

1. Completed in ex1.py
2. The definition of an $O(n \log n)$ complexity is that it divides the set of data in half then operates on the subsets of data which merge sort does do. In the merge_sort() function

```
def merge_sort(arr, low, high):  
    if low < high:  
        mid = (low + high)//2  
        merge_sort(arr, low, mid)  
        merge_sort(arr, mid+1, high)  
        merge(arr, low, mid, high)
```

We can see here that it recursively splits the array into subarrays until there are 1 element per subarray

```
merge_sort(arr, low, mid)  
merge_sort(arr, mid+1, high)
```

Then the merge() function is called and operates on those subarrays

```
n1 = mid - low + 1  
n2 = high - mid  
  
L = [0] * n1  
H = [0] * n2  
  
for i in range(n1):  
    L[i] = arr[low + i]  
for j in range(n2):  
    H[j] = arr[mid + 1 + j]
```

This will first copy the elements into left and right into the Low and High arrays respectively

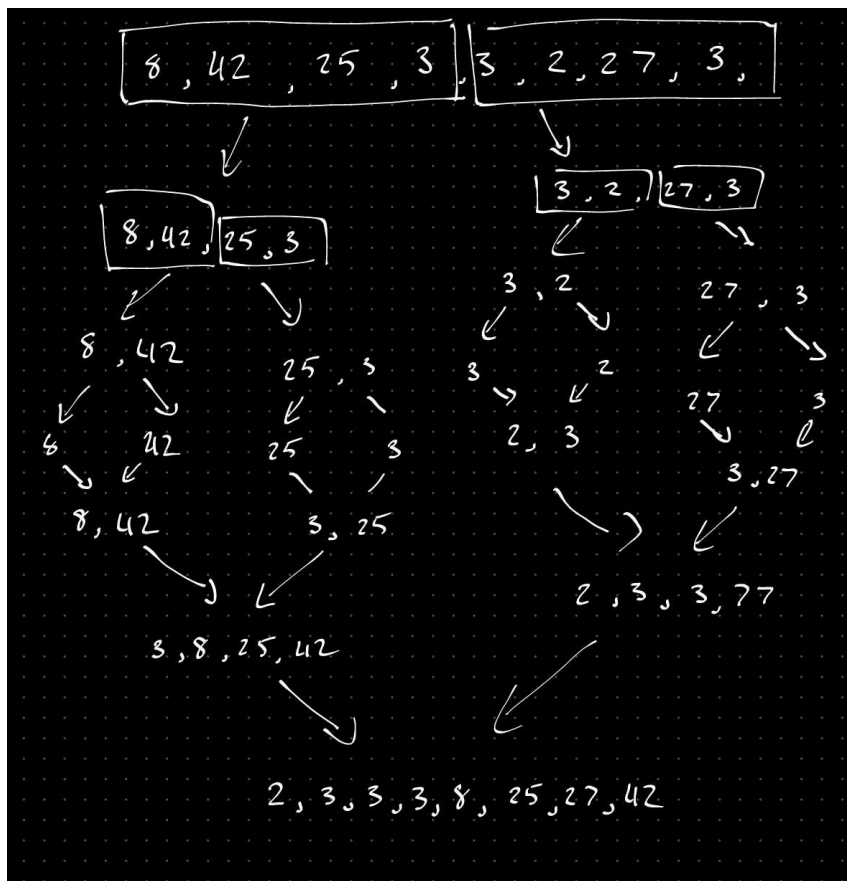
```
i = 0 # temp index  
j = 0 # temp index  
k = low # main index  
  
while i < n1 and j < n2:  
    if L[i] <= H[j]:  
        arr[k] = L[i]  
        i += 1  
  
    else:  
        arr[k] = H[j]  
        j += 1  
    k += 1
```

Then it will compare values in an $O(n)$ complexity and insert them back to the main array in a sorted manner

```
while i < n1:
    arr[k] = L[i]
    i += 1
    k += 1

while j < n2:
    arr[k] = H[j]
    j += 1
    k += 1
```

If any elements are left it will be taken care of by this segment of code



- 3.
4. Yes the number of steps follows the $O(n \log n)$ complexity. With each level of recursion the complexity is $O(\log n)$ and each time an element is processed and sorted it is $O(n)$ complexity