

21.12.(4)

假设该式不是重言式，则存在 u 和 φ , 使

$$V(\forall x(p \rightarrow q) \rightarrow (p \rightarrow (\forall xq))) = 0$$

$$\text{即} \begin{cases} V(\forall x(p \rightarrow q)) = 1 \\ V(p \rightarrow \forall xq) = 0 \end{cases}$$

$$\text{即} \begin{cases} V(\forall x(p \rightarrow q)) = 1 \\ V(p) = 1 \\ V(\forall xq) = 0 \end{cases}$$

则在 u 和 φ 下，取 $x' \notin X \cup C$

对 x' 的任一指派 $\varphi_0'(x') = \alpha \in U$ 有

$$V'(p(x')) \rightarrow q(x') = 1$$

而 $V(\forall xq)=0$ 表示存在对 x' 的某一指派 φ_0^*

使 $V'(q(x'))=0$, 而对该指派，同样有

$$V'(p(x')) \rightarrow q(x') = 1$$

$$\text{即 } 1 + V'(p(x')) = 1$$

$$\text{所以 } V'(p(x')) = 0$$

而 $V(p(x))=1$, 且由 $x \notin \text{var}(p)$

有 $p(x)=p(x')=p, V'(p(x'))=V(p)$

所以 $V(p)=0$ 与 $V(p)=1$ 矛盾

所以假设不成立

所以 $\vdash \forall x(p \rightarrow q) \rightarrow (p \rightarrow \forall xq)$, 这里 $x \notin \text{var}(p)$

21.12(5)

假设该式不是重言式，则存在 u 和 φ

$$\text{使 } V(\forall x(p(x)) \rightarrow p(t)) = 0$$

$$\text{即} \begin{cases} V(\forall xp(x)) = 1 \\ V(p(t)) = 0 \end{cases}$$

$V(\forall xp(x))=1$, 即在 u 和 φ 下，取 $x' \notin X \cup C$,

对 x' 的任一指派 $\varphi_0'(x') = \alpha \in U$ 有

$$V'(p(x'))=1$$

取 φ_0^* 使 $\varphi_0^*(x') = \varphi(t)$

则对 $p(x)$ 中所有包含自由出现的 x 的原子公式 $R_k^i(t_1, t_2, \dots, t_m, t_{m+1}, \dots, t_k)$

$$\text{其中 } t_m = f_{k_2}^j(\dots, x, \dots)$$

在 $p(x')$ 中被替换为 $t'_m = f_{k_2}^j(\dots, x, \dots)$

$$\text{则 } \varphi_0^*(t'_m) = \overline{f_{k_2}^j}(\dots, \varphi_0^*(x'), \dots)$$

$$= \overline{f_{k_2}^j}(\dots \varphi(t), \dots)$$

又 t 对 x 自由，所以 t 的所有个体变元在 $p(t)$ 中的所有项 $t''_m = f_{k_2}^j(\dots, t, \dots)$ 内是自由出现

则 $t_m'' = \overline{f_{k_2}}(\dots \varphi(t), \dots) = \varphi_0^*(t_m')$

所以 $v'(R_k^i(t_1', t_2' \dots t_k')) = v(R_k^i(t_1'', t_2'' \dots t_k''))$

并且 $P(t)$ 中所有来自替换的 t 的个体变元在替换处均为自由出现

所以 $v'(p(x')) = v(p(t))$ 这与 $v'(p(x')) = 1, v(p(t)) = 0$ 矛盾

所以 $\vdash \forall x(p(x)) \rightarrow p(t)$, 对 p 中的 x 是自由的

21.13(1)错误

取 $P(x) = R_2'(x, c)$ 取 u 和 φ 令 $U = \{1, 2, 3\}$, $\varphi_3(R_2')$ 为等于, $\varphi_1(c) = 2, \varphi_0(x) = 2$

则 $v(p(x)) = 1$

而取 $x' \notin X \cup C$, 令 $\varphi_0'(x') = 3$

则 $v'(p(x')) = 0$, 则 $v(\forall x(p(x))) = 0$

所以 $p(x) \Vdash \forall x(p(x))$ 并不正确

(2)错误

取与(1)相同的 $p(x)$ 及 u , φ 则 $v(p(x)) = 1, v(\forall x(p(x))) = 0$

所以 $v(p(x)) \rightarrow \forall x p(x) = 0$

得证

(3)错误

取 $p(x) = R_2'(x, c)$

取 u 及 φ 使 $U = \{1, 2, 3\}$, $\varphi_3(R_2')$ 为等于, $\varphi_1(c) = 2, \varphi_0(x) = 3$

则 $p(x) = 0$

而 $\exists x p(x) = \neg \forall x \neg p(x)$

$= \forall x \neg p(x) \rightarrow F$

取 $x' \notin X \cup C$, 令 $\varphi_0'(x') = 2$

则 $v'(p(x')) = 1, v'(\neg p(x')) = 0$

所以 $v(\forall x \neg p(x)) = 0$, 则 $v(\neg \forall x \neg p(x)) = 1$

则 $v(\exists x p(x)) = 1$

得证

(4)正确

假设存在 u 及 φ 使

$v((p \rightarrow \forall x q(x)) \rightarrow \forall x(p \rightarrow q)) = 0$

即 $\begin{cases} v(p \rightarrow \forall x q(x)) = 1 \\ v(\forall x(p \rightarrow q)) = 0 \end{cases}$

其中 $v(\forall x(p \rightarrow q)) = 0$ 即取 $x' \notin X \cup C$, 存在 x' 的某一指派 φ_0^*

使 $v'(p(x') \rightarrow q(x')) = 0$, 又由于 $x \notin \text{var}(p)$

所以此即 $v'(p \rightarrow q(x')) = 0$

则 $v(p) = 1, v'(q(x')) = 0$

即在该指派下, $v'(q(x')) = 0$, 则 $v(\forall x q(x)) = 0$

所以 $v(p \rightarrow \forall x q(x)) = 0$

这与 $v(p \rightarrow \forall x q(x)) = 1$ 矛盾

所以 $\vdash (p \rightarrow \forall x q(x)) \rightarrow \forall x(p \rightarrow q)$

(5)正确

假设存在 u 和 φ 使 $v((p \rightarrow \exists x q(x)) \rightarrow \exists x(p \rightarrow q)) = 0$

$$\text{即} \begin{cases} v(p \rightarrow \exists x(q(x))) = 1 \\ v(\exists x(p \rightarrow q)) = 0 \end{cases}$$

其中 $\exists x(p \rightarrow q) = \neg \forall x \neg(p \rightarrow q)$

则 $v(\forall x \neg(p \rightarrow q)) = 1$

即取 $x' \notin X \cup C$, 对 x' 的任意指派 φ'_0 有

$v'(\neg(p(x') \rightarrow q(x'))) = 1$

即 $v'(p(x') \rightarrow q(x')) = 0$

$$\text{即} \begin{cases} v'(p(x')) = 1 \\ v'(q(x')) = 0 \end{cases}$$

又 $x' \notin \text{var}(p)$, 则 $p(x') = p$ 所以 $v(p) = 1$

则由 $v(p \rightarrow \exists x(q(x))) = 1$

有 $v(\exists x q(x)) = 1$

即 $v(\forall x \neg q(x)) = 0$

即存在对 x' 的指派 φ_0^* 使 $v'(\neg q(x')) = 0$

而该指派下仍有 $v'(q(x')) = 0$

但 $v'(\neg q(x')) = 0$ 即 $v'(q(x')) = 1$ 矛盾

得证

(6)正确。假设存在解释域 u 和项解释 φ

使 $v(\forall x p(x) \rightarrow p(t)) = 0$

$\therefore v(\forall x p(x)) = 1 \quad v(p(t)) = 0$

\therefore 在 u 和 φ 下, 取 $x' \notin X \cup C$, 存在 x' 的任一指派 $\varphi_0(x')$

有 $v'(p(x')) = 1$

$\therefore \exists$ 指派 $\varphi_0(x') = \varphi(t)$

$\therefore v(p(t)) = v'(p(x')) = 1$

\therefore 矛盾 \therefore 命题得证

21.14

$\therefore \models p(x)$

\therefore 对于任意项解释 φ 的解释域 V , $V(p(x)) = 1$

\therefore 在任一 φ 和 U 下, 取 $x' \notin x \cup c$, 则对 x' 的任意指派 $\varphi_0(x')$ 有 $v'(p(x')) = 1$

$\therefore v(\forall x p(x)) = 1$

又由 21.13(1) 得 $v(\forall x p(x) \rightarrow p(t)) = 1$

$\therefore v(p(t)) = 1$

$\therefore \models p(t)$

21.18

(1) $p_1 = \forall y \forall x p(x, y)$

(A)

$p_2 = \forall y \forall x p(x, y) \rightarrow \forall x p(x, y)$

(A₅)

$p_3 = \forall x p(x, y)$

($p_2 = p_1 \rightarrow p_3$)

$$\begin{array}{ll}
p_4 = \forall x p(x, y) \rightarrow p(x, y) & (A_5) \\
p_5 = p(x, y) & (p_4 = p_3 \rightarrow p_5) \\
p_6 = \forall y p(x, y) & (p_5, y \notin \text{var}(A), G) \\
p_7 = \forall x \forall y p(x, y) & (p_6, x \notin \text{var}(A), G)
\end{array}$$

$$\begin{array}{ll}
(2) \quad p_1 = \forall x (p(x) \rightarrow q(x)) & \\
p_2 = \forall x (p(x) \rightarrow q(x)) \rightarrow (p(x) \rightarrow q(x)) & (A_5) \\
p_3 = p(x) \rightarrow q(x) & (Mp) \\
p_4 = (\forall x \rightarrow q(x) \rightarrow (\forall x \rightarrow p(x) \rightarrow F)) \\
\quad \rightarrow ((\forall x \rightarrow q(x) \rightarrow \forall x \rightarrow p(x)) \rightarrow (\forall x \rightarrow q(x) \rightarrow F)) & (A) \\
p_5 = (\forall x \rightarrow q(x) \rightarrow \forall x \rightarrow p(x)) \rightarrow (\forall x \rightarrow q(x) \rightarrow F) & (p_3 p_4 Mp) \\
p_6 = \forall x (p(x) \rightarrow q(x)) & (A) \\
p_7 = \forall x (p(x) \rightarrow q(x)) \rightarrow (p(x) \rightarrow q(x)) & (A_5) \\
p_8 = p(x) \rightarrow q(x) & (p_6 p_7 Mp) \\
p_9 = (p(x) \rightarrow q(x)) \rightarrow (q(x) \rightarrow \rightarrow p(x)) \\
p_{10} = \rightarrow q(x) \rightarrow \rightarrow p(x) & (p_8 p_9 Mp) \\
p_{11} = \forall x (\rightarrow q(x) \rightarrow \rightarrow p(x)) & (G) \\
p_{12} = \forall x (\rightarrow q(x) \rightarrow \rightarrow p(x)) \rightarrow (\forall x \rightarrow q(x) \rightarrow \forall x \rightarrow p(x)) \\
p_{13} = \forall x \rightarrow q(x) \rightarrow \forall x \rightarrow p(x) \\
p_{14} = \forall x \rightarrow q(x) \rightarrow F = \exists x q(x)
\end{array}$$

$$\begin{array}{ll}
(3) \quad A = \{\forall x (p(x) \rightarrow q(x)), \forall x p(x)\} & \\
p_1 = \forall x (p(x) \rightarrow q(x)) & \\
p_2 = \forall x (p(x) \rightarrow q(x)) \rightarrow (p(x) \rightarrow q(x)) & (A_5) \\
p_3 = p(x) \rightarrow q(x) & (Mp) \\
p_4 = \forall x p(x) & (A_5) \\
p_5 = \forall x p(x) \rightarrow p(x) & (A_5) \\
p_6 = p(x) & (p_4 p_5 Mp) \\
p_7 = q(x) & (p_6 p_3 Mp) \\
p_8 = \forall x q(x) & (G, x \notin \text{var}(\forall x (p(x) \rightarrow q(x)))) \cup \text{var}(\forall x p(x)) \\
\therefore \text{由演绎定理 } \vdash \forall x (p(x) \rightarrow q(x)) \rightarrow (\forall x p(x) \rightarrow \forall x q(x)) &
\end{array}$$

$$\begin{array}{ll}
(4) \quad p_1 = \forall x q \rightarrow q & (A_5) \\
p_2 = (\forall x q \rightarrow q) \rightarrow (p \rightarrow (\forall x q \rightarrow q)) & (A_1) \\
p_3 = p \rightarrow (\forall x q \rightarrow q) & (p_1, p_2, Mp) \\
p_4 = (p \rightarrow (\forall x q \rightarrow q)) \rightarrow ((p \rightarrow \forall x q) \rightarrow (p \rightarrow q)) & (A_2) \\
p_5 = (p \rightarrow \forall x q) \rightarrow (p \rightarrow q) & (p_3, p_4, Mp) \\
p_6 = p \rightarrow \forall x q & A \\
p_7 = p \rightarrow q & (p_6 p_5 Mp) \\
p_8 = \forall x (p \rightarrow q) \\
\text{由演绎定理 } \vdash (p \rightarrow \forall x q(x)) \rightarrow \forall x (p \rightarrow q) &
\end{array}$$

$$\begin{array}{ll}
(5) \quad A = \{p \rightarrow \exists x q\} & \\
p_1 = \neg \neg (p \rightarrow \exists x q) \rightarrow \exists x (p \rightarrow q) & \\
p_2 = p \rightarrow \exists x q & \\
p_3 = (p \rightarrow \exists x q) \rightarrow \neg \neg (p \rightarrow \exists x q) &
\end{array}$$

$$P_4 = \neg(p \rightarrow \exists xq)$$

$$P_5 = \exists x(p \rightarrow q)$$

$$A \vdash \exists x(p \rightarrow q)$$

$$\text{所以} \vdash (p \rightarrow \exists xq) \rightarrow \exists x(p \rightarrow q)$$

$$21.19 \text{ 正确 } \because A \cup \{p(x)\} \vdash q$$

$$\therefore \vdash p(x) \rightarrow q$$

$$A' = \{p \rightarrow q, \neg q \cup A$$

$$p_1 = (p \rightarrow (q \rightarrow F)) \rightarrow ((p \rightarrow q) \rightarrow (p \rightarrow F)) \quad A_2$$

$$P_2 = p \rightarrow q$$

$$P_3 = (q \rightarrow F) \rightarrow (p \rightarrow (q \rightarrow F))$$

$$P_4 = \neg q = q \rightarrow F$$

$$P_5 = p \rightarrow (q \rightarrow F)$$

$$P_6 = (p \rightarrow q) \rightarrow (p \rightarrow F)$$

$$P_7 = p \rightarrow F = \neg p$$

$$p_8 = \forall x \neg p$$

$$\therefore A \vdash (p \rightarrow q) \rightarrow (\neg q \rightarrow \forall x \neg p)$$

$$\therefore A \vdash (\neg q \rightarrow \forall x \neg p)$$

$$\therefore A \vdash (\neg \forall x \neg p \rightarrow \neg \neg q)$$

$$\therefore A \vdash \exists xq \rightarrow q \quad \therefore A \cup \{\exists xq\} \vdash q$$