



CS 152/252A Computer Architecture and Engineering



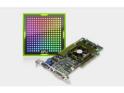
Lecture 18: GPUs



A Timeline of Innovation

GPU





Invents the GPU, the graphics processing unit,

which sets the stage to reshape the computing













Plays a foundational role in the building of the metaverse, the next stage of the internet, with the NVIDIA Omniverse™ platform.



https://www.nvidia.com/en-us/about-nvidia/corporate-timeline/

Last Time in Lecture

Vector supercomputers

- Vector register versus vector memory
- Scaling performance with lanes
- Stripmining
- Chaining
- Masking
- Scatter/Gather

Types of Parallelism

- Instruction-Level Parallelism (ILP)
 - Execute independent instructions from one instruction stream in parallel (pipelining, superscalar, VLIW)
- Thread-Level Parallelism (TLP)
 - Execute independent instruction streams in parallel (multithreading, multiple cores)
- Data-Level Parallelism (DLP)
 - Execute multiple operations of the same type in parallel (vector/SIMD execution)
- Which is easiest to program?
- Which is most flexible form of parallelism?
 - i.e., can be used in more situations
- Which is most efficient?
 - i.e., greatest tasks/second/area, lowest energy/task

Simplified Processor Evolution

Memory

Memory

Memory

Cache

Fetch/Decode

ALU

Registers

Cache

Fetch/ Decode

ALU

Registers

Single-core, singlethread, **super-scalar** processor

Single-core, singlethread, **SIMD/vector** processor

Single-core, single-thread, scalar processor

Simplified Processor Evolution

Memory

Memory

Cache

Fetch/Decode

ALU

Registers

Fetch/Decode

ALU

Registers

Registers

Single-core, single-thread, scalar processor Single-core, multi-thread, scalar processor

Simplified Processor Evolution

Memory

Fetch/Decode

Instruction Selection

ALU

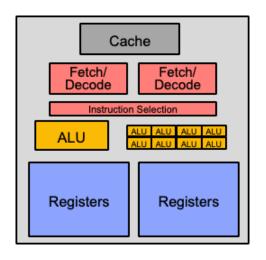
ALU ALU ALU ALU ALU
ALU ALU ALU
Registers

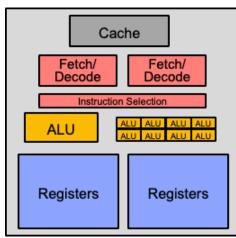
Registers

Registers

Memory

Shared Cache





Single-core, multi-thread, super-scalar, OoO, SIMD processor multi-core, multi-thread, super-scalar, OoO, SIMD processor

Resurgence of DLP

 Convergence of application demands and technology constraints drives architecture choice

 New applications, such as graphics, machine vision, speech recognition, machine learning, etc. all require large numerical computations that are often trivially data parallel

 SIMD-based architectures (vector-SIMD, subword-SIMD, SIMT/GPUs) are most efficient way to execute these algorithms

Graphics Processing Units (GPUs)

- Original GPUs were dedicated fixed-function devices for generating 3D graphics (mid-late 1990s) including highperformance floating-point units
 - Provide workstation-like graphics for PCs
 - User could configure graphics pipeline, but not really program it
- Over time, more programmability added (2001-2005)
 - E.g., New language Cg for writing small programs run on each vertex or each pixel, also Windows DirectX variants
 - Massively parallel (millions of vertices or pixels per frame) but very constrained programming model
- Some users noticed they could do general-purpose computation by mapping input and output data to images, and computation to vertex and pixel shading computations
 - Incredibly difficult programming model as had to use graphics pipeline model for general computation

General-Purpose GPUs (GP-GPUs)

- In 2006, Nvidia introduced GeForce 8800 GPU, which supported a new programming language CUDA (in 2007)
 - "Compute Unified Device Architecture"
 - Subsequently, broader industry pushing for OpenCL, a vendor-neutral version of same ideas.
- Idea: Take advantage of GPU computational performance and memory bandwidth to accelerate some kernels for general-purpose computing
- Attached processor model: Host CPU issues data-parallel kernels to GP-GPU for execution
- This lecture has a simplified version of Nvidia CUDA-style model and only considers GPU execution for computational kernels, not graphics
 - Would need whole other course to describe graphics processing

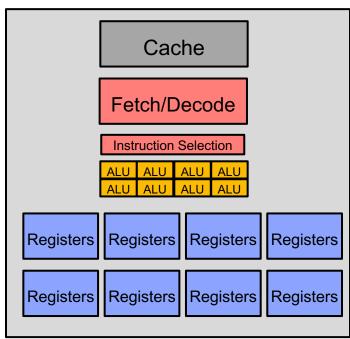
GPU SIMT

- "Single instruction multiple thread"
 - Or multi-threaded SIMD/streaming multi-processor (SM)

Memory Cache Fetch/ Fetch/ Decode Decode **Instruction Selection** ALU **ALU** Registers Registers

> Single-core, multithread, super-scalar, OoO, SIMD processor

Memory



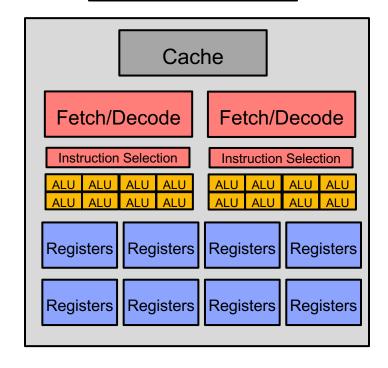
Single-core, multithread, InO, SIMD processor

GPU SIMT

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Memory

Single-core, multithread, super-scalar, OoO, SIMD processor Memory

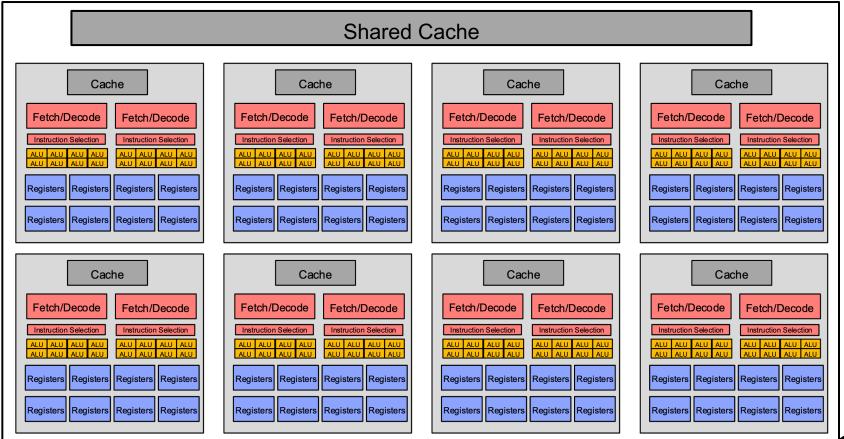


Single-core, multithread, InO, SIMD processor

GPU Architecture

 Many multi-threaded SIMD cores (streaming multiprocessors)

GPU Memory

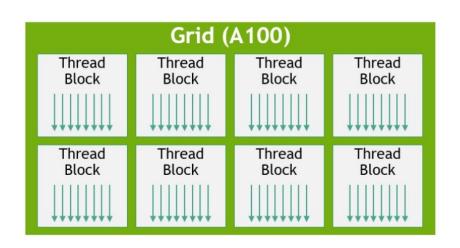


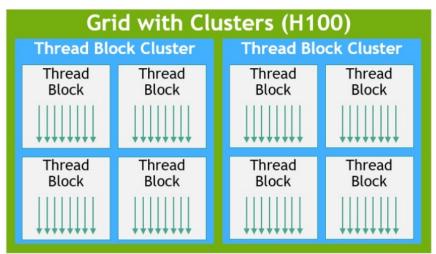
Simplified CUDA Programming Model

 Computation performed by a very large number of independent small scalar threads (CUDA threads or microthreads) grouped into thread blocks.

```
// C version of DAXPY loop.
void daxpy(int n, double a, double*x, double*y)
{ for (int i=0; i<n; i++)
      y[i] = a*x[i] + y[i];
// CUDA version.
 host // Piece run on host processor.
int nblocks = (n+255)/256; //256 CUDA threads/block
daxpy << nblocks, 256>>> (n, 2.0, x, y);
 device // Piece run on GP-GPU.
void daxpy(int n, double a, double*x, double*y)
{ int i = blockIdx.x*blockDim.x + threadId.x;
 if (i<n) y[i]=a*x[i]+y[i]; }</pre>
```

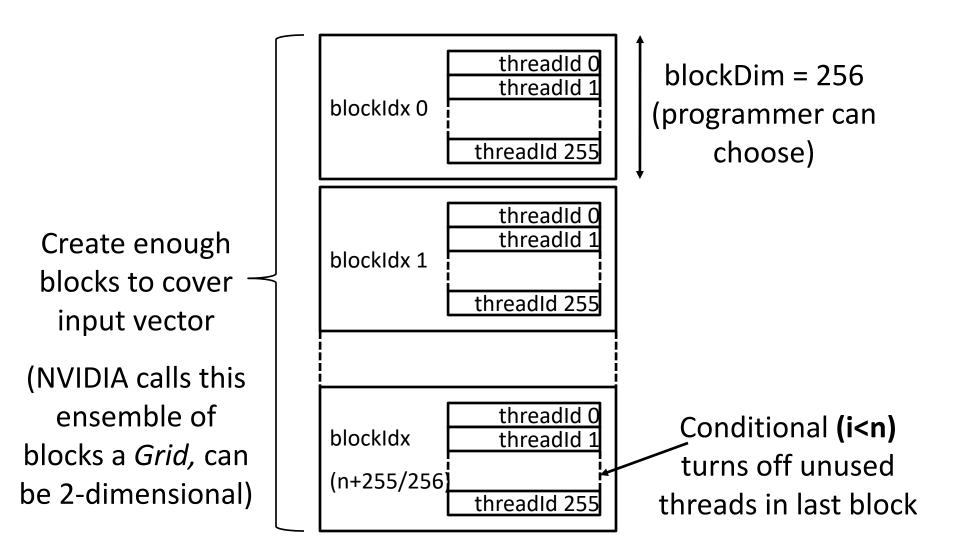
Programmer's View of Execution



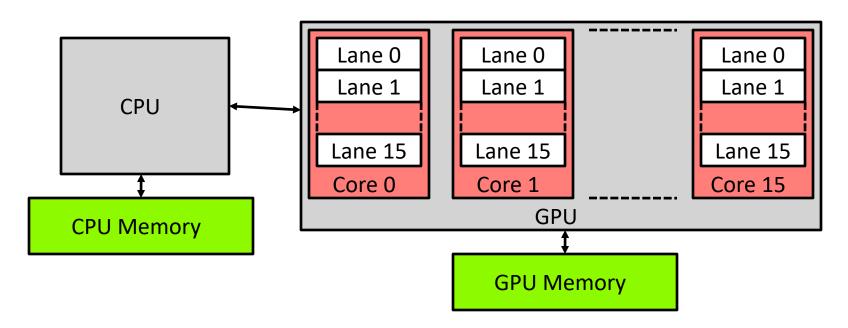


https://developer.nvidia.com/blog/nvidia-hopperarchitecture-in-depth/

Programmer's View of Execution



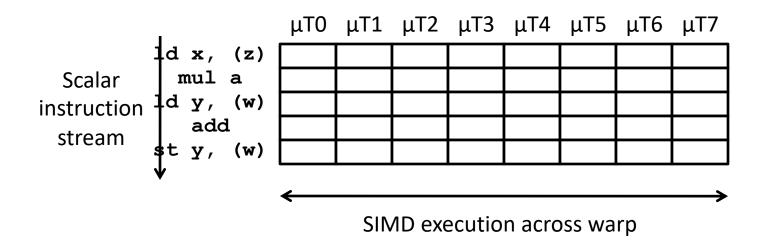
Hardware Execution Model



- GPU is built from multiple parallel cores, each core contains a multithreaded SIMD processor with multiple lanes but with no scalar processor
- CPU sends whole "grid" over to GPU, which distributes thread blocks among cores (each thread block executes on one core)
 - Programmer unaware of number of cores

"Single Instruction, Multiple Thread" (SIMT)

 GPUs use a SIMT model, where individual scalar instruction streams for each CUDA thread are grouped together for SIMD execution on hardware (NVIDIA groups 32 CUDA threads into a warp)



Implications of SIMT Model

- All "vector" loads and stores are scatter-gather, as individual μthreads perform scalar loads and stores
 - GPU adds hardware to dynamically coalesce individual μthread loads and stores to mimic vector loads and stores
- Every μthread has to perform stripmining calculations redundantly ("am I active?") as there is no scalar processor equivalent

Multithreaded SIMD Processor

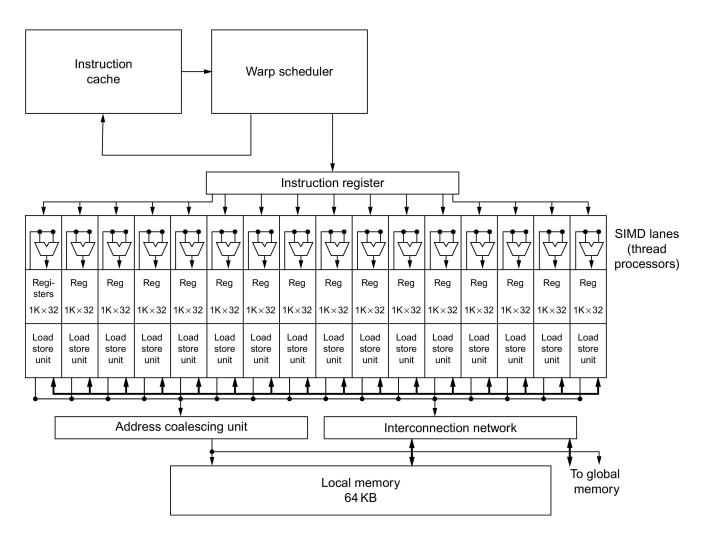
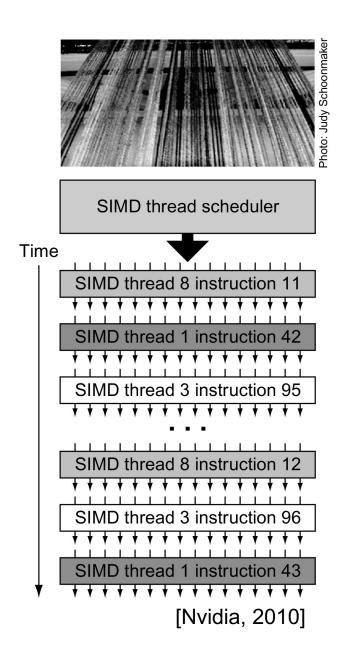


Figure 4.14 Simplified block diagram of a multithreaded SIMD Processor. It has 16 SIMD Lanes. The SIMD Thread Scheduler has, say, 64 independent threads of SIMD instructions that it schedules with a table of 64 program counters (PCs). Note that each lane has 1024 32-bit registers.

Warps are multithreaded on core



- Warp == SIMD Thread, i.e., a thread of SIMD instructions
- One warp of 32 µthreads is a single thread in the hardware
- Multiple warp threads are interleaved in execution on a single core to hide latencies (memory and functional unit)
- A single thread block can contain multiple warps (up to 512 μT max in CUDA), all mapped to single core
- Can have multiple blocks executing on one core

Vector vs Multithreaded SIMD

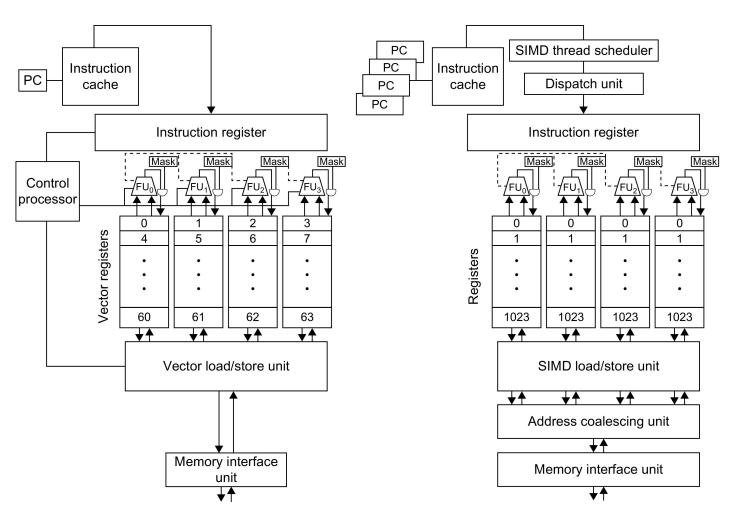


Figure 4.22 A vector processor with four lanes on the left and a multithreaded SIMD Processor of a GPU with four SIMD Lanes on the right. (GPUs typically have 16 or 32 SIMD Lanes.) The Control Processor supplies scalar operands for scalar-vector operations, increments addressing for unit and nonunit stride accesses to memory, and performs other accounting-type operations. Peak memory performance occurs only in a GPU when the Address Coalescing Unit can discover localized addressing. Similarly, peak computational performance occurs when all internal mask bits are set identically. Note that the SIMD Processor has one PC per SIMD Thread to help with multithreading.

CS152 Administrivia

- HW 4 out this week
- Lab 3 due 3/23
- Midterm feedback

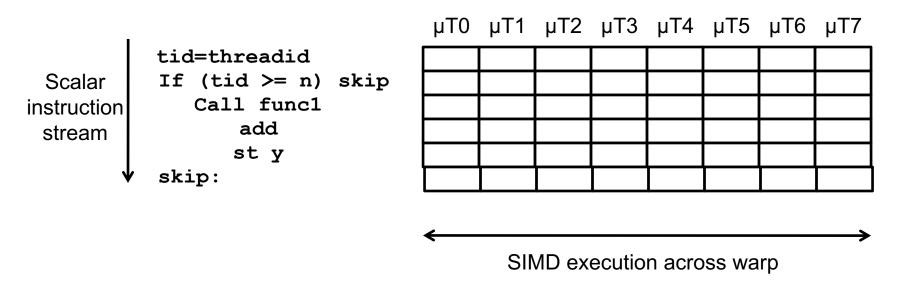
CS252 Administrivia

Readings on OoO this Wednesday

CS252 23

Conditionals in SIMT model

- Simple if-then-else are compiled into predicated execution, equivalent to vector masking
- More complex control flow compiled into branches
- How to execute a vector of branches? Vector function calls?



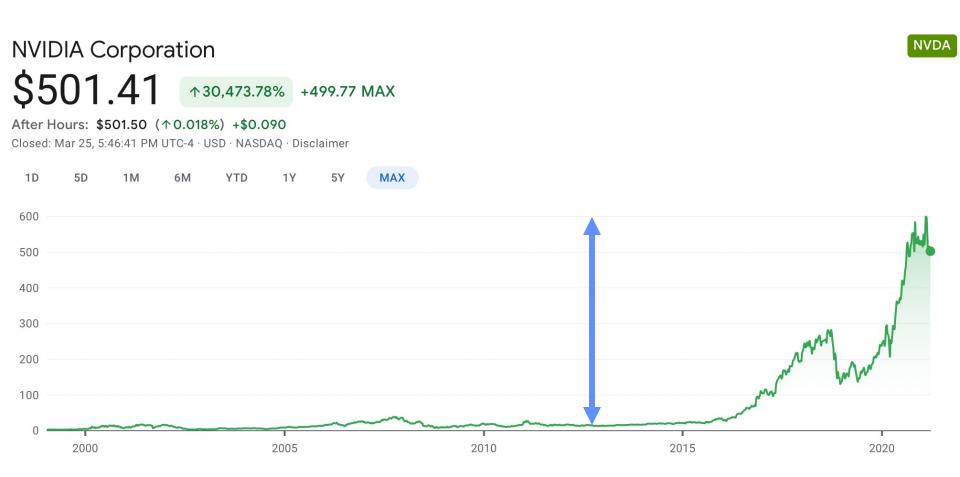
Branch Divergence

- Hardware tracks which µthreads take or don't take branch
- If all go the same way, then keep going in SIMD fashion
- If not, create mask vector indicating taken/not-taken
- Keep executing not-taken path under mask, push taken branch PC+mask onto a hardware stack and execute later
- When can execution of µthreads in warp reconverge?

SIMT

- Illusion of many independent threads
- But for efficiency, programmer must try and keep µthreads aligned in a SIMD fashion
 - Try and do unit-stride loads and store so memory coalescing kicks in
 - Avoid branch divergence so most instruction slots execute useful work and are not masked off

Important of Machine Learning for GPUs



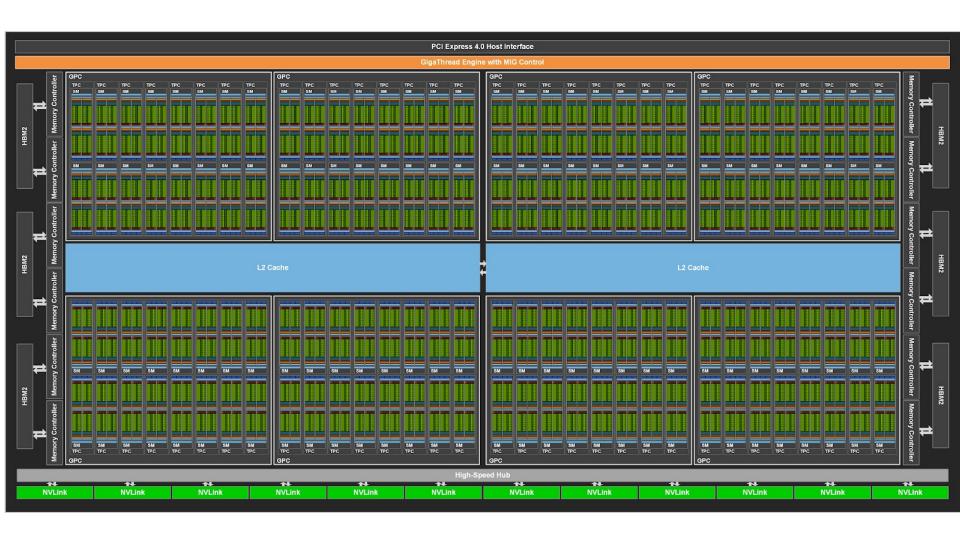
NVIDIA stock price 40x in 9 years (since deep learning became important)

Nvidia Ampere A100 GPU

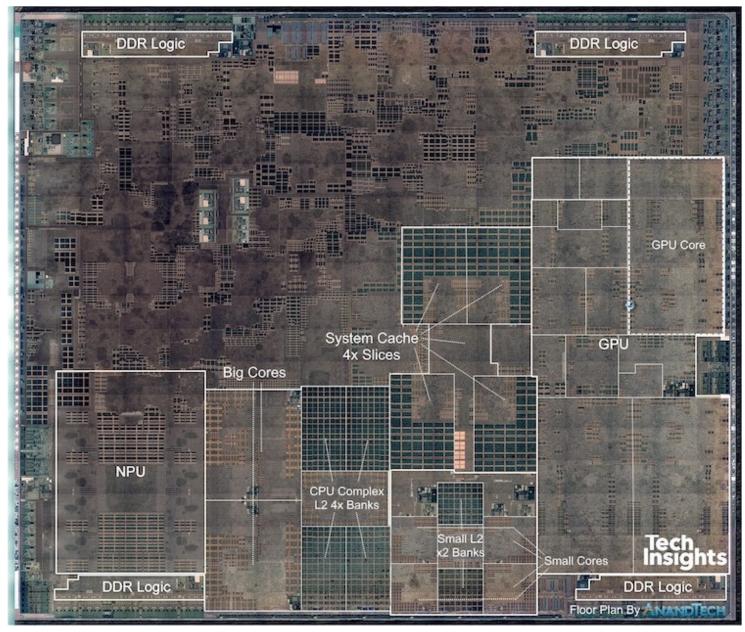


[Nvidia, 2022]

Nvidia Ampere A100 GPU



Apple A12 Processor (2018)



- 83.27mm²
- 7nm technology

[Source: Tech Insights, AnandTech]

Acknowledgements

- This course is partly inspired by previous MIT 6.823 and Berkeley CS252 computer architecture courses created by my collaborators and colleagues:
 - Arvind (MIT)
 - Krste Asanovic (MIT/UCB)
 - Joel Emer (Intel/MIT)
 - James Hoe (CMU)
 - John Kubiatowicz (UCB)
 - David Patterson (UCB)