Mathematical Foundations
of
Computer Science

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9 Hamilton Cycles, Hamilton Paths, and Non-isomorphic Trees

- Homework assignment published on Monday 2018-04-23

- Submit your first solution by Sunday, 2018-04-29, 18:00, by email

- Submit your final solution by Sunday, 2018-05-06.

9.1 Regular Orientations of a Regular Graph

We call a graph *d-regular* if every vertex has degree d. A directed graph is (d, d)-regular if every vertex has d incoming and d outgoing edges.

Exercise 9.1. Show that in every 4-regular graph, you can orient the edges such that every vertex has two incoming and two outgoing edges, i.e., such that the resulting digraph is $(2, 2)$-regular. See the picture below for an illustration.
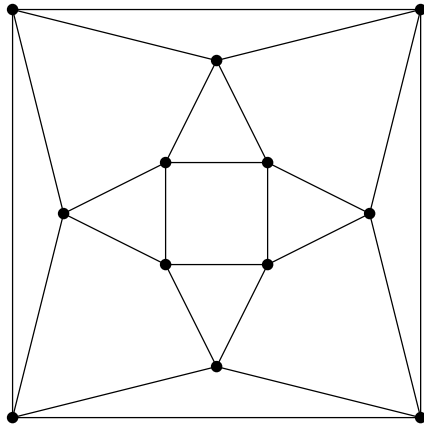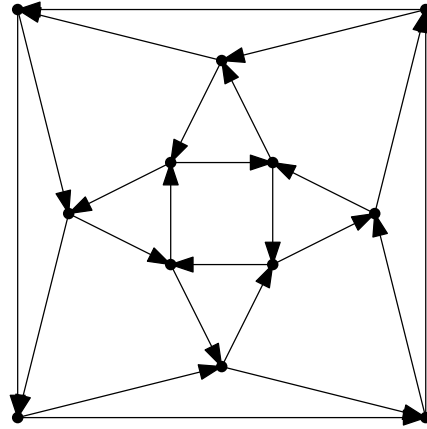
a 4-regular graph   a (2,2)-regular orientation

Proof.   □

1. Suppose that there exits a 4-regular graph, we can't orient the edges such that some vertices has two incoming and two outgoing edges.

2. Assume there are 2 vertices u, v not meeting the condition. Since $\sum\limits_{i=1}^{n} deg(v)$ is even, let u has 3 incoming edges and v has 3 outgoing edges.

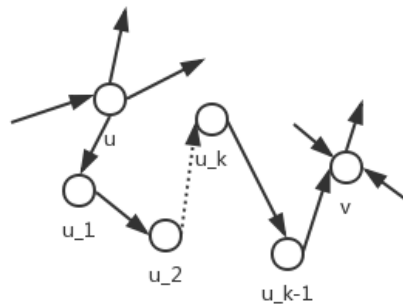   Then along the arrow find a path from u to v as Figure9-1-1 shows.

Figure9-1-1

The path is $u \rightarrow u_1 \rightarrow u_2 \rightarrow ... \rightarrow u_{k-1} \rightarrow u_k \rightarrow v$.

Then change the direction of every edges in the path.

Then the path is changed into $\leftarrow u_1 \leftarrow u_2 \leftarrow ... \leftarrow u_{k-1} \leftarrow u_k \leftarrow v.$
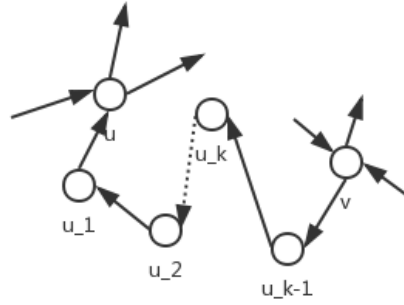
The graph is as Figure9-1-2 shows.

Figure9-1-2

Then clearly for $u_1, u_2...u_k$,one of their incoming edges is changed into outgoing edges and one of their outgoing edges is changed into incoming edges.They still have two incoming and two outgoing edges.

For u, v they now each have two incoming and two outgoing edges as well.

The resulting digraph is (2,2)-regular.

3. If there is 2 vertices and one have 4 incoming edges and another have 4 outgoing edges.Just do the operation above twice.

   If there is more than 2 vertices not meeting the condition. Just group them by 2.And do the same operation on every group.Then we get a (2,2)-regular orentation.

4. Above all,t in every 4-regular graph, you can orient the edges such that every vertex has two incoming and two outgoing edges, i.e., such that the resulting digraph is (2,2)-regular.

9.2   Hamilton Cycles and Ore's Theorem

Consider K_n, the complete graph on n vertices. For $n \geq 3$, this obviously has a Hamilton cycle. How many edges do you have to delete from K_n to

destroy all Hamilton cycles? That is, what is the smallest set S such that $\left(V, \binom{V}{2} \setminus S\right)$ has no Hamilton cycle? Let s_n denote the size of this set (this depends on n, thus the notation s_n). For example, $s_2 = 0$ since K_2 has no Hamilton path to begin with; $s_3 = 1$ since removing one edge from K_3 results in a graph without a Hamilton cycle.

Exercise 9.2. Find a closed formula for s_n and prove it! **Hint.** One part will be easy. For the other part, use Ore's Theorem.

Proof. □

1. For a graph.If it has no Hamilton cycles. There must exit two parts that we can explore one then the other but can't travel back.

2. when we have two disconnected parts, there must be no Hamilton cycles.Let the two partation be U, V,with x vertices $\{u_1, u_2...u_x\}$ in U and y vertices $\{v_1, v_2...v_y\}$ in V. Define $|V| = n$. Clearly we have

$$x + y = |V| = n.$$

3. To find the smallest set S such that $\left(V, \binom{V}{2} \setminus S\right)$ has no Hamilton cycle, we must add as many edges as possible to the graph.Then we have the following steps to do.

   (a) Add edges in U, V separately so that U, V are both complete graphs.

   (b) Add edges between U and V as many edges as possible making sure not existing Hamilton cycle.

   (c) Use x, y to expessing the number of edges we add. With $x + y = n$,computing in which case we have the largest number of edges added. Write it as M_n

   (d) $S_n = |V| - M_n$.

4. Now we have the graph bellow for an illustration.
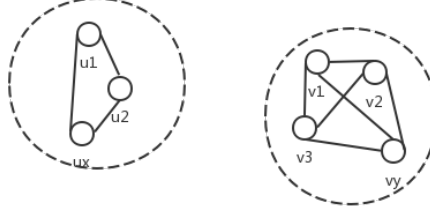
Two partiation

When U, V are both complete graph,

$$|U| = \binom{x}{2} = \frac{x(x-1)}{2} \tag{1}$$

$$|V| = \binom{y}{2} = \frac{y(y-1)}{2} \tag{2}$$

5. From the definition of Hamilton cycle.We know in a graph, if we explore from a vertex, after a cycle we will come back.

   (a) If there is only a vertex in U(or V) has edges to the other. ($|U| > 1, |V| > 1$)

   Take U as an example. There is a vertex u_i having edges to vertices in V. Just start exploring from u_i.

   If we exploring vertices in U first, we must return to u_i to elpore V.u_i is explored twice.

   If we exploring vertices in V first, after exploring all vertices in V, we must return to u_i to elpore U.u_i is explored twice as well.

   There is no Hamilton cycles.

   Now we can add $maxx, y$ edges between U, V at most.

   (b) If there are two vertices in U having edges to 2 vertices in V.

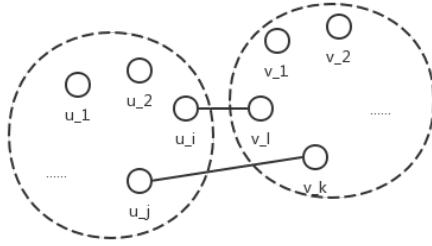   Since U, V are both complete graph,

   $$(u_i, u_j) \in |E|$$
   $$(v_i, v_k) \in |E|$$

   We can find Hamilton cycles

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$$u_j \to \ldots \to u_i \to u_j$$
$$v_l \to \ldots \to v_k \to v_l$$

As the figure bellow shows, we can easly find a cycle

$$\underbrace{u_j \to \ldots \to u_i}_{\text{Hamilton cycle of } U} \quad \to \quad \overbrace{v_l \to \ldots \to v_k}^{\text{Hamilton cycle of } V} \quad \to \quad u_j$$

Thus between U and V, if there are more than 2 vertices in one partation have edges to vertices the other, they must be connected to the same vertex.

6. So

$$M_n = max\{|U| + |V|\} + max\{x, y\} \tag{3}$$

$$S_n = |V| - M_n \tag{4}$$
$$= max\{\binom{n}{2} - \binom{x}{2} - \binom{y}{2}\} - max\{x, y\} \tag{5}$$
$$\tag{6}$$

Since

$$x + y = n \tag{7}$$

$$S_n = max\{x(n - x)\} - max\{x, y\} \tag{8}$$

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If n is even,

$$S_n = \frac{n^2 - 2n}{4} \tag{9}$$

If n is odd,

$$S_n = \frac{n^2 - 2n - 3}{4} \quad (n > 3) \tag{10}$$

7. Above all

$$S_n = \begin{cases} 0 & (n = 1 \ or \ 3) \\ \dfrac{n^2 - 2n - 3}{4} & (n > 3 \ and \ n \ is \ odd) \\ \dfrac{n^2 - 2n}{4} & (n \ is \ even) \end{cases}$$

A *tournament* is a directed graph in which, for each pair $u, v \in V$, exactly one of the directed edges (u, v) and (v, u) is in the graph. Imagine a sports tournament in which every participant plays against every other exactly once. Draw an arc from u to v if u beat v in this tournament.
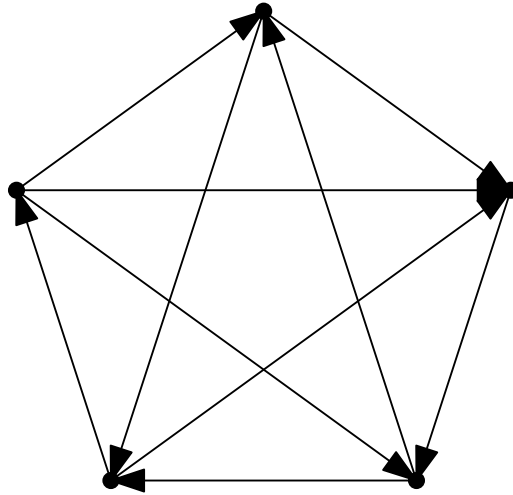
A tournament on five vertices.

Exercise 9.3. Show that every tournament has a *directed Hamilton path*, i.e., a sequence $u_1, u_2, \ldots, u_n$ such that $(u_i, u_{i+1}) \in E$ for all $i = 1, \ldots, n - 1$. See the picture below.

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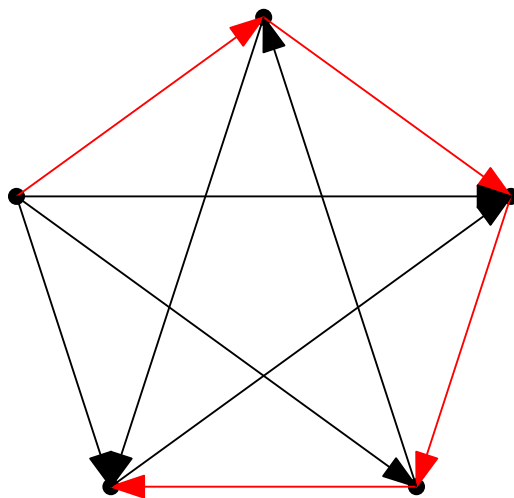
The same tournament wiht a Hamilton path.

You probably won't be able to use the proof of Ore's Theorem directly, but you can use the proof idea.

Proof. □

1. We can use an induction to prove it.

2. Base case: when there are 2 vertices, clearly there is a Hamilton path.See the picture bellow.

3. Now a graph with k vertices has a Hamilton path. Add a vertex u to it.

   Along the path, numbered the vertices as their explored order as $(v_1, v_2...v_k)$.

4. If the edge between v_k and u is (v_k, u) ,add u as the tail of the path.Then we get $(v_1, ...v_k, u)$.

5. Else turn to $v_k - 1$, if there is (v_{k-1}, u), clearly we can insert u into the path:$(v_1, v_2...v_k, u, v_{k+1}, v_k)$.

6. Else turn to the vertex before continously until there is (v_i, u), Clearly we can insert u into the path:$(v_1, v_2...v_i, u, v_{i+1}, ..., v_k)$.

7. Else all edges are from from u to v_i, add u as the head of the path.Then we get $(v_1, ...v_k, u)$ or $(u, v_1, ...v_k)$.

8. Above all, every tournament has a *directed Hamilton path*.

9.3   Isomorphism Classes of Trees

In the lecture (and in the videos) we have seen that the number of trees on vertex set $V = \{1, 2, \ldots, n\}$ is n^{n-2}. This however ignores isomorphisms. For example, there are $3^{3-2} = 3$ trees on vertex set $\{1, 2, 3\}$, but all those trees look alike (are isomorphic). On $\{1, 2, 3, 4\}$, there are 16 trees, but there are only two isomorphism classes: the path and the star. For five vertices, there are 125 trees but only three isomorphism classes: the path, the star, and the "T-shape" (see video on counting the number of trees). For $n = 6$ we get the path, the Y-shape, the Euro symbol, the Star Wars fighter, the Scandinavian cross, and the star, so six isomorphism classes (but a total of 1296 trees).
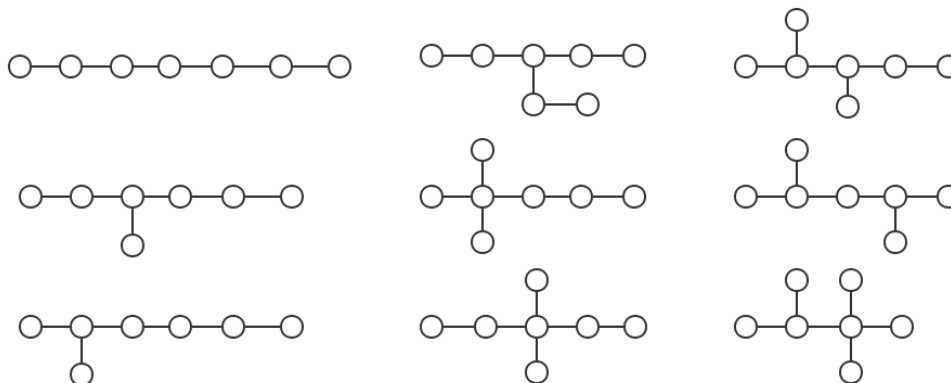
Exercise 9.4. List of isomorphism classes on seven vertices. That is, draw trees $T_1, \ldots, T_m$ on seven vertices such that no two of them are isomorphic but every tree on seven vertices is isomorphic to one of them. How many do you get?

n	1	2	3	4	5	6	7
number of isomorphism classes	1	1	1	2	3	6	?

Solution. ∎

We get 9 different trees.

$T_1, T_2 \ldots T_9$

Alright, so let's denote by t_n the number of isomorphism classes of trees on n vertices. That is, t_n is the largest number m such that we can find trees $T_1, \ldots, T_m$ on n vertices such that no two of them are isomorphic. We would like to have an exact and explicit formula for t_n, but that is probably too much to ask for. Instead, let us try to understand t_n approximately and asymptotically.
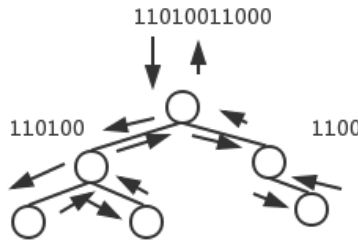
Exercise 9.5. Show that $t_n \leq 4^n$. Hint: Consider the video on the isomorphism problem on trees. It defines a way to encode a tree as a $0/1$-sequence.

Proof. □

Encode a tree as $0/1$-sequence in the way below.
If there is a vertex ,we write 1 to explore it and then write 1 to return.
Take the following tree as an example.

Then if we have n vertices, we have a sequence of 2n bits.
There are $2^{2n} = 4^n$ trees.
However, in a tree if the leftchild tree of a node is the same as the leftchild tree of the same node, they are isomorphism.BUt they have different $0/1$-sequences.
Thus $t_n \leq 4^n$

Exercise 9.6. Show that $t_n \geq \frac{e^n}{\text{poly}(n)}$, where $\text{poly}(n)$ is some polynomial in n. Hint: There are n^{n-2} trees on $V = [n]$. We group them together in "buckets" of isomorphic trees. How large can a bucket be? Answer this and then use Stirling's approximation for $n!$.

**Exercise 9.7. *Try to improve those bounds. That is, find some $a < 4$ such that $t_n \in O(a^n)$ or some $b > e$ such that $t_n \in \Omega(b^n)$. Any improvement will be kind of interesting. Aim for simple proofs!*

Remark. The "true" rate of growth is known by a result of George Pólya but apparently it is quite difficult (I write "apparently" because I have never studied this work).