

A conjecture regarding optimality of the dictator function under Hellinger distance

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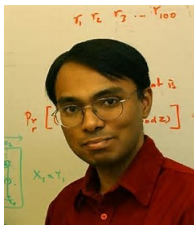
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Thanks: Simon's institute

Introduction

Starting point: the following conjecture¹ by Kumar ('12)

X: uniform on $\{-1, +1\}^n$

Y: obtained from **X** via the standard *noise-operator*, i.e.

- flip each bit (independently) with probability $\frac{1-\rho}{2}$.

Conjecture-MI

The dictator function $f_d(\mathbf{X}) = X_1$ maximizes the mutual information $I(f(\mathbf{X}); \mathbf{Y})$ among all boolean functions $f(\mathbf{X})$.

¹Thomas A Courtade and Gowtham R Kumar. “Which Boolean functions maximize mutual information on noisy inputs?” In: *IEEE Transactions on Information Theory* 60.8 (2014),



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Alternate view

Let $\Phi_{JS}(x) := 1 - H_b(x) = JS[(x, 1-x), (1-x, x)]$.

Given $f(\mathbf{X}) : \{-1, +1\}^n \mapsto \{-1, +1\}$, let $Z_f(\mathbf{Y}) := \frac{1 - (T_\rho f)(\mathbf{Y})}{2}$ where $(T_\rho f)(\mathbf{Y}) = E(f(\mathbf{X}) | \mathbf{Y})$.

Conjecture-MI (restatement)

The dictator function $f_d(\mathbf{X}) = X_1$ maximizes the Φ_{JS} -entropy, $E(\Phi_{JS}(Z_f(\mathbf{Y}))) - \Phi_{JS}(E(Z_f(\mathbf{Y})))$, among all boolean functions $f(\mathbf{X})$.

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Main Conjecture

Idea: Replace $\Phi_{JS}(x)$ by other convex functions.

Consider squared Hellinger distance between $(x, 1 - x)$ and $(1 - x, x)$

$$\Phi_{\mathcal{H}^2}(x) := 1 - 2\sqrt{x(1 - x)}.$$

As before, given $f(\mathbf{X}) : \{-1, +1\}^n \mapsto \{-1, +1\}$, let $Z_f(\mathbf{Y}) := \frac{1 - (T_\rho f)(\mathbf{Y})}{2}$.

Conjecture-SH

The dictator function $f_d(\mathbf{X}) = X_1$ maximizes the $\Phi_{\mathcal{H}^2}$ -entropy, $\mathbb{E}(\Phi_{\mathcal{H}^2}(Z_f(\mathbf{Y}))) - \Phi_{\mathcal{H}^2}(\mathbb{E}(Z_f(\mathbf{Y})))$, among all boolean functions $f(\mathbf{X})$.

Equivalently

$$\sqrt{1 - \mathbb{E}(f)^2} - \mathbb{E} \left(\sqrt{1 - (T_\rho f)^2(\mathbf{Y})} \right) \leq 1 - \sqrt{1 - \rho^2}.$$



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- ❶ Why is this interesting? (or, why should one care about this Φ ?)
- ❷ Evidence to the veracity of the conjecture
- ❸ Weaker forms



About the Hellinger conjecture

In short, two lemmas:

- 1 Conjecture-SH implies Conjecture-MI
- 2 Conjecture-SH is extremal



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Proposition

Conjecture-SH implies Conjecture-MI

Lemma

The function $H_b\left(\frac{1-\sqrt{1-x^2}}{2}\right)$ is non-negative, increasing, and convex in x for $x \in [0, 1]$.

Let $\Psi(x) = \frac{1-\sqrt{1-x^2}}{2}$. Observe that

$$H\left(\frac{1-x}{2}\right) = H\left(\Psi\left(\sqrt{1-x^2}\right)\right).$$

Conjecture-MI can be expressed as

$$H\left(\Psi\left(\sqrt{1-E(f)^2}\right)\right) - E\left(H\left(\Psi\left(\sqrt{1-(T_\rho f)^2(\mathbf{Y})}\right)\right)\right) \leq H(\Psi(1)) - H\left(\Psi\left(\sqrt{1-\rho^2}\right)\right).$$



Idea of proof

By the convexity of $H(\Psi(x))$ (lemma) suffices to show

$$H\left(\Psi\left(\sqrt{1 - \mathbb{E}(f)^2}\right)\right) - H\left(\Psi\left(\mathbb{E}\left(\sqrt{1 - (T_\rho f)^2(\mathbf{Y})}\right)\right)\right) \leq H(\Psi(1)) - H\left(\Psi\left(\sqrt{1 - \rho^2}\right)\right).$$

However Conjecture-SH implies

$$\sqrt{1 - \mathbb{E}(f)^2} - \mathbb{E}\left(\sqrt{1 - (T_\rho f)^2(\mathbf{Y})}\right) \leq 1 - \sqrt{1 - \rho^2}.$$

Apply weak-majorization inequality: in particular use convexity, non-negativeness, and increasing property of $H(\Psi(x))$.



Extremality of Hellinger Conjecture

Conjecture-SH states

$$\sqrt{1 - \mathbb{E}(f)^2} - \mathbb{E} \left(\sqrt{1 - (T_\rho f)^2(\mathbf{Y})} \right) \leq 1 - \sqrt{1 - \rho^2}.$$

Take $\lim \rho \rightarrow 1$ (clean channel).

If Conjecture-SH is true then (for balanced Boolean functions)

$$\mathbb{E} \left(\sqrt{Sen_f(\mathbf{X})} \right) \geq 1.$$

$Sen_f(\mathbf{x})$: sensitivity at \mathbf{x} , number of neighbors with opposite value of $f(\mathbf{x})$.



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Similar limit for Conjecture-MI would be equivalent to (for balanced Boolean functions)

$$\mathbb{E} (Sen_f(\mathbf{X})) \geq 1.$$

This is known to be true (Poincare's inequality, Parseval's theorem, Harper's isoperimetric inequality)



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On the other hand, best lower bound for balanced functions

$$\mathbb{E} \left(\sqrt{Sen_f(\mathbf{X})} \right) \geq \sqrt{\frac{2}{\pi}} \quad (\text{Bobkov '98}).$$

Therefore even in this limit, the conjecture would imply something new.



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Lemma

For any $\alpha < \frac{1}{2}$, let $\text{maj}(\mathbf{Y})$ denote the majority function (assume that n is odd). Then there exists large enough n such that

$$\mathbb{E} (\text{Sen}_{\text{maj}}^\alpha(\mathbf{Y})) < \mathbb{E} (\text{Sen}_{\text{dic}}^\alpha(\mathbf{Y})) = 1.$$

Evidence to the veracity of Conjecture-SH

$$\sqrt{1 - \mathbb{E}(f)^2} - \mathbb{E} \left(\sqrt{1 - (T_\rho f)^2(\mathbf{Y})} \right) \leq 1 - \sqrt{1 - \rho^2}.$$

- ❶ verified numerically until $n = 8$
- ❷ Conjecture-SH is true if

$$\sqrt{1 - \rho^2} + \sqrt{1 - \mathbb{E}(f)^2} \leq 1$$



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On numerical verification

- Issue: Number of Boolean functions is 2^{2^n}
- **Lemma:** For any convex Φ , there is a *doubly monotone* boolean function that maximizes the Φ -entropy, $\mathbb{E}(\Phi(Z_f)) - \Phi(\mathbb{E}(Z_f))$, where maximization is over all boolean functions.
- While the number of *doubly monotone* boolean functions also grows doubly exponentially, still amenable till $n = 8$ (or a bit more).

On lemma

Doubly-monotone: A boolean function is said to be doubly-monotone if it is monotone, and for any $1 \leq i < j \leq n$,

$$f(S \cup \{i, j\}) \geq f(S \cup \{i\}) \geq f(S \cup \{j\}) \geq f(S), \quad \forall S \subseteq [1 : n].$$

Proof: Follows from majorization and Karamata's inequality.
Similar argument also present in (Courtade-Kumar '14)



2nd evidence: Proof in the parameter regime

$$G(\lambda) := \sqrt{1 - (1 - \lambda) \mathbb{E}(f)^2} - \mathbb{E} \left(\sqrt{1 - \lambda \rho^2 - (1 - \lambda) g^2(\mathbf{y})} \right),$$

where $g(\mathbf{y}) = (T_\rho f)(\mathbf{y})$. Want to show $G(1) \geq G(0)$.

$$G'(\lambda) = \frac{\mathbb{E}(f)^2}{2\sqrt{1 - (1 - \lambda) \mathbb{E}(f)^2}} - \mathbb{E} \left(\frac{g^2(\mathbf{y}) - \rho^2}{2\sqrt{1 - \lambda \rho^2 - (1 - \lambda) g^2(\mathbf{y})}} \right)$$



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Lemma

For any $0 \leq \rho^2, \lambda \leq 1$ the function

$$f(u) := \frac{u - \rho^2}{\sqrt{1 - \lambda \rho^2 - (1 - \lambda) u}}$$

is convex and increasing in u when $u \in [0, 1]$.

If U is a random variable that takes values in $[0, 1]$ then

$$E(f(U)) \leq (1 - E(U))f(0) + E(U)f(1).$$



Calculus of variations..

Denoting $\alpha = E(g^2(Y))$ we obtain

$$G'(\lambda) \geq \frac{E(f)^2}{2\sqrt{1 - (1 - \lambda)E(f)^2}} + (1 - \alpha) \frac{\rho^2}{2\sqrt{1 - \lambda\rho^2}} - \alpha \frac{\sqrt{1 - \rho^2}}{2\sqrt{\lambda}}. \quad (1)$$

Thus, integrating both sides with respect to λ from 0 to 1 we obtain

$$\int_0^1 G'(\lambda) d\lambda \geq 2 - \alpha - \sqrt{1 - E(f)^2} - \sqrt{1 - \rho^2}.$$

Since $\alpha \leq E(f)^2 + \rho^2(1 - E(f)^2)$ (by Parseval) we are done if

$$1 + (1 - E(f)^2)(1 - \rho^2) \geq \sqrt{1 - E(f)^2} + \sqrt{1 - \rho^2}.$$



Weaker form

Conjecture[†]-SH-W

For any pairs of boolean functions $f(\mathbf{X}), g(\mathbf{Y})$ on the hypercube taking values in $\{-1, +1\}$,

$$\sqrt{1 - \mathbb{E}(f(\mathbf{X}))^2} - \mathbb{E} \left(\sqrt{1 - \mathbb{E}(f(\mathbf{X})|g(\mathbf{Y}))^2} \right) \leq 1 - \sqrt{1 - \rho^2}.$$



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If this is true, then it would imply the following:

Proposition (Pichler et.al. '16)

For any pairs of boolean functions $f(\mathbf{X}), g(\mathbf{Y})$ on the hypercube,

$$I(f(\mathbf{X}); g(\mathbf{Y})) \leq 1 - H_b \left(\frac{1 - \rho}{2} \right)$$

However the techniques used here is insufficient to prove Conjecture[†]-SH-W.



Why is the weak form true[†]?

Conjecture[†]

For any pair of binary random variables (U, V) with V taking values in $\{-1, 1\}$ the following holds:

$$\sqrt{1 - \mathbb{E}(V)^2} - \mathbb{E} \left(\sqrt{1 - \mathbb{E}(V|U)^2} \right) + \sqrt{1 - s_{\dagger}(U; V)} \leq 1,$$

where $s_{\dagger}(U; V) = \lim_{p \rightarrow 0} s_p(U; V)$.

$s_p(U; V)$: reverse-hypercontractivity parameter.



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Reverse-hypercontractivity

A pair of random variables (U, V) is said to be (p, q) -reverse-hypercontractive for $1 > q \geq p$ if

$$\mathbb{E}(f(U)g(V)) \geq \|f(U)\|_{p'} \|g(V)\|_q.$$

For an fixed $p < 1$ define

$$s_p(U; V) := \sup \left\{ \frac{q-1}{p-1} : (U, V) \text{ is } (p, q)\text{-reverse-hypercontractive} \right\}.$$

An inequality[†]

For any $(s, c, d) \in [0, 1]$, the inequality

$$\sqrt{1 - (s(\bar{d} - d) + \bar{s}(c - \bar{c}))^2} - s\sqrt{1 - (\bar{d} - d)^2} - \bar{s}\sqrt{1 - (\bar{c} - c)^2} + \sqrt{\frac{D(s\bar{c} + \bar{s}d\|s\bar{d} + \bar{s}c)}{sD(c\|d) + \bar{s}D(d\|c)}} \leq 1.$$

- Seems to be true (numerical simulations)
- If true, will imply Conjecture[†] (hence Conjecture[†]-SH-W)
- Can formally establish it for certain parameters, including a neighborhood of equality achieving points*.



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Remarks:

- Conjecture[†] may also be of independent interest
- Can possibly obtain a computer assisted formal proof of the inequality above



Conclusion

Proposed $\Phi_{\mathcal{H}^2}(x)$ to be *an extremal* function for which dictator maximizes the Φ -entropy.

- Gave a proof in some (limited) parameter regimes

Proposed an explicit three variable inequality that establishes the weaker form

- This is done via another inequality involving hypercontractivity



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Lots of open conjectures

Connections to deeper stuff

- Talagrand's inequality (via Bobkov)

Thank You

