

## 去花

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## 去花

新手时复现花指令很少见师傅放源码，其实对比观察源码就很清楚了，本质就是去除多余opcode，亲自动手写一遍再逆向感觉是完全不一样的。

源码如下

## 第三处

类似smc吧，一开始想整活但嫌麻烦，先加密再去加的花，所以不做太大改动，原理就是静态时是E9 jmp指令，运行后xor成5个nop，不影响程序流执行

```
... .text:00000000000024A5 1B8 48 8B 85 78 FE FF          mov    rax, [rbp+var_188]
.text:00000000000024A5 1B8 FF
.text:00000000000024AC 1B8 48 89 C7          mov    rdi, rax
.text:00000000000024AF 1B8 E8 D7 FE FF FF        call   sub_238B
.text:00000000000024B4
.text:00000000000024B4
.text:00000000000024B4 1B8 E9 11 45 14 19      loc_24B4:    jmp   near ptr 191469CAh ; DATA XREF: sub_246B+2C↑
.text:00000000000024B4
.text:00000000000024B4
.text:00000000000024B4
.text:00000000000024B4
.text:00000000000024B9 48 B8 DC C3 F5 CB          sub_246B:    endp
.text:00000000000024B9 D9 C6 CB D9
.text:00000000000024C3 48 89 85 CF FE FF        mov    rax, 0D9CBC6D9CBF5C3DCh
.text:00000000000024C3 FF
.text:00000000000024CA 48 B8 9F 9B 9F CF        mov    [rbp-131h], rax
.text:00000000000024CA
... 000000000000000000 00 00 00 00
```

```

1 static inline void* page_align(void *addr) {
2     size_t page_size = getpagesize();
3     return (void*)((uintptr_t)addr & ~(page_size - 1));
4 }
5
6 void flower_with_addr(void *junk_addr) {
7     void *page = page_align(junk_addr);
8     size_t page_size = getpagesize();
9
10    if (mprotect(page, page_size, PROT_READ | PROT_WRITE | PROT_EXEC) != 0) {
11        perror("mprotect");
12        return;
13    }
14
15    uint8_t key[] = {0x79, 0x81, 0xd5, 0x84, 0x89};
16    for (int i = 0; i < 5; i++) {
17        ((uint8_t*)junk_addr)[i] ^= key[i]; // -> 0x90 each
18    }
19
20    mprotect(page, page_size, PROT_READ | PROT_EXEC);
21 }
22
23 void *junk_addr = &&junk_label;
24 flower_with_addr(junk_addr);
25 junk_label:
26     __asm__ volatile (
27         ".byte 0xe9 \n\t"
28         ".byte 0x11 \n\t"
29         ".byte 0x45 \n\t"
30         ".byte 0x14 \n\t"
31         ".byte 0x19 \n\t"
32         ::: "memory"
33     );

```

## 第一处

标准jz/jnz

```
1 __asm__ volatile (
2         "pushq %%rbx \n\t"
3         "xorq %%rbx, %%rbx \n\t"    //ida 识别$1 $1+1
4         "testq %%rbx, %%rbx \n\t"
5         "jnz    1f \n\t"
6         "jz     2f \n\t"
7
8         "1: \n\t"
9         ".byte 0xe9 \n\t"
10
11        "2: \n\t"
12        "popq %%rbx \n\t"
13        :
14        :
15        : "rbx", "memory"
16 );
```

第二处

call+ret, rsp具体加的值根据opcode计算其实很明显

```
    .text:000000000002DA0 BE 60 00 00 00          mov    esi, 60h ; 
    .text:000000000002DA5 48 89 C7              mov    rdi, rax
    .text:000000000002DA8 E8 4C F5 FF FF        call   sub_22F9
    .text:000000000002DAD 48 8D 85 C0 FE FF      lea    rax, [rbp-140h]
    .text:000000000002DAD FF
    .text:000000000002DB4 48 89 85 A8 FE FF      mov    [rbp-158h], rax
    .text:000000000002DB4 FF
    .text:000000000002DBB B8 00 00 00 00          mov    eax, 0
    .text:000000000002DC0 E8 2C FE FF FF        call   sub_2BF1
    .text:000000000002DC5 E8 01 00 00 00          call   loc_2DCB
    .text:000000000002DC5
    .text:000000000002DCA 83                   db 83h
    .text:000000000002DCB
    .text:000000000002DCB
    .text:000000000002DCB
    .text:000000000002DCB 48 83 04 24 08       loc_2DCB:           add    qword ptr [rsp], 8 ; CODE XREF: .text:000000000002DC5↑j
    .text:000000000002DD0 C3                   retn
    .text:000000000002DD0
    .text:000000000002DD1 F3 83 BD 9C FE FF...  db 0F3h, 83h, 0BDh, 9Ch, 0FEh, 2 dup(0FFh)
    .text:000000000002DD8 01 0E 8F FD 00 00      db 48000000FD8F0E01h, 8848FFFFE90858Rh, 0FFFF88589480840h
```

```
1  __asm__ volatile (
2      "call 1f \n\t"
3      ".byte 0x83 \n\t"
4
5      "1: \n\t"
6      "addq $8, (%rsp) \n\t"
7      "ret \n\t"
8      ".byte 0xf3 \n\t"
9      :
10     :
11     : "memory"
12 );
```

## 解密

salsa20+aes\_ecb无魔改

中间加了简单的倒序xor处理去除一些内容特征

```
1 int64 __fastcall sub_22F9(int64 a1, unsigned int64 i_1, int64 a3)
2 {
3     int64 result; // rax
4     unsigned int64 i; // [rsp+20h] [rbp-8h]
5
6     for ( i = 0LL; i < i_1; ++i )
7         *(_BYTE*)(a3 + i) = *(_BYTE*)(i_1 - i - 1 + a1) ^ 0xAA;
8     result = a3 + i_1;
9     *(_BYTE*)(a3 + i_1) = 0;
10    return result;
11 }
```

eg.cyber\_chef

**Recipe**

**AES Decrypt**

Key: venom2025v... (UTF8)

IV: (HEX)

Mode: ECB

Input: Hex

Output: Raw

**Salsa20**

Key: 0ff71ec39a5... (HEX)

Nonce: salsa\_iv (UTF8)

Counter: 0

Rounds: 20

Input: Raw

Output: Raw

**Input**

```
59cc9e17b97199cad788f5f9d919531ee9f09bb233b5565233aef  
c|
```

**Output**

```
vctf{ab92105b_047e_aa38_0b25_a6989b8e5ed5}
```

eg.code

```
1 import binascii
2 import struct
3 from Crypto.Cipher import AES,Salsa20
4
5 #aes_dec
6 # enc = [
7 #     0xc9, 0xcf, 0xce, 0xc9, 0xce, 0xcb, 0x92, 0xcb, 0xc9, 0x99, 0x9f, 0x
8 c9, 0x9f, 0xcb, 0xc9, 0x93,
9 #     0xcc, 0xc9, 0x9d, 0x9e, 0xce, 0x93, 0x9b, 0x92, 0xc9, 0x99, 0x9a, 0x
10 9d, 0x9d, 0x99, 0x98, 0x9f,
11 #     0x99, 0x9c, 0x9c, 0xce, 0x9b, 0xcb, 0x93, 0x99, 0xcf, 0xce, 0x9b, 0x
12 cc, 0xcf, 0xcb, 0x99, 0x99,
13 #     0x98, 0x9f, 0x9c, 0x9f, 0x9f, 0xc8, 0x99, 0x99, 0x98, 0xc8, 0xc8, 0x
14 93, 0x9a, 0xcc, 0x93, 0xcf,
15 #     0xcf, 0x9b, 0x99, 0x9f, 0x93, 0x9b, 0x93, 0xce, 0x93, 0xcc, 0x9f, 0x
16 cc, 0x92, 0x92, 0x9d, 0xce,
17 #     0xcb, 0xc9, 0x93, 0x93, 0x9b, 0x9d, 0x93, 0xc8, 0x9d, 0x9b, 0xcf, 0x
18 93, 0xc9, 0xc9, 0x93, 0x9f
19 # ]
20 # hoge=bytes(b ^ 0xAA for b in enc[::-1])
21 # print(hoge)
22 # b'59cc9e17b97199cad788f5f9d919531ee9f09bb233b5565233aef1de39a1d663523770
23 3c819d47cf9ca5c53ca8adcdec'
24 hoge=bytes.fromhex('59cc9e17b97199cad788f5f9d919531ee9f09bb233b5565233aef1
25 de39a1d6635237703c819d47cf9ca5c53ca8adcdec')
26 aes_key=b"venom2025venom25"
27 cipher_aes=AES.new(aes_key, AES.MODE_ECB)
28 pad=cipher_aes.decrypt(hoge)
29 pad_len=pad[-1]
30 #去除PKS7填充
31 enc_tmp=pad[:-pad_len]
32
33 #salsa20_dec
34 # nonce=0xD9CBC6D9CBF5C3DC
35 # salsa=[0x999D939CCF9F9B9F,0x9A93C89CC898CF93,0x93CB9FCE9998989D,0x9ACCCC
36 9D9BCFC999,
37 #     0x999D939CCF9F9B9F,0x9A93C89CC898CF93,0x93CB9FCE9998989D,0x9ACCCC
38 9D9BCFC999]
39 # hoge1=[]
40 # hoge2=[]
41 # hoge1.extend(nonce.to_bytes(8, 'little'))
42 # for qword in salsa:
43 #     hoge2.extend(qword.to_bytes(8, 'little'))
44 # nonce=bytes(b ^ 0xAA for b in hoge1[::-1])
```

```
35 # salsa_key=bytes(b ^ 0xAA for b in hoge2[::-1])
36 # print(nonce,salsa_key)
37 # b'salsa_iv'
38 # b'0ff71ec39a5d322709b6b2e93796e5150ff71ec39a5d322709b6b2e93796e515'
39 nonce=b'salsa_iv'
40 salsa_key=bytes.fromhex('0ff71ec39a5d322709b6b2e93796e5150ff71ec39a5d32270
9b6b2e93796e515')
41 cipher_salsa=Salsa20.new(key=salsa_key,nonce=nonce)
42 flag=cipher_salsa.decrypt(enc_tmp)
43 print(flag.decode())
44 # vctff{ah92105h_047e_aa38_0b25_a6989h8e5ed5}
```