Quest Explain Process states with transition diagram in alwall.

De A process is a program in execution and it is more than a program cool called of text section and this concept works under all the operating system because all the task perform by the operating system needs a process to perform the task.

The process executes when 4 changes the state. The state of a process is defined by the awarent activity of the process. tack process may be in any one of the following states -

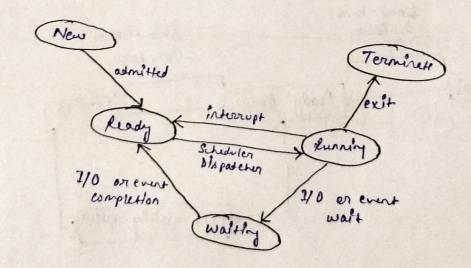
· New: The process is being created.

· Running: In this state the instructions are being executed.

· waiting: The process is in weiting state until an event occurs like I/o operation completion or receiving a signal.

· heady: The process is weithy to be assigned to a processer.

· Terminated: The process has finished execution. Dayran ->



Over 2. Explain various Schedulers in detail.

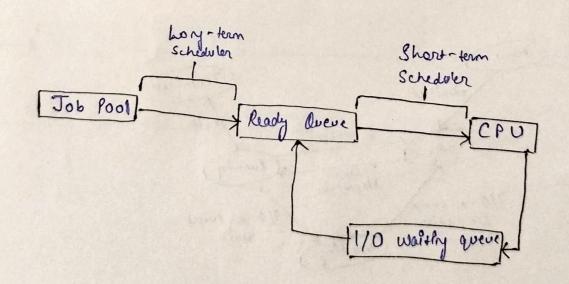
As There are three types of process schedulers:

i) Lary Term on Job Scheduler

It britis the new process to the 'Ready State'. It controls the object of Mustin programming, i.e. the number of processes present in a ready state at any point in three. It is important that the boy-term scheduler make a careful selection of both 1/0 and CPU bound processes. The job scheduler increases efficiency by maintaining a balance between the two.

ii) Shord-Term on CPU Scheduler

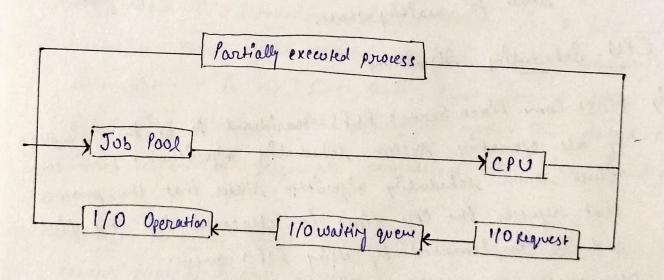
It is responsible for scheduling one process from the ready state for scheduling it on the running state. Short ferm scheduler only scheduler the process to schedule it doesn't load the process on running. The CPU scheduler is responsible for ensuring no starvation due to high burst time processes.



The dispatcher is responsible for loadly the process relieved by the Short-term schedules on the CPU (heady to Running state) Consext Switching is done by the allegatcher only.

ii) Medium - Term Scheduler

It is responsible for suspending and resuming the process. It mainly aloes swapping (moving processes from main memory to disk and vice-verse). Iwapping may be necessary to improve the process mix or because a charge in memory sugularments has overcommitted anallable memory, regularly memory to be freed up. It is helpful in maintaining a perfect balance between the 1/0 bound and the course sound. It reduces the degree of multiprogramming.



- Ques 3. what is a CPU Schoduling? Explain narlow Scheduling algorithms in oletail.
- Down the CPU to use while another process he suspended. The main function of the CPU schooling is to ensure that whenever the CPU remains idle, the OS has at least selected one of the processes available in the ready-to-use Un.

There are mainly two types of scheduling methods:

- · Breenptine Schooluling: Preenptive scheduling is used when a process switches from running state to ready state or from the westing state to the ready state.
- · Non-Breenpffin Scheduly: Non-breempfine scheduling is used when a process ferminates, or when a process switches from turning state to waiting state;

CPU Scheduling Algorithms

i) first Com Plast Serve: fcfs considered to be the simplest of all operating system scheduling algorithms. First come first serve scheduling algorithm states that the process that requests the CPU florst to allocated the CPU florst and is implemented by using fifo quewe.

Chara Charistics

- · FCFS supports non-preemptine and preemptine CPU Scheduling ago. · fCFS is easy to implement and use.
- · This algorithm is not much efficient in performence and the wait time its quite high.

- ii) Shortest Job first: SJf is a scheduling process that selects the waiting process with the smallest execution time to execute next. This scheduling method may on may not be preemptime. Significantly reduce the average waiting time for other processes waiting to be executed.
- Scheduling Algerithm that works based on the priority of a process. In this algerithm, the colitor sets the functions to be as important, meaning that the most important process must be done first. In the case of any conflict, that he, where there is more than one process with equal value, then the most important clu planning also. works on the basis of the fifs algorithm.
- Round Robin CPU Scheduling: Round Robin ils a CPU scheduling mechanism those eyelis around assigning each task a specific time slot. It is the first Come, flost Served CPU Scheduling technique with preemptine mode. The Round Robin CPU algorithm frequently emphasizes the Time-Sharing Method,
- Shortest Remainly Tim: This is similar to shortest Job First, except that if a new program is started, the OS compares the time it needs with the time the currently running program has left. If the new program would finish sooner, then the currently running program is switched out and the epu starts processly the new program.

Quesa. Explain Threads and their management in detall.

And thread is a shyle sequence struct within a process. Threads are also called lightweight processes as they possess some of the properties of processes.

Threads own in parallel improving the application performance, Each such thread has its own CPU state and stack, but they share the address space of the process and the environment.

Types of Thread in Os

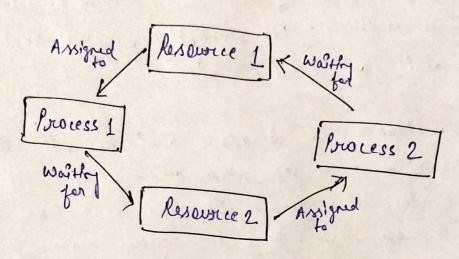
- · User Level Thread
- · Kernel Level Thread

User Level Threads: 9t is a type of thread that is not created using system calls. The kirnel has no work in the management of user-level threads. User-level threads can be easily implemented by the user. In case when user-level threads are single-handed processes, kernel-level thread manages them

Kernel Level Threads: A kernel level thread Is a type of thread that can recognize the Operating system easily. Kernel Level Threads has its own thread table where it keeps treak of the system. The Operating System Kernel helps in managing threads. Kernel Threads have somehow longer context switching time. Kernel helps in the management of threads.

Over 5. Explain Deadlock and narrious methods to handle deadlock h detall.

Aus A deadlock is a situation where a set of processes are blocked because each process is holdly a resource and waiting for another suscerce acquired by some other process.



Deadlock can wrise if the following four conditions hold simultaneously (necessary conditions)

- · Mutual Exclusion; Tous or more resources are non-shareable. (Only one process can use at a thre).
- · field and wast: A process is holding at least one resource and waiting for resources.
- · No freemption: A resource cannot be taken from a process unless the process releases the resources.
- · Chrowlar wait: A set of processes waiting for each other
- In circular form.

 If any one of the above, aloes not holds, then deadlock can be prevented.

- Deadlock Ignorance: If a deadlock is very now, then let is topped happen and restart the system. This is the approach that both Windows and UNIX take, we use the ostruck algorithm for deadlock ignorance.
- ii) Deadlock prevention: The deadlock can be prevented if any one of the necessary conditions of deadlock is violated, i.e. violate one of these conditions: No preemption, mutual exclusion, hold and wast on circular wast.
- Deadlock avoidance: In deadlock avoidance, the Os chicks whether the system is in safe state or in unsafe state at every step which the Os performs. The process continues until the system is in safe state. Once the system moves to unsafe state, the Os has to backtrack one step.

In simple words, The OS reviews each allocation to that the allocation doesn't cause the deadlock in the system.

fall he deadlock and then periodically check whether deadlock occurs in the system on not, if it occurs then it applies some of the recovery methods to the system to get vild of deadlock.