

Stack Organisation

- * Stack is a storage ~~data~~ structure that stores information in such a way that the last item stored is the first item retrieved.
- * It is based on the principle of LIFO (Last-In-First-Out).
- * The stack in digital computers is a group of memory locations with a register that holds the address of top of elements.
- * This register that holds the address of top of element of the stack is called stack pointer.
- * In stack organisation, ALU operations are performed on stack data i.e. both the operands are always required on the stack.
- * After manipulation, the result is placed in the stack.

Stack operations

The main two operations that are performed on the operators of the stack are:-

- * Push:- Insert an item on top of stack.
- * Pop:- Delete an item from top of stack.

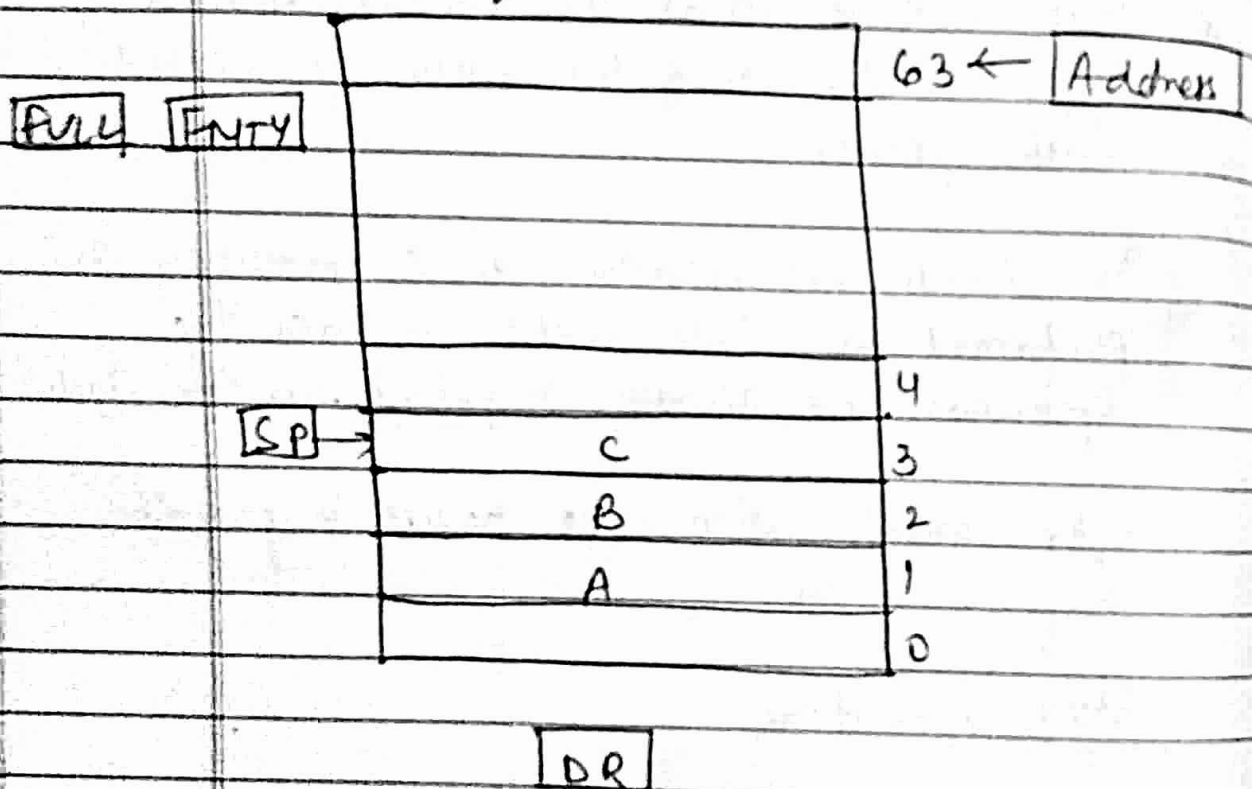
Implementation of stack

In digital computers, stack can be implemented in two ways:-

- * Register stack
- * Memory stack

① Register Stack

- * A stack can be organised as a collection of finite no. of registers that are used to store temporary information during the execution of a program.



Block Diagram of 64 word stack

- A Figure shows the 64-word register arrangement of the stack.
- * The stack pointer register holds the address of the element present at the top of the stack.
 - * Three elements A, B and C are located in the stack.
 - * Element C is at the top of the stack and SP holds the address of C i.e. 3.
 - * The top element is popped from the stack through reading memory word at address 3 and decrementing the SP by 1.
 - * Then, B is at the top of the stack and the SP holds the address of B that is 2.
 - * It can insert a new word, the stack is pushed by incrementing the stack pointer by 1 and inserting a word in that incremented location.

The PUSH operation is executed as follows:-

- * Insert an item on top of stack.

$SP \leftarrow SP + 1$

It can increment SP

$M[SP] \leftarrow DR$

It can write element

$B[SP=0] \text{ then } (FULL \leftarrow 1)$

Check if stack is full.

$EMPTY \leftarrow 0$

Mark the stack not empty

- * The SP, includes 6 bits, because $2^6 = 64$.

- * SP cannot exceed (111111 in binary).

- * After all, if 63 is incremented by 1, therefore the result is $0(111111 + 1 = 1000000)$. SP holds only the six least significant bits. 000000 is decremented by 1 thus the

Result is 111111.

- * Therefore, when the stack is full, the one-bit register 'Full' is set to 1. If the stack is NULL, then the one-bit register 'EMPTY' is set to 1.
- * The data register DR holds the binary information which is composed into readout of the stack.
- * First, the SP is set to 0, EMPTY is set to 1, and Full is set to 0.
- * Now, as the stack is not full (Full = 0), a new element is inserted using the push operation.

The POP operation is executed as follows:-

- * Delete ~~the~~ an item from top of stack.

$DR \leftarrow M[SP] \Rightarrow$ It can read an element from the top of the stack.

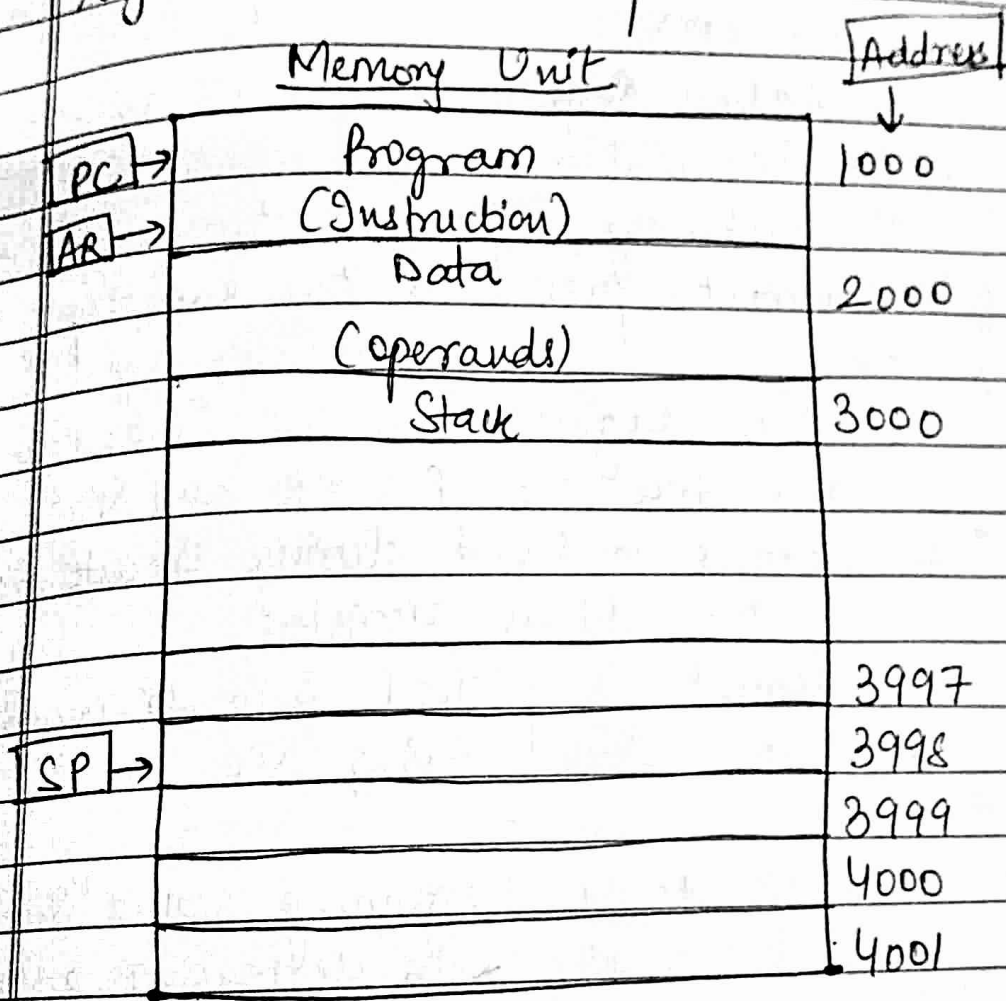
$SP \leftarrow SP - 1 \Rightarrow$ It can decrement the stack pointer

If $(SP = 0)$ then $(EMPTY \leftarrow 1) \Rightarrow$ Check if stack is empty.

$FULL \leftarrow 0 \Rightarrow$ Mark the stack not full.

(2) Memory Stack

- * A stack can exist as a standalone unit.
- * It can be implemented in a random access memory (RAM) attached to a CPU.
- * The implementation of a stack in the CPU is done by assigning a portion of memory to a stack operation and using a processor register as a stack pointer.



DR

Computer Memory with Program, data and Stack segments

An area of the computer memory is broken into three segments such as:-

- * Program
- * Data
- * Stack

* The address of the next instruction in the program is saved in the pointer Program Counter (PC)

* The Address Register (AR) points to an array of the information.

* SP continually influences the address of the element present at the ~~end~~ top of the stack.

* The three registers that are linked to the common bus are PC, AR and SP.

* An operand is read during execute stage using the address register.

* An element is pushed into or popped from the stack using SP.

* SP points to a beginning value '4001' and increases with decreasing addresses.

* The first element is saved at address 4000, the next element is saved at address 3999 and the last element is saved at address 1000.

Push operation

- * The data register can read an element into or from the stack. It can use push operation to insert a new element into the stack.

$$SP \leftarrow SP - 1$$

$$M[SP] \leftarrow DR.$$

Pop operation

- * It can delete an element from the stack. It can use the pop operation which is as follows:-

$$DR \leftarrow M[SP]$$

$$SP \leftarrow SP + 1$$

- * Two processor registers are used to check the stack limits.
- * One processor register for the upper limit (1000) and one for the lower limit (4001).
- * During push operation, the SP is compared with the upper limit to check if the stack is full.
- * During pop operation, the SP is compared with the lower limit to check if the stack is empty.