Q1:

```
1. Each recursive call is n-1 so each accurate can be showed as:

T(n) = T(n-1) + O(n)

We can firther say:

T(n-1) = T(n-2) + O(n-1)
T(n-2) = \dots + O(n-2) = \text{and so on}

So T(n) can be represented as the series:

T(n) = O(n) + O(n-1) + \dots + O(n)

and the sum of this series is

t(n) = O\left(\frac{n(n+1)}{2}\right) = \frac{n}{2}

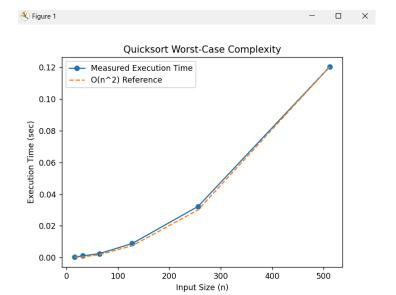
taken from math stack exchange
= O\left(\frac{n^2}{2}\right)
```

Q3:

This program runs on a number of differing size arrays which all incur a worst case scenario for quick sort(that being already sorted arrays).

Q4:

As you can see from the graph below as the number of inputs increase the time taken starts to go up quadratically which fits the previous formula of $O(n^*n)$ complexity



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