average end-end throughput?



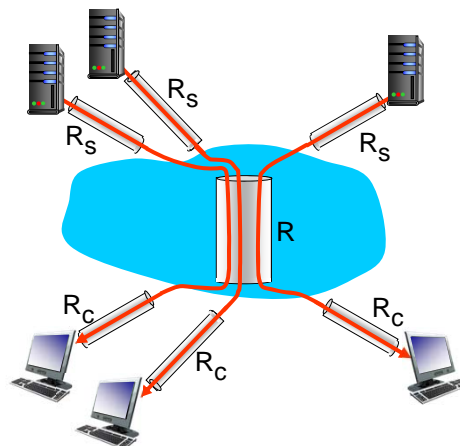
*bottleneck link*

link on end-end path that constrains end-end throughput

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## Throughput: Internet scenario

- ❖ per-connection end-end throughput:  $\min(R_c, R_s, R/10)$
- ❖ in practice:  $R_c$  or  $R_s$  is often bottleneck



10 connections (fairly) share backbone bottleneck link  $R$  bits/sec

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# Addressing

- ❖ Physical address (aka link address)
  - E.g.: 07:01:02:01:2C:4B
- ❖ Logical address: IP
  - No two publicly addressed and visible hosts on the Internet can have the same IP
  - ARP/RARP: physical address  $\leftrightarrow$  IP address
- ❖ Port address
  - E.g.: 80 web port
- ❖ Others: URL, easy to remember
  - DNS: URL  $\rightarrow$  IP