Subformula Linking

**Idea:** bring matching subformulas through switch rules

$$\begin{cases} \underline{A} \wedge B \otimes B \wedge (\underline{A} \vee C) \wedge D \\ \rightarrow B \wedge (\underline{A} \wedge B \otimes (\underline{A} \vee C) \wedge D) \\ \rightarrow B \wedge (\underline{A} \wedge B \otimes \underline{A} \vee C) \wedge D \\ \rightarrow B \wedge ((\underline{A} \wedge B \otimes \underline{A}) \vee C) \wedge D \\ \rightarrow B \wedge ((B \Rightarrow (\underline{A} \otimes \underline{A})) \vee C) \wedge D \\ \rightarrow B \wedge ((B \Rightarrow \top) \vee C) \wedge D \\ \end{cases}$$
 identity 
$$\{ \rightarrow B \wedge (T \vee C) \wedge D \\ \rightarrow B \wedge T \wedge D \\ \rightarrow B \wedge D \end{cases}$$

Variant of the **Calculus of Structures** (Guglielmi 1999)

 $\exists y. \ \forall x. \underline{R(x,y)} \otimes \forall a. \exists b. \underline{R(a,b)}$ 

• Unify linked subformulas

$$x \longmapsto a$$

$$\exists y. \ \forall x. \underline{R(x,y)} \otimes \forall a. \exists b. \underline{R(a,b)}$$
 
$$y \longleftarrow b$$

- Unify linked subformulas
- Check for ∀∃ dependency cycles

$$x \longmapsto a$$
 
$$\exists y. \ \forall x. \underline{R(x,y)} \otimes \forall a. \exists b. \underline{R(a,b)}$$
 
$$y \longleftarrow b$$

 $x \longmapsto a$ 

- Unify linked subformulas
- Check for ∀∃ dependency cycles
- **Switch** uninstantiated quantifiers

$$\exists y. \ \forall x. \underline{R(x,y)} \otimes \forall a. \exists b. \underline{R(a,b)}$$

$$\rightarrow \forall y. (\forall x. \underline{R(x,y)} \otimes y \longleftarrow b$$

$$\forall a. \ \exists b. \underline{R(a,b)}$$

$$\rightarrow \forall y. \forall a. (\forall x. \underline{R(x,y)} \otimes \exists b. \ \underline{R(a,b)})$$

- Unify linked subformulas
- Check for ∀∃ dependency cycles
- **Switch** uninstantiated quantifiers
- Instantiate unified variables

$$x \longmapsto a$$

$$\exists y. \ \forall x. \underline{R(x,y)} \otimes \forall a. \exists b. \underline{R(a,b)}$$

$$\rightarrow \forall y. (\forall x. \underline{R(x,y)} \otimes y \longleftarrow b$$

$$\forall a. \ \exists b. \underline{R(a,b)})$$

$$\rightarrow \forall y. \forall a. (\forall x. \underline{R(x,y)} \otimes \exists b. \underline{R(a,b)})$$

$$\rightarrow \forall y. \forall a. (\forall x. \underline{R(x,y)} \otimes \underline{R(x,y)} \otimes \underline{R(a,y)})$$

$$\rightarrow \forall y. \forall a. (\underline{R(a,y)} \otimes \underline{R(a,y)})$$

• Unify linked subformulas

$$\forall a. \exists b. \underline{R(a,b)} \otimes \exists y. \forall x. \underline{R(x,y)} \qquad x \longleftarrow a$$

$$y \longmapsto b$$

- Unify linked subformulas
- Check for ∀∃ dependency cycles

$$\forall a. \exists b. \underline{R(a,b)} \otimes \exists y. \forall x. \underline{R(x,y)} \left( \begin{array}{c} x & \longleftarrow & a \\ \\ y & \longleftarrow & b \end{array} \right)$$







- Quelques features manquantes, mais déjà utilisable
- Mathis a simplifié l'install -> n'hésitez pas à installer !
- Bientôt prêt pour évaluation sur des étudiants, voire experts

## **Bibliography**

- Chaudhuri, Kaustuv. 2013. "Subformula Linking as an Interaction Method". Edited by Sandrine Blazy, Christine Paulin-Mohring, David Pichardie, David Hutchison, Takeo Kanade, Josef Kittler, Jon M. Kleinberg, et al.. *Interactive Theorem Proving*. Berlin, Heidelberg: Springer Berlin Heidelberg. https://doi.org/10.1007/978-3-642-39634-2\_28
- Donato, Pablo, Pierre-Yves Strub, and Benjamin Werner. 2022. "A Drag-and-Drop Proof Tactic". In *Proceedings of the 11th ACM SIGPLAN International Conference on Certified Programs and Proofs*, 197–209. CPP 2022. Philadelphia, PA, USA: Association for Computing Machinery. https://doi.org/10.1145/3497775.3503692
- Guglielmi, Alessio. 1999. "A Calculus of Order and Interaction". https://www.researchgate.net/publication/2807151\_A\_Calculus\_of\_Order\_and\_Interaction