

stake.link Security Review



Lead Auditors



PeterSR



0x539.eth

Table of Contents

- Table of Contents
- Protocol Summary
- Disclaimer
- · Risk Classification
- Audit Details
 - Scope
 - Roles
- Executive Summary
 - Issues found
- Findings
 - Medium

Protocol Summary

stake.link is a first-of-its-kind liquid delegated staking platform delivering DeFi composability for Chainlink Staking. Built by premier Chainlink ecosystem developer LinkPool, powered by Chainlink node operators, and governed by the stake.link DAO, stake.link's extensible architecture is purpose-built to support Chainlink Staking and to extend participation in the Chainlink Network.

Disclaimer

The ChainDefenders team makes all effort to find as many vulnerabilities in the code in the given time period, but holds no responsibilities for the findings provided in this document. A security audit by the team is not an endorsement of the underlying business or product. The audit was time-boxed and the review of the code was solely on the security aspects of the Solidity implementation of the contracts.

Risk Classification

Likelihood/Impact	High	Medium	Low
High	Н	H/M	M
Medium	H/M	M	M/L
Low	M	M/L	L

Audit Details

Scope

Id	Files in scope
1	LSTRewardsSplitter.sol
2	LST Rewards Splitter Controller. sol
3	PriorityPool.sol
4	WithdrawalPool.sol
5	StakingPool.sol
6	Vault.sol
7	VaultControllerStrategy.sol
8	CommunityVault.sol
9	CommunityVCS.sol
10	FundFlowController.sol
11	OperatorStakingPool.sol
12	OperatorVault.sol
13	OperatorVCS.sol

Roles

Id	Roles	
1	Staker	
2	Queued Staker	
3	Queued Withdrawer	
4	Operator	
5	Operator Rewards Receiver	
6	Owner	

Executive Summary

Issues found

Severity	Count	Description
High	0	Critical vulnerabilities
Medium	3	Significant risks
Low	0	Minor issues with low impact
Informational	0	Best practices or suggestions
	0	Optimization opportunities

Findings

Medium

Mid 01 Malicious User Can DoS Withdraw

Summary

The WithdrawalPool contains an optimization for managing withdrawal batches. The function updateWithdrawalBatchIdCutoff is responsible for updating the withdrawalBatchIdCutoff value. However, if a user has an incomplete withdrawal for a specific batch, this update cannot proceed until the user completes their withdrawal. A malicious user could exploit this situation to indefinitely prevent the update, leading to a denial-of-service (DoS) condition

Vulnerability Details

Look at the following scenario:

- 1. A user submits a withdrawal request, which is assigned an ID of 130 and is associated with batch 35.
- 2. The withdrawalBatchIdCutoff remains at 35 until the user's withdrawal request is finalized.
- 3. If the withdrawal never completes, the getBatchIds function will iteratively process all batches for every withdrawal ID, potentially causing significant delays or halting other processes.

```
1 function getBatchIds(uint256[] memory _withdrawalIds) public view
     returns (uint256[] memory) {
          uint256[] memory batchIds = new uint256[](_withdrawalIds.
     length);
          for (uint256 i = 0; i < withdrawalIds.length; ++i) {</pre>
              uint256 batchId;
              uint256 withdrawalId = _withdrawalIds[i];
              for (uint256 j = withdrawalBatchIdCutoff; j <</pre>
     withdrawalBatches.length; ++j) {
                  uint256 indexOfLastWithdrawal = withdrawalBatches[j].
     indexOfLastWithdrawal;
                  if (withdrawalId ≤ indexOfLastWithdrawal) {
                      batchId = j;
                      break;
              }
              batchIds[i] = batchId;
          }
          return batchIds;
      }
1 function updateWithdrawalBatchIdCutoff() external { // audit-mid
     possible dos
          uint256 numWithdrawals = queuedWithdrawals.length;
          uint256 newWithdrawalIdCutoff = withdrawalIdCutoff;
          for (uint256 i = newWithdrawalIdCutoff; i < numWithdrawals; ++</pre>
     i) {
              newWithdrawalIdCutoff = i;
              Withdrawal memory withdrawal = queuedWithdrawals[i];
              if (withdrawal.sharesRemaining \neq 0 || withdrawal.
     partiallyWithdrawableAmount \neq 0) {
                  break;
          } // this depends does the withdraw is already excuted
          uint256 numBatches = withdrawalBatches.length;
          uint256 newWithdrawalBatchIdCutoff = withdrawalBatchIdCutoff;
```

```
remaining

for (uint256 i = newWithdrawalBatchIdCutoff; i < numBatches;

++i) {

if (withdrawalBatches[i].indexOfLastWithdrawal ≥

newWithdrawalIdCutoff) {

break;

}

newWithdrawalBatchIdCutoff = i;

}

withdrawalIdCutoff = uint128(newWithdrawalIdCutoff);

withdrawalBatchIdCutoff = uint128(newWithdrawalBatchIdCutoff);

withdrawalBatchIdCutoff = uint128(newWithdrawalBatchIdCutoff);

}
```

Impact

This vulnerability could allow a malicious user to block the withdrawal process for other users by preventing the system from advancing the withdrawalBatchId-Cutoff. This can lead to a DoS attack where no new withdrawals can be processed until the malicious withdrawal is completed, affecting the overall availability of the system.

Tools Used

Manual review

Recommendations

To mitigate this issue, consider implementing a timeout or fallback mechanism where an admin can finalize or cancel old withdrawal requests after a certain period of inactivity.

Mid 02 VaultMapping Is Not Correctly Updated

Summary

When a vault is added to the OperatorVCS the function addVault is used:

```
1 function addVault(
      address _operator,
      address _rewardsReceiver,
      address _pfAlertsController
  ) external onlyOwner {
      bytes memory data = abi.encodeWithSignature(
          "initialize(address,address,address,address,address,
     address)",
          address(token),
          address(this),
          address(stakeController),
          stakeController.getRewardVault(),
          _pfAlertsController,
          _operator,
          rewardsReceiver
      );
      deployVault(data);
      vaultMapping[address(vaults[vaults.length - 1])] = true;
      emit VaultAdded(_operator);
19 }
```

We can see that the vaultMapping for the concrete vault is set to true. When removing a vault the function queueVaultRemoval is used and then the function

removeVault:

```
1 function queueVaultRemoval(uint256 _index) external {
      address vault = address(vaults[_index]);
      if (!IVault(vault).isRemoved()) revert OperatorNotRemoved();
      for (uint256 i = 0; i < vaultsToRemove.length; ++i) {</pre>
          if (vaultsToRemove[i] = vault) revert
     VaultRemovalAlreadyQueued();
      vaultsToRemove.push(address(vaults[_index]));
      // update group accounting if vault is part of a group
      if (_index < globalVaultState.depositIndex) {</pre>
          uint256 group = _index % globalVaultState.numVaultGroups;
          uint256[] memory groups = new uint256[](1);
          groups[0] = group;
          fundFlowController.updateOperatorVaultGroupAccounting(groups);
          // if possiible, remove vault right away
          if (vaults[_index].claimPeriodActive()) {
              removeVault(vaultsToRemove.length - 1);
```

```
}
  function removeVault(uint256 _queueIndex) public {
      address vault = vaultsToRemove[_queueIndex];
      vaultsToRemove[_queueIndex] = vaultsToRemove[vaultsToRemove.length
      - 1];
      vaultsToRemove.pop();
      _updateStrategyRewards();
      (uint256 principalWithdrawn, uint256 rewardsWithdrawn) =
      IOperatorVault(vault).exitVault();
      totalDeposits -= principalWithdrawn + rewardsWithdrawn;
      totalPrincipalDeposits -= principalWithdrawn;
      uint256 numVaults = vaults.length;
      uint256 index;
      for (uint256 i = 0; i < numVaults; ++i) {
          if (address(vaults[i]) = vault) {
              index = i;
              break;
          }
44
      for (uint256 i = index; i < numVaults - 1; ++i) {
          vaults[i] = vaults[i + 1];
      vaults.pop();
      token.safeTransfer(address(stakingPool), token.balanceOf(address(
      this)));
51 }
```

However, we can see that in both function this vaultMapping that was set to true is never set to false. The mapping is used in withdrawOperatorRewards:

```
function withdrawOperatorRewards(
   address _receiver,
   uint256 _amount

/ Dexternal returns (uint256) {
   if (!vaultMapping[msg.sender]) revert SenderNotAuthorized();

/ IERC20Upgradeable lsdToken = IERC20Upgradeable(address(stakingPool));

uint256 withdrawableRewards = lsdToken.balanceOf(address(this));

uint256 amountToWithdraw = _amount > withdrawableRewards?
```

```
withdrawableRewards : _amount;

unclaimedOperatorRewards -= amountToWithdraw;
lsdToken.safeTransfer(_receiver, amountToWithdraw);

return amountToWithdraw;
}
```

This means that even if a vault is "removed" it will still be able to withdraw operator rewards.

Vulnerability Details

Let's have the following scenario:

- 1. Vault A is added via addVault.
- 2. Vault A is removed via queueVaultRemoval -> removeVault.
- 3. However, Vault A is able to call withdrawOperatorRewards setting _receiver and _amount and withdrawing the rewards even though it should not be able to.

Impact

Such an issue has a critical impact to the security of the OperatorVCS contract as no one is able to stop a vault from operating even after removing it.

Tools Used

Manual Review

Recommendations

Update the vaultMapping to reflect that a given vault was removed:

```
function removeVault(uint256 _queueIndex) public {
   address vault = vaultsToRemove[_queueIndex];
   vaultMapping[vault] = false;
   vaultsToRemove[_queueIndex] = vaultsToRemove[vaultsToRemove.length - 1];
   vaultsToRemove.pop();
```

```
_updateStrategyRewards();
      (uint256 principalWithdrawn, uint256 rewardsWithdrawn) =
      IOperatorVault(vault).exitVault();
      totalDeposits -= principalWithdrawn + rewardsWithdrawn;
      totalPrincipalDeposits -= principalWithdrawn;
      uint256 numVaults = vaults.length;
      uint256 index;
      for (uint256 i = 0; i < numVaults; ++i) {</pre>
          if (address(vaults[i]) = vault) {
              index = i;
              break;
          }
      for (uint256 i = index; i < numVaults - 1; ++i) {
          vaults[i] = vaults[i + 1];
      vaults.pop();
      token.safeTransfer(address(stakingPool), token.balanceOf(address(
      this)));
27
```

Mid 03 Remove Splitter Will Not Work In Certain Situations

Summary

The LSTRewardsSplitterController:: removeSplitter function is responsible for removing a splitter and withdrawing all associated funds. The funds are transferred to an account specified as a parameter. Before this happens, the function calls splitRewards on the splitter to distribute any outstanding rewards to fee receivers. However, there is a flaw in this process: after splitting rewards, the function attempts to withdraw the original balance (before the split) instead of the updated balance, which will likely cause the transaction to revert if the balance has changed.

Vulnerability Details

Consider the following scenario to illustrate the issue:

- 1. A splitter has a balance of 100 tokens.
- 2. After splitting the rewards, part of the balance (e.g., 10 tokens) is distributed to fee receivers, leaving a remaining balance of 90 tokens.
- 3. The removeSplitter function, however, tries to withdraw the original balance of 100 tokens, which exceeds the actual available balance of 90 tokens. This mismatch causes a failure or revert when the withdrawal is attempted.

This issue arises due to the fact that the function does not update the balance after calling splitRewards. It attempts to withdraw the pre-split balance, which is incorrect.

The vulnerability can be observed in the code snippet below:

```
function removeSplitter(address _account) external onlyOwner {
    ILSTRewardsSplitter splitter = splitters[_account];
    if (address(splitter) = address(0)) revert SplitterNotFound();

uint256 balance = IERC20(lst).balanceOf(address(splitter)); //
Original balance is captured
uint256 principalDeposits = splitter.principalDeposits();
if (balance ≠ 0) {
    if (balance ≠ principalDeposits) splitter.splitRewards(); //
    Split rewards, changing the balance
    splitter.withdraw(balance, _account); // Attempts to withdraw the old balance, likely causing a revert
}
```

In the above function, the balance is initially fetched before the rewards are split. This results in an outdated balance being passed to the withdraw function, leading to a revert if the balance has changed due to the rewards distribution.

Impact

This vulnerability can cause the removeSplitter function to revert when attempting to withdraw funds, preventing the successful removal of the splitter and the withdrawal of remaining funds. This could halt operations or cause disruptions, especially if multiple splitters are being managed within the system.

Tools Used

Manual review

Recommendations

To fix this issue, the balance should be updated after the splitRewards call to ensure that the correct amount is withdrawn. Below is the recommended change to the removeSplitter function:

By re-fetching the balance after the splitRewards function call, the correct (post-split) balance will be withdrawn, preventing any reverts and ensuring smooth operation.