

去花

第三处

第一处

第二处

解密

去花

新手时复现花指令很少见师傅放源码，其实对比观察源码就很清楚了，本质就是去除多余opcode，亲自动手写一遍再逆向感觉是完全不一样的。

源码如下

第三处

类似smc吧，一开始想整活但嫌麻烦，先加密再去加的花，所以不做太大改动，原理就是静态时是E9 jmp指令，运行后xor成5个nop，不影响程序流执行

```
.text:00000000000024A5 1B8 48 8B 85 78 FE FF      mov     rax, [rbp+var_188]
.text:00000000000024A5 1B8 FF
.text:00000000000024AC 1B8 48 89 C7                mov     rdi, rax
.text:00000000000024AF 1B8 E8 D7 FE FF FF        call    sub_238B
.text:00000000000024B4
.text:00000000000024B4      loc_24B4:                ; DATA XREF: sub_246B+2C↑o
.text:00000000000024B4 1B8 E9 11 45 14 19        jmp     near ptr 191469CAh
.text:00000000000024B4      sub_246B                endp
.text:00000000000024B9      ; -----
.text:00000000000024B9 48 B8 DC C3 F5 CB          mov     rax, 0D9CBC6D9CBF5C3DCh
.text:00000000000024B9 D9 C6 CB D9
.text:00000000000024C3 48 89 85 CF FE FF          mov     [rbp-131h], rax
.text:00000000000024C3 FF
.text:00000000000024CA 48 B8 9F 9B 9F CF          mov     rax, 999D939CCF9F9B9Fh
```

```
1 static inline void* page_align(void *addr) {
2     size_t page_size = getpagesize();
3     return (void*)((uintptr_t)addr & ~(page_size - 1));
4 }
5
6 void flower_with_addr(void *junk_addr) {
7     void *page = page_align(junk_addr);
8     size_t page_size = getpagesize();
9
10    if (mprotect(page, page_size, PROT_READ | PROT_WRITE | PROT_EXEC) !=
11        0) {
12        perror("mprotect");
13        return;
14    }
15
16    uint8_t key[] = {0x79, 0x81, 0xd5, 0x84, 0x89};
17    for (int i = 0; i < 5; i++) {
18        ((uint8_t*)junk_addr)[i] ^= key[i]; // -> 0x90 each
19    }
20
21    mprotect(page, page_size, PROT_READ | PROT_EXEC);
22
23    void *junk_addr = &&junk_label;
24    flower_with_addr(junk_addr);
25    junk_label:
26    __asm__ volatile (
27        ".byte 0xe9 \n\t"
28        ".byte 0x11 \n\t"
29        ".byte 0x45 \n\t"
30        ".byte 0x14 \n\t"
31        ".byte 0x19 \n\t"
32        ::: "memory"
33    );
```

第一处

标准jz/jnz

```

.text:0000000000002CB5 48 89 45 E8      mov     [rbp-18h], rax
.text:0000000000002CB7 31 C0           xor     eax, eax
.text:0000000000002CB9 53             push    rbx
.text:0000000000002CBA 48 31 DB       xor     rbx, rbx
.text:0000000000002CBD 48 85 DB       test    rbx, rbx
.text:0000000000002CC0 75 02          jnz     short loc_2CC4
.text:0000000000002CC2 74 01          jz      short near ptr loc_2CC4+1
.text:0000000000002CC4
.text:0000000000002CC4
.text:0000000000002CC4 E9 5B 48 B8 C9   jmp     near ptr 0FFFFFFFFC9B87524h
.text:0000000000002CC4
.text:0000000000002CC9 CF CE C9 CE CB 92... db 0CFh, 0CEh, 0C9h, 0CEh, 0CBh, 92h, 0CBh
.text:0000000000002CD0
.text:0000000000002CD0 48 BA C9 99 9F C9 mov     rdx, 93C9CB9FC99F99C9h

```

```

1  __asm__ volatile (
2      "pushq %%rbx \n\t"
3      "xorq  %%rbx, %%rbx \n\t" //ida 识别$1 $1+1
4      "testq %%rbx, %%rbx \n\t"
5      "jnz  1f \n\t"
6      "jz   2f \n\t"
7
8      "1: \n\t"
9      ".byte 0xe9 \n\t"
10
11     "2: \n\t"
12     "popq  %%rbx \n\t"
13     :
14     :
15     : "rbx", "memory"
16 );

```

第二处

call+retn, rsp具体加的值根据opcode计算其实很明显

```

.text:0000000000002DA0 BE 60 00 00 00      mov     esi, 60h ; ...
.text:0000000000002DA5 48 89 C7           mov     rdi, rax
.text:0000000000002DA8 E8 4C F5 FF FF      call    sub_22F9
.text:0000000000002DAD 48 8D 85 C0 FE FF    lea     rax, [rbp-140h]
.text:0000000000002DAD FF
.text:0000000000002DB4 48 89 85 A8 FE FF    mov     [rbp-158h], rax
.text:0000000000002DB4 FF
.text:0000000000002DBB B8 00 00 00 00      mov     eax, 0
.text:0000000000002DC0 E8 2C FE FF FF      call    sub_2BF1
.text:0000000000002DC5 E8 01 00 00 00      call    loc_2DCB
.text:0000000000002DC5
; -----
.text:0000000000002DCA 83             db 83h
.text:0000000000002DCB
; -----
.text:0000000000002DCB
; -----
loc_2DCB:
; CODE XREF: .text:0000000000002DC5↑j
add     qword ptr [rsp], 8
retn
; -----
.text:0000000000002DD0 48 83 04 24 08      db 0F3h, 83h, 0BDh, 9Ch, 0FEh, 2 dup(0FFh)
.text:0000000000002DD0 C3
; -----
.text:0000000000002DD1 F3 83 BD 9C FE FF... db 48000000FD8F0F01h, 8B48FFFFFFF90858Bh, 0FFFFFF88589480840h
.text:0000000000002DD8 01 0F 8F FD 00 00

```

```

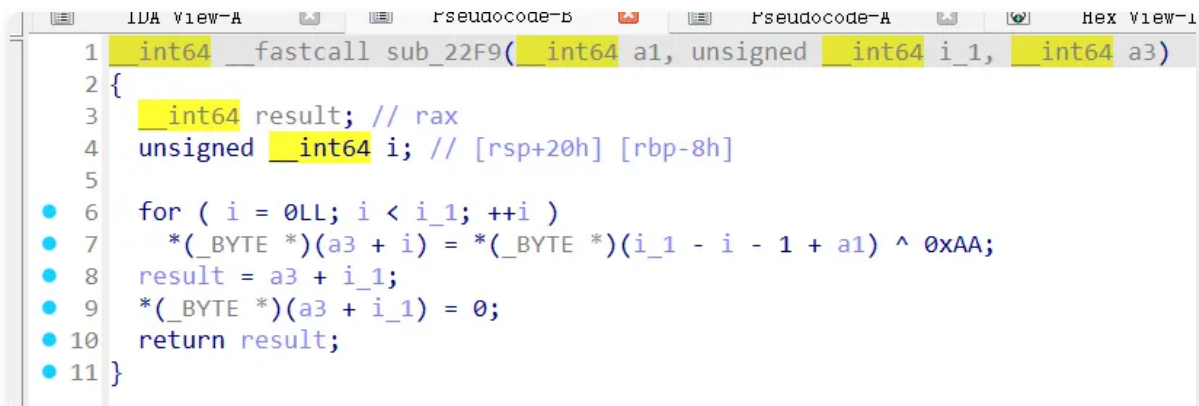
1  __asm__ volatile (
2      "call 1f \n\t"
3      ".byte 0x83 \n\t"
4
5      "1: \n\t"
6      "addq $8, (%rsp) \n\t"
7      "ret \n\t"
8      ".byte 0xf3 \n\t"
9      :
10     :
11     : "memory"
12 );

```

解密

salsa20+aes_ecb无魔改

中间加了简单的倒序xor处理去除一些内容特征



```

1  int64 __fastcall sub_22F9( int64 a1, unsigned int64 i_1, int64 a3)
2  {
3      int64 result; // rax
4      unsigned int64 i; // [rsp+20h] [rbp-8h]
5
6      for ( i = 0LL; i < i_1; ++i )
7          *(_BYTE *)(a3 + i) = *(_BYTE *)(i_1 - i - 1 + a1) ^ 0xAA;
8      result = a3 + i_1;
9      *(_BYTE *)(a3 + i_1) = 0;
10     return result;
11 }

```

eg.cyber_chef

Recipe

^

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🗑️

AES Decrypt

^

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⏸️

Key

venom2025v...

UTF8 ▾

IV

HEX ▾

Mode

ECB

Input

Hex

Output

Raw

Salsa20

^

🛑

⏸️

Key

0ff71ec39a5...

HEX ▾

Nonce

salsa_iv

UTF8 ▾

Counter

0

Rounds

20

Input

Raw

Output

Raw

Input

59cc9e17b97199cad788f5f9d919531ee9f09bb233b5565233aef
c

abc 96

≡

1

Output

vctf{ab92105b_047e_aa38_0b25_a6989b8e5ed5}

eg.code

```

1  import binascii
2  import struct
3  from Crypto.Cipher import AES,Salsa20
4
5  #aes_dec
6  # enc = [
7  #     0xc9, 0xcf, 0xce, 0xc9, 0xce, 0xcb, 0x92, 0xcb, 0xc9, 0x99, 0x9f, 0x
8  #     c9, 0x9f, 0xcb, 0xc9, 0x93,
9  #     0xcc, 0xc9, 0x9d, 0x9e, 0xce, 0x93, 0x9b, 0x92, 0xc9, 0x99, 0x9a, 0x
10 #     9d, 0x9d, 0x99, 0x98, 0x9f,
11 #     0x99, 0x9c, 0x9c, 0xce, 0x9b, 0xcb, 0x93, 0x99, 0xcf, 0xce, 0x9b, 0x
12 #     cc, 0xcf, 0xcb, 0x99, 0x99,
13 #     0x98, 0x9f, 0x9c, 0x9f, 0x9f, 0xc8, 0x99, 0x99, 0x98, 0xc8, 0xc8, 0x
14 #     93, 0x9a, 0xcc, 0x93, 0xcf,
15 #     0xcf, 0x9b, 0x99, 0x9f, 0x93, 0x9b, 0x93, 0xce, 0x93, 0xcc, 0x9f, 0x
16 #     cc, 0x92, 0x92, 0x9d, 0xce,
17 #     0xcb, 0xc9, 0x93, 0x93, 0x9b, 0x9d, 0x93, 0xc8, 0x9d, 0x9b, 0xcf, 0x
18 #     93, 0xc9, 0xc9, 0x93, 0x9f
19 # ]
20 # hoge=bytes(b ^ 0xAA for b in enc[::-1])
21 # print(hoge)
22 # b'59cc9e17b97199cad788f5f9d919531ee9f09bb233b5565233aef1de39a1d663523770
23 # 3c819d47cf9ca5c53ca8adcdec '
24 hoge=bytes.fromhex('59cc9e17b97199cad788f5f9d919531ee9f09bb233b5565233aef1
25 de39a1d6635237703c819d47cf9ca5c53ca8adcdec')
26 aes_key=b"venom2025venom25"
27 cipher_aes=AES.new(aes_key, AES.MODE_ECB)
28 pad=cipher_aes.decrypt(hoge)
29 pad_len=pad[-1]
30 #去除PKS7填充
31 enc_tmp=pad[:-pad_len]
32
33 #salsa20_dec
34 # nonce=0xD9CBC6D9CBF5C3DC
35 # salsa=[0x999D939CCF9F9B9F,0x9A93C89CC898CF93,0x93CB9FCE9998989D,0x9ACCCC
36 # 9D9BCFC999,
37 #     0x999D939CCF9F9B9F,0x9A93C89CC898CF93,0x93CB9FCE9998989D,0x9ACCCC
38 # 9D9BCFC999]
39 # hoge1=[]
40 # hoge2=[]
41 # hoge1.extend(nonce.to_bytes(8, 'little'))
42 # for qword in salsa:
43 #     hoge2.extend(qword.to_bytes(8, 'little'))
44 # nonce=bytes(b ^ 0xAA for b in hoge1[::-1])

```

```

35 # salsa_key=bytes(b ^ 0xAA for b in hoge2[::-1])
36 # print(nonce,salsa_key)
37 # b'salsa_iv'
38 # b'0ff71ec39a5d322709b6b2e93796e5150ff71ec39a5d322709b6b2e93796e515'
39 nonce=b'salsa_iv'
40 salsa_key=bytes.fromhex('0ff71ec39a5d322709b6b2e93796e5150ff71ec39a5d32270
9b6b2e93796e515')
41 cipher_salsa=Salsa20.new(key=salsa_key,nonce=nonce)
42 flag=cipher_salsa.decrypt(enc_tmp)
43 print(flag.decode())
44 # vctf{ah92105h_047e_aa38_0b25_a6Q8Qh8e5ed5}

```