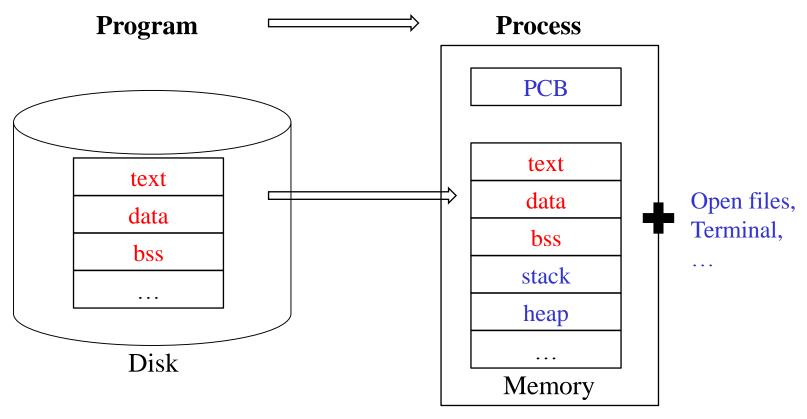
Process control

Process

- a program in execution
- program image + environment
 - text(code) /data / stack / heap segment
 - kernel data structure, address space, open files, ...

```
/* test_program.c */
#include <stdio.h>
int a,b;
int glob_var = 3;
void test_func(void) {
   int local_var, *buf;
   buf = (int *) malloc(10,000 * sizeof(int));
int main(int argc, char *argv[]) {
   int i = 1;
   int local_var;
   a = i + 1;
                                           $ gcc -o test_program test_program.c
   printf("value of a = \%d\n", a);
   test_func();
```

\$./test_program



* bss: Block Started by Symbol

PCB: Process Control Block

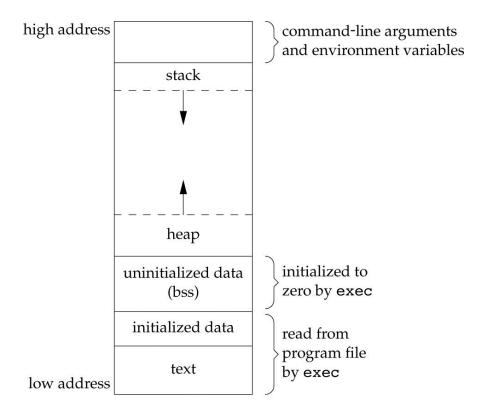
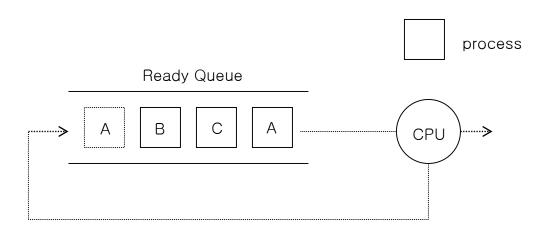


Figure 7.6 Typical memory arrangement

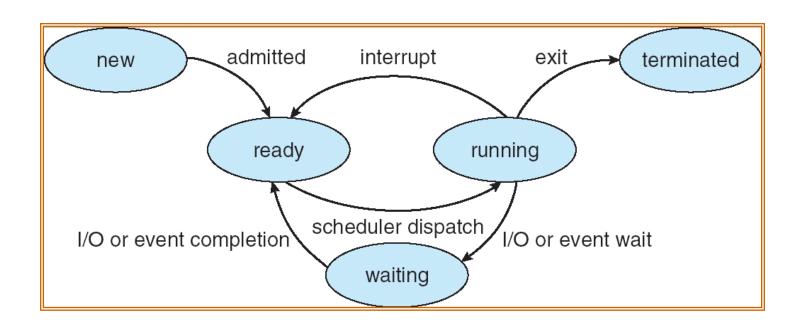
UNIX is a multitasking system.

```
$ ps –e
PID TTY
              TIME CMD
1?
       00:00:00 init
       00:00:00 migration/0
3?
       00:00:00 ksoftirqd/0
4?
       00:00:00 watchdog/0
5?
       00:00:00 migration/1
       00:00:00 ksoftirqd/1
6?
       00:00:00 watchdog/1
8?
       00:00:00 migration/2
29122 pts/9
            00:00:00 bash
29151?
           00:00:00 sshd
29170?
           00:00:00 sshd
29171 pts/10 00:00:00 bash
30465?
           00:00:00 httpd
$
```

UNIX is a time sharing system.



State of a process



Process identifiers

process id

- Every process has a unique process ID.
- As processes terminate, IDs become candidates for reuse.

```
$ ps –е
              TIME CMD
       00:00:00 init
       00:00:00 migration/0
3?
       00:00:00 ksoftirqd/0
       00:00:00 watchdog/0
5?
       00:00:00 migration/1
29122 pts/9
             00:00:00 bash
29151 ?
            00:00:00 sshd
29170 ?
            00:00:00 sshd
29171 pts/10 00:00:00 bash
30465 ?
            00:00:00 httpd
```

Process identifiers

- Process 0 (scheduler process or swapper)
 - Part of the kernel.
 - System process.
- Process 1 (init process)
 - /sbin/init
 - Invoked at the end of the bootstrap procedure.
 - Initialize the UNIX system with /etc/rc* or /etc/inittab.
 - Never dies.
 - A user process with superuser privilege.

Process identifiers

```
#include <unistd.h>
pid_t getpid(void);
                     Returns: process ID of calling process
pid_t getppid(void);
                     Returns: parent process ID of calling process
uid_t getuid(void);
                     Returns: real user ID of calling process
uid_t geteuid(void);
                     Returns: effective user ID of calling process
gid_t getgid(void);
                     Returns: real group ID of calling process
gid_t getegid(void);
                     Returns: effective group ID of calling process
```

```
#include <unistd.h>
```

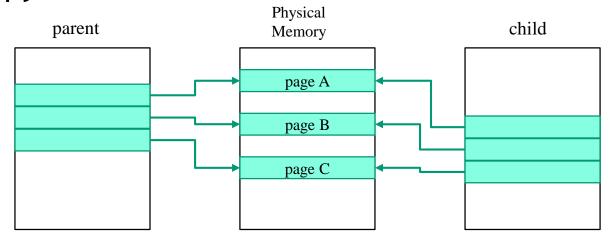
pid_t fork(void);

Returns: 0 in child, process ID of child in parent, -1 on error

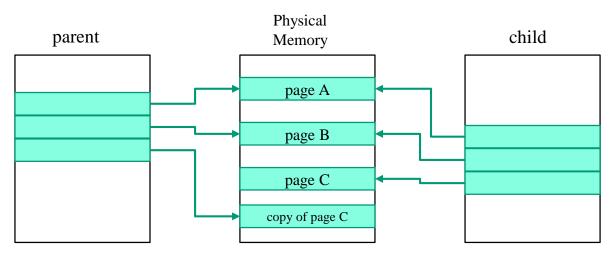
Create a child process.

- It is called once but returns twice.
 - Return value of child: 0
 - Return value of parent: PID of child
- Both child and parent continue executing with the instruction that follows the call to fork().
- The child is a copy of parent.
 - Child gets a copy of the parent's data, heap, and stack.
 - Parent and child often share the text segment.

Copy on write



• After parent process write a page C.



Example

```
#include "apue.h"
     glob = 6; /* external variable in initialized data */
int
char buf[] = "a write to stdout\n";
int
main(void)
                /* automatic variable on the stack */
  int
         var:
  pid_t
         pid;
                                         What about strlen() instead of sizeof()?
  var = 88;
  if (write(STDOUT_FILENO, buf, sizeof(buf)-1) != sizeof(buf)-1)
     err_sys("write error");
  printf("before fork\n"); /* we don't flush stdout */
```

Example(cont.)

```
#include "apue.h"
int glob = 6;
char buf[] = "a write to stdout\n";
int main(void)
  int
         var;
  pid_t pid;
  var = 88:
  if (write(STDOUT_FILENO, buf, sizeof(buf)-1) != sizeof(buf)-1)
     err_sys("write error");
  printf("before fork\n"); /* we don't flush stdout */
  if ((pid = fork()) < 0) {
                                  /* pid is child's pid(non-0). */
     err_sys("fork error");
  } else if (pid == 0) {
                                  /* child */
                                  /* modify variables */
     glob++;
     var++;
   } else {
                                  /* parent */
     sleep(2);
  printf("pid = %d, glob = %d, var = %d\n", getpid(), glob, var);
  exit(0);
```

```
#include "apue.h"
int glob = 6;
char buf[] = "a write to stdout\n";
int main(void)
  int
         var;
  pid_t pid;
  var = 88;
  if (write(STDOUT_FILENO, buf, sizeof(buf)-1) != sizeof(buf)-1)
     err_sys("write error");
  printf("before fork\n"); /* we don't flush stdout */
                                   /* pid is 0. */
  if ((pid = fork()) < 0) {
     err_sys("fork error");
  } else if (pid == 0) {
                                  /* child */
                                  /* modify variables */
     glob++;
     var++;
   } else {
     sleep(2);
                                  /* parent */
  printf("pid = \%d, glob = \%d, var = \%d\n", getpid(), glob, var);
  exit(0);
```

Parent process

Child process



execution

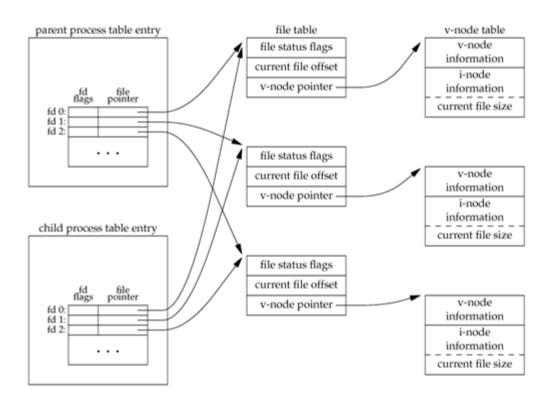
```
$ ./a.out
a write to stdout
before fork
pid = 430, glob = 7, var = 89
pid = 429, glob = 6, var = 88
$ ./a.out > temp.out
$ cat temp.out
a write to stdout
before fork
pid = 432, glob = 7, var = 89
before fork
pid = 431, glob = 6, var = 88
```

We never know if the child starts executing before the parent or vice versa.

- It depends on the process scheduling algorithm.

File sharing

Sharing of open files between parent and child after fork()



- Information shared by child and parent.
 - real-uid(gid), effective-uid(gid)
 - controlling terminal
 - current working directory, root directory
 - signal handlers
 - environment
 - resource limits
- differences between child and parent.
 - the return value from fork.
 - PID and PPID
 - child's resource utilizations are set to 0.
 - pending signals

Two main reasons for fork to fail

- if there are already too many processes in the system.
- if the total number of processes for this real user ID exceeds the system's limit.

Two uses for fork

- A process wants to duplicate itself.
 - Parent and child can each execute different sections of code at the same time.
 - Common for network servers.
- A process wants to execute a different program.
 - Common for shells.
 - Child does exec() right after it returns from fork().

```
#include <unistd.h>
```

pid_t vfork(void);

Returns: 0 in child, process ID of child in parent, -1 on error

- same calling sequence and same return values as fork().
- intended to create a new process when the purpose of the new process is to exec() a new program.
 - Does not copy the address space of parent into the child.
 - The child calls exec() or exit() right after the vfork().
 - The child runs in the address space of the parent.
 - Provides an efficiency.
- vfork() guarantees that the child runs first.

Example

```
#include "apue.h"
int glob = 6; /* external variable in initialized data */
int
main(void)
              /* automatic variable on the stack */
  int var;
  pid_t pid;
  var = 88;
  printf("before vfork\n"); /* we don't flush stdio */
  if ((pid = vfork()) < 0) {
    err_sys("vfork error");
  ellipse = 0  { /* child */
    glob++; /* modify parent's variables */
    var++;
              /* child terminates */
    _exit(0);
```

Example(cont.)

```
/*
 * Parent continues here.
 */
 printf("pid = %d, glob = %d, var = %d\n", getpid(), glob, var);
 exit(0);
}
```

Execution

```
$ ./a.out
before vfork
pid = 29039, glob = 7, var = 89
$
```

exit()

```
#include <stdlib.h>
void exit(int status);
void _Exit(int status);

#include <unistd.h>
void _exit(int status);
```

- causes normal program termination and the value of status is returned to the parent.
- _exit and _Exit
 - return to the kernel immediately.
 - _exit : System call
- exit
 - C standard library
 - performs cleanup processing and then returns to the kernel.
 - all open streams are flushed and closed.

atexit()

```
#include <stdlib.h>
int atexit(void (*func)(void));
```

Returns: 0 if OK, nonzero on error

exit handler

- Function automatically called by exit
- Registered by atexit
- exit() calls the registered handlers in reverse order of their registration with repetition.

Process termination

- Two types of process termination
 - (voluntarily) normal termination
 - return from main (equivalent to calling exit)
 - calling exit
 - calling _exit or _Exit
 - Returning of the last thread from its start routine
 - Calling pthread_exit from the last thread
 - (involuntarily) abnormal termination
 - calling abort
 - Receipt of a signal
 - Response of the last thread to a cancellation request
- Regardless of how a process terminate, the same code is eventually executed.
 - Close open descriptors, release the memory, ...

Process termination

Termination status

- Terminating process notify its parent how it terminated.
- normal termination
 - pass an exit status as argument to exit() or _exit().
 - return value from main()
 - exit status is converted into a termination status.
- abnormal termination
 - Kernel generates a termination status to indicate the reason for the abnormal termination.
- The parent of the terminated process can obtain the termination status from wait() or waitpid().

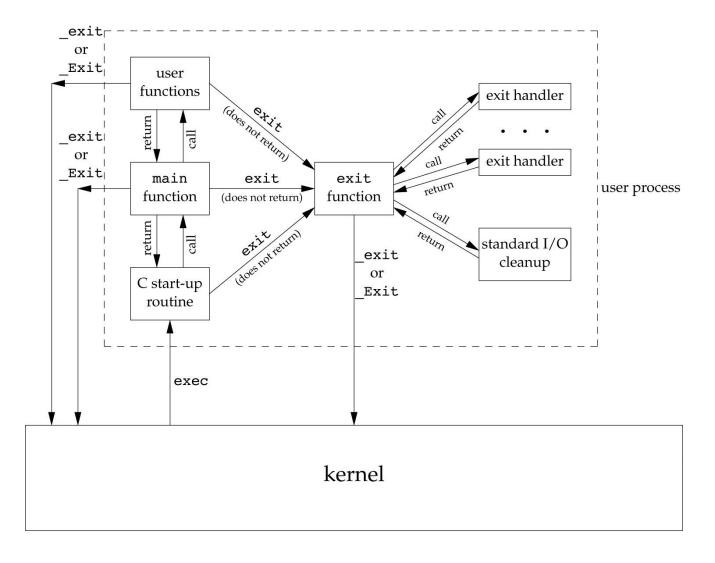


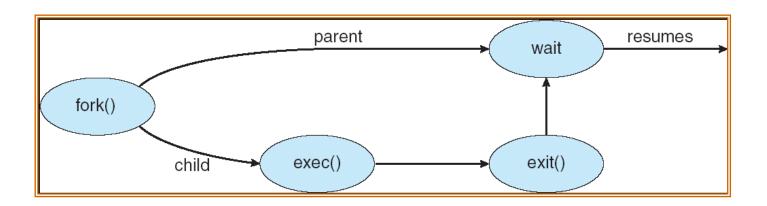
Figure 7.2 How a C program is started and how it terminates

Process termination

zombie state

- Kernel has to keep some information
 - PID, termination status, and CPU usage time
 - This information is available when parent calls wait().
- the process that has terminated, but whose parent has not yet waited for it, is called a zombie.
- If the parent terminates before the child?
 - init process becomes the parent of the orphaned process.
- If an orphaned process is terminated?
 - init calls wait() to fetch the termination status.

- A process that calls wait() or waitpid()
 - Block, if all of its children are still running.
 - Return immediately if a child has terminated.
 - Return immediately with an error, if it doesn't have any child processes.



```
#include <sys/wait.h>
pid_t wait(int *statloc);
Both return: process ID if OK, 0 (see later), or -1 on error
```

wait()

- If a child has already terminated and is a zombie, wait returns immediately with that child's status.
- Otherwise, it blocks the caller until a child terminates.
- If the caller blocks and has multiple children, wait returns when one terminates.

statloc argument

Store termination status in the location pointed to by staloc.

Macro to examine the termination status

- WIFEXITED(status)
 - True if a child terminated normally
- WEXITEDSTATUS(status) (only if WIFEXITED(status) is true)
 - Fetch the low-order 8 bits of the argument the child passed
- WIFSIGNALED(status)
 - True if a child terminated abnormally by receipt of a signal
- WTERMSIG(status) (only if WIFSIGNALED(status) is true)
 - Fetch the signal number
- WIFSTOPPED(status)
 - True if a child is currently stopped
- WSTOPSIG(status) (only if WIFSTOPPED(status) is true)
 - Fetch the signal number

Example

```
#include <sys/types.h>
#include <sys/wait.h>
void pr_exit(int status)
    if (WIFEXITED(status))
         printf("normal termination, exit status = \% d n",
              WEXITSTATUS(status));
    else if (WIFSIGNALED(status))
         printf("abnormal termination, signal number = %d\n",
              WTERMSIG(status));
    else if (WIFSTOPPED(status))
         printf("child stopped, signal number = %d\n",
              WSTOPSIG(status));
```

Example(cont.)

```
#include <sys/types.h>
#include <sys/wait.h>
int main(void)
     pid_t pid;
     int
           status;
     if ((pid = fork()) < 0)
          printf("fork error");
     else if (pid == 0)
                                               /* child */
          exit(7);
     if (wait(&status) != pid)
                                               /* wait for child */
          printf("wait error");
     pr_exit(status);
                                               /* and print its status */
```

Example(cont.)

```
if (\text{pid} = \text{fork}()) < 0)
      printf("fork error");
else if (pid == 0)
                                /* child */
      abort();
                                /* generates SIGABRT */
if (wait(&status) != pid)
                                /* wait for child */
      printf("wait error");
pr_exit(status);
                                /* and print its status */
if (\text{pid} = \text{fork}()) < 0)
      printf("fork error");
else if (pid == 0)
                                /* child */
      status = 0;
                                /* divide by 0 generates SIGFPE */
if (wait(&status) != pid)
                                /* wait for child */
      printf("wait error");
pr_exit(status);
                                /* and print its status */
```

Execution

```
$ ./a.out
normal termination, exit status = 7
abnormal termination, signal number = 6
abnormal termination, signal number = 8
$
```

Each signal number is defined in <signal.h> SIGABRT: 6

SIGFPE: 8

#include <sys/wait.h>

pid_t waitpid(pid_t pid, int *statloc, int options);

Both return: process ID if OK, 0 (see later), or -1 on error

Why waitpid() is required?

- If we have more than one child, wait() returns on termination of any of the children.
- What if we want to wait for a specific process to terminate?

waitpid()

- Waits for a <u>specific</u> process with *pid* argument.
- Provides some controls with options argument.

- pid argument
 - pid == -1
 - Waits for any child process.
 - In this case, waitpid() is equivalent to wait().
 - pid > 0
 - Waits for the child whose process ID equals pid.
 - pid == 0
 - Waits for any child whose process group ID equals that of the calling process.
 - pid < -1</p>
 - Waits for any child whose process group ID equals the absolute value of pid.

options argument

- WNOHANG
 - Not block(return immediately) if a child specified by pid is not terminated.
- WUNTRACED
 - Not block if a child specified by pid has stopped.

Example

```
#include "apue.h"
#include <sys/wait.h>
int main(void)
  pid_t pid;
  if ((pid = fork()) < 0) {
     err_sys("fork error");
  } else if (pid == 0) {
                         /* first child */
     if ((pid = fork()) < 0)
       err_sys("fork error");
     else if (pid > 0)
       exit(0);
                                  /* parent from second fork == first child */
     /* We're the second child; our parent becomes init. */
     sleep(2);
     printf("second child, parent pid = \%d\n", getppid());
     exit(0);
```

Example

```
if (waitpid(pid, NULL, 0) != pid) /* wait for first child */
    err_sys("waitpid error");

/*
    * We're the parent (the original process); we continue executing,
    * knowing that we're not the parent of the second child.
    */
    exit(0);
}
```

Execution

```
$ ./a.out
$ second child, parent pid = 1
```

Execute a program.

- When a process calls one of the exec functions, the process is completely replaced by the new program.
- New program starts executing at its main() function.
- exec() merely replaces the current process(text, data) with a brand new program from disk.

- pathname(filename) argument
 - pathname(filename) of a file to be executed.
 - execl, execv, execle, execve take a pathname argument.
 - execlp, execvp take a filename argument.
 - If filename contains a '/', it is taken as a pathname.
 - Otherwise, the executable file is searched for in the directories specified by PATH.
 - E.g. PATH=/bin:/usr/bin:/usr/local/bin/:.

Argument list

- arg0, arg1, ..., argn in the execl, execlp, and execle.
- The list of arguments is terminated by a NULL pointer.
 - execl("/bin/ls", "ls", "-al", 0);

Argument vector

- *argv[] in execv, execve and execvp.
- The array of pointers is terminated by a NULL pointer.
 - char *argv[3] = {"ls", "-al", 0};
 - execv("/bin/ls", argv);

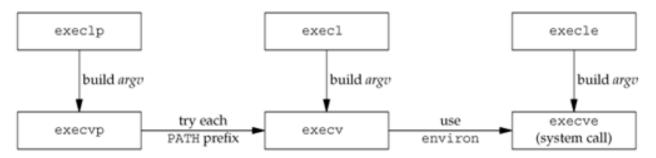
Environment variable

- *envp[] in execve and execle
- A pointer to an array of pointers to the environment strings.
- The other functions use the environ variable.
 - char *env[2] = {"TERM=vt100", 0};
 - execle("/bin/ls", "ls", 0, env);

Differences among the six exec functions

function	pathname	filename	arg list	argv[]	environ	envp[]
execl	0		0		0	
execlp		0	0		0	
execle	0		0			0
execv	0			0	0	
execvp		0		0	0	
execve	0			0		0
(letter in name)		p	l	٧		e

- execve is a system call.
 - execl, execv, execle, execlp, execvp are library functions.
 - Relationship of the six exec functions.



Example

```
#include "apue.h"
#include <sys/wait.h>
      *env_init[] = { "USER=unknown", "PATH=/tmp", NULL };
char
int
main(void)
  pid_t pid;
  if ((pid = fork()) < 0) {
     err_sys("fork error");
  } else if (pid == 0) { /* specify pathname, specify environment */
     if (execle("/home/sar/bin/echoall", "echoall", "myarg1",
          "MY ARG2", (char *)0, env_init) < 0)
       err_sys("execle error");
```

Example(cont.)

```
if (waitpid(pid, NULL, 0) < 0)
    err_sys("wait error");

if ((pid = fork()) < 0) {
    err_sys("fork error");
} else if (pid == 0) { /* specify filename, inherit environment */
    if (execlp("echoall", "echoall", "only 1 arg", (char *)0) < 0)
        err_sys("execlp error");
}

exit(0);
}</pre>
```

echoall

```
#include "apue.h"
int
main(int argc, char *argv[])
  int
           **ptr;
  char
  extern char **environ;
  for (i = 0; i < argc; i++)
                                             /* echo all command-line args */
     printf("argv[%d]: %s\n", i, argv[i]);
  for (ptr = environ; *ptr != 0; ptr++) /* and all env strings */
     printf("%s\n", *ptr);
  exit(0);
```

Execution

```
$ ./a.out
argv[0]: echoall
argv[1]: myarg1
argv[2]: MY ARG2
USER=unknown
PATH=/tmp
$ argv[0]: echoall
argv[1]: only 1 arg
USER=sar
LOGNAME=sar
SHELL=/bin/bash
...
HOME=/home/sar
```

```
#include <stdlib.h>
int system(const char *cmdstring);
Returns: (see below)
```

- Consists of fork(), exec(), and waitpid().
 - If fork() fails or waitpid() returns an error other than EINTR, returns -1.
 - If exec() fails, return value is 127.
 - If successful, return the termination status of the shell.

An implementation of system()

```
#include
          <sys/wait.h>
#include
           <errno.h>
#include <unistd.h>
int
system(const char *cmdstring)
  pid_t pid;
  int
        status;
  if (cmdstring == NULL)
     return(1);
                                     -c option: commands are read from 'cmdstring'
  if ((pid = fork()) < 0) {
                                             /* probably out of processes */
     status = -1;
  } else if (pid == 0) {
                                             /* child */
     execl("/bin/sh", "sh", "-c", cmdstring, (char *)0);
                                             /* execl error */
     _exit(127);
```

An implementation of system() (cont.)

Example

```
#include "apue.h"
#include <sys/wait.h>
int main(void)
  int
         status;
                                          pr_exit() is defined in "wait() and waitpid()" section.
  if ((status = system("date")) < 0)
     err_sys("system() error");
  pr exit(status);
  if ((status = system("nosuchcommand")) < 0)
    err_sys("system() error");
  pr_exit(status);
 if ((status = system("who; exit 44")) < 0)
    err_sys("system() error");
  pr_exit(status);
 exit(0);
                                                                                          55
```

Execution

```
$./a.out
Sun Mar 21 18:41:32 EST 2004
                                                     for date
normal termination, exit status = 0
sh: nosuchcommand: command not found
normal termination, exit status = 127
                                                     for nosuchcommand
            Mar 18 19:45
      :0
sar
      pts/0 Mar 18 19:45 (:0)
sar
            Mar 18 19:45 (:0)
      pts/1
sar
      pts/2 Mar 18 19:45 (:0)
sar
      pts/3 Mar 18 19:45 (:0)
sar
normal termination, exit status = 44
                                                     for exit
$
```