CS2308 - SYSTEM SOFTWARE LABORATORY

LABORATORY MANUAL

FOR FIFTH SEMESTER B.TECH-IT

ACADEMIC YEAR 2013-2014 (ODD)

(FOR PRIVATE CIRCULATION ONLY)

ANNA UNIVERSITY, CHENNAI

DEPARTMENT OF INFORMATION TECHNOLOGY

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GENERAL INSTRUCTIONS FOR LABORATORY CLASSES

- Enter Lab with **CLOSED FOOTWEAR**
- Boys should "TUCK IN" the shirts
- Students should wear uniform only
- LONG HAIR should be protected, let it not be loose especially near ROTATING MACHINERY.
- Any other machines/ equipments **should not be operated** other than the prescribed one for that day.
- POWER SUPPLY to your test table should be obtained only through the LAB
 TECHNICIAN
- Do not **LEAN** and do not be **CLOSE** to the rotating components.
- TOOLS, APPARATUS & GUAGE Sets are to be returned before leaving the Lab.
- **HEADINGS & DETAILS** should be neatly written
 - 1. Aim of the experiment
 - 2. Apparatus / Tools/ Instruments required
 - 3. Procedure / Theory / Algorithm / Program
 - 4. Model Calculations
 - 5. Neat Diagram/ Flow charts
 - 6. Specifications/ Designs details
 - 7. Tabulation
 - 8. Graph
 - 9. Result/ Discussions
- Before doing the experiment, the student should get the circuit/ Program approval by the
 FACULTY-IN-CHARGE
- **Experiment date** should be written int the appropriate place
- After completing the experiments, the answer to the VIVA-VOCE Questions should be neatly written in the workbook
- Be PATIENT, STEADY, SYSTEMATIC, & REGULAR

HARDWARE REQUIREMENTS:

Processors - 2.0 GHz or Higher
RAM - 256 MB or Higher
Hard Disk - 20 GB or Higher

Operating System - Windows 2000/XP/NT

SOFTWARE REQUIREMENTS:

TURBO C (Freeware)

UNIVERSITY PRACTICAL EXAMINATION

ALLOTMENT OF MARKS

Internal assessment - 20 marks Practical assessment - 80 marks

Total - 100 marks

INTERNAL ASSESSMENT (20 marks)

Staff should maintain the assessment Register and the Head of the Department should monitor it.

SPLIT UP OF INTERNAL MARKS

Record Note - 10 marks

Model Exam - 5 marks Attendance - 5 marks

Total - 20 marks

UNIVERSITY EXAMINATION

The exam will be conducted for 100 marks. Then the marks will be calculated to 80 marks.

SPLIT UP OF PRACTICAL EXAMINATION MARKS

Aim and Algorithm - 20 marks
Program - 40 marks
Output - 20 marks
Result - 10 marks
Viva-voce - 10 marks

Total - 100 marks

CS2308 - SYSTEM SOFTWARE LAB

LIST OF EXPERIMENTS

- 1. Implement a symbol table with functions to create, insert, modify, search, and display.
- 2. Implement pass one of a two pass assembler.
- 3. Implement pass two of a two pass assembler.
- 4. Implement a single pass assembler.
- 5. Implement a two pass macro processor
- 6. Implement a single pass macro processor.
- 7. Implement an absolute loader.
- 8. Implement a relocating loader.
- 9. Implement pass one of a direct-linking loader.
- 10. Implement pass two of a direct-linking loader.
- 11. Implement a simple text editor with features like insertion / deletion of a character, word, and sentence.
- 12. Implement a symbol table with suitable hashing

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Exercise Number: 1

Title of the Exercise : IMPLEMENTATION OF A SYMBOL TABLE

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

• To write a program to implement symbol table.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no	Details of the step
1	Initialize all the variables.
2	Design a menu through which we can create a symbol Table and perform operations as insert, modify, search and display.
3	Create a symbol table with fields as variable and value using create option. Entries may be added to the table while it is created itself.
4	Append new contents to the symbol table with the constraint that there is no duplication of entries, using insert option.
5	Modify existing content of the table using modify option.
6	Use display option to display the contents of the table.
7	End of the program.

```
#include<stdio.h>
                                                  #include<stdlib.h>
                            #include<conio.h>
                                                                        #include<string.h>
Struct table
char var[10]; int value;
struct table tb1[20];
              void create();
                                    void modify();
int i,j,n;
int search(char variable[],int n);
                                    void insert();
                                                         void display();
void main()
       int ch,result=0;
char v[10];
clrscr();
do
printf("\n\n1.CREATE\n2.INSERT\n3.MODIFY\n4.SEARCH\n5.DISPLAY\n6.EXIT:\t");
scanf("%d",&ch);
```

```
switch(ch)
{
case 1:create();
                       break;
case 2:insert();
                       break;
case 3:modify();
                       break;
case 4:printf("\nEnter the variable to be searched:");
scanf("%s",&v);
result=search(v,n);
if(result==0)
printf("\nThe variable is not present\n");
printf("\nThe location of the variable is %d \n The value of %s is
%d.",result,tb1[result].var,tb1[result].value);
break;
case 5:display();
                       break;
case 6:exit(1);
}while(ch!=6);
getch();
void create()
       printf("\nEnter the no. of entries:");
scanf("%d",&n);
printf("\nEnter the variable and the values:-\n");
for(i=1;i \le n;i++)
{ scanf("%s%d",tb1[i].var,&tb1[i].value);
check:
if(tb1[i].var[0] >= '0' && tb1[i].var[0] <= '9')
{ printf("\nVariable should start with alphabet\nEnter correct name\n");
scanf("%s%d",tb1[i].var,&tb1[i].value);
goto check;
check1:
for(j=1;j< i;j++)
\{ if(strcmp(tb1[i].var,tb1[i].var)==0 \}
{ printf("\nThe variable already present. Enter another:");
scanf("%s%d",&tb1[i].var,&tb1[i].value);
goto check1;
printf("\nThe table after creation is:\n");
display();
void insert()
if(i \ge 20)
printf("\nCannot insert.table is full\n");
else
n++;
```

```
printf("\nEnter the variable and the value:");
scanf("%s%d",&tb1[n].var,&tb1[n].value);
if(tb1[i].var[0] > = '0' & & tb1[i].var[0] < = '9')
printf("\nVariable should start with alphabet\nEnter correct name\n");
scanf("%s%d",tb1[i].var,&tb1[i].value);
goto check;
check1:
for(j=1;j< n;j++)
if(strcmp(tb1[i].var,tb1[i].var)==0)
printf("\nThe variable already present. Enter another:");
scanf("%s%d",&tb1[i].var,&tb1[i].value);
goto check1;
printf("\nThe table after insertion is:");
display();
void modify()
char variable[10];
int result=0;
printf("\nEnter the variable to be modified:");
scanf("%s",&variable);
result=search(variable,n);
if(result==0)
printf("%s not present\n",variable);
else
printf("\nThe current value of the variable %s is %d.\nEnter the new variable and its
value",tb1[result].var,tb1[result].value);
scanf("%s%d",tb1[result].var,&tb1[result].value);
check:
if(tb1[i].var[0] > = '0' \&\& tb1[i].var[0] < = '9')
printf("\nVariable should start with alphabet\nEnter correct name\n");
scanf("%s%d",tb1[i].var,&tb1[i].value);
goto check;
printf("\nThe table after modification is:");
display();
}
int search(char variable[],int n)
```

```
IT/VSem
Dr.NNCE
int flag;
for(i=1;i \le n;i++)
if(strcmp(tb1[i].var,variable)==0)
flag=1;
break;
if(flag==1)
return i;
else
return 0;
void display()
printf("\nVariable\tvalue\n");
for(i=1;i \le n;i++)
printf("%s\t\t%d\n",tb1[i].var,tb1[i].value);
}
d) Output:
1.CREATE
2.INSERT
3.MODIFY
4.SEARCH
5.DISPLAY
6.EXIT:1
Enter the no. of entries:2
Enter the variable and the values:-
a 23
c 45
The table after creation is:
Variable
          value
          23
          45
c
1.CREATE 2.INSERT 3.MODIFY 4.SEARCH 5.DISPLAY 6.EXIT:2
Enter the variable and the value:b 34
The table after insertion is:
Variable
           value
          23
a
         45
c
```

SS Lab - LM

1.CREATE 2.INSERT 3.MODIFY 4.SEARCH 5.DISPLAY 6.EXIT: 3

Enter the variable to be modified:c

34

The current value of the variable c is 45.

Enter the new variable and its valuec 44

The table after modification is:

Variable	value
a	23
c	44
b	34

1.CREATE 2.INSERT 3.MODIFY 4.SEARCH 5.DISPLAY 6.EXIT: 4

Enter the variable to be searched:a

The location of the variable is 1

The value of a is 23.

1.CREATE 2.INSERT 3.MODIFY 4.SEARCH 5.DISPLAY 6.EXIT: 5

Variable	value
a	23
c	44
b	34

1.CREATE 2.INSERT 3.MODIFY 4.SEARCH 5.DISPLAY 6.EXIT: 6

e) Result:

Thus the symbol table is created and operations are verified successfully.

VIVA – QUESTION AND ANSWER:

1. What is system software?

System software consists of variety of programs that supports the operations of a computer. This makes it possible for the user to focus on an application or other problem to be solved,

without needing to know the details of how the machine works internally. Examples of system software are text-editors, compilers, loaders or linkers, debuggers,

2. Give some applications of operating system.

- > to make the computer easier to use
- > to manage the resources in computer
- > process management
- > data and memory management

assemblers, and operating systems.

- > to provide security to the user.
- Operating system acts as an interface between the user and the system Eg: windows, Linux, UNIX, dos

3. What is SIC machine?

SIC refers to Simplified Instruction Computer which is a hypothetical computer that has been designed to include the hardware features most often found on real machines, while avoiding unusual and irrelevant complexities. This allows to clearly separating the central concepts of system software from the implementation details associated with a particular machine

Exercise Number: 2

Title of the Exercise : IMPLEMENTATION OF PASS ONE OF A TWO PASS ASSEMBLER

Date of the Exercise:

OBJECTIVE (AIM) OF THE EXPERIMENT

To write a program to implement pass one of a two pass assembler.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no.	Details of the step			
1	begin			
2	read first input line;			
	if OPCODE = 'START' then			
	begin			
	 i. save #[OPERAND] as starting address 			
3	ii. initialized LOCCTR to starting address			
	iii. write line to intermediate file			
	iv. read next input line			
	end {if START}			
4	else			
7	initialized LOCCTR to 0			
	while OPCODE != 'END' do			
	begin			
	a. i. if this is not a comment line then			
	ii. begin			
	iii. if there is a symbol in the LABEL field then			
	begin			
	 search SYMTAB for LABEL 			
	2. if found then			
5	3. set error flag (duplicate symbol)			
	4. else			
	5. insert (LABEL, LOCCTR) into SYMTAB			
	end {if symbol}			
	iv. search OPTAB for OPCODE			
	v. if found then			
	add 3 {instruction length} to LOCCTR			
	vi. else if OPCODE = 'WORD' then			
	add 3 to LOCCTR			
	vii. else if OPCODE = 'RESW' then			

F			
	add 3 * #[OPERAND] to LOCCTR		
	viii. else if OPCODE = 'RESB' then		
	add #[OPERAND] to LOCCTR		
	ix. else if OPCODE = 'BYTE' then		
	begin		
	1. find length of constant in bytes		
	2. add length to LOCCTR		
	end {if BYTE}		
	x. else		
	set error flag (invalid operation code)		
	end {if not a comment}		
6	write line to intermediate file		
7	read next input line		
8	end {while not END}		
9	write last line to intermediate file		
10	save (LOCCTR - starting address) as program length		
11	end		

```
#include<stdio.h>
#include<conio.h>
#include<string.h>
void main()
char opcode[10],mnemonic[3],operand[10],label[10],code[10];
int locctr, start, length;
FILE *fp1,*fp2,*fp3,*fp4;
clrscr();
fp1=fopen("input.txt","r");
fp2=fopen("symtbl.txt","w");
fp3=fopen("out.txt","w");
fp4=fopen("optab.txt","r");
fscanf(fp1,"%s%s%s",label,opcode,operand);
if(strcmp(opcode, "START")==0)
start=atoi(operand);
locctr=start;
fprintf(fp3,"\t%s\t%s\t%s\n",label,opcode,operand);
fscanf(fp1,"%s%s%s",label,opcode,operand);
}
else
locctr=0;
while(strcmp(opcode,"END")!=0)
fprintf(fp3,"%d\t",locctr);
if(strcmp(label,"**")!=0)
```

```
fprintf(fp2,"%s\t%d\n",label,locctr);
fscanf(fp4,"%s%s",code,mnemonic);
while(strcmp(code, "END")!=0)
if(strcmp(opcode,code)==0)
locctr+=3;
break;
fscanf(fp4,"%s%s",code,mnemonic);
if(strcmp(opcode,"WORD")==0)
locctr+=3;
else if(strcmp(opcode, "RESW")==0)
locctr+=(3*(atoi(operand)));
else if(strcmp(opcode, "RESB")==0)
locctr+=(atoi(operand));
else if(strcmp(opcode, "BYTE")==0)
++locctr:
fprintf(fp3,"%s\t%s\t\n",label,opcode,operand);
fscanf(fp1,"%s%s%s",label,opcode,operand);
fprintf(fp3,"%d\t%s\t%s\t%s\n",locctr,label,opcode,operand);
length=locctr-start;
printf("The length of the program is %d",length);
fclose(fp1);
fclose(fp2);
fclose(fp3);
fclose(fp4);
getch();
}
```

d) Output:

INPUT FILES:

input.txt

```
**
     START
                2000
**
     LDA FIVE
**
     STA ALPHA
**
     LDCH CHARZ
**
     STCH C1
          RESW 1
ALPHA
FIVE WORD
CHARZ
          BYTE C'Z'
C1
     RESB 1
**
     END **
```

optab.txt

START
LDA 03
STA 0f
LDCH 53
STCH 57
END *

OUTPUT FILES:

The length of the program is 20

svmtab.txt

ALPHA 2012 FIVE 2015 CHARZ 2018 C1 2019

ouput.txt

START 2000 2000 ** LDA FIVE ** 2003 STA ALPHA 2006 ** LDCH CHARZ ** 2009 STCH C1 2012 ALPHA RESW 1 2015 FIVE WORD BYTE C'Z' 2018 CHARZ 2019 C1 RESB 1 2020 ** END **

e) Result:

Thus pass one of two passes assembler is implemented and the result is verified successfully.

VIVA – QUESTION AND ANSWER:

1. What are the instruction set for SIC/XE?

Instruction Set

- a. New registers: LDB, STB, etc.
- b. floating-point arithmetic: ADDF, SUBF, MULF, DIVF
- c. register move: RMO
- d. register-register arithmetic: ADDR, SUBR, MULR, DIVR
- e. supervisor call: SVC- generates an interrupt for OS (Chap 6)

Input/Output

f. SIO, TIO, HIO: start, test, halt the operation of I/O device

2. Define loader.

Loader is a set of program that loads the machine language translated by the translator into the main memory and makes it ready for execution.

Exercise Number: 3

Title of the Exercise : IMPLEMENTATION OF PASS TWO OF A TWO PASS

ASSEMBLER

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

To write program to implement pass two of a two pass assembler.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no.	Details of the step			
1	begin			
2	read first input file {from intermediate file}			
3	if OPCODE = 'START' then			
	begin			
4	a) write listing line			
	b) read next input line			
5	end {if START}			
6	write header record to object program			
7	initialized first Text record			
8	while OPCODE != 'END' do			
9	a) if this is not a comment line then i) begin ii) search OPTAB for OPCODE iii) if found then iv) begin (1) if there is a symbol in OPERAND field then (2) begin (a) search SYMTAB for OPERAND (b) if found then (c) store symbol value as operand address (d) else (e) begin (f) store 0 as operand address (g) set error flag (undefined symbol)			

	(h) end
	(3) end {if symbol}
	(4) else
	(5) store 0 as operand address
	(6) assemble the object code instruction
	i) end {if opcode found}
	ii) else if OPCODE = 'BYTE' or 'WORD' then
	iii) convert constant to object code
	iv) if object code not fit into the current Text record
	then
	v) begin
	(7) write Text record to object program
	(8) initialized new Text record
	vi) end
	vii) add object code to Text record
	viii) end {if not comment}
	b) write listing line
	read next input line
10	end {while not END}
11	write last Text record to object program
12	write End record to object program
13	write last listing line
14	end

```
#include<stdio.h>
                        #include<conio.h>
                                                    #include<string.h>
void main()
{
char
opcode[10],operand[10],symbol[10],label[10],code[10],mnemonic[5],character,add[10],objectco
de[10];
int flag,flag1,locetr,location,loc;
FILE *fp1,*fp2,*fp3,*fp4;
clrscr();
fp1=fopen("out.txt","r");
                               fp2=fopen("twoout.txt","w");
fp3=fopen("optab.txt","r");
                               fp4=fopen("symtbl.txt","r");
fscanf(fp1,"%s%s%s",label,opcode,operand);
if(strcmp(opcode,"START")==0)
fprintf(fp2,"%s\t%s\n",label,opcode,operand);
   fscanf(fp1,"%d%s%s%s",&locctr,label,opcode,operand);
while(strcmp(opcode,"END")!=0)
\{ flag=0;
   fscanf(fp3,"%s%s",code,mnemonic);
while(strcmp(code,"END")!=0)
{ if((strcmp(opcode,code)==0) && (strcmp(mnemonic,"*"))!=0)
          flag=1;
          break;
fscanf(fp3,"%s%s",code,mnemonic);
```

```
if(flag==1)
   { flag1=0;
                     rewind(fp4);
   while(!feof(fp4))
   fscanf(fp4,"%s%d",symbol,&loc);
   if(strcmp(symbol,operand)==0)
   flag1=1;
                     break;
   if(flag1==1)
   itoa(loc,add,10);
   strcpy(objectcode,strcat(mnemonic,add));
   else if(strcmp(opcode,"BYTE")==0 || strcmp(opcode,"WORD")==0)
   if((operand[0]=='C') \parallel (operand[0]=='X'))
   character=operand[2];
   itoa(character,add,16);
   strcpy(objectcode,add);
   }
   else
   itoa(atoi(operand),add,10);
   strcpy(objectcode,add);
   }
   else
   strcpy(objectcode,"\0");
   fprintf(fp2,"%s\t%s\t%s\t%\d\t%s\n",label,opcode,operand,locctr,objectcode);
   fscanf(fp1,"%d%s%s%s",&locctr,label,opcode,operand);
   fprintf(fp2,"%s\t%s\t%d\n",label,opcode,operand,locctr);
   fclose(fp1);
                            fclose(fp2);
   fclose(fp3);
                            fclose(fp4);
   getch();
d) Output:
INPUT FILES:
out.txt
                                   2000
                     START
2000
                                   FIVE
                     LDA
2003
       **
                     STA
                                   ALPHA
       **
2006
                     LDCH
                                   CHARZ
      **
2009
                     STCH
                                   C1
2012
      ALPHA
                     RESW
                                   1
2015
      FIVE
                     WORD
                                   5
```

Dr.NNCE			IT/VSem	SS Lab - LM
2018 2019 2020	CHARZ C1 **	BYTE RESB END	C'Z' 1 **	
	optab.txt			
STAR	T *			
LDA	03			
STA	0f			
LDCH	53			
STCH	57			
END	*			

sym	tbl.txt
ALPHA	2012
FIVE	2015
CHARZ	2018
C1	2019

OUTPUT FILES:

two	out.txt					
**	START		2000			
**	LDA	FIVE	2000		0320	15
**	STA	ALPHA	2003		0f201	2
**	LDCH	CHARZ	2006		5320	18
**	STCH	C1	2009		5720	19
ALPHA	RESW	1	2012			
FIVE	WORD	5		2015		5
CHARZ	BYTE	C'Z'	2018		5a	
C1	RESB	1	2019			
**	END	**	2020			

e) Result:

Thus pass two of two pass assembler is implemented and the result is verified successfully. **VIVA – QUESTION AND ANSWER:**

1. Define the basic functions of assembler.

- * Translating mnemonic operation codes to their machine language equivalents.
- * Assigning machine addresses to symbolic labels used by the programmer.

2. What is meant by assembler directives? Give example.

These are the statements that are not translated into machine instructions, but they Provide instructions to assembler itself.

Example START, END, BYTE, WORD, RESW and RESB

3. What is the need of SYMTAB (symbol table) in assembler?

The symbol table includes the name and value for each symbol in the source program, together with flags to indicate error conditions. Some times it may contain details about the data area. SYMTAB is usually organized as a hash table for efficiency of insertion and retrieval.

Exercise Number: 4

Title of the Exercise : IMPLEMENTATION OF A SINGLE PASS ASSEMBLER

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

To write a program to implement a single pass assembler.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no	Details of the step		
1	Begin		
2	Read first input line		
2	if OPCODE='START' then a. save #[operand] as starting address		
3	b. initialize LOCCTr as starting addressc. read next input lineend		
4	else initialize LOCCTR to 0		
5	while OPCODE != 'END' do d. if there is not a comment line then e. if there is a symbol in the LABEL field then i. search SYMTAB for LABEL ii. if found then 1. if symbol value as null 2. set symbol value as LOCCTR and search the linked list with the corresponding operand 3. PTR addresses and generate operand addresses as corresponding symbol values 4. set symbol value as LOCCTR in symbol table and delete the linked list iii. end iv. else insert (LABEL,LOCCTR) into SYMTAB v. end		
6	search OPTAB for OPCODE		
7	if found then		
	search SYMTAB for OPERAND address		
8	if found then		

	f. if symbol value not equal to null then	
	i) store symbol value as operand address	
	else insert at the end of the linked list with a node with address as LOCCTR	
9	else insert (symbol name, null) add 3 to LOCCTR.	
	elseif OPCODE='WORD' then	
10	add 3 to LOCCTR & convert comment to object code	
11	elseif OPCODE = 'RESW' then add 3 #[OPERND] to LOCCTR	
	elseif OPCODE = 'RESB' then	
12	add #[OPERND] to LOCCTR	
	elseif OPCODE = 'BYTE' then	
12	g. find length of the constant in bytes	
13	h. add length to LOCCTR	
	convert constant to object code	
	if object code will not fit into current text record then	
1.4	i. write text record to object program	
14	j. initialize new text record	
	o. add object code to text record	
15	write listing line	
16	read next input line	
17	write last text record to object program	
18	write end record to object program	
19	write last listing line	
20	End	

```
#include<stdio.h>
                         #include<conio.h>
#include<string.h>
                         #include<stdlib.h>
                                                       #define MAX 10
struct input
char label[10],opcode[10],operand[10],mnemonic[5];
                                                                      int loc;
struct input table[MAX];
struct symtab
char sym[10];
                         int f,val,ref;
struct symtab symtbl[MAX];
void main()
int f_{i,i=1,j=1,f} ag, locetr, x;
char add[10],code[10],mnemcode[5];
FILE *fp1,*fp2,*fp3;
clrscr();
fp1=fopen("input1.txt","r");
                                 fp2=fopen("optab1.txt","r"); fp3=fopen("spout.txt","w");
fscanf(fp1,"%s%s%s",table[i].label,table[i].opcode,table[i].operand);
if(strcmp(table[i].opcode, "START")==0)
locctr=atoi(table[i].operand);
```

```
i++;
fscanf(fp1,"%s%s%s",table[i].label,table[i].opcode,table[i].operand);
else
locctr=0;
while(strcmp(table[i].opcode,"END")!=0)
{ if(strcmp(table[i].label,"**")!=0)
for(x=1;x<j;x++)
{ f1=0;
if((strcmp(symtbl[x].sym,table[i].label)==0) && (symtbl[x].f==1))
symtbl[x].val=locetr;
                         symtbl[x].f=0;
table[symtbl[x].ref].loc=locctr;
                                        f1=1;
break;
if(f1==0)
strcpy(symtbl[j].sym,table[i].label); symtbl[j].val=locctr; symtbl[j].f=0;
j++;
fscanf(fp2,"%s%s",code,mnemcode);
while(strcmp(code,"END")!=0)
   if(strcmp(table[i].opcode,code)==0)
   { strcpy(table[i].mnemonic,mnemcode);
    locctr+=3;
for(x=1;x<=j;x++)
flag=0;
if(strcmp(table[i].operand,symtbl[x].sym)==0)
flag=1;
if(symtbl[x].f==0)
table[i].loc=symtbl[x].val;
break;
if(flag!=1)
strcpy(symtbl[i].sym,table[i].operand);
symtbl[i].f=1;
symtbl[j].ref=i;
j++;
fscanf(fp2,"%s%s",code,mnemcode);
rewind(fp2);
```

```
if(strcmp(table[i].opcode,"WORD")==0)
locctr+=3;
strcpy(table[i].mnemonic,'\0');
table[i].loc=atoi(table[i].operand);
else if(strcmp(table[i].opcode,"RESW")==0)
locctr+=(3*(atoi(table[i].operand)));
strcpy(table[i].mnemonic,'\0');
table[i].loc=atoi('\0');
else if(strcmp(table[i].opcode,"RESB")==0)
locctr+=(atoi(table[i].operand));
strcpy(table[i].mnemonic,'\0');
table[i].loc=atoi('\0');
else if(strcmp(table[i].opcode,"BYTE")==0)
++locctr;
if((table[i].operand[0]=='C') \parallel (table[i].operand[0]=='X'))
table[i].loc=(int)table[i].operand[2];
table[i].loc=locctr;
i++:
fscanf(fp1,"%s%s%s",table[i].label,table[i].opcode,table[i].operand);
for(x=1;x<=i;x++)
fprintf(fp3,"%s\t%s\t%s\t%s\n",table[x].label,table[x].opcode,table[x].operand,strcat(table[x].mn
emonic,itoa(table[x].loc,add,10)));
for(x=1;x<j;x++)
printf("%s\t%d\n",symtbl[x].sym,symtbl[x].val);
getch();
```

d) Output:

INPUT FILES:

<u>input1.txt</u>		
**	START	6000
**	JSUB	CLOOP
**	JSUB	RLOOP
ALPHA	WORD	23
BETA	RESW	3
GAMMA	BYTE	C'Z'
DELTA	RESB	4
CLOOP	LDA	ALPHA
RLOOP	STA	BETA
**	LDCH	GAMMA

**	STCH	DELTA
**	END	**

optab1.txt

START	*
JSUB	48
LDA	14
STA	03
LDCH	53
STCH	57
FND	*

OUTPUT FILES:

CLOOP 6023 RLOOP 6026 ALPHA 6006 BETA 6009 GAMMA 6018 DELTA 6019

spout.txt

**	START	6000	0
**	JSUB	CLOOP	486023
**	JSUB	RLOOP	486026
ALPHA	WORD	23	23
BETA	RESW	3	0
GAMMA	BYTE	C'Z'	90
DELTA	RESB	4	0
CLOOP	LDA	ALPHA	146006
RLOOP	STA	BETA	036009
**	LDCH	GAMMA	536018
**	STCH	DELTA	576019
**	END	**	0

e) Result:

Thus the program for single pass assembler is implemented and the output is verified accordingly.

VIVA – QUESTION AND ANSWER:

1. Define load and go assembler.

One pass assembler that generates their object code in memory for immediate execution is known as load and go assembler. Here no object programmer is written out and hence no need for loader.

2. What are the two different types of jump statements used in MASM assembler?

- Near jump: A near jump is a jump to a target in the same segment and it is assembled by using a current code segment CS.
- Far jump: A far jump is a jump to a target in a different code segment and it is assembled by using different segment registers.

Exercise Number: 5

Title of the Exercise : IMPLEMENTATION OF A MACRO PROCESSOR

Date of the Exercise:

OBJECTIVE (AIM) OF THE EXPERIMENT

To write a program to implement a macro processor.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no.	Details of the step	
1	Start the macro processor program.	
2	Include the necessary header files and variable.	
	Open the three files	
3	a.fl=macin.dat with read privilege	
3	b. f2=macout.dat with write privilege	
	c. f3= deftab.dat with write privilege	
4	Get the variable form f1 file macin.dat for label,opcode,operand	
5	Read the variable until the opcode is not is equal to zero	
	Then check if the opcode is equal to Macro if Macro	
	Then Copy macroname=label	
	a. Get the variable label, opcode, operand	
	b. In these if condition perform the while loop until opcode	
6	is not equal to MEND	
	c. Copy the variable	
	d. close while loop and if condition	
	e. else if opcode is equal to macro name	
	Perform the for loop from 0 to length	
7	Finally terminate the program.	

```
{ fscanf(fp1,"%s%s%s",ilab,iopd,oper);
if(stremp(iopd,"macro")==0)
n++;
printf("No. of macros=%d\n",n);
printf("Enter the text filename\n");
for(i=0;i< n;i++)
{ scanf("%s",fn[i]);
p[i]=fopen(fn[i],"w");
}
n=0;
rewind(fp1);
while(!feof(fp1))
{ fscanf(fp1,"%s%s%s",ilab,iopd,oper);
if(strcmp(iopd,"macro")==0)
{ strcpy(m[n],oper);
fscanf(fp1,"%s%s%s",ilab,iopd,oper);
while(strcmp(iopd,"mend")!=0)
{ fprintf(p[n],"%s %s %s\n",ilab,iopd,oper);
fscanf(fp1,"%s%s%s",ilab,iopd,oper);
fclose(p[n]);
n++;
}
for(i=0;i<n1;i++)
p[i]=fopen(fn[i],"r"); fp2=fopen("outmac.txt","w");
                                                           rewind(fp1);
fscanf(fp1,"%s%s%s",ilab,iopd,oper);
while(!feof(fp1))
{ if(strcmp(iopd,"call")==0)
\{ for(i=0;i< n1;i++) \}
\{ if(strcmp(m[i],oper)==0) \}
{ rewind(p[i]);
fscanf(p[i],"%s%s%s",ilab,iopd,oper);
while(!feof(p[i]))
{ fprintf(fp2,"%s %s %s\n",ilab,iopd,oper);
c1=1;
fscanf(p[i],"%s%s%s",ilab,iopd,oper);
break;
    }
if(c1!=1)
fprintf(fp2,"%s %s %s\n",ilab,iopd,oper);
c1=0:
fscanf(fp1,"%s%s%s",ilab,iopd,oper);
fprintf(fp2,"%s %s %s\n",ilab,iopd,oper);
```

d) Output:

INPUT FILE:

```
macin.txt
       macro m1
**
       move
                     a,b
**
       mend ---
**
       macro m2
**
       lda
                     b
**
       mend ---
**
       start
              1000
**
       lda
              a
**
       call
              m1
**
       call
              m2
**
       add
                     a,b
```

OUPUT FILE:

No. of macros=2

Enter the text filename

outmac

macin

outmac.txt

```
macro m1
**
       move a,b
**
      mend ---
**
      macro m2
**
      lda
                     b
**
      mend ---
**
             1000
      start
**
      lda
                     a
**
      move
                     a,b
**
      lda
                     b
**
      add
                     a,b
```

e) Result:

Thus the macro processor is implemented and the result is verified successfully.

VIVA – QUESTION AND ANSWER:

1. What are the basic functions of loaders?

- Loading brings the object program into memory for execution
- Relocation modifies the object program so that it can be loaded at an address different from the location originally specified
- Linking combines two or more separate object programs and also supplies the information needed to reference them.

2. Define absolute loader.

The loader, which is used only for loading, is known as absolute loader. e.g. Bootstrap loader

3. What is meant by bootstrap loader?

This is a special type of absolute loader which loads the first program to be run by the computer. (usually an operating system)

Exercise Number: 6

Title of the Exercise : IMPLEMENTATION OF AN ABSOLUTE LOADER

Date of the Exercise:

OBJECTIVE (AIM) OF THE EXPERIMENT

To write a program to implement an absolute loader.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no.	Details of the step		
1	Begin		
2	Read Header record		
3	Verify program name and length		
4	Read first Text record		
5	while record type ≠ 'E' do		
6	Begin a. {if object code is in character form, convert into		
7	End		
8	Jump to address specified in End record		
9	End		

```
#include<stdio.h>
#include<conio.h>
#include<string.h>
#include<stdlib.h>
struct object_code
{
int locctr;
char byte[5];
};
struct object_code code[200];
```

```
void main()
FILE *fp1,*fp2;
char input[15];
int i,len,n=0,count=0,inc=0,textloc,tlen,tloc=0,num=0,loc;
clrscr();
fp1=fopen("loadin.txt","r");
fp2=fopen("loadout.txt","w");
rewind(fp1);
rewind(fp2);
fscanf(fp1,"%s",input);
if(strcmp(input,"H")==0)
for(i=0;i<4;i++)
if(i==1)
fscanf(fp1,"%x",&loc);
else
fscanf(fp1,"%s",input);
}
tloc=loc;
while(strcmp(input,"E")!=0)
if(strcmp(input,"T")==0)
fscanf(fp1,"%x",&textloc);
for(i=0;i<(textloc-(tloc+tlen));i++)
strcpy(code[inc].byte,"xx");
code[inc++].locctr=loc++;
fscanf(fp1,"%x",&tlen);
tloc=textloc;
}
else
len=strlen(input);
for(i=0;i<len;i++)
code[inc].byte[num++]=input[i];
if(num > 1)
code[inc].locctr=loc;
loc++;
inc++;
num=0;
```

```
fscanf(fp1,"%s",input);
}
n=0;
i=0;
count=0;
fprintf(fp2,"%x\t",code[i].locctr);
for(i=0;i<inc;i++)
fprintf(fp2,"%s",code[i].byte);
if(n > 3)
fprintf(fp2,"\t");
n=0;
count++;
if(count > 3)
fprintf(fp2,"\n%x\t",code[i+1].locctr);
count=0;
}
getch();
```

d) Output:

INPUT FILE:

loadin.txt

Н	COPY 00200	00	00107	a				
T	002000	1e	14203	3 48	3039	102036	282030	302015
	483061	3c200	300202	a0c203900	202d			
T	00201e15	2C20	36	483062	18203	3 40	C0000	454f46
	200003	10000	00					
T	002039 1e	24203	30	302030	e0305	d30303fd	8305d	
	282030	30305	57	53a0392c	305e38303	f		
T	002057	a	10203	6 4c	0000F1	201000		
T	002071	19	34203	0 e03	307930306	44 41	fa039 dc3079	
	2c203638306	54	4c0000	015				
E	002000							

OUTPUT FILE:

loadout.txt

2000	14203348	30391020	36282030	30201548
2010	30613c20	0300202a	0c203900	202d2C20
2020	36483062	1820334C	0000454f	46200003
2030	100000xx	XXXXXXX	xx242030	302030e0
2040	305d3030	3fd8305d	28203030	305753a0
2050	392c305e	38303f10	20364c00	00F12010
2060	00xxxxxx	XXXXXXX	XXXXXXX	xxxxxxx
2070	xx342030	e0307930	30644fa0	39dc3079
	2080 2c20	3638 30644	4c00 0015	

e) Result:

Thus the absolute loader is implemented and the result is verified successfully.

VIVA – QUESTION AND ANSWER:

1. Define bit mask.

The relocation bits are gathered together following the length indicator in each text record and which is called as bit mask. For e.g. the bit mask FFC (111111111100) specifies that the first 10 words of object code are to be modified during relocation.

2. What is the need of ESTAB?

It is used to store the name and address of the each external symbol. It also indicates in which control section the symbol is defined.

3. Write the two passes of a linking loader.

Pass1: assigns address to all external symbols

Pass2: it performs actual loading, relocation and linking.

4. Define automatic library search.

In many linking loaders the subroutines called by the program being loaded are automatically fetched from the library, linked with the main program and loaded. This feature is referred to as automatic library search.

5. Define dynamic linking.

If the subroutine is loaded and linked to the program during its first call (run time), then it is called as dynamic loading or dynamic linking.

6. Write the advantage of dynamic linking.

- It has the ability to load the routine only when they are needed.
- The dynamic linking avoids the loading of entire library for each execution.

Exercise Number: 7

Title of the Exercise : IMPLEMENTATION OF A RELOCATING LOADER

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

To write a program to implement a relocating loader.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no	Details of the step
1	Start the program.
2	Include the necessary header file and variable
3	Open the two file for
	a.fp1= relinput.dat and give read
	b.fp2= reloutput.dat and give write
4	Read the content using while loop perform the loop Until character is not
	equal to E
	a.while (strcmp (input,"E")!=0)
	i. If the character is H Get the variable add, length, and input
	ii.Else if the character is T Get the variable address and bitmask
	And perform the for loop for starting zero to up to len
	Get the opcode ,addr and assign relocbit to bitmask
	iii. If relocabit is zero Then actualadd=addr;
	iv. Else
	Add the addr and star value
5	Finally terminate the program.

```
printf("Enter the location where the present program has to be loaded:");
scanf("%s",stloc);
                       start=atoi(stloc);
location=start;
tloc=start;
fscanf(fp1,"%s",input[i]);
while(strcmp(input[i],"T")!=0)
strcpy(output[i],input[i]);
i++;
fscanf(fp1,"%s",input[i]);
strcpy(output[i],input[i]);
itoa(start,output[2],10);
while(strcmp(input[i],"E")!=0)
strcpy(output[i],input[i]);
if(strcmp(input[i],"T")==0)
for(j=0;j<3;j++)
i++;
fscanf(fp1,"%s",input[i]);
strcpy(output[i],input[i]);
bitmask=atoi(output[i]);
itoa(bitmask,binary,2);
strcpy(output[i],NULL);
textloc=atoi(output[i-2]);
textloc=textloc+start;
itoa(textloc,output[i-2],10);
for(n=0;n<(textloc-(tloc+tlen));n++)
strcpy(obcode[inc].add,"xx");
obcode[inc++].locetr=location++;
tlen=atoi(output[i-1]);
tloc=textloc;
k=0;
else
if(binary[k]=='1')
num=0;
len=strlen(output[i]);
strcpy(address,NULL);
for(j=2;j<len;j++)
address[num]=output[i][j];
output[i][j]='0';
```

```
num++;
loc=atoi(address);
                             loc=loc+start;
                                                   itoa(loc,address,10);
strcat(output[i],address);
}
k++;
len=strlen(output[i]);
                             num=0;
for(n=0;n<len;n++)
{obcode[inc].add[num++]=output[i][n];
if(num>1)
{obcode[inc++].locctr=location++;
num=0;
                             }
i++;
fscanf(fp1,"%s",input[i]);
strcpy(output[i],input[i]);
i++;
fscanf(fp1,"%s",input[i]);
loc=atoi(input[i]);
loc=loc+start;
strcpy(output[i],itoa(loc,address,10));
count=0;
i=0;
n=0:
fprintf(fp2,"%d\t",obcode[n].locctr);
for(n=0;n<inc;n++)
fprintf(fp2,"%s",obcode[n].add);
i++;
if(i > 3)
fprintf(fp2,"\t");
                             i=0;
                                            count++;
if(count > 3)
fprintf(fp2,"\n%d\t",obcode[n+1].locctr);
                                                   count=0;
getch();
```

d) Output:

INPUT FILE:

relin.txt

H COPY 000000 001073

T 000000 10 015 140033 481039 100036 280030 300015 481061 311003 200030 211033 200033

T 000011 19 045 412036 481061 380033 412000 454196 100003 200000

T 000031 15 135 140030 430030 141013 301044 241064 210030 301057

543039 212064 381045

T 000058 05 056 100036 520000 151 301000

T 000065 19 080 340030 141079 301064 503039 152079 220036 381064

430000 25

E 000000

OUTPUT FILE:

relout.txt

3000	14303348	40391030	36283030	30001548
3016	10613110	03200030	21103320	0033xx41
3032	50364810	61383033	41500045	41961030
3048	03200000	xx143030	43003014	10133010
3064	44241064	21303030	40575460	39212064
3080	381045xx	XXXXXXXX	XXXXXXXX	xxxxxx10
3096	30365230	00154000	301000xx	xx343030
3112	14107930	40645030	39152079	22003638
3128	10644300	0025		

e) Result:

Thus the relocating loader is implemented and the result is verified successfully.

VIVA – QUESTION AND ANSWER:

1. Define macro processor.

Macro processor is system software that replaces each macroinstruction with the corresponding group of source language statements. This is also called as expanding of macros

2. What do macro expansion statements mean?

These statements give the name of the macroinstruction being invoked and the arguments to be used in expanding the macros. These statements are also known as macro call

3. What is the use of macro time variable?

Macro time variable can be used to store working values during the macro expansion. Any symbol that begins with the character & and then is not a macro instruction parameter is assumed to be a macro time variable.

4. What is meant by line by line processor?

This macro processor reads the source program statements, process the statements and then the output lines are passed to the language translators as they are generated, instead of being written in an expanded file.

5. Give the advantages of general-purpose macro processors.

- The programmer does not need to learn about a macro facility for each compiler.
- Overall saving in software development cost and maintenance cost.

6. What is the symbol used to generate unique labels?

\$ Symbol is used in macro definition to generate unique symbols. Each macro expansion the \$ symbol is replaced by \$XX, where XX is the alpha numeric character.

Exercise Number: 8

Title of the Exercise : IMPLEMENTATION OF PASS 1 OF A DIRECT LINKING

LOADER

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

• To write programs to implement pass 1 of a direct linking loader.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step	Details of the step
no.	
1	Begin
2	get PROGADDR from operating system
3	set CSADDR to PROGADDR {for first control section}
4	while not end of input do
5	i. read next input record {Header record for control section} ii. set CSLTH to control section length iii. search ESTAB for control section name iv. if found then set error flag {duplicate external symbol} v. else enter control section name into ESTAB with value CSADDR
6	while record type ≠ 'E' do vi. begin 1. read next input record 2. if record type = 'D' then a. for each symbol in the record do b. begin i. search ESTAB for symbol name ii. if found then set error flag (duplicate external symbol) iii. else enter symbol into ESTAB with value (CSADDR + indicated address) end {for}

7	end {while ≠ 'E'}
8	add CSLTH to CSADDR {starting address for next control section
9	end {while not EOF}
10	end {Pass 1}

```
#include<stdio.h>
#include<conio.h>
#include<stdlib.h>
#include<string.h>
#define MAX 10
struct estab
char csect[10];
char sym name[10];
long int add;
int length;
}table[MAX];
void main()
FILE *fp1,*fp2;
char input[10];
long int i,count=0,start,length,loc;
clrscr();
fp1=fopen("dlinkin.txt","r");
fp2=fopen("dlinkout.txt","w");
printf("\nEnter the location where the program has to be loaded:");
scanf("%lx",&start);
fprintf(fp2,"CSECT\tSymname\tAddress\tLength\n");
rewind(fp1);
while(!feof(fp1))
fscanf(fp1,"%s",input);
if(strcmp(input,"H")==0)
fscanf(fp1,"%s",input);
strcpy(table[count].csect,input);
strcpy(table[count].sym name,"\0");
fscanf(fp1,"%s",input);
table[count].add=atoi(input)+start;
fscanf(fp1,"%s",input);
length=atoi(input);
table[count++].length=atoi(input);
fscanf(fp1,"%s",input);
if(strcmp(input,"D")==0)
fscanf(fp1,"%s%lx",input,&loc);
while((strcmp(input,"R")!=0))
```

```
strcpy(table[count].csect,"\0");
strcpy(table[count].sym name,input);
table[count].add=loc+start;
table[count++].length=0;
fscanf(fp1,"%s%lx",input,&loc);
while(strcmp(input,"T")!=0)
fscanf(fp1,"%s",input);
if(strcmp(input,"T")==0)
while(strcmp(input,"E")!=0)
fscanf(fp1,"%s",input);
fscanf(fp1,"%s",input);
start=start+length;
for(i=0;i < count;i++)
fprintf(fp2,"%s\t%s\t%d\n",table[i].csect,table[i].sym_name,table[i].add,table[i].length);
getch();
d) Output:
```

INPUT FILE:

dlinkin.txt

```
H PROGA
            000000
                        000070
D LISTA 000040 ENDA 000054
R LISTB ENDB LISTC ENDC
T 000020 10 03201D 77100004 150014
T 000054 16 100014 15100006 00002F 100014 FFFFC0
M 000024 05 +LISTB
M 000054 06 +LISTC
M 000058 06 +ENDC
M 000064 06 +LISTB
E 000000
H PROGB 000000 000088
D LISTB 000060 ENDB 000070
R LISTA ENDA LISTC ENDC
T 000036 11 03100000 772027 05100000
T 000070 18 100000 05100006 05100020 05100030 100000
M 000037 05 +LISTA
M 000044 05 +ENDA
M 000070 06 +ENDA
M 000074 06 +ENDC
M 000078 06 +ENDC
M 000082 06 +ENDA
E 000000
H PROGC 000000 000057
D LISTC 000030 ENDC 000042
R LISTA ENDA LISTB ENDB
T 000018 12 03100000 77100004 05100000
```

T 000042 15 100030 100008 100011 100000 100000

M 000019 05 +LISTA

M 000023 05 +LISTB

M 000027 05 +ENDA

M 000048 06 +LISTA

M 000051 06 +ENDA

M 000054 06 +LISTB

E 000000

OUTPUT FILE:

dlinkout.txt

CSECT	Symname	Addres	S	Length
PROGA		2000		70
	LISTA 2040		0	
	ENDA	2054		0
PROGB		2046		88
	LISTB 20a6		0	
	ENDB	20b6		0
PROGC		209e		57
	LISTC 20ce		0	
	ENDC	20e0		0

e) Result:

Thus the pass 1 of a direct linking loader is implemented and the result is verified successfully.

VIVA – QUESTION AND ANSWER:

1. Define interactive editor?

An interactive editor is a computer program that allows a user to create and revise a target document. The term document includes objects such as computer programs, text, equations, tables, diagrams, line art, and photographs any thing that one might find on a printed page.

2. What are the tasks performed in the editing process?

4 tasks

- Select the part of the target document to be viewed and manipulated.
- Determine how to format this view on-line and how to display it.
- Specify and execute operations that modify the target document.
- Update the view appropriately.

3. What are the three categories of editor's devices?

- Text device/ String devices
- Button device/Choice devices
- Locator device

4. What is the function performed in editing phase?

In the actual editing phase, the target document is created or altered with a set of operations such as insert, delete, replace, move and copy.

5. Define Locator device?

Locator devices are two-dimensional analog-to-digital converters that position a cursor symbol on the screen by observing the user's movement of the device. The most common such devices for editing applications are the mouse and the data tablet.

Exercise Number: 9

Title of the Exercise : IMPLEMENTATION OF PASS 2 OF A DIRECT LINKING LOADER

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

• To write a program to implement pass 2 of a direct linking loader.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step	Details of the step		
no.			
1	Begin		
2	set CSADDR to PROGADDR		
3	set EXECADDR to PROGADDR		
4	while not end of input do		
5	begin		
	i. read next input record {Header record}		
	ii. set CSLTH to control section length		
	iii. while record type ≠'E' do		
	iv. begin		
	1. read next input record		
	2. if record type = 'T' then		
	a. begin		
	b. {if object code is in character form, convert into internal		
	representation}		
	c. move object code from record to location		
	i. (CSADDR + specified address)		
	end {if 'T'}		
6	else if record type = 'M' then d. begin		
0	e. search ESTAB for modifying symbol name		
	f. if found then		
	i. add or subtract symbol value at location		
	ii. (CSADDR + specified address)		
	g. else		
	i. set error flag (undefined external symbol)		
	h. end {if 'M'}		
	v. end{while \neq'E'}		

	vi. if an address is specified {in End record} then set EXECADDR to
	(CSADDR + specified address)
	add CSLTH to CSADDR.
7	end {while not EOF}
	jump to location given by EXECADDR {to start execution of loaded program}
8	
9	end {Pass 2}

```
#include<stdio.h>
                             #include<conio.h>
#include<string.h>
                             #include<stdlib.h>
struct ext table
{char csect[10];
                             char sname[10];
                                                           int padd;
                                                                          int plen;
}estab[20];
struct object code
char code[15];
                      int add;
}obcode[500];
void main()
FILE *fp1,*fp2,*fp3;
int i,j,x,y,n=0,num=0,inc=0,count=0,record=0,pstart,exeloc,start,textloc;
int loc,mloc[30],textlen,mlen[30],length,location;
long int newadd;
char *add1,operation,lbl[10],input[10],label[30][10],
address[10];
clrscr();
printf("\n this is pass 2 of direct linking loader\n");
fp1=fopen("dlink2in.txt","r");
                                            fp2=fopen("dlink1out.txt","r");
fp3=fopen("dlink2out.txt","w");
                                            rewind(fp1);
rewind(fp2);
                      rewind(fp3);
while(!feof(fp2))
fscanf(fp2,"%s%s%d%d",estab[num].csect,estab[num].sname,&estab[num].padd,&estab[num].plen);
num++;
exeloc=estab[0].padd;
                                     loc=exeloc:
                                                           start=loc;
while(!feof(fp1))
fscanf(fp1,"%s",input);
if(stremp(input,"H")==0)
fscanf(fp1,"%s",input);
for(i=0;i \le num;i++)
if(strcmp(input,estab[i].csect)==0)
pstart=estab[i].padd;
                                     break;
```

```
while(strcmp(input,"T")!=0)
fscanf(fp1,"%s",input);
do
       if(strcmp(input,"T")==0)
              fscanf(fp1,"%d",&textloc);
       textloc=textloc+pstart;
for(i=0;i<(textloc-loc);i++)
strcpy(obcode[inc].code,"xx");
obcode[inc++].add=start++;
fscanf(fp1,"%d",&textlen);
loc=textloc+textlen;
else if(strcmp(input,"M")==0)
fscanf(fp1,"%d",&mloc[record]);
                                            mloc[record]=mloc[record]+pstart;
                                            fscanf(fp1,"%s",label[record++]);
fscanf(fp1,"%d",&mlen[record]);
}
else
length=strlen(input);
for(i=0;i<length;i++)
obcode[inc].code[x++]=input[i];
if(x>1)
obcode[inc++].add=start++;
x=0;
fscanf(fp1,"%s",input);
} while(strcmp(input, "E")!=0);
if(strcmp(input,"E")==0)
fscanf(fp1,"%s",input);
for(n=0;n<record;n++)
operation=label[n][0];
length=strlen(label[n]);
for(i=1;i<length;i++)
lbl[i-1]=label[n][i];
lbl[length-1]='\0';
length=0;
strcpy(address,"\0");
location=mloc[n]-exeloc;
loc=location;
```

```
count=0;
while(length<mlen[n])</pre>
strcat(address,obcode[location++].code);
count++;
length+=2;
for(i=0;i < num;i++)
if(strcmp(lbl,estab[i].sname)==0)
break;
switch(operation)
case '+':newadd=strtol(address,&add1,10)+(long int)estab[i].padd;
                                                                                   break;
case '-':newadd=strtol(address,&add1,10)-(long int)estab[i].padd;
                                                                           break;
ltoa(newadd,address,10);
x=0;
                      y=0;
while(count > 0)
       obcode[loc].code[x++]=address[y++];
if(x > 1)
x=0;
loc++;
count--;
count=0;
n=0:
fprintf(fp3,"%d\t",obcode[0].add);
for(i=0;i\leq inc;i++)
fprintf(fp3,"%s",obcode[i].code);
n++;
if(n>3)
fprintf(fp3,"\t");
n=0;
count++;
if(count>3)
fprintf(fp3,"\n%d\t",obcode[i+1].add);
count=0;
}
getch();
```

d) Output:

INPUT FILES:

dlink1in.txt

H PROGA 000000 000070

D LISTA 000040 ENDA 000054

R LISTB ENDB LISTC ENDC

T 000020 10 03201D 77100004 150014

T 000054 16 100014 15100006 00002F 100014 FFFFC0

M 000024 05 +LISTB

M 000054 06 +LISTC

M 000058 06 +ENDC

M 000064 06 +LISTB

E 000000

H PROGB 000000 000088

D LISTB 000060 ENDB 000070

R LISTA ENDA LISTC ENDC

T 000036 11 03100000 772027 05100000

T 000070 18 100000 05100006 05100020 05100030 100000

M 000037 05 +LISTA

M 000044 05 +ENDA

M 000070 06 +ENDA

M 000074 06 +ENDC

M 000078 06 +ENDC

M 000082 06 +ENDA

E 000000

H PROGC 000000 000057

D LISTC 000030 ENDC 000042

R LISTA ENDA LISTB ENDB

T 000018 12 03100000 7100004 05100000

T 000042 15 100030 100008 100011 100000 100000

M 000019 05 +LISTA

M 000023 05 +LISTB

M 000027 05 +ENDA

M 000048 06 +LISTA

M 000051 06 +ENDA

M 000054 06 +LISTB

E 000000

dlink1out.txt

PROG.	A	**	2000	70
**	LISTA		2040	0
**	ENDA	2054	0	
PROG:	В	**	2070	88
**	LISTB		2130	0
**	ENDB	2140	0	
PROG	C	**	2158	57
**	LISTC		2188	0
**	ENDC	2200	0	

OUTPUT FILE:

Dlink2out.txt

2000	XXXXXXX	XXXXXXX	XXXXXXX	XXXXXXX
2016	XXXXXXX	03201D77	10213415	0014xxxx
2032	XXXXXXX	XXXXXXX	XXXXXXXX	XXXXXXXX
2048	XXXXXXX	xxxx1022	02151022	0600002F
2064	102144FF	FFC0xxxx	XXXXXXXX	XXXXXXXX
2080	XXXXXXX	XXXXXXX	XXXXXXXX	XXXXXXX
2096	XXXXXXX	XXXXXXX	xxxx0310	20407720
2112	27051020	54xxxxxx	XXXXXXXX	XXXXXXXX
2128	XXXXXXX	XXXXXXX	XXXXXXXX	10205405
2144	10220605	10222005	10208410	0000xxxx
2160	XXXXXXX	XXXXXXX	XXXXXXXX	XXXXXXXX
2176	03102040	77102134	05102054	XXXXXXXX
2192	XXXXXXX	xxxxxxxx	10003010 000810	020
2208	51102054	102130		

e) Result:

Thus the pass 2 of a direct linking loader is implemented and the result is verified successfully.

VIVA – QUESTION AND ANSWER:

1. What is the function performed in voice input device?

Voice-input devices, which translate spoken words to their textual equivalents, may prove to be the text input devices of the future. Voice recognizers are currently available for command input on some systems.

2. What are called tokens?

The lexical analyzer tracks the source program one character at a time by making the source program into sequence of atomic units is called tokens.

3. Name some of typical tokens.

Identifiers, keywords, constants, operators and punctuation symbols such as commas and parentheses are typical tokens.

4. Mention the main disadvantage of interpreter.

The main disadvantage of interpreter is that the execution time of interpreted program is slower than that of a corresponding compiled object program.

5. What is meant by code optimization?

The code optimization is designed to improve the intermediate code, which helps the object program to run faster and takes less space.

Exercise Number: 10

Title of the Exercise : IMPLEMENTATION OF A SIMPLE TEXT EDITOR

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

• To Create a Component for retrieving stock market exchange information using CORBA

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no.	Details of the step	
1	Begin	
2	Provide a blank screen for the user to type the document.	
3	Instruct the user to enter the text using the editor.	
4	Characters displayed on the screen are stored in a character array.	
5	If SAVE option is selected	
	a. Get the data file name from the user.	
	b. Open the file in write mode.	
	c. Move the content of the character array to the file and	
	close the file.	
	1. If OPEN option is selected	
	a. Open the data file in the read mode.	
	b. Read one character at a time from the file and moved it to	
	the character array.	
	c. Repeat the above steps until the end of file is reached.	
	d. Display the character from the character array on the	
	screen.	
	Finally close the text editor and exit.	

c) Program:

#include<stdio.h> #include<conio.h> #include<dos.h> #include<fstream.h> #include<fstream.h>

char filename[15]; char buff[1000]; int curx,cury,count;

```
void cur pos()
curx=wherex();
cury=wherey();
textcolor(12);
textbackground(9);
gotoxy(35,20);
cout << "\n";
cout << "\n";
                                   TEXT EDITOR
                                                           n'';
cout<<"
cout << "\n";
cout<<"\n Type ur text and then press escape key\n";
gotoxy(100,100);
cprintf("%2d%2d",curx,cury);
gotoxy(curx,cury);
void main()
char ch,c;
ofstream outfile;
ifstream infile;
clrscr();
cur pos();
curx=wherex();
cury=wherey();
while(c!=27)
c=getch();
switch(c)
case 80:gotoxy(curx,cury+1);
break;
case 77:gotoxy(curx+1,cury);
break:
case 72:gotoxy(curx,cury-1);
break;
case 75:gotoxy(curx-1,cury);
break;
case 32:printf(" ");
buff[count++]='';
break;
case 13:gotoxy(1,cury+1);
buff[count++]='\n';
break;
default:textcolor(13);
if((c \ge 65) \&\& (c \le 122) || (c \ge 48 \&\& c \le 57))
printf("%c",c);
```

```
break;
}
buff[count++]=c;
cur pos();
cprintf("\n\nDo u want to save?");
scanf("%c",&c);
if(c == 'y')
cprintf("\n\nEnter the filename with extension in 8 char only: ");
scanf("%s",filename);
outfile.open(filename,ios::out);
outfile << buff;
outfile.close();
cout << "\nDo u want to open\n";
ch=getch();
if(ch=='y')
cprintf("\n\nEnter the filename to open:
                                           ");
scanf("%s",filename);
infile.open(filename,ios::in);
infile.get(buff,count,'*');
gotoxy(90,1);
printf("%s",buff);
getch();
infile.close();
```

d) Output:

hi welcome

This is a sample text editor program for ECE data structures lab

Do u want to save? y

Enter the filename with extension in 8 char only: sample.txt

Do u want to open

Enter the filename to open: sample.txt hi welcome

This is a sample text editor program for ECE data structures lab

TEXT EDITOR

Type ur text and then press escape key

28 5

e) Result:

Thus the text editor is implemented and the operations are verified successfully.

VIVA – QUESTION AND ANSWER:

1. What is error handler?

The error handler is used to check if there is an error in the program. If any error, it should warn the programmer by instructions to proceed from phase to phase.

2. Name some of text editors.

- line editors
- stream editors
- screen editors
- word processors
- structure editors

3. Mention the features of word processors.

- moving text from one place to another
- merging of text
- searching
- word replacement

4. Define traveling phase.

The phase specifies the region of interest. Traveling is achieved using operations such as next screen, bottom, find pattern.

5. Give the components of editor structure

4 components

- Editing component
- Traveling component
- Viewing component
- Display component

6. What are the basic types of computing environments used in editor's functions?

Editor's function in three basic types of computing environments

i. Time sharing ii. Stand-alone iii. Distributed

Exercise Number: 11

Title of the Exercise : LINE AND SCREEN EDITOR

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

• To write a program to implement line and screen editor **FACILITIES REQUIRED AND PROCEDURE**

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no	Details of the step
1	Initialize all the variables.
Design a menu through which we can create a symbol Table	
2	perform operations as insert, modify, search and display.
3	Create a symbol table with fields as variable and value using create
3	option. Entries may be added to the table while it is created itself.
4	Append new contents to the symbol table with the constraint that there
	is no duplication of entries, using insert option.
5	Modify existing content of the table using modify option.
6	Use display option to display the contents of the table.
7	End of the program.

```
#include<stdio.h>
#include<conio.h>
#include<graphics.h>
void main()
int i,j,k,q;
clrscr();
textattr(10);
gotoxy(1,1);
cprintf("LINE AND SCREEN EDITOR\n\n");
gotoxy(5,5);
cprintf("at 5,5");
gotoxy(20,5);
cprintf("at 5,20");
gotoxy(5,7);
cprintf("at 7,5");
gotoxy(20,7);
```

```
cprintf("at 7,20");
gotoxy(5,9);
cprintf("at 9,5");
gotoxy(20,9);
cprintf("at 9,20");
for(j=0;j<q;j+2)
{
    gotoxy(30,i);
    cprintf("at %d 30",i);
    gotoxy(40,i);
    cprintf("at %d 40",i);
}
getch();
}</pre>
```

d) Output:

LINE AND SCREEN EDITOR

```
at 5,5 at 5,20 at 7,5 at 7,20 at 9,5 at 9,20
```

e) Result:

Thus the line and screen editor is created and operations are verified successfully.

VIVA – QUESTION AND ANSWER:

1. What is system software?

System software consists of variety of programs that supports the operations of a computer. This makes it possible for the user to focus on an application or other problem to be solved, without needing to know the details of how the machine works internally. Examples of system software are text-editors, compilers, loaders or linkers, debuggers, assemblers, and operating systems.

2. Give some applications of operating system.

- > to make the computer easier to use
- > to manage the resources in computer
- > process management
- data and memory management
- > to provide security to the user.
- Operating system acts as an interface between the user and the system Eg: windows, Linux, UNIX, dos

3. What is SIC machine?

SIC refers to Simplified Instruction Computer which is a hypothetical computer that has been designed to include the hardware features most often found on real machines, while avoiding unusual and irrelevant complexities. This allows to clearly separating the central concepts of system software from the implementation details associated with a particular machine

Exercise Number: 12

Title of the Exercise : BOOTSTRAP LOADERS

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

To write a program to implement bootstrap loaders.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no	Details of the step
1	Initialize all the variables.
2	Design a menu through which we can create a symbol Table and perform operations as insert, modify, search and display.
3	Include the necessary header files and variable.
4	Open the three files a.fl=macin.dat with read privilege b. f2=macout.dat with write privilege c. f3= deftab.dat with write privilege
5	Get the variable form f1 file macin.dat for label,opcode,operand
6	Read the variable until the opcode is not is equal to zero
7	End of the program.

```
#include<stdio.h>
                             #include<string.h>
char line[5][20];
                             void numtok();
int cur=0;
                      int lth;
void main()
FILE *fin;
int i=0,m=0,num=0,p=0,k=0,1,j;
                                            int prgstart,prglen,textaddr,disp;
char str[70];
struct loccontent
int addr;
              char content[3];
}mem[200];
clrscr();
for(i=0;i<200;i++)
mem[i].content[0]='x';
                             mem[i].content[1]='x';
                                                           mem[i].content[2]='/0';
fin=fopen("loadin.txt","r");
```

```
while(!feof(fin))
fgets(str,70,fin);
                              numtok(str);
if(strcmp(line[0],"H")==0)
printf("Program name is%s\n",line[1]);
                                                    prgstart=atoi(line[2]);
prglen=atoi(line[3]);
for(i=prgstart;i<=prgstart+prglen;i++)
mem[j].addr=i;
                              j ++;
else if(strcmp(line[0],"T")==0)
textaddr=atoi(line[1]);
                                     disp=textaddr-prgstart;
for(m=3;m<lth;++m)
for(i=0;i<=5;i+=2)
mem[disp].addr=textaddr++;
                                             mem[disp].content[0]=line[m][i];
mem[disp].content[1]=line[m][i+1];
                                             mem[disp].content[2]='\0';
disp++;
else if(strcmp(line[0],"E")==0)
break;
k=0;
for(i=0;iprglen;i+=16)
{printf("%d\t",mem[k].addr);
for(1=0;1<4;1++)
\{for(j=0;j<4;j++)\}
{printf("%c%c",mem[k].content[0],mem[k].content[1]);
printf("\t");
printf("\n");
getch();
void numtok(char str[70])
int i=0, j=0, k=0;
while(str[j]!='\0')
\{line[i][k]=str[j];
if(str[j]=='^')
\{line[i][k]='\0';
k=-1;
i++;
```

```
k++;
j++;
}
line[i][k]='\0';
lth=i+1;
}
```

d) Output:

INPUT FILE:

loadin.txt:

H^COPY^1000^115

T^1000^12^141003^482039^001036^281030

T^1012^9^345890^581056^454645

T^1012^12^141033^482049^001036^281030

E^1000

OUTPUT FILE:

Progra	am name isCC)PY		
1000	14100348	20390010	36281030	14103348
1016	20490010	36281030	XXXXXXXX	XXXXXXXX
1032	XXXXXXX	XXXXXXXX	XXXXXXXX	XXXXXXX
1048	XXXXXXX	XXXXXXXX	XXXXXXXX	XXXXXXX
1064	XXXXXXX	XXXXXXXX	XXXXXXXX	XXXXXXXX
1080	XXXXXXX	XXXXXXXX	XXXXXXXX	XXXXXXXX
1096	XXXXXXX	XXXXXXXX	XXXXXXXX	XXXXXXXX
1112	XXXXXXXX	XXXXXXXX	XXXXXXXX	XXXXXXXX

e) Result:

Thus the bootstrap loaders is created and operations are verified successfully

VIVA – QUESTION AND ANSWER:

1. What is system software?

System software consists of variety of programs that supports the operations of a computer. This makes it possible for the user to focus on an application or other problem to be solved, without needing to know the details of how the machine works internally. Examples of system software are text-editors, compilers, loaders or linkers, debuggers, assemblers, and operating systems.

2. Give some applications of operating system.

to make the computer easier to use
 process management
 to manage the resources in computer
 data and memory management

- > to provide security to the user.
- Operating system acts as an interface between the user and the system Eg: windows, Linux, UNIX, dos

3. What is SIC machine?

SIC refers to Simplified Instruction Computer which is a hypothetical computer that has been designed to include the hardware features most often found on real machines, while avoiding unusual and irrelevant complexities. This allows to clearly separating the central concepts of system software from the implementation details associated with a particular machine

Exercise Number: 13

Title of the Exercise : MULTI PASS ASSEMBLERS

Date of the Exercise :

OBJECTIVE (AIM) OF THE EXPERIMENT

To write a program to implement symbol table.

FACILITIES REQUIRED AND PROCEDURE

a) Facilities Required:

S.No.	Facilities required	Quantity
1	System	1
2	O/S	Windows XP
3	S/W name	C or C++

b) Procedure:

Step no	Details of the step		
1	Initialize all the variables.		
2	write line to intermediate file		
3	read next input line		
4	Append new contents to the symbol table with the constraint that there		
4	is no duplication of entries, using insert option.		
5	Modify existing content of the table using modify option.		
6	Use display option to display the contents of the table.		
7	End of the program.		

```
#include<stdlib.h>
#include<conio.h>
#include<stdio.h>
#include<string.h>
int search(char sym[10]);
char praddr[10];
int k,data[100],mem[100],i=0,q,d=0, value,j;
int csaddr,exaddr,cslth;
char type[10],addr[10],obj[10],mflag[10],op;
void main()
FILE *in, *address;
in=fopen("OBJECT.TXT","r");
address=fopen("ADDRESS.TXT","r");
clrscr();
fscanf(address,"%s",praddr);
fclose(address);
csaddr=atoi(praddr);
```

```
exaddr=atoi(praddr);
fscanf(in,"%s\t%s\t%s\t%s\t",type,addr,obj,mflag);
while(!feof(in))
cslth=atoi(obj);
while(strcmp(type, "E")!=0)
fscanf(in,"%s\t%s\t%s\t%s\t",type,addr,obj,mflag);
if(strcmp(type, "T")==0)
mem[i]=atoi(addr)+csaddr;
data[i]=atoi(obj);
i++;
else if(strcmp(type,"M")==0)
op=mflag[0];
for(j=1;j<strlen(mflag);j++)
mflag[j-1]=mflag[j];
mflag[j-1]='\0';
value=search(mflag);
q=d+atoi(addr)/3;
if(op=='+')
data[q]=data[q]+value;
else
data[q]=data[q]-value;
d=cslth/3;
exaddr=csaddr+atoi(addr);
csaddr=csaddr+cslth:
fscanf(in,"%s\t%s\t%s\t%s\n",type,addr,obj,mflag);
fclose(in);
q=(exaddr-(atoi(praddr)))/3;
printf("Program Loaded in memory is ready for execution\n");
printf("Execution Starting Address:%d\n",exaddr);
printf("ADDRESS\t DATA\n");
for(j=0;j< i;j++)
printf("%d\t%d\n",mem[j],data[j]);
gotoxy(1,q+4);
getch();
int search(char sym[10])
FILE *fp;
char s1[10],s2[10],s3[10],s4[10];
fp=fopen("ESTAB.TXT","r");
while(!feof(fp))
```

```
fscanf(fp,"%s\t%s\t%s\t%s\n",s1,s2,s3,s4);
if(strcmp(s2,sym)==0)
{
fclose(fp);
return(atoi(s3));
}
}
fclose(fp);
return 0;
}
```

d) Output:

Program Loaded in memory is ready for execution Execution Starting Address:9009 ADDRESS DATA 9000 6 9003 10006 9006 9 9009 6 9012 10006

e) Result:

9015 9

Thus the multi pass assemblers is created and operations are verified successfully.

VIVA – QUESTION AND ANSWER:

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3. What is the important machine structures used in the design of system software?

Memory structure Registers Data formats

> Instruction formats Addressing modes Instruction set

4. Define compiler and interpreter.

Compiler is a set of program which converts the whole high level language program to machine language program. Interpreter is a set of programs which converts high level language program to machine language program line by line.

CS2308 - SYSTEM SOFTWARE LAB

MODEL QUESTION SET

- 1. Write a C program to implement a symbol table with functions to create, insert, modify, search, and display.
- 2. Write a C program to implement pass one of a two pass assembler.
- 3. Write a C program to implement pass two of a two pass assembler.
- 4. Write a C program to implement a single pass assembler.
- 5. Write a C program to implement a two pass macro processor
- 6. Write a C program to implement a single pass macro processor.
- 7. Write a C program to implement an absolute loader.
- 8. Write a C program to implement a relocating loader.
- 9. Write a C program to implement pass one of a direct-linking loader.
- 10. Write a C program to implement pass two of a direct-linking loader.
- 11. Write a C++ program to implement a simple text editor with features like insertion / deletion of a character, word, and sentence.
- 12. Write a C program to implement a symbol table with suitable hashing