MLIR Side Effect Modelling

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```
func @foo(%a: i32, %b: i32) {
    %c = add %a, %b
    %d = add %c, %b
    return %d
}
```

```
func @foo(%a: i32, %b: i32) {
    %d = add %c, %b
    %c = add %a, %b
    return %d
}
```

Legal

no side effects ~ "Can be rearranged freely"

```
func @bar(%p: i32, %q: i32) {
    %r = add %a, %b
    ... (LOTS OF CODE)
    %s = add %a, %b
    return %s
}
```

```
func @bar(%p: i32, %q: i32) {
    %r = add %p, %q
    ... (LOTS OF CODE)

    %s = %r
    return %s
}
```

Legal

no side effects ~ "output depends only on SSA inputs"

```
func @baz(%a: memref<?xi32>, %i: index) {
    %v1 = memref.load %a[%i] : i32
    ...
    ...
    ...
    return %v2
}
```

```
func @baz(%a: memref<?xi32>, %i: index) {
    %v1 = memref.load %a[%i] : i32
    %c42 = constant 42 : i32
    memref.store %c42, %a[%i] : i32
    ...
    ...
    return %v2
}
```

```
func @baz(%a: memref<?xi32>, %i: index) {
    %v1 = memref.load %a[%i] : i32
    %c42 = constant 42 : i32
    memref.store %c42, %a[%i] : i32
    %c100 = constant 100 : i32
    memref.store %c100, %a[%i] : i32
    ...
    return %v2
}
```

```
func @baz(%a: memref<?xi32>, %i: index) {
                                             func @baz(%a: memref<?xi32>, %ix: index) {
 %v1 = memref.load %a[%i] : i32
                                               %v1 = memref.load %a [%ix] : i32
 %c42 = constant 42 : i32
                                               %c100 = constant 100 : i32
 memref.store %c42, %a[%i] : i32
                                               memref.store %c100, %a[%ix] : i32
 %c100 = constant 100 : i32
                                               %c42 = constant 42 : i32
                                               memref.store %c42, %a[%ix] : i32
 memref.store %c100, %a[%i] : i32
 return %v2
                                               return %v2
```

Illegal

```
func @baz(%a: memref<?xi32>, %i: index) {
    %v1 = memref.load %a[%i] : i32
    %c42 = constant 42 : i32
    memref.store %c42, %a[%i] : i32
    %c100 = constant 100 : i32
    memref.store %c100, %a[%i] : i32
    %v2 = memref.load %a[%i] : i32
    return %v2
}
```

```
func @baz(%a: memref<?xi32>, %i: index) {
                                           func @baz(%a: memref<?xi32>, %i: index) {
 %v1 = memref.load %a[%i] : i32
                                             %v1 = memref.load %a[%i] : i32
 %c42 = constant 42 : i32
                                             %c42 = constant 42 : i32
 memref.store %c42, %a[%i] : i32
                                             memref.store %c42, %a[%i] : i32
 %c100 = constant 100 : i32
                                             %c100 = constant 100 : i32
 memref.store %c100, %a[%i] : i32
                                             memref.store %c100, %a[%i] : i32
 %v2 = memref.load %a[%i] : i32
                                             %v2 = %v1
 return %v2
                                             return %v2
```

Illegal

```
func @baz(%mem: memref<?xi32>, %ix: index) {
    %v1 = memref.load %a [%ix] : i32
    %c42 = constant 42 : i32
    memref.store %c42, %a[%ix] : i32
    %c100 = constant 100 : i32
    memref.store %c100, %a[%ix] : i32
    %v2 = memref.load %mem [%ix] : i32
    return %v2
}
```

side effects ~= "output depends on implicit state (e.g. memory)"

Side effect modeling in MLIR

SideEffect

EffectKind (TypeID) Resource (singleton) Affected (Value)

Side effect modeling in MLIR

SideEffect EffectKind Resource Affected (Value) (TypeID) (singleton) SideEffect EffectOpInterface It's **super** abstract! SideEffect %2:2 = "i have.side effects"(%0, %1) : (pointer, index) -> pointer SideEffect

Example: MemoryEffect

EffectKind

- Alloc
- Free
- Read
- Write

MemoryEffect

Resource

- GPU shared
- Runtime handles
- JIT contexts
- . . .

Affected

- memref
- alloca
- Global symbol
- . . .

Example: MemoryEffect

```
func @baz(%a: memref<?xi32>, %i: index) {
                                        store(%a)
 %v1 = memref.load %a[%i] : i32
 %c42 = constant 42 : i32
 memref.store %c42, %a[%i] : i32
                                         load(%a)
 %c100 = constant 100 : i32
 memref.store %c100, %a[%i] : i32
                                        store(%a)
 %v2 = memref.load %a[%i] : i32
                                         load(%a)
 return %v2
```

CSE: MemoryEffects in action!

```
func @cse reads(%a: memref<?xi32>, %i: index) {
                                        load(%a)
 %v1 = memref.load %a[%i] : i32
 %v2 = arith.addi %v1, %v1 : i32
 %v3 = memref.load %a[%i] : i32
                                        load(%a)
 %c100 = constant 100 : i32
 memref.store %c100, %a[%i] : i32
                                       store(%a)
 %v4 = arith.addi %v2, %v3
 return %v4 : i32
```

No other effects

CSE: MemoryEffects in action!

```
func @cse reads(%a: memref<?xi32>, %i: index) {
                                        load(%a)
 %v1 = memref.load %a[%i] : i32
 %v2 = arith.addi %v1, %v1 : i32
 %v3 = memref.load %a[%i] : i32
 %c100 = constant 100 : i32
 memref.store %c100, %a[%i] : i32
                                      store(%a)
 %v4 = arith.addi %v2, %v3
 return %v4 : i32
```

Is MemoryEffect::Alloc elidable?

applyPatternsAndFoldGreedily happily deletes this!

What about Leak Sanitizer?

Result: memory leak is "optimized away" by DCE

Other effect stuff

effectOnFullRegion: "if this side effect acts on every single value of resource" (2)

- 2 google search results (https://shorturl.at/kvwzB)
- O GitHub search results outside of llvm-project and forks
 (https://github.com/search?q=getEffectOnFullRegion+&type=code)

stage: Order of effects on the same operation

 An operation can Read and then Free a resource, if the stage of the Read is earlier than the Free

RecursiveMemoryEffects: Operation effects := union of all effects in region

ConditionallySpeculatable

```
TL;DR: can the op be hoisted?
                                       func @divs_safe(%lhs: index, %rhs: index) -> index {
Example: div
                                         %idx0 = index.constant 0
                                         %0 = index.cmp eq(%rhs, %idx0)
                                         %1 = scf.if %0 -> index {
                                           scf.yield %idx0 : index
                                         } else {
                   Potential UB
                                           %2 = index.divs %lhs, %rhs
                                           scf.yield %2 : index
                                         return %1 : index
```

What's missing?

Memory effects on functions and function arguments

- Enables propagating memory effects to callers
- E.g. readnone, readonly, read+write (i.e. local to the caller)

MemoryEffectOpInterface "black hole"

- Super-privileged in MLIR infra
- Users are forced into an often substandard model to reuse upstream infra

Example: DCE

```
// Trait for enforcing that a side-effecting op is executed, even if it would be
// considered dead by MLIR (see b/195782952).

// The trait is implemented as a write effect for a fake resource which is
// ignored by side effect analysis, so it does not affect execution order

// constraints and control dependencies at all (for example, multiple ops with
// this trait do not have to execute in order).

def TF_MustExecute : MemoryEffects<[MemWrite<TF_MustExecuteResource>]>;
```

"Dear MLIR, please don't delete this op"

Example: DCE

Whenever an op in the loop is modified, canonicalizer checks if the loop can be DCE'd

⇒ RecursiveMemoryEffects rewalks the whole body



```
func @big_loop(%lb: index, %ub: index) {
    %step = index.constant 1
    scf.for %i = %lb to %ub step %step {
        // Huge inlined loop body with more loops inside
    }
    return
}
```

Some Researchy Ideas for Side Effect Modelling'

Key Idea: Move side effects into the def-use chain, have interfaces for the Ops

Eg. MemoryEffect \rightarrow MemorySSA, Async \rightarrow Async.token, poison

Pro: Reuses MLIR's strengths! MemorySSA analysis becomes another dialect.

Con: UB does not fit into this model:(

Synthesis: def-use chain for those that can, Interfaces that which cannot.

[1] Gordon D Plotkin and Matija Pretnar. Handling Algebraic Effects. Logical Methods in Computer Science, Volume 9, Issue 4. December 2013.

[2]Novillo, Diego. "Memory SSA-a unified approach for sparsely representing memory operations." *Proc of the GCC Developers' Summit*. 2007.

Even More Bleeding Edge Ideas! (Koka)



fun sqr fun divide

fun turing

fun print

fun rand

The precise effective well-studied categories and compilers

Japanese word



Conclusion

EffectKind (TypeID)

SideEffect

Resource (singleton)

Affected (Value)

```
EffectOpInterface
%2:2 = "i_have.side_effects"(%0, %1)
: (pointer, index) -> pointer
SideEffect
SideEffect
```