

## (a) Aggregate Method:

=> In the aggregate method, we analyze the total cost for a sequence of operations and then divide it by the no. of operations to get the average cost (amortized cost) per operation.

1. Inserting an element without resizing: For each insertion where there is enough space in the table, the cost is constant,  $O(1)$ .
2. Resizing: Every time the table doubles in size, the cost is proportional to the no. of elements being copied to the new table. If the table size is  $k$  before doubling, copying all  $k$  elements takes  $O(k)$  time.

Total cost of  $n$  Insertions:

- => The cost of inserting  $n$  elements:  $O(n)$  for the insertions.
- => The cost of resizing: The first doubling involves copying 1 element, the next involves copying 2, then 4 and so on.
- => The sum of these doubling operations is:  

$$O(1+2+4+8+\dots+2^k) \sim O(n)$$

=> Amortized cost per Insertion:  $\frac{O(n)+O(n)}{n} = O(1)$

So, using aggregate method, the amortized time complexity for inserting  $n$  elements is  $O(1)$ .

## (b) Accounting Method:

=> The accounting method assigns "credits" to each operation to account for the costs of future expensive operations.

### 1. Assigning Credits:

Each insertion will be charged 3 credits

=> 2 credits for the insertion itself, which pays constant time  $O(1)$  operation.

=> 1 credit to help pay for the cost of future resizing operations.

### 2. Cost of Insertion:

=> When no resizing happens, the cost is exactly 1 credit for the insertion.



$\Rightarrow$  When a resizing happens, it costs  $O(c)$  for copying  $c$  elements, but since we have 1 credit saved for each previous insertion, we have enough credits to cover the resizing.

### Resizing & cost:

$\Rightarrow$  When the table doubles the cost of copying elements doubles as well.

$\Rightarrow$  The total no. of credits that we collect is 3 credits per insertion, for  $n$  insertions, resulting in  $3n$  credits.

$\Rightarrow$  Each resizing is covered by the saved credits. The total no. of resizing operations is proportional to the no. of doublings (about  $\log n$  times).

### Final Amortized cost:

- Inserting  $n$  elements costs 1 credit each.
- Total credits collected  $= 3n$
- Cost of each resizing operation is already covered by saved credits.

Amortized cost per Insertion:

$$= \frac{3n}{n}$$

$$= O(1)$$

$\therefore O(1)$  is the amortized time complexity for inserting  $n$  elements using accounting method for a dynamic table that doubles in size.