

18. Spoofing & Firewalls & IDS

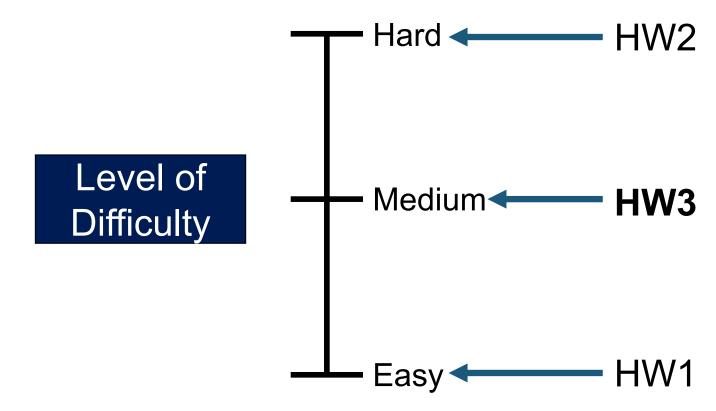
Seongil Wi



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HW3 is Released!

- Related to web security
- CTF-style homework (5 problems)
- Last homework! ☺



Recap: Computer Network

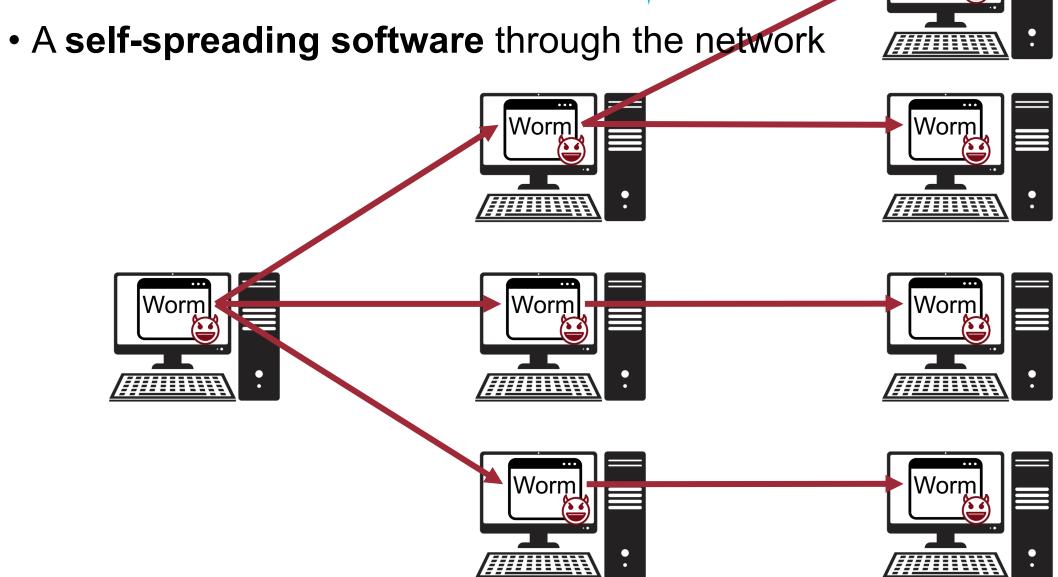
- A telecommunications network that allows computers to exchange data
- Networked computing devices pass data to each other along data connections



Worm

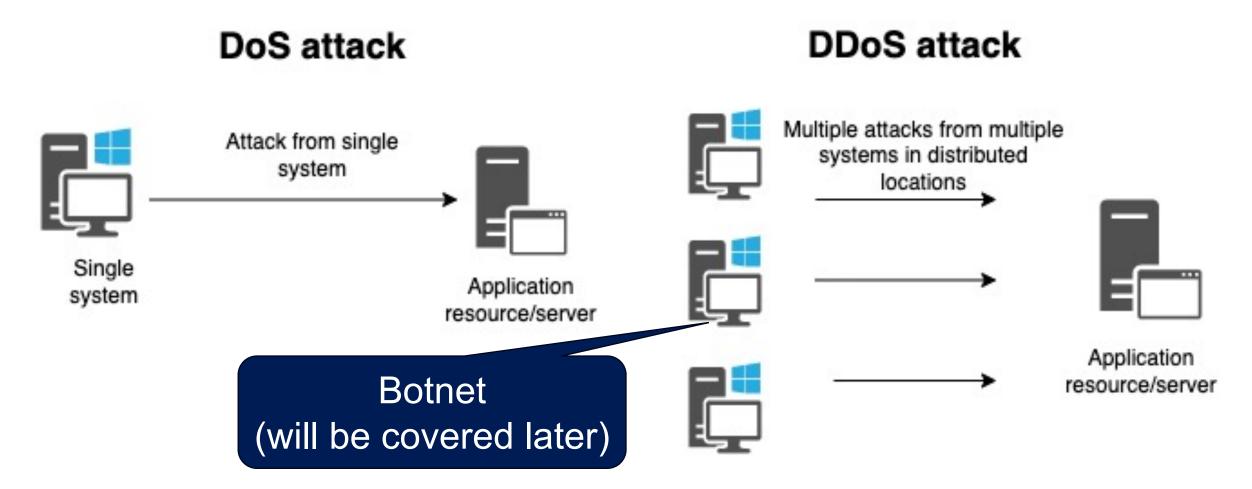
Recap: Worm





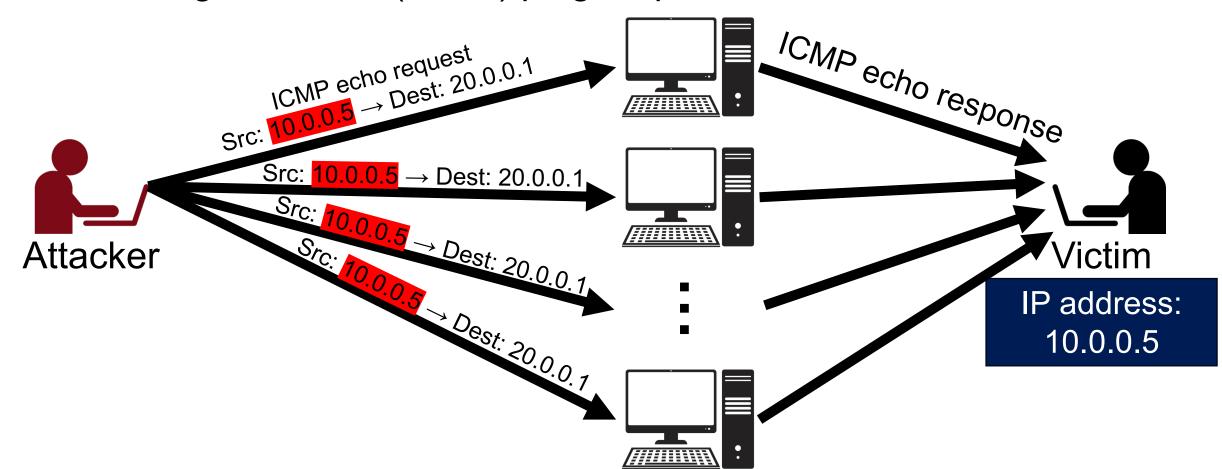
Recap: Distributed Denial-of-Service (DDoS)

 Employ multiple (compromised) computers to perform a coordinated and widely distributed DoS attack



Recap: Ping Flood Attack

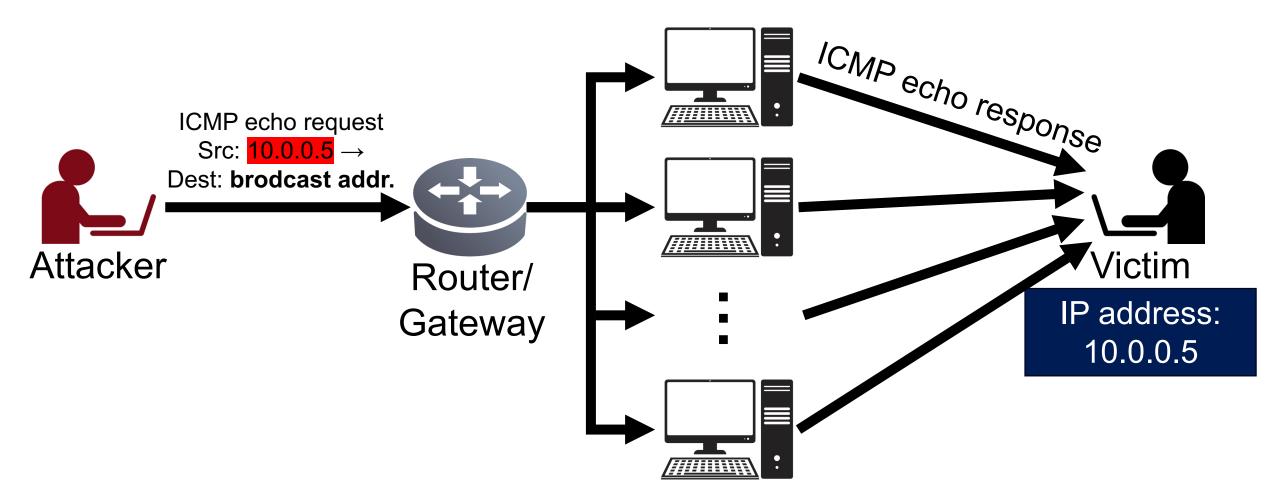
 The computing device is flooded with tons of Internet Control Message Protocol (ICMP) ping response



Recap: Smurf Attack

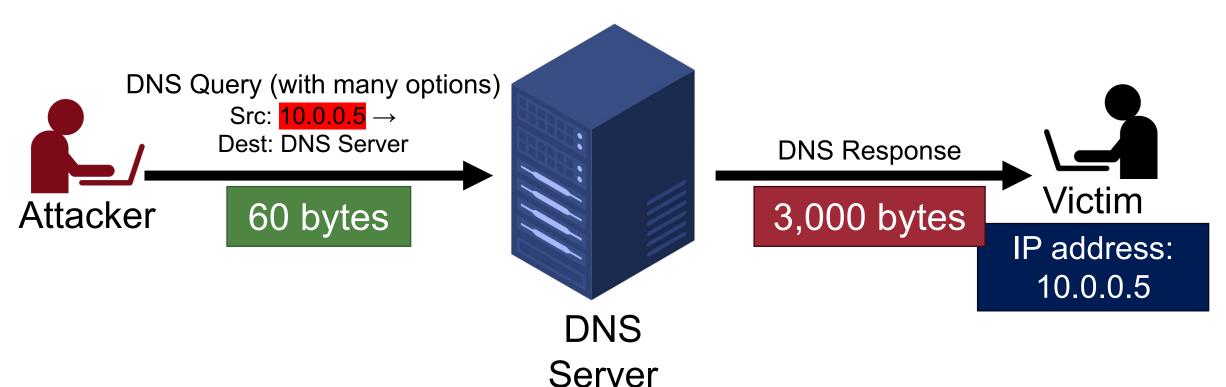
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• Idea: sending ping request to broadcast address



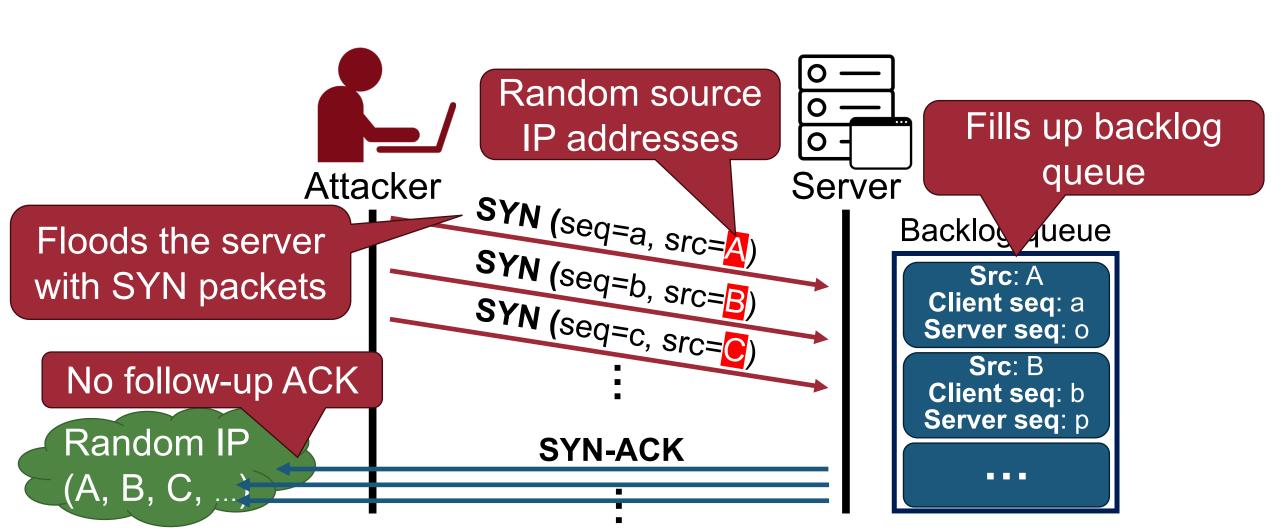
Recap: Amplification Attack

- Idea: controlling the size of responses, not the number of responses
- Example: DNS Amplification Attack



Recap: TCP SYN Flooding Attack

Floods the server with SYN Packets



Recap: Mitigation



- Set the Queue Size for TCP Backlog
 - \$ sysctl -w net.ipv4.tcp_max_syn_backlog=1024
 - ✓ Limitation: Arms race! Attackers can easily win
- Set the Firewalls
 - E.g., Blocks if similar packets exceed 10 per second

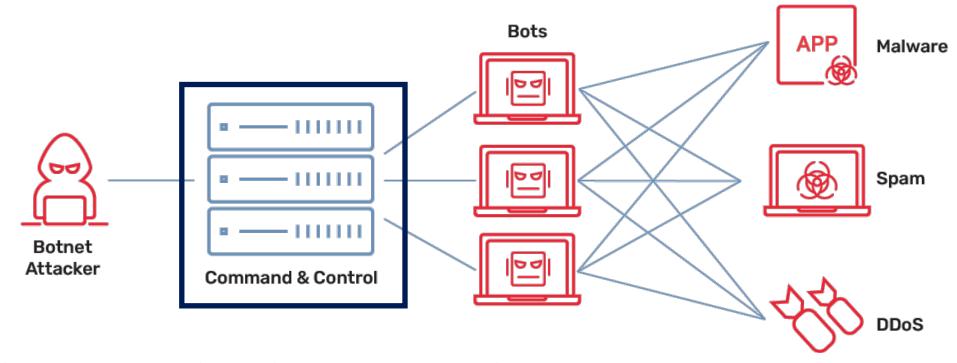
```
$ iptables -A INPUT -p TCP --dport 80 --syn -m limit 10/second -j ACCEPT
$ iptables -A INPUT -p TCP --dport 80 --syn -j DROP
```

✓ Limitation: Performance

SYN Cookie

Recap: Command and Control (C&C) Server

- Essential for operation and support of botnet
- Two styles
 - Centralized
 - -P2P



Spoofing

Spoofing



 A situation in which a person or program is successfully identified as another by modifying data

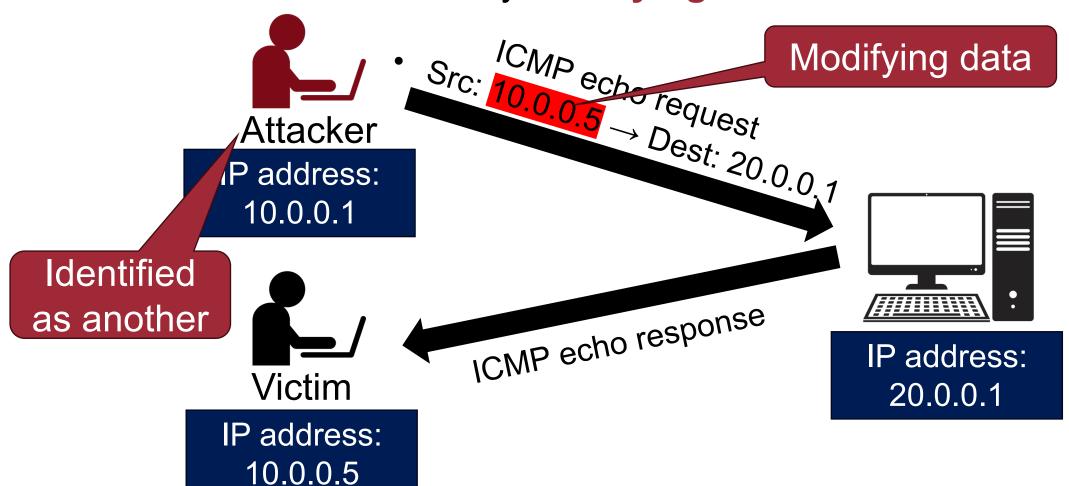
- IP spoofing
- ARP spoofing
- DNS spoofing

•

Recap: IP Spoofing



 A situation in which a person or program is successfully identified as another by modifying data



Recap: IP Spoofing Method

- How to
 - Network RAW socket programming
 - Nmap
 - My own IP address is 10.0.0.1
 - \$ nmap -e eth0 -S 10.0.0.5 20.0.0.1
 - Use the network interface eth0 to send a spoofed packet (10.0.0.5) to 20.0.0.1

– . . .

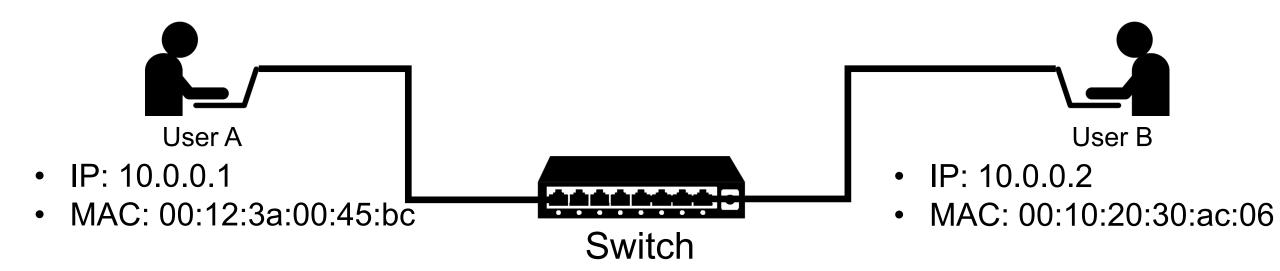
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An attacker sends spoofed Address Resolution Protocol (ARP) response messages

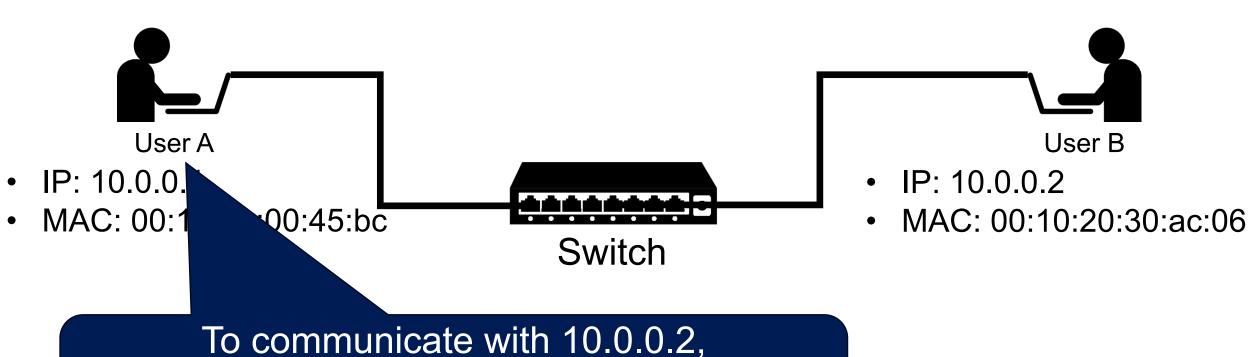
Also known as ARP cache poisoning attack

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 A protocol that maps an ever-changing IP address to a fixed physical machine address (MAC address)



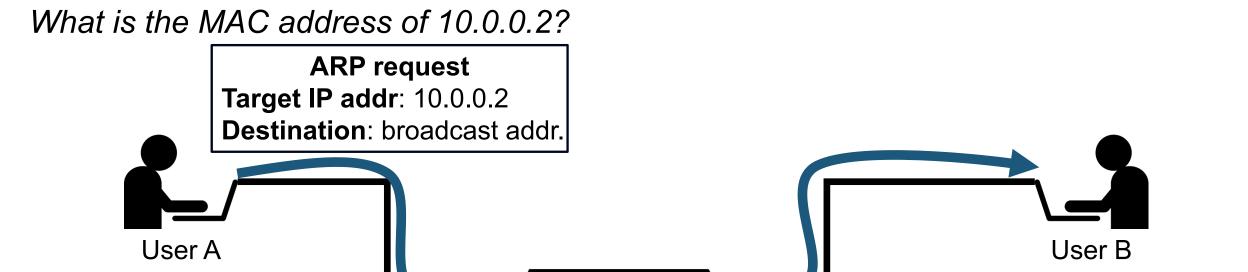
 A protocol that maps an ever-changing IP address to a fixed physical machine address (MAC address)



To communicate with 10.0.0.2, User A needs to know User B's MAC address → Use ARP protocol

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Address Resolution Protocol (ARP)



Switch

- IP: 10.0.0.1
- MAC: 00:12:3a:00:45:bc

• IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

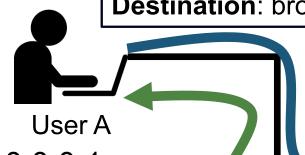


What is the MAC address of 10.0.0.2?



Target IP addr: 10.0.0.2

Destination: broadcast addr.



• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc



IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

User B

ARP response

My IP Addr: 10.0.0.2

My MAC addr: 00:10:20:30:ac:06

Destination: User A

My MAC address is 00:10:20:30:ac:06

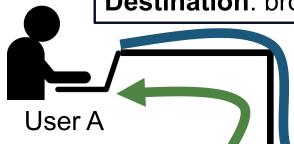
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What is the MAC address of 10.0.0.2?

ARP request

Target IP addr: 10.0.0.2

Destination: broadcast addr.



• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06

Switch

IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

User B

ARP response

My IP Addr: 10.0.0.2

My MAC addr: 00:10:20:30:ac:06

Destination: User A

My MAC address is 00:10:20:30:ac:06

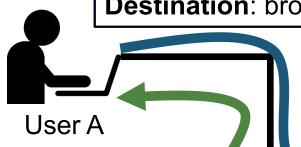


What is the MAC address of 10.0.0.2?

ARP request

Target IP addr: 10.0.0.2

Destination: broadcast addr.



IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr Mac Addr

10.0.0.2 00:10:20:30:ac:06 **Switch**

IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

User B

ARP response

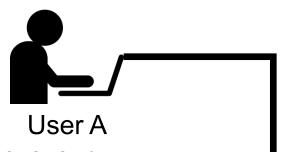
10.0.0.2 User A knows the

dr: 00:10:20:30:ac:06 IP-MAC addr. mapping info.

: User A

My MAC address is 00:10:20:30:ac:06



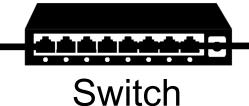


• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06



User B

• IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc







Through WIFI or Ethernet

• IP: 10.0.0.3

• MAC: 00:01:12:44:3a:6c

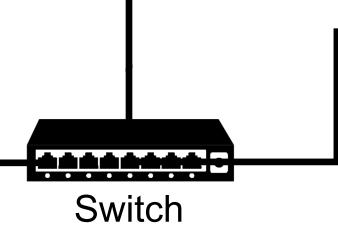


• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06



• IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

User B

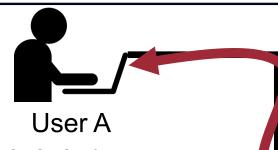
IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc

ARP response

My IP Addr: 10.0.0.2

My MAC addr: 00:01:12:44:3a:6c

Destination: User A



• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06



IP: 10.0.0.3

Switch

MAC: 00:01:12:44:3a:6c



User B

IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

User B - ARP	cacne

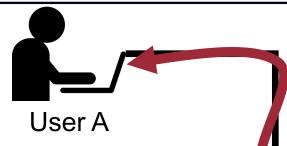
IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc

ARP response

My IP Addr: 10.0.0.2

My MAC addr: 00:01:12:44:3a:6c

Destination: User A



• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06



IP: 10.0.0.3

MAC: 00:01:12:44:3a:6c



• IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

Switch

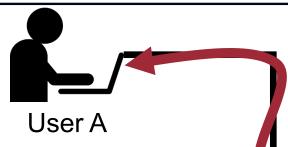
IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc

ARP response

My IP Addr: 10.0.0.2

My MAC addr: 00:01:12:44:3a:6c

Destination: User A



• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06
	00:01:12:44:3a:6c



IP: 10.0.0.3

Switch

MAC: 00:01:12:44:3a:6c



• IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

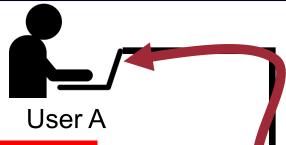
IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc

ARP response

My IP Addr: 10.0.0.2

My MAC addr: 00:01:12:44:3a:6c

Destination: User A



• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06 00:01:12:44:3a:6c



IP: 10.0.0.3

MAC: 00:01:12:44:3a:6c



My IP Addr: 10.0.0.1

My MAC addr: 00:01:12:44:3a:6c

Destination: User B



User B

IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

Switch

IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc



ARP response

My IP Addr: 10.0.0.2

My MAC addr: 00:01:12:44:3a:6c

Destination: User A



• IP: 10.0.0.1

MAC: 00:12:3a:00:45:bc

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06
	00:01:12:44:3a:6c



IP: 10.0.0.3

MAC: 00:01:12:44:3a:6c



My IP Addr: 10.0.0.1

My MAC addr: 00:01:12:44:3a:6c

Destination: User B



User B

• IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

Switch

IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc
	00:01:12:44:3a:6c





Man-in-the-Middle (MITM) attack

• IP: 10.0.0.3

• MAC: 00:01:12:44:3a:6c

User B

• IP: 10.0.0.1

MAC: 00:12:3a:00:45

User A

• IP: 10.0.0.2

MAC: 00:10:20:30:ac:06

User A - ARP cache

IP Addr	Mac Addr
10.0.0.2	00:10:20:30:ac:06 00:01:12:44:3a:6c

Switch

IP Addr	Mac Addr
10.0.0.1	00:12:3a:00:45:bc 00:01:12:44:3a:6c

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ARP Spoofing – Use Case

Home > 전체기사

Public WIFI is vulnerable to MITM attacks (via ARP spoofing)

카페서 와이파이 쓰다 스마트폰 가로채기 당한다?

입력: 2015-10-29 11:30



모바일에서 공유기 통해 외부 접근시 공격자가 스마트폰 정보 가로채기 가능

[보안뉴스 김경애] 올해 1분기 주요 보안 이슈로 떠오른 바 있는 공유기 보안 문제가 좀처럼 해결되지 않고 있다. 사람들이 이용하는 카페 등 공공장소에서 공유기를 통해 무료 와이파이를 이용할 경우 모바일을 이용해 정보를 가로챌 수 있는 ARP 스푸핑(Spoofing) 공격이 가능한 것으로 드러났기 때문이다. 이에 따라 공공장소나 카페에서 와이파이 접속시 이용자들의 주의가 요구된다.

Not recommended to use public WIFI!

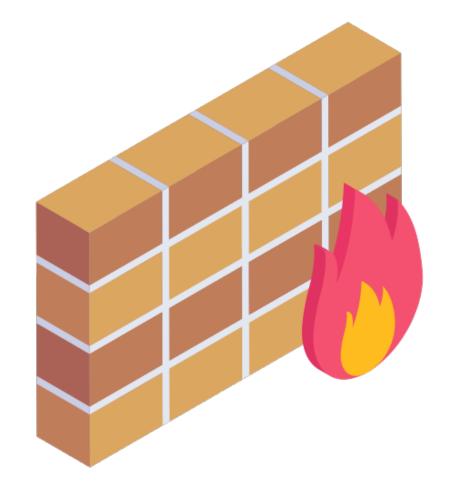
How to Mitigate ARP Spoofing Attacks?



- Packet encryption
 - E.g., Visit websites via HTTPs protocol
- Static ARP caches
 - IP-MAC address mappings in the local ARP cache may be statically entered

- Packet filtering
 - Check if many ARP responses are sent without a request

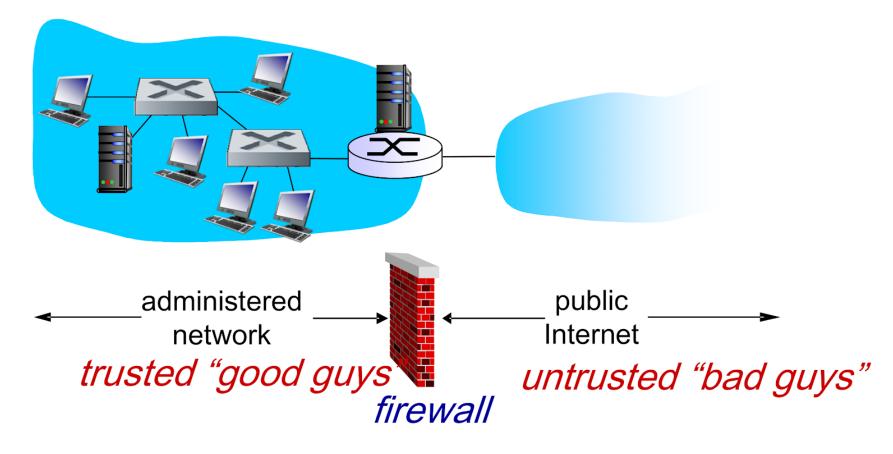
Firewalls **①**



Firewalls

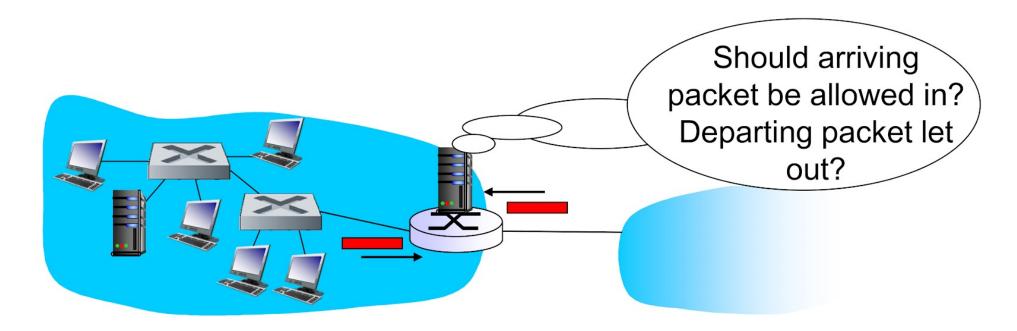


 Isolate organization's internal net from larger Internet, allowing some packets to pass, blocking others



Main Goals of Firewall

- *
- **Isolate** organization's <u>internal net</u> from <u>larger Internet</u>, allowing some packets to pass, blocking others
- Restrict incoming and outgoing traffic by IP address, ports, or users
 - Block invalid traffic and only authorized traffic is allowed



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Types of Firewall



- Three types of firewalls:
 - -Packet filters
 - Application gateways

Packet Filters





- Internal network connected to Internet via router firewall
- Router filters packet-by-packet from the network layer, decision to forward/drop packet based on:
 - Source/destination IP address
 - Source/destination port numbers
 - ICMP message type
 - Protocol status (e.g., TCP SYN, TCP ACK, ...)

Packet Filters via Router Firewall



Application

Presentation

Session

Transport

Network

Data Link

Physical

Check IP address and port numbers

Frame IP TCP HTTP Frame header header header eques footer

Network

Data Link

Physical

Network 2

Application

Presentation

Session

Transport

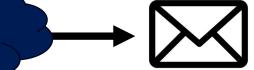
Network

Data Link

Physical



Network 1



Packet Filters – High-level Example

Policy	Firewall Setting		
No incoming web access	Drop all incoming packets to any IP address, port 80		
No incoming TCP connections, except those for institution's public web server only	Drop all incoming TCP SYN packets to any IP except 130.207.244.203, port 80		
Prevent your network from being used for a Smurf DoS attack	Drop all ICMP packets going to a "broadcast" address		

Packet Filters – Detailed Example

Action	Source Address	Dest. Address	Protocol	Source port	Dest. port	Flag bit
Allow	222.22/16	Outside of 222.22/16	TCP	> 1023	80	any
Allow	Outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK
Allow	222.22/16	Outside of 222.22/16	UDP	> 1023	53	
Allow	Outside of 222.22/16	222.22/16	UDP	53	> 1023	
Deny	all	all	all	all	all	all

Packet Filters – Pros and Cons

- Pros
 - Simple to implement
 - Low impact on network performance
- Cons
 - Malicious content in application data cannot be filtered out

Packer Filter: Limitation







Application

Presentation

Session

Transport

Network

Data Link

Physical

Check IP address and port numbers

HTTP **TCP** Frame header header reques footer

Network

Data Link

Physical

Application

Presentation

Session

Transport

Network

Data Link

Physical



Network 1



Network 2

If encrypted, it is impossible to observe

the application-level payload contents



Application Gateways

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Filter packets on application data as well as on IP/TCP/UDP fields

Application Gateway

Filter packets on application data as well as on IP/TCP/UDP fields





Frame IP TCP HTTP Frame header header header request footer



Application

Presentation

Session

Transport

Network

Data Link

Physical

Application

Presentatic n

Session

Transport

Network

Data Link

Physical

Application

Presentation

Session

Transport

Network

Data Link

Physical



Network 1

Network 2



Application Gateways – Pros and Cons

- Pros
 - Tend to be more secure than packet filters because they can examine every layer of the communication
- Cons
 - Complex to implement
 - High impact on network performance

Intrusion Detection System (IDS) 🛈

Intrusion Detection



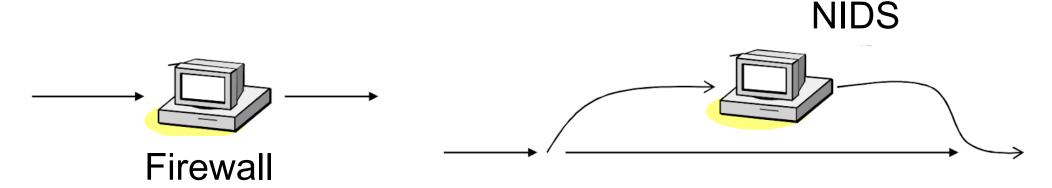


- Intrusion
 - A set of actions aimed to compromise the security goals
- Intrusion detection
 - The process of identifying and responding to intrusion activities



Firewall vs. Intrusion Detection System (IDS)

- Firewall
 - Active filtering (prevent intrusion)
 - Location: Between networks (if an attack is from inside the network it doesn't signal)
- IDS
 - Passive monitoring (detect intrusion)
 - Location: Inside the network



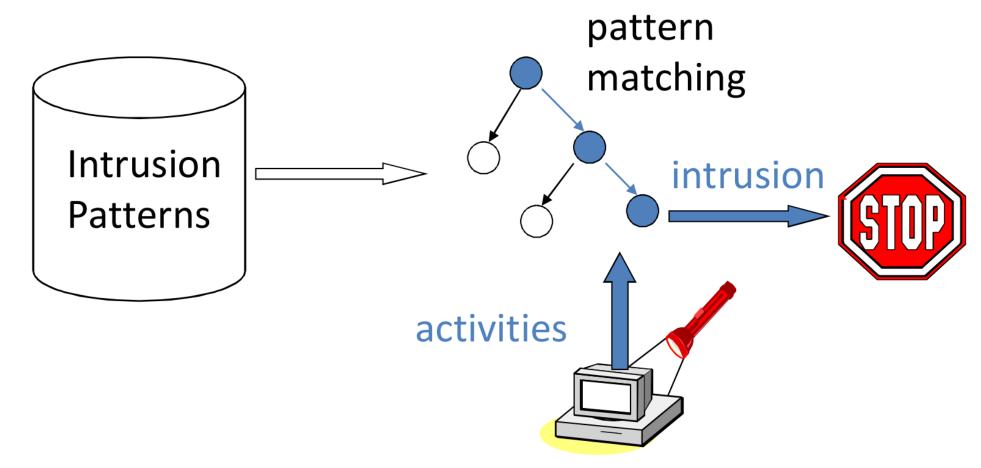
Detection Methods of IDS

1. Signature-based detection

2. Anomaly-based detection

IDS Type #1. Signature-based Detection

- Detects the attacks on the basis of the specific patterns
- E.g., Snort, Bro



IDS Type #1. Signature-based Detection







Example: Snort





Signature example:

```
alert tcp 192.168.2.0/24 23 -> any any \
     (content: "naver" msg: "NAVER Detected!!")
```

Example: Snort





```
Action Protocol src. IP src. port dest. IP dest. port

ignatur example:

alert tcp 192.168.2.0/24 23 -> any any \
(content: "naver" msg: "NAVER Detected!!")
```

If the packet contains the string "naver", ...

...log this message

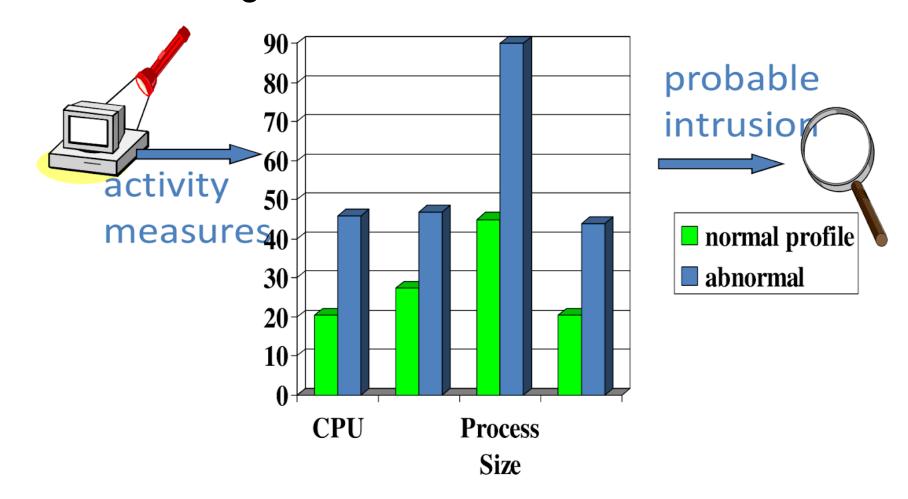
IDS Type #1. Signature-based Detection

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- Detects the attacks on the basis of the specific patterns
- Limitation: difficult to detect new malware attacks as their pattern is not known

IDS Type #2. Anomaly-based Detection

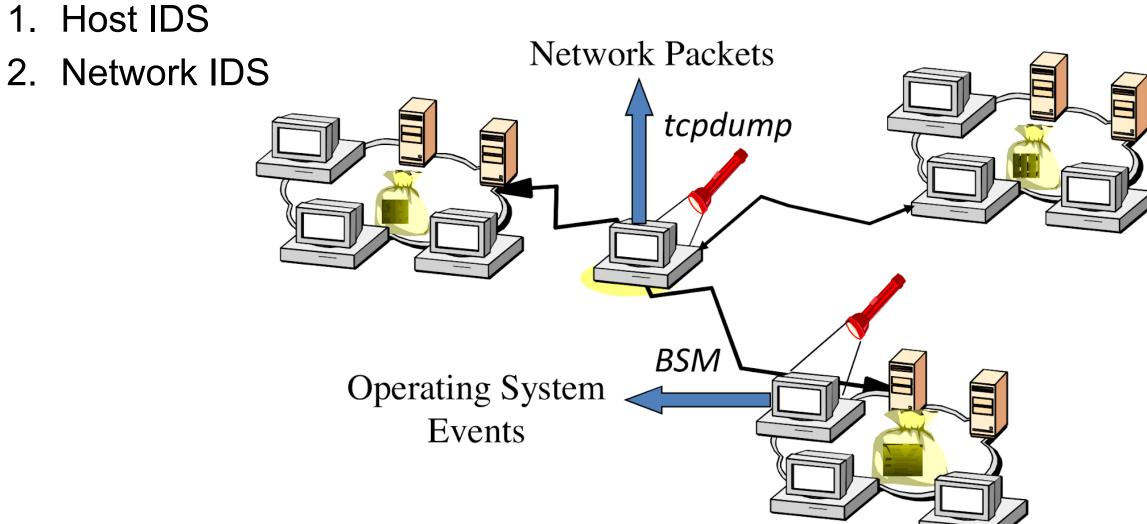
 Detects the attacks on the basis of the statistical models or machine learning models



IDS Type #2. Anomaly-based Detection

- Detects the attacks on the basis of the statistical models or machine learning models
- Limitation: relatively high false positive rate anomalies can just be new normal activities

Deployment of IDS



HIDS (Host IDS)



- Using OS auditing mechanisms
 - e.g., BSM on Solaris: logs all direct or indirect events generated by a user
 - strace for system calls made by a program
- Monitoring user activities
 - e.g., Analyze shell commands
- Monitoring executions of system programs
 - e.g., Analyze abnormal system calls

Limitations:

- Occupies a certain portion of host resources
- An IDS must be installed and operated on every individual host

NIDS (Network IDS)



- Deploying sensors at strategic locations
 - -e.g., Packet sniffing via tcpdump at routers
- Inspecting network traffic
 - Watch for violations of protocols and unusual connection patterns
- Monitoring user activities
 - Look into the data portions of the packets for malicious command sequences

Limitations:

Easily defeated by encryption

Summary



- Spoofing:
 - A person is successfully identified as another by modifying data
- Firewalls:
 - Isolate internal net from larger Internet

- Intrusion Detection System (IDS)
 - The process of identifying and reporting to intrusion activities

Question?