



# CS425 – Fall 2017

## Boris Glavic

### Chapter 6: Advanced SQL

modified from:

Database System Concepts, 6<sup>th</sup> Ed.

©Silberschatz, Korth and Sudarshan

See [www.db-book.com](http://www.db-book.com) for conditions on re-use



# Chapter 6: Advanced SQL

- Accessing SQL From a Programming Language
  - Dynamic SQL
    - ▶ JDBC and ODBC
  - Embedded SQL
- Functions and Procedural Constructs
- Triggers



**Textbook: Chapter 5**



# Accessing SQL From a Programming Language



# JDBC and ODBC

- API (application-program interface) for a program to interact with a database server
- Application makes calls to
  - Connect with the database server
  - Send SQL commands to the database server
  - Fetch tuples of result one-by-one into program variables
- ODBC (Open Database Connectivity) works with C, C++, C#, and Visual Basic
  - Other API's such as ADO.NET sit on top of ODBC
- JDBC (Java Database Connectivity) works with Java



# Native APIs

- Most DBMS also define DBMS specific APIs
  - Oracle: OCI
  - Postgres: libpg
- ...



# JDBC

- JDBC is a Java API for communicating with database systems supporting SQL.
- JDBC supports a variety of features for querying and updating data, and for retrieving query results.
- JDBC also supports metadata retrieval, such as querying about relations present in the database and the names and types of relation attributes.
- Model for communicating with the database:
  - Open a connection
  - Create a “statement” object
  - Execute queries using the Statement object to send queries and fetch results
  - Exception mechanism to handle errors



# JDBC Code

```
public static void JDBCexample(String dbid, String userid, String passwd)
{
    try {
        Class.forName ("oracle.jdbc.driver.OracleDriver"); // load driver
        Connection conn = DriverManager.getConnection( // connect to server
            "jdbc:oracle:thin:@db.yale.edu:2000:univdb", userid, passwd);
        Statement stmt = conn.createStatement(); // create Statement object
        ... Do Actual Work ....
        stmt.close(); // close Statement and release resources
        conn.close(); // close Connection and release resources
    }
    catch (SQLException sqle) {
        System.out.println("SQLException : " + sqle); // handle exceptions
    }
}
```



# JDBC Code (Cont.)

- Update to database

```
try {
    stmt.executeUpdate(
        "insert into instructor values(' 77987' , ' Kim' , ' Physics' ,
98000)");
} catch (SQLException sqle)
{
    System.out.println("Could not insert tuple. " + sqle);
}
```

- Execute query and fetch and print results

```
ResultSet rset = stmt.executeQuery(
    "select dept_name, avg (salary)
     from instructor
     group by dept_name");
while (rset.next()) {
    System.out.println(rset.getString("dept_name") + " " +
                       rset.getFloat(2));
}
```



# JDBC Code Details

- Result stores the current row position in the result
  - Pointing before the first row after executing the statement
  - `.next()` moves to the next tuple
    - ▶ Returns false if no more tuples
- Getting result fields:
  - `rs.getString("dept_name")` and `rs.getString(1)` equivalent if `dept_name` is the first attribute in select result.
- Dealing with Null values
  - `int a = rs.getInt("a");`  
`if (rs.wasNull()) System.out.println("Got null value");`



# Prepared Statement

- ```
PreparedStatement pStmt = conn.prepareStatement(
    "insert into instructor values(?, ?, ?, ?, ?)");
pStmt.setString(1, "88877");    pStmt.setString(2, "Perry");
pStmt.setString(3, "Finance");  pStmt.setInt(4, 125000);
pStmt.executeUpdate();
pStmt.setString(1, "88878");
pStmt.executeUpdate();
```
- For queries, use `pStmt.executeQuery()`, which returns a `ResultSet`
- WARNING: always use prepared statements when taking an input from the user and adding it to a query
  - NEVER create a query by concatenating strings which you get as inputs
  - "insert into instructor values(' " + ID + " ', ' " + name + " ', " +
 " ' + dept name + " ', " ' balance +
 ")"
  - What if name is “D’ Souza”?



# SQL Injection

- Suppose query is constructed using
  - "select \* from instructor where name = '" + name + "' "
- Suppose the user, instead of entering a name, enters:
  - X' or 'Y' = 'Y
- then the resulting statement becomes:
  - "select \* from instructor where name = '" + "X' or 'Y' = 'Y" + "' "
  - which is:
    - ▶ select \* from instructor where name = 'X' or 'Y' = 'Y'
  - User could have even used
    - ▶ X'; update instructor set salary = salary + 10000; --
- Prepared statement internally uses:  
"select \* from instructor where name = 'X\'' or '\''Y\'' = '\''Y'"
  - **Always use prepared statements, with user inputs as parameters**



# Metadata Features

- ResultSet metadata
- E.g., after executing query to get a ResultSet rs:
  - ```
ResultSetMetaData rsmd = rs.getMetaData();
for(int i = 1; i <= rsmd.getColumnCount(); i++) {
    System.out.println(rsmd.getColumnName(i));
    System.out.println(rsmd.getColumnTypeName(i));
}
```
- How is this useful?



# Metadata (Cont)

- Database metadata

```
■ DatabaseMetaData dbmd = conn.getMetaData();  
ResultSet rs = dbmd.getColumns(null, "univdb", "department", "%");  
// Arguments to getColumns: Catalog, Schema-pattern, Table-pattern,  
// and Column-Pattern  
// Returns: One row for each column; row has a number of attributes  
// such as COLUMN_NAME, TYPE_NAME  
while( rs.next()) {  
    System.out.println(rs.getString("COLUMN_NAME"),  
                      rs.getString("TYPE_NAME"));  
}  
}
```

- And where is this useful?



# Transaction Control in JDBC

- By default, each SQL statement is treated as a separate transaction that is committed automatically
  - bad idea for transactions with multiple updates
- Can turn off automatic commit on a connection
  - `conn.setAutoCommit(false);`
- Transactions must then be committed or rolled back explicitly
  - `conn.commit();`    or
  - `conn.rollback();`
- `conn.setAutoCommit(true)` turns on automatic commit.



# Other JDBC Features

## ■ Calling functions and procedures

- `CallableStatement cStmt1 = conn.prepareCall("{? = call some function(?)}");`
- `CallableStatement cStmt2 = conn.prepareCall("{call some procedure(?,?)?}");`

## ■ Handling large object types

- `getBlob()` and `getClob()` that are similar to the `getString()` method, but return objects of type Blob and Clob, respectively
- get data from these objects by `getBytes()`
- associate an open stream with Java Blob or Clob object to update large objects
  - ▶ `blob.setBlob(int parameterIndex, InputStream inputStream).`



# SQLJ

- JDBC is dynamic, errors cannot be caught by compiler
- SQLJ: embedded SQL in Java

```
● #sql iterator deptInfoIter ( String dept name, int avgSal);  
deptInfoIter iter = null;  
  
#sql iter = { select dept_name, avg(salary) from instructor  
              group by dept name };  
  
while (iter.next()) {  
    String deptName = iter.dept_name();  
    int avgSal = iter.avgSal();  
    System.out.println(deptName + " " + avgSal);  
}  
  
iter.close();
```



# ODBC

- Open DataBase Connectivity(ODBC) standard
  - standard for application program to communicate with a database server.
  - application program interface (API) to
    - ▶ open a connection with a database,
    - ▶ send queries and updates,
    - ▶ get back results.
- Applications such as GUI, spreadsheets, etc. can use ODBC
- Was defined originally for Basic and C, versions available for many languages.



# ODBC (Cont.)

- Each database system supporting ODBC provides a "driver" library that must be linked with the client program.
- When client program makes an ODBC API call, the code in the library communicates with the server to carry out the requested action, and fetch results.
- ODBC program first allocates an SQL environment, then a database connection handle.
- Opens database connection using SQLConnect(). Parameters for SQLConnect:
  - connection handle,
  - the server to which to connect
  - the user identifier,
  - password
- Must also specify types of arguments:
  - SQL\_NTS denotes previous argument is a null-terminated string.



# ODBC Code

```
■ int ODBCexample()
{
    RETCODE error;
    HENV env; /* environment */
    HDBC conn; /* database connection */
    SQLAllocEnv(&env);
    SQLAllocConnect(env, &conn);
    SQLConnect(conn, "db.yale.edu", SQL_NTS, "avi", SQL_NTS,
               "avipasswd", SQL_NTS);
    { .... Do actual work ... }

    SQLDisconnect(conn);
    SQLFreeConnect(conn);
    SQLFreeEnv(env);
}
```



# ODBC Code (Cont.)

- Program sends SQL commands to database by using SQLExecDirect
- Result tuples are fetched using SQLFetch()
- SQLBindCol() binds C language variables to attributes of the query result
  - When a tuple is fetched, its attribute values are automatically stored in corresponding C variables.
  - Arguments to SQLBindCol()
    - ▶ ODBC stmt variable, attribute position in query result
    - ▶ The type conversion from SQL to C.
    - ▶ The address of the variable.
    - ▶ For variable-length types like character arrays,
      - The maximum length of the variable
      - Location to store actual length when a tuple is fetched.
      - Note: A negative value returned for the length field indicates null value
  - Good programming requires checking results of every function call for errors; we have omitted most checks for brevity.



# ODBC Code (Cont.)

## ■ Main body of program

```
char deptname[80];
float salary;
int lenOut1, lenOut2;
HSTMT stmt;
char * sqlquery = "select dept_name, sum (salary)
                    from instructor
                    group by dept_name";
SQLAllocStmt(conn, &stmt);
error = SQLExecDirect(stmt, sqlquery, SQL_NTS);
if (error == SQL_SUCCESS) {
    SQLBindCol(stmt, 1, SQL_C_CHAR, deptname , 80, &lenOut1);
    SQLBindCol(stmt, 2, SQL_C_FLOAT, &salary, 0 , &lenOut2);
    while (SQLFetch(stmt) == SQL_SUCCESS) {
        printf ("%s %g\n", deptname, salary);
    }
}
SQLFreeStmt(stmt, SQL_DROP);
```



# ODBC Prepared Statements

## ■ Prepared Statement

- SQL statement prepared: compiled at the database
- Can have placeholders: E.g. insert into account values(?, ?, ?)
- Repeatedly executed with actual values for the placeholders

## ■ To prepare a statement

`SQLPrepare(stmt, <SQL String>);`

## ■ To bind parameters

`SQLBindParameter(stmt, <parameter#>,  
... type information and value omitted for simplicity..)`

## ■ To execute the statement

`retcode = SQLExecute( stmt);`

## ■ To avoid SQL injection security risk, do not create SQL strings directly using user input; instead use prepared statements to bind user inputs



# More ODBC Features

## ■ Metadata features

- finding all the relations in the database and
- finding the names and types of columns of a query result or a relation in the database.

## ■ By default, each SQL statement is treated as a separate transaction that is committed automatically.

- Can turn off automatic commit on a connection
  - ▶ `SQLSetConnectOption(conn, SQL_AUTOCOMMIT, 0)`
- Transactions must then be committed or rolled back explicitly by
  - ▶ `SQLTransact(conn, SQL_COMMIT)` or
  - ▶ `SQLTransact(conn, SQL_ROLLBACK)`



# ODBC Conformance Levels

- Conformance levels specify subsets of the functionality defined by the standard.
  - Core
  - Level 1 requires support for metadata querying
  - Level 2 requires ability to send and retrieve arrays of parameter values and more detailed catalog information.
- SQL Call Level Interface (CLI) standard similar to ODBC interface, but with some minor differences.



# ADO.NET

- API designed for Visual Basic .NET and C#, providing database access facilities similar to JDBC/ODBC

- Partial example of ADO.NET code in C#

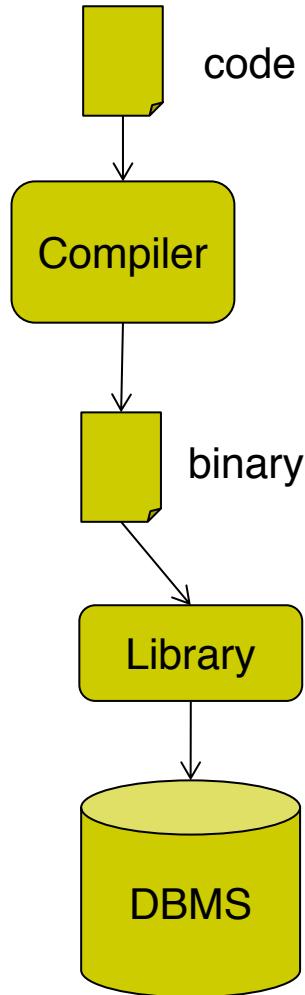
```
using System, System.Data, System.Data.SqlClient;
SqlConnection conn = new SqlConnection(
    "Data Source=<IPaddr>, Initial Catalog=<Catalog>");
conn.Open();
SqlCommand cmd = new SqlCommand("select * from students",
                                conn);
SqlDataReader rdr = cmd.ExecuteReader();
while(rdr.Read()) {
    Console.WriteLine(rdr[0], rdr[1]); /* Prints result attributes 1 & 2 */
}
rdr.Close(); conn.Close();
```

- Can also access non-relational data sources such as
  - OLE-DB, XML data, Entity framework

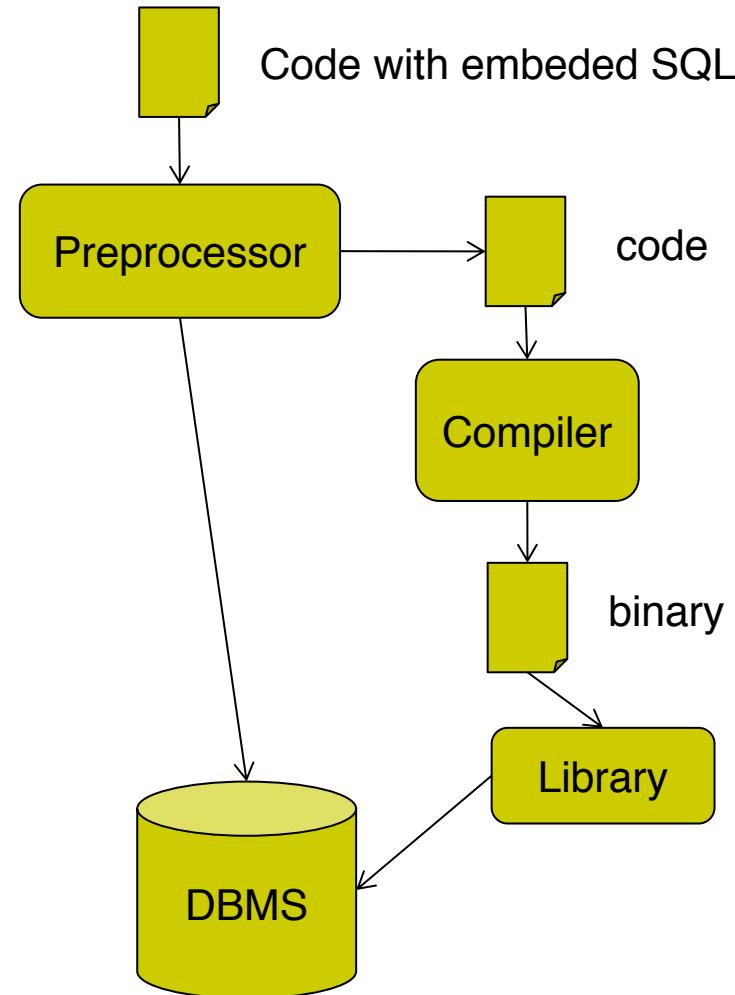


# Dynamic vs. Embedded SQL

## Dynamic SQL



## Embedded SQL





# Embedded SQL

- The SQL standard defines embeddings of SQL in a variety of programming languages such as C, Java, and Cobol.
- A language to which SQL queries are embedded is referred to as a **host language**, and the SQL structures permitted in the host language comprise *embedded* SQL.
- The basic form of these languages follows that of the System R embedding of SQL into PL/I.
- **EXEC SQL** statement is used to identify embedded SQL request to the preprocessor

```
EXEC SQL <embedded SQL statement> END_EXEC
```

Note: this varies by language (for example, the Java embedding uses # SQL { .... }; )



# Example Query

- From within a host language, find the ID and name of students who have completed more than the number of credits stored in variable `credit_amount`.
- Specify the query in SQL and declare a *cursor* for it

EXEC SQL

```
declare c cursor for
select ID, name
from student
where tot_cred > :credit_amount
```

END\_EXEC



# Embedded SQL (Cont.)

- The **open** statement causes the query to be evaluated

**EXEC SQL open c END\_EXEC**

- The **fetch** statement causes the values of one tuple in the query result to be placed on host language variables.

**EXEC SQL fetch c into :si, :sn END\_EXEC**

Repeated calls to **fetch** get successive tuples in the query result

- A variable called SQLSTATE in the SQL communication area (SQLCA) gets set to ‘02000’ to indicate no more data is available
- The **close** statement causes the database system to delete the temporary relation that holds the result of the query.

**EXEC SQL close c END\_EXEC**

Note: above details vary with language. For example, the Java embedding defines Java iterators to step through result tuples.



# Updates Through Cursors

- Can update tuples fetched by cursor by declaring that the cursor is for update

```
declare c cursor for
    select *
        from instructor
        where dept_name = 'Music'
    for update
```

- To update tuple at the current location of cursor *c*

```
update instructor
    set salary = salary + 100
    where current of c
```



# Procedural Constructs in SQL



# Procedural Extensions and Stored Procedures

- SQL provides a **module** language
  - Permits definition of procedures in SQL, with if-then-else statements, for and while loops, etc.
- Stored Procedures
  - Can store procedures in the database
  - then execute them using the **call** statement
  - permit external applications to operate on the database without knowing about internal details
- Object-oriented aspects of these features are covered in Chapter 22 (Object Based Databases) in the textbook



# Why have procedural extensions?

- Shipping data between a database server and application program (e.g., through network connection) is costly
- Converting data from the database internal format into a format understood by the application programming language is costly
- Example:
  - Use Java to retrieve all users and their friend-relationships from a friends relation representing a world-wide social network with 10,000,000 users
  - Compute the transitive closure
    - ▶ All pairs of users connects through a path of friend relationships. E.g., (Peter, Magret) if Peter is a friend of Walter who is a friend of Magret
  - Return pairs of users from Chicago – say 4000 pairs
  - 1) cannot be expressed (efficiently) as SQL query, 2) result is small
    - ▶ -> save by executing this on the DB server



# Functions and Procedures

- SQL:1999 supports functions and procedures
  - Functions/procedures can be written in SQL itself, or in an external programming language.
  - Functions are particularly useful with specialized data types such as images and geometric objects.
    - ▶ Example: functions to check if polygons overlap, or to compare images for similarity.
  - Some database systems support **table-valued functions**, which can return a relation as a result.
- SQL:1999 also supports a rich set of imperative constructs, including
  - Loops, if-then-else, assignment
- Many databases have proprietary procedural extensions to SQL that differ from SQL:1999.



# SQL Functions

- Define a function that, given the name of a department, returns the count of the number of instructors in that department.

```
create function dept_count (dept_name varchar(20))  
returns integer  
begin  
    declare d_count integer;  
    select count (*) into d_count  
    from instructor  
    where instructor.dept_name = dept_name;  
    return d_count;  
end
```

- Find the department name and budget of all departments with more than 12 instructors.

```
select dept_name, budget  
from department  
where dept_count (dept_name) > 1
```



# Table Functions

- SQL:2003 added functions that return a relation as a result
- Example: Return all accounts owned by a given customer

```
create function instructors_of(dept_name char(20)
```

```
returns table ( ID varchar(5),  
                    name varchar(20),  
                    dept_name varchar(20),  
                    salary numeric(8,2))
```

```
return table
```

```
(select ID, name, dept_name, salary  
from instructor  
where instructor.dept_name = instructors_of.dept_name)
```

- Usage

```
select *  
from table (instructors_of(‘Music’))
```



# SQL Procedures

- The *dept\_count* function could instead be written as procedure:

```
create procedure dept_count_proc (in dept_name varchar(20),
                                  out d_count integer)
begin
    select count(*) into d_count
    from instructor
    where instructor.dept_name = dept_count_proc.dept_name
end
```

- Procedures can be invoked either from an SQL procedure or from embedded SQL, using the **call** statement.

```
declare d_count integer;
call dept_count_proc( 'Physics' , d_count);
```

Procedures and functions can be invoked also from dynamic SQL

- SQL:1999 allows more than one function/procedure of the same name (called name **overloading**), as long as the number of arguments differ, or at least the types of the arguments differ



# Procedural Constructs

- Warning: most database systems implement their own variant of the standard syntax below
  - read your system manual to see what works on your system
- Compound statement: **begin ... end**,
  - May contain multiple SQL statements between **begin** and **end**.
  - Local variables can be declared within a compound statements
- **While** and **repeat** statements :

```
declare n integer default 0;  
while n < 10 do  
    set n = n + 1  
end while  
  
repeat  
    set n = n - 1  
until n = 0  
end repeat
```



# Procedural Constructs (Cont.)

## ■ For loop

- Permits iteration over all results of a query
- Example:

```
declare n integer default 0;  
for r as  
    select budget from department  
        where dept_name = 'Music'  
do  
    set n = n - r.budget  
end for
```



# Procedural Constructs (cont.)

- Conditional statements (**if-then-else**)  
SQL:1999 also supports a **case** statement similar to C case statement
- Example procedure: registers student after ensuring classroom capacity is not exceeded
  - Returns 0 on success and -1 if capacity is exceeded
  - See book for details
- Signaling of exception conditions, and declaring handlers for exceptions

```
declare out_of_classroom_seats condition  
declare exit handler for out_of_classroom_seats  
begin  
...  
.. signal out_of_classroom_seats  
end
```

- The handler here is **exit** -- causes enclosing **begin..end** to be exited
- Other actions possible on exception



# External Language Functions/Procedures

- SQL:1999 permits the use of functions and procedures written in other languages such as C or C++
- Declaring external language procedures and functions

```
create procedure dept_count_proc(in dept_name varchar(20),  
                                out count integer)
```

```
language C  
external name '/usr/avi/bin/dept_count_proc'
```

```
create function dept_count(dept_name varchar(20))  
returns integer  
language C  
external name '/usr/avi/bin/dept_count'
```



# External Language Routines (Cont.)

- Benefits of external language functions/procedures:
  - more efficient for many operations, and more expressive power.
- Drawbacks
  - Code to implement function may need to be loaded into database system and executed in the database system's address space.
    - ▶ risk of accidental corruption of database structures
    - ▶ security risk, allowing users access to unauthorized data
  - There are alternatives, which give good security at the cost of potentially worse performance.
  - Direct execution in the database system's space is used when efficiency is more important than security.



# Security with External Language Routines

- To deal with security problems
  - Use **sandbox** techniques
    - ▶ E.g., use a safe language like Java, which cannot be used to access/damage other parts of the database code.
  - Or, run external language functions/procedures in a separate process, with no access to the database process' memory.
    - ▶ Parameters and results communicated via inter-process communication
- Both have performance overheads
- Many database systems support both above approaches as well as direct executing in database system address space.



# Triggers



# Triggers

- A **trigger** is a statement that is executed automatically by the system as a side effect of a modification to the database.
- To design a trigger mechanism, we must:
  - Specify the conditions under which the trigger is to be executed.
  - Specify the actions to be taken when the trigger executes.
- Triggers introduced to SQL standard in SQL:1999, but supported even earlier using non-standard syntax by most databases.
  - Syntax illustrated here may not work exactly on your database system; check the system manuals



# Trigger Example

- E.g. *time\_slot\_id* is not a primary key of *timeslot*, so we cannot create a foreign key constraint from *section* to *timeslot*.
- Alternative: use triggers on *section* and *timeslot* to enforce integrity constraints

```
create trigger timeslot_check1 after insert on section
referencing new row as nrow
for each row
when (nrow.time_slot_id not in (
    select time_slot_id
    from time_slot) /* time_slot_id not present in time_slot */
begin
    rollback
end;
```



# Trigger Example Cont.

```
create trigger timeslot_check2 after delete on timeslot
referencing old row as orow
for each row
when (orow.time_slot_id not in (
    select time_slot_id
    from time_slot)
/* last tuple for time slot id deleted from time slot */
and orow.time_slot_id in (
    select time_slot_id
    from section)) /* and time_slot_id still referenced from section*/
begin
    rollback
end;
```



# Triggering Events and Actions in SQL

- Triggering event can be **insert**, **delete** or **update**
- Triggers on update can be restricted to specific attributes
  - E.g., **after update of takes on grade**
- Values of attributes before and after an update can be referenced
  - **referencing old row as** : for deletes and updates
  - **referencing new row as** : for inserts and updates
- Triggers can be activated before an event, which can serve as extra constraints. E.g. convert blank grades to null.

```
create trigger setnull_trigger before update of takes  
referencing new row as nrow  
for each row  
when (nrow.grade = ' ')  
begin atomic  
    set nrow.grade = null;  
end;
```



# Trigger to Maintain credits\_earned value

- **create trigger** *credits\_earned* **after update** of *takes* **on** (*grade*)  
**referencing new row as** *nrow*  
**referencing old row as** *orow*  
**for each row**  
**when** *nrow.grade*  $\neq$  'F' **and** *nrow.grade* **is not null**  
    **and** (*orow.grade* = 'F' **or** *orow.grade* **is null**)  
**begin atomic**  
    **update** *student*  
    **set** *tot\_cred*= *tot\_cred* +  
        (**select** *credits*  
         **from** *course*  
         **where** *course.course\_id*= *nrow.course\_id*)  
    **where** *student.id* = *nrow.id*;  
**end;**



# Statement Level Triggers

- Instead of executing a separate action for each affected row, a single action can be executed for all rows affected by a transaction
  - Use **for each statement** instead of **for each row**
  - Use **referencing old table** or **referencing new table** to refer to temporary tables (called ***transition tables***) containing the affected rows
  - Can be more efficient when dealing with SQL statements that update a large number of rows



# When Not To Use Triggers

- Triggers were used earlier for tasks such as
  - maintaining summary data (e.g., total salary of each department)
  - Replicating databases by recording changes to special relations (called **change** or **delta** relations) and having a separate process that applies the changes over to a replica
- There are better ways of doing these now:
  - Databases today provide built in materialized view facilities to maintain summary data
  - Databases provide built-in support for replication
- Encapsulation facilities can be used instead of triggers in many cases
  - Define methods to update fields
  - Carry out actions as part of the update methods instead of through a trigger



# When Not To Use Triggers

- Risk of unintended execution of triggers, for example, when
  - loading data from a backup copy
  - replicating updates at a remote site
  - Trigger execution can be disabled before such actions.
- Other risks with triggers:
  - Error leading to failure of critical transactions that set off the trigger
  - Cascading execution



# Recursive Queries



# Recursion in SQL

- SQL:1999 permits recursive view definition
- Example: find which courses are a prerequisite, whether directly or indirectly, for a specific course

```
with recursive rec_prereq(course_id, prereq_id) as (
    select course_id, prereq_id
    from prereq
    union
    select rec_prereq.course_id, prereq.prereq_id,
    from rec_rereq, prereq
    where rec_prereq.prereq_id = prereq.course_id
)
select *
from rec_prereq;
```

This example view, *rec\_prereq*, is called the *transitive closure* of the *prereq* relation



# Recursion in SQL - Syntax

- General form

```
with recursive R as (
    init_query
    union
    recursive_step)
select *
from R;
```

- init\_query returns the initial content of R
- recursive\_step is a query that mentions R exactly once in the FROM clause



# Recursion in SQL - Semantics

## ■ General form

```
with recursive R as (
    init_query
    union
    recursive_step)
select *
from R;
```

## ■ Fixpoint computation

- $R_0$  = result of init\_query
- In step  $i$ :  $R_i$  is computed as
  - ▶  $R_{i-1}$  union recursive\_step( $R_{i-1}$ )
- The computation stops when recursive\_step( $R_{i-1}$ ) is the empty set, i.e.,  $R_{i-1} = R_i$



# The Power of Recursion

- Recursive views make it possible to write queries, such as transitive closure queries, that cannot be written without recursion or iteration.
  - Intuition: Without recursion, a non-recursive non-iterative program can perform only a fixed number of joins of *prereq* with itself
    - ▶ This can give only a fixed number of levels of managers
    - ▶ Given a fixed non-recursive query, we can construct a database with a greater number of levels of prerequisites on which the query will not work
    - ▶ Alternative: write a procedure to iterate as many times as required
      - See procedure *findAllPrereqs* in book



# The Power of Recursion

- Computing transitive closure using iteration, adding successive tuples to *rec\_prereq*
  - The next slide shows a *prereq* relation
  - Each step of the iterative process constructs an extended version of *rec\_prereq* from its recursive definition.
  - The final result is called the *fixed point* of the recursive view definition.
- Recursive views are **monotonic**. That is,
  - if we add tuples to *prereq* the view *rec\_prereq* contains all of the tuples it contained before, plus possibly more



# Example of Fixed-Point Computation

course_id	prereq_id
BIO-301	BIO-101
BIO-399	BIO-101
CS-190	CS-101
CS-315	CS-101
CS-319	CS-101
CS-347	CS-101
EE-181	PHY-101

Iteration Number	Tuples in cl
0	
1	(CS-301)
2	(CS-301), (CS-201)
3	(CS-301), (CS-201)
4	(CS-301), (CS-201), (CS-101)
5	(CS-301), (CS-201), (CS-101)



# Another Recursion Example

- Given relation

*manager(employee\_name, manager\_name)*

- Find all employee-manager pairs, where the employee reports to the manager directly or indirectly (that is manager's manager, manager's manager's manager, etc.)

```
with recursive empl (employee_name, manager_name) as (
    select employee_name, manager_name
    from manager
    union
    select manager.employee_name, empl.manager_name
    from manager, empl
    where manager.manager_name = empl.employee_name)
select *
from empl
```

This example view, *empl*, is the *transitive closure* of the *manager* relation



# Recap

- Programming Language Interfaces for Databases
  - Dynamic SQL (e.g., JDBC, ODBC)
  - Embedded SQL
  - SQL Injection
- Procedural Extensions of SQL
  - Functions and Procedures
- External Functions/Procedures
  - Written in programming language (e.g., C)
- Triggers
  - Events (insert, ...)
  - Conditions (WHEN)
  - per statement / per row
  - Accessing old/new table/row versions
- Recursive Queries



# End of Chapter

modified from:

**Database System Concepts, 6<sup>th</sup> Ed.**

©Silberschatz, Korth and Sudarshan

See [www.db-book.com](http://www.db-book.com) for conditions on re-use



# Outline

- Introduction
- Relational Data Model
- Formal Relational Languages (relational algebra)
- SQL - Advanced
- **Database Design – ER model**
- Transaction Processing, Recovery, and Concurrency Control
- Storage and File Structures
- Indexing and Hashing
- Query Processing and Optimization